# Number Theory and Cryptography Chapter 4: Part I

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## Chapter motivations

- Number theory is the part of mathematics devoted to the study of the integers and their properties.
- The key ideas in number theory include <u>divisibility</u> and the primality of integers.
- Representations of integers, including binary and hexadecimal representations, are part of number theory and essential to computer science.
- Mumber theory has long been studied because of the beauty of its ideas, its accessibility, and its wealth of open questions.
- We will use many ideas developed in Chapter 1 about proof methods and proof strategies in our exploration of number theory.
- Mathematicians have long considered number theory to be pure mathematics, but it has important applications to computer science and cryptography studied in the second part of this Chapter

- 1. Divisibility and Modular Arithmetic
- 1.1 Divisibility
- 1.2 Division
- 1.3 Congruence Relation
- 2. Integer Representations and Algorithms
- 2.1 Representations of Integers
- 2.2 Base conversions
- 2.3 Binary Addition and Multiplication
- 3. Prime Numbers
- 3.1 The Fundamental Theorem of Arithmetic
- 3.2 The Sieve of Erastosthenes
- 3.3 Infinitude of Primes
- 4. Greatest Common Divisors
- 4.1 Definition
- 4.2 Least common multiple
- 4.3 The Euclidean Algorithm

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# Divisibility

#### **Definition**

If a and b are integers with  $a \neq 0$ , then we say that a divides b if there exists an integer c such that b = ac holds.

- ① When a divides b we say that a is a <u>factor</u> or <u>a divisor</u> of b and we say that b is a multiple of a.
- ② The notation a b denotes the fact that a divides b.
- **3** If  $a \mid b$ , then  $\frac{b}{a}$  is an integer.
- 4 If a does not divide b, then we write a + b.

#### Example

Determine whether 3 | 7 holds and whether 3 | 12 holds.

**Solution**:  $3 \nmid 7$  but  $3 \mid 12$ 

## Properties of divisibility

#### **Theorem**

Let a, b, and c be integers, where  $a \neq 0$ .

- 1) If  $a \mid b$  and  $a \mid c$ , then  $a \mid (b+c)$ ;  $\frac{b}{a} = x + y$
- 2 If  $a \mid b$ , then  $a \mid b$  c for all integers c;  $\frac{b}{a} = \frac{1}{2}$
- 3 If  $a \mid b$  and  $b \mid c$ , then  $a \mid c$ .  $\frac{1}{6} = \sqrt{5} = \sqrt{3}$   $\frac{1}{6} = \sqrt{3}$ .

#### Proof.

- 1 We prove the first property. Suppose  $a \mid b$  and  $a \mid c$ , then it follows that there are integers s and t with b = as and c = at. Hence, b + c = as + at = a(s + t). Hence,  $a \mid (b + c)$ .
- 2 Parts (2) & (3) can be proven similarly. Try it as an exercise.

#### Corollary

If a, b, and c are integers, where  $a \neq 0$ , such that  $a \mid b$  and  $a \mid c$ , then  $a \mid mb + nc$  for any integers m and n. (Proof left as exercise)

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# The division $a = d \cdot (a \operatorname{div} d) + (a \operatorname{mod} d)$

## Theorem ("Division Algorithm")

If a is an integer and d is a positive integer, then there are unique integers q and r with  $0 \le r < d$ , such that a = dq + r (proved in the tutorial.

- a is called the dividend.
- 2 d is called the divisor.
- q is called the quotient.

Definitions div and mod:

- $\int_{a}^{6} r = a \mod d$

We have: a div  $d = \lfloor \frac{a}{d} \rfloor$ .

adiv  $d = \frac{\alpha}{1}$ :77 as cd

Example

- ① Quotient and remainder when 101 is divided by 11? We have 101 div 11 = 9 and 101 mod 11 = 2.
- **Q** Quotient and remainder when 11 is divided by 3? We have  $11 \operatorname{div} 3 = 3$  and  $11 \operatorname{mod} 3 = 2$ .
- **3** Quotient and remainder when -11 is divided by 3? We have  $-11 \operatorname{div} 3 = -4$  and  $-11 \operatorname{mod} 3 = 1$ .

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## Congruence relation

#### **Definition**

If a and b are integers and m is a positive integer, then a is congruent to b modulo m if m divides a - b.

- 1 The notations  $a \equiv b \pmod{m}$  and  $a \equiv b \pmod{m}$  say that a is congruent to b modulo m.
- 2 We say that  $a \equiv b \mod m$  is a <u>congruence</u> and that m is its <u>modulus</u>.
- 3 Two integers are congruent mod m if and only if they have the same remainder when divided by m. (to be proved later)
- If a is not congruent to b modulo m, then we write  $a \not\equiv b \mod m$ .

#### Example

- ① Determine whether 17 is congruent to 5 modulo 6. ((7-5)%6 = 0)  $17 = 5 \mod 6$  because 6 divides 17 5 = 12.
- 20 Determine whether 24 and 14 are congruent modulo 6.  $(24 14) \% 6 \stackrel{1}{>}0$   $24 \not\equiv 14 \mod 6$  since 24 14 = 10 is not divisible by 6.

## More on congruences

#### Theorem

Let m be a positive integer. The integers a and b are congruent modulo m if and only if there is an integer k such that a = b + km.  $\sim x mt = c \sim b = (x - y) m$ 

Proof.

- If  $a \equiv b \mod m$  holds, then (by the definition of congruence) we have:  $m \mid a b$ .  $a \equiv b \mod m$  holds =>  $m \mid a b => (a b) \% m = 0$
- 2 Hence, there is an integer k such that a b = km holds and equivalently a = b + km.
- **3** Conversely, if there is an integer k such that a = b + km, then we have: km = a b.
- 4 Hence, we have  $m \mid a b$ . Thus,  $a \equiv b \mod m$  holds.

b= ym+ 3

## Relationship between the mod m and mod m notations

The use of "mod" in  $a \equiv b \mod m$  is different from its use in  $a = b \mod m$ .

