Lesson 2

Lists and the Collections Framework

Chapter Objectives

- □ The List interface
- Writing an array-based implementation of List
- Linked list data structures:
 - □ Singly-linked
 - □ Doubly-linked
 - □ Circular
- Big-O notation and algorithm efficiency
- Implementing the List interface as a linked list
- The Iterator interface
- Implementing Iterator for a linked list
- Generics
- The Java Collections framework (hierarchy)

3 Day - 1

Introduction to data structures

Section 2.0

What is data?

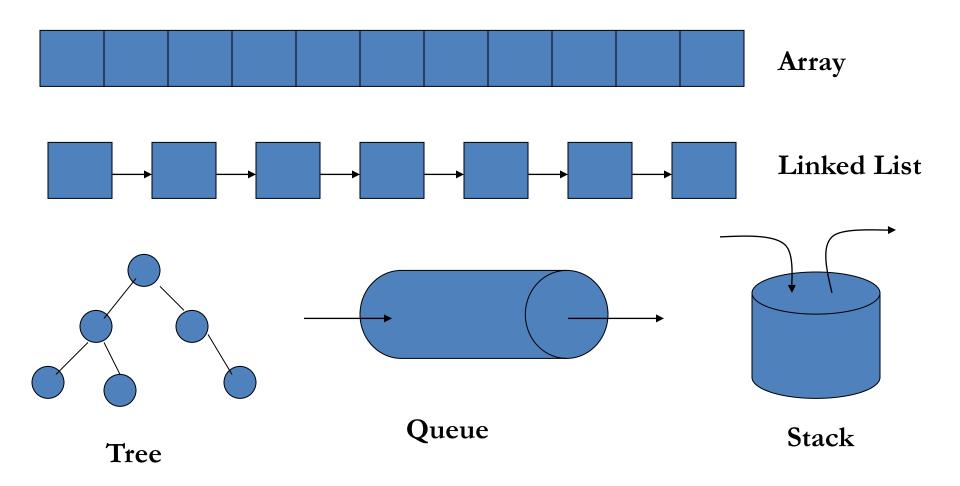
- Data
 - A collection of facts from which conclusion may be drawn
 - e.g. Data: Temperature 35°C; Conclusion: It is hot.
- Types of data
 - Textual: For example, your name (John)
 - Numeric: For example, your ID (090254)
 - Audio: For example, your voice
 - Video: For example, your voice and picture

What is data structure?

- □ A particular way of storing and organizing data in a computer so that it can be used efficiently and effectively.
 □ Data structure is the logical or mathematical model of a particular organization of data.
 □ A group of data elements grouped together under
 - For example, an array of integers

one name.

Types of data structures



There are many, but we named a few. We'll learn these data structures in great detail!

The need for data structures

- ☐ Goal: to **organize data**
- Criteria: to facilitate efficient
 - **storage** of data
 - retrieval of data
 - manipulation of data
- Design Issue:
 - select and design appropriate data types

(This is the main motivation to learn and understand data structures)

Data structure operations

□ Traversing

Accessing each data element exactly once so that certain items in the data may be processed

Searching

Finding the location of the data element (key) in the structure

lefta Insertion

Adding a new data element to the structure

Data structure operations (cont.)

☐ Deletion

Removing a data element from the structure

☐ Sorting

Arrange the data elements in a logical order (ascending/descending)

☐ Merging

 Combining data elements from two or more data structures into one

Linear or sequential search

- The sequential search (also called the linear search) is the simplest search algorithm.
- It is also the least efficient.
- It simply examines each element sequentially, starting with the first element, until it finds the key element or it reaches the end of the array.
- Demo : TestLinearSearch.java
- Animation : https://visualgo.net/en/list

Binary search

- The binary search is the standard algorithm for searching through a sorted sequence.
- It is much more efficient than the sequential search, but it does require that the elements be in order.
- It repeatedly divides the sequence in two, each time restricting the search to the half that would contain the element.
- You might use the binary search to look up a word in a dictionary.
- Demo : TestBinarySearch.java
- Animation : https://visualgo.net/en/bst

Ordered array & advantages

- used a binary search to locate the position where a new item will be inserted with the help of find() method.
- □ Items should be stored in an ordered manner.
- The major advantage is that search times are much faster than in an unordered array.
- The disadvantage is that insertion and deletion takes longer because all the data items must be moved up to make room.

Ordered array applications

- Ordered arrays are therefore useful in situations in which searches are frequent, but insertions and deletions are not.
- An ordered array might be appropriate for a database of company employees
- A retail store inventory, on the other hand, would not be a good candidate for an ordered array because the frequent insertions and deletions, as items arrived in the store and were sold, would run slowly.

Algorithm Efficiency and Big-O

Section 2.1

What is Algorithm?

An algorithm is a set of instructions to be followed to solve a problem.

- There can be more than one solution (more than one algorithm) to solve a given problem.
- An algorithm can be implemented using different programming languages on different platforms.

Algorithmic Performance

There are two aspects of algorithmic performance:

Time

Instructions take time.

How fast does the algorithm perform?

What affects its runtime?

Space

Data structures take space
What kind of data structures can be used?
How does choice of data structure affect the runtime?

We will focus on time:

How to estimate the time required for an algorithm

Analysis of Algorithms

When we analyze algorithms, we should employ mathematical techniques that analyze algorithms independently of specific implementations, computers, or data.

To analyze algorithms:

- First, we start to count the number of significant operations in a particular solution to assess its efficiency.
- Then, we will express the efficiency of algorithms using growth functions.

The Execution Time of Algorithms

Each operation in an algorithm (or a program) has a cost. Each operation takes a certain of time.

```
count = count + 1; \rightarrow take a certain amount of time, but it is constant
```

A sequence of operations:

Total Cost =
$$c1 + c2$$

The Execution Time of Algorithms (cont.)

Example: Simple Loop

```
i = 1;
sum = 0;
while (i <= n) {
    i = i + 1;
    sum = sum + i;
}</pre>
Cost
c1
1
1
c2
1
n
n
sum = c4
n
sum = sum + i;
c5
n
```

Total Cost =
$$c1 + c2 + (n+1)*c3 + n*c4 + n*c5$$

→ The time required for this algorithm is proportional to n

The Execution Time of Algorithms (cont.)

Example: Nested Loop

```
Cost
                                                                   <u>Times</u>
i=1;
                                                 с1
sum = 0;
                                                 с2
while (i \le n) {
                                                 с3
                                                                      n+1
         j=1;
                                                 c4
                                                                      n
         while (i \le n)
                                                 С5
                                                                      n*(n+1)
              sum = sum + i;
                                                 С6
                                                                      n*n
              \dot{j} = \dot{j} + 1;
                                                 с7
                                                                      n*n
   i = i + 1;
                                                 С8
                                                                      n
```

Total Cost = c1 + c2 + (n+1)*c3 + n*c4 + n*(n+1)*c5+n*n*c6+n*n*c7+n*c8

→ The time required for this algorithm is proportional to n²

Algorithm Efficiency and Big-O

- Getting a precise measure of the performance of an algorithm is difficult
- Big-O notation expresses the performance of an algorithm as a function of the number of items to be processed
- This permits algorithms to be compared for efficiency
- For more than a certain number of data items, some problems cannot be solved by any computer

Big-O Notation

- The O() can be thought of as an abbreviation of "order of magnitude" to describe the performance of an Algorithm.
- A simple way to determine the big-O notation of an algorithm is to look at the loops and to see whether the loops are nested and how n is increasing/decreasing.
- Assuming a loop body consists of only simple statements,
 - \square a single loop is O(n)
 - \square a pair of nested loops is $O(n^2)$
 - \square a nested pair of loops inside another is $O(n^3)$
 - \square and so on . . .

