Computer Science Fundamentals II   Sunday, December 8, 2019

Final Exam

CS1027A   *University of Western Ontario*

Last Name: ———————————————

First Name: ———————————————

Student Number: ————————————

Section Number (1 - Sarlo, 2 - Bloch-Hansen): ———

Exam Code: 111

PART I		10
PART II		20
PART III		24
PART IV		35
PART V		24
PART VI		24
BONUS		8
Total		137

Instructions

- Fill in your name, student number, and section.

- You are allowed a 8.5 by 11 inch one-sided (hand-written) cheat sheet.

- Verify that your exam paper has 33 pages and 57 questions.

- The exam is 3 hours long and is scored out of 137 marks. There is a bonus question worth 8 marks, which can bring your score above perfect.

- **PART I:** True/false questions. For each question circle only **one** answer.

- **PART II:** Multiple choice questions. For each question circle only **one** answer.

- **PART III:** Analysis of Algorithms questions.

- **PART IV:** Binary tree questions.

- **PART V:** Sorting questions.

- **PART VI:** Algorithm design questions.

- **BONUS:** Bonus question.

- When you are done, raise your hand and one of the TA's will collect your exam.

1 (10 marks) Part I. True/False Questions

Each true/false question is worth 1 point; circle **only one** answer.

1. Insertion sort can work on lists that contains duplicate values.
   (A) True   (B) False

2. Binary tree nodes can have at most 2 descendants.
   (A) True   (B) False

3. Unordered Lists must be implemented with linked lists and Ordered Lists must be implemented with arrays.
   (A) True   (B) False

4. Stacks are typically used for implementing the level-order traversal of a tree.
   (A) True   (B) False

5. An inorder traversal of a binary search tree will display the nodes in ascending order.
   (A) True   (B) False

6. Postorder traversal of binary trees means looking at the right child before the left child.
   (A) True   (B) False

7. The *Comparable<T>* interface's method, *compareTo(T element)*, returns a *boolean* variable to indicate if the current object (*this*) is smaller than element.
   (A) True   (B) False

8. The best case of a stack-based insertion sort is when the input array is already ordered but in reverse (descending) order.
   (A) True   (B) False

9. The time complexity of *push(T element)* is always the same for any implementation of a stack.
   (A) True   (B) False

10. The following diagram depicts a legal step partway through the process of performing a stack-based insertion sort on the array, A = [5, 8, 2, 6, 9, 1].
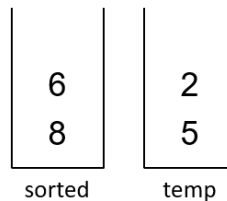   (A) True   (B) False

Figure 1

2 (20 marks) Part II. Multiple Choice Questions

Each multiple choice question is worth 1 point; circle **only one** answer.

11. Consider the following circular array queue. Indicate what would be printed when executing the following code (assuming array is the variable containing the elements):
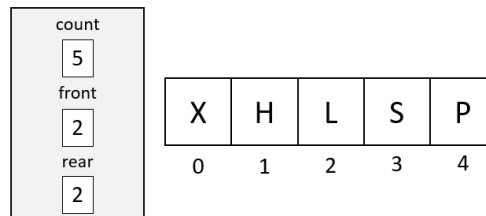

Figure 2

```
1  for (int i = 0; i < count; i++) {
2    System.out.print(array[i] + " ");
3  }
```

(A) X H L S P   (B) L S P   (C) L S P X H   (D) S P X H L
(E) None of the above

12. There is an office Christmas party and all the attendees are encouraged to shake hands with one another. Given an arbitrary number of people n at the party, where each pair of people shake hands exactly one time, calculate the exact number of handshakes that happen at the party.
(A) 1   (B) n (C) n^2 (D) $\log(n)$ (E) $1 + 2 + ... + n - 1$

13. There is an office Christmas party and all the attendees are encouraged to shake hands with one another. Given an arbitrary number of people n at the party, where each pair of people shake hands exactly one time, determine the Big-Oh time complexity for this sequence of handshakes.
(A) $O(1)$ (B) $O(n)$ (C) $O(n^2)$ (D) $O(n \cdot \log(n))$

14. Rank the following time complexities, labelled as i) through iv), in order from smallest to greatest.
i) $O(n^2)$
ii) $O(n \cdot \log(n))$
iii) $O(2^n)$
iv) $O(n)$
(A) iv < i < iii < ii
(B) iv < ii < i < iii
(C) iii < ii < i < iv
(D) ii < iv < i < iii

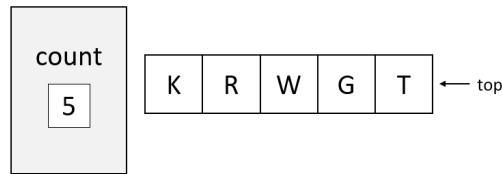15. Consider the following stack. What would be returned when executing *pop()*?