**border(1)**  $a \equiv b \mod m$  denotes a relation in the Cartesian product  $\mathbb{Z} \times \mathbb{Z}$ 

mt  $2^{n} = b \mod m$  denotes a function from  $\mathbb{Z} \times \mathbb{Z}$  to  $\mathbb{Z}$ .

The relationship between the two notions is stated below:

#### Theorem

Let a and b be integers, and let m be a positive integer. Then  $a \equiv b \mod m$  if and only if  $a \mod m = b \mod m$  (See Tutorial.)

## Congruences of sums and products

#### Theorem

Let a, b, c, d be integers. Let m be a positive integer. If  $a \equiv b \mod m$  and  $c \equiv d \mod m$  both hold, then we have:  $a + c \equiv b + d \mod m$  and  $ac \equiv bd \mod m$ .

Proof. 
$$\frac{(a-b)/m^2\pi}{(c-d)/m^2\eta}. \qquad a=m\infty+b \qquad \text{atc}=m(x+\eta)+b+d.$$
1 Since we have  $a\equiv b \mod m$  and  $c\equiv d \mod m$ , there exist

- integers s and t with b = a + sm and d = c + tm.
- 2 Therefore, we have:
  - a b + d = (a + sm) + (c + tm) = (a + c) + m(s + t) and
  - **b** bd = (a + sm)(c + tm) = ac + m(at + cs + stm).
- Hence, we have:
  - $a + c \equiv b + d \mod m$ , and
  - **b**  $ac \equiv bd \mod m$ .

Because  $7 \equiv 2 \mod 5$  and  $11 \equiv 1 \mod 5$ , it follows that:

 $18 = 7 + 11 \equiv 2 + 1 = 3 \mod 5$  and  $77 = 7 \cdot 11 \equiv 2 \cdot 1 = 2 \mod 5$ .

## Algebraic manipulation of congruences

Multiplying both sides of a valid congruence by an integer preserves the congruence.

If  $a \equiv b \mod m$  holds, then  $c \cdot a \equiv c \cdot b \mod m$ , where c is any integer, holds from the previous slide with d = c.

Adding an integer to both sides of a valid congruence

Adding an integer to both sides of a valid congruence preserves the congruence.

If  $a \equiv b \mod m$  holds, then  $c + a \equiv c + b \mod m$ , where c is any integer, holds from the previous slide with d = c.

- 3 NOTE: dividing a congruence by an integer may not produce a valid congruence.
  - a The congruence  $14 \equiv 8 \mod 6$  holds.
  - b Dividing both sides by 2 gives an invalid congruence since  $\frac{14}{2} = 7$  and  $\frac{8}{2} = 4$ , but  $7 \not\equiv 4 \mod 6$ .
  - Congruence.
    Congruence
    Congruence

# Computing the **mod** *m* function of products and sums

Given integers a, b, c, d and a positive integer m, recall the following properties:

- 2  $(a \equiv b \mod m) \land (c \equiv d \mod m) \rightarrow (a+c \equiv b+d \mod m) \land (ac \equiv bd \mod m)$

From there, we deduce the following properties:

- (ab)  $mod m = ((a mod m) \times (b mod m)) mod m$ .

See the tutorial for a proof.

- (1) a=xm+a (atb) mod m=[(x+y) m+(a+B)] mod m.=(a+B) mod m b=ym+B. [ca mod m)+cb mod m)] mod m=(a+B) mod m. (ab) mod m=aB mod m.
- [a.B] mod m= aB mod m.

### Arithmetic modulo *m*

#### **Definition**

Let  $\mathbb{Z}_m = \{0, 1, ..., m-1\}$  be the set of non-negative integers less than m. Assume  $a, b \in \mathbb{Z}_m$ .

- 1 The operation  $+_m$  is defined as  $a +_m b = a + b \mod m$ . This is the addition modulo m.
- 2 The operation  $\cdot_m$  is defined as  $a \cdot_m b = a \cdot b \mod m$ . This is the multiplication modulo m.
- 1 Using these operations is said to be doing arithmetic modulo m.

## Example

- ① Using the definitions above, find  $7 +_{11} 9$
- **Solution**:  $7 +_{11} 9 = (7 + 9) \mod 11 = 16 \mod 11 = 5$
- 3 Using the definitions above, find  $7 \cdot_{11} 9$ .
- **Solution**:  $7 \cdot_{11} 9 = (7 \cdot 9) \mod 11 = 63 \mod 11 = 8$

#### Arithmetic modulo *m*

The operations  $+_m$  and  $\cdot_m$  satisfy many of the same properties as ordinary addition and multiplication:

- 1) Closure: If a and b belong to  $\mathbb{Z}_m$ , then  $a +_m b$  and  $a \cdot_m b$  belong to  $\mathbb{Z}_m$ .
- **2** Associativity: If a, b, and c belong to  $\mathbb{Z}_m$ , then  $(a+_m b)+_m c=a+_m (b+_m c)$  and  $(a\cdot_m b)\cdot_m c=a\cdot_m (b\cdot_m c)$ .
- **3** Commutativity: If a and b belong to  $\mathbb{Z}_m$ , then  $\underline{a +_m b = b +_m a}$  and  $a \cdot_m b = b \cdot_m a$ .
- 4 Identity Elements: The elements 0 and 1 are identity elements for addition and multiplication modulo m, respectively.
  - a If a belongs to  $\mathbb{Z}_m$ , then  $a +_m 0 = a$  and  $a \cdot_m 1 = a$ .

