

# **Operating System**

## **MP4: File System**

Team 73

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# Part I - Understanding NachOS file system

## 1. How does the NachOS FS manage and find free block space? Where is this information stored on the raw disk (which sector)?

FileSystem建構式中, 如果需要format, 會建立PersistentBitmap \*freeMap物件以記錄sector的使用情形, 並建立FileHeader \*mapHdr配置freeMap所需空間, 以一個bit表示一個block的是否被使用, 1024個blocks共1024個bits, 正好為一個sector的大小, FileHeader::Allocate配置時呼叫Bitmap::NumClear檢查是否有足夠的sector可用, 如果有, 在迴圈中使用Bitmap::FindAndSet取得第一個可用的sector, 將此block的位置紀錄mapHdr的dataSectors中, 配置成功後將mapHdr寫入硬碟上索引0的sector, mapHdr寫入成功後, 從硬碟開啟freeMapFile, 將freeMap的變更寫入該檔案。

最終, mapHdr存放在sector 0、dirHdr存放在sector 1, 而freeMap的data block則在使用mapHdr配置的sector 2之中。

FileSystem建構式中, 如果不需要format, 則直接從sector 0開啟freeMapFile。

## 2. What is the maximum disk size that can be handled by the current implementation? Explain why.

FileSystem建構式中建立的freeMap共有1024個sectors (32 sectors/track \* 32 tracks), 每個sector大小為128 bytes (定義在disk.h), 故模擬的硬碟大小為128 KB。

值得注意的是, disk.cc定義的DiskSize是128 KB加上額外的4 bytes, 原因是為了避免誤認其他檔案為模擬的硬碟, 然而disk.cc中的操作並不會使用這4 bytes, 故實際能儲存的空間仍是128 KB。

### 3. How does the NachOS FS manage the directory data structure? Where is this information stored on the raw disk (which sector)?

FileSystem建構式中, 如果需要format, 會建立Directory \*directory, 其中包含10個DirectoryEntry, 每個DirectoryEntry紀錄這個entry是否在使用中、在硬碟上哪個sector與其檔名, 隨後建立FileHeader \*dirHdr以配置directory所需空間, 每個DirectoryEntry大小為20 bytes, directory大小為200 bytes, 需要配置2個sectors, 配置成功後dirHdr寫入硬碟上索引1的sector, dirHdr寫入成功後, 從硬碟開啟directoryFile, 將directory的變更寫入該檔案。

最終, dirHdr存放在sector 1, 而directory的data block則在使用dirHdr配置的sector 3、4之中。

FileSystem建構式中, 如果不需要format, 則直接從sector 1開啟directoryFile。

FileSystem的公開方法都與directory管理相關, 以下是各個公開方法如何使用操作directory的細節:

- a. bool FileSystem::Create(char \*name, int initialSize)
  - 從directoryFile讀取directory物件
  - 如果該directory有同名的檔案, 創建檔案失敗
  - 從freeMapFile讀取freeMap物件
  - 嘗試從freeMap取得第一個可用的sector存放file header, 若沒有可用的sector, 則創建檔案失敗
  - 嘗試將檔名及file header的sector位置資訊寫入directory, 若沒有可用的空間 (預設一個directory只能放10個檔案), 創建檔案失敗
  - 建立FileHeader物件, 嘗試從freeMap配置檔案所需空間, 如果沒有足夠的空間, 創建檔案失敗
  - 如果上述操作都成功, 將file header、directory、freeMap寫回硬碟, 回傳檔案創建成功
- b. OpenFile \*FileSystem::Open(char \*name)
  - 從directoryFile讀取directory物件

- 尋找該檔案的header在directory下哪個sector, 若找不到, 回傳NULL, 否則使用找到的sector開啟檔案並回傳
- c. `bool FileSystem::Remove(char *name)`
- 從directoryFile讀取directory物件
  - 尋找該檔案的header在directory下哪個sector, 若找不到, 回傳檔案刪除失敗
  - 建立fileHdr及freeMap物件, 將header及data blocks歸還freeMap, 從directory移除該檔案
  - 將directory、freeMap寫回硬碟, 回傳檔案刪除成功
- d. `void FileSystem::List()`
- 從directoryFile讀取directory物件
  - 迴圈中列出directory每一個使用中的DirectoryEntry的檔名
- e. `void FileSystem::Print()`
- 從directoryFile與freeMapFile讀取directory和freeMap物件
  - 從sector 0讀取freeMap的file header並印出
  - 從sector 1讀取directory的file header並印出
  - 印出freeMap中每個bit的使用狀況
  - 迴圈中針對directory每一個使用中的DirectoryEntry, 印出其file header與檔案內容

#### 4. What information is stored in an inode? Use a figure to illustrate the disk allocation scheme of the current implementation.