Figure 3

(A) 5   (B) K   (C) R   (D) T   (E) An exception would be thrown

16. Consider the following code fragment:

```
1  public class Foo {
2     private static Person p;
3     public static bar() {
4        String[] arr = new String[1];
5        try { arr[0] = p.getLastName(); }
6        catch (ArrayIndexOutOfBoundsException e) { System.out.println("A"); }
7        finally { System.out.println("B"); }
8        System.out.println("C");
9     }
10 public static void main(String[] args) {
11    try { bar(); }
12    catch (NullPointerException e) { System.out.println("D"); }
13 }
```

What does the code fragment above output when it is executed?
(A) "A"   (B) "B"   (C) "C"   (D) "D"
(E) "AB"   (F) "ABC"   (G) "BD"   (H) "BCD"

17. The following values are inserted, in the given order, into an initially empty binary search tree: 7, 3, 5, 4, 11, 6. Which value is stored in a node in level 2 of the tree?
(A) 3   (B) 5   (C) 4   (D) 6

18. Consider the following algorithm:

```
1  private void foo(int size) {
2     if (size == 0) return;
3     else foo(size - 3);
4  }
```

How many activation records (or call frames) are created in the execution stack (or runtime stack, or call stack) for method *foo()* when the statement *foo(12);* is executed? Do not include activation records for other methods such as *main()*.
(A) 0   (B) 4   (C) 5   (D) 6

4

19. Consider the following code fragment:

```
1  curr = front;
2  next = curr.getNext();
3  prev = null;
4  while (curr != null) {
5     (*)
6  }
7  front = prev;
```

Which code must be added in the part marked (*) so the above code correctly inverts a non-empty singly linked list? See Figure 4 to understand what "invert" means.
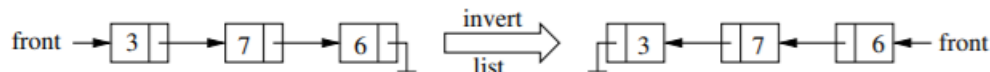

Figure 4

(A) *next.setNext(prev); prev = curr; curr = next; if (next != null) next = next.getNext();*
(B) *curr.setNext(prev); prev = curr; curr = next; if (next != null) next = next.getNext();*
(C) *next.setNext(curr); prev = curr; curr = next; if (next != null) next = next.getNext();*
(D) *prev = curr; curr = next; curr.setNext(prev); if (next != null) next = next.getNext();*

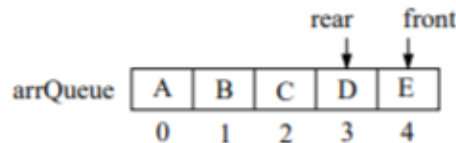20. Consider the following queue implemented using a circular array:

Figure 5

Which code fragment correctly implements the *dequeue()* operation on this particular queue? The name of the array storing the data items in the queue is *arrQueue*.

(A) *r = arrQueue[front]; front = front + 1; count--; return r;*
(B) *r = arrQueue[rear]; rear = rear - 1; count--; return r;*
(C) *r = arrQueue[front]; front = (front + 1) % arrQueue.length; count--; return r;*
(D) *r = arrQueue[front]; rear = (rear + 1) % arrQueue.length; count--; return r;*
(E) *r = arrQueue[front]; front = (front - 1) % arrQueue.length; count--; return r;*

5

21. The largest value in a binary search tree is always located in:
    (A) The left-most node
    (B) The right-most node
    (C) The bottom-most node
    (D) The top-most node
    (E) It depends on the order the elements were added

22. Which operation always has a time complexity of O(1)?
    (A) Enqueueing an item to a linked queue
    (B) Removing an item from a linked list
    (C) Adding an item to an ordered list
    (D) Copying elements from one array to another
    (E) Both A and B

23. Consider the following code fragment:

```
1  public class classA<T> {
2    public T smaller(T elem1, T elem2) {
3      if (elem1.compareTo(elem2) < 0) return elem1;
4      else return elem2;
5    }
6  }
```

Which of the following statements is correct?
(A) Line 1 has a compilation error
(B) Line 3 has a compilation error
(C) There are no compilation errors, but the above code will produce a runtime error.
(D) There are no errors in this code.