(a+o) mod 
$$m = a$$

To be continued  $\rightarrow$ 

#### Arithmetic modulo *m*

**6** Additive inverses: If  $a \neq 0$  belongs to  $\mathbb{Z}_m$ , then m - a is the additive inverse of a modulo m and 0 is its own additive inverse. [at (m-a)] mod m = 0 + m = 0. a + m (m-a) = 0 and 0 + m = 0

6 Distributivity: If a, b, and c belong to  $\mathbb{Z}_m$ , then

$$\frac{a \cdot_m (b +_m c) = (a \cdot_m b) +_m (a \cdot_m c)}{(a +_m b) \cdot_m c = (a \cdot_m c) +_m (b \cdot_m c)}$$
 and

Multiplicative inverses have not been included since they do not always exist. For example, there is no multiplicative inverse of 2 modulo 6, i.e.

$$2 \cdot_m a \neq 1$$
 for any  $a \in \mathbb{Z}_6$ 

(optional) Using the terminology of abstract algebra,  $\mathbb{Z}_m$  with  $+_m$  is a commutative group and  $\mathbb{Z}_m$  with  $+_m$  and  $\cdot_m$  is a commutative ring.

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## Representations of integers

- 1 In the modern world, we use *decimal*, or *base* 10, to represent integers. For example when we write 965, we mean  $9 \cdot 10^2 + 6 \cdot 10^1 + 5 \cdot 10^0$ .
- ② We can represent numbers using any base b, where b is a positive integer greater than 1.
- **3** The bases b = 2 (binary), b = 8 (octal), and b = 16 (hexadecimal) are important for computing and communications
- The ancient Mayas used base 20 and the ancient Babylonians used base 60.

## Base b representations

① We can use any positive integer b greater than 1 as a base, because of this theorem:

#### Theorem

- a Let b be a positive integer greater than 1.
- **b** Then if n is a positive integer, it can be expressed uniquely in the form:

$$n = a_k b^k + a_{k-1} b^{k-1} + \dots + a_1 b + a_0$$

where k is a non-negative integer, such that  $a_0, a_1, \ldots a_k$  are non-negative integers less than b, and  $a_k \neq 0$ .

- The  $a_j$ , for j = 0, ..., k are called the base-b digits of the representation.
- d We will prove this using mathematical induction in Chapter 5.
- ② The representation of n given in the theorem is called the base b expansion of n and is denoted by  $(a_k a_{k-1} \dots a_1 a_0)_b$ .
- 10 We usually omit the subscript 10 for base 10 expansions.

## Binary expansions

Most computers represent integers and do arithmetic with binary (base 2) expansions of integers. In these expansions, the only digits used are 0 and 1.

#### Example

① What is the decimal expansion of the integer that has  $(1\ 0101\ 1111)_2$  as its binary expansion?

**Solution**: 
$$(1\ 0101\ 1111)_2 = 1 \cdot 2^8 + 0 \cdot 2^7 + 1 \cdot 2^6 + 0 \cdot 2^5 + 1 \cdot 2^4 + 1 \cdot 2^3 + 1 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 = 351.$$

What is the decimal expansion of the integer that has  $(1\ 1011)_2$  as its binary expansion?

**Solution**:  $(1\ 1011)_2 = 1 \cdot 2^4 + 1 \cdot 2^3 + 0 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 = 27$ .

# Octal expansions

The octal expansion (base 8) uses the digits  $\{0, 1, 2, 3, 4, 5, 6, 7\}$ .

### Example

① What is the decimal expansion of the number with octal expansion  $(7016)_8$ ?

**Solution**:  $7 \cdot 8^3 + 0 \cdot 8^2 + 1 \cdot 8^1 + 6 \cdot 8^0 = 3598$ 

② What is the decimal expansion of the number with octal expansion  $(111)_8$ ?

**Solution**:  $1 \cdot 8^2 + 1 \cdot 8^1 + 1 \cdot 8^0 = 64 + 8 + 1 = 73$ 

## Hexadecimal expansions

The hexadecimal expansion needs 16 digits, but our decimal system provides only 10. So letters are used for the additional symbols. The hexadecimal system uses the digits  $\{0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F\}$ . The letters A through F represent the decimal numbers 10 through 15.

## Example

① What is the decimal expansion of the number with hexadecimal expansion  $(2AE0B)_{16}$ ?

**Solution**:  $2 \cdot 16^4 + 10 \cdot 16^3 + 14 \cdot 16^2 + 0 \cdot 16^1 + 11 \cdot 16^0 = 175627$ 

② What is the decimal expansion of the number with hexadecimal expansion  $(1E5)_{16}$ ?

**Solution**:  $1 \cdot 16^2 + 14 \cdot 16^1 + 5 \cdot 16^0 = 256 + 224 + 5 = 485$ 

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### Base conversion

To construct the base b expansion of an integer n (given in base 10):

Divide n by b to obtain the quotient  $q_0$  and remainder  $a_0$ :

$$n = bq_0 + a_0, \quad 0 \le a_0 < b$$

- The remainder,  $a_0$ , is the rightmost digit in the base b expansion of n.
- **3** If  $q_0 = 0$ , then  $n = (a_0)_b$ .
- 4 If  $0 < q_0 < b$ , then  $n = (q_0 a_0)_b$ .
- **6** If  $b \le q_0$ , then divide  $q_0$  by b to obtain the quotient  $q_1$  and remainder a<sub>1</sub>:

$$q_0 = bq_1 + a_1, \quad 0 \le a_1 < b$$

- **6** The remainder,  $a_1$ , is the second digit from the right in the base b expansion of n.
- Continuing in this manner (by successively dividing the quotients by b) we obtain the additional base b digits as remainders. The process terminates when a quotient is 0.

