linux中断

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一、发生中断时的执行流程

W(b) vector_irq + stubs_offset -> vector_irq -> __irq_usr或者__irq_svc -> irq_handler -> gic_handle_irq -> handle_IRQ -> generic_handle_irq -> generic_handle_irq_desc -> desc->handle_irq

发生中断时,cpu进入irq异常模式,会跑到异常向量表运行。linux在之前已经初始化了异常向量表和异常处理代码,异常向量表在虚拟地址0xffff0000处。异常向量表和处理代码在arch/arm/kernel/entry_armv.S中定义。

```
/* 异常处理函数 */
   .globl __stubs_start
__stubs_start: // 异常处理函数起始
   // irq、data abort、prefetch abort和udef异常处理函数都用vector_stub宏定义
   vector_stub irq, IRQ_MODE, 4
   .long __irq_usr
                               @ 如果进入中断前是用户态则跳到___irq_usr
   .long __irq_invalid
                               a 1
   .long __irq_invalid
                               @ 如果进入中断前是内核态则跳到__irq_svc
   .long __irq_svc
                             a 4
   .long __irq_invalid
   .long __irq_invalid
                               a 5
   .long __irq_invalid
   .long __irq_invalid
                               @ 7
   .long __irq_invalid
                               a 8
                               @ 9
   .long __irq_invalid
   .long __irq_invalid
                               @ a
   .long __irq_invalid
                              @ b
   .long __irq_invalid
   .long __irq_invalid
   .long __irq_invalid
   .long __irq_invalid
   vector_stub dabt, ABT_MODE, 8
   .long __dabt_usr
                           @ 0 (USR_26 / USR_32)
   .long __dabt_invalid
                             @ 1 (FIQ_26 / FIQ_32)
          __dabt_invalid
   .long
                              @ 2 (IRQ_26 / IRQ_32)
   .long
          __dabt_svc
                           @ 3 (SVC_26 / SVC_32)
```

```
.long __dabt_invalid @ 4
. . .
// fiq、swi和address异常单独使用特定编写的处理函数
vector_fiq:
   subs pc, lr, #4
vector_addrexcptn:
   b vector_addrexcptn
   .align 5
.LCvswi:
   .word vector_swi
   .globl __stubs_end
......
/* 异常向量表 */
   .equ stubs_offset, __vectors_start + 0x200 - __stubs_start
   .globl __vectors_start
__vectors_start: // 异常向量表起始
ARM( swi SYS_ERROR0 )
   W(b) vector_und + stubs_offset
   W(ldr) pc, .LCvswi + stubs_offset
   w(b) vector_pabt + stubs_offset
   W(b) vector_dabt + stubs_offset
   w(b) vector_addrexcptn + stubs_offset
   W(b)
         vector_irq + stubs_offset
   W(b)
        vector_fiq + stubs_offset
   .globl __vectors_end
__vectors_end: // 异常向量表结束
......
/* 除了fiq、swi和address异常之外的异常处理函数宏 */
vector_\name:
   .if \correction
   sub lr, lr, #\correction //修正返回地址lr_<exception>
   .endif
   @ Save r0, lr_<exception> (parent PC) and spsr_<exception>
   ( (parent CPSR)
   /* 将r0,修正后的1r_<exception>, spsr_<exception>保存到sp_<exception>中,
   * spsr_<exception>是进入中断前的cpsr, sp_<exception>空间很小
   */
   stmia sp, {r0, lr}
                      @ save r0, 1r
   mrs lr, spsr @ lr = spsr_<exception>
   str lr, [sp, #8]
                   @ save spsr
   a
```

```
@ Prepare for SVC32 mode. IRQs remain disabled.
   /* 将cpsr_<exception>读到r0中,然后将模式域切换为SVC模式,
   * 再将r0写到spsr_<exception>,因此现在还是<exception>不是svc,
   * 同时未更改前的spsr_<exception>也保存着副本在sp_<exception>中
   mrs r0, cpsr
   eor r0, r0, #(\mode ^ SVC_MODE | PSR_ISETSTATE)
   msr spsr_cxsf, r0
   @ the branch table must immediately follow this code
   /*
   * 保存sp_<exception>到r0,
   *根据发生中断前的模式进入对应的子处理程序,如irq模式下的__irq_usr或者__irq_svc,
   * 同时将以修改的spsr_<exception>恢复到cpsr中,因此等于将中断前的cpsr的模式域切换为svc再恢复
   and lr, lr, #0x0f //获取修改前的spsr_<exception>即中断前的cpsr的模式
   mov r0, sp //将sp_<exception>保存到r0,然后可以在子程序中使用
   movs pc, 1r
                 // 根据1r跳转到适合的子处理程序如irq模式下的__irq_usr或者__irq_svc,
                   //因为movs,所以同时cpsr = 修改过的spsr_<exception>
ENDPROC(vector_\name)
   .align 2
   @ handler addresses follow this label
1:
   endm
   .align 5
__irq_usr:
   usr_entry // 保存usr中断现场
   kuser_cmpxchg_check
   irq_handler //
   get_thread_info tsk
mov why, #0
//wh
                     //why = 0
   b ret_to_user_from_irq //
ENDPROC(__irq_usr)
   .align 5
__irq_svc:
   svc_entry// 保存svc中断现场irq_handler// 调用gic_handle_irq
#ifdef CONFIG_PREEMPT
   get_thread_info tsk //
   teq r8, #0 @ if preempt count != 0
   movne r0, #0 @ force flags to 0
   tst r0, #_TIF_NEED_RESCHED
```

```
blne svc_preempt //
#endif
  svc_exit r5 @ return from exception
ENDPROC(__irq_svc)
.macro irg_handler
#ifdef CONFIG_MULTI_IRQ_HANDLER
  ldr r1, =handle_arch_irq //对于socfpga会设置为gic_handle_irq
  mov r0, sp //传给gic_handle_irq的参数r0 = sp_svc,即pt_regs
  adr lr, BSYM(9997f)
  ldr pc, [r1] //调用gic_handle_irq
#else
  arch_irq_handler_default
#endif
9997:
  .endm
```

