

# 信号量测试

除了互斥锁外，信号量是另外一种对共享资源进行保护的机制。互斥锁其实是特殊的信号量，只有0和1，而一般信号量的值可以超过1。下面展示有无信号量的测试对比。

无信号量:

The screenshot displays a debugger interface with the following components:

- Debug Control:** Shows the status of two CPUs. Both `cpu0` and `cpu1` are connected and have stopped in SVC mode.
- Commands:** Shows the execution flow. It indicates that the program stopped in SVC mode at `S:0x1E001448` in `main.c`. The command `ret = verify((u8 *)p, (u8 *)&cnt, 16);` is highlighted.
- Locals:** A table showing the values of local variables for the selected CPU.
- Disassembly:** A table showing the instructions being executed, including `verify` and `ret`.
- App Console:** Shows the output of the program, including the message `core0 wrong cnt %d\r\n`.

Name	Value	Type
i	Optimised away	s32
j	Optimised away	s32
cpu_flag	0	volatile u32*
p	0x20000000	u32*
cnt	13095432	u32
wrong_cnt	1079	u32
ret	0	s32

Address	Opcode	Disassembly
S:0x1E001638	E24B3018	SUB r3, r11, #0x18
S:0x1E00163C	E1A01003	MOV r1, r3
S:0x1E001640	E3A02010	MOV r2, #0x10
S:0x1E001644	E0FFFF6F	BL verify ; 0x1E001648
S:0x1E001648	E51B3010	STR r0, [r11, #-0x10]
S:0x1E00164C	E3530000	LDR r3, #0
S:0x1E001650	0AFFFFEC	BEQ {pc}-0x48 ; 0x1E001658
S:0x1E001658	E51B3008	LDR r3, [r11, #-8]
S:0x1E00165C	E2833001	ADD r3, r3, #1
S:0x1E001660	E50B3008	STR r3, [r11, #-8]
S:0x1E001664	E30C01B4	MOV r0, #0xc1b4
S:0x1E001668	E3410E00	MOVT r0, #0x1e00

Debug Control | Project Explorer | Remote System... | Commands | History | Scripts | Linked: cpu1

cpu0 connected  
Cortex-A9\_0 #1 stopped (SVC)  
cpu1 connected  
Cortex-A9\_1 #1 stopped (SVC)

Status: connected

Command: Press (Alt+/) for Content Assist Submit

main.c

```

156
157 #if 1
158 /*
159  * this part is for test without mutex and semaphore
160  */
161
162 u32 *p = 0x20000000; // 读写数据总线
163 u32 cnt = 0xffffffff;
164 u32 wrong_cnt = 0;
165 s32 ret;
166 while (1) {
167     cnt--;
168     // led15_blink();
169     // yeshe_mdelay_nobreak(100);
170
171     /* 读写数据总线+16字节(16字节) 读写数据 */
172     memset(p, cnt & 0xff, 16);
173
174     /* 读写数据+16字节 读写 数据 */
175     ret = verify((u8 *)p, (u8 *)&cnt, 16);
176     if (ret) {
177         wrong_cnt++;
178         uartprintf("core1 wrong cnt %d\n", wrong_cnt);
179     }
180 }
181 #endif
182
183

```

Execution stopped in SVC mode at S:0x1F001538  
S:0x1F001538 175,0 ret = verify((u8 \*)p, (u8 \*)&cnt, 16);  
wait  
continue  
interrupt  
Execution stopped in SVC mode at S:0x1F0013EC  
S:0x1F0013EC 67,0 if (memcmp(ptr+i, val, 1))  
wait  
finish  
Execution stopped in SVC mode at S:0x1F001538  
S:0x1F001538 175,0 ret = verify((u8 \*)p, (u8 \*)&cnt, 16);

Linked: cpu1

Name	Value	Type
i	Optimised away	s32
j	Optimised away	s32
cpu_flag	1	volatile u32*
p	0x20000000	u32*
cnt	4285694136	u32
wrong_cnt	1089	u32
ret	0	s32