## Algorithm: constructing base b expansions

### **Algorithm 1** base\_b\_expansion(n, b)

```
Require: n, b \in \mathbb{Z}^+, b > 1

Ensure: base b expansion of n: (a_{k-1} \cdots a_1 a_0)_b.

1: q \leftarrow n

2: k \leftarrow 0

3: while q \neq 0 do

4: a_k \leftarrow q \mod b

5: q \leftarrow q \operatorname{div} b

6: k \leftarrow k + 1

7: end while

8: return (a_{k-1} \cdots a_1 a_0)
```

- q represents the quotient obtained by successive divisions by b, starting with q = n.
- ② The digits in the base b expansion are the remainders of the division given by q mod b.
- 3 The algorithm terminates when q = 0 is reached.

#### Base conversion

## Example

Find the octal expansion of  $(12345)_{10}$ 

**Solution**: Successively dividing by 8 gives:

$$12345 = 8 \cdot 1543 + 1$$

$$21543 = 8 \cdot 192 + 7$$

$$\mathbf{3} \ 192 = 8 \cdot 24 + 0$$

$$4 \cdot 24 = 8 \cdot 3 + 0$$

$$\mathbf{6} \ 3 = 8 \cdot 0 + 3$$

The remainders are the digits from right to left yielding  $(30071)_8$ .

# Comparison of the hexadecimal, octal, and binary representations

TABLE 1 Hexadecimal, Octal, and Binary Representation of the Integers 0 through 15.																
Decimal	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Hexadecimal	0	1	2	3	4	5	6	7	8	9	A	В	С	D	Е	F
Octal	0	1	2	3	4	5	6	7	10	11	12	13	14	15	16	17
Binary	0	1	10	11	100	101	110	111	1000	1001	1010	1011	1100	1101	1110	1111

Initial Os are not shown

- Each octal digit corresponds to a block of 3 binary digits.
- ② Each hexadecimal digit corresponds to a block of 4 binary digits.
- 6 So, conversion between binary, octal, and hexadecimal is easy.

# Conversion between the binary, octal, and hexadecimal expansions

#### Example

- ① Find the octal expansion of  $(1111110101111100)_2$ .
  - **Solution**: To convert to octal, we group the digits into blocks of three  $(011\ 111\ 010\ 111\ 100)_2$ , adding initial 0s as needed. The blocks from left to right correspond to the digits 3,7,2,7, and 4. Hence, the solution is  $(37274)_8$ .
- ② Find the hexadecimal expansions of  $(1111110101111100)_2$ .
  - **Solution**: To convert to hexadecimal, we group the digits into blocks of four  $(0011\ 1110\ 1011\ 1100)_2$ , adding initial 0s as needed. The blocks from left to right correspond to the digits 3,E,B, and C. Hence, the solution is  $(3EBC)_{16}$ .

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# Binary addition of integers

Algorithms for performing operations with integers using their binary expansions are important as computer chips work with binary numbers. Each digit is called a *bit*.

## **Algorithm 2** add (a, b)

```
Require: a, b \in \mathbb{Z}^+, {the binary expansions of a and b are (a_{n-1}, a_{n-2}, \dots, a_0)_2 and
     (b_{n-1}, b_{n-2}, \dots, b_0)_2, respectively
Ensure: (s_n, \ldots, s_1, s_0), the addition of a and b. {the binary expansion of the sum is
(s_n, s_{n-1}, \dots, s_0)_2 } 1: c_{prev} \leftarrow 0
                                                 > represents carry from the previous bit addition
2: for j \leftarrow 0, n-1 do
3: c \leftarrow \left| \frac{(a_j + b_j + c_{prev})}{2} \right|
                                                   > quotient (carry for the next digit of the sum)
 4: s_i \leftarrow a_i + b_i + c_{prev} - 2c
                                                                   > remainder (j-th digit of the sum)
5: c_{prev} \leftarrow c
                                                                                         a_0 + b_0 = c_0 \cdot 2 + s_0
 6: end for
                                                                                   a_1 + b_1 + c_0 = c_1 \cdot 2 + s_1
7: s_n \leftarrow c
                                                                                 a_i + b_i + c_{i-1} = c_j \cdot 2 + s_j
 8: return (s_n, ..., s_1, s_0)
```

The number of additions of bits used by the algorithm to add two *n*-bit integers is  $\mathcal{O}(n)$ .

## Binary multiplication of integers

Algorithm for computing the product of two n bit integers.

```
a \cdot b = a \cdot (b_k 2^k + b_{k-1} 2^{k-1} + \dots + b_1 2 + b_0
= ab_k 2^k + ab_{k-1} 2^{k-1} + \dots + ab_1 2 + ab_0
shift by k shift by k-1 shift by 1 no shift
```

### **Algorithm 3** multiply (a, b)

```
Require: a, b \in \mathbb{Z}^+, {the binary expansions of a and b are (a_{n-1}, a_{n-2}, \dots, a_0)_2 and
     (b_{n-1}, b_{n-2}, \ldots, b_0)_2, respectively
Ensure: p, the value of ab.
1: for i \leftarrow 0, n-1 do
    if b_i = 1 then
    c<sub>i</sub> ← a
                                                                                                   \triangleright shifted j places
     else
     c_i \leftarrow 0
                                                                \triangleright { c_0, c_1, \ldots, c_{n-1} are the partial products}
      end if
7: end for
                                                                                                            110
                                                                                                                     a
8: p \leftarrow 0
                                                                                                          x 101
9: for i \leftarrow 0, n-1 do
10: p \leftarrow p + c_j
                                                                                                            110
                                                                                                                    ab<sub>∩</sub>
11: end for
                                                                                                          000
                                                                                                                    ab_1
12: return p \{ p \text{ is the value of } ab \}
                                                                                                         110
                                                                                                                    ab_2
```

The number of additions of bits used by the algorithm to multiply two n-bit integers is  $\mathcal{O}(n^2)$ .