O(n) - Linear Growth Rate

□ If processing time increases in proportion to the number of inputs n, the algorithm grows at a linear rate which is described as O(n) or "a growth rate of order n"

```
public static int search(int[] x, int target)
{
  for(int i=0; i < x.length; i++) {    if
      (x[i]==target)
      return i;
  }
  return -1; // target not found
}</pre>
```

O(n²) - Quadratic Growth Rate

If processing time is proportional to the square of the number of inputs n, the algorithm grows at a quadratic rate

```
public static boolean areUnique(int[] x) {
  for(int i=0; i < x.length; i++) {
    for(int j=0; j < x.length; j++) {
      if (i != j && x[i] == x[j])
        return false;
    }
  }
  return true;
}</pre>
```

O(log n) Logarithms

- The inverse of raising something to a power is called a logarithm. Here logarithms are used to calculate the number of steps necessary to perform certain operations.
- □ Refer the next slide
- □ Example : If the size is 100 requires the growth rate of 6.64.
 - ☐ First find log(100) for base 10 and multiply by 3.322(helps to convert base 10 to base 2) or else directly calculate for logarithms base 2.
 - \square Log(100) = 2 (base 10); 2*3.322 = 6.644 (base 2)

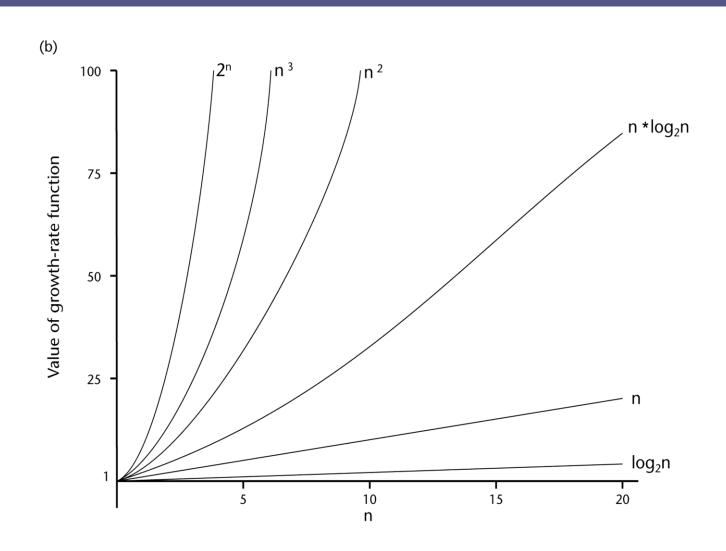
Common Growth Rates

Big-O	Name	
O(1)	Constant	
$O(\log n)$	Logarithmic	
O(n)	Linear	
$O(n \log n)$	Log-linear	
$O(n^2)$	Quadratic	
$O(n^3)$	Cubic	
$O(2^n)$	Exponential	
O(n!)	Factorial	

Effects of Different Growth Rates

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O(f(<i>n</i>))	f(50)	f(100)	f(100)/f(50)
O(1)	1	1	1
$O(\log n)$	5.64	6.64	1.18
O(n)	50	100	2
$O(n \log n)$	282	664	2.35
$O(n^2)$	2500	10,000	4
$O(n^3)$	12,500	100,000	8
$O(2^n)$	1.126×10^{15}	1.27×10^{30}	1.126×10^{15}
O(n!)	3.0×10^{64}	9.3×10^{157}	3.1×10^{93}
+7 t Wh			

Comparison of Growth Rates



What to Analyze

An algorithm can require different times to solve different problems of the same size.

Eg. Searching an item in a list of n elements using sequential search. \rightarrow Cost: 1,2,...,n

Worst-Case Analysis —The maximum amount of time that an algorithm require to solve a problem of size n.

This gives an upper bound for the time complexity of an algorithm.

Normally, we try to find worst-case behavior of an algorithm.

Best-Case Analysis —The minimum amount of time that an algorithm require to solve a problem of size n.

The best case behavior of an algorithm is NOT so useful.

Average-Case Analysis —The average amount of time that an algorithm require to solve a problem of size n.

Sometimes, it is difficult to find the average-case behavior of an algorithm.

We have to look at all possible data organizations of a given size n, and their distribution probabilities of these organizations.

Worst-case analysis is more common than average-case analysis.

Performance of KWArrayList

- The set and get methods execute in constant time:O(1)
- □ Inserting or removing general elements is linear time: O(n)
- □ Adding at the end is (usually) constant time: O(1)
 - □ With our reallocation technique the average is O(1)
 - □ The worst case is O(n) because of reallocation

Generic Collections

The statement

```
List<String> myList = new
ArrayList<String>();
```

uses a language feature called generic collections or generics

- The statement creates a List of String; only references of type String can be stored in the list
- String in this statement is called a type parameter
- The type parameter sets the data type of all objects stored in a collection

Generic Collections (cont.)

The general declaration for generic collection is

```
CollectionClassName<E> variable =
    new CollectionClassName<E>();
```

- □ The <E> indicates a type parameter
- Adding a noncompatible type to a generic collection will generate an error during compile time
- □ However, primitive types will be autoboxed:

Why Use Generic Collections?

- Type-safety: can hold only a single type of objects in generics. It doesn't allow to store other objects.
- Type casting is not required: Avoids the need to downcast from Object.
- 3. Compile-Time Checking: It is checked at compile time so problem will not occur at runtime.

^{*} More details in week 2

Day - 2

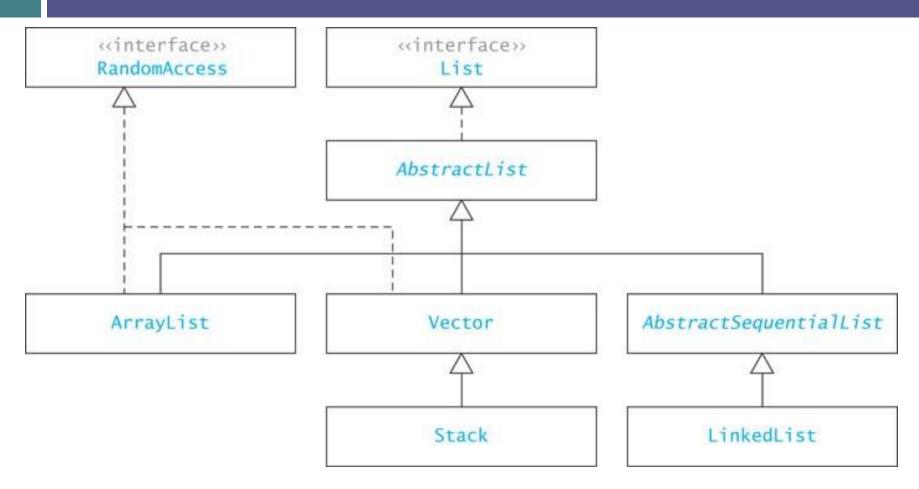
The List Interface and ArrayList Class

Section 2.2

Introduction

- A list is a collection of elements, each with a position or index
- Iterators facilitate sequential access to lists
- Classes ArrayList, Vector, and LinkedList are subclasses of abstract class AbstractList and implement the List interface

java.util.List Interface and its Implementers



List Interface and ArrayList

Class

- □ An array is an indexed structure
- In an indexed structure,
 - □ elements may be accessed in any order using subscript values
 - elements can be accessed in sequence using a loop that increments the subscript
- With the Java Array object, you cannot:
 - □ increase or decrease its length (length is fixed)
 - □ add an element at a specified position without shifting elements to make room
 - remove an element at a specified position and keep the elements contiguous without shifting elements to fill in the gap

List Interface and ArrayList

Class (cont.)

