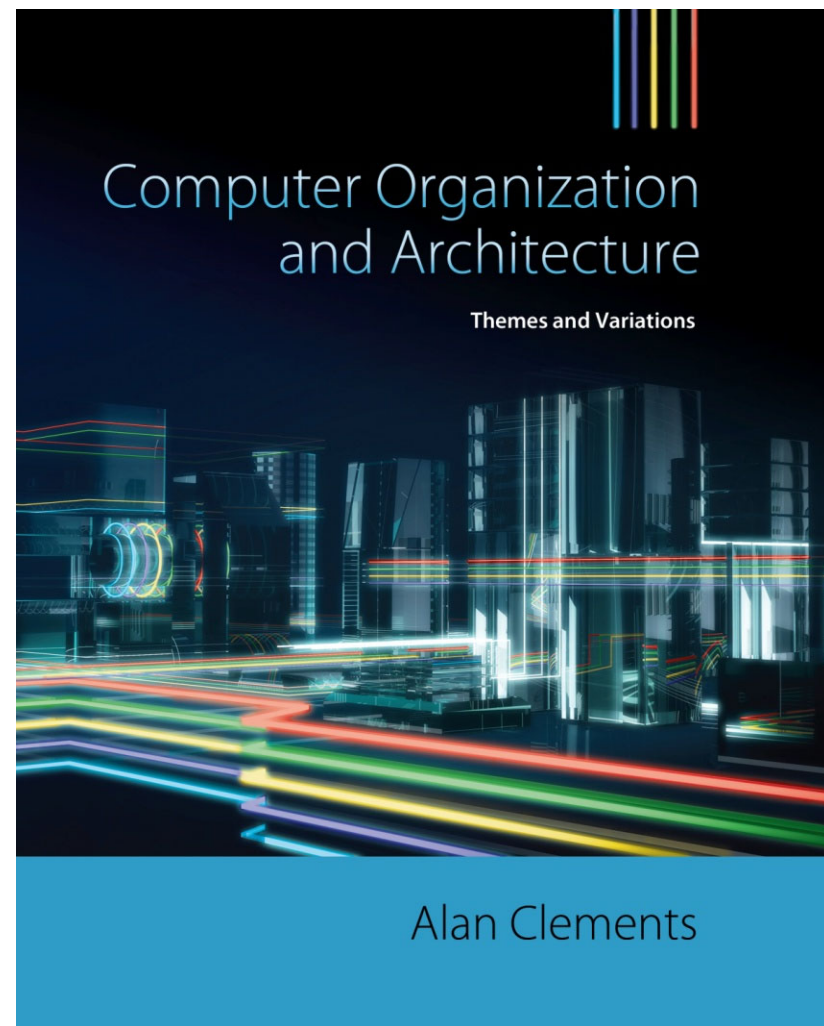


Part 1

CHAPTER 2

Computer Arithmetic and Digital Logic



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Bits and Bytes

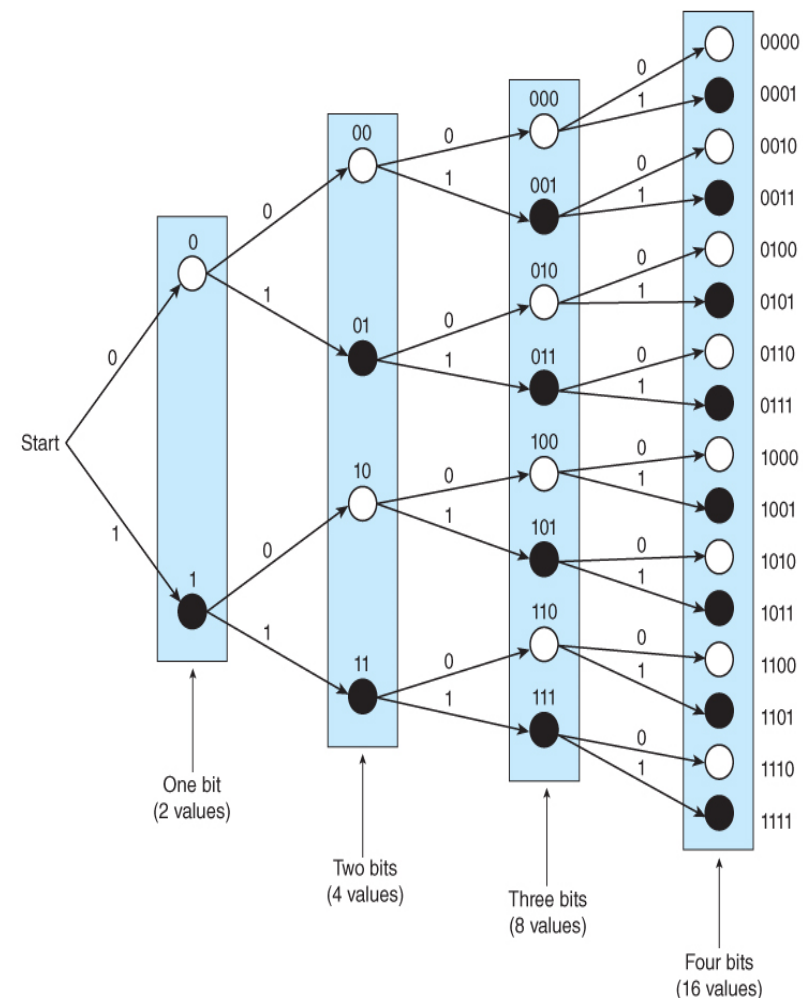
- ❑ In digital computers, data is represented using *Bits* (*Binary digiT*)s. \Rightarrow store binary value.
- ❑ A *bit* has *two values* that we call 0 and 1, low and high, false and true, clear and set, and so on.
- ❑ Using bits, it is easy to represent real-world quantities.
 - Sound and images can easily be converted to bits.
 - Strings of bits can be converted back to sound or images.
- ❑ We call a *unit of 8 bits* a *byte*. This is a *convention*.

Bit Patterns

- ❑ One bit can have two values, 0 or 1.
- ❑ Two bits can have four values, 00, 01, 10, 11.
- ❑ Each time you introduce a bit, you double the number of possible combinations, as Figure 2.1 demonstrates.

- ❑ 3 bits can have (2^3) 8 values
- ❑ 4 bits can have (2^4) 16 values
- ❑ 5 bits can have (2^5) 32 values
- ❑ 6 bits can have (2^6) 64 values
- ❑ 7 bits can have (2^7) 128 values
- ❑ 8 bits can have (2^8) 256 values
- ❑ 9 bits can have (2^9) 512 values