24. Consider the following code fragment:

```
1  public class A {
2    private int x;
3    public String s = "";
4    public A(int y) { x = y }
5  }
```

Which of the following statements is correct?
(A) The empty string referenced by s is allocated memory in the heap, while x, s, and y are allocated memory in the execution stack (or runtime stack or call stack).
(B) The empty string referenced by s and s are allocated memory in the heap, while x and y are allocated memory in the execution stack.
(C) The empty string referenced by s, s, and x are allocated memory in the heap, while y is allocated memory in the execution stack.
(D) x and s are allocated memory in the heap, while y is allocated memory in the execution stack. No memory needs to be allocated for the empty string.

25. What is the last node visited in a postorder traversal?
    (A) Left-most node at the bottom
    (B) Right-most node at the bottom
    (C) Root node
    (D) It depends on the structure of the tree

26. Array-based stacks are often implemented with a default of *top* = 0. Which of the following are correct interpretations of what *top* represents (in general) when an empty stack starts with *top* at 0?
    (A) The stack is empty
    (B) The number of elements in the stack
    (C) The maximum capacity of the stack
    (D) The index in the array where the next element will be pushed
    (E) Both B and D
    (F) None of the above

27. Consider the following *toString()* method. What data structure or ADT does this method correspond to?

```
1  public String toString() {
2    String str = "";
3    Iterator<T> iter = this.iterator();
4    while (iter.hasNext()) { str += iter.next() + " "; }
5    return str;
6  }
```

    (A) Stack   (B) Queue   (C) Linked List   (D) Tree   (E) All of the above

    Consider the following code being executed on an existing array-based stack called s:

```
1  for (int i = 0; i < s.size(); i++)
2    s.push(s.pop());
```

28. Use the stack code shown above. What happens to s as a result of executing the above loop? (Compare s before the algorithm to s after) Assume that s is not empty.
    (A) Nothing   (B) s is reversed   (C) s is deleted   (D) An exception is thrown

29. Use the stack code shown above. What happens to s as a result of executing the above loop? (Compare s before the algorithm to s after) Assume that s is empty.
    (A) Nothing   (B) s is reversed   (C) s is deleted   (D) An exception is thrown

30. Use the stack code shown above. What happens if s was a linked stack instead of an array stack? Does it change the outcome?
    (A) Yes   (B) No

3 (24 marks) Part III. Analysis of Algorithms

For each one of the following 4 questions, compute the time complexity of the given code. You **must explain how you computed the time complexity** and you must give the order (big-Oh) of the time complexity. For method *mystery()*, assume it was called on the root of a tree containing n nodes.

31. (3 marks)

```
1  int i, j, k, x = 0;
2  for (i = n; i > 0; i--) {
3    k = i * n;
4    for (j = k; j < n; j++) {
5      x = x + j;
6    }
7    i = i - k;
8  }
```

32. (3 marks)

```
1  for (int i = 0; i < n; i++) {
2    for (int j = 0; j < 5; j++) {
3      for (int k = 0; k < n; k++) {
4        if ((i != j) && (i != k)) {
5          System.out.println(k);
6        }
7      }
8    }
9  }
```

33. (3 marks)

```
1  public String modify(String str) {
2    if (str.length() <= 1) return "";
3    int half = str.length() / 2;
4    modify(str.substring(half));
5  }
```

34. (3 marks)

```
1  public void mystery(BinaryTreeNode<T> node) {
2    if (node == null) return;
3    if (node.getLeft() != null || node.getRight() != null)
4      System.out.println(node.getElement());
5    mystery(node.getLeft());
6    mystery(node.getRight());
7  }
```

35. (3 marks) Write **exactly one** for loop whose start condition is $int\ i = 0;$ and whose end condition is $i < n;$ such that your code fragment has a time complexity of $O(n^2)$.

36. (3 marks) Write **exactly three** for loops that are nested inside of each other whose start conditions are $int\ i = 0;$, $int\ j = 0;$, and $int\ k = 0;$, respectively, and whose end conditions are $i < n;$, $j < n;$, and $k < n;$, respectively, such that your code fragment has a time complexity of $O(n)$.

37. (3 marks) Consider a sorted list implemented using a circular array. Let *first* be the index of the first element in the list and *last* be the index of the last element in the list. Let the list store n integer values, where n is an odd number. Operation *median()* returns the value of the median or the $\frac{n+1}{2}$-th smallest element in the list. For example, the median of this sorted list: 1, 4, 5, 8, 9, 12, 15, is 8. What is the time complexity of the best possible implementation of operation *median()*? Explain how it would work.

