Review Session

COMPSCI 3331

Fall 2022

1. Let G be the CFG defined by the following set of productions.

$$S
ightarrow bbAaA \mid SSaa \mid aa \mid ABC$$

 $A
ightarrow Ab \mid Ac \mid CC$
 $B
ightarrow BA \mid bb \mid Dd$
 $C
ightarrow DA \mid \varepsilon$
 $D
ightarrow a$

Give an equivalent grammar to G that has no ε -productions.

2. Let G be the CFG defined by the following set of productions.

$$S
ightarrow SbaB \mid aa \mid ABC$$

 $A
ightarrow Ba \mid DaDd$
 $B
ightarrow BA \mid ca \mid Dd$
 $C
ightarrow DA \mid \varepsilon$
 $D
ightarrow a$

Convert the grammar to CNF.

- 3. Let $G = (V, \Sigma, P, S)$ be a CFG in CNF. Give an $O(n^3)$ algorithm for the following problem:
 - ▶ Input: A word w and a nonterminal $A \in V$.
 - Output: the value $n_A = \max\{|u| : u \text{ is a suffix of } w \text{ and } A \Rightarrow^* u\}.$

That is, n_A is the length of the longest suffix of w that is generated by A in the grammar.

- 4. Construct PDAs for the following languages:
- (a) $L = \{a^n b^m xy : x, y \in \{0, 1, 2\}, x \equiv n \pmod{3} \text{ and } y \equiv m \pmod{3} \}.$
- (b) $L = \{ w \# x : w, x \in \{a, b\}^*, |w|_a = |x|_a, |w|_b \equiv |x|_b \pmod{3} \}.$

Be sure to indicate what the starting stack symbol is for your PDA and how your PDA accepts words.

5. Let *C* be a fixed integer. Extend the language from Assignment 3 as follows:

$$L_C = \{x \# 1^n : n \ge 0, x \in \{a, b\}^* \text{ and } n - C \le |x|_a \le n + C\}$$

Give a context-free grammar for L_C . The productions in your grammar will depend on the value of C. Describe them using a uniform notation (e.g., by using consistently named variables or consistently defined productions, for instance).

6. Consider the following modified language from Assignment 3:

$$L = \{u \# v : u, v \in \{0, 1\}^* \text{ and } bin(v^R) = bin(u) + 2\}$$

Give a PDA that accepts L.

7. A CFG *G* is in Griebach Normal Form (GNF) if every production has the form

$$A \rightarrow aB_1B_2 \cdots B_n$$

for some letter a and nonterminals B_1, B_2, \cdots, B_n (where $n \ge 0$). Any grammar (that does not derive ε) can be converted to GNF. Given this fact, show that for any CFG L that does not include ε , you can construct a PDA M that accepts L in the following additional conditions:

- The PDA accepts by empty stack.
- ▶ The PDA M does not have any ε -transitions. That is, there are no rules of the form $\delta(q, \varepsilon, \gamma) = \{(q', \beta), \dots\}$ for any stack symbol γ .

- 8. Prove that the following languages are not context-free:
 - (a) $L = \{a^p : p \text{ is a prime number}\}.$
 - (b) $L = \{a^n b^{n^3} : n \ge 0\}.$
 - (c) $L = \{ w \in \{a, b\}^* : |w|_b = 2^{|w|_a} \}.$

- 9. For each of the languages in the previous question, give an informal description of a multi-tape TM that recognizes the language.
 - (a) $L = \{a^p : p \text{ is a prime number}\}.$
 - (b) $L = \{a^n b^{n^3} : n \ge 0\}.$
 - (c) $L = \{ w \in \{a, b\}^* : |w|_b = 2^{|w|_a} \}.$

10. Show that the following language is r.e.:

$$L = \{e(M_1) \# e(M_2) : L(M_1) \cap L(M_2) \neq \emptyset\}$$

11. Show that the following problem is undecidable by reduction: Given a TM M, is L(M) a finite language?

12. Show that the following language is decidable:

 $L_{ND} = \{ e(M) : M \text{ is a nondeterministic TM } \}.$

13. Show that the following problem is either decidable or undecidable: Given a CFG G is L(G) infinite? (Hint: review the proof of the pumping lemma for CFLs.)