

Sept 20

COMPSCI 3331

Fall 2022

What's next?

- ▶ Current location: up to Lecture 3, Part 1
- ▶ Tomorrow: more on Lecture 3.
- ▶ Assignment 1: out by Sept 27 (at the latest), due Oct 11.
- ▶ Quiz 1: Sept 28 (including Lecture 3 material).
- ▶ Updates - this file available before class, “unanswered questions” file.

Software tools

- ▶ For assignments, you can use a software tool to draw automata.
- ▶ Two that I know of:
 - ▶ JFLAP: <https://www.jflap.org/> - self-contained software.
 - ▶ FAdo: <https://fado.dcc.fc.up.pt/> - py library
- ▶ Not necessary for the course.
- ▶ Fine to use for assignments.

Reversal

- Inductive definition.

- Proof of $(xy)^R = y^R x^R$

$\varepsilon^R \subseteq \varepsilon \leftarrow \begin{matrix} \text{empty} \\ \text{string} \end{matrix}$

$$(xa)^R = ax^R \quad \forall x \in \varepsilon^R \\ \forall a \in \Sigma$$

$$(abcd)^R = dcba$$

induction on length of y :

base case: $|y| = 0 \Rightarrow y = \varepsilon$

$$(xy)^R = (x\varepsilon)^R = x^R = \varepsilon \quad x^R = y^R x^R$$

inductive case:

assume the statement holds for all $y \in \Sigma^*$ with $|y| = n$.

let y be a word with $|y| = n+1$

$$\Rightarrow y = wa \text{ for } w \in \Sigma^*, a \in \Sigma \quad \text{definition}$$

$$(xy)^R = (x(wa))^R = ((xw)a)^R = a(xw)^R$$

\uparrow \uparrow
 letter words.

$\frac{w}{\uparrow} \frac{a}{\uparrow}$
subword

the last letter definition.

Since a is a single letter, no need of reverse itself

Σ^* : a set of languages.

Language Examples

- ▶ $L_1 = \{w \in \{a, b\}^* : |w|_a > |w|_b\}$.
1. w is made up from a and b .
2. w has more a than b .
- ▶ $L_2 = \{x \# y : x, y \in \{0, 1, 2, \dots, 9\}^* \text{ and } x^2 = y\}$
 \uparrow
 just a symbol,
 which has no meaning
 aka. placeholder. e.g. $0\#0, 1\#1, 2\#4$
 Since the order matter, so $4\#2$ is not
 in the language.

Language Identities

- ▶ $L_1(L_2 \cup L_3) = L_1L_2 \cup L_1L_3$? ✓
- ▶ $(L_1L_2)^R = L_2^R L_1^R$? ✓
- ▶ $\overline{L^*} = (\overline{L})^*$? ✓

$$L_1(L_2 \cup L_3) \subseteq L_1L_2 \cup L_1L_3$$

$$\text{let } x \in L_1(L_2 \cup L_3)$$

$$\exists y \in L_1, z \in L_2 \cup L_3 \text{ so}$$

$$x = yz$$

$$L = \emptyset \quad L^* = \{\epsilon\} \quad (\overline{L^*}) = \overline{\{\epsilon\}} = \Sigma^* \setminus \{\epsilon\}$$

$$(\overline{L})^* = \overline{L^*} \leftarrow$$

\overline{L} : complement of L .

L^* : words not in L

$$\Sigma = \{a, b\} \quad L = \{a^i; i \geq 0\}$$

$$L^* = \bigcup_{i=0}^{\infty} L^i = L \text{ (in this case)}$$

$$\overline{L^*} = \{w \in \{a, b\}^*, |w|_b > 0\} = (\overline{L})^*$$

a set of words of $\{a, b\}^*$ which has at least one b in each word.

DFAs

- ▶ $L_1 = \{x \in \{a, b, c\}^* : |x|_b \equiv 0 \pmod{3}\}.$
- ▶ $L_2 = \{(aabc^i)^j : i \geq 0, j \geq 0\}$
- ▶ $L_3 = \{(abbd^i)^j : i \geq 0, j \geq 0\}$
- ▶ $L_4 = L_2 \cup L_3$