

# (BINARY) RELATIONS

## OUTLINE:

- 1) Introduction to binary relations
- 2) Properties of relations
- 3) Combining relations
- 4) Representing relations
- 5) Equivalence relations
- 6) Partial orderings

# 1. INTRODUCTION TO BINARY RELATIONS

# Binary relations

- A **binary relation** from a set  $A$  to a set  $B$  is a subset  $R \subseteq A \times B$ .
  - EX:  $A = \{-2, -1, 0, 1, 2, 3\}$ ,  $B = \{0, 1, 2, 3\}$ ,  $R = \{(a, b) \in A \times B \mid a > b\} = \{(1, 0), (2, 0), (2, 1), (3, 0), (3, 1), (3, 2)\}$
  - EX:  $A = \{-2, -1, 0, 1, 2, 3\}$ ,  $B = \{0, 1, 2, 3\}$ ,  $R = \{(a, b) \in A \times B \mid b = a^2\} = \{(-1, 1), (0, 0), (1, 1)\}$
  - EX:  $A = \{-2, -1, 0, 1, 2, 3\}$ ,  $B = \{0, 1, 2, 3\}$ ,  $R = \{(a, b) \in A \times B \mid b < a^2 + 3\} = \{(-2, 0), (-2, 1), (-2, 2), (-2, 3), (-1, 0), (-1, 1), (-1, 2), (-1, 3), (0, 0), (0, 1), (0, 2), (1, 0), \dots\} = A \times B \setminus \{(0, 3)\}$
  - EX:  $A = \{-2, -1, 0, 1, 2, 3\}$ ,  $B = \{0, 1, 2, 3\}$ ,  $R = \{(a, b) \in A \times B \mid b > a^2 + 3\} = \emptyset$
  - EX:  $A = \{-2, -1, 0, 1, 2, 3\}$ ,  $B = \{0, 1, 2, 3\}$ ,  $R = \{(-1, 3), (-1, 0), (1, 2), (3, 0)\} \subseteq A \times B$  (kinda “random” relation, constructed specifying the couples one by one.)

# Binary relations

- A **binary relation** on a set  $A$  is a subset  $R \subseteq A \times A$ .
  - EX:  $A = \{0,1,2,3\}$ ,  $R = \{(a,b) \in A \times A \mid a \text{ is a multiple of } b\} = \{(0,0), (0,1), (0,2), (0,3), (1,1), (2,1), (2,2), (3,1), (3,3)\}$
  - EX:  $A = \mathbf{N}$ ,  $R = \{(a,b) \in \mathbf{N} \times \mathbf{N} \mid a = b\} = \{(0,0), (1,1), (2,2), \dots\}$   
(infinite set)
  - EX:  $A = \mathbf{N}$ ,  $R = \{(a,b) \in \mathbf{N} \times \mathbf{N} \mid a+b \text{ is even}\} = \{(0,0), (0,2), (0,4), \dots, (1,1), (1,3), \dots\}$  (infinite set)
  - EX:  $A = \mathbf{N}$ ,  $R = \{(a,b) \in \mathbf{N} \times \mathbf{N} \mid a \leq b\} = \{(0,0), (0,1), (0,2), \dots, (1,1), (1,2), \dots\}$  (infinite set)

# Notation for binary relations

- If  $R \subseteq A \times B$  is a binary relation, there are several ways to denote its elements:
  - $(a,b) \in R$  (set-theoretic notation);
  - $R(a,b)$  (logical notation, using the predicate corresponding to the set  $R$ ); EX:  $\text{Equal}(a,b)$ ;
  - $aRb$  (logical infix notation); EX:  $a=b$ ;

# Binary relations on a finite set

- If  $A$  is a finite set with  $|A| = n$ , then how many distinct binary relations are there on  $A \times A$ ?

# Binary relations on a finite set

- If  $A$  is a finite set, then how many distinct binary relations are there on  $A \times A$ ?
- Theorem 1: a set with  $m$  elements has  $2^m$  subsets (prove this by induction on  $m$ )
- Theorem 2: for any finite sets  $S$  and  $T$ ,  $|S \times T| = |S| \cdot |T|$  (prove this using the definition of cartesian product)
- The binary relations on  $A$  are the subsets of  $A \times A$ . Since by Theorem 1  $|A \times A| = |A|^2$ , by Theorem 2 the number of distinct binary relations on  $A$  is

$$2^{|A|^2}$$


## 2. PROPERTIES OF RELATIONS



# Reflexivity

- A relation  $R \subseteq A \times A$  is **reflexive** if  $\forall a (a \in A \rightarrow (a, a) \in R)$
- EX: on the integers  $\mathbf{Z}$ , which of the following relations are reflexive?
  - $R = \{(x, y) \in \mathbf{Z} \mid x = y\}$  ✓
  - $R = \{(x, y) \in \mathbf{Z} \mid x \leq y\}$  ✓ *for any  $a \in \mathbf{Z}$ ,  $a \leq a$ , so  $(a, a) \in R$ .*
  - $R = \{(x, y) \in \mathbf{Z} \mid |x| = |y|\}$  ✓
  - $R = \{(x, y) \in \mathbf{Z} \mid x < y\}$  ✗
  - $R = \mathbf{Z} \times \mathbf{Z}$  ✓
  - $R = \emptyset$  ✗  *$\Rightarrow$  it does not contain the shape of any  $(a, a)$ .*
  - $R = \{(x, y) \in \mathbf{Z} \mid x \text{ is a multiple of } y\}$  ✗

