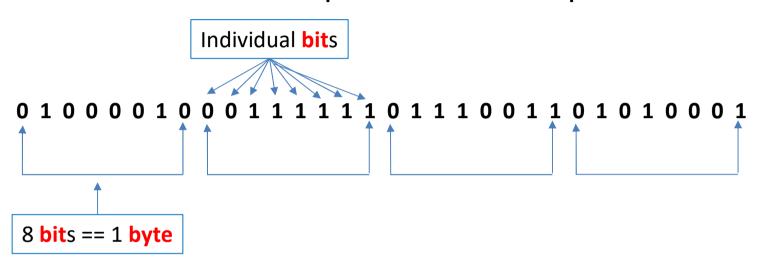
CS 2211 Systems Programming

Part Four:

Memory Maps

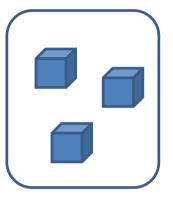
computer stores ons and offs the ons are 1 and offs are 0

BUT! what data does this sequence of ons and offs represent ????



Direct translation of systems:

DECIMAL (B	ASF 1	0) BINARY (B <mark>ASE 2</mark>)
	0	0
	1	1
	2	10
	3	11
	4	100
	5	101
	6	110
	7	111
	8	1000
	9	1001
	10	1010
	11	1011
1	000	1111101000



3 boxes

3

11 boxes

11

<u>128</u> <u>64</u> <u>32</u> <u>16</u> <u>8</u> <u>4</u> <u>2</u> <u>1</u>

with this, you can decode any binary number (up to 255 [28] decimal)

10010010 = ?

Computer

Computer

<u>128</u> <u>64</u> <u>32</u> <u>16</u> <u>8</u> <u>4</u> <u>2</u> <u>1</u>

with this, you can decode any binary number (up to 255 [28] decimal)

10010010 = **146**

$$= 128 + 0 + 0 + 16 + 0 + 0 + 2 + 0$$

= 146

American Standard Code for Information Interchange (ASCII)

Binary representation of characters:

Dec	Symbol	Binary	Dec	Symbol	Binary
65	A	0100 0001	83	S	0101 0011
66	В	01000001	84	T	0101 0100
67	C	01000011	85	Ü	0101 0101
68	D	0100 0100	86	V	0101 0110
69	E	01000101	87	W	0101 0111
70	F	0100 0110	88	X	0101 1000
71	G	0100 0111	89	Υ	0101 1001
72	Н	0100 1000	90	Z	0101 1010
73	I	0100 1001	91	[0101 1011
74	J	0100 1010	92	\	0101 1100
75	K	0100 1011	93]	0101 1101
76	L	0100 1100	94	^	0101 1110
77	М	0100 1101	95	_	0101 1111
78	Ν	0100 1110	96	×	01 0 0000
79	0	0100 1111	97	а	0110 0001
80	Р	0101 0000	98	b	0110 0010
81	Q	0101 0001	99	С	0110 0011
82	R	0101 0010	100	d	0110 0100

UNICODE: UTF-8 8 bits = 2^7 = 128 possible characters.

UNICODE (UTF-8 UTF-16 UTF-32)

Binary representation of characters:

Dec	Din	Hex	Char	Dog	Bin	Ном	Char	Dec	Din	Ном	Char	Dec	Din	Пох	Char
0	0000 0000	00	[NUL]	32	0010 0000	20	space	64	0100 0000	40	@	96	0110 0000	60	Char
1	0000 0000	01	[SOH]	33	0010 0000	21	space	65	0100 0000	41	A	97	0110 0000	61	a
2	0000 0001	02	[STX]	34	0010 0001	22		66	0100 0001	42	В	98	0110 0001	62	a b
3	0000 0010	03	[ETX]	35	0010 0010	23	#	67	0100 0010	43	C	99	0110 0010	63	C
4	0000 0011	04	[EOT]	36	0010 0011	24	\$	68	0100 0011	44	D	100	0110 0011	64	d
5	0000 0100	05	[ENO]	37	0010 0100	25	8	69	0100 0100	45	E		0110 0100	65	e e
6	0000 0101	06	[ACK]	38	0010 0101	26	£	70	0100 0101	46	F	102	0110 0101	66	f
7	0000 0110	07	[BEL]	39	0010 0110	27		71	0100 0110	47	G		0110 0110	67	100
8	0000 1000	08	[BS]	40	0010 0111	28	(72	0100 0111	48	н		0110 0111	68	g h
9	0000 1000	09	[TAB]	41	0010 1000	29)	73	0100 1000	49	I		0110 1000	69	n i
10	0000 1001	0A	[LF]	42	0010 1001	2 A	*	74	0100 1001	4.A	J		0110 1001	6A	j
11	0000 1010	0B	[VT]	43	0010 1010	2B	+	75	0100 1010	4B	K	107		6B) k
12	0000 1011	OC	[FF]	44	0010 1011	2C		76	0100 1011	4C	L		0110 1011	6C	1
13	0000 1100	OD	[CR]	45	0010 1100	2D	_	77	0100 1100	4D	м		0110 1100	6D	-
14	0000 1101	0E	[SO]	46	0010 1101	2E	-	78	0100 1101	4E	M N		0110 1101	6E	m n
15	0000 1110	0F	[SI]	47	0010 1110	2F	,	79	0100 1110	4E	0		0110 1110	6F	0
16	0000 1111	10	[DLE]	48	0010 1111	30	0	80	0100 1111	50	P		0111 0000	70	
17	0001 0000	11	[DC1]	49	0011 0000	31	1	81	0101 0000	51	0		0111 0000	71	p q
18	0001 0001	12	[DC1]	50	0011 0001	32	2	82	0101 0001	52	R		0111 0001	72	q r
19	0001 0010	13	[DC3]	51	0011 0010	33	3	83	0101 0010	53	S		0111 0010	73	s
20	0001 0011	14	[DC4]	52	0011 0011	34	4	84	0101 0011	54	T T	0.00077507.000	0111 0111	74	t.
21	0001 0100	15	[NAK]	53	0011 0100	35	5	85	0101 0100	55	U		0111 0100	75	u
22		16	[SYN]	54	0011 0101	36	6	86	0101 0101	56	v		0111 0101	76	v
23	0001 0110	17	[ETB]	55	0011 0111	37	7	87	0101 0110	57	w		0111 0111	77	w
24	0001 1000	18	[CAN]	56	0011 1000	38	8	88	0101 1000	58	x		0111 1000	78	x
25	0001 1000	19	[EM]	57	0011 1001	39	9	89	0101 1000	59	Y		0111 1001	79	у
26	0001 1001	1A	[SUB]	58	0011 1010	3A	:	90	0101 1001	5A	z		0111 1010	7A	y Z
27	0001 1010	1B	[ESC]	59	0011 1011	3B	;	91	0101 1010	5B	1		0111 1011	7B	{
28	0001 1011	1C	[FS]	60	0011 1100	3C	<	92	0101 1011	5C	1		0111 1100	7C	ì
29	0001 1101	1D	[GS]	61	0011 1101	3D	=	93	0101 1101	5D	ì		0111 1101	7D	i
30	0001 1110	1E	[RS]	62	0011 1110	3E	>	94	0101 1110	5E	,		0111 1110	7E	~
31	0001 1111		[US]	63	0011 1111		?	95	0101 1111				0111 1111	7F	[DEL]
21	0001 1111		[co]	100	0011 1111	JE		33	0101 1111	J.E	_	127	VIII 1111	, E	التجدا

UNICODE is a computing industry standard for the consistent encoding, representation, and handling of text expressed in most of the world's writing systems.

UNICODE: UTF-16 16 bits = 2^{15} = 1,112,064 possible characters **UTF-32** 32 bits = 2^{32} = 2,147,483,648 possible characters

Binary Representation

END OF PART 1

+25:0011001

Magnitude Form

- unsigned (positive only) whole numbers

Sign-Magnitude Form

- signed (positive and negative) whole numbers -25: 1011001.