- ❑ 10 bits can have (2^{10}) 1024 values
- ❑ 10 bits can have (2^{10}) 1 K values $\hookrightarrow 1024$
- ❑ 11 bits can have (2^{11}) 2 K values
- ❑ 12 bits can have (2^{12}) 4 K values
- ❑ 13 bits can have (2^{13}) 8 K values
- ❑ 20 bits can have (2^{20}) 1 Mega values $\hookrightarrow 1024^2$
- ❑ 23 bits can have (2^{23}) 8 Mega values
- ❑ 30 bits can have (2^{30}) 1 Giga values $\hookrightarrow 1024^3$
- ❑ 32 bits can have (2^{32}) 4 Giga values
- ❑ 34 bits can have (2^{34}) 16 Giga values

FIGURE 2.1 The binary tree



Bit Patterns

- ❑ One of the **first** quantities **to be represented** *digitally* were characters (*letters*, *numbers*, and *symbols*).
- ❑ This was necessary in order to transmit text across the networks that were developed as a result of the invention of the telegraph.
- ❑ This led to a **standard code** for characters called **ASCII** (*American Standard Code for Information Interchange*)
 - **7-bit code** *extended ASCII : 8-bit code.*
 - **representing** up to $2^7 = 128$ characters of **Latin alphabet**.
- ❑ Today, the **16-bit unicode** has been devised to represent a much greater range of characters including *non-Latin alphabets*.

