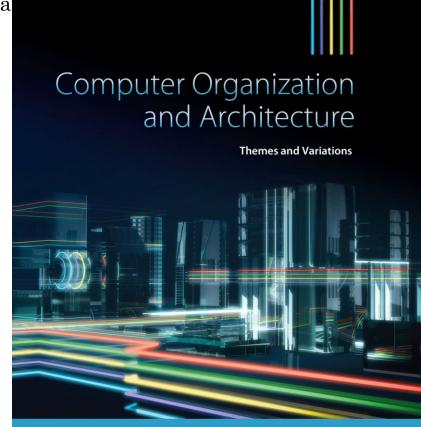
Computer Organization and Architecture: Themes and Varia

Part 2

CHAPTER 4

Computer Organization and Architecture



Alan Clements

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```
AREA TestProg, CODE, READONLY
ENTRY ; This is the calling environment.
; subroutine code is on the next slide.
```

Missing the post update value in page 237

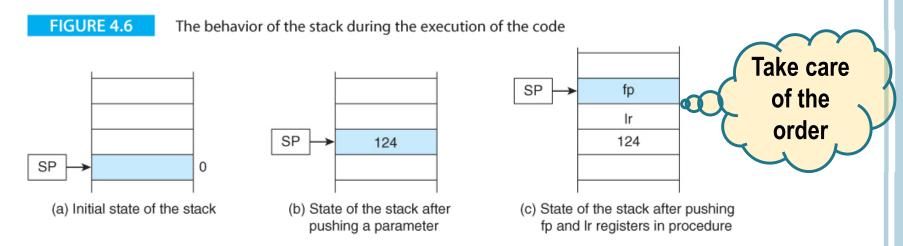
Bold is not correct in page 237