One's Compliment Form (1's compliment) たんろ しいいい

- negative represented by the compliment (flipping bits) of the positive number

Two's Compliment Form (2's compliment) えんなう。 じしい キャルラニ (を含まれ)

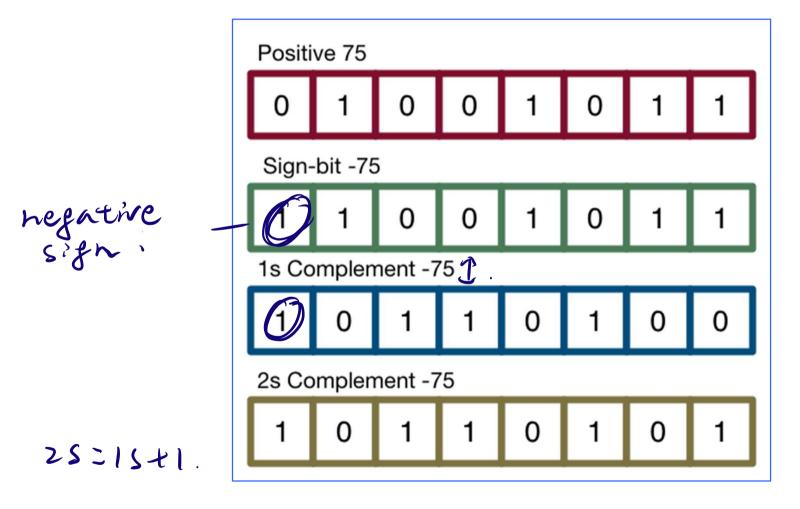
- negative represented by adding the values of 1 to the compliment of the positive number

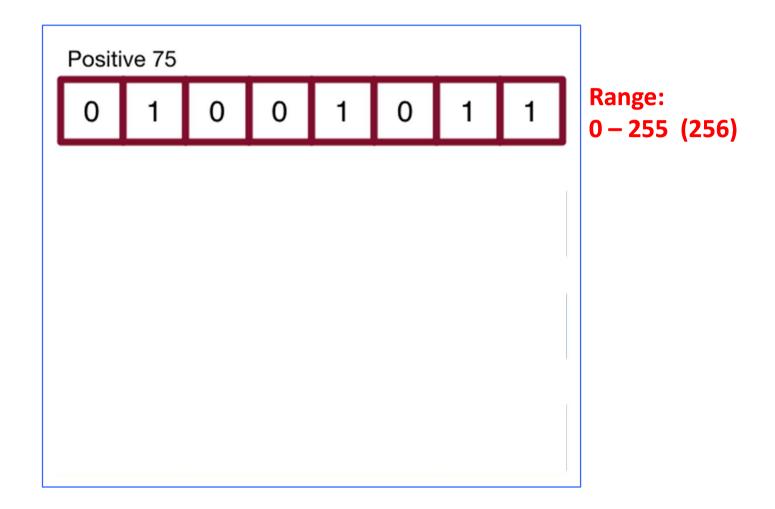
Single Precision Floating Point Form

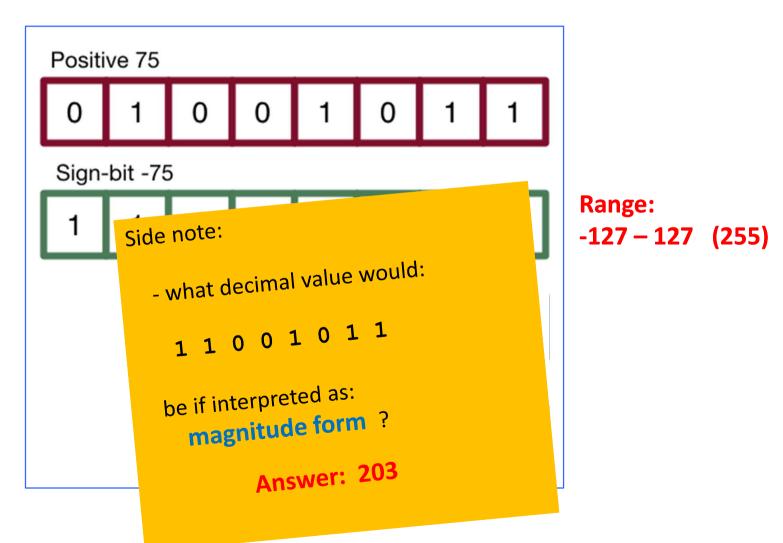
- fixed point model using Scientific Notation to represent values (32 bit form)