- Java provides a List interface as part of its API java.util
- Classes that implement the List interface provide the functionality of an indexed data structure and offer many more operations
- A sample of the operations:
 - □ Obtain an element at a specified position
 - □ Replace an element at a specified position
 - ☐ Find a specified target value
 - □ Add an element at either end
 - □ Remove an element from either end
 - □ Insert or remove an element at any position
 - □ Traverse the list structure without managing a subscript
- All classes introduced in this chapter support these operations, but they do not support them with the same degree of efficiency

List Interface and ArrayList Class

- Unlike the Array data structure, classes that implement the List interface cannot store primitive types
- □ Classes must store values as objects
- This requires you to wrap primitive types, such an int and double in object wrappers, in these cases, Integer and Double

ArrayList Class

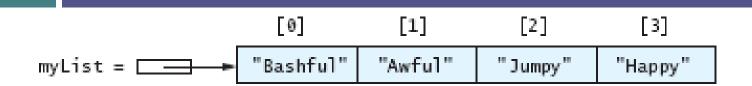
- □ The simplest class that implements the List interface
- □ An improvement over an array object
- □ Use when:
 - □ you will be adding new elements to the end of a list
 - □ you need to access elements quickly in any order

To declare a List "object" whose elements will reference String objects:

```
List<String> myList = new ArrayList<String>();
```

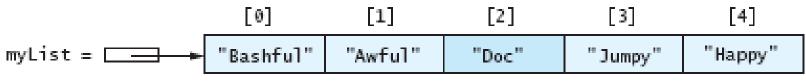
- The initial List is empty and has a default initial capacity of 10 elements
- □ To add strings to the list,

```
myList.add("Bashful");
myList.add("Awful");
myList.add("Jumpy");
myList.add("Happy");
```



Adding an element with subscript 2:

```
myList.add(2, "Doc");
```

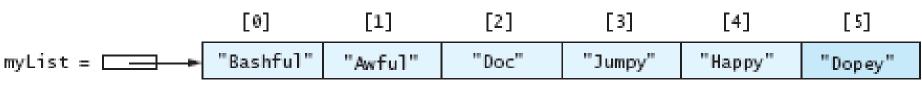


After insertion of "Doc" before the third element

Notice that the subscripts of "Jumpy" and "Happy" have changed from [2],[3] to [3],[4]

When no subscript is specified, an element is added at the end of the list:

```
myList.add("Dopey");
```

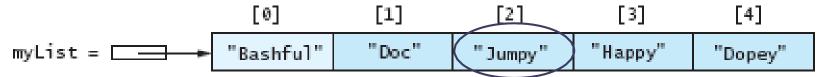


After insertion of "Dopey" at the end

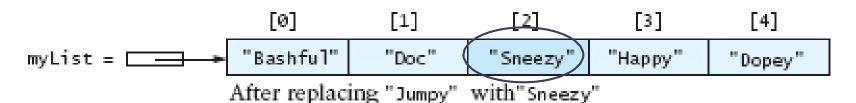
Removing an element: [2] 0 [1] [3] [4] [5]"Bashful" "Awful" "Doc" "Happy" myList = = "Jumpy" "Dopey" myList.remove(1); [4] 0 $\lceil 1 \rceil$ [2] 37 "Doc" myList = ["Bashful" "Jumpy" "Happy" "Dopey" After removal of "Awful"

The strings referenced by [2] to [5] have changed to [1] to [4]

You may also replace an element:



myList.set(2, "Sneezy");

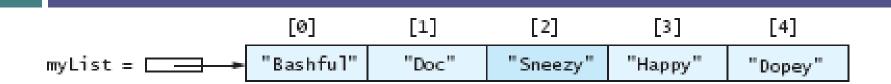




- You cannot access an element using a bracket index as you can with arrays (array[1])
- □ Instead, you must use the get() method:

```
String dwarf = myList.get(2);
```

☐ The value of dwarf becomes "Sneezy"



□ You can also search an ArrayList:

```
myList.indexOf("Sneezy");
```

☐ This returns 2 while

```
myList.indexOf("Jumpy");
```

returns -1 which indicates an unsuccessful search

Applications of ArrayList

Section 2.3

Example Application of ArrayList

```
ArrayList<Integer> someInts = new ArrayList<Integer>();
int[] nums = {5, 7, 2, 15};
for (int i = 0; i < nums.length; i++) {
  someInts.add(nums[i]);
// Display the sum
int sum = 0;
for (int i = 0; i < someInts.size(); i++) {
  sum += someInts.get(i);
System.out.println("sum is " + sum);
```

Example Application of ArrayList (cont.)

```
ArrayList<Integer> someInts = new ArrayList<Integer>();
int[] nums = {5, 7, 2, 15};
for (int i = 0; i < nums.length; i++) {
  someInts.add(nums[i]);
// Display the sum
int sum = 0;
for (int i = 0; i < someInts.size();
  sum += someInts.get(i);
System.out.println("sum is " + sum);
```

nums[i] is an int; it is automatically wrapped in an Integer object

Phone Directory Application

```
public class DirectoryEntry {
   String name;
   String number;
}
```

Create a class for objects stored in the directory

```
public class DirectoryEntry {
   String name;
   String number;
}

private ArrayList<DirectoryEntry> theDirectory =
        new ArrayList<DirectoryEntry>();
```

Create the directory

```
public class DirectoryEntry {
                                        Add a DirectoryEntry
  String name;
                                                object
  String number;
private ArrayList<DirectoryEntry> theDirectory =
          new ArrayList<DirectoryEntry>();
theDirectory.add(new DirectoryEntry("Jane Smith",
                                      "555-1212"));
```

```
Method indexOf searches
public class DirectoryEntry {
                                 theDirectory by applying the equals
  String name;
                                   method for class DirectoryEntry.
                                  Assume DirectoryEntry's equals
  String number;
                                     method compares name fields.
private ArrayList<DirectoryEntry> theDirectory =
          new ArrayList<DirectoryEntry>();
theDirectory.add(new DirectoryEntry("Jane Smith"
                                       "555-12/2"));
int index = theDirectory.indexOf(new DirectoryEntry(aName,
```

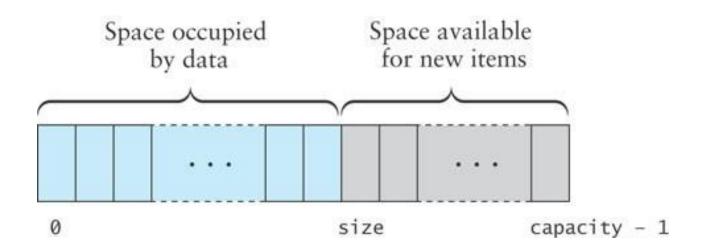
```
public class DirectoryEntry {
  String name;
  String number;
private ArrayList<DirectoryEntry> theDirectory =
          new ArrayList<DirectoryEntry>();
theDirectory.add(new DirectoryEntry("Jane Smith", "555-1212"));
int index = theDirectory.indexOf(new DirectoryEntry(aName, ""));
if (index != -1)
  dE = theDirectory.get(index);
else
  dE = null:
```

Implementation of an ArrayList Class

Section 2.4

Implementing an ArrayList Class

- KWArrayList: a simple implementation of ArrayList
 - □ Physical size of array indicated by data field capacity
 - □ Number of data items indicated by the data field size



KWArrayList Fields

```
import java.util.*;
/** This class implements some of the methods of the Java ArrayList class
* /
public class KWArrayList<E> {
  // Data fields
  /** The default initial capacity */
  private static final int INITIAL CAPACITY = 10;
  /** The underlying data array */
  private E[] theData;
  /** The current size */
  private int size = 0;
  /** The current capacity */
  private int capacity = 0;
```

KWArrayList Constructor

```
public KWArrayList () {
    capacity = INITIAL_CAPACITY;
    theData = (E[]) new Object[capacity];
}
```

This statement allocates storage for an array of type Object and then casts the array object to type E[]

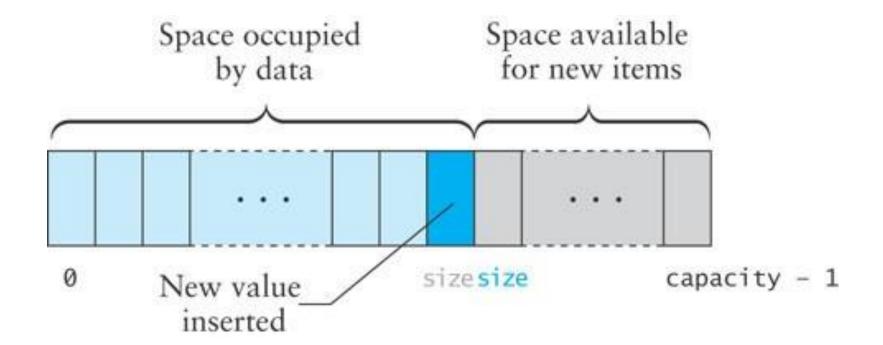
Although this may cause a compiler warning, it's ok

Implementing ArrayList.add(E)

- □ We will implement two add methods
- □ One will append at the end of the list
- □ The other will insert an item at a specified position

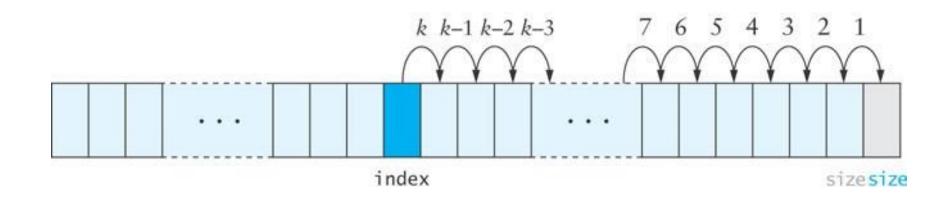
Implementing ArrayList.add(E)(cont.)

