WEEK 3

ADDING RECORDS TO THE HARD DRIVE USING A HEAP ORGANIZATION

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STUDENT OBJECTIVES

- Upon completion of this video, you should be able to:
 - Explain how the records are added to the disk when using a heap organization
 - Given adding, modifying and deleting records, determine which operations are efficient and which operations are costly
 - Determine when a heap organization is appropriate
 - Given a number of records, record size and block size, figure out the average number of searches needed to find a record and the worst case scenario for searching for a given record

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CONSIDER HOW "EXPENSIVE" EACH FILE ORGANIZATION WOULD BE:

- A heap (or a pile or unordered file) → each time you get a new record, just add it to then end.
- Consider inserting the following items (25, 12, 89, 64, 1, 13, 19) into the disk below:

25 | 12 | 89 | 64 | 1 | 13 | 19 |

- Simplest type of organization
- QUESTION: Using a heap, of the 3 operations (insert, modify, delete), which ones would be cheap (fast and easy), which ones would be expensive?

deleting is the most expensive: no not have a gap, you have so more all the John newsods one position forwards

Good in the following situations:

- Data is being bulk loaded into the relation.
- Relation is only a few pages long (searches don't have lots of data to go through)
- When every tuple has to be retrieved every time the relation is used
- When the relation has an additional access structure such as an indexed key

Bad in the following situations:

- Search (linear search, on average b/2 where b is the number of blocks that must be brought into main memory)
- Deletions, same as searches and it leaves open blocks (wasted storage) or require reorganization occasionally to clean up.
- Sorting

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EXAMPLE

QUESTION: Find the average search time to find a record if you use a heap organization for the following scenario:

- r = 100,000 records stored on a disk with block size B = 2048 bytes.
- Each record (R) is a fixed size of R = 500 bytes.
- Blocking Factor (bfr) = 2048/500 = 4 records per block (fill in the blank)
- # of blocks needed is 100,000/4 = 25,000 blocks
- Linear Search: on average b/2 = 25,000 / 2 = 12,500 block accesses

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100 revords too bytes long

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