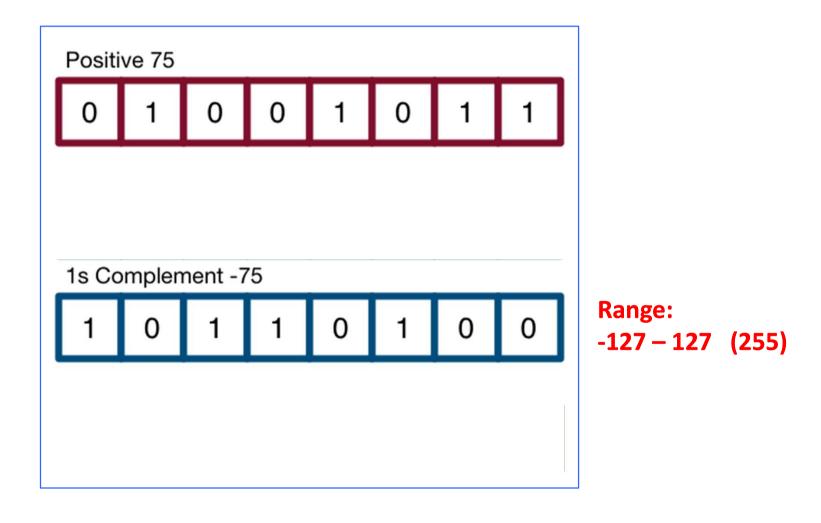
Double Precision Floating Point Form

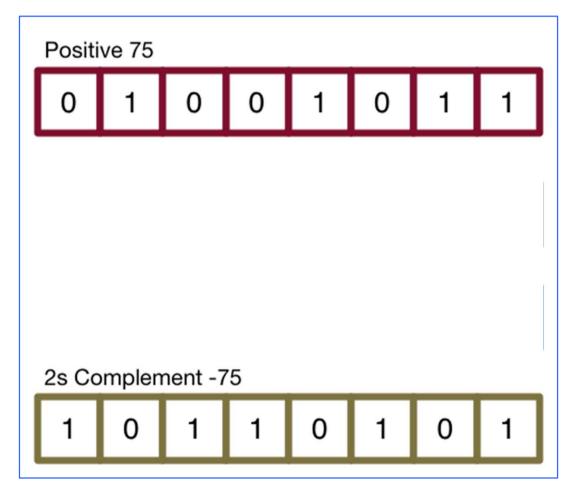
- fixed point model using Scientific Notation to represent values (64 bit form)











Range: -128 – 127 (256)

Binary Representation

END OF PART 2

Fixed Point Form

method for storing real (fractional) numbers

for example: **18.375**

18 -> whole number

375 -> fractional part

- can use negative powers of two

remember: $2^3 -> 8$ $2^{-3} -> 0.125$

 2^4 2^3 2^2 2^1 2^0 2^{-1} 2^{-2} 2^{-3} 16 8 4 2 1 0.5 0.25 0.125

1 0 0 1 0 0 1 1 -> 18.375

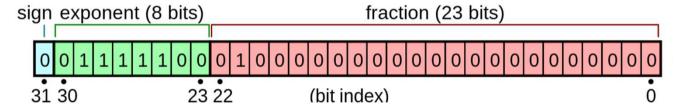
can be written as: **10010.011** (this is FIXED POINT MODEL) limited range of values due to forced position