- If size is less than capacity, then to append a new item
 - insert the new item at the position indicated by the value of size
 - increment the value of size
 - 3. return true to indicate successful insertion



Implementing ArrayList.add(int index, E anEntry)

To insert into the middle of the array, the values at the insertion point are shifted over to make room, beginning at the end of the array and proceeding in the indicated order



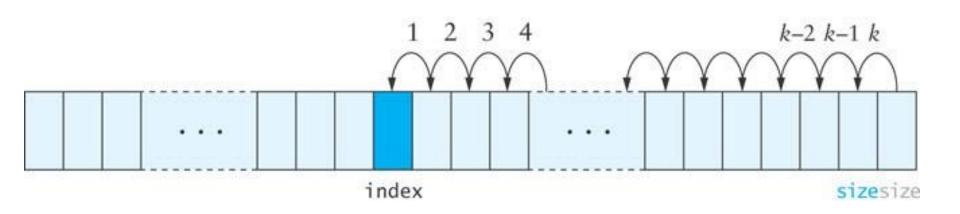
Implementing ArrayList.add(index, E)

```
public void add (int index, E anEntry) {
  // check bounds
  if (index < 0 \mid | index > size) {
    throw new ArrayIndexOutOfBoundsException(index);
  // Make sure there is room
  if (size >= capacity) {
    reallocate();
  // shift data
  for (int i = size; i > index; i--) {
    theData[i] = theData[i-1];
  // insert item
  theData[index] = anEntry;
  size++;
```

set and get Methods

```
public E get (int index) {
  if (index < 0 \mid | index >= size) {
    throw new ArrayIndexOutOfBoundsException(index);
  return theData[index];
public E set (int index, E newValue) {
  if (index < 0 \mid | index >= size) {
    throw new ArrayIndexOutOfBoundsException(index);
  E oldValue = theData[index];
  theData[index] = newValue;
  return oldValue;
```

remove Method



- When an item is removed, the items that follow it must be moved forward to close the gap
- Begin with the item closest to the removed element and proceed in the indicated order

remove Method (cont.)

```
public E remove (int index) {
  if (index < 0 \mid | index >= size) {
    throw new ArrayIndexOutOfBoundsException(index);
  E returnValue = theData[index];
  for (int i = index + 1; i < size; i++) {
    theData[i-1] = theData[i];
  size--;
  return return Value;
```

reallocate Method

 Create a new array that is twice the size of the current array and then copy the contents of the new array

```
private void reallocate () {
  capacity *= 2;
  theData = Arrays.copyOf(theData, capacity);
}
```

reallocate Method (cont.)

```
private void reallocate () {
  capacity *= 2;
  theData = Arrays.copyOf(theData, capacity);
}
```

The reason for doubling is to spread out the cost of copying;

The java.util.Arrays.copyOf(int[] original,int newLength) method copies the specified array, truncating or padding with zeros (if necessary) so the copy has the specified length.

KWArrayList as a Collection of Objects

- Earlier versions of Java did not support generics; all collections contained only Object elements
- □ To implement KWArrayList this way,
 - \square remove the parameter type <E> from the class heading,
 - □ replace each reference to data type E by Object
 - ☐ The underlying data array becomes private Object[] the Data;
- See Demo code : w112 package arraylist and
 arraylistapi (sub package)

Vector Class

- The Java API java.util contains two very similar classes, Vector and ArrayList
- New applications normally use ArrayList rather
 than Vector as ArrayList is generally more efficient
- Vector class is synchronized, which means that multiple threads can access a Vector object without conflict

Specification of the ArrayList Class

Method	Behavior
<pre>public E get(int index)</pre>	Returns a reference to the element at position index.
<pre>public E set(int index, E anEntry)</pre>	Sets the element at position index to reference anEntry. Returns the previous value.
<pre>public int size()</pre>	Gets the current size of the ArrayList.
public boolean add(E anEntry)	Adds a reference to anEntry at the end of the ArrayList. Always returns true.
<pre>public void add(int index, E anEntry)</pre>	Adds a reference to anEntry, inserting it before the item at position index.
int indexOf(E target)	Searches for target and returns the position of the first occurrence, or -1 if it is not in the ArrayList.
<pre>public E remove(int index)</pre>	Returns and removes the item at position index and shifts the items that follow it to fill the vacated space.

Day - 3 & 4

Single-Linked Lists

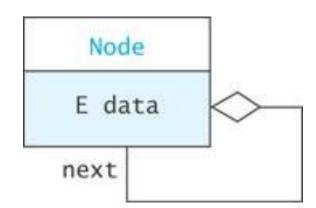
Section 2.5

Single-Linked Lists

- A linked list is useful for inserting and removing at arbitrary locations
- □ The ArrayList is limited because its add and remove methods operate in linear (O(n)) time requiring a loop to shift elements
- A linked list can add and remove elements at a known location in O(1) time
- In a linked list, instead of an index, each element is linked to the following element

A List Node

- □ A node can contain:
 - □ a data item
 - □ one or more links
- A link is a reference to a list node
- In our structure, the node contains a data field named data of type E
- and a reference to the next node,
 named next



List Nodes for Single-Linked Lists (cont.)

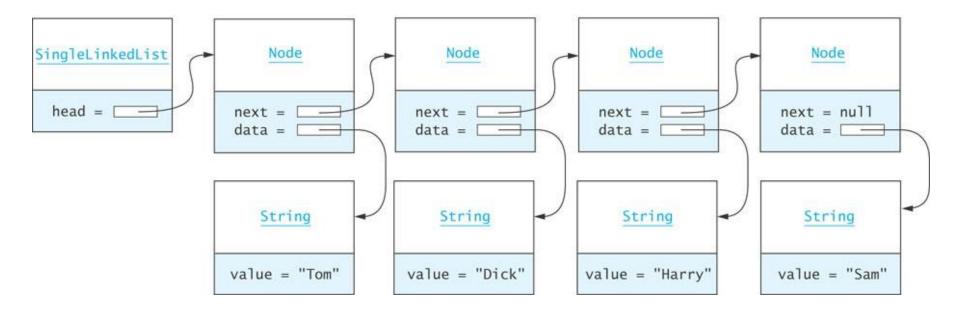
```
private static class Node<E> {
 private E data;
 private Node<E> next;
 /** Creates a new node with a null next field
      @param dataItem The data stored
  * /
  private Node(E dataItem) {
    data = dataItem:
    next = null:
 /** Creates a new node that references another node
      @param dataItem The data stored
      @param nodeRef The node referenced by new node
  * /
  private Node(E dataItem, Node<E> nodeRef) {
    data = dataItem;
    next = nodeRef;
```

The keyword static indicates that the Node<E> class will not reference its outer class

Static inner classes are also called *nested classes*

Generally, all details of the Node class should be private.
This applies also to the data fields and constructors.

