# Digital Systems L2 - Computer Systems

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#### DEVICE DRIVERS

- Introduction to Device Drivers
- Building and 'Running' a Module
- Character Device Drivers
- Debugging Techniques
- Kernel Timer





### Role of Device Drivers

All devices work under the control of certain softwares. These softwares are device drivers. In Linux, drivers can be part of kernel image, or independent code. This code is also called a module.

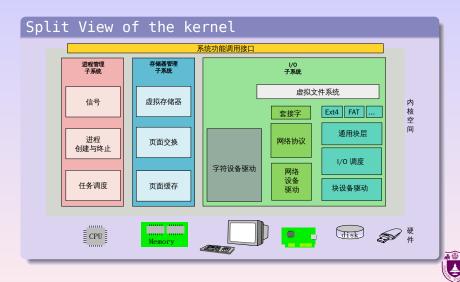
A flexible driver provides mechanism rather than policy.

- Mechanism
  What capabilities are to be
  provided
- Policy
  How those capabilities can be used





# Splitting the kernel



### Functions of the OS

- Process management
  - Creating, destroying and scheduling processes
  - Handling the Communication among different processes
  - Handling their connection to the outside world
- Memory management
   Virtual addressing space for all processes
   on top of the limited available resources.
- File systems
   Structured filesystem on top of unstructured hardware
- Device control
   Device drivers operate all physical devices.







# Types of Devices and Modules

Modules are part of kernel codes that can be added to or removed from dynamically.

- Character devices
  - A stream of bytes
  - Accessed by means of filesystem nodes
- Block devices can host a filesystem
- Network interfaces sending and receiving data packets





### Comments to the First Module

```
#include linux/module.h> // head files for kernel
int init module(void) // called when insmod
{
    printk(KERN_EMERG "Hello\n");
        // kernel version of printf
    return 0;
}
void cleanup module(void) // called when rmmod
{
    printk(KERN_EMERG "module_removed\n");
}
```



```
# Makefile (!not makefile)
ifneq ($(KERNELRELEASE),)
obj-m
            := hello.o
else
KDIR
            := /lib/modules/`uname -r`/build
            := $(shell pwd)
PWD
default:
        $(MAKE) -C $(KDIR) SUBDIRS=$(PWD) modules
clean:
        $(MAKE) -C $(KDIR) SUBDIRS=$(PWD) clean
endif
```



### Some Macros Unsed in Modules

o replace init module() and cleanup module():

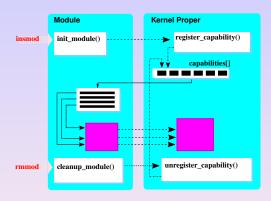
```
#include <linux/module.h>
module init(user init func)
module_exit(user_cleanup_func)
```

Extra infos

```
MODULE_LICENSE("GPL/LGPL/private")
MODULE_AUTHOR("zhang3u<zhang3@gmail.com>");
MODULE_DESCRIPTION("Ausimple_helloworld");
MODULE_VERSION("0:0.99");
```



# Kernel Modules vs. Applications



An application performs a single task from beginning to end, whereas a module registers itself in order to serve future requests, and its "main" function(init\_module) terminates immediately.



#### I/O ports

```
int check_region(unsigned long start,
    unsigned long len);
struct resource *request_region(
    unsigned long start,
    unsigned long len, char *name);
void release_region(unsigned long start,
    unsigned long len);
```

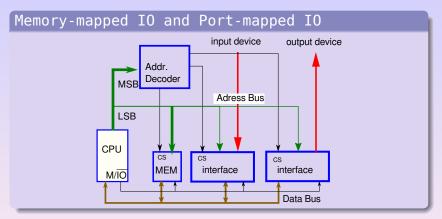


# Passing Arguments Kernel Module

```
In command line:
# insmod hello.ko val=3600
In module code:
#include <linux/moduleparam.h>
module_param(val, int, 0666);
available variable types: byte, short, ushort,
int, uint, long, ulong, charp, bool, invbool.
permission types:
S IWUSR S IRGRP
S IRUSR S IWGRP
 S IXUSR S IXGRP
```



# Different Memmory Spaces



CISC computers use different instructions (produce  $M/\overline{10}$ ) to access memory or IO.