Floating Point Form

```
remember (in Decimal (base 10))
           18.375 = 1.8375 \times 10^{1} -> move it one (1) place
          123.456 = 1.23456 \times 10^2 -> move it two (2) places
  we can do the same thing in Binary:
      10010.011 = 1.0010011 \times 2^4 \rightarrow \text{move it four (4) places}
this allows for FLOATING POINT representation
 => ignore the 2 (always present – assumed)
 => ignore the leading mantissa (number to the left of decimal)
 just need to store the sign (- or +), the fractional part and the exponent
 8 bits does not give enough to play with: assume 32 bits (4 bytes)
   - store exponent using 8 bits - store fraction using 23 bits
```

Single Precision Floating Point Form

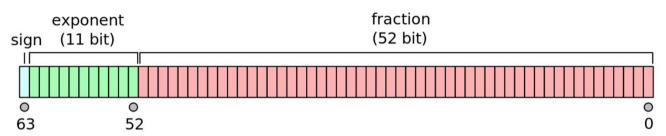
bit: 1 8 23 sign exponent (e) fraction (f) position 31 30 ... 23 22 ... 0



Double Precision Floating Point Form

bit: 1 11 53 sign exponent (e) fraction (f)

position 63 62 ... 53 52 ... 0



Binary Representation

END OF PART 3

ALL DATA IS STORED AS A TYPE

characters - 1 byte
short - 2 bytes
int - 4 bytes
float - 4 bytes
double -8 bytes

C's built-in data types that are similar to ones in Java:

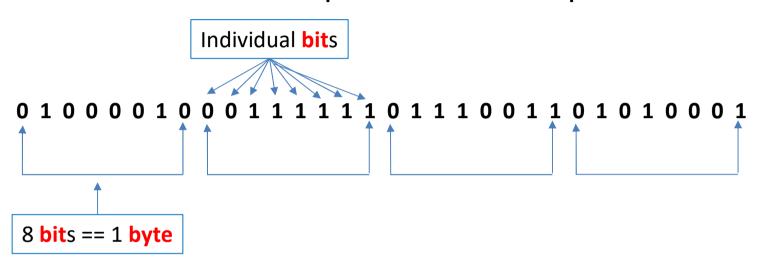
Syntax	Name	Java's counterpart	use					
char character		byte	Stores an ASCII code (character) or it can also store a very short integer (-128127)					
short short integer		short	uses 2 byte memory, value between -32768 and 32767					
int	ordinary integer	int	uses 4 byte memory, value between -2147483648 and 2147483647					
long integer		long	uses 8 bytes memory, value between -9223372036854775808 and 9223372036854775807					
float single precision float		float	uses 4 byte memory, absolute value between 1.4E-45 and 3.4E38					
double double precision float		double	uses 8 byte memory, absolute value between 4.9E-324 and 1.8E308					
Bool	boolean	boolean	true (1) or false (0)					

C's built-in data types that do not have an equivalent in Java:

Syntax	Name	use
unsigned char	Unsigned character	Very short positive integer (0255)
unsigned short	Unsigned short integer	uses 2 byte memory, value between 0 and 65535
unsigned int	Unsigned ordinary integer	uses 4 byte memory, value between 0 and 4294967295
unsigned long	Unsigned long integer	uses 8 bytes memory, value between 0 and 18446744073709551615
*	Reference type	Contains a memory address (usually 4 bytes, but 64 bits machines will use 8 bytes)

computer stores ons and offs the ons are 1 and offs are 0

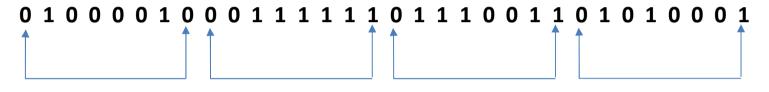
BUT! what data does this sequence of ons and offs represent ????



computer stores ons and offs the ons are 1 and offs are 0

BUT! what data does this sequence of ons and offs represent ????

4 bytes [1 byte = 8 bits]



B : ? : S : Q

66 : 63 : 115 : 81

65 A.

char

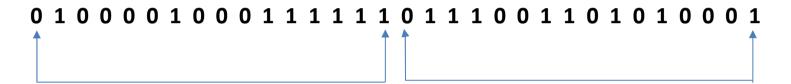
char

*very short integer

computer stores ons and offs the ons are 1 and offs are 0

BUT! what data does this sequence of ons and offs represent ????

