# **HPC File Systems**

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## Overview

**HPC Environment** 

File Systems

Lustre

**Best Practices** 

File Striping

File I/O

## **HPC Environment**

The typical HPC system has multiple file systems intended for different uses.

#### ► Home

Individual user home directories. Backed-up with quotas.

#### Scratch

Fast temporary access, not backed-up.

- Shared file system across all nodes.
- Local each node has it's own.
   (A lot of clusters do not put disks in the nodes)

#### Mass Storage

Long term storage, typically a tape system.

## File Systems

- NFS Network File System
- CXFS Clustered XFS
- PanFS Panasas ActiveScale File System
- GPFS General Parallel File System
- PVFS Parallel Virtual File System
- Lustre

#### Why?

- Spinning disks are slow.
- Serial I/O is even slower.

## Key Features

- Scalability.
   Can scale out to tens of thousands of nodes and petabytes of storage.
- Performance.
   Throughput of a single stream ~GB/s and parallel I/O ~TB/s.
- High availability.
- ► POSIX compliance.
- Supports ROMIO

#### Lustre

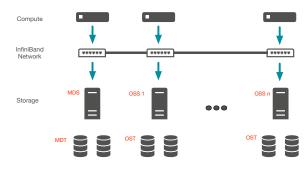
Approximately 50% of top 100 HPC's use Lustre. It consists of four components:

MDS Metadata Server

MDT Metadata Target

OSS Object Storage Server

OST Object Storage Target



## Metadata Server and Target

The MDS is a single service node that assigns and tracks all of the storage locations associated with each file in order to direct file I/O requests to the correct set of OSTs and corresponding OSSs.

The MDT stores the metadata, filenames, directories, permissions and file layout.

# Object Storage Servers and Targets

An OSS managers a small set of OSTs by controlling I/O access and handling network requests to them.

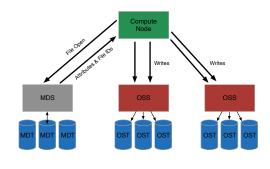
An OST is a block storage device. Often several disks in a RAID configuration.

## Typical Setup

- All nodes (login, compute, compile, ...) have the lustre file-system mounted at /lustre/.
- The number of servers and targets.
  - 2 MDSs (active, standby)
  - ▶ 10's OSSs
  - 100's OSTs
- Often giving a total of ~1PB of usable disk space.

## File Operations

- When a compute node needs to create or access a file, it requests the associated storage locations from the MDS and the associated MDT.
- I/O operations then occur directly with the OSSs and OSTs associated with the file bypassing the MDS.
- For read operations, file data flows from the OSTs to the compute node.



### **User Commands**

Lutre provides a utility to query and set access to the file system. They are all sub commands to the program 1fs.

For a complete list of available options.

```
login01 $ lfs help
```

▶ To get more information on a specific option.

```
login01 $ lfs help option
```

## Checking Diskspace

The lfs df command displays the file system disk space usage. Additional parameters can be specified to display inode usage of each MDT/OST or a subset of OSTs. The usage for the lfs df command is:

```
\label{login01 } $ lfs \ help \ df $$ Usage: \ df \ [-i] \ [-h] \ [--lazy|-1] \ [--pool|-p < fsname > [.< pool >] \ [path] $$ $$
```

#### Example, get a summary of the disk usage:

## Finding Files

The lfs find command is more efficient the GNU find. Example, finding fortran source files accessed within the last day.

```
login01 $ lfs find . -atime -1 -name '*.f90'
```

### Other 1fs Commands

- ▶ lfs cp to copy files.
- Ifs 1s to list directories and files.

These commands are often quicker as they reduce the number of stat and rpc calls needed.

#### **Avoid Wild Cards**

- tar and rm are inefficient when operating on a large class of files on lustre.
- The reason lies in the time it takes to expand the wild card.
- rm -rf \* on millions of files could take days, and impact all other users.
- Generate a list of files to be removed or tar-ed, and to act them one at a time, or in small sets.

```
login01 $ lfs find old/ -t f -print0 | xargs -0 rm
```

#### **Avoid Colons**

MPI-IO will split a path and file name on a colon to obtain a file system hint.

▶ If the colon is used in the time format specification, the interpretation of the hint will cause an error.

```
/lustre/tbrown/wrfout_2015_06_23_00:00:00.nc
```

Provide a hint that the underlying file system is lustre.

```
lustre:/lustre/tbrown/wrfout_2015_06_23_00_00_00.nc
```

# Limit Files Per Directory

- It is best to limit the number of files per directory.
- Writing thousands of files to a single directory produces a massive load on the MDSs, this often takes the file system offline.
- If you need to create a large number of files. Use a directory structure.
- ▶ A suggested approach is a two-level directory structure with  $\sqrt{N}$  directories each containing  $\sqrt{N}$  files, where N is the number of tasks.

## Read Only Access

- If a file is only going to be read, open it as O\_RDONLY.
- If you don't care about the access time, open it as O\_RDONLY|O\_NOATIME.
- If you need access time information and your doing parallel IO, let the master open it as 0\_RDONLY and all other ranks as 0\_RDONLY|0\_NOATIME.