Connecting Nodes



Connecting Nodes (cont.)

```
Node<String> tom = new Node<String>("Tom");
Node<String> dean = new Node<String>("Dean");
Node<String> harry = new Node<String>("Harry");
Node<String> sam = new Node<String>("Sam");

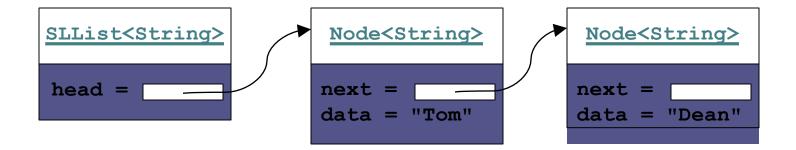
tom.next = dean;
dick.next = harry;
harry.next = sam;
```

A Single-Linked List Class

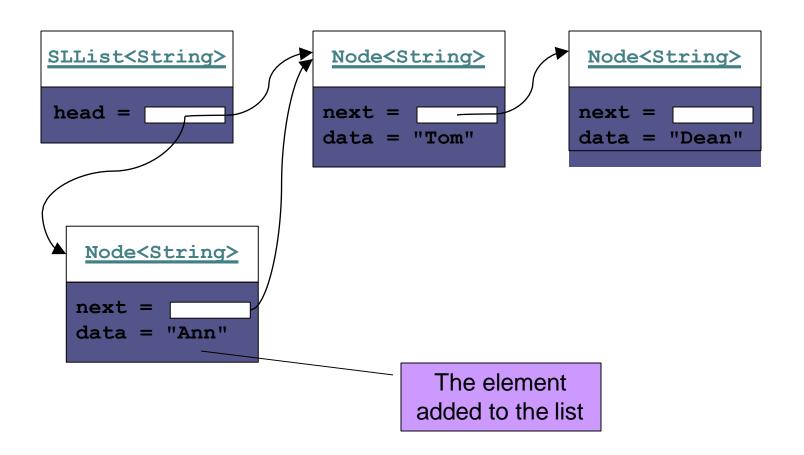
- Generally, we do not have individual references to each node.
- A SingleLinkedList object has a data field head, the list head, which references the first list node

```
public class SingleLinkedList<E> {
  private Node<E> head = null;
  private int size = 0;
  ...
}
```

SLList: An Example List



Implementing sllist.addFirst(E item)

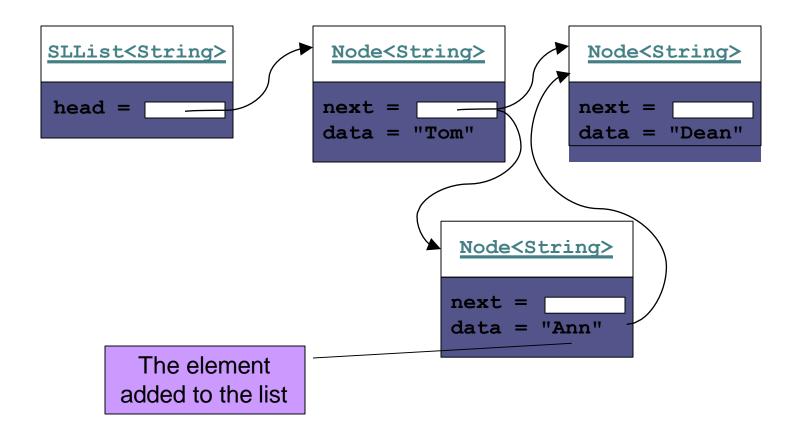


Implementing SLList.addFirst(E item) (cont.)

```
private void addFirst (E item) {
  Node<E> temp = new Node<E>(item, head);
  head = temp;
  size++;
or, more simply ...
private void addFirst (E item) {
  head = new Node < E > (item, head);
  size++;
```

This works even if head is null

Implementing addAfter(Node<E> node, E item)



item) (cont.)

```
private void addAfter (Node<E> node, E item) {
  Node<E> temp = new Node<E> (item, node.next);
  node.next = temp;
  size++;
}

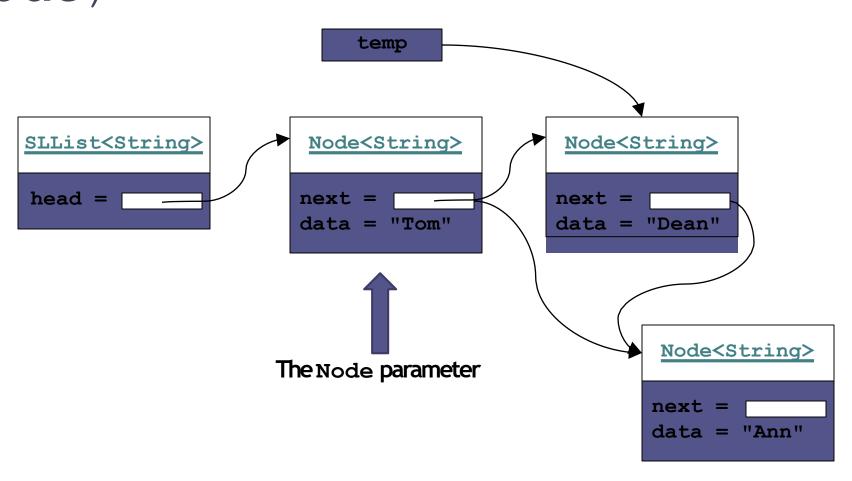
We declare this method
  since it should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class. Later we should not be contained to the class.
```

We declare this method private since it should not be called from outside the class. Later we will see how this method is used to implement the public add methods.

or, more simply.

```
private void addAfter (Node<E> node, E item) {
  node.next = new Node<E>(item, node.next);
  size++;
}
```

Implementing removeAfter (Node<E> node)

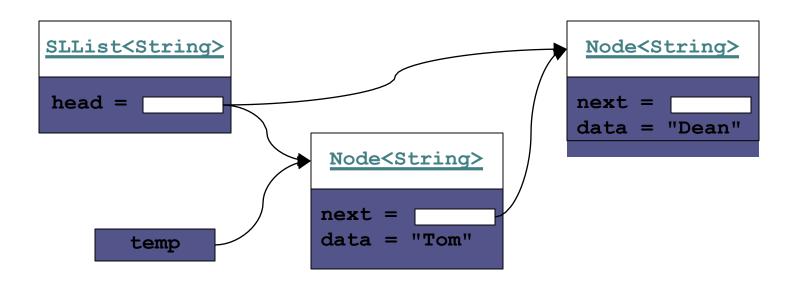


Implementing removeAfter(Node<E> node) (cont.)

```
private E removeAfter (Node<E> node) {
  Node<E> temp = node.next;
  if (temp != null) {
    node.next = temp.next;
    size--;
    return temp.data;
  } else {
    return null;
  }
}
```

Implementing

SLList.removeFirst()

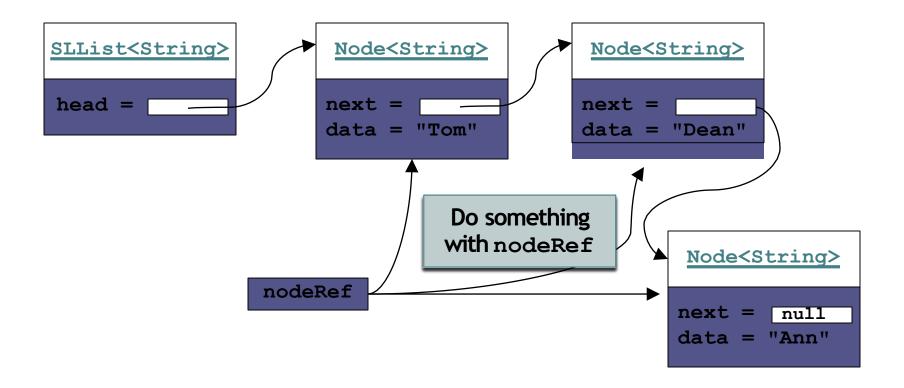


Implementing

SLList.removeFirst() (cont.)

```
private E removeFirst () {
  Node<E> temp = head;
  if (head != null) {
    head = head.next;
    size--;
    return temp.data
  else {
    return null;
```

Traversing a Single-Linked List



Traversing a Single-Linked List (cont.)

toString() can be implemented with a traversal:

```
public String toString() {
  Node<String> nodeRef = head;
  StringBuilder result = new StringBuilder();
  while (nodeRef != null) {
    result.append(nodeRef.data);
    if (nodeRef.next != null) {
      result.append(" ==> ");
    nodeRef = nodeRef.next;
  }
  return result.toString();
```

SLList.getNode(int)

In order to implement methods required by the List interface, we need an additional helper method:

```
private Node<E> getNode(int index) {
  Node<E> node = head;
  for (int i=0; i<index && node != null; i++) {
     node = node.next;
  }
  return node;
}</pre>
```

Completing the SingleLinkedList Class

Method	Behavior
<pre>public E get(int index)</pre>	Returns a reference to the element at position index.
<pre>public E set(int index, E anEntry)</pre>	Sets the element at position index to reference anEntry. Returns the previous value.
public int size()	Gets the current size of the List.
public boolean add(E anEntry)	Adds a reference to anEntry at the end of the List. Always returns true.
<pre>public void add(int index, E anEntry)</pre>	Adds a reference to anEntry, inserting it before the item at position index.
int indexOf(E target)	Searches for target and returns the position of the first occurrence, or -1 if it is not in the List.

