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$$\begin{array}{lcl} \text{I. (i)} & P \vee Q & \neg Q \equiv T \quad P \vee F \equiv T \\ & P \rightarrow r & \\ & \neg Q & \\ \hline & \therefore r & \end{array} \quad \begin{array}{l} Q \equiv F \\ P \equiv T \text{ (domination laws)} \end{array}$$

$$P \rightarrow r \equiv \neg P \vee r \text{ (conditional identity)}$$

$$\equiv F \vee r$$

$$F \vee r \equiv T$$

$$r \equiv T \text{ (domination laws)}$$

$\therefore$  the proof is valid.

$$\text{II) } \forall x (P(x) \rightarrow (Q(x) \wedge (R(x))))$$

$$\forall x (S(x) \vee P(x))$$

$$\neg \forall x Q(x)$$

$$\exists x S(x).$$

$$\neg \forall x Q(x) \equiv \exists x \neg Q(x).$$

$$\therefore \forall x Q(x) \equiv F.$$

$$\forall x (P(x) \rightarrow (Q(x) \wedge (R(x)))) \equiv \neg (\forall x (P(x))) \vee ((Q(x) \wedge (R(x)))) \quad \text{(condition identity)}$$

$$\equiv \neg (\forall x (P(x))) \vee (F \wedge (R(x)))$$

$$\equiv \neg (\forall x (P(x))) \vee F \text{ (domination laws).}$$

$$\neg (\forall x (P(x))) \vee F \equiv T$$

$$\neg (\forall x (P(x))) \equiv T \text{ (domination laws).}$$

$$\forall x (P(x)) \equiv F.$$

$$\forall x (S(x) \vee P(x)) \equiv (\forall x S(x)) \vee (\forall x P(x))$$

$$\equiv (\forall x S(x)) \vee F.$$

$$(\forall x S(x)) \vee F \equiv T.$$

$$\forall x S(x) \equiv T \text{ (domination laws).}$$

$$\therefore \exists x S(x) \equiv T$$

$\therefore$  the proof is valid.

2. (i) Some student don't have an official email address.

(ii) Exist a real number  $x$  that there is no real number  $y$  satisfies  $x = y^2$

3. Theorem:  $P_1, \dots, P_n$  is a collection of propositions.

$$\neg(P_1 \vee \dots \vee P_n) \equiv (\neg P_1 \wedge \dots \wedge \neg P_n).$$

Proof:  $(P_1 \vee \dots \vee P_n)$  is false iff all propositions from  $P_1$  to  $P_n$  are false, so  $\neg(P_1 \vee \dots \vee P_n)$  is true iff all propositions are false.

Assume a new collection  $q_1, \dots, q_n$  where  $q_n = \neg P_n$ .

$q_1 \wedge \dots \wedge q_n$  is true iff all propositions from  $q_1$  to  $q_n$  are true, so  $\neg P_1 \wedge \dots \wedge \neg P_n$  is true iff all propositions are false.

the only true condition is the same from both sides, so  $\neg(P_1 \vee \dots \vee P_n) \equiv (\neg P_1 \wedge \dots \wedge \neg P_n)$  ■

4. Theorem:  $P$  a predicate depends on variables  $x_1, \dots, x_n$

$Q$  denote  $\forall x_n \forall x_{n-1} \dots \forall x_1 P$  for all  $n \geq 1$ .

$$\neg Q = \exists x_n \exists x_{n-1} \exists \dots \exists x_1 \neg P$$

Proof:  $P$  is false iff  $\exists x_n \exists x_{n-1} \dots \exists x_1 P$  is false, so

$\neg P$  is true iff  $\exists x_n \exists x_{n-1} \dots \exists x_1 \neg P$  is true.

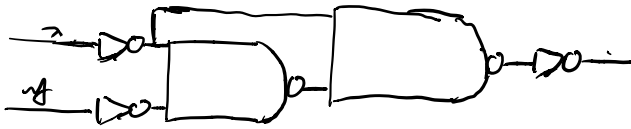
$Q$  is false iff  $\exists x_n \exists x_{n-1} \dots \exists x_1 P$  is false. so

$\neg Q$  is true iff  $\exists x_n \exists x_{n-1} \dots \exists x_1 \neg P$  is true. ■

5. (i)  $f(x, y) = xy + x\bar{y} + \bar{x}\bar{y}$ .

(ii)  $x \uparrow y = \overline{x \cdot y} = \bar{x} + \bar{y}$ .

$$f(x, y) = x + \bar{x}\bar{y} = x + \overline{\bar{x}\bar{y}} = \overline{\bar{x} \cdot \bar{x}\bar{y}}.$$



(iii).  $g(x, y) = x(\bar{x} + \bar{y})$   
 $= x\bar{x} + x\bar{y}$   
 $= 0 + x\bar{y}$   
 $= x\bar{y}$ .

Not the same.

6 (1) One which can be used to express all possible truth tables by combining members of the set into a Boolean expression.

(2)  $xy = (x+y) \cdot \overline{x \cdot y}$   
 $= \overline{\bar{x}\bar{y}} \cdot \overline{x \cdot y}$   
 $= \overline{\bar{x}\bar{y} \cdot xy}$

$\bar{a} + \bar{b} = \overline{ab}$ .

7. To prove  $f(x, y) \cdot g(x, y)$  is satisfiable when both  $f(x, y)$ ,  $g(x, y)$  are satisfiable, is equal to prove its contrapositive: If either  $f(x, y)$ ,  $g(x, y)$  is not satisfiable or both  $f(x, y)$  and  $g(x, y)$  are unsatisfiable, then  $f(x, y) \cdot g(x, y)$  is not satisfiable.

when either  $f(x, y)$ ,  $g(x, y)$  is not satisfiable, the truth

value of  $f(x,y) \cdot g(x,y)$  is 0

when both  $f(x,y), g(x,y)$  are not satisfiable, the truth value of  $f(x,y) \cdot g(x,y)$  is 0.

Thus, the contrapositive is true. So  $f(x,y) \cdot g(x,y)$  is satisfiable when both  $f(x,y), g(x,y)$  are satisfiable.

8. Assume  $F(n) = 5^n - 1$ .

Base case:  $n=1$   $5^1 - 1 = 4$  can be divided by 4.

when  $n \geq 2$ ,  $F(n) - F(n-1) = 5^n - 1 - 5^{n-1} + 1 = 4 \cdot 5^{n-1}$ .

so  $k$  can be  $5^{n-1}$  that  $5^n - 1$  can be divided by 4 for any  $n \geq 2$ .

So  $5^n - 1$  can be divided by 4.