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### **Primes**

### **Definition**

- ① A positive integer p greater than 1 is said prime if the only positive factors of p are 1 and p.
- ② A positive integer that is greater than 1 and is not prime is called composite.

## Example

The integer 7 is prime because its only positive factors are 1 and 7, but 9 is composite because it is divisible by 3.

# The fundamental theorem of arithmetic (prime factorization )

#### Theorem

- Every positive integer greater than 1 can be written uniquely as a prime or as the product of two or more primes where the prime factors are written in order of nondecreasing size.
- **2** More formally, for every positive integer a greater than 1, there exists a positive integer n such that there exist prime numbers  $p_1, \ldots, p_n$  and positive integers  $a_1, \ldots, a_n$  such that:

$$a = p_1^{a_1} p_2^{a_2} \cdots p_n^{a_n}$$
 and  $p_1 < p_2 < \cdots < p_n$ .

### Example

- $100 = 2 \cdot 2 \cdot 5 \cdot 5 = 2^2 \cdot 5^2$
- **2** 641 = 641
- $999 = 3 \cdot 3 \cdot 3 \cdot 37 = 3^3 \cdot 37$

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# The sieve of Erastosthenes



The *Sieve of Erastosthenes* can be used to <u>find all primes</u> not exceeding a specified positive integer.

### Example

- Consider the list of integers between 1 and 100:
  - Delete all the integers, other than 2, divisible by 2.
  - **(b)** Delete all the integers, other than 3, divisible by 3.
  - Next, delete all the integers, other than 5, divisible by 5.
  - d Next, delete all the integers, other than 7, divisible by 7.

all remaining numbers between 1 and 100 are prime:  $\{2,3,7,11,19,23,29,31,37,41,43,47,53,59,61,67,71,73,79,83,89,97\}$ 

Why does this work?

continued →

### The sieve of Erastosthenes

	TABLE 1 The Sieve of Eratosthenes.																				
Г	Inte	egers	divisi	ble b	y 2 ot	her ti	han 2			In	Integers divisible by 3 other than 3										
	receive an underline.											receive an underline.									
Г	1	2	3	4	5	6	7	8	9	10	1	2	3	4	5	<u>6</u>	7	8	9	10	
	11	12	13	14	15	16	17	18	19	20	11	<u>12</u>	13	14	15	16	17	18	19	20	
	21	22	23	<u>24</u>	25	<u>26</u>	27	28	29	<u>30</u>	21	<u>22</u>	23	<u>24</u>	25	<u> 26</u>	<u>27</u>	<u>28</u>	29	<u>30</u>	
	31	<u>32</u>	33	<u>34</u>	35	<u>36</u>	37	<u>38</u>	39	<u>40</u>	31	<u>32</u>	<u>33</u>	34	35	<u>36</u>	37	<u>38</u>	<u>39</u>	<u>40</u>	
	41	<u>42</u>	43	<u>44</u>	45	<u>46</u>	47	<u>48</u>	49	<u>50</u>	41	<u>42</u>	43	<u>44</u>	<u>45</u>	<u>46</u>	47	<u>48</u>	49	<u>50</u>	
	51	<u>52</u>	53	<u>54</u>	55	<u>56</u>	57	<u>58</u>	59	<u>60</u>	<u>51</u>	<u>52</u>	53	<u>54</u>	55	<u>56</u>	<u>57</u>	<u>58</u>	59	<u>60</u>	
	61	<u>62</u>	63	<u>64</u>	65	<u>66</u>	67	<u>68</u>	69	<u>70</u>	61	<u>62</u>	<u>63</u>	<u>64</u>	65	<u>66</u>	67	<u>68</u>	<u>69</u>	<u>70</u>	
	71	<u>72</u>	73	<u>74</u>	75	<u>76</u>	77	<u>78</u>	79	80	71	<u>72</u>	73	<u>74</u>	<u>75</u>	<u>76</u>	77	<u>78</u>	79	<u>80</u>	
l	81	<u>82</u>	83	<u>84</u>	85	<u>86</u>	87	<u>88</u>	89	<u>90</u>	81	<u>82</u>	83	<u>84</u>	85	<u>86</u>	<u>87</u>	<u>88</u>	89	90	
	91	<u>92</u>	93	<u>94</u>	95	<u>96</u>	97	<u>98</u>	99	<u>100</u>	91	<u>92</u>	<u>93</u>	<u>94</u>	95	<u>96</u>	97	<u>98</u>	<u>99</u>	<u>100</u>	
Г	Integers divisible by 5 other than 5											Integers divisible by 7 other than 7 receive									
	receive an underline.										an underline; integers in color are prime.										
Γ	1				-	6	7	0													
		2	3	<u>4</u>	5	0	/	8	9	<u>10</u>	1	2	3	<u>4</u>	5	6	7	8	9	<u>10</u>	
	11	2 12	3 13	<u>4</u> <u>14</u>		<u>6</u> <u>16</u>	17	<u>8</u> <u>18</u>	9 19		1 11	2 12	3 13			<u>6</u> <u>16</u>	7 17	<u>8</u> <u>18</u>	9 19	<u>10</u> <u>20</u>	
	11 21				15 25					<u>20</u>				4 14 24	5 <u>15</u> <u>25</u>		7 17 <u>27</u>			<u>20</u>	
		<u>12</u>	13	<u>14</u>	<u>15</u>	<u>16</u>	17	<u>18</u>	19	<u>20</u> <u>30</u>	11	<u>12</u>	13	<u>14</u>	15 25	16 26		<u>18</u>	19	<u>20</u> <u>30</u>	
	<u>21</u>	<u>12</u> <u>22</u>	13 23	14 24	<u>15</u> <u>25</u>	16 26	17 <u>27</u>	18 28	19 29	<u>20</u>	11 <u>21</u>	12 22 32	13 23	<u>14</u> <u>24</u>	15 25 35	<u>16</u>	<u>27</u>	<u>18</u> <u>28</u>	19 29	<u>20</u>	
	<u>21</u> 31	$\frac{\underline{12}}{\underline{22}}$ $\underline{32}$	13 23 <u>33</u>	14 24 34	$\frac{\underline{15}}{\underline{25}}$ $\underline{35}$	16 26 36	17 <u>27</u> 37	$\frac{18}{28}$ $\underline{38}$	19 29 <u>39</u>	20 30 40	11 <u>21</u> 31	<u>12</u> <u>22</u>	13 23 <u>33</u>	14 24 34	15 25	16 26 36	<ul><li>27</li><li>37</li></ul>	$\frac{\underline{18}}{\underline{28}}$ $\underline{38}$	19 29 39	<u>20</u> <u>30</u> <u>40</u>	
	21 31 41	$\frac{12}{22}$ $\frac{32}{42}$	13 23 <u>33</u> 43	14 24 34 44	$\frac{15}{25}$ $\frac{35}{45}$	16 26 36 46 56	17 <u>27</u> 37 47	18 28 38 48	19 29 39 49	20 30 40 50 60	11 <u>21</u> 31 41	<u>12</u> <u>22</u> <u>32</u> <u>42</u>	13 23 33 43 53	14 24 34 44	15 25 35 45	16 26 36 46	27 37 47	18 28 38 48	19 29 39 49	20 30 40 50 60	
	21 31 41 51	12 22 32 42 52 62	13 23 <u>33</u> 43 53	14 24 34 44 54	15 25 35 45 55 65	16 26 36 46	17 <u>27</u> 37 47 <u>57</u>	18 28 38 48 58	19 29 39 49 59	20 30 40 50	11 21 31 41 51	$     \begin{array}{r}                                     $	13 23 33 43	14 24 34 44 54	15 25 35 45 55 65	16 26 36 46 56	<ul><li>27</li><li>37</li><li>47</li><li>57</li></ul>		19 29 39 49 59	20 30 40 50	
	<ul> <li>21</li> <li>31</li> <li>41</li> <li>51</li> <li>61</li> </ul>	$\frac{12}{22}$ $\frac{32}{42}$ $\frac{42}{52}$	13 23 33 43 53 63	14 24 34 44 54 64	15 25 35 45 55	16 26 36 46 56 66	17 <u>27</u> 37 47 <u>57</u> 67		19 29 39 49 59	20 30 40 50 60 70	11 21 31 41 51 61	$     \begin{array}{r}                                     $	13 23 33 43 53 <u>63</u>	14 24 34 44 54 64	15 25 35 45 55	16 26 36 46 56 66	<ul> <li>27</li> <li>37</li> <li>47</li> <li>57</li> <li>67</li> </ul>	18 28 38 48 58	19 29 39 49 59	20 30 40 50 60 70	