位于__vectors_start和__vectors_end之间的是真正的异常向量表,位于__stubs_start和__stubs_end之间的是处理代码。对于irq中断先进入W(b) vector_irq + stubs_offset然后进入vector_irq , 根据进入irq前的模式进入__irq_usr或者__irq_svc , 无论是__irq_usr还是__irq_svc都会进入irq_handler , 然后调用handle_arch_irq即gic_handle_irq。其中__irq_usr和__irq_svc分别是保存从用户和内核态进入中断的保存中断现场。

```
.macro svc_entry, stack_hole=0
    /* sp_svc = sp_svc - (S_FRAME_SIZE - 4), S_FRAME_SIZE是struct pt_regs的大小
    */
    //sp_svc指向pt_regs中的ARM_r1
   sub sp, sp, #(S_FRAME_SIZE + \stack_hole - 4)
   stmia sp, {r1 - r12} //将r1到r12保存到sp_svc中,因为r1到r12在vector_\name中没有改变,因此
                      //这里的r1到r12是发生中断异常前的r1到r12
   ldmia r0, {r3 - r5} //r0保存着sp_<exception>,因此这里把vector_\name中保存到
sp_<exception>
                       //中的r0,1r_<exception>, spsr_<exception>恢复到r3,r4,r5
   /* S_SP是offsetof(struct pt_regs, ARM_sp)
   * 将pt_regs中ARM_sp处的地址给r7
   add r7, sp, #S_SP - 4 @ here for interlock avoidance
                        a "" "" ""
   mov r6, #-1
   add r2, sp, #(S_FRAME_SIZE + \stack_hole - 4) //r2保存原来未发生异常前的sp_svc
   // 将发生中断异常前的r0保存到sp_svc中的pt_regs中
   str r3, [sp, #-4]! @ save the "real" r0 copied
                   @ from the exception stack
   mov r3, lr //将lr_svc保存在r3中
   @ We are now ready to fill in the remaining blanks on the stack:
```

```
@ r2 - sp_svc
   @ r3 - 1r_svc
   @ r4 - lr_<exception>, already fixed up for correct return/restart
   @ r5 - spsr_<exception>
   @ r6 - orig_r0 (see pt_regs definition in ptrace.h)
   stmia r7, {r2 - r6} //将最后几个保存到pt_regs中
   .endm
.macro usr_entry
   sub sp, sp, #S_FRAME_SIZE //sp_svc指向pt_regs中的ARM_r0
   stmib sp, {r1 - r12} //将发生中断前的r1到r12保存到sp_svc的pt_regs中
   //将中断前的r0, pc(lr_<exception>)和cpsr(spsr_<exception>)恢复到r3,r4,r5
   1dmia r0, \{r3 - r5\}
   //r0指向pt_regs中的ARM_pc
   add r0, sp, #S_PC @ here for interlock avoidance
                    @ "" ""
   mov r6, #-1
   // 将发生中断异常前的r0保存到pt_regs中
   str r3, [sp]
                @ save the "real" r0 copied
                 @ from the exception stack
   @ We are now ready to fill in the remaining blanks on the stack:
   @ r4 - lr_<exception>, already fixed up for correct return/restart
   @ r5 - spsr_<exception>
   @ r6 - orig_r0 (see pt_regs definition in ptrace.h)
   @ Also, separately save sp_usr and lr_usr
   stmia r0, {r4 - r6} //将这几个保存到pt_regs中,和svc不同的是没有保存sp_usr和1r_usr,放到下一
步
   stmdb r0, {sp, 1r}^ //将sp_usr和1r_usr保存到pt_regs中,因为^,所以不是sp_svc和1r_svc
   .endm
```