#### I/O memory

```
int check_mem_region(unsigned long start,
    unsigned long len);
int request_mem_region(unsigned long start,
    unsigned long len, char *name);
int release_mem_region(unsigned long start,
    unsigned long len);
```



#### Register

```
int register_chrdev(unsigned int major,
    const char *name,
    struct file_operations *fops);
```

- major: a specified major number, or dynamic allocation(major=0)
- name:device filename, should be created by mknod
- fops: a global pointer to the struct file operations





- Char devices are accessed through names in the filesystem.
- The major number identifies the driver associated with the device.
- The minor number is used only by the driver specified by the major number.
- (with or without devfs)



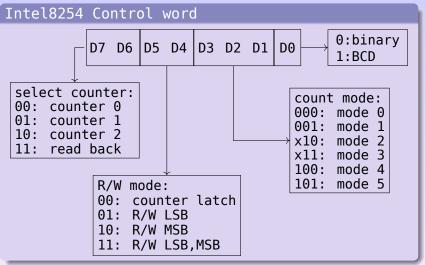


Each field in the structure file operations must point to the function in the driver that implements a specific operation, or be left NULL for unsupported operations. An example:

```
struct file_operations scull_fops = {
        lseek: scull_llseek,
        read: scull read,
        write: scull_write,
        ioctl: scull_ioctl,
        open: scull_open,
        release: scull release,
};
```

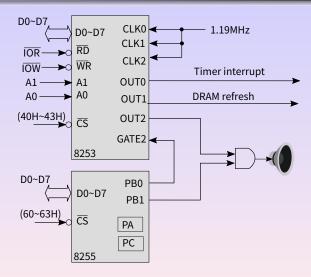


```
int (*open)(struct inode *node, struct file *filp);
If this entry is NULL, open() is always
succeeds without notified.
Device number node->i rdev in
include/linux/kdev t.h
struct file {
    . . . . . .
    struct path
                     f path;
    unsigned int
                   f_flags;
    fmode t
                     f mode;
    loff t
                     f pos;
    void
                     *private data;
    . . . . . .
```





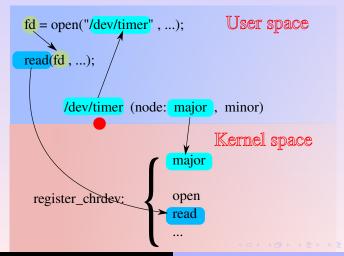
# Speaker in PC





#### Device Node

A device node connects between user space and kernel.





#### Removing a driver from the system

You also need to remove the device files. Create and remove device node manually, or using udev.





# Memory Usage

```
unsigned long copy_to_user(void *to,
      const void *from, unsigned long count);
unsigned long copy_from_user(void *to,
      const void *from, unsigned long count);
ssize_t dev_read(struct file *file, char *buf, size_t count, loff_t *ppos);
  struct file
                    Ruffer
                                               Buffer
  f count
                  (in the driver)
                                               (in the
  c flags
                                              application
  f mode
                               copy to user()
                                               or libc )
  f pos
```

#### Special Operations on I/O Devices

- Most devices can perform operations beyond simple data transfers
- User space must often be able to request operations other than read/write.
- These operations are usually supported via the ioctl (I/O Control) method.



### About \_IO macros (linux/ioctl.h)

- \_IOC\_TYPEBITS, 8 bits, just choose a number(magic number)
- \_IOC\_NRBITS, 8 bits, sequential number
- o \_IOC\_DIRBITS, 2 bits
  - \_IO | an ioctl with no parameters
  - \_IOW | an ioctl with write parameters
  - \_IOR | an ioctl with read parameters
  - \_IOWR | an ioctl with both w/r parameters
- \_ IOC\_SIZEBITS, 14 bits





Unlike X86, RISC machines have memory mapped I/O. I/O addresses of memory must be converted into virtual address before use.

```
#include <linux/io.h>
void *ioremap(unsigned long phys_addr,
    unsigned long size);
void iounmap(void * addr);
```





### Block I/O

Read/write I/O device may block when data is unavailable. The process must hang up in order to do other tasks until data transferring is available, then wake up to finish read/write.
#include linux/wait.h>

```
DECLARE_WAIT_QUEUE_HEAD(wq);
...
wait_event_interruptible(wait_queue_head_t queue,
        int condition);
...
void wake_up_interruptible(wait_queue_head_t *queue
```



Create proc file

The /proc is a special, software-created filesystem that is used by the kernel to export information to user space.

```
const char *name,
umode_t mode,
struct proc_dir_entry *parent,
const struct file_operations *proc_fops
```

The default proc file \*name is under /proc when \*parent is NULL.



Each file under /proc is tied to a kernel function that generates the file's "contents" on-the-fly when the file is read.

```
Remove proc file
void remove_proc_entry(const char *,
            struct proc_dir_entry *);
```





# Interrupt handling

```
Register and unregister an interrupt handler:
    int request_irq(unsigned int irq
        irqreturn_t (*handler)(int , void *,
    struct pt_regs*)
        unsigned long flags,
        const char *dev_name,
        void *dev_id);
    void free_irq(int irq, void *dev_id);
```



## Implementing a Handler

- 'Fast' interrupt (SA\_INTERRUPT) handlers are executed with interrupts disabled on the current processor.
- SA\_SHIRQ signals that the interrupt can be shared between devices.
- Interrupt handlers need to finish up quickly and not keep interrupts blocked for long. The interrupt handler is splitted into two halves:
  - The top half is the route that actually responds to the interrupt.
  - The bottom half is scheduled by the top half to be executed later at a safer time, with all interrupts enabled.