Disassembly

Address	Opcode	Disassembly
S:0x1F001528	E24B3018	SUB r3, r11, #0x18
S:0x1F00152C	E1A01003	MOV r1, r3
S:0x1F001530	E3A02010	MOV r2, #0x10
S:0x1F001534	E8FFFF9C	BL verify ; 0x1F001
S:0x1F001538	E51B3010	STR r0, [r11, #-0x10]
S:0x1F00153C	E51B3010	LDR r3, [r11, #-0x10]
S:0x1F001540	E3530000	CMP r3, #0
S:0x1F001544	0AFFFFEC	BEQ {pc} -0x48 ; 0x1F
S:0x1F001548	E51B3008	LDR r3, [r11, #-8]
S:0x1F00154C	E2833001	ADD r3, r3, #1
S:0x1F001550	E50B3008	STR r3, [r11, #-8]
S:0x1F001554	E30B032C	MOV r0, #0xb32c
S:0x1F001558	E3410F00	MOVT r0, #0x1f00

App Console | Target Console... | Error Log

Linked: cpu1

A suitable Debug Server is already running

以上展示的是两个核在没有任何同步保护机制的情况下对同一块共享内存进行写入然后读取验证。wrong\_cnt记录验证错误的次数，也就是读出来的值不是之前写进去的值的次数，从图中看到两个核在进行一定次数的测试后分别发生了1079次和1089次的验证数据错误，因此没有信号量等同步机制的情况下访问共享资源是不安全的。

**有信号量:**

**阻塞版:**



上图是两个核利用信号量对共享内存进行读写的测试，每个核在操作共享内存时都要通过init\_semaphore初始化信号量，传入的参数为初始化的信号量的值，第一个调用init\_semaphore的线程成功的初始化信号量，第二个调用init\_semaphore的线程无法初始化信号量为传入的参数值，只是单纯返回信号量地址，信号量值由第一个线程是否已经释放信号量来决定。每个核在操作共享内存时候通过sem\_dec将信号量减1，操作完后通过sem\_inc将信号量加1，当信号量大于0时，sem\_dec才能成功返回并将信号量减1，否则一直阻塞直到信号量大于0。从图中可以发现两个核的测试中的wrong\_cnt一直为0，即一直没有出错，因此本文实现的信号量能很好的保护双核通信中共享资源的访问。