# Irreflexivity

- A relation  $R \subseteq A \times A$  is **irreflexive** if  $\forall a (a \in A \rightarrow (a, a) \notin R)$
  - EX: on the integers  $\mathbf{Z}$ , which of the following relations are irreflexive?
    - $R = \{(x, y) \in \mathbf{Z} \mid x = y\}$  ✗
    - $R = \{(x, y) \in \mathbf{Z} \mid x \leq y\}$  ✗
    - $R = \{(x, y) \in \mathbf{Z} \mid x \neq y\}$  ✓
    - $R = \{(x, y) \in \mathbf{Z} \mid x < y\}$  ✓
    - $R = \mathbf{Z} \times \mathbf{Z}$  ✗
    - $R = \emptyset$  ✓  $\forall a (a \in \mathbf{Z} \rightarrow (a, a) \notin \emptyset)$
    - $R = \{(x, y) \in \mathbf{Z} \mid x \text{ is a multiple of } y\}$  ✓
- 

# Question

- Is the following relation on the integers reflexive or irreflexive? *Neither.*

$$R = \{(x,y) \in \mathbf{Z} \mid x \text{ and } y \text{ are coprime (i.e., } \gcd(x,y)=1)\}$$

$$\gcd(1,1) = 1 \Rightarrow R \text{ is not irreflexive}$$

$$\gcd(2,2) = 2 \Rightarrow R \text{ is not reflexive.}$$


# Question

- Is the following relation on the integers reflexive or irreflexive?

$$R = \{(x,y) \in \mathbf{Z} \mid x \text{ and } y \text{ are coprime (i.e., } \gcd(x,y)=1)\}$$

- Neither! In fact,
  - $\gcd(1,1) = 1$ , so  $(1,1) \in R$ , therefore  $R$  cannot be irreflexive
  - $\gcd(2,2) = 2$ , so  $(2,2) \notin R$ , therefore  $R$  cannot be reflexive

# Symmetry

- A relation  $R \subseteq A \times A$  is **symmetric** if  $\forall a, b \in A ((a, b) \in R \rightarrow (b, a) \in R)$
- EX: on the integers  $\mathbf{Z}$ , which of the following relations are reflexive?
  - $R = \{(x, y) \in \mathbf{Z} \mid x = y\}$  ✓
  - $R = \{(x, y) \in \mathbf{Z} \mid x \leq y\}$  ✗  $(0, 1)$  ✓  $(1, 0)$  ✗.
  - $R = \{(x, y) \in \mathbf{Z} \mid |x| = |y|\}$  ✓ 
  - $R = \{(x, y) \in \mathbf{Z} \mid x < y\}$  ✗
  - $R = \mathbf{Z} \times \mathbf{Z}$  ✓
  - $R = \emptyset$  ✗.
  - $R = \{(x, y) \in \mathbf{Z} \mid x \text{ is a multiple of } y\}$  ✗
  - $R = \{(x, y) \in \mathbf{Z} \mid x + y < 2\}$  ✓

# Antisymmetry

- A relation  $R \subseteq A \times A$  is **antisymmetric** if  $\forall a, b \in A ((a, b) \in R \wedge (b, a) \in R \rightarrow a = b)$  that is, the only way for both  $(a, b)$  and  $(b, a)$  to belong to  $R$  is if  $a = b$

- EX: on the integers  $\mathbf{Z}$ , which of the following relations are antisymmetric?

*no symm case  
except diagonal*

- $R = \{(x, y) \in \mathbf{Z} \mid x = y\}$  ✓
- $R = \{(x, y) \in \mathbf{Z} \mid x \leq y\}$   *$a \leq b \wedge b \leq a \rightarrow a = b$*
- $R = \{(x, y) \in \mathbf{Z} \mid |x| = |y|\}$  ✗  *$|a| = |b| \wedge |b| = |a| \rightarrow a = b \vee a = -b$*
- $R = \{(x, y) \in \mathbf{Z} \mid x < y\}$   *$a < b \wedge b < a \rightarrow ?$*  ✗
- $R = \mathbf{Z} \times \mathbf{Z}$  ✗
- $R = \emptyset$
- $R = \{(x, y) \in \mathbf{Z} \mid x \text{ is a multiple of } y\}$  ✓
- $R = \{(x, y) \in \mathbf{Z} \mid x + y < 2\}$  ✗

# Asymmetry

no symm in all cases .



- A relation  $R \subseteq A \times A$  is **asymmetric** if  $\forall a, b \in A ((a, b) \in R \rightarrow (b, a) \notin R)$
- EX: on the integers  $\mathbf{Z}$ , which of the following relations are asymmetric?
  - $R = \{(x, y) \in \mathbf{Z} \mid x = y\}$  ✗
  - $R = \{(x, y) \in \mathbf{Z} \mid x \leq y\}$
  - $R = \{(x, y) \in \mathbf{Z} \mid |x| = |y|\}$
  - $R = \{(x, y) \in \mathbf{Z} \mid x < y\}$
  - $R = \mathbf{Z} \times \mathbf{Z}$
  - $R = \emptyset$
  - $R = \{(x, y) \in \mathbf{Z} \mid x \text{ is a multiple of } y\}$
  - $R = \{(x, y) \in \mathbf{Z} \mid x + y < 2\}$