### Printing

### printk loglevel (linux/kernel.h)

- KERN\_EMERG: emergency messages(crash)
- KERN\_ALERT: requiring immediate action
- KERN\_CRIT: critical conditions(serious hardware or software failures)
- KERN\_ERR: error conditions
- KERN\_WARNING: warnings
- KERN\_NOTICE: security-related conditions
- KERN\_INFO: informational messages
- KERN\_DEBUG: debugging messages.

printk slows down the system noticeably, as printing causes a disk operation.



### Debuggers and Related Tools

- gdb gdb /usr/src/linux/vmlinux /proc/kcore compiling the kernel with debugging support (-g) results a huge vmlinux file.
- kdb The kernel debugger
- The integrated kernel debugger patch
- The kgdb patch
- Kernel crash dump analyzers
- The user-mode Linux port
- The Linux trace toolkit
- Dynamic probes





- Understanding kernel timing
   Timer interrupt, which is the mechanism the kernel uses to keep track of time intervals.
   The interval of timer interrupts is set by the value HZ.
- Using the jiffies counter(modified during timer interrupt)
- Processor-specific registers, rdtsc(low32, high32) (x86)
- Delaying operation for a specified amount of time: udelay(), mdelay()
- Scheduling asynchronous functions to happen after a specified time lapse



# Delay in Kernel

- udelay(), mdelay() and even ndelay() are software loop (busy-wait functions).
- while (jiffies < j) ;</pre>
- wait event timeout(jit delay\*HZ)





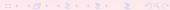
#### Calendar date and time

```
/* time in user space */
#include <sys/time.h>
int gettimeofday(struct timeval *tv,
        struct timezone *tz);
int settimeofday(const struct timeval *tv,
        const struct timezone *tz);
/* time in kernel space */
do_gettimeofday(struct timeval *tv);
```



```
/* Number of seconds since
    1970-01-01 00:00:00 +0000 (UTC) */
time_t time(time_t *tloc);
```





- Kernel Layout
- Kernel Modules
- Char Device Drivers
- PROCFS
- Kernel Timing





```
#include <linux/module.h>
int init module(void);
void cleanup_module(void);
module_init(x)
modile exit(x)
MODULE_LICENSE(string)
MODULE AUTHOR (string)
MODULE DESCRIPTION(string)
MODULE VERSION(string)
```





### Function List 2

```
#include <linux/fs.h>
struct file_operations {
    struct module *owner;
    loff t (*llseek) (struct file *, loff_t, int);
    ssize t (*read) (struct file *, char user *,
    ssize t (*write) (struct file *, const char u
    unsigned int (*poll) (struct file *, struct pol
    long (*unlocked ioctl) (struct file *, unsigned
    long (*compat ioctl) (struct file *, unsigned i
    int (*mmap) (struct file *, struct vm area stru
    int (*open) (struct inode *, struct file *);
    int (*flush) (struct file *, fl_owner_t id);
    int (*release) (struct inode *, struct file
    . . . ;
```

### Function List 3

```
#include <linux/fs.h>
int register chrdev(int major, char *name,
                    struct file_operations *);
/* extract major and minor */
MAJOR (kdev t dev);
MINOR (kdev t dev);
MKDEV (int major, int minor);
void unregister chrdev(int major, char *name);
```



### Function List 4

```
#include <linux/slab.h>
void *kmalloc(size_t size, gfp_t flags);
void kfree(const void *);
#include <linux/types.h>
typedef enum {
    GFP_KERNEL,
    GFP ATOMIC,
    __GFP_HIGHMEM,
    __GFP_HIGH
} gfp_t;
```



```
#include <linux/uaccess.h>
unsigned long copy_to_user(void *to,
                            const void *from,
                            unsigned long n);
unsigned long copy_from_user(void *to,
                              const void *from,
                              unsigned long n);
#include <linux/io.h>
void *ioremap (unsigned long phys_addr,
               unsigned long size);
void iounmap(void *addr);
```



```
#include <linux/proc fs.h>
struct proc_dir_entry *proc_create(
        const char *name,
        umode_t mode,
        struct proc_dir_entry *parent ,
        const struct file_operations *proc_fops);
void remove_proc_entry(
        const char *.
        struct proc dir entry *);
```



```
/* Platform Specific, may only on X86 available */
#include <linux/io.h>
unsigned char inb(unsigned long addr);
void outb(unsigned short value,
          unsigned long addr);
void insb(unsigned long addr,
          void *buffer,
          unsigned int count);
void outsb(unsigned long addr,
          const void *buffer,
          unsigned int count);
```

Suffix b, w, l specify 8bits, 16bits or 32bits



port access. s for string I/O.

```
Userspace port mapped IO (priviledge required!)
```

```
#include <sys/io.h>
int ioperm(unsigned long from,
           unsigned long num,
           int turn_on);
unsigned char inb(unsigned long addr);
```