看了保存中断现场,现在来看看恢复中断现场。

```
/* 恢复svc中断现场 */
.macro svc_exit, rpsr
msr spsr_cxsf, \rpsr //将spsr_<exception>恢复, 因为传参是r5
...
//恢复所有的pt_regs, 同时spsr_<exception>写到cpsr, 如果发生中断前中断是打开,那么
ldmia sp, {r0 - pc}^ @ load r0 - pc, cpsr
.endm
/* 恢复usr中断现场 */
```

```
ENTRY(ret_to_user_from_irq)
   ldr r1, [tsk, #TI_FLAGS] //从svc栈的thread_info获取发生中断前的进程的task_struct中的flags
   tst r1, #_TIF_WORK_MASK
                         //获取flags中的work部分
   bne work_pending //如果不为0说明有工作要做(有信号或者要被调度)
no_work_pending:
   //如果没有工作就正常恢复用户中断现场
   arch_ret_to_user r1, lr //什么都不做
   restore_user_regs fast = 0, offset = 0 //将保存的用户寄存器恢复
ENDPROC(ret_to_user_from_irq)
   /* 将保存的用户寄存器恢复 */
   .macro restore_user_regs, fast = 0, offset = 0
   ldr r1, [sp, #\offset + S_PSR] //获取发生异常前的cpsr到r1
   ldr lr, [sp, #\offset + S_PC]! //获取异常后返回的地址到Lr_svc, 同时更新sp_svc指向
   msr spsr_cxsf, r1 //将发生异常前的cpsr写到spsr_svc中
   . . .
   .if \fast
   ldmdb sp, \{r1 - lr\}^{\wedge} @ get calling r1 - lr
   .else
                             //将sp_svc指向的pt_regs恢复到r0到r12以及sp_usr和lr_usr
   ldmdb sp, {r0 - lr}∧
   .endif
   mov r0, r0
                       @ ARMv5T and earlier require a nop
                    @ after ldm {}^
   add sp, sp, #S_FRAME_SIZE - S_PC //将sp_svc指回没发生异常前的地方
   movs pc, lr
                    //返回发生异常前的地址,同时恢复异常前的cpsr
   .endm
/* 处理pending事物,调度或者处理信号 */
work_pending:
                     @ 'regs'
   mov r0, sp
   mov r2, why
                       @ 'syscall'
   bl do_work_pending //调度或者处理信号(没打开中断,为什么能调度?因为中断已经处理完)
   cmp r0, #0 //如果do_work_pending返回值为0则跳到no_work_pending,否则要restar还是strace什么的
   beq no_work_pending
   movlt scno, #(__NR_restart_syscall - __NR_SYSCALL_BASE)
   ldmia sp, \{r0 - r6\} @ have to reload r0 - r6
   b local_restart @ ... and off we go
/* 调度或处理信号 */
asmlinkage int do_work_pending(struct pt_regs *regs, unsigned int thread_flags, int
syscall)
{
   do {
      if (likely(thread_flags & _TIF_NEED_RESCHED)) { //调度
          schedule();
      } else { //处理信号
         if (unlikely(!user_mode(regs)))
             return 0;
          local_irq_enable();
          if (thread_flags & _TIF_SIGPENDING) {
             int restart = do_signal(regs, syscall);
```

```
if (unlikely(restart)) {
                    /*
                     * Restart without handlers.
                     * Deal with it without leaving
                     * the kernel space.
                    return restart;
                }
                syscall = 0;
            } else {
                clear_thread_flag(TIF_NOTIFY_RESUME);
                tracehook_notify_resume(regs);
            }
        }
        local_irq_disable();
        thread_flags = current_thread_info()->flags;
    } while (thread_flags & _TIF_WORK_MASK);
    return 0;
}
```

```
asmlinkage void __exception_irq_entry gic_handle_irq(struct pt_regs *regs)
{
   u32 irqstat, irqnr;
   struct gic_chip_data *gic = &gic_data[0];
   void __iomem *cpu_base = gic_data_cpu_base(gic);
   do {
      irqstat = readl_relaxed(cpu_base + GIC_CPU_INTACK);
      irqnr = irqstat & ~0x1c00; /* 读出硬件中断号 */
*************
                               Embest Tech co., 1td
                               www.embest-tech.com
**************
*chage the IPI interrupt handler, cause the another cpu not secheule the linux thread any
more.
*/
      if (likely(irqnr < 1021)) {</pre>
          if(irqnr<16)</pre>
          {
             /* 如果是sgi中断则写EOI */
             writel_relaxed(irqstat, cpu_base + GIC_CPU_EOI);
          irqnr = irq_find_mapping(gic->domain, irqnr); /* 硬件中断号转化为虚拟中断号 */
          handle_IRQ(irqnr, regs); /* 进入irqnr的通用层处理 */
          continue;
      }
```

```
break:
   } while (1);
}
void handle_IRQ(unsigned int irq, struct pt_regs *regs)
{
   struct pt_regs *old_regs = set_irq_regs(regs);
   irq_enter(); /* 增加preempt count HARDOFF */
   if (unlikely(irq >= nr_irqs)) {
       if (printk_ratelimit())
           printk(KERN_WARNING "Bad IRQ%u\n", irq);
       ack_bad_irq(irq);
   } else {
       generic_handle_irq(irq); /* 进入通用层 */
   }
   irq_exit(); /*减少preempt count HARDOFF,同时检查是否可以和需要执行软中断 */
   set_irq_regs(old_regs);
}
int generic_handle_irq(unsigned int irq)
   struct irq_desc *desc = irq_to_desc(irq); //取出irq对应的irq_desc
   if (!desc)
       return -EINVAL;
   generic_handle_irq_desc(irq, desc);
   return 0;
}
 static inline void generic_handle_irq_desc(unsigned int irq, struct irq_desc *desc)
{
   desc->handle_irq(irq, desc); //调用流控层函数,对于gic的spi调用的是fast_eoi,在下文中的初始化中
设置
}
```