# Asymmetric vs antisymmetric

- A relation  $R \subseteq A \times A$  is **asymmetric** iff it is both **antisymmetric and irreflexive** [homework: prove this]
- EX: on the integers  $\mathbf{Z}$ ,
  - $R = \{(x,y) \in \mathbf{Z} \mid x \leq y\}$  is antisymmetric but not irreflexive, hence not asymmetric
  - $R = \{(x,y) \in \mathbf{Z} \mid x < y\}$  is antisymmetric and also irreflexive, hence asymmetric



# Transitivity

- A relation  $R \subseteq A \times A$  is **transitive** if  $\forall a, b, c \in A ((a, b) \in R \wedge (b, c) \in R \rightarrow (a, c) \in R)$
- EX: on the integers  $\mathbf{Z}$ , which of the following relations are transitive?
  - $R = \{(x, y) \in \mathbf{Z} \mid x = y\}$
  - $R = \{(x, y) \in \mathbf{Z} \mid x \leq y\}$
  - $R = \{(x, y) \in \mathbf{Z} \mid |x| = |y|\}$
  - $R = \{(x, y) \in \mathbf{Z} \mid x < y\}$
  - $R = \mathbf{Z} \times \mathbf{Z}$
  - $R = \emptyset$
  - $R = \{(x, y) \in \mathbf{Z} \mid x \text{ is a multiple of } y\}$
  - $R = \{(x, y) \in \mathbf{Z} \mid x + y < 2\}$

# 3. COMBINING RELATIONS

# Set-theoretic operations

- Relations are sets, therefore they can be combined using the set operations  $\cap, \cup, ^c, \setminus$
- EX: on  $S = \{0,1,2,3\}$ , let  $R_1 = \{(0,0),(1,1),(2,1),(2,2),(3,1),(3,3)\}$  and  $R_2 = \{(1,1),(1,2),(1,3)\}$ . Then
  - $R_1 \cup R_2 = \{(0,0),(1,1),(1,2),(1,3),(2,1),(2,2),(3,1),(3,3)\}$
  - $R_1 \cap R_2 = \{(1,1)\}$
  - $R_1^c = S \times S \setminus R_1 = \{(0,1),(0,2),(0,3),(1,0),(1,2),(1,3),(2,0),(2,3),(3,0),(3,2)\}$   
note that the universe is the cartesian product of the set on which the relation is defined
  - $R_2 \setminus R_1 = \{(1,2),(1,3)\}$

# Inverse relation

- The inverse of a relation  $R \subseteq A \times B$  is the relation  $R^{-1} = \{(b,a) \in B \times A \mid (a,b) \in R\} \subseteq B \times A$ 
  - EX: Let  $A = \{a,b,c\}$  and  $B = \{0,1\}$  and let  $R = \{(a,0), (b,1), (c,1)\} \subseteq A \times B$  be a binary relation from  $A$  to  $B$ . Then  $R^{-1} = \{(0,a), (1,b), (1,c)\} \subseteq B \times A$
  - EX: Let  $R = \{(0,1), (1,1), (1,2), (1,3)\}$  be a binary relation on  $S = \{0,1,2,3\}$ . Then  $R^{-1} = \{(1,0), (1,1), (2,1), (3,1)\} \subseteq S \times S$

# Composition of relations

- The composition of a relation  $R_2 \subseteq B \times C$  with a relation  $R_1 \subseteq A \times B$  is the relation  $R_2 \circ R_1 \subseteq A \times C$  defined as

$$R_2 \circ R_1 = \{(a, c) \in A \times C \mid \exists b \in B ((a, b) \in R_1 \wedge (b, c) \in R_2)\}$$

- EX: Let  $A = \{0, 1, 2\}$ ,  $B = \{m, n, o, p\}$ ,  $C = \{w, x, y\}$ .  
Let  $R_1 = \{(1, p), (2, m)\} \subseteq A \times B$ ,  $R_2 = \{(m, x), (m, y), (o, w)\} \subseteq B \times C$ .  
Then  $R_2 \circ R_1 = \{(2, x), (2, y)\}$

# powers of relations

- A binary relation  $R \subseteq S \times S$  can be composed with itself
- EX: If  $S$  is the set of humans and  $R = \{(x,y) \in S \times S \mid x \text{ is a child of } y\}$ , then
$$R^2 = R \circ R = \{(x,y) \in S \times S \mid x \text{ is a grandchild of } y\},$$
$$R^3 = R \circ R \circ R = \{(x,y) \in S \times S \mid x \text{ is a great-grandchild of } y\},$$
and also  $R^{-1} = \{(x,y) \in S \times S \mid y \text{ is a child of } x\} = \{(x,y) \in S \times S \mid x \text{ is a parent of } y\}$ 
$$R^{-2} = (R^{-1})^2 = R^{-1} \circ R^{-1} = (R^2)^{-1} = \{(x,y) \in S \times S \mid x \text{ is a grandparent of } y\}$$

## 4. REPRESENTING RELATIONS

# Representation via matrices

- A relation between finite sets can be represented using a matrix of 0s and 1s:
- If  $R \subseteq A \times B$  with  $A = \{a_1, \dots, a_n\}$  and  $B = \{b_1, \dots, b_k\}$ , then the matrix of  $R$  is the  $n \times k$  matrix  $M_R = [m_{ij}]$  with
$$m_{ij} = 1 \text{ if } (a_i, b_j) \in R$$
$$m_{ij} = 0 \text{ if } (a_i, b_j) \notin R$$
- Note that the matrix depends on the choice of an ordering of the elements of  $A$  and an ordering of the elements of  $B$ . Any ordering is acceptable, but **when  $A = B$  we use the same ordering.**