38. (3 marks) Consider the same previous question, but this time the sorted list is implemented using a doubly linked list. Reference variable *first* points to the first node in the list and *last* points to the last node in the list. What is the time complexity of the best possible implementation of operation *median()*? Explain how it would work.

4 (35 marks) Part V. Binary Trees

39. For each of the following tree diagrams, determine whether or not the tree is a binary search tree. If it is not a binary search tree, explain why.
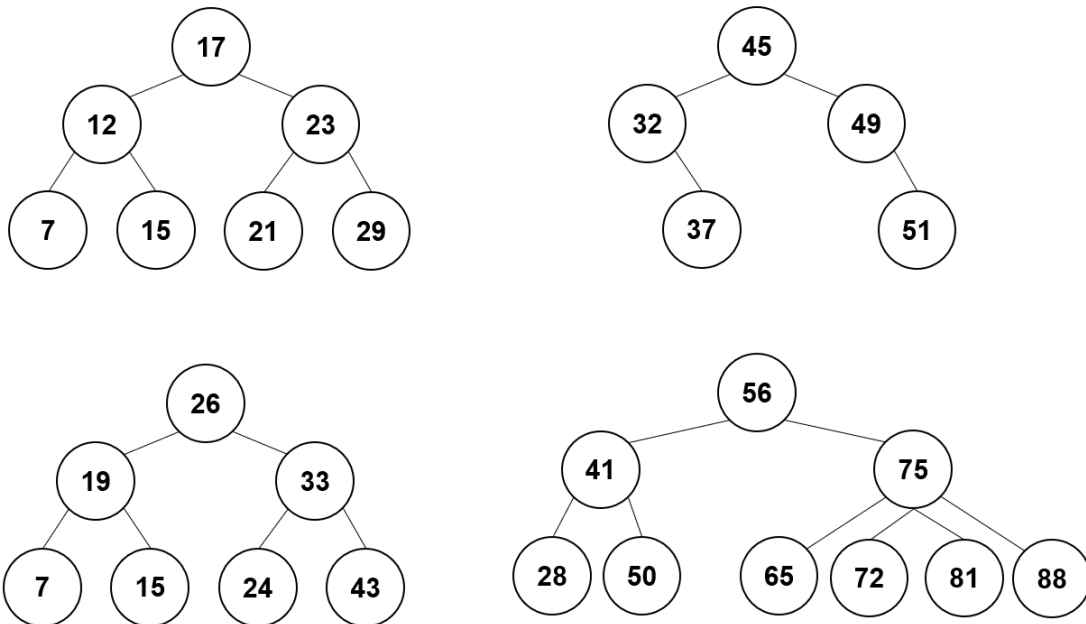
Figure 6

(Top-left) (1 mark)

(Top-right) (1 mark)

(Bottom-left) (1 mark)

(Bottom-right) (1 mark)

40. Consider the following tree:
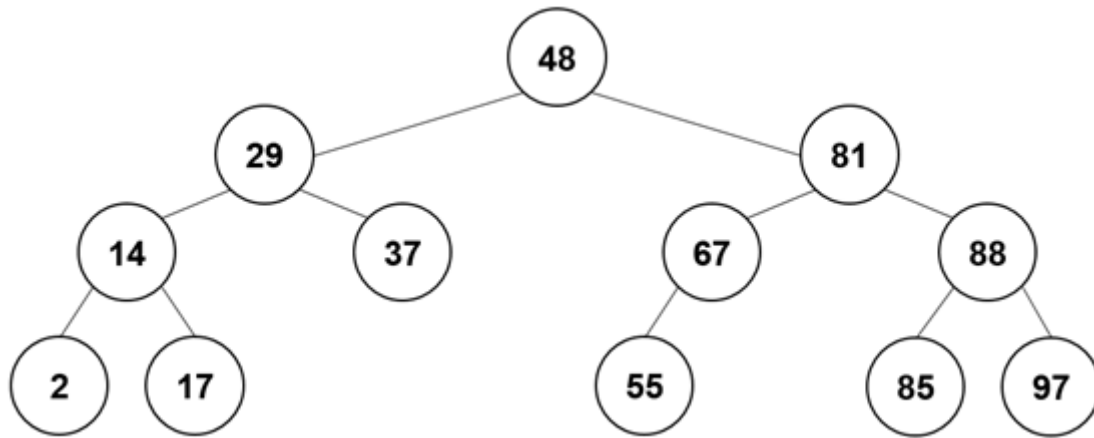
Figure 7

a) (1 mark) What is the height of this tree?

