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$$\begin{array}{lll} \text{I. (i)} & p \vee q & \neg q \equiv T \quad p \vee F \equiv T \\ & p \rightarrow r & \\ & \neg q & q \equiv F \quad p \equiv T \text{ (domination laws)} \\ \hline & \therefore r & \end{array}$$

$$p \rightarrow r \equiv \neg p \vee r \text{ (conditional identity)}$$

$$\equiv F \vee r$$

$$F \vee r \equiv T$$

$$r \equiv T \text{ (domination laws)}$$

\therefore the proof is valid.

$$\text{(ii)} \quad \forall x (P(x) \rightarrow (Q(x) \wedge (R(x))))$$

$$\forall x (S(x) \vee P(x))$$

$$\neg \forall x Q(x)$$

$$\exists x S(x).$$

$$\neg \forall x Q(x) \equiv \exists x \neg Q(x).$$

$$\therefore \forall x Q(x) \equiv F.$$

(condition identity)

$$\forall x (P(x) \rightarrow (Q(x) \wedge (R(x)))) \equiv \neg (\forall x (P(x))) \vee ((Q(x) \wedge (R(x))))$$

$$\equiv \neg (\forall x (P(x))) \vee (F \wedge (R(x)))$$

$$\equiv \neg (\forall x (P(x))) \vee F \text{ (domination laws).}$$

$$\neg (\forall x (P(x))) \vee F \equiv T$$

$$\neg (\forall x (P(x))) \equiv T \text{ (domination laws).}$$

$$\forall x (P(x)) \equiv F.$$

$$\forall x (S(x) \vee P(x)) \equiv (\forall x S(x)) \vee (\forall x P(x))$$

$$\equiv (\forall x S(x)) \vee F.$$

$$(\forall x S(x)) \vee F \equiv T.$$

$$\forall x S(x) \equiv T \text{ (domination laws).}$$

$$\therefore \exists x S(x) \equiv T$$

\therefore the proof is valid.

2. (i) Some student don't have an official email address.

(ii) Exist a real number x that there is no real number y satisfies $x = y^2$

3. Theorem: P_1, \dots, P_n is a collection of propositions.

$$\neg(P_1 \vee \dots \vee P_n) \equiv (\neg P_1 \wedge \dots \wedge \neg P_n).$$

Proof: $(P_1 \vee \dots \vee P_n)$ is false iff all propositions from P_1 to P_n are false, so $\neg(P_1 \vee \dots \vee P_n)$ is true iff all propositions are false.

Assume a new collection q_1, \dots, q_n where $q_n = \neg P_n$.

$q_1 \wedge \dots \wedge q_n$ is true iff all propositions from q_1 to q_n are true, so $\neg P_1 \wedge \dots \wedge \neg P_n$ is true iff all propositions are false.

the only true condition is the same from both sides, so $\neg(P_1 \vee \dots \vee P_n) \equiv (\neg P_1 \wedge \dots \wedge \neg P_n)$ ■

4. Theorem: P a predicate depends on variables x_1, \dots, x_n

Q denote $\forall x_n \forall x_{n-1} \dots \forall x_1 P$ for all $n \geq 1$.

$$\neg Q = \exists x_n \exists x_{n-1} \dots \exists x_1 \neg P$$

Proof: P is false iff $\exists x_n \exists x_{n-1} \dots \exists x_1 P$ is false, so

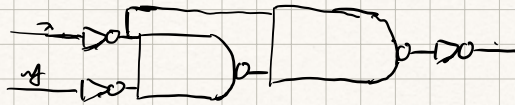
$\neg P$ is true iff $\exists x_n \exists x_{n-1} \dots \exists x_1 \neg P$ is true.

Q is false iff $\exists x_n \exists x_{n-1} \dots \exists x_1, P$ is false. so
 $\neg Q$ is true iff $\exists x_n \exists x_{n-1} \dots \exists x_1, \neg P$ is true. ■

5. (i) $f(x, y) = xy + x\bar{y} + \bar{x}y$.

(ii) $x \uparrow y = \overline{x \cdot y} = \bar{x} + \bar{y}$.

$$f(x, y) = x + \bar{x}y = x + \overline{\bar{x} \cdot \bar{y}} = \overline{\bar{x} \cdot \bar{x} \cdot \bar{y}}.$$



(iii). $g(x, y) = x(\bar{x} + \bar{y})$
 $= x\bar{x} + x\bar{y}$
 $= F + x\bar{y}$
 $= x\bar{y}$.

Not the same.

6 (1) One which can be used to express all possible truth tables by combining members of the set into a Boolean expression.

(2) $xy = (x+y) \cdot \overline{x \cdot y}$
 $= \overline{\bar{x} \cdot \bar{y}} \cdot \overline{x \cdot y}$
 $= \overline{\bar{x} \cdot \bar{y} \cdot xy}$
 $\bar{a} + \bar{b} = \overline{ab}$

7. To prove $f(x, y) \cdot g(x, y)$ is satisfiable when both $f(x, y)$, $g(x, y)$ are satisfiable, is equal to prove its contrapositive:
 If either $f(x, y)$, $g(x, y)$ is not satisfiable or both $f(x, y)$ and $g(x, y)$ are unsatisfiable, then $f(x, y) \cdot g(x, y)$ is not satisfiable

when either $f(x, y)$, $g(x, y)$ is not satisfiable, the truth

value of $\neg f(x,y) \cdot g(x,y)$ is 0

when both $\neg f(x,y), g(x,y)$ are not satisfiable, the truth value of $\neg f(x,y) \cdot g(x,y)$ is 0.

Thus, the contrapositive is true. So $\neg f(x,y) \cdot g(x,y)$ is satisfiable when both $\neg f(x,y), g(x,y)$ are satisfiable.

8. Assume $F_n = 5^n - 1$.

Base case: $n=1$ $5^1 - 1 = 4$ can be divided by 4.

when $n \geq 2$, $F_n - F_{n-1} = 5^n - 1 - 5^{n-1} + 1 = 4 \cdot 5^{n-1}$.

so k can be 5^{n-1} that $5^n - 1$ can be divided by 4 for any $n \geq 2$.

So $5^n - 1$ can be divided by 4.