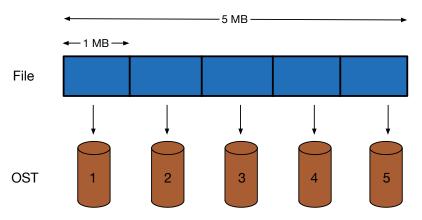
#### **Broadcast Stat**

- If many processes need the information from stat().
- Have the master process perform the stat() call.
- Then broadcast it to all processes.

```
#include <stdio h>
#include <stdlib.h>
#include <sys/stat.h>
#include <err h>
#include <mpi.h>
main(int argc, char **argv)
        int rank = 0;
        int len = 0:
        struct stat sbuf = {0}:
        MPI_Init(&argc, &argv);
        MPI Comm rank(MPI COMM WORLD, &rank);
        if (rank == 0) {
                if (stat(argv[0], &sbuf)) {
                        warn("Unable to stat %s", argv[0]);
                }
        len = sizeof(sbuf);
        MPI Bcast(&sbuf. len. MPI BYTE. 0. MPI COMM WORLD):
        MPI_Finalize();
        return(EXIT_SUCCESS);
```

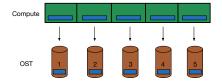
# File Striping

▶ A file is split into segments and consecutive segments are stored on different physical storage devices (OSTs).

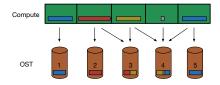


# Aligned vs Unaligned Stripes

Aligned stripes is where each segment fits fully onto a single OST. Processes accessing the file do so at corresponding stripe boundaries.



Unaligned stripes means some file segments are split across OSTs.



## Stripe Sizes

- You can get/set the stripe size, number of OSTs and which OST to start at.
- The stripe size must be a multiple of the maximum page size (64 KB).
- The typical default is
  - stripe count: 1
  - stripe size: 1048576 (1 MB)
  - stripe offset: -1 (MDS selects)

```
login01 $ lfs getstripe .
login01 $ lfs setstripe -s 32m -c 4 .
```

# Large File Stripe Sizes

- Set the stripe count of the directory to a large value.
- ► This spreads the reads/writes across more OSTs, therefore balancing the load and data.

```
login01 $ lfs setstripe -c 30 \
    /lustre/tbrown/large_files/
```

# Small File Stripe Sizes

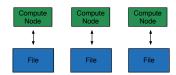
- Place small files on a single OST.
- This causes the small files not to be spread out/fragmented across OSTs.

```
login01 $ lfs setstripe -s 1m -c 1 \
    /lustre/tbrown/small_files/
```

## File I/O

Three cases of file I/O:

▶ Single stream.

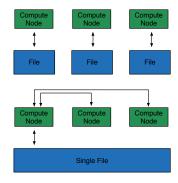


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Three cases of file I/O:

Single stream.

Single stream through a master.



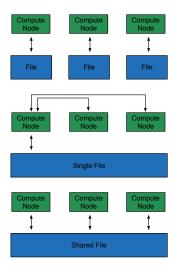
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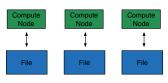
Single stream through a master.

Parallel.



## Single Stream IO

- Set the stripe count to 1 on a directory.
- Write all files in this directory.

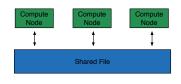


Otherwise set the stripe count to 1 for the file.

```
login01 $ lfs setstripe -s 1m -c 1 \
    /lustre/tbrown/serial/
```

## Parallel IO Stripe Count I

Single shared files should have a stripe count equal to the number of processes which access the file.



- If the number of processes is >160, set the count to -1, this will stripe across all OSTs (lustre has a max of 160).
- ► The stripe size should be set to allow as much stripe alignment as possible.
- Try to keep each process accessing as few OSTs as possible.

```
login01 $ lfs setstripe -s 32m -c 24 \
    /lustre/tbrown/parallel/
```

## Parallel IO Stripe Count II

You can specify the stripe count and size programmatically, by creating an MPI info object.

```
use mpi
use hdf5
implicit none
                  :: info
                                         ! MPI IO Info
integer
integer
                  ·· ierr
                                           ! Error status
integer(kind=hid_t) :: p_id, f_id
                                          ! Property and file id
character(len=256) :: filename,lcount,lsize ! Filename
! Init the HDF5 library
call h5open_f(ierr)
! Create an MPI object setting the strip size and count
call mpi info create(info, ierr)
write(lcount, '(I4)') 4
write(lsize, '(I8)') 4 * 1024 * 1024
call mpi_info_set(info, "striping_factor", trim(lcount), ierr)
call mpi_info_set(info, "striping_unit", trim(lsize), ierr)
! Set up the access properties
call h5pcreate_f(H5P_FILE_ACCESS_F, p_id, ierr)
call h5pset_fapl_mpio_f(p_id, MPI_COMM_2D, info, ierr)
! Open the file
call h5fcreate f(filename, H5F ACC TRUNC F, f id, ierr, &
                 access_prp = p_id)
```

# Questions? Survey

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