二、系统启动时的中断初始化

```
.reserve = socfpga_ucosii_reserve,
   .dt_compat = altera_dt_match,
MACHINE_END
#define DT_MACHINE_START(_name, _namestr)
static const struct machine_desc __mach_desc_##_name \
__used
__attribute__((__section__(".arch.info.init"))) = {
                  \
   .nr = \sim 0
           = _namestr,
   .name
asmlinkage void __init start_kernel(void)
   setup_arch(&command_line);
   early_irq_init();
   init_IRQ();
   . . .
}
```

```
void __init setup_arch(char **cmdline_p)
{
   struct machine_desc *mdesc;
   setup_processor();
   mdesc = setup_machine_fdt(__atags_pointer); /* 获取机器描述符 */
   if (!mdesc)
       mdesc = setup_machine_tags(__atags_pointer, machine_arch_type);
   machine_desc = mdesc;
   machine_name = mdesc->name;
   setup_dma_zone(mdesc);
   if (mdesc->restart_mode)
        reboot_setup(&mdesc->restart_mode);
   init_mm.start_code = (unsigned long) _text;
   init_mm.end_code = (unsigned long) _etext;
   init_mm.end_data = (unsigned long) _edata;
   init_mm.brk = (unsigned long) _end;
   /* populate cmd_line too for later use, preserving boot_command_line */
   strlcpy(cmd_line, boot_command_line, COMMAND_LINE_SIZE);
    *cmdline_p = cmd_line;
   parse_early_param();
   sort(&meminfo.bank, meminfo.nr_banks, sizeof(meminfo.bank[0]), meminfo_cmp, NULL);
   sanity_check_meminfo();
   arm_memblock_init(&meminfo, mdesc);
```

```
paging_init(mdesc); /* 初始化一些页表映射和异常向量表 */
    request_standard_resources(mdesc);
   if (mdesc->restart)
       arm_pm_restart = mdesc->restart;  /* */
   unflatten_device_tree();
#ifdef CONFIG_SMP
   if (is_smp()) {
       smp_set_ops(mdesc->smp);
       smp_init_cpus();
#endif
   if (!is_smp())
       hyp_mode_check();
   reserve_crashkernel();
   tcm_init();
#ifdef CONFIG_MULTI_IRQ_HANDLER
   handle_arch_irq = mdesc->handle_irq;  /* */
#endif
#ifdef CONFIG_VT
#if defined(CONFIG_VGA_CONSOLE)
   conswitchp = &vga_con;
#elif defined(CONFIG_DUMMY_CONSOLE)
   conswitchp = &dummy_con;
#endif
#endif
   if (mdesc->init_early)
       mdesc->init_early();
}
```

```
if (WARN_ON(initcnt > IRQ_BITMAP_BITS))
    initcnt = IRQ_BITMAP_BITS;

if (initcnt > nr_irqs)
    nr_irqs = initcnt;

for (i = 0; i < initcnt; i++) {
    desc = alloc_desc(i, node, NULL);
    set_bit(i, allocated_irqs);
    irq_insert_desc(i, desc);
}
return arch_early_irq_init();
}</pre>
```

三、中断映射

四、中断注册