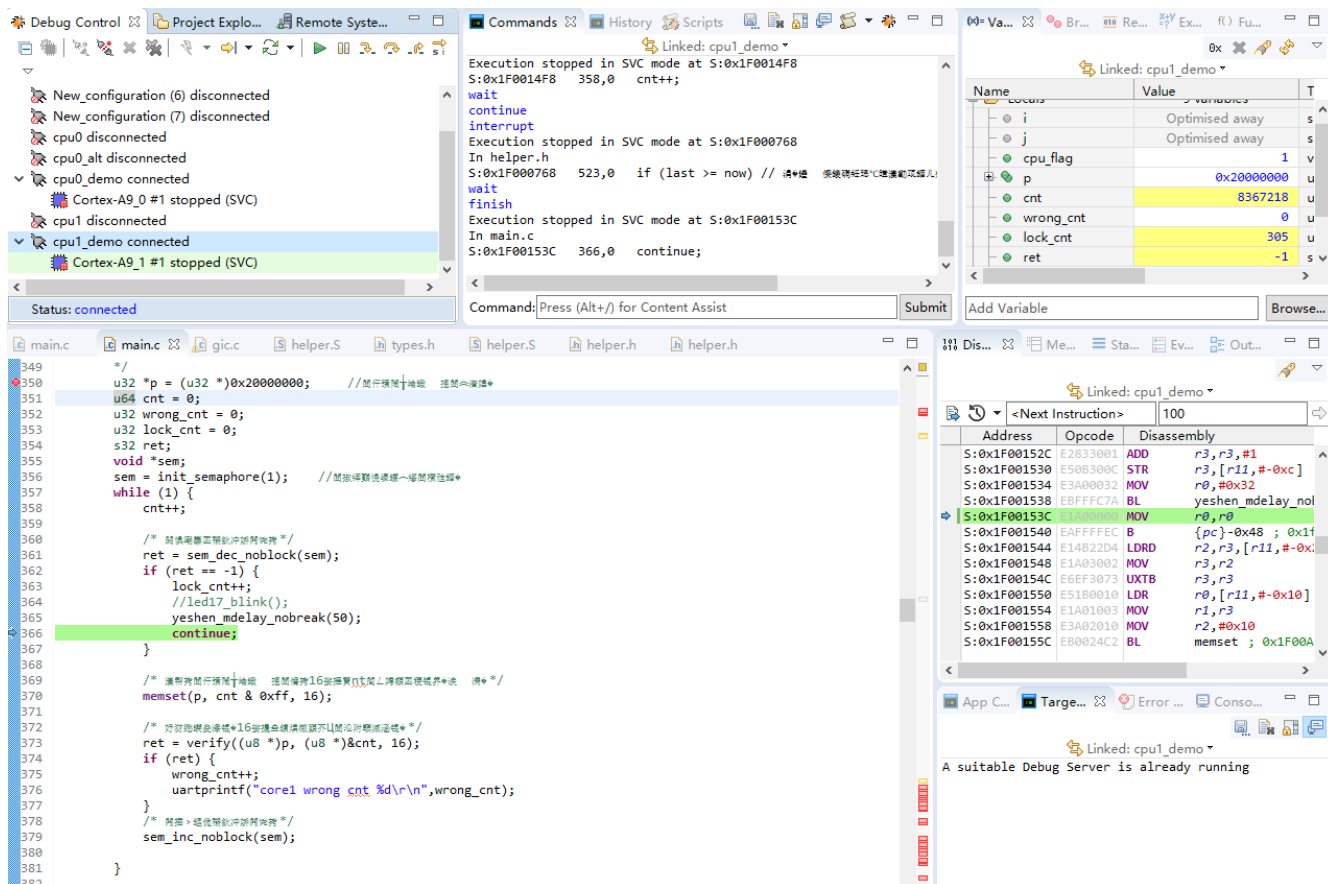
## 非阻塞版：

上述版本的信号量是阻塞版的，就是当本线程调用sem\_dec获取信号量时发现信号量等于0的时候会阻塞在临界区，直到信号量被其他线程通过sem\_inc加1后sem\_dec获取了信号量并减1后才能返回访问共享资源。在这阻塞期间本线程什么都不能做，本文提供了另外一种非阻塞方案，可以让用户选择在信号量不大于0被占用的情况下是否做其他工作。

The screenshot displays a development environment with three main panels:

- Left Panel (Debug Control):** Shows the status of various components. 'Cortex-A9\_0 #1 stopped (SVC)' and 'Cortex-A9\_1 #1 stopped (SVC)' are highlighted in green, indicating the current execution state.
- Center Panel (Commands):** Displays the execution log. It shows the program execution stopping in SVC mode at different memory addresses (S:0x1E001650, S:0x1E001424, S:0x1E0016CC) and the corresponding assembly instructions being executed, such as 'wait', 'continue', 'interrupt', and 'finish'.
- Right Panel (Variable Watch):** Shows a list of variables and their values. The 'p' variable is highlighted in yellow, showing its value as 0x20000000. Other variables like 'cnt', 'wrong\_cnt', 'lock\_cnt', 'ret', and 'sem' are also listed.

Below the command window, the source code for 'main.c' is visible, showing the implementation of the semaphore and the test logic. The code includes comments in Chinese and uses standard C syntax for memory management and thread synchronization.



非阻塞版和阻塞版其中行为不同的只是获取并将信号量减1的操作，初始化和释放信号量的行为一致。而获取信号量并减1在不成功的时候直接返回-1,表示信号量已经为0，返回0表示获取信号量并减1成功，用户可以根据返回值作出相应的动作，不必一直阻塞直到信号量值大于0。从图中看到两个核的测试中wrong\_cnt也是一直为0，说明共享资源访问没有冲突，而另外用一个lock\_cnt的变量记录获取信号量时不大于0的次数，两个核分别有108和305次申请获取信号量时信号量不大于0。

## 总结

在没有同步机制保护的情况下双核同时访问共享资源会导致数据错乱冲突，在有信号量保护的情况下双核可以同时访问共享资源而不发生冲突，且根据应用的需求提供了阻塞版和非阻塞版两种信号量保护机制。