# Example

- Let  $A = \{a, b, c\}$  and  $B = \{\text{'vowel'}, \text{'consonant'}\}$
- Let  $R = \{(a, \text{'vowel'}), (b, \text{'consonant'}), (c, \text{'consonant'})\}$

- The matrix of  $R$  is

$$M_R = \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 0 & 1 \end{bmatrix}$$

- The ordering  $A = \{b, a, c\}$ ,  $B = \{\text{'vowel'}, \text{'consonant'}\}$  would produce a different matrix:

$$M_R = \begin{bmatrix} 0 & 1 \\ 1 & 0 \\ 0 & 1 \end{bmatrix}$$

# Example

- Let  $A = \{a,b,c\}$  and  $B = \{0,1,2,3\}$
- Let  $R$  be the relation on  $A \times B$  represented by the matrix

$$M_R = \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Describe  $R$  with the roster method

# Example

- Let  $A = \{a,b,c\}$  and  $B = \{0,1,2,3\}$
- Let  $R$  be the relation on  $A \times B$  represented by the matrix

$$M_R = \begin{matrix} & \overset{A}{\begin{matrix} 1 & 0 & 1 & 0 \end{matrix}} \\ \underset{B}{\begin{matrix} 0 \\ 1 \\ 2 \end{matrix}} & \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

$A \times B$   $B$

Describe  $R$  with the roster method

$$R = \{(a,0), (a,2), (b,1), (b,2)\}$$

# Matrices and relation properties

- Remember that for binary relations  $R \subseteq S \times S$  over a set  $S$  we use the same ordering on the 2 copies of  $S$
- A relation  $R \subseteq S \times S$  is reflexive iff all the elements on the main diagonal of  $M_R$  are 1
- A relation  $R \subseteq S \times S$  is irreflexive iff all the elements on the main diagonal of  $M_R$  are 0
- A relation  $R \subseteq S \times S$  is symmetric iff  $M_R$  is a symmetric matrix (i.e.,  $m_{ij} = m_{ji}$  for all indices  $i$  and  $j$ )
- A relation is antisymmetric iff, for any indices  $i \neq j$ ,  $(m_{ij} = 0 \vee m_{ji} = 0)$

# Example

- Let  $S = \{0,1,2,3\}$  and  $R \subseteq S \times S$  be defined by the matrix

$$M_R = \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{bmatrix}$$

- What properties of  $R$  can we deduce from  $M_R$ ?

# Example

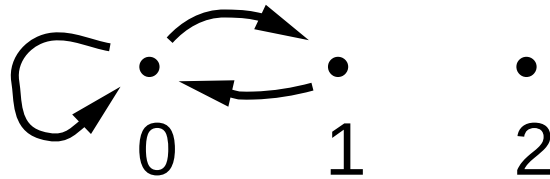
- Let  $S = \{0,1,2,3\}$  and  $R \subseteq S \times S$  be defined by the matrix

$$M_R = \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{bmatrix}$$

- $m_{33} = 0$ , so  $R$  is not reflexive
- $m_{11} = 1$ , so  $R$  is not irreflexive
- $m_{13} = 1$  and  $m_{31} = 0$ , so  $R$  is not symmetric
- $(m_{13} = 1 \text{ and } m_{31} = 0)$ ,  $(m_{23} = 1 \text{ and } m_{32} = 0)$ ,  $(m_{42} = 1 \text{ and } m_{24} = 0)$ , so  $R$  is antisymmetric (whenever we have a 1 off the main diagonal, in the symmetric position we have a 0)

# Representation via graphs

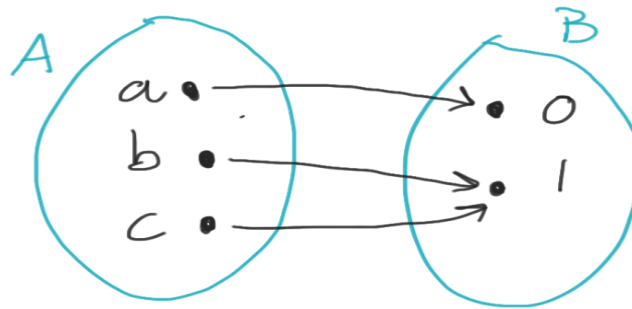
- A **directed graph** consists of a set  $V$  of **vertices** (aka **nodes** or **points**) and a set  $E \subseteq V \times V$  of **edges**. If  $(a,b) \in E$ , then  $a$  is the **initial vertex** and  $b$  is the **terminal vertex** of the edge  $(a,b)$ . An edge of the form  $(a,a)$  is a **loop**. Edges are drawn as arrows from their initial to their terminal vertex.
- EX: the graph  $G=(V,E)$  with  $V = \{0,1,2\}$  and  $E = \{(0,0),(0,1), (1,0)\}$  is