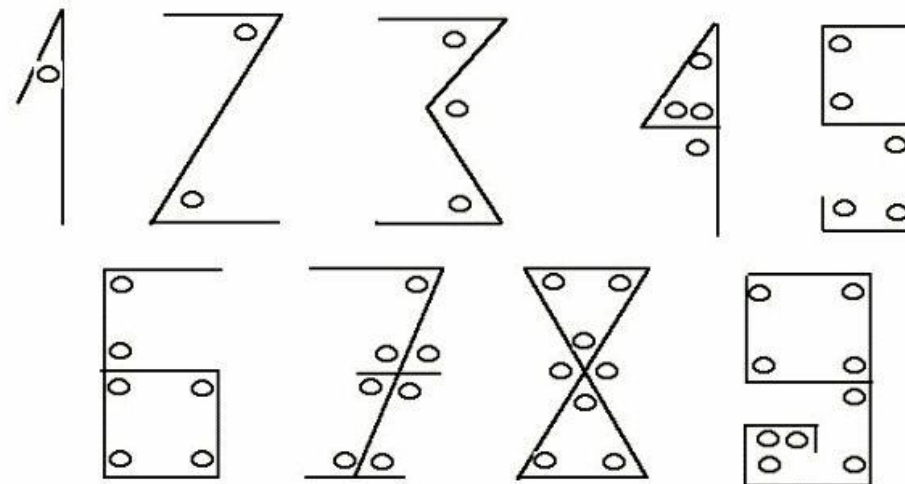
any A-Z : add 30 to uppercase
to get lowercase

ASCII Code

code	value	code	value	code	value	code	value	code	value	code	value	code	value	code	value
0	NUL	1	SOH	2	STX	3	ETX	4	EOT	5	ENQ	6	ACK	7	BEL
8	BS	9	HT	10	NL	11	VT	12	NP	13	CR	14	SO	15	SI
16	DLE	17	DC1	18	DC2	19	DC3	20	DC4	21	NAK	22	SYN	23	ETB
24	CAN	25	EM	26	SUB	27	ESC	28	FS	29	GS	30	RS	31	US
32	SP	33	!	34	"	35	#	36	\$	37	%	38	&	39	'
40	(41)	42	*	43	+	44	,	45	-	46	.	47	/
48	0	49	1	50	2	51	3	52	4	53	5	54	6	55	7
56	8	57	9	58	:	59	;	60	<	61	=	62	>	63	?
64	@	65	A	66	B	67	C	68	D	69	E	70	F	71	G
72	H	73	I	74	J	75	K	76	L	77	M	78	N	79	O
80	P	81	Q	82	R	83	S	84	T	85	U	86	V	87	W
88	X	89	Y	90	Z	91	[92	\	93]	94	^	95	_
96	`	97	a	98	b	99	c	100	d	101	e	102	f	103	g
104	h	105	i	106	j	107	k	108	l	109	m	110	n	111	o
112	p	113	q	114	r	115	s	116	t	117	u	118	v	119	w
120	x	121	y	122	z	123	{	124		125	}	126	~	127	DEL

Numbers and Binary Arithmetic

- ❑ One of the great advances in history was the **move away from Roman numerals** (I, II, III, IV, V, VI, VII, VIII, IX, X, ..., L, ..., C, ..., D, ..., M, ... where **I=1**, **V=5**, **X=10**, **L=50**, **C=100**, **D=500**, and **M=1000**) to the **Hindu-Arabic** notation (**0, 1, 2, 3, 4, 5, 6, 7, 8, 9**) that we use today.
- ❑ Invented by **Muhammad Musa Al-Khwarizmi** (Born 780—Died 850)
- ❑ **Numerals** represent the number of angles in a digit



