

Last day of class: Dec 8

§ 6.4. Strong Induction

Strategy: To prove a goal: $\forall n \in \mathbb{N}, \text{prove: } \forall k \in \mathbb{N} (k < n, P(k)) \rightarrow P(n)$.

$$n=0: (\forall k < 0, P(k)) \rightarrow P(0) \Rightarrow \text{True for } n=0.$$

Empty \Rightarrow True True

$$n=1: (\forall k < 1, P(k)) \rightarrow P(1) \Rightarrow (\underbrace{(k=0, P(k))}_{\text{true}}) \rightarrow \underbrace{P(1)}_{\text{True}}.$$

$$n=2: (\forall k < 2, P(k)) \rightarrow P(2) \Rightarrow (\underbrace{(k \in \{0, 1\}, P(k))}_{\text{true}}) \rightarrow P(2).$$

\vdots

\vdots

Thm 6.4.1 (Division Algorithm):

$$\forall m \in \mathbb{Z}^+, \forall n \in \mathbb{N}, \exists q, r \in \mathbb{N} \text{ such that } n = qm + r, r \in [0, m).$$

Proof: Let $m \in \mathbb{Z}^+$, We use strong induction on n , Let $n \in \mathbb{N}$, and assume for all $k < n$, there exist $q, r \in \mathbb{N}$, $k = qm + r$ and $r < m$.

Case 1: $n < m$: then $n = 0 \cdot m + n$ so $q = 0, r = n < m$.

Case 2: $n \geq m$: by inductive hypothesis, $q', r', n' \in \mathbb{N}$ such that

$$n - m = q'm + r', r' < m. \text{ Then } n = (q' + 1)m + r',$$

$$\text{so take } q = q' + 1 \text{ and } r = r' \quad \square.$$

Recall: $n > 1$ is prime $\iff \neg \exists a \in \mathbb{N}, \exists b \in \mathbb{N} (n = ab \wedge a < n \wedge b < n)$

Thm 6.4.2: $\forall n > 1$ is either a prime or a product of prime.

Proof: We use strong induction on n .

Let $n > 1$: Suppose that for all k with $1 < k < n$, k is either prime or a product of primes.

If n is prime, we're done.

If n is not prime: there exist $a, b \in \mathbb{N}$ such that $ab = n$, $a < n$, $b < n$.

Since $a \in (1, n)$, $b \in (1, n)$. So by the induction hypothesis.

a, b are either or a product of primes. In any case,
 $n \mid ab$ is a product of primes \square .

Thm 6.4.4: (Well-ordered principle): Every non-empty set of natural number has a smallest element.

Proof: Let $S \subseteq \mathbb{N}$, We'll prove by contrapositive.

S does not have a smallest element. We'll show that $S = \emptyset$.

i.e. $\forall n \in \mathbb{N} (n \notin S)$. We'll use strong induction.

Assume $\forall k < n (k \notin S)$. The contrapositive is if $k \in S$, $n \leq k$.

Therefore, if $n \in S$, it is the smallest element in S .

But in the assumption, S is empty. $n \notin S$.

Ex: Every $q \in \mathbb{Q}^+$ can be written $\frac{m}{n}$ such that $m, n \in \mathbb{Z}^+$ having no common divisor > 1 .

Proof: Let $q \in \mathbb{Q}^+$, Let $S = \{m \in \mathbb{Z}^+ \mid \exists n \in \mathbb{Z}^+ q = \frac{m}{n}\}$.

This is non-empty, so it has a smallest element m ,

then $q = \frac{m}{n}$ for some $n \in \mathbb{Z}^+$,

If $\frac{m}{n}$ is not in lowest term, then $\exists d > 1$ such that $d \mid m$,

$d \mid n$, But then $q = \frac{m/d}{n/d}$, $m/d \in \mathbb{Z}^+$, $n/d \in \mathbb{Z}^+$, so $m/d \in S$.

but $m/d < m$, it is a contradiction.