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Tutorial 04: Rounding and Normalization

Computer Science Department

CS2208: Introduction to Computer Organization and Architecture

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Rounding

- □ The rounding mechanisms include
 - Truncation (i.e., dropping unwanted bits) by rounding towards zero; a.k.a., rounding down
 - Rounding towards positive or negative infinity, the nearest valid floating-point number in the direction positive or negative infinity, respectively, is chosen to decide the rounding; a.k.a., rounding up.
 - o *Rounding to nearest*, the closest floating-point representation to the actual number is used.

Rounding

Example 1: Round to the nearest the following numbers value to 8 digits after the binary point.

```
0.110101011001000 ==> 0.11010101
                                           0.110101001001000 ==> 0.11010100
                        0.0000001
                                                                + 0.00000001
                         0.11010110
                                                                  = 0.11010101
                                                                   If it is == case.
                                                                   and this hit = 0.
                       you round up.
                                                                  vou round down
0.110101011000000
                                           0.110101001000000
                    ==> 0.11010101
                                                               ==> 0.11010100
                        0.0000001
                                                                 + 0.0000000
                                                 1000000
                        0.11010110
                                                                  = 0.11010100
                   Mid-way → round to even significand
                                                 100000
0.110101010_{xxxxxx} ==> 0.11010101
                                           0.1101010000xxxxxx ==> 0.11010100
                        0.0000000
                                                                   0.0000000
   xxxxxx0
                         0.11010101
                                                                    0.11010100
                                                 \mathbf{0}XXXXXX
   1000000
                                                 1000000
```

Normalization

■ Example 2: Convert the unsigned value AB.BA₁₆ to binary. Normalize your answer.

AB.BA₁₆

 \rightarrow 10101011.10111010₂

After normalization,

 \rightarrow 1.010101110111010₂ × 2⁺⁷

In base b, a normalized number will have the form $\pm b_0 \cdot b_1 b_2 b_3 ... \times b^n$ where $b_0 \neq 0$, and b_0 , b_1 , b_2 , b_3 ... are integers between 0 and b -1

0 = 0000

1 = 0001

2 = 0010

3 = 0011

4 = 0100

5 = 0101

6 = 0110

7 = 0111

8 = 1000

9 = 1001

A = 1010

B = 1011

C = 1100

D = 1101

E = 1110

F = 1111

Normalization and Rounding

Example 3: Consider the unsigned normalized binary value $1.0101011110111010_2 \times 2^{+7}$

```
limit it (using truncation / rounding down) to 6 bits (1 +
                                                            5 bits) in total
                                               6 bits (1 + 5 bits) in total 4 = 0100
limit it (using rounding up)
limit it (using rounding to the nearest) to 6 bits (1 + 5 \text{ bits}) in total 5 = 0101
limit it (using truncation / rounding down) to 9 bits (1 + 8 bits) in total
                                    to 9 bits (1 + 8 bits) in total
limit it (using rounding up)
                                            to 9 bits (1 + 8 \text{ bits}) in total 8 = 1000
limit it (using rounding to the nearest)
limit it (using truncation / rounding down) to 14 bits (1 + 13 bits) in total
                                            to 14 bits (1 + 13 bits) in total
limit it (using rounding up)
                                            to 14 bits (1 + 13 bits) in total
limit it (using rounding to the nearest)
```

Calculate the rounding error in each case.

```
Note that: The binary value 1.010101110111010<sub>2</sub> \times 2<sup>+7</sup>
                                 = 10101011.101111010_{2} = AB.BA_{16}
```

$$2 = 0010$$

$$3 = 0011$$

$$\frac{1}{1}4 = 01$$

$$6 = 0110$$

$$C = 1100$$

$$E = 1110$$



Normalization and Rounding

- Limiting the answer to 6 bits (1 + 5) in total, $\rightarrow 1.010101110111010_2 \times 2^{+7}$
- \rightarrow 1.01010₂ × 2⁺⁷ (using truncation / rounding down)
- ⇒ 10101000_{2} ⇒ $A8_{16}$ *Truncation* error = $AB.BA_{16} A8_{16} = 3.BA_{16}$
- $\rightarrow 1.01011_2 \times 2^{+7}$ (using rounding up) $\rightarrow 10101100_2 \rightarrow AC_{16}$
- **Rounding up** error = $AB.BA_{16} AC_{16} = -0.46_{16}$
- \blacksquare As $11101111010_2 > 1000000000_2$
- \rightarrow 1.01011₂ × 2⁺⁷ (using rounding to the nearest) → $101011\overline{00}_{2}$ → AC_{16}
- **Rounding to the nearest** error = $AB.BA_{16} AC_{16} = -0.46_{16}$

```
0 = 0000
1 = 0001
2 = 0010
3 = 0011
4 = 0100
5 = 0101
6 = 0110
7 = 0111
8 = 1000
9 = 1001
A = 1010
B = 1011
C = 1100
```

D = 1101

E = 1110

F = 1111



Normalization and Rounding

- Limiting the answer to 9 bits (1 + 8) in total, $\rightarrow 1.0101011101111010_2 \times 2^{+7}$
- \rightarrow 1.01010111₂ × 2⁺⁷ (using truncation / rounding down)
 - $\rightarrow 10101011.1_{2} \rightarrow AB.8_{16}$
- **Truncation** error = $AB.BA_{16} AB.8_{16} = 0.3A_{16}$
- \rightarrow 1.01011000₂ × 2⁺⁷ (using rounding up) → 10101100.0_{2} → AC_{16}
- **Rounding up** error = $AB.BA_{16} AC_{16} = -0.46_{16}$
- \blacksquare As $01111010_2 < 1000000_2$
- \rightarrow 1.01010111₂ × 2⁺⁷ (using rounding to the nearest)
 - $\rightarrow 10101011.1_{2}^{2} \rightarrow AB.8_{16}$
- **Rounding to the nearest** error = $AB.BA_{16} AB.8_{16} = 0.3A_{16}$

- 0 = 0000
- 1 = 0001
- 2 = 0010
- 3 = 0011
- 4 = 0100
- 5 = 0101
- 6 = 0110
- 7 = 0111
- 8 = 1000
- 9 = 1001
- A = 1010
- B = 1011
- C = 1100
- D = 1101
- E = 1110
- F = 1111



- Limiting the answer to 14 bits (1 + 13) in total, $\rightarrow 1.010101110111010_2 \times 2^{+7}$

- ⇒ 10101011.1011110_{2}^{2} ⇒ $AB.B8_{16}^{16}$ *Truncation* error = $AB.BA_{16} AB.B8_{16} = 0.02_{16}$
- \rightarrow 1.01010111011111₂ × 2⁺⁷ (using rounding up)
- ⇒10101011.1011111 $_{2}$ → AB.BC $_{16}$ Rounding up error = AB.BA $_{16}$ AB.BC $_{16}$ = -0.02 $_{16}$
- \blacksquare As $10_2 == 10_2$
- \rightarrow 1.0101011110110₂ × 2^{+7} (using rounding to the nearest)
- ⇒ 10101011.110110 $_{2}$ ⇒ AB.B8 $_{16}$ Rounding to the nearest error = AB.BA $_{16}$ AB.B8 $_{16}$ = 0.02 $_{16}$
- Which rounding mechanism produces less error?

- 0 = 0000
- 1 = 0001

- 4 = 0100
- 5 = 0101
- 6 = 0110
- 7 = 0111
- 8 = 1000
- 9 = 1001
- A = 1010
- B = 1011
- C = 1100
- D = 1101
- E = 1110
- F = 1111