- Graphs will be studied in detail in next episodes

# Representation via graphs

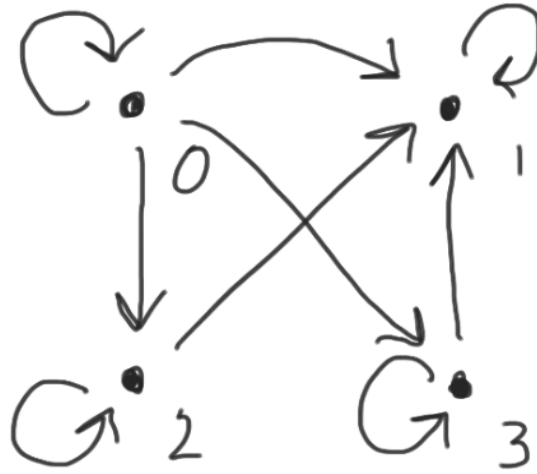
- A relation  $R \subseteq A \times B$  can be represented as a graph with vertex set  $V = A \cup B$  and edge set  $R$ . If  $A \neq B$ , then the elements of  $A$  are kept “separate” from the elements of  $B$  (usually, Venn diagrams for  $A$  and  $B$  are also included).
- EX: if  $A = \{a, b, c\}$  and  $B = \{0, 1\}$ , the relation  $R = \{(a, 0), (b, 1), (c, 1)\} \subseteq A \times B$  can be represented by the graph





# Representation via graphs

- EX: if  $A = \{0,1,2,3\}$ , the relation  $R = \{(a,b) \in A \times A \mid \text{a is a multiple of b}\}$  can be represented by the graph



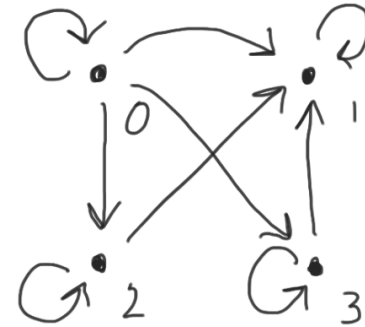
$(0,0)$   $(1,1)$   
 $(2,1)$   $(3,1)$   
 $(0,1)$   $(2,0)$   
 $(3,0)$ , -

# Graphs and relation properties

- A relation is reflexive iff all vertices have a loop
- A relation is irreflexive iff no vertex has a loop
- A relation is symmetric iff whenever  $(x,y)$  is an edge, then so is  $(y,x)$
- A relation is antisymmetric iff whenever  $(x,y)$  is an edge with  $x \neq y$ , then  $(y,x)$  is not an edge
- A relation is transitive iff whenever  $(x,y)$  and  $(y,z)$  are edges, then so is  $(x,z)$

# Previous example, revisited

- EX: if  $A = \{0,1,2,3\}$ , the relation  $R = \{(a,b) \in A \times A \mid a \text{ is a multiple of } b\}$  can be represented by the graph
- Each vertex has a loop, so  $R$  is reflexive
- $(0,1)$  is an edge, but  $(1,0)$  is not, so  $R$  is not symmetric
- Whenever  $(a,b)$  is an edge with  $a \neq b$ , then  $(b,a)$  is not an edge (check  $(0,1)$  vs  $(1,0)$ ,  $(0,2)$  vs  $(2,0)$ ,  $(0,3)$  vs  $(3,0)$ ,  $(2,1)$  vs  $(1,2)$ ,  $(3,1)$  vs  $(1,3)$ ), so  $R$  is antisymmetric
- Whenever  $(a,b)$  and  $(b,c)$  are edges, so is  $(a,c)$  (e.g.,  $(0,2)$ ,  $(2,1)$  and  $(0,1)$ ), so  $R$  is transitive



# Another example (mind trivial cases!)

- EX: if  $A = \{a, b, c\}$  and  $R$  is the relation represented by the graph



- $R$  is neither reflexive nor irreflexive (only  $a$  has a loop)
- $R$  is **vacuously** symmetric (there are no edges  $(x, y)$  with  $x \neq y$ )
- $R$  is **vacuously** antisymmetric (same reason)
- $R$  is **vacuously** transitive (same reason)

Remember vacuous conditionals? (Conditionals with false premise are true)

# 5. EQUIVALENCE RELATIONS

# Equivalence relations

- A relation  $R \subseteq A \times A$  (same set) is called an **equivalence relation** if it is reflexive, symmetric, and transitive.
- If  $R$  is an equivalence relation, two elements  $a$  and  $b$  such that  $aRb$  are called equivalent. In this case, the notation  $a \sim b$  is often used.
- EX: For any set  $A$ , the **identity relation**  $I_A = \{(a,b) \in A \times A \mid a=b\} = \{(a,a) \mid a \in A\}$  is an equivalence relation. In fact, it is
  - Reflexive, because any  $a \in A$  is equal to itself ( $a=a$ )
  - Symmetric, because if  $a=b$  then  $b=a$
  - Transitive, because if  $a=b$  and  $b=c$ , then  $a=c$
- In fact, the identity is the archetypical equivalence relation: the definition of equivalence relation is modelled on the properties of the identity relation

# Equivalence classes

- If  $R \subseteq A \times A$  is an equivalence relation, for any fixed  $x \in A$ , the subset  $\{a \in A \mid a \sim x\} \subseteq A$  of the elements in relation with  $x$  is called the **equivalence class** of  $x$ , and denoted  $[x]_R$ , or just  $[x]$  if  $R$  is clear from the context.
- CAUTION!  $[x] = [y]$  for any  $y$  such that  $y \sim x$ , thus in general  $[x] = [y]$  does not imply  $x = y$ .
- When we write an equivalence class as  $[x]$ , we say that  $x$  is a **representative** of that class. Any element of a class can be used as representative.

