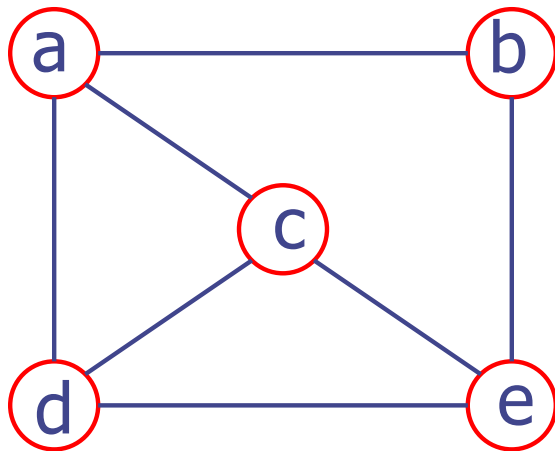


Graphs

A graph is a pair (V, E) , where

- V is a set of **nodes** or **vertices**
- E is a collection of pairs of vertices (u,v) , called **edges, links, or arcs**



$$V = \{a, b, c, d, e\}$$

$$E = \{(a,b), (a,c), (a,d), (b,e), (c,d), (c,e), (d,e)\}$$

Edge Types

- Directed edge (*vector*) .
 - ordered pair of vertices (u, v)
 - first vertex u is the origin
 - second vertex v is the destination



Edge Types

- Directed edge
 - ordered pair of vertices (u,v)
 - first vertex u is the origin
 - second vertex v is the destination
- Undirected edge
 - unordered pair of vertices (u,v)



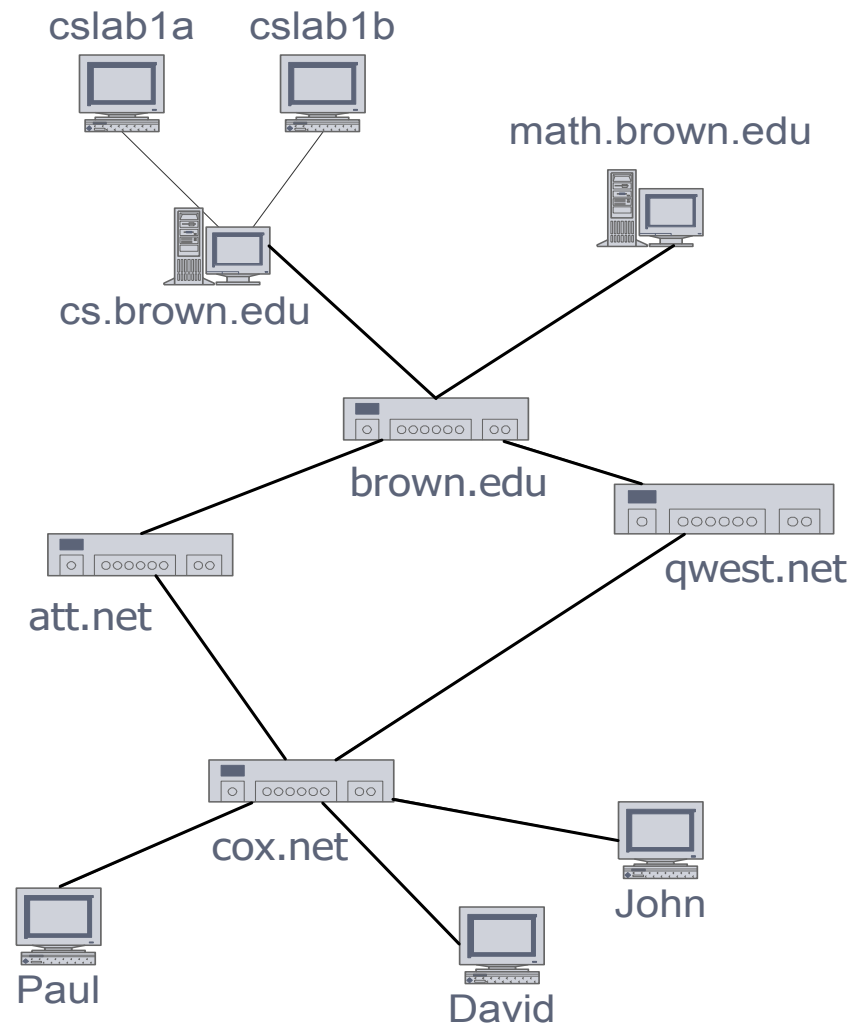
Edge Types

- ❑ Directed edge
 - ordered pair of vertices (u,v)
 - first vertex u is the origin
 - second vertex v is the destination
- ❑ Undirected edge
 - unordered pair of vertices (u,v)
- ❑ Directed graph or digraph
 - all the edges are directed
- ❑ Undirected graph
 - all the edges are undirected
- ❑ Mixed graph
 - directed and undirected edges