- 1 If an integer n is a composite integer, then it must have a prime divisor less than or equal to  $\sqrt{n}$ .
- 2 To see this, note that if n = ab, then  $a \le \sqrt{n}$  or  $b \le \sqrt{n}$ .
- 3 For n = 100,  $\sqrt{n} = 10$ , thus any composite integer  $\leq 100$  must have prime factors less than 10, that is 2,3,5,7. The remaining integers  $\leq 100$  are prime.
- 4 Trial division, a very inefficient method of determining if a number n is prime, is to try every integer  $i \le \sqrt{n}$  and see if n is divisible by i.

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B.C)

### **Theorem**

There are infinitely many primes.

#### Proof.

- **1** Assume finitely many primes:  $p_1, p_2, \ldots, p_n$ .
- ② Let  $q = p_1 p_2 \cdots p_n + 1$
- - a If a prime  $p_j$  divides q, and since  $p_j \mid p_1 p_2 \cdots p_n$  holds as well, then  $p_j$  divides  $q p_1 p_2 \cdots p_n = 1$ .
  - **b** Thus, if a prime  $p_j$  divides q, then  $p_j = 1$ , which is a contradiction with  $p_i > 1$ .
- 4 Hence, there is no prime on the list  $p_1, p_2, \ldots, p_n$  dividing q, that is, q is a prime.
- **5** This contradicts the assumption that  $p_1, p_2, \ldots, p_n$  are all the primes.
- 6 Consequently, there are infinitely many primes.

This proof was given by Euclid in *The Elements*.

# Generating primes

- The problem of generating large primes is of both theoretical and practical interest.
- ② Finding large primes with hundreds of digits is important in cryptography.
- 3 So far, no useful closed formula that always produces primes has been found. There is no simple function f(n) such that f(n) is prime for all positive integers n.
- $f(n) = n^2 n + 41$  is prime for all integers 1, 2, ..., 40. Because of this, we might conjecture that f(n) is prime for all positive integers n. But  $f(41) = 41^2$  is not prime.
- 6 More generally, there is no polynomial with integer coefficients such that f(n) is prime for all positive integers n.
- 6 Fortunately, we can generate large integers which are almost certainly primes.



(1588 - 1648)

### **Definition**

Prime numbers of the form  $2^p - 1$ , where p is prime, are called *Mersenne primes*.