Numbers and Binary Arithmetic

- ❑ Arithmetic calculations are remarkably *difficult* using *Roman numerals*, but they are far *simpler* using *Hindu-Arabic positional notation* system.
- ❑ In the positional notation system, an *n*-digit integer number, *N*, is written as a sequence of digits in the form

$$a_{n-1} a_{n-2} \dots a_i \dots a_1 a_0$$

- For example, when *N* is 278, then $a_2 = 2$, $a_1 = 7$, and $a_0 = 8$.
- ❑ The value of this number, expressed in the positional notation system in the *base b*, is defined as

$$N = a_{n-1} \times b^{n-1} \dots + a_1 \times b^1 + a_0 \times b^0$$

i.e.,

$$N = 2 \times 10^2 + 7 \times 10^1 + 8 \times 10^0 = 200 + 70 + 8 = 278$$

Numbers and Binary Arithmetic

- ❑ The positional notation system can be extended to express real values by using a *radix point* to separate the integer and fractional part, e.g.,
 - decimal point in base ten arithmetic or
 - binary point in base two arithmetic.
- ❑ A real value in decimal arithmetic is written in the form 1234.567.

integer ↓ ↓ *fraction.*
- ❑ To generalize, if we have
 - n digits to the left of the radix point and
 - m digits to the right of the radix point,
 we can write the number as $a_{n-1} a_{n-2} \dots a_i \dots a_1 a_0 . a_{-1} a_{-2} \dots a_{-m}$
- ❑ The value of this number, expressed in the positional notation system in the base b , is defined as

$$N = a_{n-1} \times b^{n-1} \dots + a_1 \times b^1 + a_0 \times b^0 + a_{-1} \times b^{-1} + a_{-2} \times b^{-2} \dots + a_{-m} \times b^{-m}$$

$$= \sum_{i=-m}^{n-1} a_i b^i$$

concept.
 \Rightarrow *example in tutorial.*

Warning!

- ❑ Any *integer* number can be *accurately* converted from a base to the other without any error.
- ❑ Some *fractions* that can be represented in a base cannot be represented in another base
 - for example, 0.1_{10} cannot be *accurately* converted into a binary form.

ex. Cannot get $2/3$ accurately in base 10.

options: $\left[\begin{array}{l} \text{(1) Round} \rightarrow 0.66667 \\ \text{(2) Truncate (cut)} \rightarrow 0.66666 \end{array} \right.$

Binary Arithmetic

□ These tables cover the fundamental arithmetic operations.

Produces
sum & carry

Addition

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

$$1 + 1 = 0 \text{ (carry 1)}$$

Produces
difference & borrow

Subtraction

$$0 - 0 = 0$$

$$0 - 1 = 1 \text{ (borrow 1)}$$

$$1 - 0 = 1$$

$$1 - 1 = 0$$

Multiplication

$$0 \times 0 = 0$$

$$0 \times 1 = 0$$

$$1 \times 0 = 0$$

$$1 \times 1 = 1$$

fundamental.
1 bit.
gah damn.
shaah. ooh lala.
music taste
is something.

Binary Arithmetic

→ similar to modular arithmetic.

□ These tables cover the fundamental arithmetic operations.

Produces
sum & carry

Addition

$$0 + 0 = 0 \text{ (carry 0)}$$

$$0 + 1 = 1 \text{ (carry 0)}$$

$$1 + 0 = 1 \text{ (carry 0)}$$

$$1 + 1 = 0 \text{ (carry 1)}$$

Produces
difference & borrow

Subtraction

$$0 - 0 = 0 \text{ (borrow 0)}$$

$$0 - 1 = 1 \text{ (borrow 1)}$$

$$1 - 0 = 1 \text{ (borrow 0)}$$

$$1 - 1 = 0 \text{ (borrow 0)}$$

Multiplication

$$0 \times 0 = 0$$

$$0 \times 1 = 0$$

$$1 \times 0 = 0$$

$$1 \times 1 = 1$$

* needs to be memorized.