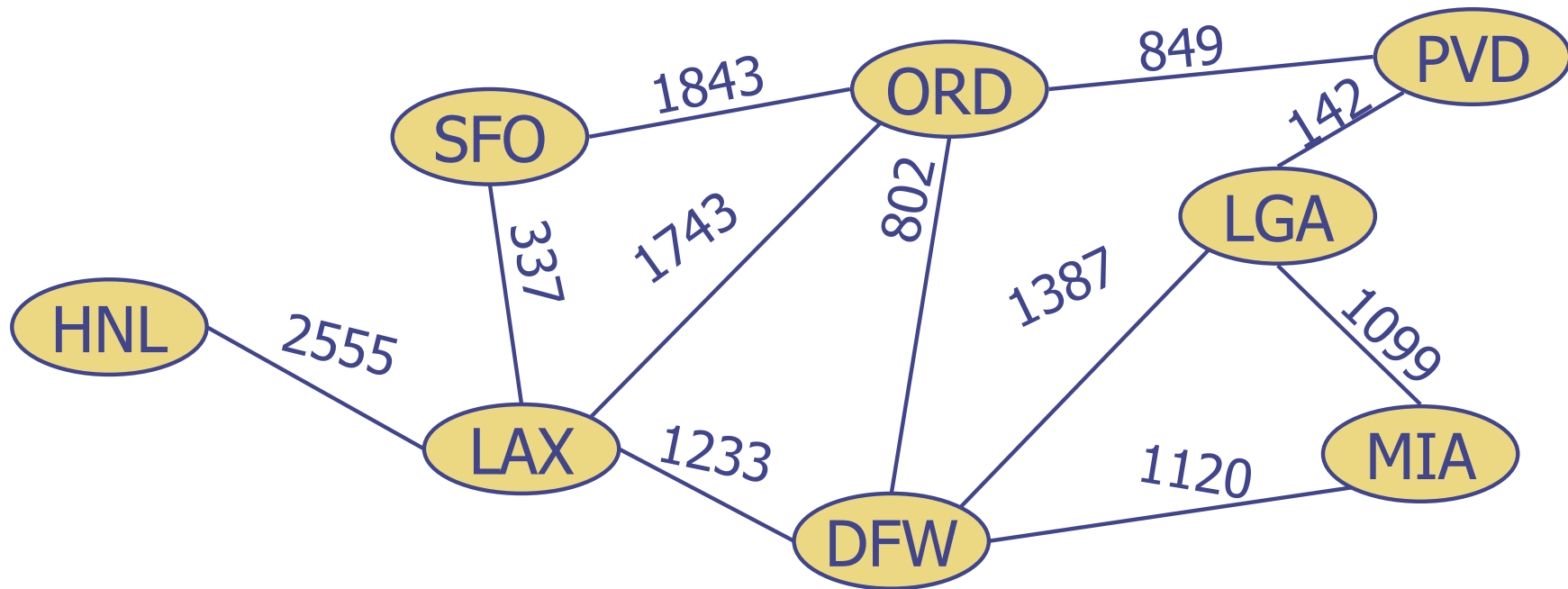
Applications

□ Computer networks



Applications

- Transportation networks

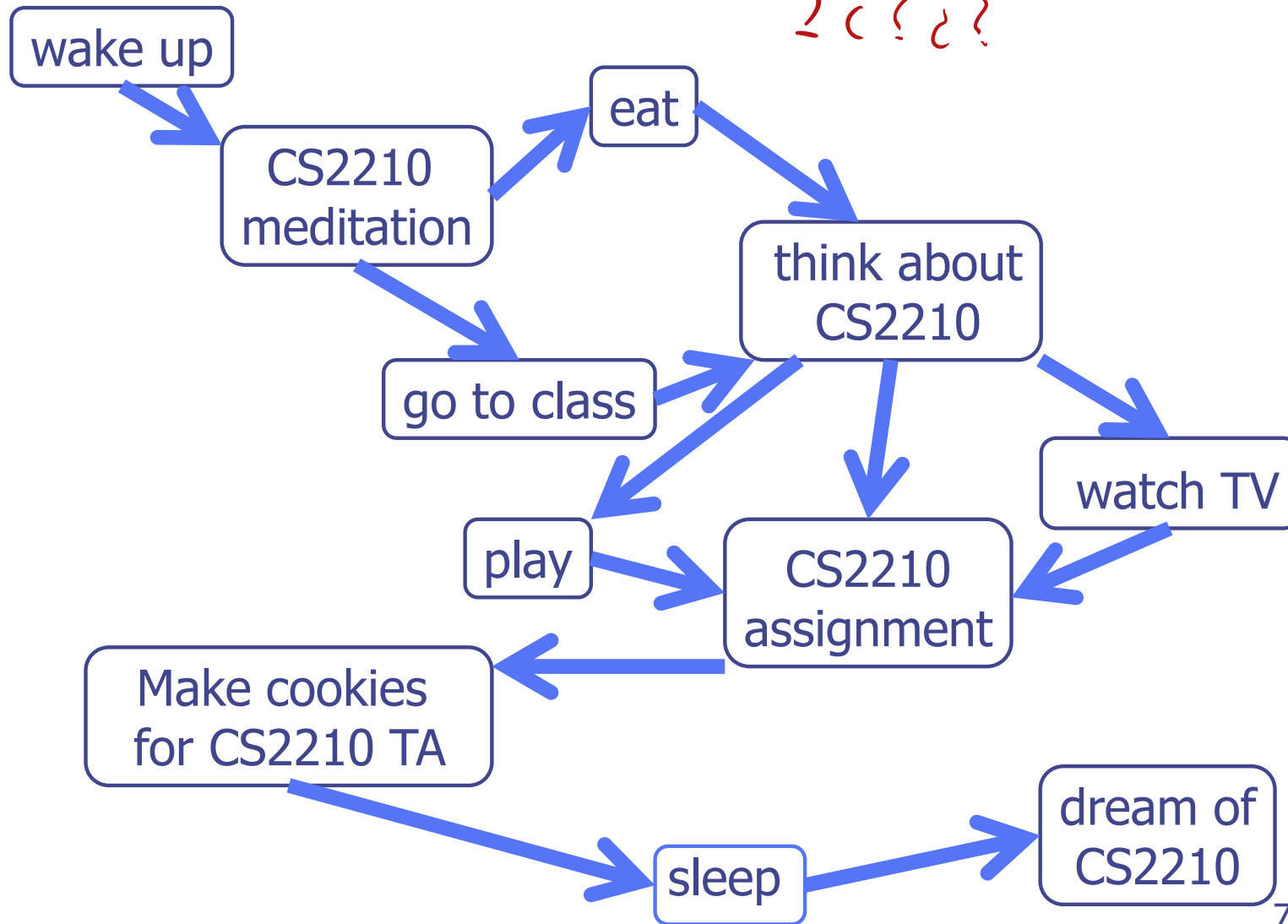


Applications

- Scheduling tasks

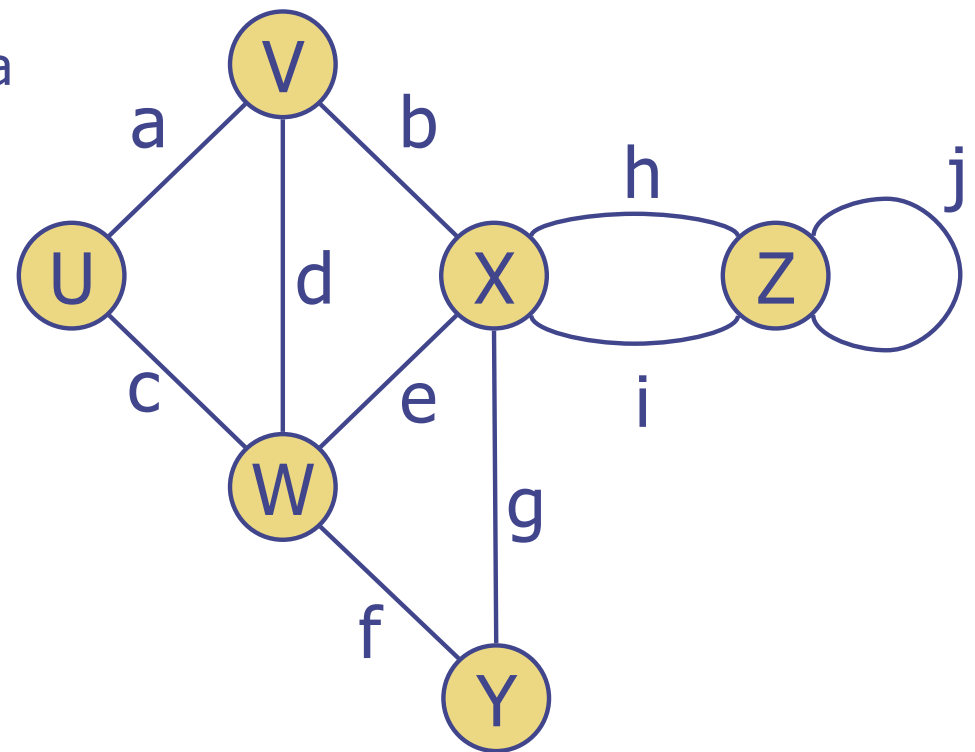
A typical student day

2 0 0 0 0



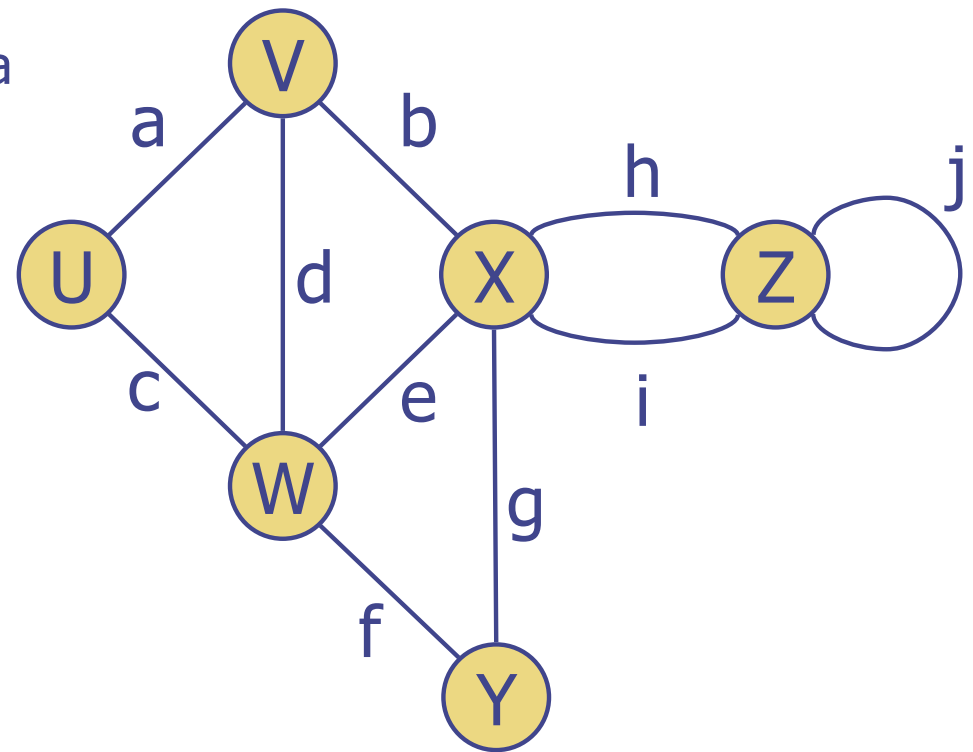
Terminology

- End vertices (or endpoints) of an edge
 - U and V are the endpoints of a



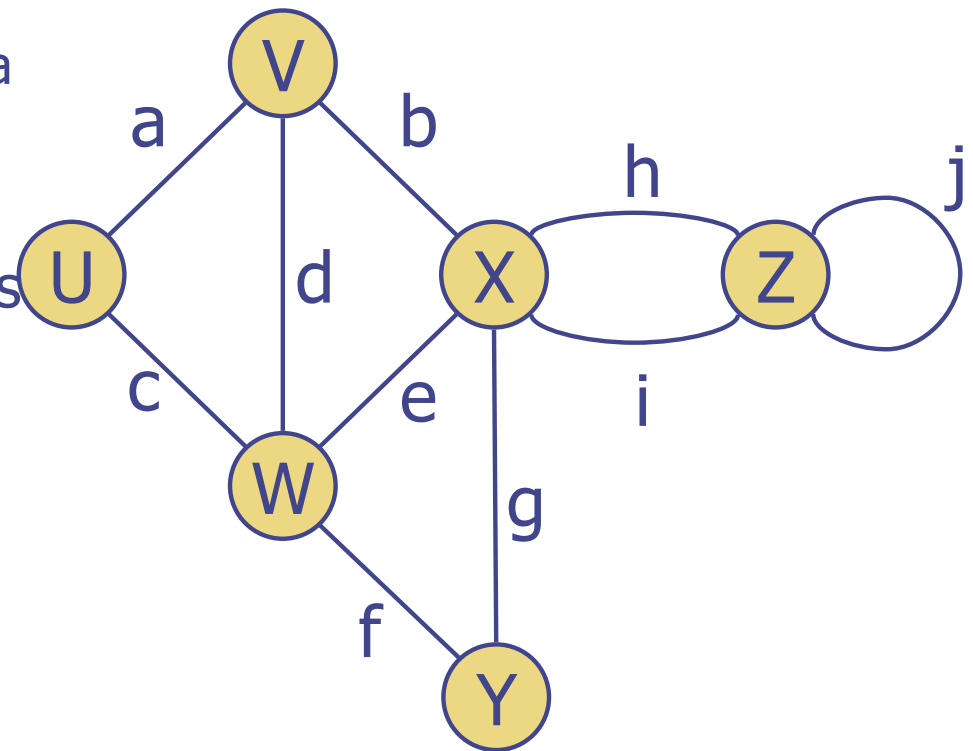
Terminology

- End vertices (or endpoints) of an edge
 - U and V are the endpoints of a
- Edges incident on a vertex