- ①  $2^2 1 = 3$ ,  $2^3 1 = 7$ ,  $2^5 1 = 37$ , and  $2^7 1 = 127$  are Mersenne primes.
- $2^{11} 1 = 2047$  is not a Mersenne prime since  $2047 = 23 \cdot 89$ .
- 3 There is an efficient test for determining if  $2^p 1$  is prime.
- The largest known prime numbers are Mersenne primes.
- **6** On December 26 2017, the 50-th Mersenne primes was found, it is  $2^{77,232,917} 1$ , which is the largest Marsenne prime known. It has more than 23 million decimal digits.
- **6** The *Great Internet Mersenne Prime Search* (*GIMPS*) is a distributed computing project to search for new Mersenne Primes.

http://www.mersenne.org/

# Conjectures about primes

Even though primes have been studied extensively for centuries, many conjectures about them are unresolved, including:

- **Goldbach's conjecture**: Every even integer n, n > 2, is the sum of two primes. It has been verified by computer for all positive even integers up to 1.  $6 \cdot 10^{18}$ . The conjecture is believed to be true by most mathematicians.
- **Landau's conjecture**: There are infinitely many primes of the form  $n^2 + 1$ , where n is a positive integer. But it has been shown that there are infinitely many numbers of the form  $n^2 + 1$  which are the product of at most two primes.
- **The Twin Prime Conjecture**: there are infinitely many pairs of twin primes. Twin primes are pairs of primes that differ by 2. Examples are 3 and 5, 5 and 7, 11 and 13, etc. The current world's record for twin primes (as of mid 2011) consists of numbers  $65,516,468,355 \cdot 23^{33,333} \pm 1$ , which have 100,355 decimal digits.

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# Greatest common divisor (GCD)

From *primes* to *relative primes* 

### **Definition**

Let a and b be integers, not both zero.

- 1 The largest integer d such that  $d \mid a$  and also  $d \mid b$  is called the greatest common divisor of a and b.
- ② The greatest common divisor (GCD) of a and b is denoted by gcd(a, b).

One can find GCDs of small numbers by inspection.

## Example

- What is the greatest common divisor of 24 and 36?
  - **Solution**: gcd(24, 26) = 12
- What is the greatest common divisor of 17 and 22?

**Solution**: 
$$gcd(17, 22) = 1$$

# Greatest common divisor (GCD) From primes to relative primes

#### **Definition**

The integers a and b are relatively prime if their greatest common divisor is gcd(a, b) = 1.

### Example

17 and 22

#### **Definition**

The integers  $a_1, a_2, \ldots, a_n$  are pairwise relatively prime if  $gcd(a_i, a_j) = 1$  whenever  $1 \le i < j \le n$ .

### Example

① Determine whether the integers 10, 17 and 21 are pairwise relatively prime.

**Solution**: Because gcd(10, 17) = 1, gcd(10, 21) = 1, and gcd(17, 21) = 1, 10, 17, and 21 are pairwise relatively prime.

2 Determine whether the integers 10, 19, and 24 are pairwise relatively prime.

**Solution**: No, since gcd(10, 24) = 2.

# Finding GCDs using prime factorizations

Suppose that the prime factorizations of a and b are:

$$a = p_1^{a_1} p_2^{a_2} \cdots p_n^{a_n}, \qquad b = p_1^{b_1} p_2^{b_2} \cdots p_n^{b_n}$$

where each exponent is non-negative, and where all primes occurring in either prime factorization are included in both.

2 Then:

$$gcd(a,b) = p_1^{min(a_1,b_1)} p_2^{min(a_2,b_2)} \cdots p_n^{min(a_n,b_n)}$$

- This formula is valid since
  - a the integer on the right-hand side divides both a and b,
  - **b** No larger integer can divide both *a* and *b*.

# Example

Since  $120 = 2^3 \cdot 3 \cdot 5$  and  $500 = 2^2 \cdot 5^3$ , we have:

$$\gcd(120,500) = 2^{\min(3,2)} \cdot 3^{\min(1,0)} \cdot 5^{\min(1,3)} = 2^2 \cdot 3^0 \cdot 5^1 = 20$$

Remark: finding the GCD of two positive integers using their prime factorizations is not efficient because there is no efficient algorithm for finding the prime factorization of a positive integer.

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# Least common multiple (LCM)

### **Definition**

- 1 The least common multiple (LCM) of the positive integers a and b is the smallest positive integer that is divisible by both a and b. It is denoted by lcm(a, b).
- The least common multiple can also be computed from the prime factorizations.

$$lcm(a,b) = p_1^{max(a_1,b_1)} p_2^{max(a_2,b_2)} \cdots p_n^{max(a_n,b_n)}$$

This number is divided by both a and b and no smaller number is divided by a and b.

### Example

$$lcm(2^{3}3^{5}7^{2}, 2^{4}3^{3}) = 2^{max(3,4)}3^{max(5,3)}7^{max(2,0)} = 2^{4}3^{5}7^{2}$$

#### **Theorem**

Let a and b be positive integers. Then, we have:

$$a \cdot b = \gcd(a, b) \cdot \operatorname{lcm}(a, b)$$

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# The Euclidean Algorithm

- The Euclidean Algorithm is an <u>efficient method</u> for computing the GCD of two integers.
- It is based on the idea that

$$gcd(a, b) = gcd(b, r)$$

when a > b and r is the remainder when a is divided by b.

1 Indeed, since a = bq + r, then r = a - bq. Thus, if  $d \mid a$  and  $d \mid b$  then  $d \mid r$ .

## Example

- **1** Find gcd(287, 91):
  - (a)  $287 = 91 \cdot 3 + 14$  Divide 287 by 91
  - **b**  $91 = 14 \cdot 6 + 7$  Divide 91 by 14
  - c  $14 = 7 \cdot 2 + 0$  Divide 14 by 7 Zero remainder is our stopping condition.

$$gcd(287,91) = gcd(91,14) = gcd(14,7) = gcd(7,0) = 7$$

# The Euclidean Algorithm

The Euclidean algorithm expressed in pseudo-code is:

# **Algorithm 4** gcd(a, b)

```
Require: a, b \in \mathbb{Z}^+, a > b
```

**Ensure:** x, the GCD of a and b.

- 1:  $x \leftarrow a$
- 2:  $y \leftarrow b$
- 3: **while**  $y \ne 0$  **do**
- 4:  $r \leftarrow x \mod y$
- 5:  $x \leftarrow y$
- 6:  $y \leftarrow r$
- 7: end while
- 8: **return** *X*

Note: the time complexity of the algorithm is  $\mathcal{O}(\log^2 a)$ , where a > b.