```
STMFD sp!, {fp, lr} ; push frame-pointer and link-register
Sub
     MOV
                           ; frame pointer at the bottom of
            fp,sp
                           ; the frame
     SUB
            sp, sp, #4
                          ; create the stack frame (one word)
     LDR \mathbf{r2}, [fp, #8]; get the pushed parameter
            r2, r2, #120
     ADD
                          ; do a dummy operation on
                           ; the parameter
            r2,[fp,#-4]
     STR
                           ; store it in the stack frame
            sp, sp, #4
     ADD
                           ; clean up the stack frame
     LDMFD sp!, {fp,pc}
                          ; restore frame pointer and return
      DCD
             0x0000
                           ; clear memory
             0 \times 0 0 0 0
      DCD
                                 Bold is not correct in page 238
      DCD
             0x0000
             0x0000
      DCD
             0 \times 0 0 0 0
Stack DCD
                           ; start of the stack
```

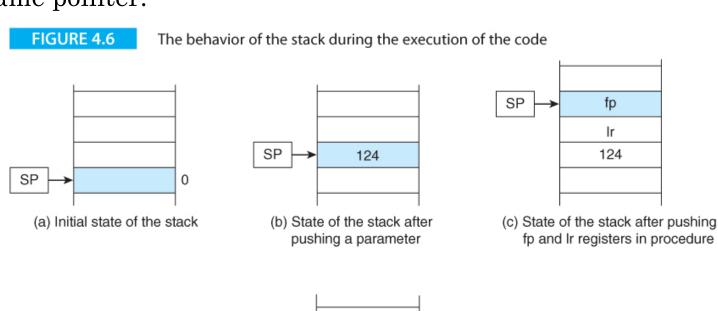
END

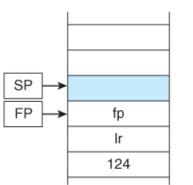
R2 has been changed inside the subroutine. Hence, it should be saved at the beginning of the subroutine and restored at the end.

□ Figure 4.6 demonstrates the behavior of the stack during the code's execution. Figure 4.6a depicts the stack's initial state. In Figure 4.6b the parameter has been pushed onto the stack. In Figure 4.6c the frame pointer and link register have been stacked by STMFD sp!, {fp,lr}.

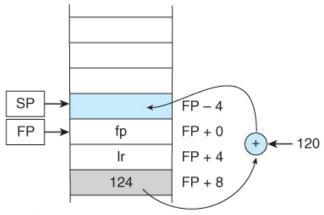


☐ In Figure 4.6d a 4-byte word has been created at the top of the stack. Finally, Figure 4.6e demonstrates how the pushed parameter is accessed and moved to the new stack frame using register indirect addressing with the frame pointer.





(d) State of stack after creating 4-byte space on the stack



(e) State of stack after the sequence
 LDR r2, [fp, #8] ;get parameter

ADD r2,r2,#120 ;add 120

STR r2, [fp, #-4] ; store sum in stack frame

25

Calling a Subroutine Step-by-Step

- ☐ To call a subroutine, the following steps need to be performed:
 - *Parameters* need to be *passed* from the caller to the subroutine.

 This can be performed via the stack.
 - The *address* of the instruction immediately after the calling instruction needs to be *saved in a safe place BEFORE* branching to the subroutine.

This can be performed by using BL instruction or via the stack, or both.

- o *Inside the subroutine*, we need to:
 - Push the values of all registers to be used inside the subroutine, as well as the FP (R11) and LR (R14).
 - Make the FP (R11) point to the bottom of the frame by copying the value of the SP (R13) to the FP (R11).
 - Create a space inside the stack for local variables.
 - Perform the subroutine instructions.
 - The addresses of parameters and local variables are calculated relative to the value of the FP (R11).
 - At the end of the subroutine, deallocate all created local variables.
 - Pop all pushed registers but use PC (R15) instead of LR (R14).
 - At the caller program, all pushed parameters need to be popped.

- ☐ You can pass a parameter to a subroutine
 - o by value
 - o by reference
- ☐ When passed *by value*, the subroutine receives a <u>copy</u> of the parameter.
 - Passing a parameter by value causes the *parameter to be cloned* and the *cloned version of the parameter* to be used by the subroutine.
 - o If the parameter is modified by the subroutine, the new value does not affect the value of the parameter elsewhere in the program.
- □ When passed *by reference*, the subroutine receives a <u>pointer</u>, (i.e., an <u>address</u>) to the parameter.
 - o *There is only one copy of the parameter*, and the subroutine can access this value because it knows the address of the parameter.
 - o If the subroutine modifies the parameter, it is modified the original value.

- ☐ The subroutine swap (int a, int b) *intends* to exchange two values.
- □ Let's examine how parameters are passed to this subroutine.



```
ENTRY
ADR sp,STACK
MOV fp,#0xFFFFFFFF
B main

SPACE 0x20
STACK DCD 0

ENTRY

; set up stack pointer
; set up dummy fp for tracing
; jump to the function main
```

You need to re-do it yourself using the other stack types.

```
; void swap (int a, int b)
; Parameter a is at [fp]+4
; Parameter b is at [fp]+8
; Variable temp is at [fp]-4
```

FD Stack

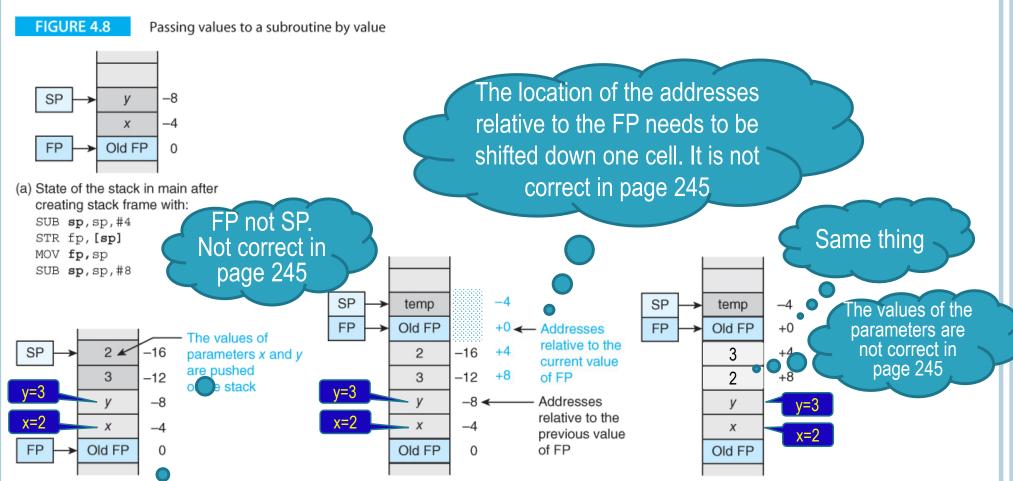
Passing Parameters via the Stack

You need to re-do it yourself using the other stack types.