 - a, d, and b are incident on V



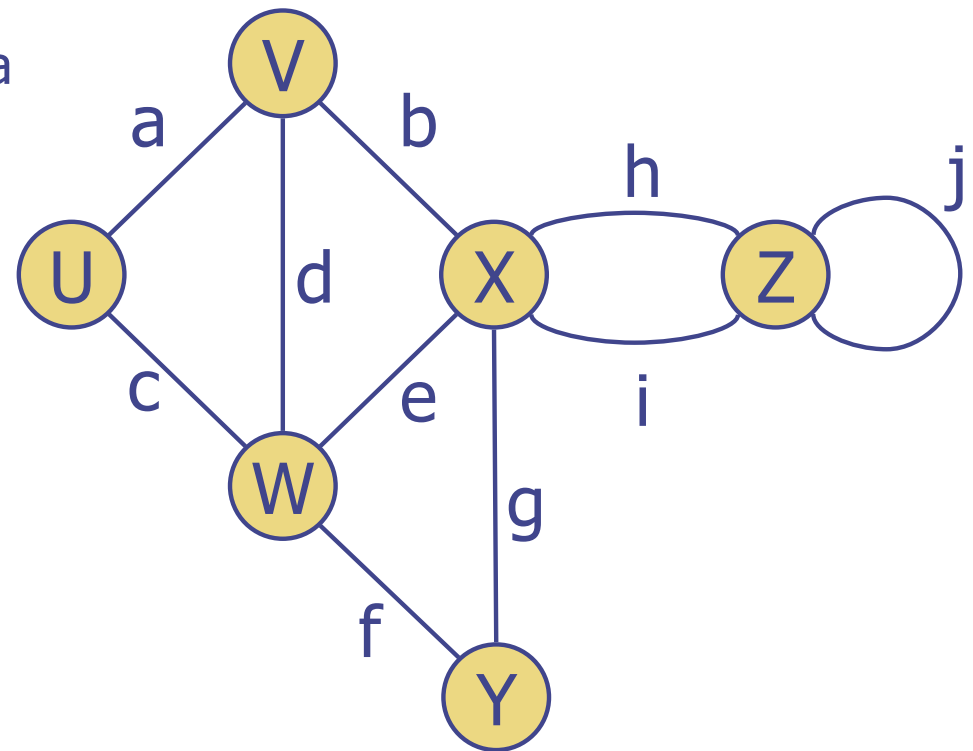
Terminology

- End vertices (or endpoints) of an edge
 - U and V are the endpoints of a
- Edges incident on a vertex
 - a, d, and b are incident on V
- Adjacent vertices or neighbours
 - U and V are adjacent



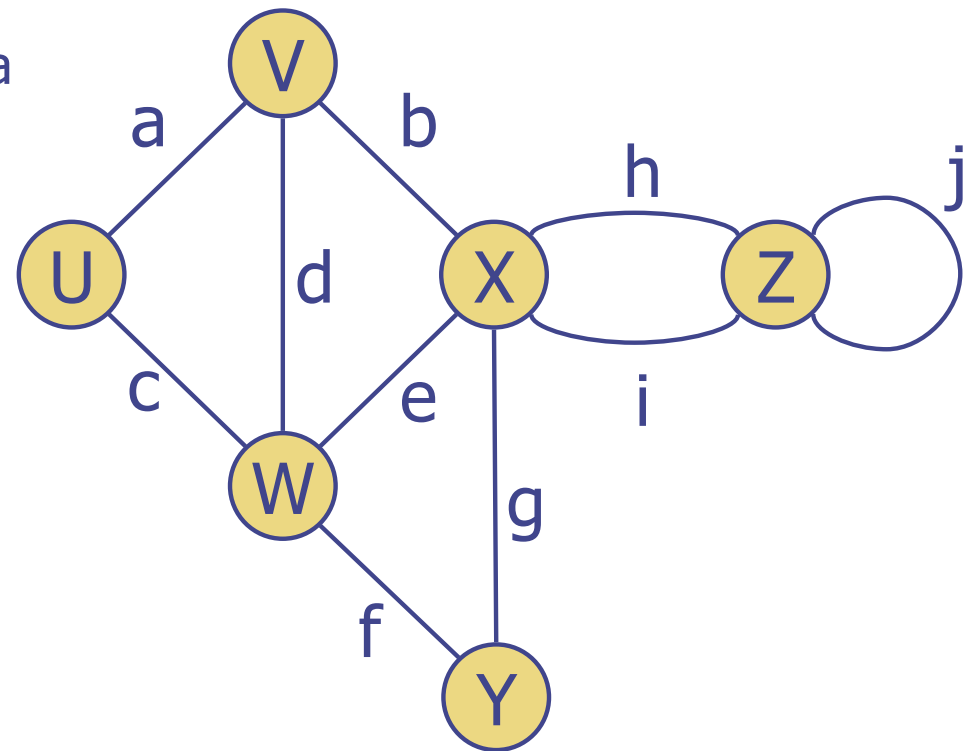
Terminology

- End vertices (or endpoints) of an edge
 - U and V are the endpoints of a
- Edges incident on a vertex
 - a, d, and b are incident on V
- Adjacent vertices
 - U and V are adjacent
- Degree of a vertex
 - X has degree 5