b) (1 mark) What is the level of the node containing "17"?

c) (1 mark) What is the value contained in this tree's root node?

d) (1 mark) What are the ancestors of the node containing "81"?

e) (1 mark) How many leaf nodes are contained in this tree?

f) (1 mark) What is the degree of the node containing "67"?

g) (1 mark) What is the degree of this tree?

h) (1 mark) Trace out the pre-order traversal of this tree.

i) (1 mark) Trace out the in-order traversal of this tree.

j) (1 mark) Trace out the post-order traversal of this tree.

k) (1 mark) Trace out the level-order traversal of this tree.

41. (3 marks) Draw the binary tree from the given level-order traversal. We have not included "null" for every empty position in every level of the tree; instead, we have written null when an **existing** node has no child in that position.

   - 27, 14, 11, null, 5, 21, 32, 19, null, null, 7, 36, null, null, null, null, null, 10, null, null, null

42. (3 marks) Draw the binary tree such that:

   - a pre-order traversal visits the nodes in this order: E, A, D, B, C, and
   - a post-order traversal visits the nodes in this order: A, B, C, D, E.

43. (4 marks) Consider the binary tree shown in Figure 8. Perform the iterative pre-order tree traversal using a stack as shown to you in class. Write your answers in the U-shaped diagrams in Figure 9. For the first stack, show the value on the stack before entering the loop. For the rest of the stacks, show the values on the stack at the end of each loop iteration.
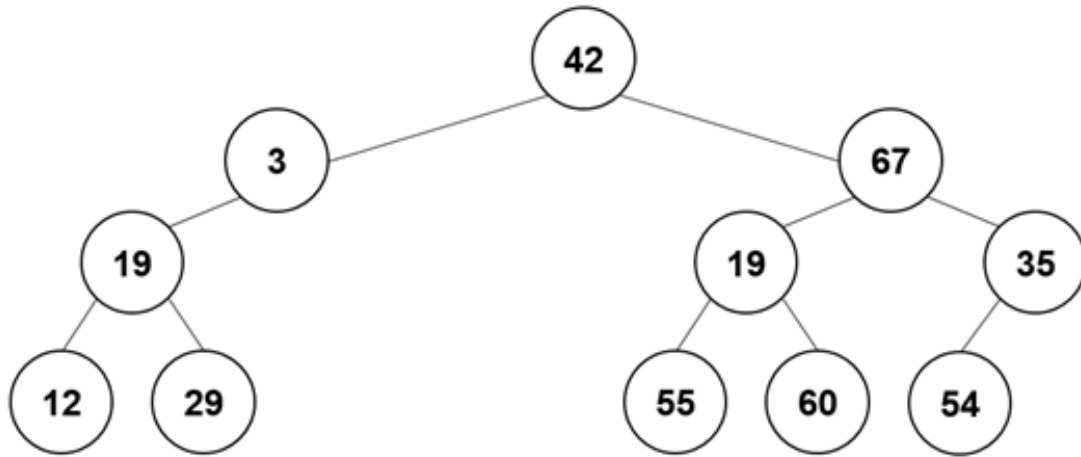
Figure 8