```
swap SUB
         sp, sp, #4
                       ; Create stack frame: decrement sp
    STR fp,[sp]
                      ; push the frame pointer onto the stack
                      ; frame pointer points at the base
          fp,sp
    VOM
    int temp;
    SUB
          sp, sp, #4
                   ; move sp up 4 bytes for temp
    temp = a;
    LDR r0,[fp,#4] ;get parameter a from the stack
    STR r0, [fp,#-4]; copy a to temp onto the stack frame
         = b;
    a
    LDR r0, [fp, #8] ; get parameter b from the stack
    STR r0, [fp, #4] ; copy b to a
    b
         = temp;
    LDR \mathbf{r0}, [fp, #-4]; get temp from the stack frame
         r0,[fp,#8] ;copy temp to b
    STR
                       ; Collapse stack frame created for swap
                      ; restore the stack pointer
    VOM
          sp, fp
    LDR
          fp, [sp]
                       ; restore old frame pointer from stack
                      ; move stack pointer down 4 bytes
          sp, sp, #4
    ADD
                                                             29
          pc, lr
    VOM
                      ; return by loading LR into PC
```

```
void main(void)
main
                       ;Create stack frame in main for x, y
          sp, sp, #4
     SUB
                       ; move the stack pointer up
     STR fp,[sp].
                       ; push the frame pointer onto the stack
                      ;the frame pointer points at the base;
          fp,sp
     MOV
     int x = 2, y = 3; Bold is not correct in page 244
     SUB sp, sp, #8 ; move sp up 8 bytes for 2 integers
     MOV r0, #2
                     ; x = 2
     STR r0,[fp,\#-4]; put x in stack frame
     MOV r0, #3
                       ; y = 3
          r0, [fp,#-8] ; put y in stack frame
     STR
     swap(x, y);
     LDR
          r0,[fp,#-8]; get y from stack frame
     STR
          r0,[sp,#-4]! ;push y on stack
          r0, [fp, #-4]; get x from stack frame
     LDR
     STR
          r0, [sp, #-4]!]; push x on stack
     BL
                       ; call swap, save return address in LR
          swap
          sp, sp, #8
     ADD
                       ; Clean the stack from the parameters
     VOM
          sp,fp
                       ; restore the stack pointer
     LDR
                       ; restore old frame pointer from stack
          fp, [sp]
          sp, sp, #4
                       ; move stack pointer down 4 bytes
     ADD
Loop B
          Loop
                       ;Stop
```

- ☐ This code swaps the variables inside the stack frame
- ☐ When the return is made, the stack frame will be collapsed, and the effect of the swap will be lost.
- \Box The variables in the calling environment are not affected.



(b) The stack in main after putting two parameters in the stack frame with:

```
MOV r0,#2
```

Then pushing two parameters on the stack

(c) The stack after the creation of a stack frame in swap. The new stack frame is four bytes deep and holds the variable temp. The frame is created by:

SUB sp,sp,#4

(d) The stack after executing the body of swap. Note that all data is referenced to FP.

☐ In the next example, we pass parameters by reference

ENTRY
ADR sp, STACK ;set up stack pointer
MOV fp,#0xFFFFFFFF ;set up dummy fp for tracing
B main ;jump to main function

SPACE 0x20

STACK DCD 0

; void swap (int *a, int *b)
; Parameter *a is at [fp]+4

; Parameter *b is at [fp]+8

Variable temp is at [fp]-4