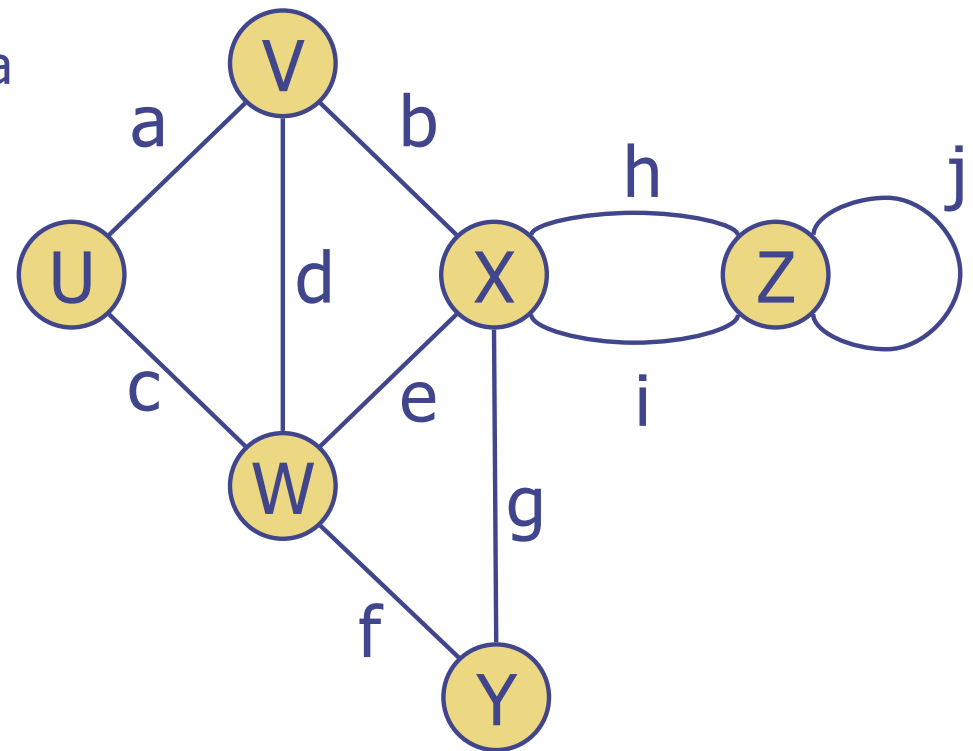
Terminology

- ❑ End vertices (or endpoints) of an edge
 - U and V are the endpoints of a
- ❑ Edges incident on a vertex
 - a, d, and b are incident on V
- ❑ Adjacent vertices
 - U and V are adjacent
- ❑ Degree of a vertex
 - X has degree 5
- ❑ Parallel edges
 - h and i are parallel edges
- ❑ Self-loop
 - j is a self-loop



Terminology

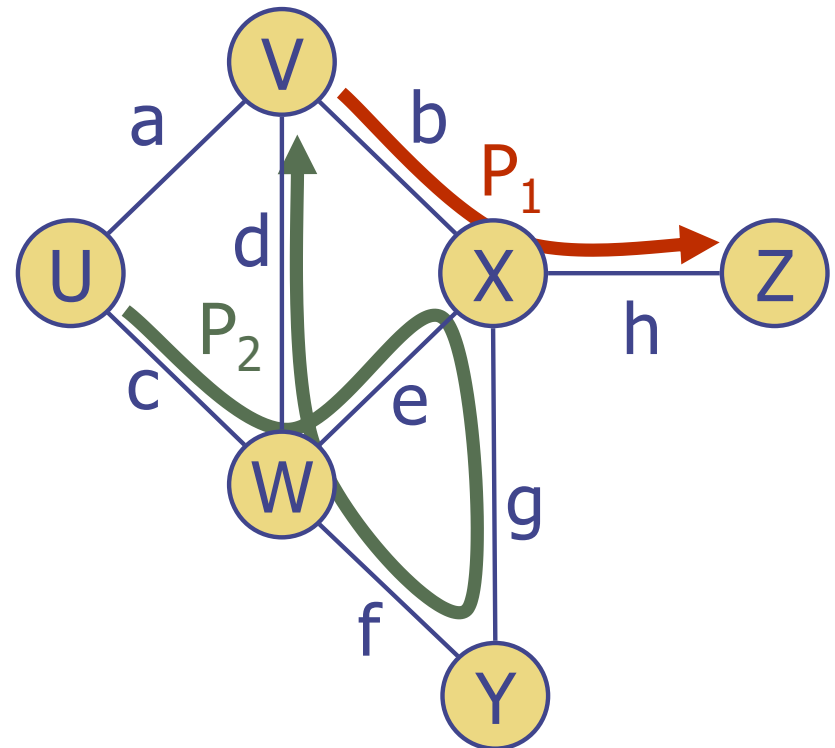
- ❑ End vertices (or endpoints) of an edge
 - U and V are the endpoints of a
- ❑ Edges incident on a vertex
 - a, d, and b are incident on V
- ❑ Adjacent vertices
 - U and V are adjacent
- ❑ Degree of a vertex
 - X has degree 5
- ❑ Parallel edges
 - h and i are parallel edges
- ❑ Self-loop
 - j is a self-loop



A graph without parallel edges or self-loops is called a simple graph

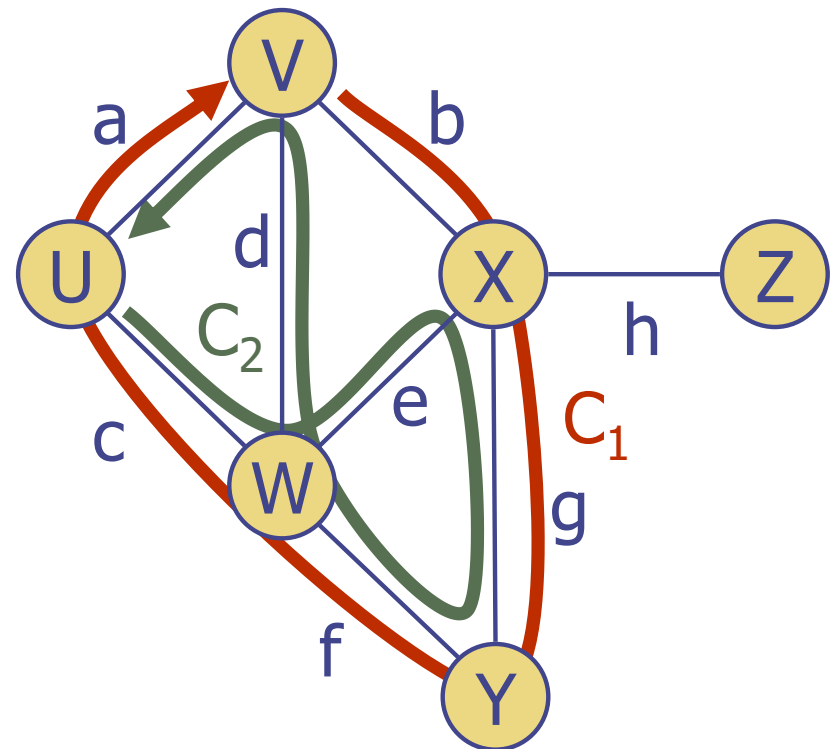
Terminology

- Path
 - sequence of adjacent vertices
- Simple path
 - path such that all its vertices and edges are distinct
- Examples
 - $P_1 = (V, X, Z)$ is a simple path
 - $P_2 = (U, W, X, Y, W, V)$ is a path that is not simple



Terminology

- Cycle
 - circular sequence of adjacent vertices
- Simple cycle
 - cycle such that all its vertices are distinct (except first and last)
- Examples
 - $C_1 = (V, X, Y, W, U, V)$ is a simple cycle
 - $C_2 = (U, W, X, Y, W, V, U)$ is a cycle that is not simple