Visited nodes: ____ ____ ____ ____ ____ ____ ____ ____ ____ ____ ____

Figure 9

15

44. (4 marks) Evaluate the following expression tree. Show your work by drawing on the Figure (write numbers above nodes where appropriate, as shown in class).
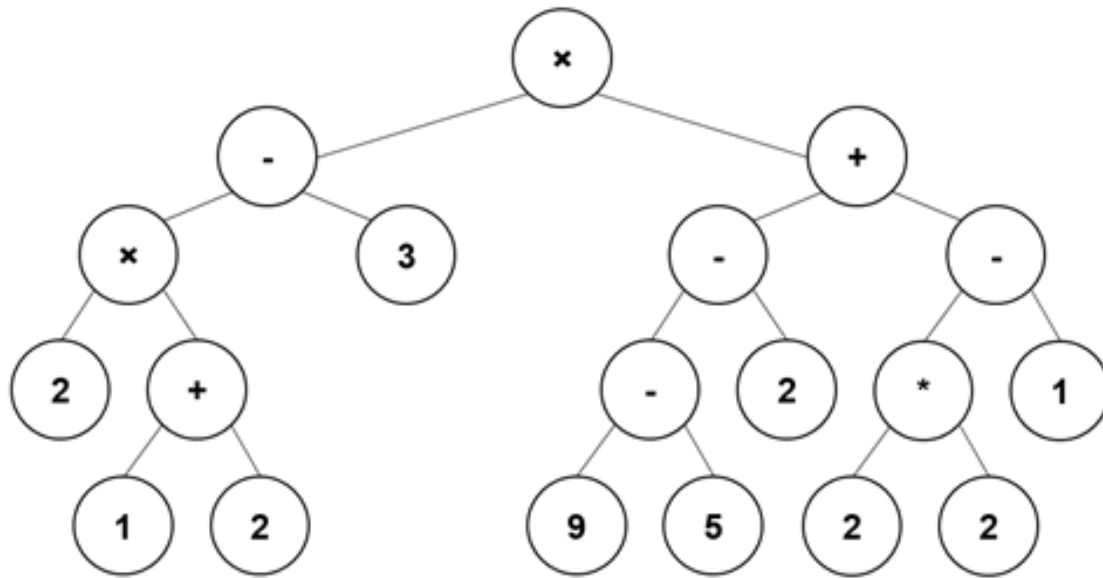
Figure 10

45. (1 mark) What kind of traversal did you use to evaluate the expression tree?

46. (1 mark) Will the root of an expression tree always be the final operator?

In the following 3 questions, explain the general approach to building an expression tree using a stack.

47. (1 mark) What is the input to the algorithm that builds an expression tree?

48. (1 mark) What variable type is being stored in the stack?

49. (2 marks) When building an expression tree, explain what happens when you come across an operator from the input.

16

5 (24 marks) Part VI. Sorting

50. (7 marks) Consider the array $A = [7, 9, 2, 5]$. Trace through a stack-based implementation of selection sort on this array. Assume that during each step, you scan through the entire array to find the minimum value, then replace that index of the array with *null*. After finding the minimum value in each step (write down what this number is), draw 3 stacks: (1) the *sorted* stack with the new number in it, (2) the *temp* stack holding whatever numbers are necessary to have drawn the previous *sorted* stack, and (3) the *sorted* stack with everything from *temp* in it.

51. (1 mark) Analyze this algorithm using Big-Oh notation.

52. (8 marks) Consider the array $A = [18, 9, 7, 2, 15, 12, 5]$. Trace through the quicksort algorithm on this array. Use the **last** element as the pivot. For each recursive call, draw the execution stack (don't worry about other methods such as $main()$), show the pivot, and show the arrays smaller, equal, and larger. During some of the recursive calls some of the sub-arrays will be adjusted so that they become sorted, make sure this happens in the correct recursive call!

53. (8 marks) *MergeSort* is a divide-and-conquer algorithm based on the idea of breaking down a list into several sub-lists until each sub-list consists of a single element. These sub-lists are then merged together in a manner that forms a larger sorted sub-list. Eventually, the algorithm will have sorted the list and will terminate.

Consider the following pseudocode for mergesort:

```
1  MergeSort(int[] A, int left, int right) {
2    if (right > left) {
3      int middle = (left + right) / 2;
4      MergeSort(A, left, middle);
5      MergeSort(A, middle + 1, right);
6      Merge(A, left, middle, right);
7    }
8  }
```

We first call *MergeSort()* with the parameters $A = [3, 1, 7, 2]$, *left = 0*, and *right = A.length - 1*. Therefore, this first method call represents sorting the entire array; however, as this is a recursive method, our first step is to divide the full array into two sub-arrays. Thus, on line 4, we call *MergeSort()* again but adjust the value of *right* so that this second method call represents sorting only the first half of the entire array.

This process of dividing the array into smaller sub-arrays is continued until we are considering sub-arrays of size 1. This is the base case. Only after we have encountered the base case will the method *Merge* be called to piece the sub-arrays back together in a sorted manner.

You can assume that calling method *Merge()* will modify the array A such that the elements in the range of *left* to *right* are now a sorted sub-set.

Fill in the Figures on the following 2 pages as you trace through the algorithm on the input array $A = [3, 1, 7, 2]$. We have provided the state of the execution stack after every method invocation and after every method termination. You need to track the values of A, *left*, and *right* (and sometimes *middle*).

Hint. You might find it useful to write comments to yourself beside each activation record to remind yourself where in the function you are returning to. For example, the first activation record (provided for you) represents calling *MergeSort* the first time. The second activation record represents calling *MergeSort* from line 4 of the previous activation record.