Addition (three bits)

$$0 + 0 + 0 = 0 \text{ (carry 0)}$$

$$0 + 0 + 1 = 1 \text{ (carry 0)}$$

$$0 + 1 + 0 = 1 \text{ (carry 0)}$$

$$0 + 1 + 1 = 0 \text{ (carry 1)}$$

$$1 + 0 + 0 = 1 \text{ (carry 0)}$$

$$1 + 0 + 1 = 0 \text{ (carry 1)}$$

$$1 + 1 + 0 = 0 \text{ (carry 1)}$$

$$1 + 1 + 1 = 1 \text{ (carry 1)}$$

Subtraction (three bits)

$$0 - 0 - 0 = 0 \text{ (borrow 0)}$$

$$0 - 0 - 1 = 1 \text{ (borrow 1)}$$

$$0 - 1 - 0 = 1 \text{ (borrow 1)}$$

$$0 - 1 - 1 = 0 \text{ (borrow 1)}$$

$$1 - 0 - 0 = 1 \text{ (borrow 0)}$$

$$1 - 0 - 1 = 0 \text{ (borrow 0)}$$

$$1 - 1 - 0 = 0 \text{ (borrow 0)}$$

$$1 - 1 - 1 = 1 \text{ (borrow 1)}$$

Subtractions are difficult for computers.
↳ clever way to convert to addition

□ The digital logic necessary to implement bit-level arithmetic operations is trivial.

Binary Arithmetic

- ❑ When you add two binary numbers, you add same position bits together, one column at a time, starting with the least-significant bit.
- ❑ Any carry-out is added to the next column on the left.

Example 1

1

$$\begin{array}{r} 00101010 \\ +01001101 \\ \hline 01110111 \end{array}$$

Example 2

11111

$$\begin{array}{r} 10011111 \\ +00000001 \\ \hline 10100000 \end{array}$$

Example 3

$$\begin{array}{r} 00110011 \\ +11001100 \\ \hline 11111111 \end{array}$$

Example 4

111 11

$$\begin{array}{r} 01110011 \\ +01110011 \\ \hline 11100110 \end{array}$$

Addition

$0 + 0 = 0$ (carry 0)
 $0 + 1 = 1$ (carry 0)
 $1 + 0 = 1$ (carry 0)
 $1 + 1 = 0$ (carry 1)

Addition (three bits)

$0 + 0 + 0 = 0$ (carry 0)
 $0 + 0 + 1 = 1$ (carry 0)
 $0 + 1 + 0 = 1$ (carry 0)
 $0 + 1 + 1 = 0$ (carry 1)
 $1 + 0 + 0 = 1$ (carry 0)
 $1 + 0 + 1 = 0$ (carry 1)
 $1 + 1 + 0 = 0$ (carry 1)
 $1 + 1 + 1 = 1$ (carry 1)

Binary Arithmetic

❑ When subtracting binary numbers, you have to remember that *0 – 1 results in a difference 1 and a borrow from the column on the left.*

The borrow is not correct in the book.

Example 1 Example 2 Example 3 Example 4 Example 5

	1	1	1 111	
01101001	10011111	10111011	10110000	01100011
-01001001	-01000001	-10000100	-01100011	-10110000
00100000	01011110	00110111	01001101	

We have reversed the subtraction (smaller from larger) as we do in conventional mathematic.

1 1111
10110000
-01100011
01001101

Confused.

01100011
-10110000
-01001101

Subtraction
0 – 0 = 0 (borrow 0)
0 – 1 = 1 (borrow 1)
1 – 0 = 1 (borrow 0)
1 – 1 = 0 (borrow 0)

Subtraction (three bits)
0 – 0 – 0 = 0 (borrow 0)
0 – 0 – 1 = 1 (borrow 1)
0 – 1 – 0 = 1 (borrow 1)
0 – 1 – 1 = 0 (borrow 1)
1 – 0 – 0 = 1 (borrow 0)
1 – 0 – 1 = 0 (borrow 0)
1 – 1 – 0 = 0 (borrow 0)
1 – 1 – 1 = 1 (borrow 1)

Computers do not operate in this way!!