Properties

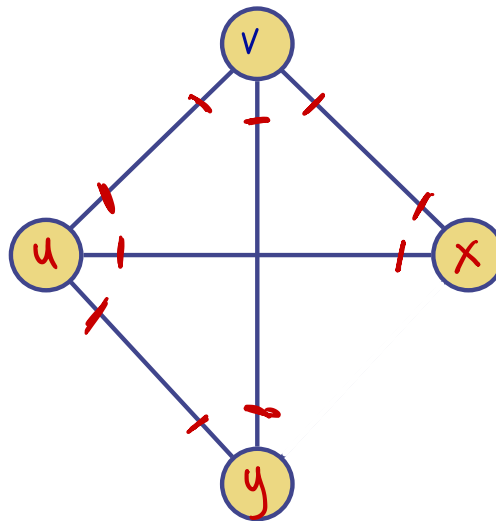
Notation

n	number of vertices
m	number of edges
$\deg(v)$	degree of vertex v

Property 1

$$\sum_v \deg(v) = 2m$$

$$m = \frac{1}{2} \sum_v \deg(v)$$

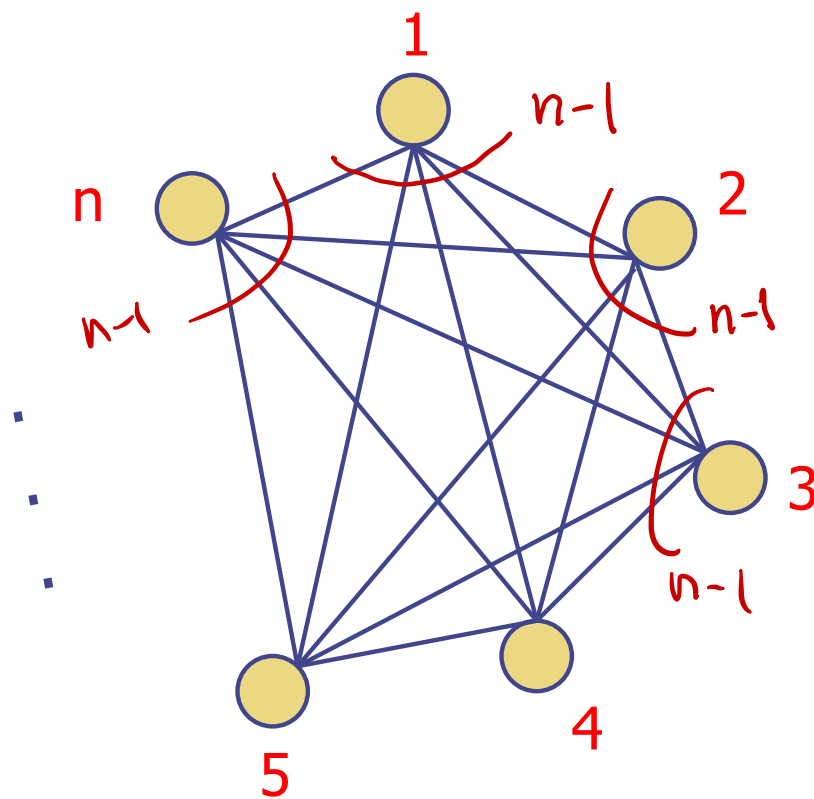


Example

- $n = 4$
- $m = 5$
- $\deg(v) = 3$
- $\sum_v \deg(v) = 10$

Complete Graph or Clique

Each vertex is connected to every other vertex.

