

Puzzle:  $S = \{A \mid A \notin A\}$ .  $S \in S$ ?  $S \notin S$ ?

$A \in S \iff A \notin A$

$S \in S \iff S \notin S$ .

Russell's paradox, 1901.

Zermelo, 1908: only allow  $\{x \in B \mid P(x)\}$   
in some set.

Ex 3.4.5:  $A \cap (B \setminus C) = (A \cap B) \setminus C$ .

Given

Goal

$\forall x (x \in A \cap (B \setminus C) \iff x \in (A \cap B) \setminus C)$

1.  $x \in A \cap (B \setminus C)$

1.  $x \in (A \cap B) \setminus C$

$x \in A \wedge x \in (B \setminus C)$

$x \in A \wedge x \in B \wedge x \notin C$

$x \in A \wedge (x \in B \wedge x \notin C)$

$x \in A \wedge x \in B \wedge x \notin C$

( $\rightarrow$ ) Proof:

Let  $x$  be arbitrary. We have to show that  $A \cap (B \setminus C) \iff x \in (A \cap B) \setminus C$ . Both of them are equivalent to  $x \in A \wedge x \in B \wedge x \notin C$  by associativity. So they are equivalent to each other. So  $A \cap (B \setminus C) = (A \cap B) \setminus C$ .  $\square$

§ 3.5.

Proofs involving

$\vee$

Strategy: To use the given of  $P \vee Q$ , break the proof into two parts: 1) Case 1: suppose  $P$  is true

2) Case 2:  $Q$

Form:

Case 1: Assume  $P$ . [Proof]

Case 2: Assume  $Q$ .

Since  $P \vee Q$ , these cases are exhaustive, so we have proof our goals.

Ex 2:  $(A \cup B) \setminus C \subseteq A \cup (B \setminus C)$

Given

Goal

$x \in (A \cup B) \setminus C$

$\forall x (x \in (A \cup B) \setminus C \rightarrow x \in A \cup (B \setminus C))$ .  
 $\hookrightarrow x \in A \cup (B \setminus C)$



$\neg(P \vee Q) \rightarrow \neg P \wedge \neg Q$

Strategy: To prove a goal of the form  $P \vee Q$ , add  $\neg P$  to given and try to prove  $Q$ .

Case:  $P \checkmark$

Case:  $\neg P \rightarrow Q \rightarrow P \vee Q$

Ex:  $\forall x \in \mathbb{R}$ , if  $x^2 \geq x$  then  $x \leq 0$  or  $x \geq 1$ .

Proof: Let  $x \in \mathbb{R}$  assume  $x^2 \geq x$ . If  $x \leq 0$ , then we're done. So assume  $x > 0$ . Dividing both side by  $x$  giving  $x \geq 1$ . So either  $x \leq 0$  or  $x \geq 1$ .  $\square$ .