2 words [1 word = 16 bits][1 word = 2 bytes]



16,959 : 259,521

 short

computer stores ons and offs the ons are 1 and offs are 0

BUT! what data does this sequence of ons and offs represent ????

1 double word [1 double word = 32 bits][1 double word = 4 bytes]

 $0\;1\;0\;0\;0\;0\;1\;0\;0\;0\;1\;1\;1\;1\;1\;1\;0\;1\;1\;1\;0\;0\;1\;1\;0\;1\;0\;1\;0\;0\;0\;1$

1,111,454,545

47.862613677978515625

integer

float

computer stores ons and offs the ons are 1 and offs are 0

BUT! what data does this sequence of ons and offs represent ????

1 quad word [1 quad word = 64 bits] [1 quad word = 8 bytes]

-3,203,931,608,977,542,319

long

-2.6437373164653993971898000075 E⁹⁴

double

Binary Representation

END OF PART 4

SIMPLE C PROGRAM

```
#include <stdio.h>
                                   Variable declaration
int main(int argc, char *argv[])
{
  char a; /* 1 byte
                          */
   int b; /* 4 bytes */
                                      OUTPUT:
  float c; /* 4 bytes */
                                       1st value of a is: K
  double d; /* 8 bytes */
                                       2nd value of b is: 37
                                       3rd value of c is: 2.50000000
   a = 'K';
                                       4th value of d is: 75.50000000
  b = 37;
   c = 2.5;
   d = 75.3;
                         Variable definition
  printf( "1st value of a is : %c \n" , a );
  printf( "2nd value of b is : %d \n" , b );
  printf( "3rd value of c is : %f \n" , c );
  printf( "4rd value of d is : %lf \n" , d );
   return 0 ;
}
```

ALL DATA IS STORED AS A TYPE

```
characters - 1 byte
short - 2 bytes
int - 4 bytes
float - 4 bytes
double -8 bytes
```

For a computer to perform an arithmetic operation:

- the operands must usually be of the same size i.e. (the same number of bits)

and

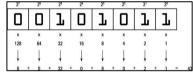
- be stored in the same way.
 - i.e. (int and float are both 4 bytes, but different usage of the 1's and 0's)

ALL DATA IS STORED AS A TYPE

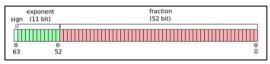
MIXED USE:

for example: adding an

int (4 bytes – double precision)



double (8 bytes floating point)



computer MUST convert to one type of data representation.

in C, the computer will always convert to the more complex representation (variable promotion)

ALL DATA IS STORED AS A TYPE

When operands of different types are mixed in expressions, the C compiler may have to generate instructions that change the types of some operands so that hardware will be able to evaluate the expression.

IMPLICIT CONVERSION

If we add a 16-bit short and a 32-bit int, the compiler will arrange for the short value to be converted to 32 bits.

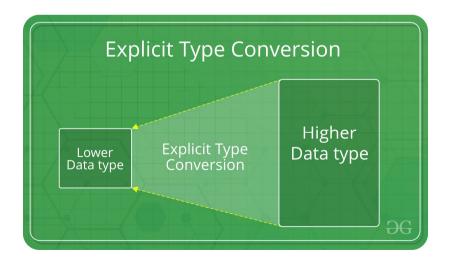
If we add an int and a float, the compiler will arrange for the int to be converted to float format.

IMPLICIT CONVERSION

Generally takes place when in an expression more than one data type is present.

In such condition type conversion (type promotion) takes place to avoid lose of data.

All the data types of the variables are upgraded to the data type of the variable with largest data type.