public E get(int index)

```
public E get (int index) {
  if (index < 0 || index >= size) {
    throw new
        IndexOutOfBoundsException(Integer.toString(index));
  }
  Node<E> node = getNode(index);
  return node.data;
}
```

public E set(int index, E newValue)

```
public E set (int index, E anEntry) {
  if (index < 0 \mid | index >= size) {
    throw new
       IndexOutOfBoundsException(Integer.toString(index));
  Node<E> node = getNode(index);
  E result = node.data;
  node.data = newValue;
  return result;
```

public void add(int index, E item)

```
public void add (int index, E item) {
 if (index < 0 \mid | index > size) {
   throw new
     IndexOutOfBoundsException(Integer.toString(index));
 if (index == 0) {
   addFirst(item);
 } else {
   Node<E> node = getNode(index-1);
   addAfter(node, item);
```

public boolean add (E item)

□ To add an item to the end of the list

```
public boolean add (E item) {
  add(size, item);
  return true;
}
```

□Demo : SingleLinkedList.java

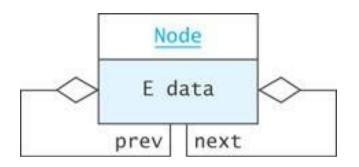
Double-Linked Lists and Circular Lists

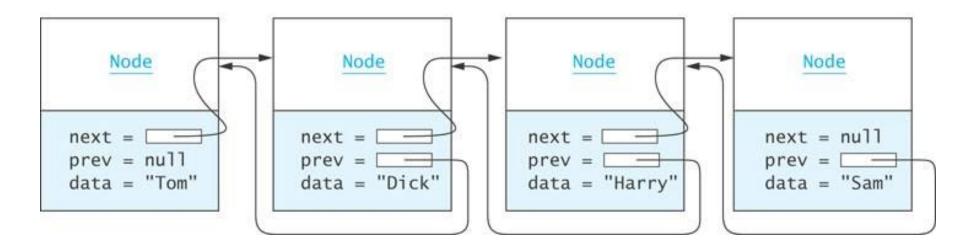
Section 2.6

Double-Linked Lists

- □ Limitations of a singly-linked list include:
 - □ Removing a node requires a reference to the previous node
 - □ We can traverse the list only in the forward direction
- □ We can overcome these limitations:
 - □ Add a reference in each node to the previous node, creating a double-linked list

Double-Linked Lists (cont.)

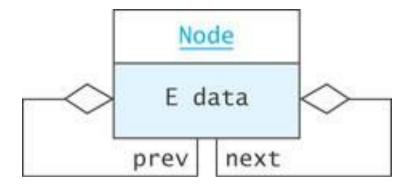




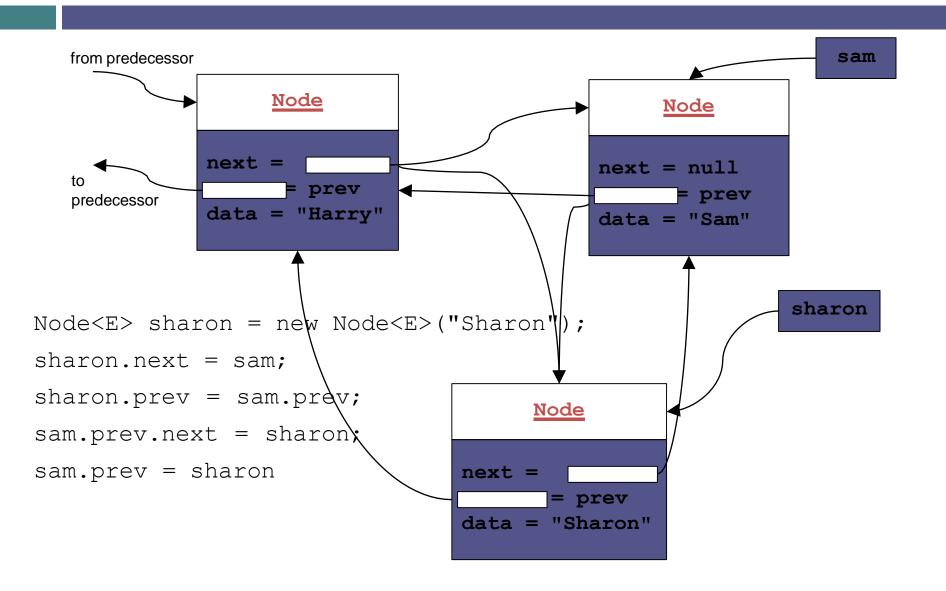
Node Class

```
private static class Node<E> {
  private E data;
  private Node<E> next = null;
  private Node<E> prev = null;

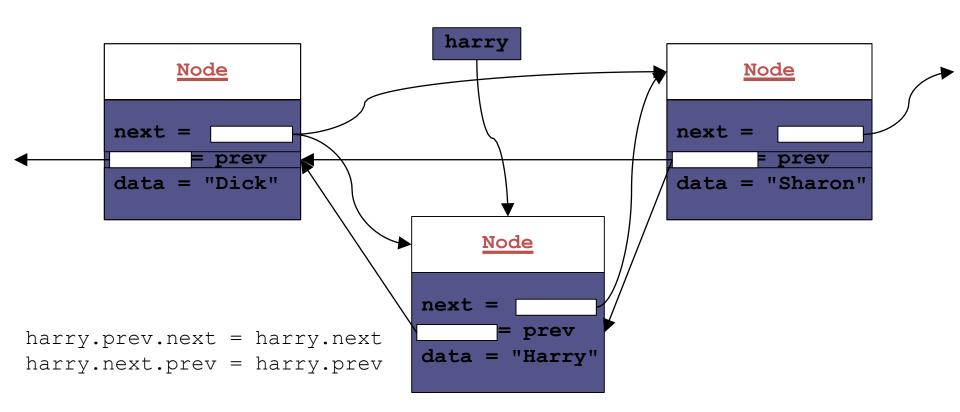
  private Node(E dataItem) {
    data = dataItem;
  }
}
```



Inserting into a Double-Linked List

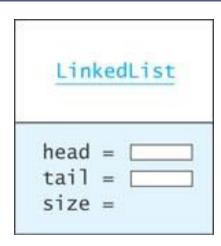


Removing from a Double-Linked List



A Double-Linked List Class

- So far we have worked only with internal nodes
- As with the single-linked class, it is best to access the internal nodes with a double-linked list object

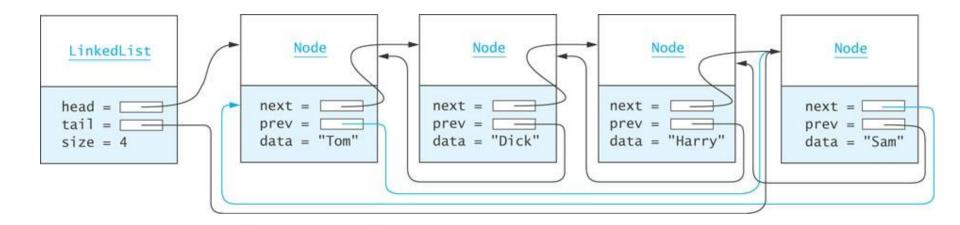


- A double-linked list object has data fields:
 - head (a reference to the first list Node)
 - tail (a reference to the last list Node)
 - □ size
- \square Insertion at either end is O(1); insertion elsewhere is still O(n)
- □ Demo: MyStringLinkedList.java, DoublyLinkedApp.java

Circular Lists

- Circular double-linked list:
 - □ Link last node to the first node, and
 - ☐ Link first node to the last node
- We can also build singly-linked circular lists:
 - □ Traverse in forward direction only
- Advantages:
 - □ Continue to traverse even after passing the first or last node
 - □ Visit all elements from any starting point
 - □ Never fall off the end of a list
- Disadvantage: Code must avoid an infinite loop!
- We are not discussing the implementation of Circular Linked
 List

Circular Lists (cont.)