```
sp, sp, #4
    SUB
                         ; Create stack frame: decrement sp
swan
     STR fp,[sp]
                         ; push the frame pointer onto the stack
     VOM
                         ; frame pointer points at the base
           fp,sp
     int temp;
         sp, sp, #4
     SUB
                        ; move sp up 4 bytes for temp
     temp = *a;
           r1, [fp, #4] ; get address of parameter a
     LDR
     LDR r2,[r1] ; get value of parameter a (i.e., *a)
     STR r2, [fp, #-4]; store *a in temp in stack frame
     *a = *b;
     LDR r0, [fp, #8] ; get address of parameter b
     LDR r3,[r0] ; get value of parameter b (i.e., *b)
     STR r3, [r1] ; store *b in *a
                                            Missing the *
     *b) =
          temp;
                                            in page 247
     LDR
           r3, [fp, #-4] ; get temp
     STR
           r3,[r0]
                          ;store temp in *b
                         ; Collapse stack frame created for swap
     VOM
                         ; restore the stack pointer
           sp, fp
                         ;restore old frame pointer from stack^{35}
           fp, [sp]
     LDR
           sp, sp, #4
                         ; move stack pointer down 4 bytes
                         ; return by loading LR into PC Clements and used with permission. New content added and copyrighted by © Mahmoud R. El-Sakka.
           pc, lr
```

```
void main(void)
main
                       ;Create stack frame in main for x, y
          sp, sp, #4
     SUB
                       ; move the stack pointer up
     STR fp,[sp].
                       ; push the frame pointer onto the stack
          fp,sp
                       ;the frame pointer points at the base;
     MOV
     int x = 2, y = 3; Bold is not correct in page 244
     SUB sp, sp, #8 ; move sp up 8 bytes for 2 integers
     MOV r0, #2
                     ; x = 2
     STR r0,[fp,#-4]
                       ; put x in stack frame
                       ; y = 3
     MOV r0, #3
          r0,[fp,#-8] ;put y in stack frame
     STR
     swap(&x, &y);
     SUB
          r0, fp, #8
                    ; get address of y in stack frame
     STR
          r0,[sp,#-4]! ;push address of y on stack
     SUB
          r0, fp, #4
                   ; get address of x in stack frame
     STR
          r0, [sp, #-4]! ; push address of x on stack
     BL
                       ; call swap, save return address in LR
          swap
          sp, sp, #8
     ADD
                       ; Clean the stack from the parameters
     MOV
          sp,fp
                       ; restore the stack pointer
                       ; restore old frame pointer from stack
     LDR
          fp, [sp]
          sp,sp,#4
                       ; move stack pointer down 4 bytes
     ADD
Loop B
          Loop
                       ;Stop
```

☐ In the function main, the addresses of the *parameters are pushed onto the stack* by means of the following instructions:

```
SUB r0,fp,#8 ;get address of y in stack frame
STR r0,[sp,#-4]! ;push address of y on stack
SUB r0,fp,#4 ;get address of x in stack frame
STR r0,[sp,#-4]! ;push address of x on stack
```

☐ In the function swap, the addresses of *parameters are read from the stack* by means of

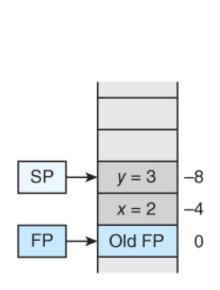
```
temp = *a;
LDR r1,[fp,#4] ;get address of parameter a
LDR r2,[r1] ;get value of parameter a (i.e., *a)
STR r2,[fp,#-4] ;store *a in temp in stack frame

; *a = *b;
LDR r0,[fp,#8] ;get address of parameter b
LDR r3,[r0] ;get value of parameter b (i.e., *b)
STR r3,[r1] ;store *b in *a