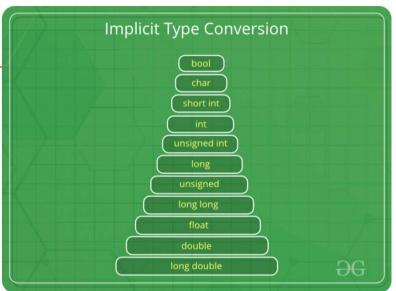
IMPLICIT CONVERSION

```
// An example of implicit conversion
#include<stdio.h>
int main()
{
   int x = 10; // integer x
   char y = 'a'; // character c

   // y implicitly converted to int.
   // value of 'a' is 97
   x = x + y;

   // x is implicitly converted to float
   float z = x + 1.0;

   printf("x = %d, z = %f", x, z);
   return 0;
}
```



implicit conversion is **automatic** it requires no user intervention

```
char ta = 'b';
long tb = 343437;
long tc;
double td;

tc = ta + tb;

td = tc / ta;

printf("value of ta: 'd\n", ta);
printf("value of tb: %d\n", tb);
printf("value of tc: %d\n", tc);
printf("value of td: %lf\n", td);
printf("value of td: %lf\n", td);
printf("\n\n");
```

IMPLICIT CONVERSION

value of ta: b value of ta: 98

value of tb: 343437 value of tc: 343535

value of td: 3505.00000000000

TOP to BOTTOM (anything above will converted to the type below it)

IMPLICIT CONVERSION

```
char c;
short int s:
int i;
unsigned int u;
long int 1;
unsigned long int ul;
float f;
                                  char/short int
double d;
long double ld;
                                                      * /
i = i + c; /* c is converted to int
i = i + s; /* s is converted to int
                                                      * /
u = u + i;  /* i is converted to unsigned int
                                                      * /
l = l + u; /* u is converted to long int
                                                      * /
ul = ul + l; /* l is converted to unsigned long int
                                                      * /
f = f + ul; /* ul is converted to float
                                                      * /
d = d + f; /* f is converted to double
                                                      * /
ld = ld + d; /* d is converted to long double
                                                      * /
```

ALL DATA IS STORED AS A TYPE

When operands of different types are mixed in expressions, the C compiler may have to generate instructions that change the types of some operands so that hardware will be able to evaluate the expression.

EXPLICIT CONVERSION

Basically, the programmer forces an expression to be of a specific type.

Explicit type conversion is also called type casting.

The general format of explicit type conversion is as follows:

(data_type)(expression);

```
// C program to demonstrate explicit type casting
#include<stdio.h>

int main()
{
    double x = 1.2;

    // Explicit conversion from double to int
    int sum = (int)x + 1;

    printf("sum = %d", sum);

    return 0;
}
```

EXPLICIT CONVERSION

notice we are *forcing* the computer to cast a higher complexity type (double) to a lower complexity (int)

C regards (*type-name*) as a **unary operator**. Unary operators have higher **precedence** than binary operators,

so the compiler interprets

```
(float) dividend / divisor
as
((float) dividend) / divisor
```

EXPLICIT CONVERSION

Casts are sometimes necessary to avoid overflow:

```
long i;
int j = 1000;
i = j * j;    /* overflow may occur */
```

Using a cast avoids the problem:

```
i = (long) j * j;
```

The statement

```
i = (long) (j * j); /*** WRONG ***/
```

wouldn't work, since the overflow would already have occurred by the time of the cast.

Binary Representation

END OF PART 5

TYPE DEFINITIONS in C

The #define directive can be used to create a "Boolean type" macro:

#define BOOL int

There's a better way using a feature known as a *type definition*:

typedef int BOOL;

Bool can now be used in the same way as the built-in type names. Example:

BOOL flag; /* same as int flag; */

typedef (known data type) (alias to be used instead of) is nothing more than creating an 'alias'

typedeffloat Dollars // Dollars now is an alias for float
Dollars cash_in, cash_out;
// is more informative than float cash_in, cash_out;

TYPE DEFINITIONS in C

Type definitions can also make a program easier to modify.

To redefine Dollars as double, only the type definition need be changed:

```
typedef double Dollars;
```

Without the type definition, we would need to locate all float variables that store dollar amounts and change their declarations.

Type definitions are an important tool for writing portable programs.

One of the problems with moving a program from one computer to another is that types may have different ranges on different machines.

If i is an int variable, an assignment like

```
i = 100000;
```

is fine on a machine with 32-bit integers, but fails on a machine with 16-bit integers.

Binary Representation

END OF PART 6