The LinkedList Class and the Iterator, ListIterator, and Iterable Interfaces

Section 2.7

The LinkedList Class

Method	Behavior
<pre>public void add(int index, E obj)</pre>	Inserts object obj into the list at position index.
public void addFirst(E obj)	Inserts object obj as the first element of the list.
<pre>public void addLast(E obj)</pre>	Adds object obj to the end of the list.
<pre>public E get(int index)</pre>	Returns the item at position index.
<pre>public E getFirst()</pre>	Gets the first element in the list. Throws NoSuchElementException if the list is empty.
<pre>public E getLast()</pre>	Gets the last element in the list. Throws NoSuchElementException if the list is empty.
public boolean remove(E obj)	Removes the first occurrence of object obj from the list. Returns true if the list contained object obj; otherwise, returns false.
public int size()	Returns the number of objects contained in the list.

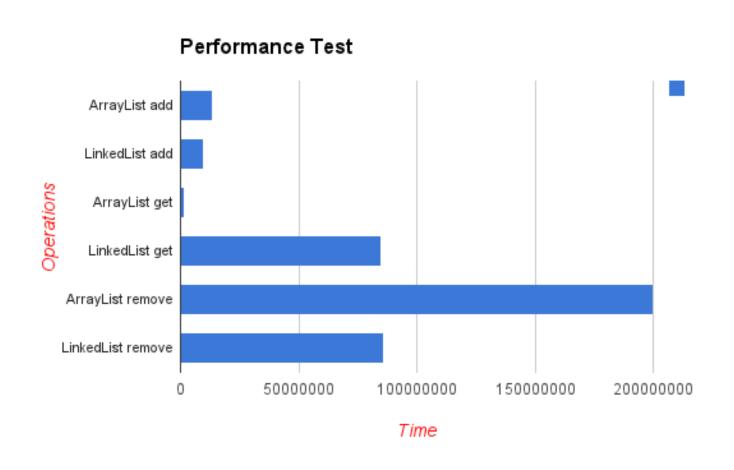
Java Collection Framework

Creation of Linked List

```
LinkedList <Integer>list = new LinkedList<Integer>();
list.add(10);
list.add(20);
int size = list.size();
```

ListMethods.doc

When to use LinkedList?



Day - 5

The Iterator

- An iterator can be viewed as a moving place marker that keeps track of the current position in a particular linked list
- An Iterator object for a list starts at the first node
- The programmer can move the Iterator by calling its next method
- The Iterator stays on its current list item until it is needed
- □ An Iterator traverses in O(n) while a list traversal using get() calls in a linked list is $O(n^2)$

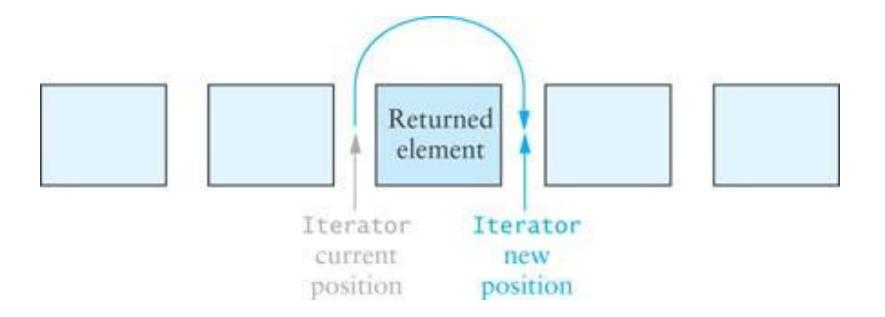
Iterator Interface

- □ The Iterator interface is defined in java.util
- The List interface declares the method iterator which returns an Iterator object that iterates over the elements of that list

Method	Behavior
boolean hasNext()	Returns true if the next method returns a value.
E next()	Returns the next element. If there are no more elements, throws the NoSuchElementException.
void remove()	Removes the last element returned by the next method.

Iterator Interface (cont.)

An Iterator is conceptually between elements; it does not refer to a particular object at any given time



Iterator Interface (cont.)

In the following loop, we process all items in List<Integer> through an Iterator

```
Iterator<Integer> iter = aList.iterator();
while (iter.hasNext()) {
  int value = iter.next();
  // Do something with value
  ...
}
```

Iterators and Removing Elements

- You can use the Iterator remove () method to remove items from a list as you access them
- □ remove() deletes the most recent element returned
- You must call next() before each remove(); otherwise, an IllegalStateException will be thrown
- □ LinkedList.remove **VS.** Iterator.remove:
 - LinkedList.remove must walk down the list each time, then remove.
 - Iterator.remove removes items without starting over at the beginning.

Iterators and Removing Elements (cont.)

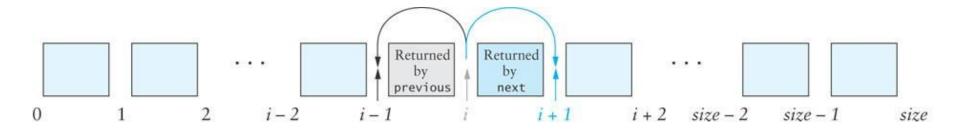
To remove all elements from a list of type Integer that are divisible by a particular value:

ListIterator Interface

- Iterator limitations
 - ☐ Traverses List only in the forward direction
 - □ Provides a remove method, but no add method
 - ☐ You must advance the Iterator using your own loop if you do not start from the beginning of the list
- ListIterator extends Iterator,
 overcoming these limitations

ListIterator Interface (cont.)

- As with Iterator, ListIterator is conceptually positioned between elements of the list
- ListIterator positions are assigned an index
 from 0 to size



ListIterator Interface (cont.)

Method	Behavior Company of the Company of t
void add(E obj)	Inserts object obj into the list just before the item that would be returned by the next call to method next and after the item that would have been returned by method previous. If method previous is called after add, the newly inserted object will be returned.
boolean hasNext()	Returns true if next will not throw an exception.
boolean hasPrevious()	Returns true if previous will not throw an exception.
E next()	Returns the next object and moves the iterator forward. If the iterator is at the end, the NoSuchElementException is thrown.
<pre>int nextIndex()</pre>	Returns the index of the item that will be returned by the next call to next. If the iterator is at the end, the list size is returned.
E previous()	Returns the previous object and moves the iterator backward. If the iterator is at the beginning of the list, the NoSuchElementExcepton is thrown.
int previousIndex()	Returns the index of the item that will be returned by the next call to previous. If the iterator is at the beginning of the list, -1 is returned.
void remove()	Removes the last item returned from a call to next or previous. If a call to remove is not preceded by a call to next or previous, the IllegalStateException is thrown.
void set(E obj)	Replaces the last item returned from a call to next or previous with obj. If a call to set is not preceded by a call to next or previous, the IllegalStateException is thrown.

ListIterator Interface (cont.)

Method		Behavior
<pre>public ListIterator<e></e></pre>	listIterator()	Returns a ListIterator that begins just before the first list element.
public ListIterator <e></e>	listIterator(int index)	Returns a ListIterator that begins just before position index.

Refer: DemoCode\w1I2\linkedlist\api

Comparison of Iterator and ListIterator

- ListIterator is a subinterface of Iterator
 - □ Classes that implement ListIterator must provide the features of both
- □ Iterator:
 - □ Requires fewer methods
 - □ Can iterate over more general data structures
- Iterator is required by the Collection interface
 - ListIterator is required only by the List
 interface

Conversion Between ListIterator and an Index

- □ ListIterator:
 - nextIndex() returns the index of item to be returned
 by next()
 - previousIndex() returns the index of item to be returned by previous()
- LinkedList has method
 listIterator(int index)
 - next() will return the item at position index

Conversion Between ListIterator and an Index (cont.)

The listIterator (int index) method creates a new ListIterator that starts at the beginning, and walks down the list to the desired position – generally an O(n) operation

Enhanced for Statement

- \square Java 5.0 introduced an enhanced for statement
- The enhanced for statement creates an Iterator object and implicitly calls its hasNext and next methods
- Other Iterator methods, such as remove, are not available

Enhanced for Statement (cont.)

 The following code counts the number of times

```
target OCCURS in myList (type
 LinkedList<String>)
count = 0;
for (String nextStr : myList) {
  if (target.equals(nextStr)) {
    count++;
```

Enhanced for Statement (cont.)

In list myList of type LinkedList<Integer>,
each

Integer object is automatically unboxed:

```
sum = 0;
for (int nextInt : myList) {
  sum += nextInt;
}
```

Enhanced for Statement (cont.)