# EX: Congruence modulo $m$

- Let  $m > 1$  be an integer. Remember that, for two integers  $a$  and  $b$ ,  $a \equiv b \pmod{m}$  means that  $a$  and  $b$  have the same remainder in the integer division by  $m$
- The relation  $\{(a,b) \in \mathbf{Z} \times \mathbf{Z} \mid a \equiv b \pmod{m}\}$  is an equivalence relation on the integers



# EX: Congruence modulo $m$

- Reflexivity: clearly for any  $a \in \mathbb{Z}$  ( $a \equiv a \pmod{m}$ )
- Symmetry: if  $a \equiv b \pmod{m}$  (that is,  $a$  and  $b$  have the same remainder when divided by  $m$ ), then  $b \equiv a \pmod{m}$  (that is,  $b$  and  $a$  have the same remainder when divided by  $m$ )
- Transitivity: if  $a \equiv b \pmod{m}$  (that is,  $a$  and  $b$  have the same remainder, say  $r$ , when divided by  $m$ ), and  $b \equiv c \pmod{m}$  (that is,  $b$  and  $c$  have the same remainder, which must be  $r$  again, when divided by  $m$ ), then  $a \equiv c \pmod{m}$  (that is,  $a$  and  $c$  have the same remainder, still  $r$ , when divided by  $m$ )

# EX: Congruence modulo $m$

- The equivalence class of an integer  $a$  modulo  $m$  is  $[a]_m = \{ \dots, a-3m, a-2m, a-m, a, a+m, a+2m, a+3m, \dots \}$
- The difference between consecutive elements in  $[a]_m$  is  $m$
- There are exactly  $m$  distinct equivalence classes modulo  $m$ :  $[0]_m, [1]_m, \dots, [m-1]_m$
- Of course, other choices of representatives are possible

# Concrete ex: Congruence modulo 3

- There are 3 equivalence classes modulo 3:
  - $[0]_3 = \{\dots, -9, -6, -3, 0, 3, 6, 9, \dots\}$
  - $[1]_3 = \{\dots, -8, -5, -2, 1, 4, 7, 10, \dots\}$
  - $[2]_3 = \{\dots, -7, -4, -1, 2, 5, 8, 11, \dots\}$
- Notice that the equivalence classes are mutually disjoint, nonempty, and their union is the whole  $\mathbb{Z}$ . This is a general property of equivalence relations.

# Equivalence relations and partitions

- A **partition** of a set  $S$  is a collection  $\{A_j \mid j \in J\}$  (where  $J$  is a set of indices) of subsets of  $S$  which are
  - **mutually disjoint** (for all  $j, k \in J$  with  $j \neq k$ ,  $A_j \cap A_k = \emptyset$ ),
  - **nonempty** (for all  $k \in J$ ,  $A_k \neq \emptyset$ ),
  - and **whose union is  $S$**  ( $\bigcup_{j \in J} A_j = S$ )
- If on a set  $S$  there is an equivalence relation, the equivalence classes form a partition of  $S$ .
- Viceversa, if a set  $S$  has a partition  $\{A_j \mid j \in J\}$ , then the relation  $R = \{(x, y) \in S \times S \mid x \text{ and } y \text{ belong to the same } A_k (k \in J)\}$  is an equivalence relation with the  $A_j (j \in J)$  as the equivalence classes.

## 6. PARTIAL ORDERINGS

# Partial orderings

- A relation  $R \subseteq A \times A$  (same set) is called a **partial ordering**, or **partial order**, if it is reflexive, antisymmetric, and transitive.
- A set  $S$  together with a partial ordering  $R$  is called a **partially ordered set**, or **poset**, and is denoted by  $(S, R)$ .
- EX: On  $\mathbf{Z}$ , the relation  $\leq$  (“less than or equal to”), i.e.  $\{(a, b) \in \mathbf{Z} \times \mathbf{Z} \mid a \leq b\}$  is a partial order. In fact, it is
  - Reflexive, because any  $a \in \mathbf{Z}$  is less than or equal to itself ( $a \leq a$ )
  - Antisymmetric, because if  $a \leq b$  and  $b \leq a$ , then  $a = b$
  - Transitive, because if  $a \leq b$  and  $b \leq c$ , then  $a \leq c$
- Therefore,  $(\mathbf{Z}, \leq)$  is a poset. The same reasoning works for the relation  $\geq$
- In fact,  $\leq$  (or  $\geq$ ) is the archetypical partial order: the definition of partial order is modelled on the properties of  $\leq$  (or  $\geq$ )

# Strict orderings

- A relation  $R \subseteq A \times A$  (same set) is called a **strict (partial) ordering** (or **order**) if it is asymmetric (or equivalently (irreflexive and antisymmetric)), and transitive.
- A set  $S$  together with a partial ordering  $R$  is called a **strict partially ordered set**, or **strict poset**, and is denoted by  $(S, R)$ .
- EX: On  $\mathbf{Z}$ , the relation  $<$  (“less than”), i.e.  $\{(a, b) \in \mathbf{Z} \times \mathbf{Z} \mid a < b\}$  is a partial order. In fact, it is
  - Asymmetric, because if  $a < b$ , then  $\neg(b < a)$
  - Transitive, because if  $a < b$  and  $b < c$ , then  $a < c$
- Therefore,  $(\mathbf{Z}, <)$  is a strict poset. The same reasoning works for the relation  $>$
- In fact,  $<$  (or  $>$ ) is the archetypical partial order: the definition of partial order is modelled on the properties of  $<$  (or  $>$ )