	0	1	2	3
A	3	1	7	2

MergeSort | **main**
A [3] [1] [7] [2]   **left** [0]   **right** [3]

MergeSort | **MergeSort** | **main**
A [ ] [ ] [ ] [ ]   **left** [ ]   **right** [ ]

MergeSort | **MergeSort** | **MergeSort** | **main**
A [ ] [ ] [ ] [ ]   **left** [ ]   **right** [ ]

MergeSort | **MergeSort** | **main**
A [ ] [ ] [ ] [ ]   **left** [ ]   **right** [ ]

MergeSort | **MergeSort** | **MergeSort** | **main**
A [ ] [ ] [ ] [ ]   **left** [ ]   **right** [ ]

MergeSort | **MergeSort** | **main**
A [ ] [ ] [ ] [ ]   **left** [ ]   **right** [ ]

Merge | **MergeSort** | **MergeSort** | **main**
A [ ] [ ] [ ] [ ]   **left** [ ]   **middle** [ ]   **right** [ ]

MergeSort | **MergeSort** | **main**
A [ ] [ ] [ ] [ ]   **left** [ ]   **right** [ ]

MergeSort | **main**
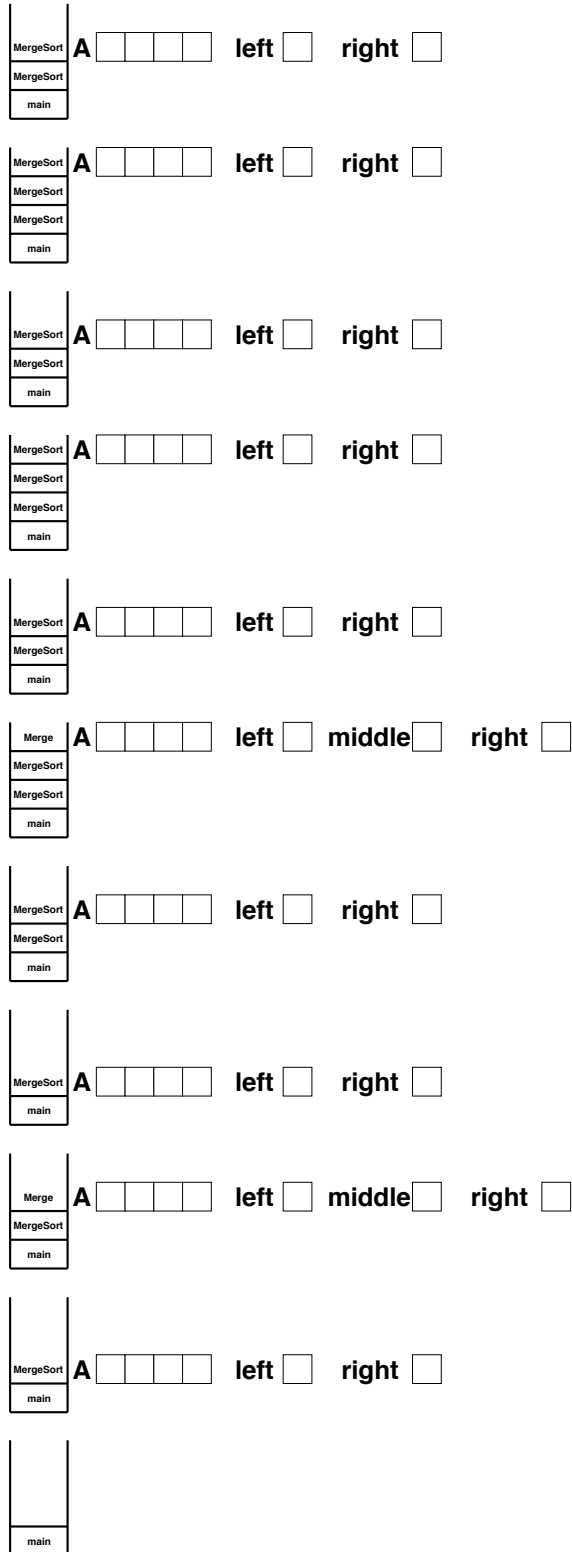A [ ] [ ] [ ] [ ]   **left** [0]   **right** [3]

Figure 11

Figure 12

6 (24 marks) Part VII. Algorithm Design

54. "Simon" is a memory-based game that requires its player to remember a given sequence of colors. It's an electronic game on a device that contains 4 colored areas that can light up. A typical game plays like this:

- A new color is added to the existing sequence of colors.
- The colors in the sequence are displayed for the player one at a time.
- The player then has to enter the sequence of colors that they can remember.
- If the player's answer was correct, repeat this process.

a) (1 mark) Which of the collection types that were taught in this course would be best suited for tracking the sequence of colors in the "Simon" game? Give a brief explanation for your response.

b) (6 marks) Write an algorithm *public void playSimon(int n)* in Java or in **detailed** Java-like pseudocode that implements the game as described above, where *n* is the maximum length of the color sequence for a particular game. You can make the following assumptions:

- Whichever collection you decide to use has already been declared as an instance variable, you just need to initialize it and use it.
- There exists a method *public Color getNextColor()* that returns the next color in the sequence.
- There is a method called *public void showColor(Color col)* that you can call to display the color for the player. You can assume that this method displays the color for an appropriate amount of time to the player.
- There exists a method *public boolean checkAnswer()* that will handle everything to do with checking the players response; if the player was correct, this method will return true. Otherwise, this method will return false.

You can use the following page to write your answer.

(This page is for the algorithm *public void playSimon(int n)*)

c) (1 mark) Analyze your algorithm using Big-Oh notation.

55. (8 marks) Consider a binary search tree formed by linking together node objects of class *BinaryTreeNode*. The *BinaryTreeNode* class provides methods *getLeft()*, *setLeft()*, *getRight()*, *setRight()*, and *getElement()*. You can create a new node by calling the constructor of this class using *BinaryTreeNode<>(element)* to create a new node storing value *element* whose left and right children are *null*.

Write an algorithm *public void addReverse(BinaryTreeNode<T> node, T element)* in Java or in **detailed** Java-like pseudocode that adds a new node storing *element* into the binary search tree so that the ordering of the binary search tree is backwards compared to how we described it in class. More specifically, in this binary search tree, every node should have the following properties:

- Nodes in the left sub-tree should contain larger values.
- Nodes in the right sub-tree should contain smaller values.

You can assume that the variable *element* is of a class that implements *Comparable<T>*.

Additionally, you must:

- Increment an instance variable called *size* when you add a node to the tree.
- Adjust the instance variable called *root* if adding the first node to the tree.
- Throw a *DuplicatedKeyException* if *element* is already in the tree. Assume that *DuplicatedKeyException.java* already exists.

You can use the following page to write your answer.

(This page is for the algorithm *public void addReverse(BinaryTreeNode<T> node, T element)*)

56. (8 marks) Consider a binary search tree formed by linking together node objects of class *BinaryTreeNode*. The *BinaryTreeNode* class provides methods *getLeft()*, *setLeft()*, *getRight()*, *setRight()*, and *getElement()*. You can create a new node by calling the constructor of this class using *BinaryTreeNode<>(element)* to create a new node storing value *element* whose left and right children are *null*.

Write an algorithm *public int countSingleChild(BinaryTreeNode<T> node)* in Java or in **detailed** Java-like pseudocode that will traverse this tree and keep count of how many nodes have **exactly 1 child**. The algorithm must work on any binary search tree.

For example, in the following binary search tree, there are 3 nodes which have exactly 1 child. Remember that your algorithm must work on **any** binary search tree.

Figure 13

You can use the following page to write your answer.

(This page is for the algorithm *public int countSingleChild(BinaryTreeNode<T> node)*)

7 (8 marks) BONUS

57. (8 marks) Let q be a queue storing n integer values. Write an algorithm *public void partition(Queue q, int target)* in Java or in **detailed** Java-like pseudocode that receives as parameters the queue q and an integer value *target* and it re-arranges the values in the queue so that all the values smaller than *target* appear before the value of *target* and all the values larger than *target* appear after the value of *target*. Your algorithm must handle the possibility of duplicate entries of the value of *target*.

For this algorithm you can use **one** auxiliary stack s. You **cannot** use any other auxiliary data structures. You can use the following queue methods: *dequeue()*, *enqueue()*, *isEmpty()*, and *size()*. You can also use the following stack methods: *push()*, *pop()*, and *isEmpty()*.

For example, for the following queue on the left side and *target = 6*, your algorithm should produce a queue like the one on the right.
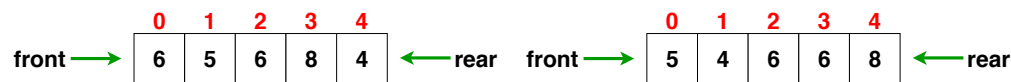

Figure 14

You can use the following page to write your answer.

(This page is for the algorithm *public void partition(Queue q, int target)*)

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