□ The enhanced for statement also can be used with arrays, in this case, chars or type char[]

```
for (char nextCh : chars) {
   System.out.println(nextCh);
}
```

Iterable Interface

- □ Each class that implements the List interface must provide an iterator method
- □ The Collection interface extends the Iterable interface
- All classes that implement the List interface (a subinterface of Collection) must provide an iterator method
- □ Allows use of the Java 5.0 for-each loop

```
public interface Iterable<E> {
    /** returns an iterator over the elements in this
    collection. */
    Iterator<E> iterator();
}
```

Inner Classes: Static and Nonstatic

- KWLinkedList contains two inner classes:
 - Node<E> is declared static: there is no need for it to access the data fields of its parent class, KWLinkedList
 - KWListIter cannot be declared static because its methods access and modify data fields of KWLinkedList's parent object which created it
- An inner class which is not static contains an implicit reference to its parent object and can reference the fields of its parent object
- Since its parent class is already defined with the parament <E>, KWListIter cannot be declared as KWListIter<E>; if it were, an incompatible types syntax error would occur

The Collection Framework Design

Section 2.9

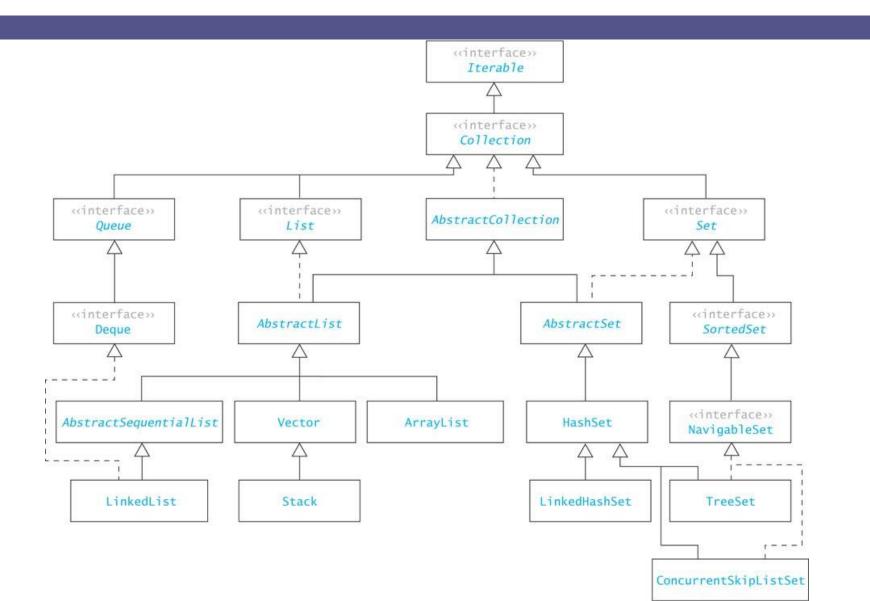
The Collection Interface

- □ Specifies a subset of methods in the List interface, specifically excluding
 - □ add(int, E)
 - □ get(int)
 - □ remove(int)
 - □ set(int, E)

but including

- □ add (E)
- □ remove(Object)
- ☐ the iterator method

The Collection Framework



Common Features of Collection

- Collection
 - □ grow as needed
 - □ hold references to objects
 - □ have at least two constructors: one to create an empty collection and one to make a copy of another collection

Common Features of Collection (cont.)

Method	Behavior
boolean add(E obj)	Ensures that the collection contains the object obj. Returns true if the collection was modified.
boolean contains(E obj)	Returns true if the collection contains the object obj.
<pre>Iterator<e> iterator()</e></pre>	Returns an Iterator to the collection.
int size()	Returns the size of the collection.

In a general
 Collection the order
 of elements is not
 specified(can't use index
 for accessing)

For collections
 implementing the List
 interface, the order of
 the elements is
 determined by the index

Common Features of Collections (cont.)

Method	Behavior
boolean add(E obj)	Ensures that the collection contains the object obj. Returns true if the collection was modified.
boolean contains(E obj)	Returns true if the collection contains the object obj.
Iterator <e> iterator()</e>	Returns an Iterator to the collection.
int size()	Returns the size of the collection.

In a general
 Collection, the
 position where an
 object is inserted is not
 specified

In ArrayList and
LinkedList, add(E)
always inserts at the
end and always returns
true

AbstractCollection, AbstractList, and AbstractSequentialList

- The Java API includes several "helper" abstract classes to help build implementations of their corresponding interfaces
- By providing implementations for interface methods not used, the helper classes require the programmer to extend the AbstractCollection class and implement only the desired methods

Implementing a Subclass of

Collection < E >

- □ Extend AbstractCollection<E>, which implements most operations
- □ You need to implement only:
 - □ add (E)
 - □ size()
 - □ iterator()
 - □ an inner class that implements Iterator<E>

Implementing a Subclass of List<E>

- □ Extend AbstractList<E>
- □ You need to implement only:
 - □ add(int, E)
 - □ get(int)
 - □ remove(int)
 - □ set(int, E)
 - □ size()
- abstractList implements Iterator<E> using
 the index

AbstractCollection, AbstractList, and AbstractSequentialList

□ Another more complete way to declare KWArrayList is:

public class KWArrayList<E> extends AbstractList<E> implements List<E>

□ Another more complete, way to declare

KWLinkedLinkedList is:

List and RandomAccess Interfaces

- □ Accessing a LinkedList using an index requires an O(n) traversal of the list until the index is located
- □ The RandomAccess interface is applied to list implementations in which indexed operations are efficient (e.g. ArrayList)
- An algorithm can test to see if a parameter of type List is also of type RandomAccess and, if not, take appropriate measures to optimize indexed operations

COMPARABLE AND COMPARATOR INTERFACE

Comparable Inteface

The Comparable interface defines the compareTo method for comparing objects.

```
package java.lang;
public interface Comparable < E > {
public int compareTo(E o);
}
```

- The **compareTo** method determines the order of this object with the specified object **o** and returns a negative integer, zero, or a positive integer if this object is less than, equal to, or greater than object **o**.
- Refer : <u>DemoCode\ w1I2\ sorting\ComparableDemo.java</u>

Comparator Interface

- Comparator can be used to compare the objects of a class that doesn't implement Comparable.
- To do so, define a class that implements the java.util.Comparator<T> interface.
- The Comparator<T> interface has two methods, compare and equals.

public int compare(T element1, T element2)

Returns a negative value if **element1** is less than **element2**, a positive value if **element1** is greater than **element2**, and zero if they are equal.

Refer: DemoCode\w1I2\ sorting\ArrayListSort.java, TestSorting.java

- If any class implements comparable interface then collection of that object can be sorted automatically using Collection.sort() or Arrays.sort().Object will be sort on the basis of compareTo method in that class.
- □ Using Comparator interface, we can write different sorting based on different attributes of objects to be sorted. You can use anonymous comparator to compare at particular line of code or other class can implement this interface to sort.
 - public void sort(Collection obj,Comparator c): is used to sort the elements of List by the given comparator.

Sorting

- To accomplish this, you specify your own ordering on a class using the Comparator interface, whose only method is compare(). Like lists, in ¡2se5.0, Comparators are parameterized.
- The compare() method is expected to behave in the following way (so it can be used in conjunction with the Collections API):

For objects a and b,

- □ compare(a,b) returns a negative number if a is "less than" b
- compare(a,b) returns a positive number if a is "greater than"
- □ compare(a,b) returns 0 if a "equals" b

Sorting (Consist with equals)

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If compare is not used in a "sensible" way, it will lead to unexpected results when used by utilities like Collections.sort.

The compare contract It must be true that:

- a is "less than" b if and only if b is "greater than" a
- \square if a is "less than" b and b is "less than" c, then a must be "less than" c.

It should also be true that the Comparator is consistent with equals; in other words:

 \Box compare(a,b) == 0 if and only if a.equals(b)

If a Comparator is not consistent with equals, problems can arise when using different container classes. For instance, the contains method of a Java List uses equals to decide if an object is in a list. However, containers that maintain the order relationship among elements (like TreeSet — more on this one later) check whether the output of compare is 0 to implement contains.

refer : DemoCode\ w1I2\ sorting\employee