# Strict vs non-strict orderings

- Let  $A$  be a set. Recall the identity relation on  $A$ :  $I_A = \{(a,b) \in A \times A \mid a=b\}$
- (1) Given a (non-strict) partial order  $P \subseteq A \times A$ , there is an induced strict partial order  $Q \subseteq A \times A$ , defined by  $Q = P \setminus I_A = \{(a,b) \in P \mid \neg(a=b)\} = \{(a,b) \in P \mid a \neq b\}$ .
- (2) Viceversa, given a strict order  $R \subseteq A \times A$ , there is an induced partial order  $T \subseteq A \times A$ , defined by  $T = R \cup I_A = \{(a,b) \in A \times A \mid (a,b) \in R \vee a=b\}$
- Homework: prove the above points. That is, for (1), show that  $Q$  is irreflexive and transitive; for (2), show that  $T$  is reflexive, antisymmetric and transitive.



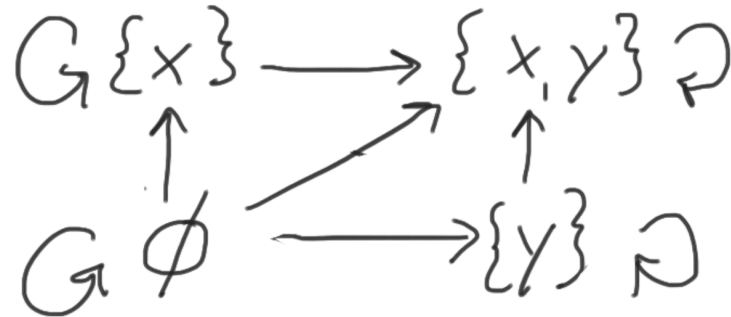
# EX: The power set poset

- Let  $A$  be a set. The power set  $P(A)$  together with the inclusion relation  $\subseteq$  is a poset. In fact, **by the definition of set inclusion**,
  - Reflexivity: every subset  $S$  of  $A$  is included in itself ( $S \subseteq S$ )
  - Antisymmetry: if 2 subsets  $S$  and  $T$  of  $A$  satisfy  $S \subseteq T$  and  $T \subseteq S$ , then  $S = T$  (this is our favourite technique to show a set equality)
  - Transitivity: if 3 subsets  $B, C, D \in P(A)$  satisfy  $B \subseteq C$  and  $C \subseteq D$ , then also  $B \subseteq D$
- EX: show that  $P(A)$  with the proper inclusion relation  $\subset$  is a strict poset

# Concrete ex: The power set of $\{x,y\}$

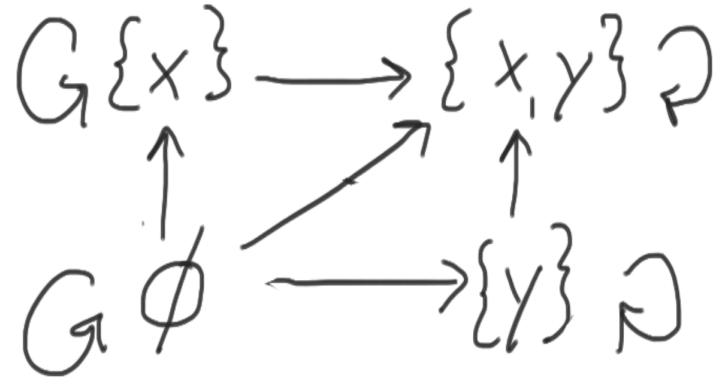
- Let  $A = \{x,y\}$
- Then  $P(A) = \{\emptyset, \{x\}, \{y\}, \{x,y\}\}$
- The inclusion relation is  $\{(\emptyset, \emptyset), (\emptyset, \{x\}), (\emptyset, \{y\}), (\emptyset, \{x,y\}), (\{x\}, \{x\}), (\{x\}, \{x,y\}), (\{y\}, \{y\}), (\{y\}, \{x,y\}), (\{x,y\}, \{x,y\})\}$ ...correct but not the clearest. In matrix and graph representation:

$$\begin{bmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$



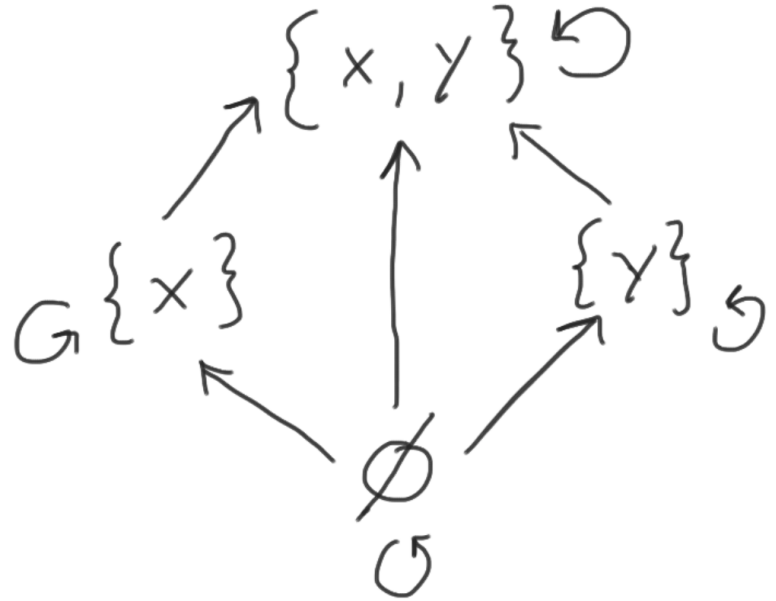
# Hasse diagrams

- Hasse diagrams are another type of graph representation **specific for partial orders**.  
Suppose you have a partial order  $R$  on a set  $S$  (we will use  $\subseteq$  on  $P(\{x,y\}) = \{\emptyset, \{x\}, \{y\}, \{x,y\}\}$  as an example).
- Start with a “normal” graph representation of  $R$