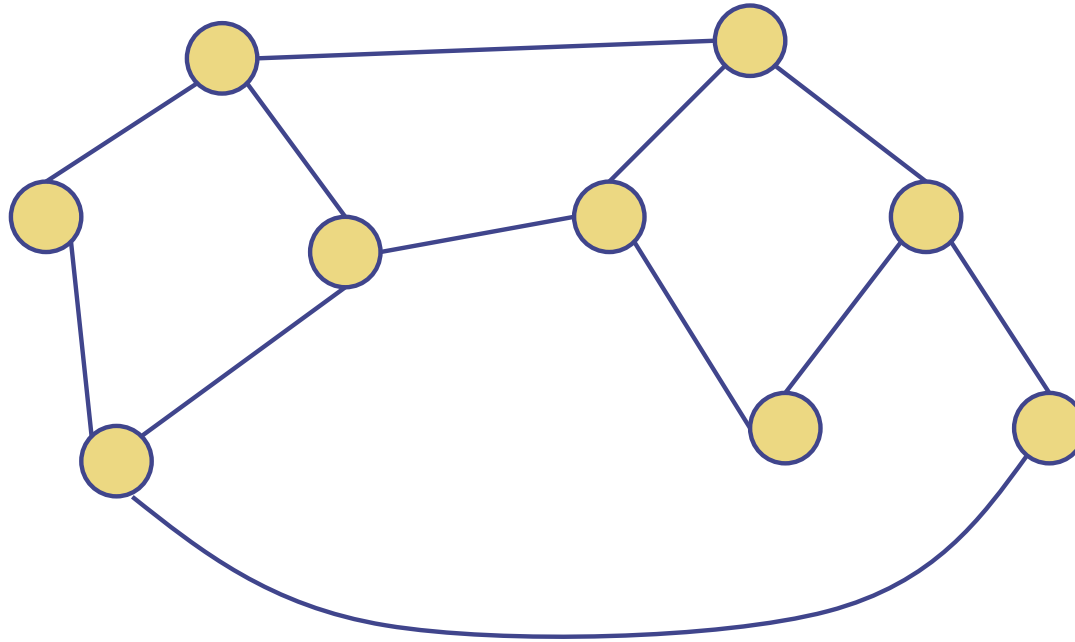


$$\sum_u \frac{\text{degree}(u)}{n-1} = n(n-1) = 2m$$

$$m = n(n-1)/2$$

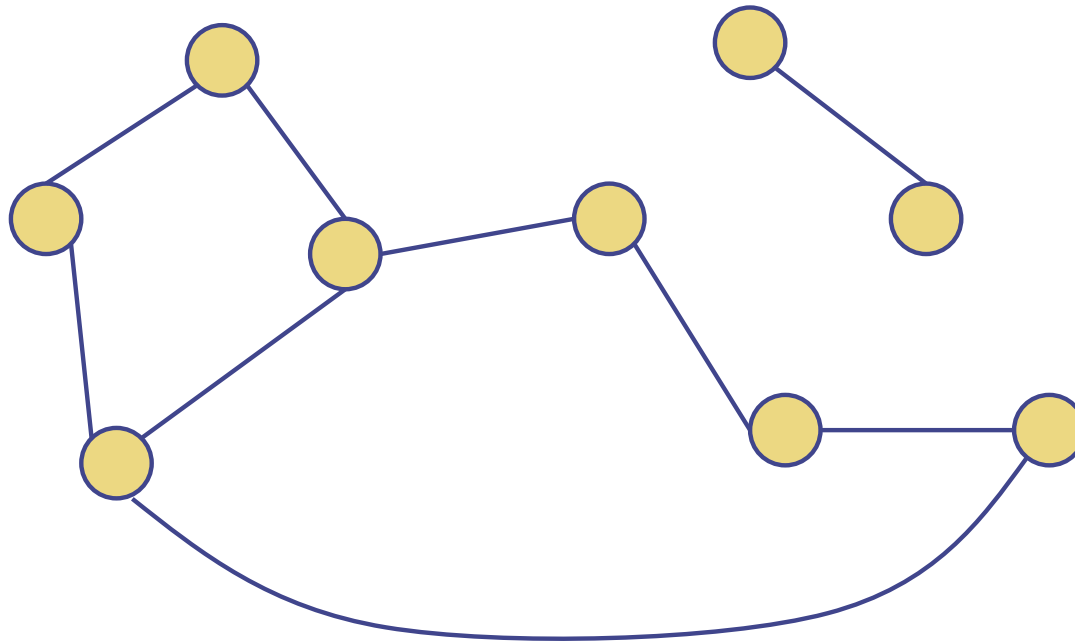
Connected Graph

A graph is connected if there is a path from each vertex to every other vertex



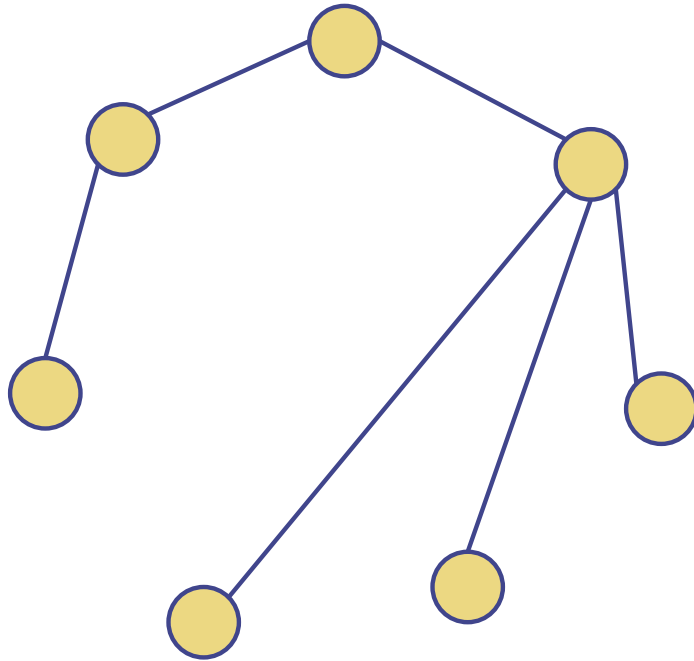
Connected Graph

A graph is connected if there is a path from each vertex to every other vertex



Trees

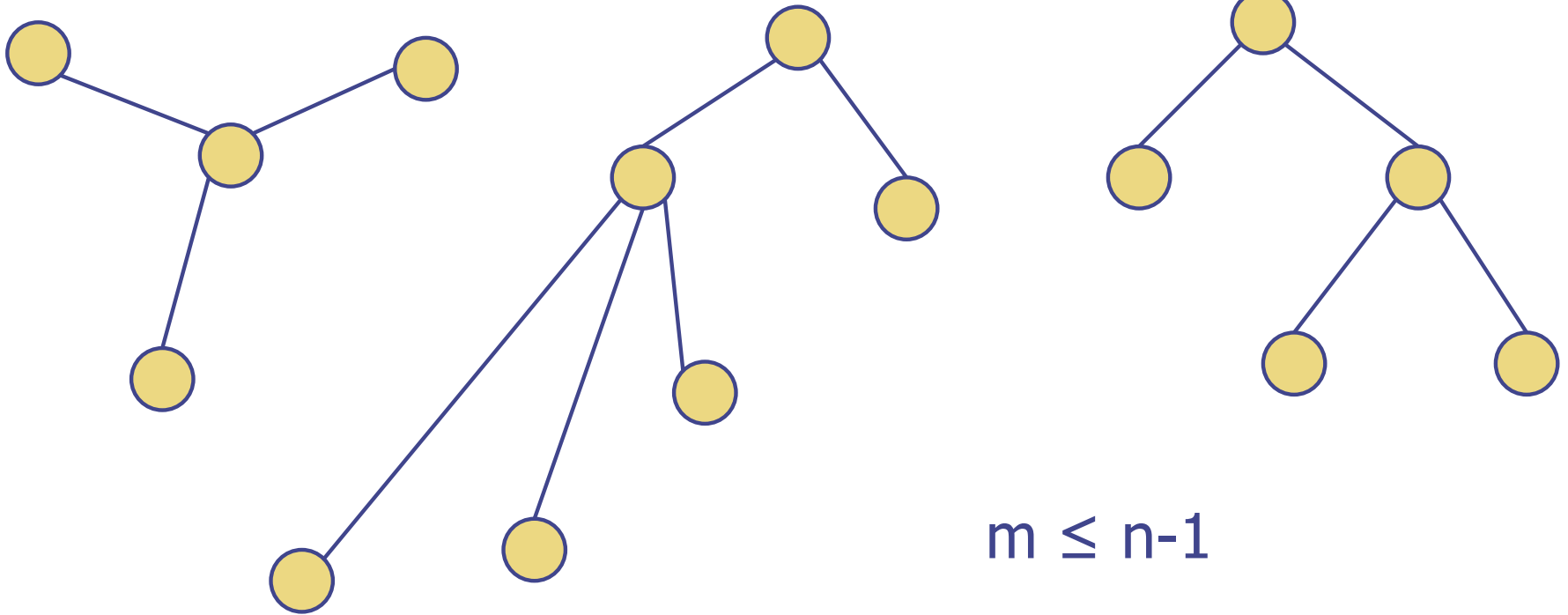
A tree is a graph without cycles.



$$m = n - 1$$

Forest

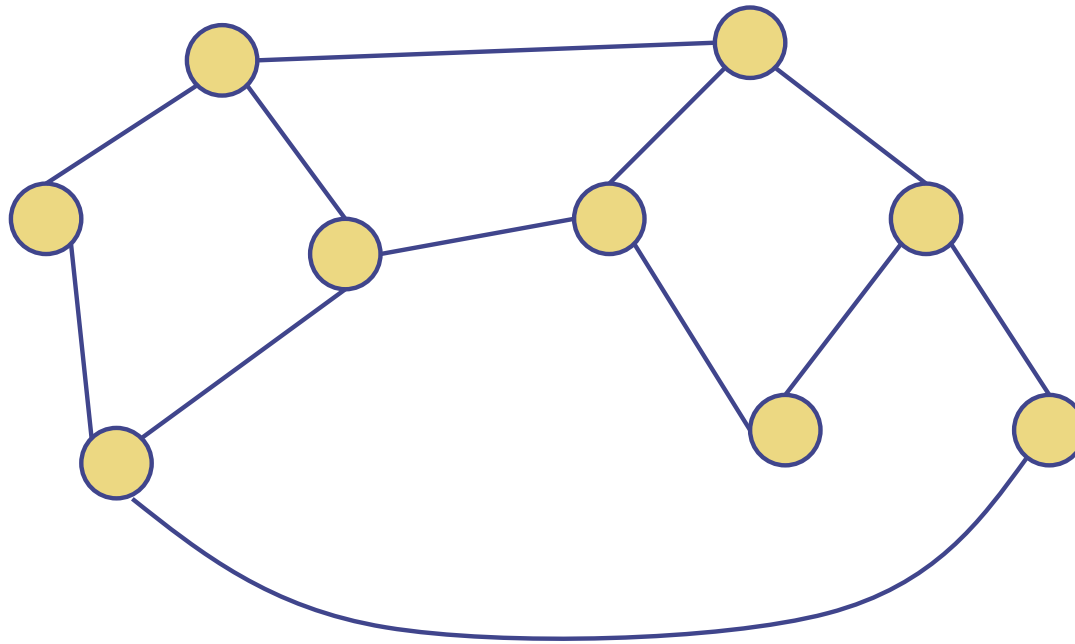
A forest is a set of trees.



$$m \leq n-1$$

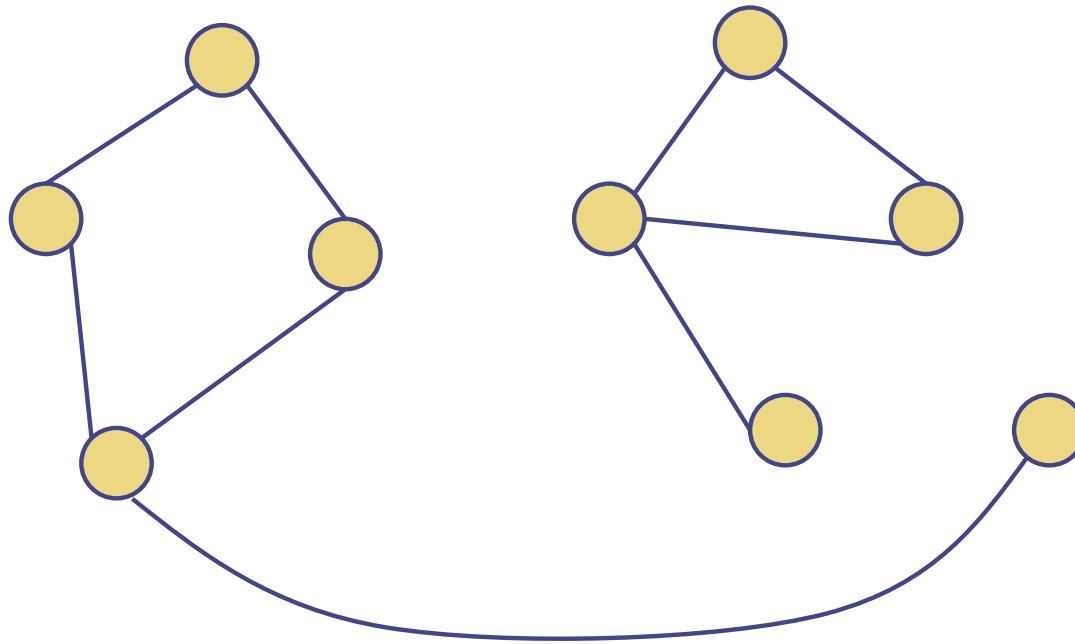
Subgraph

A subgraph is a subset of vertices and edges that forms a graph.