Binary Arithmetic

- ❑ In multiplication, **Multiplicand** × **Multiplier** = **Product**
- ❑ The following demonstrates the multiplication of 01101001_2 (the **multiplicand**) by 01001001_2 (the **multiplier**).
- ❑ You start with the least-significant bit of the **multiplier** and test whether it is a **0** or a **1**. If it is a **0**, you write down *n zeros*; if it is a **1**, you write down the **multiplicand** (this value is called a **partial product**).
- ❑ You then test the next bit of the **multiplier** to the left and carry out the same operation—in this case you write either *n zeros* or the **multiplicand** one place to the left (i.e., the **partial product** is shifted left).
- ❑ The process is continued until you have examined each bit of the **multiplier** in turn.
- ❑ Finally, you add together the *n partial products* to generate the **product** of the **multiplicand** times the **multiplier**.

get same result when multiplier × multiplicand.

Multiplicand	Multiplier	Step	Partial products
01101001	01001001 1	1	0 1 1 0 1 0 0 1
01101001	0100100 1	2	0 0 0 0 0 0 0 0
01101001	01001 0 01	3	0 0 0 0 0 0 0 0
01101001	0100 1 001	4	0 1 1 0 1 0 0 1
01101001	010 0 1001	5	0 0 0 0 0 0 0 0
01101001	01 0 01001	6	0 0 0 0 0 0 0 0
01101001	0 1 001001	7	0 1 1 0 1 0 0 1
01101001	0 1001001	8	0 0 0 0 0 0 0 0
Result			0 0 1 1 1 0 1 1

Arrows indicate the alignment of partial products for addition.

Binary Arithmetic

□ Note that,

- A computer does not perform multiplication operations in this way, as this would require *storing* the *n* partial products, followed by the simultaneous addition of *n* words.
- A better technique is to add up the partial products as they are formed.

Multiplicand	Multiplier	Step	Partial products															
01101001	0100100 1	1									0	1	1	0	1	0	0	1
01101001	0100100 1	2							0	0	0	0	0	0	0	0	0	0
01101001	01001 0 01	3						0	0	0	0	0	0	0	0	0	0	
01101001	0100 1 001	4				0	1	1	0	1	0	1	0	0	1			
01101001	010 0 1001	5			0	0	0	0	0	0	0	0	0	0				
01101001	01 0 01001	6		0	0	0	0	0	0	0	0	0	0					
01101001	0 1 001001	7		0	1	1	0	1	0	0	1							
01101001	0 1001001	8	0	0	0	0	0	0	0	0								
		Result	0	0	1	1	1	0	1	1	1	1	1	0	0	0	1	

Range, Precision, Accuracy and Errors

Range:

- ❑ The variation between the smallest and largest values that can be represented in a location
 - An n bits binary number has a range from 0 to $2^n - 1$.

For example, a 1 byte has a range from 0 to 255.

↳ 8 bits

0 to $2^8 - 1$

0 to $256 - 1$

0 to 255

- ❑ Note that: the number of possible values that can be encoded in an n bits binary number is 2^n

Range, Precision, Accuracy and Errors

Precision:

- The precision of a number is a measure of *how well (precise) we can represent the number*;
 - π cannot be exactly represented by a binary or a decimal real number – no matter how many bits we take.
 - If we use **5 decimal digits** to represent π (i.e., 3.1415), we say that its **precision is 1 in 10^5** , or you can say **5 significant figures**
 - If we use **10 decimal digits** to represent π (i.e., 3.141592653), we say that its **precision is 1 in 10^{10}** , or you can say **10 significant figures**

Range, Precision, Accuracy and Errors

Accuracy:

- ❑ The difference between *a representation* and *its actual value*
 - If we measure the temperature of a liquid as 51.32° and its actual temperature is 51.34°, the **error** is 0.02°.
 - The **error** (*true value – measured value*) is a one way to measure the **accuracy**.
- ❑ It is tempting to mix **accuracy** and **precision**.
- ❑ Be careful, they are *not the same*.
 - The temperature of the liquid may be measured as 51.320001° which has a **precision** of 8 significant figures, but if its actual temperature is 51.34° the **error** will be 0.019999°.
- ❑ What matters to us as computer designers, programmers, and users is
 - how errors *arise*,
 - how they are *controlled*, and
 - how their effects are *minimalized*.

} *relevance?*