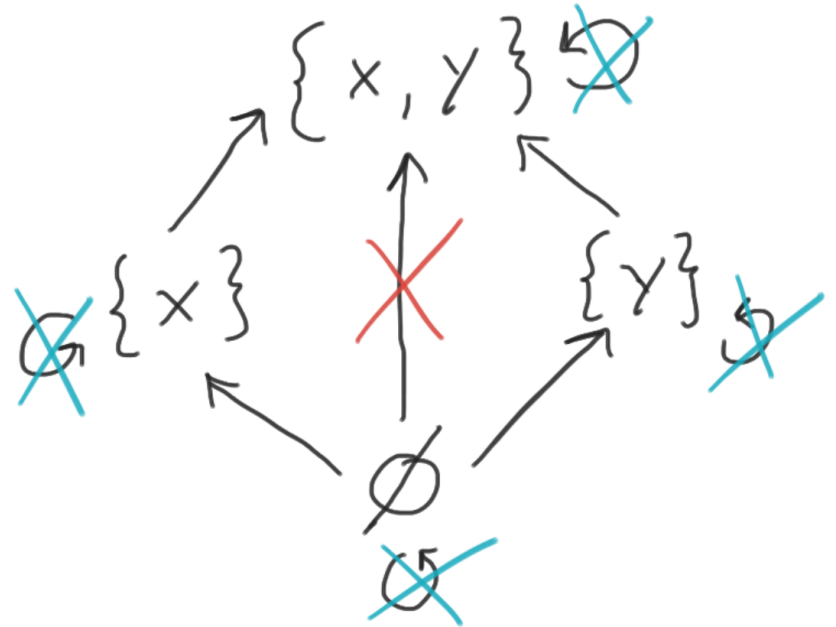
# Hasse diagrams

- Rearrange the vertices in such a way that, if vertices  $a$  and  $b$  satisfy  $aRb$ , then  $b$  is higher up on the page than  $a$  (if  $a$  and  $b$  are not related, their relative height can be whatever). In our case,  $\{x,y\}$  has to be at the top,  $\{x\}$  and  $\{y\}$  below it (and their relative height does not matter) and  $\emptyset$  must be at the bottom



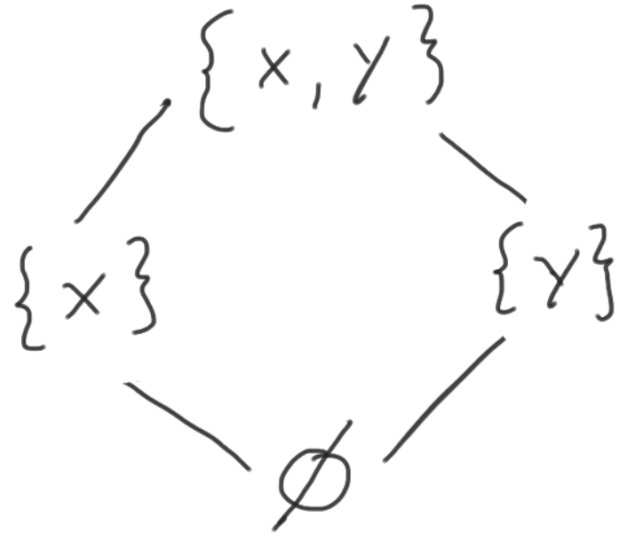
# Hasse diagrams

- Remove the **edges due to reflexivity** (i.e., the loops) and those **implied by transitivity**. In our case, the edge  $(\emptyset, \{x, y\})$  is implied by transitivity from the edges  $(\emptyset, \{x\})$  and  $(\{x\}, \{x, y\})$ , so we remove it



# Hasse diagrams

- Remove the arrow tips (the direction of edges is implied by the relative height of the vertices)
- That's your Hasse diagram



# Total orders

- Let  $(A, R)$  be a poset. Two elements  $a, b$  of  $A$  are said to be **comparable** if  $aRb$  or  $bRa$ . The elements are called **incomparable** if neither  $aRb$  nor  $bRa$ .
- A poset  $(A, R)$  in which all elements are comparable is said to be a **totally ordered set** (other names: **linearly ordered set**, **chain**) and  $R$  is called a **total** (or **linear**) **order**.
- A totally ordered set such that every nonempty subset has a minimum is called a **well-ordered set**.
  - EX:  $(P(\{x, y\}), \subseteq)$  is not a totally ordered set because  $\{x\}$  and  $\{y\}$  are incomparable.
  - EX:  $(\mathbb{Z}, \leq)$  is a totally ordered set, but not a well-ordered set ( $\mathbb{Z}$  itself has no minimum).
  - EX:  $(\mathbb{N}, \leq)$  is a well-ordered set.

# Hasse diagrams of total orders

- EX: consider the poset  $A = \{0,1,2,3,4\}$  with the total order  $\leq$ .
- Its Hasse diagram is a “line”.
- This is true for any totally ordered set.

