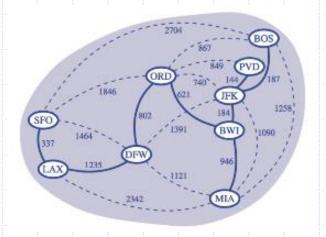
Presentation for use with the textbook Data Structures and Algorithms in Java, 6th edition, by M. T. Goodrich, R. Tamassia, and M. H. Goldwasser, Wiley, 2014. With notes from R. Solis-Oba

Minimum Spanning Trees



Minimum Spanning Trees

Spanning subgraph

Subgraph of a graph G
 containing all the vertices of G

Spanning tree

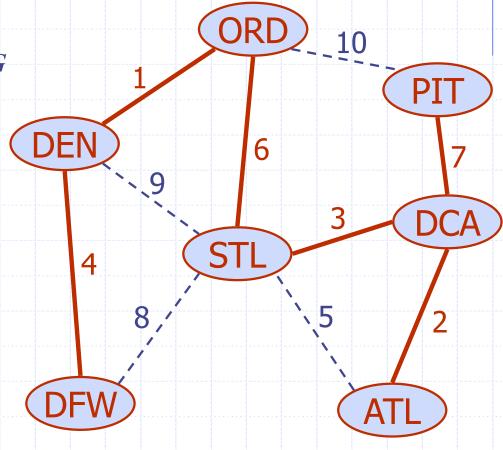
Spanning subgraph that is itself a tree

Minimum spanning tree (MST)

 Spanning tree of a weighted graph with minimum total edge weight

Applications

- Communications networks
- Transportation networks



Prim's Algorithm

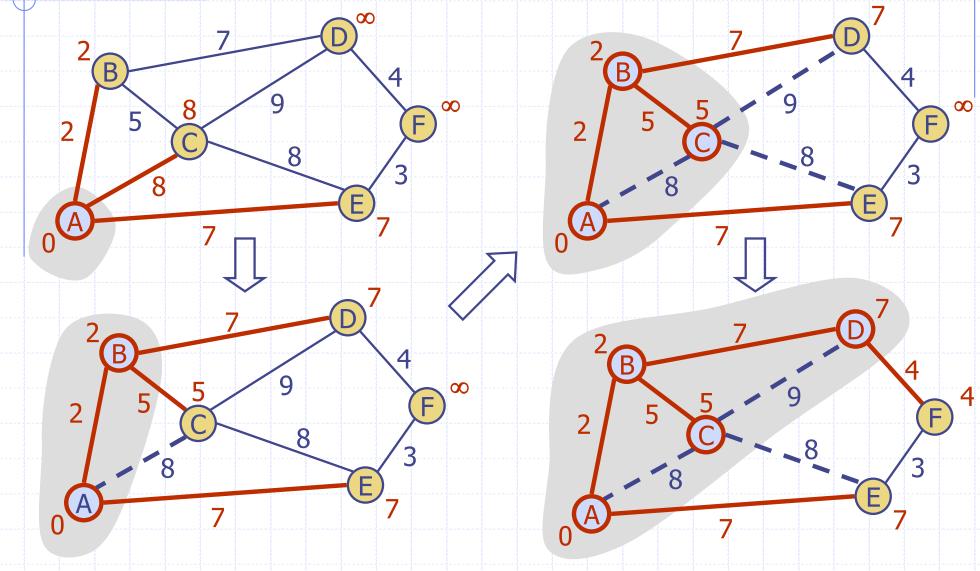
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```
Algorithm Prim (G,s)
In: weighted connected graph G and vertex s
Out: {compute a minimum spanning tree}
for each vertex u of G do {
    u.d \leftarrow \infty // distance from vertex s to vertex u
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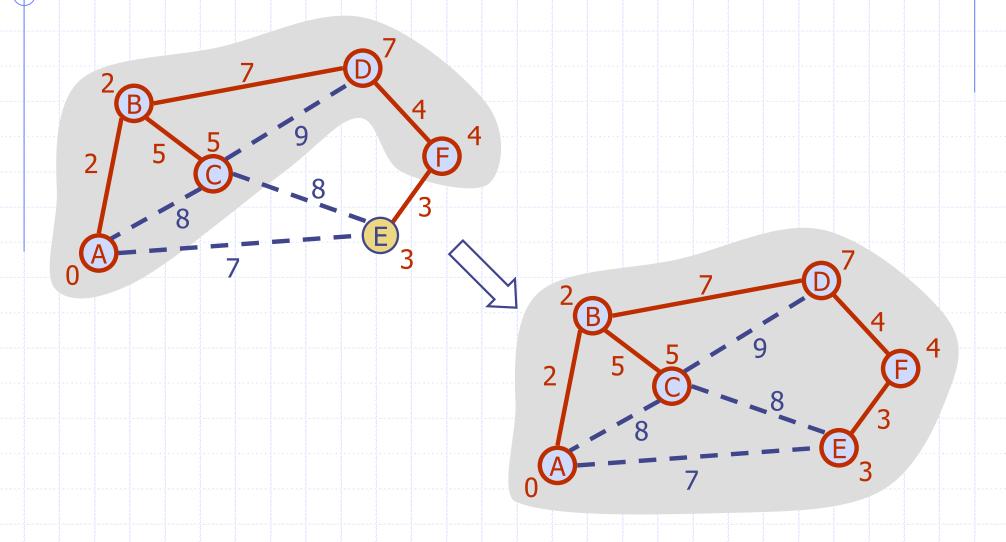
u.p \leftarrow null // predecessor or parent of vertex u in a shortest paths tree

u.marked \leftarrow false -c_1
s.d \leftarrow 0 // Distance from s to itself is 0
for i \leftarrow 0 to n-1 do {
    \min \leftarrow \infty // Find unmarked vertex u with minimum distance to s
    for each vertex v of G do
                                                                                               \begin{array}{c|c} & \sum_{u \in G} (c_4 + c_2 n + c_3 n) \\ \hline \\ & \text{(adj. matrix)} \end{array}
        if (v.marked = false) and (v.d \leq min) then {
            min \leftarrow v.d
             u \leftarrow v
    u.marked ← true // Relax all edges incident on vetex u
    for each edge (u,v) incident on u do
                                                                                               \begin{array}{c} c_3 \text{n (adj.} \\ \text{matrix)} \end{array} \quad \begin{array}{c} \sum_{u \in G} (c_4 + c_2 n + c_3 \deg(u)) \\ \text{(adj. list)} \end{array}
        if length(u,v) < v.d then {
             v.d \leftarrow length(u,v)
             v.p \leftarrow u
```

Example



Example (contd.)



Analysis of Prim's Algorithm

Using an adjacency matrix:

$$f(n,m) = c_1 n + \sum_{u \in G} (c_4 + c_2 n + c_3 n) = c_1 n + c_4 n + c_2 n^2 + c_3 n^2 \text{ is } O(n^2)$$

Using an adjacency list:

$$f(n,m) = c_1 n + \sum_{u \in G} (c_4 + c_2 n + c_3 deg(u)) = c_1 n + c_4 n + c_2 n^2 + 2c_3 m \text{ is } O(n^2)$$