



WEEK 4

INTRODUCTION TO QUERIES AND RELATIONAL LANGUAGES

CS3319

STUDENT OBJECTIVES

- Upon completion of this video, you should be able to:
 - Given some tables, give examples of queries that you would ask of the tables.
 - Distinguish between procedural languages and non procedural languages
 - Determine, given operators for a data type, if the operators enforce closure.
 - List the 5 basic operators for relational algebra and determine if each of those operations works on one table (unary) or two tables (binary).

CONSIDER...

- Fred needs to know how many people worked more than 40 hours on a project
- Sue needs to know who works in the Safety Department
- Homer wants to know how many supervising employees make over 2000 dollars

CONSIDER...

- Fred needs to know
more than 40 hours
- Sue needs to know
Department
- Homer wants to know
employees make over

Project : Table

| ProjectName | ProjectNumber | ProjectLocation | DeptNumber |
|----------------|---------------|-----------------|------------|
| Accounting Upd | A1 | Toronto | S7G |
| Acc3 | A3 | Springfield | G8H |
| Acct6 | A6 | Toronto | S7G |
| Inventory | I1 | Toronto | G8H |
| Inventory2 | I2 | London | S7G |
| Payroll | P1 | Springfield | G8H |
| Payroll2 | P2 | London | G8H |
| Payroll3 | P3 | London | G8H |

Department : Table

| DeptNumber | DeptName | ManagerSSN | ManagerStartDate |
|------------|---------------------|------------|------------------|
| G8H | Head Office | 4 | 2/2/95 |
| S7G | Safety Department | 3 | 1/1/95 |
| Y5J | Research Department | 6 | 3/3/95 |

Employee : Table

| SSN | LastName | MiddleInitial | FirstName | BDate | Address | Sex | Salary | SuperSSN | DeptNumber |
|-----|----------|---------------|-----------|--------|-----------|-----|------------|----------|------------|
| 1 | Simpson | P | Bart | 2/2/95 | London | M | \$1,000.00 | 2 | G8H |
| 2 | Smithers | J | Waylan | 1/1/60 | Springfie | M | \$2,000.00 | 4 | S7G |
| 3 | Beuvieau | P | Patty | 3/3/59 | Toronto | F | \$4,000.00 | 6 | Y5J |
| 4 | Burns | P | Montgomer | 7/7/20 | Toronto | M | \$5,000.00 | | S7G |
| 6 | Simpson | J | Lisa | 6/6/90 | London | F | \$1,000.00 | 2 | S7G |
| 12 | Simpson | J | Homer | 8/8/61 | Toronto | M | \$2,000.00 | 2 | G8H |

Record: 1 of 6

Works_On : Table

| SSN | ProjectNumber | Hours |
|-----|---------------|-------|
| 1 | A3 | 45 |
| 2 | A3 | 56 |
| 3 | A1 | 3 |
| 3 | A6 | 45 |
| 3 | I1 | 43 |
| 3 | P1 | 9 |
| 4 | A1 | 6 |
| 4 | A3 | 5 |
| 4 | A6 | 6 |
| 4 | I1 | 43 |
| 4 | I2 | 8 |
| 4 | P1 | 67 |
| 4 | P2 | 77 |
| 4 | P3 | 67 |
| 6 | I2 | 6 |
| 12 | A3 | 56 |

Record: 1

RELATIONAL LANGUAGES

- Once we have our relational model, we need to manipulate the data within the model, we use a relational language
- Some relational languages are **Procedural**, they tell us how to get the data (e.g. **Relational Algebra**)
- Some relational languages are **Non-Procedural**, they tell us what data is needed (e.g. **Relational Calculus**)
- Formally, Relational Algebra is equivalent to Relational Calculus, i.e. every expression in the algebra can be written also in the calculus and vice versa

- A language that can produce any relation that can be derived using relational calculus is relationally complete.
- Most relational query languages are relationally complete and more (i.e. they have additional power to do calculations, ordering, etc.)
- Relational algebra is a theoretical language with **operations** that perform on one or more relations and in turn produce relations based on the operations, thus both the input (**operands**) and output (results) are relations, i.e. a **closed language**, i.e. integer - integer produces an integer and a relation - relation produces a relation

EXAMPLE OF CLOSURE WITH RELATIONS

- Pretend that the symbol  represents an operation for the operand tables (i.e. relations) (just like + represents an operation with the operands integers)
- Then would be a closed operation if and only if:

| | | |
|-----|-------|------|
| 34 | Pig | Red |
| 445 | Horse | Red |
| 34 | Cat | Blue |



| | | | | | |
|----|---|-----|------|---------|------|
| 33 | X | Mar | 2018 | London | 22.2 |
| 44 | Y | Jun | 1964 | Toronto | 45.1 |
| 55 | X | Feb | 1982 | Windsor | 23.8 |
| 22 | B | Jan | 1977 | Arva | 0.1 |

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- In this case is a binary operation (2 operands).
- Could be unary operation like this:



| | | |
|-----|-------|------|
| 34 | Pig | Red |
| 445 | Horse | Red |
| 34 | Cat | Blue |

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QUESTION: Do the operations $-, +, /, *$ produce a closed language on integers? YES or NO? WHY?

NO! Because of $/ \rightarrow$ Division e.g. $9 / 2$

- Relational algebra is a set language, in which all tuples are manipulated in one statement, thus we don't use looping
- Because relations are produced, we can use the results in further operations, thus nesting our results, this is called **closure**; relations are closed under relational algebra

- **QUESTION:** What are unary operations with integers that insure closure?
- **ANSWER:** Power, Absolute Value...eg. 77^2 or $\text{abs}(77)$
- **QUESTION:** What are unary operations with integers that do not allow for closure?
- **ANSWER:** Square Root...eg. $\sqrt{77}$
- **QUESTION:** What are some unary operations when working with bits?
- **ANSWER:** FLIP...eg $\text{FLIP}(1011) = 0100$

BASIC OPERATIONS CAN BE USED TO BUILD OTHER MORE COMPLICATED OPERATIONS

- Eg. For integers, in order to create exponents, we could use the basic operation of multiplication:

3^4 is the same as $3 * 3 * 3 * 3$

so $*$ is a basic operation in arithmetic.

5 BASIC OPERATIONS IN RELATIONAL ALGEBRA

- **Selection:** Unary (works on 1 Relation (TABLE) only), returns only the tuples from a relation that satisfy the specified condition. (i.e. returns a row subset), written as:

$$\sigma_{\text{condition}} (R)$$

- **Projection:** Unary (works on 1 Relation only), returns only the requested attributes (with no duplicates) (returns a column subset), written as:

$$\pi_{\text{attribute1, attribute2}} (R)$$

- **Cartesian Product:** Binary (requires 2 Relations), returns a relation that is the concatenation of every tuple of relation R with every tuple of relation S, Symbol: $S \times R$
- **Union:** Binary (requires 2 Relations), union of relations R and S with I and J tuples, respectively, is the concatenation of them into one relation with a maximum of I+J tuples, duplicate tuples being eliminated, R and S must be union compatible (i.e. R and S must have the same columns or attribute domains). Symbol: $R \cup S$
- **Set Difference:** (requires 2 Relations), $R - S =$ a relation consisting of the tuples that are in relation R but not in S, R and S must be union compatible. $S - R$