# Correctness of the Euclidean Algorithm

#### Lemma

Let  $r = a \mod b$ , where  $a \ge b > r$  are integers. Then, we have:

$$gcd(a, b) = gcd(b, r).$$

### Proof.

- ① Any divisor of a and b must also be a divisor of b and r since r = a bq (with  $q = a \operatorname{div} b$ .)
- ② Similarly, any divisor of b and r is also a divisor of a and b.
- **3** Therefore, the set of common divisors of a and b is equal to the set of common divisors of b and r.
- 4 Therefore, gcd(a, b) = gcd(b, r).

# Correctness of the Euclidean Algorithm

① Suppose that a and b are positive integers with  $a \ge b$ . Let  $r_0 = a$  and  $r_1 = b$ . Successive applications of the division algorithm yields:

$$r_0 = q_1 r_1 + r_2$$
  $0 \le r_2 < r_1 \le r_0$   
 $r_1 = q_2 r_2 + r_3$   $0 \le r_3 < r_2$   
 $\vdots$   
 $r_{n-2} = r_{n-1} q_{n-1} + r_n$   $0 \le r_n < r_{n-1}$   
 $= r_n q_n$  (gcd)

- 2 Eventually, a remainder of zero occurs in the sequence of terms:  $a = r_0 \ge r_1 > r_2 > \cdots \ge 0$ . The sequence can not contain more than (a + 1) terms.
- 3 Then, the Lemma implies:  $gcd(a, b) = gcd(r_0, r_1) = \cdots = gcd(r_{n-1}, r_n) = gcd(r_n, 0) = r_n$ .
- 4 Hence the GCD is the last nonzero remainder in the sequence of divisions.



Étienne

Bézout

(1730 - 1783)

# Theorem (Bézout's Theorem)

If a and b are positive integers, then there exist integers s and t such that

$$gcd(a, b) = sa + tb.$$

### **Definition**

- ① If a and b are positive integers, then integers s and t such that gcd(a, b) = sa + tb are called Bézout coefficients of a and b.
- ② The equation gcd(a, b) = sa + tb is called *Bézout's identity*.
- 3 The expression sa + tb is also called a *linear combination* of a and b with coefficients of s and t.

## Example

$$gcd(6, 14) = 2 = (-2) \cdot 6 + 1 \cdot 14$$

# Finding GCD(s) as linear combinations

### Example

Express gcd(252, 198) = 18 as a linear combination of 252 and 198.

**Solution**: First use the Euclidean algorithm to show gcd(252, 198) = 18

- $252 = 1 \cdot 198 + 54$
- **b**  $198 = 3 \cdot 54 + 36$
- $\mathbf{6} \ 54 = 1 \cdot 36 + \mathbf{18}$
- **d**  $36 = 2 \cdot 18$
- Working backwards , from c and b above

$$18 = 54 - 1 \cdot 36$$
$$36 = 198 - 3 \cdot 54$$

② Substituting the  $2^{nd}$  equation into the  $1^{st}$  yields:

$$18 = 54 - 1 \cdot (198 - 3 \cdot 54) = 4 \cdot 54 - 1 \cdot 198$$

Substituting  $54 = 252 - 1 \cdot 198$  (from a above) yields:

$$18 = 4 \cdot (252 - 1 \cdot 198) - 1 \cdot 198 = 4 \cdot 252 - 5 \cdot 198$$

This method illustrated above is a two pass method. It first uses the Euclidean algorithm to find the GCD and then works backwards to express the GCD as a linear combination of the original two integers. There is a one pass method, called the *extended Euclidean algorithm*.

# Consequences of Bézout's Theorem

#### Lemma

If a, b, c are positive integers such that a and b are relatively prime (that is, gcd(a, b) = 1) and  $a \mid bc$ , then we have  $a \mid c$ .

#### PROOF:

- ① Assume gcd(a, b) = 1 and  $a \mid bc$  both hold.
- ② Since gcd(a, b) = 1, by Bézout's Theorem there are integers s and t such that sa + tb = 1 holds.
- 1 Multiplying both sides of the equation by c, yields sac + tbc = c.
- 4 Since  $a \mid bc$ , we have  $a \mid tbc$ , that is, there exists q so that we have tbc = qa.
- 6 With sac + tbc = c, it follows that a(sc + q) = c, that is,  $a \mid c$  holds.

A generalization of the above lemma is important in practice:

#### Lemma

If p is prime and  $p \mid a_1 a_2 \dots a_n$  where  $a_i$  are integers then  $p \mid a_i$  for some i.

# Dividing congruences by an integer

- ① Dividing both sides of a valid congruence by an integer does not always produce a valid congruence, as illustrated earlier.
- ② But dividing by an integer relatively prime to the modulus does produce a valid congruence.

#### Theorem

Let m be a positive integer and let a, b, and c be integers. If gcd(c, m) = 1 and  $ac \equiv bc \mod m$ , then  $a \equiv b \mod m$ .

#### Proof.

① Since  $ac \equiv bc \mod m$  holds, we have

$$m \mid ac - bc = c(a - b).$$

- ② With the previous lemma and since gcd(c, m) = 1 holds, it follows that  $m \mid a b$ ..
- **3** Hence,  $a \equiv b \mod m$ .