Connected Component

A connected component is a **maximal** connected subgraph.



How many connected components does this graph have?

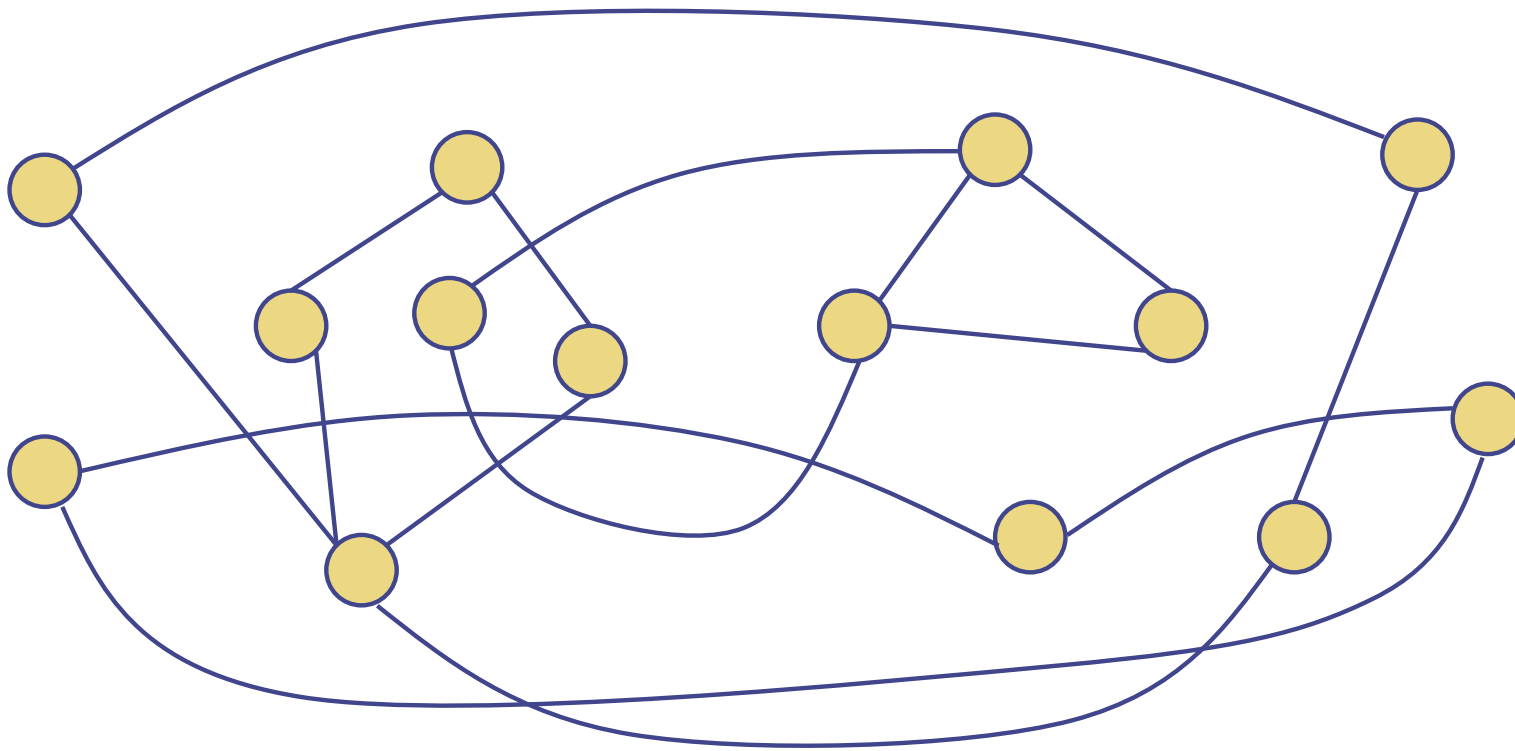
$$G = (V, E)$$

$$V = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14\}$$

$$E = \{(1, 6), (1, 12), (2, 11), (3, 4), (3, 7), (4, 6), (5, 8), \\ (5, 9), (6, 7), (6, 13), (8, 9), (8, 10), (9, 10), (11, 14), (12, 13)\}$$

Connected Component

A connected component is a **maximal** connected subgraph.



How many connected components? 3

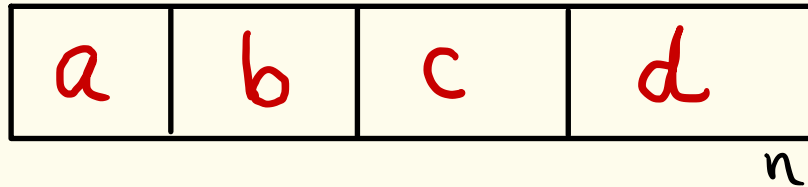
Graph ADT

- numVertices(): number of vertices of the graph
- getEdge(u,v): returns the edge between vertices u and v
- opposite(v,e): returns the vertex other than v that is incident on e
- insertVertex(x): creates and returns a new vertex storing value x
- insertEdge(u,v,x): creates an edge between u and v storing value x
- removeVertex(v): removes vertex v and all edges incident on it
- removeEdge(e): removes edge e
- areAdjacent(u,v): returns true if u and v are adjacent; false otherwise
- incidentEdges(u): returns an iterator of all edges incident on vertex u.

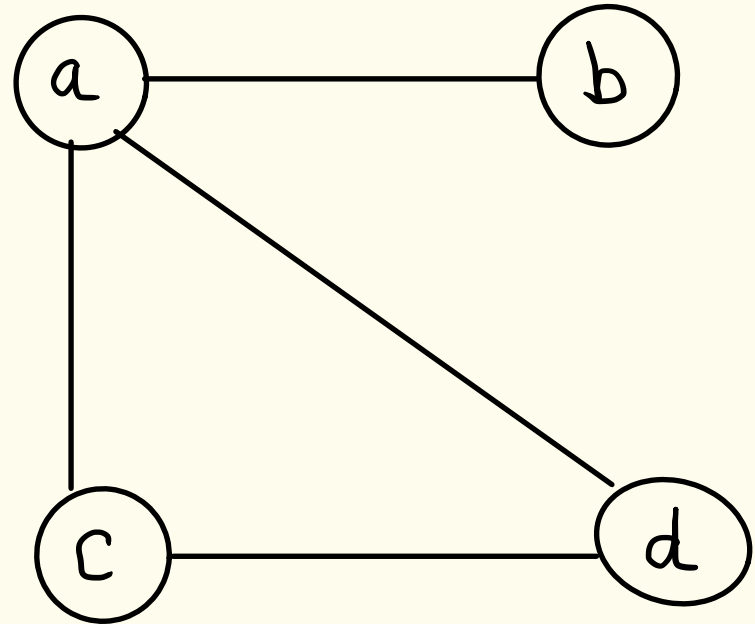
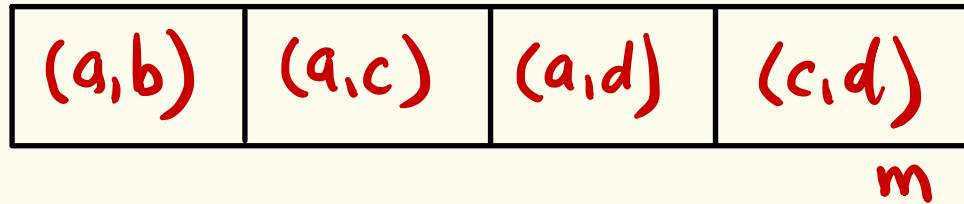
Data Structures to Store Graphs

