Hotel Reservation System

EDS: 5.000 hotels chain, I will 200Ms. Customer pay in full.

Reservation through http/app. Cancel: 4es a Allow 10'l. over 600 king.

Room price dynamically change

NFR: Kigh Concurrency Latercy (high response)

BEE: $\frac{1 \text{ will } \times 0.7}{3} = 133,333 \approx 140.000$ Daily Reservations

240.000 = ~3 Reservations per second. Not high.

QPS + 00MS = 3 Order booking page = 30 View hotel /2001 detail = 300

HLD: API deign

Kotel / Room / Reservation related API's

Post /11/22 servations

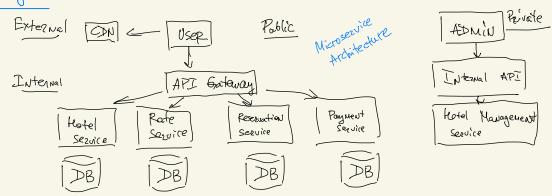
Lestart
end
Lote[Id
room Id = room Type
reservation I) = ide upotency key?

- Data Model

Relational DB (read-heavy, ACID, clear 2elational structure between entitles).

User reserves type of room: Standard, King-size, Queen lize, etc.

High Level Design:



DDD: 1. Improved data model

2. Concurrency Issues

3, Sealing

4. Data inconsistency in the Microservice architecture

1 (RDBMS 5,000 hotels x 20 types of 200MS x 2 years x 36T days = 73 mill 20WS, Enough for single DB but SPF!

HA: DB replication across multiple AZIS.

Composite Key (hotel-id, 2004-type-id, data)

- Store current and future data

- Sharding by hash (lotel-id) 1/2 wanter_of-servers

2. Avoid double booking.

· click the book btw Muttiple times

· book the same 2004 at the same line

- client-side implement. Hide button

- Idenpotent API (key as double action restriction)
idenpotency-key (= reservation_id

Serializable isolation level.

- Pessinistic Locking - lock - User2

- Optimistic Locking - version number without any timestamp lock

- DB constraints -> easy to implement DB constraint

3. DB Sharding: Lote[_id % number_of db]s

Caching: Redis -> TTL, LRU

Each shand do Asgrc Update Cache (Redis)

Data inconsistency in Microsonvices approach.

But we can use centralized DB among with Several

Services.

Two-phase commit. (2PC). DB protocol to guarantee atomic transaction commit. Plocking protocol.

- Saga. Sequence of local transactions, Provide eventual consistency.