; *b = temp;
LDR r3,[fp,#-4] ;get temp
STR r3,[r0] ;store temp in *b
```

FIGURE 4.9

Passing values to a subroutine by reference

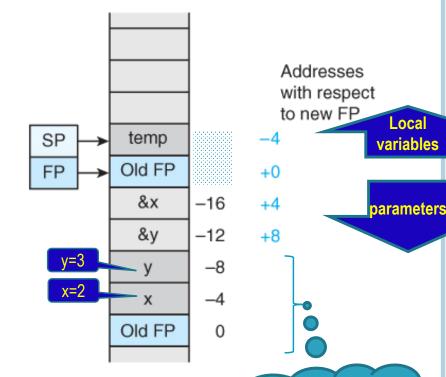


(a) State of the stack after

MOV r0, #3

SP &x -16 &y -12 -8 x=2 x Old FP x=2 x=2 x=2 x=2 x=3 x=4 x=4

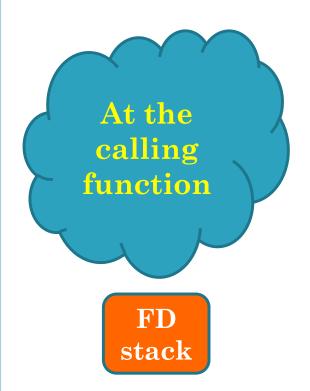
(b) State of the stack after pushing parameter addresses by



(c) State of the stack after subroutine call and stack frame created by

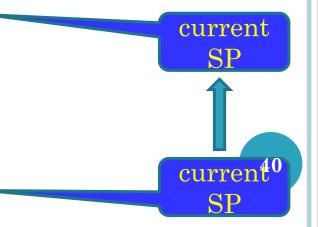
The swap function should not have a direct access to x and y

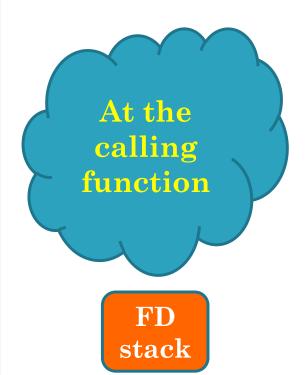
You need to re-do it yourself using the other At the stack types. calling function FD curren³⁹ stack



The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack





The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP

41

The subroutine to store inside the stack the value of all registers to be utilized during the function.

These registers, including

FP

LR

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP current FP

current

42

At the beginning of the function

FD stack

The subroutine to store inside the stack the value of all registers to be utilized during the function.

These registers, including

FP

LR

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP current FP

At the beginning of the function

FD stack

The function calculates the addresses of the local variables relative to the current FP value.

The function calculates the addresses of the parameters and the returning value relative to the current FP value.

At the beginning of the function

FD stack The subroutine to allocate memory inside the stack for the local variables

The subroutine to store inside the stack the value of all registers to be utilized during the function.

These registers, including

FP

LR

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP

current SP

current

The function calculates the addresses of the <u>local variables</u> relative to the current FP value.

call by value vs call by reference

The function calculates the addresses of the parameters and the returning value relative to the current FP value.

At the beginning of the function

FD stack The subroutine to allocate memory inside the stack for the local variables

The subroutine to store inside the stack the value of all registers to be utilized during the function.

These registers, including

FP

LR

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP

current FP



The function calculates the addresses of the <u>local variables</u> relative to the current FP value.

The function calculates the addresses of the parameters and the returning value relative to the current FP value.

At the end of the function

> FD stack

The subroutine to allocate memory inside the stack for the local variables

The subroutine to store inside the stack the value of all registers to be utilized during the function.

These registers, including

FP

LR

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP

> current SP

current FP

The function calculates the addresses of the parameters and the returning value relative to the current FP value.

At the end of the function

> FD stack

The subroutine to store inside the stack the value of all registers to be utilized during the function.

These registers, including

FP

LR

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP current FP

The function calculates the addresses of the parameters and the returning value relative to the current FP value.

At the end of the function

> FD stack

The subroutine to store inside the stack the value of all registers to be utilized during the function.

These registers, including

FP

LR

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP current FP

> current SP

LDM all the stored register values, where the LR value to be loaded as PC. Hence, returning to the caller function



FD stack The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP

LDM all the stored registers, where the LR is loaded as PC. Hence, returning to the caller function

The returned value to be accessed and popped from the stack, as well as the parameters.



stack

The caller to allocate memory inside the stack for the returning value

The caller to push the parameters on the stack

current SP 50

current



current SP 51