Edge List

V



E



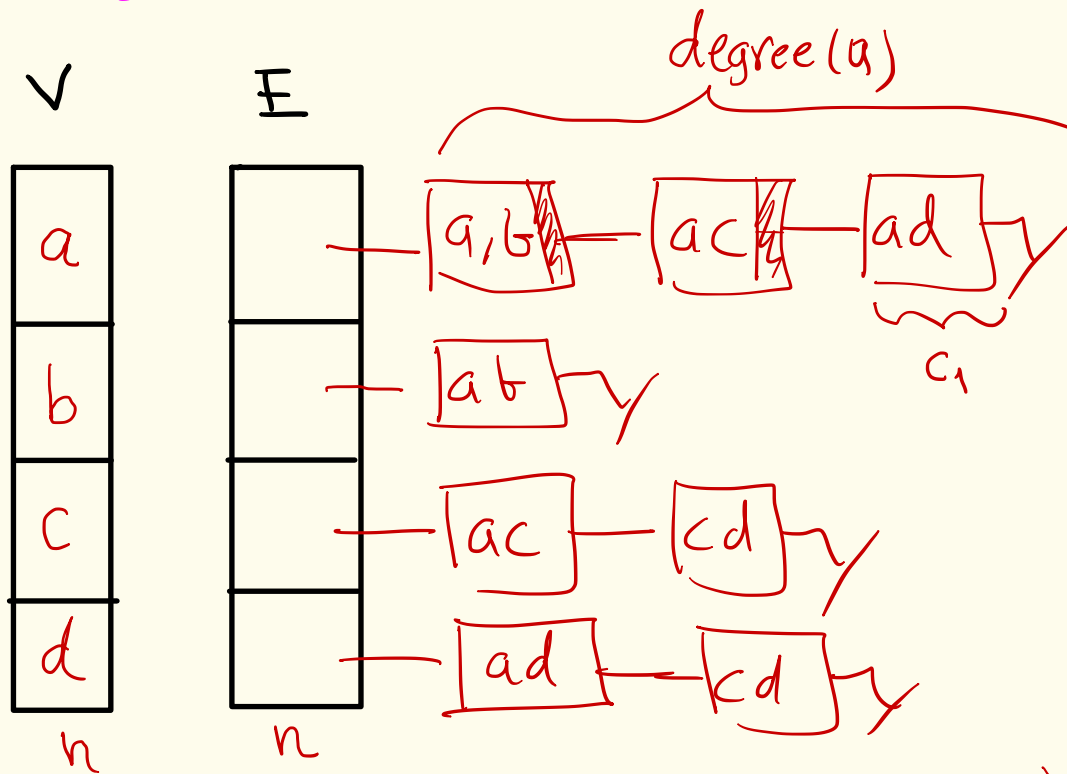
Space: $E_1 + E_2 + \dots + E_n$ is $O(n+m)$

Are Adjacent (u,v) : $O(m)$

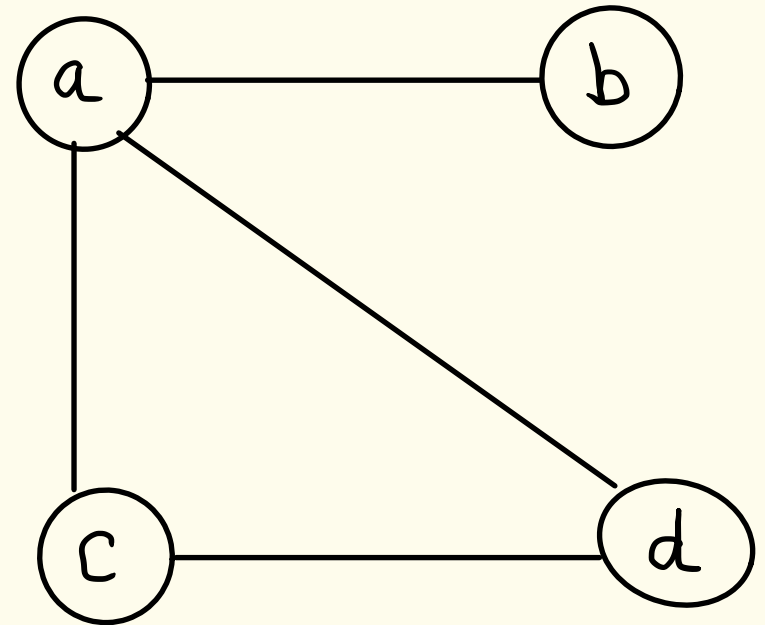
Incident Edges (u) : $O(m)$

Data Structures to Store Graphs

Adjacency List



Sparse graphs
("few" edges)



Space: $c_2 n + c_1 (2m)$ is $O(n + m)$

Are adjacent (u, v) : $O(\min \{\text{degree}(u), \text{degree}(v)\})$

Incident Edges (u) : $O(\text{degree}(u))$

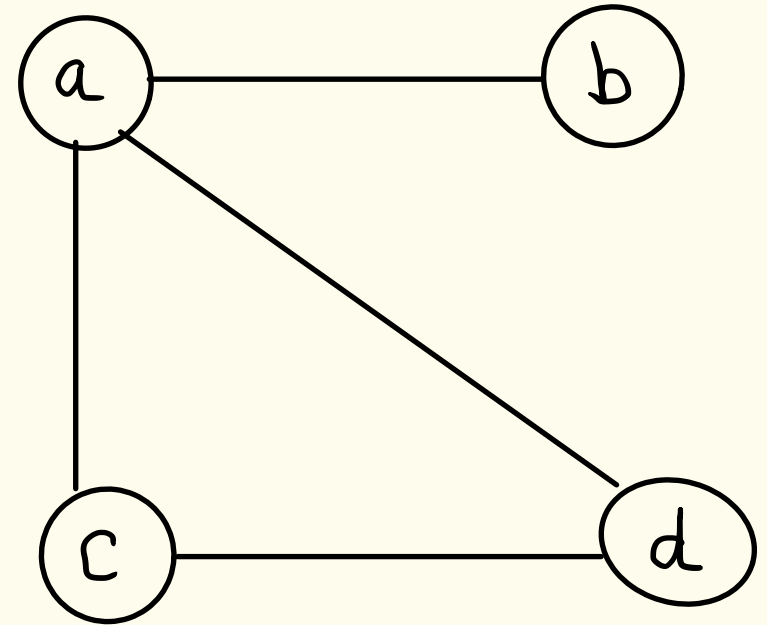
Data Structures to Store Graphs

Adjacency Matrix

V	E			
	a	b	c	d
a		(a,b)	(a,c)	(a,d)
b	(b,a)			
c	(c,a)			(c,d)
d	(d,a)		(d,c)	

$O(u)$ $O(n \times n)$

Dense graphs
("Many edges")



Space: $O(n^2)$

Are Adjacent(u,v): $O(1)$

Incident Edges(u): $O(n)$

Performance

▪ n vertices, m edges	Edge List	Adjacency List	Adjacency Matrix
Space	$O(n + m)$	$O(n + m)$	$O(n^2)$
incidentEdges(v)	$O(m)$	$O(\deg(v))$	$O(n)$
areAdjacent(v, w)	$O(m)$	$O(\min\{\deg(v), \deg(w)\})$	$O(1)$
insertVertex(o)	$O(1)$	$O(1)$	$O(n^2)$
insertEdge(v, w, o)	$O(1)$	$O(1)$	$O(1)$
removeVertex(v)	$O(m)$	$O(\deg(v))$	$O(n^2)$
removeEdge(v, w)	$O(m)$	$O(\deg(u) + \deg(v))$	$O(1)$