FileHeader即為inode, 以下為FileHeader的私有成員, 紀錄檔案大小、檔案所使用的sector數量及data blocks使用的sectors, FileHeader的大小為128 bytes, 正好可放進一個sector, 扣除numBytes和numSectors後, 剩餘120 bytes用於存放長度為30的dataSectors, 這表示一個檔案只能有30個data blocks, 即3840 bytes。

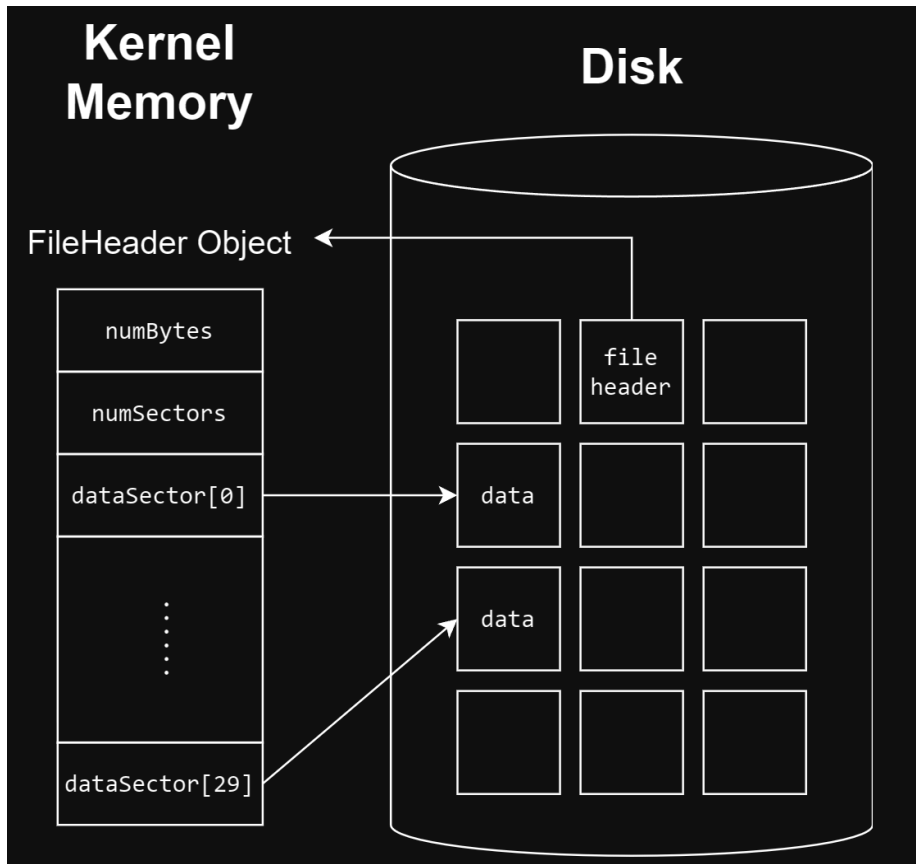
```
#define NumDirect ((SectorSize - 2 * sizeof(int)) / sizeof(int))
#define MaxFileSize (NumDirect * SectorSize)
class FileHeader
{
    // .....
private:
```

```

// Number of bytes in the file
int numBytes;
// Number of data sectors in the file
int numSectors;
// Disk sector numbers for each data block in the file
int dataSectors[NumDirect];
};

```

FileHeader使用下圖方式存取data blocks



5. What is the maximum file size that can be handled by the current implementation? Explain why.

[承上題](#), 一個檔案最大3840 bytes (3.75 KB)。

# Part II - Implementation

## 1. File I/O system calls

首先比照MP1, 在ksyscall.h加上System call介面

a. code/userprog/ksyscall.h

```
int SysCreate(char *filename, int initialSize)
{
    return kernel->fileSystem->Create(filename, initialSize);
}
OpenFileId SysOpen(char *name)
{
    return kernel->fileSystem->OpenAFile(name);
}
int SysWrite(char *buffer, int size, OpenFileId id)
{
    return kernel->fileSystem->WriteFile_(buffer, size, id);
}
int SysRead(char *buffer, int size, OpenFileId id)
{
    return kernel->fileSystem->ReadFile(buffer, size, id);
}
int SysClose(OpenFileId id)
{
    return kernel->fileSystem->CloseFile(id);
}
```

在fileSYS.h加入OpenFileTable操作方法及相關參數宣告

b. code/filesys/filesys.h

```
// .....
#define FILE_OPEN_LIMIT 20

class FileSystem
{
public:
    // .....
```



```

    OpenFileId OpenAFile(char *name);
    int WriteFile_(char *buffer, int size, OpenFileId id);
    int ReadFile(char *buffer, int size, OpenFileId id);
    int CloseFile(OpenFileId id);
private:
    OpenFile *OpenFileTable[FILE_OPEN_LIMIT];
    bool isValidFileId(OpenFileId id);
};

```

在filesys.cc加上OpenFileTable操作方法，實作基本與MP1相同，都會先檢驗OpenFileTable和OpenFileId的狀態是否合法再進行操作

c. code/filesys/filesys.cc

```

OpenFileId FileSystem::OpenAFile(char *name)
{
    OpenFileId id;
    for (id = 0; id < 20; ++id)
    {
        if (OpenFileTable[id] == NULL)
        {
            break;
        }
    }
    // exceed the opened file limit
    if (id == 20)
    {
        return -1;
    }
    OpenFileTable[id] = Open(name);
    return id;
}

int FileSystem::WriteFile_(char *buffer, int size, OpenFileId id)
{
    if (buffer != NULL && size >= 0 && isValidFileId(id))
    {
        return OpenFileTable[id]->Write(buffer, size);
    }
    return -1;
}

```

```

int FileSystem::ReadFile(char *buffer, int size, OpenFileId id)
{
    if (buffer != NULL && size >= 0 && isValidFileId(id))
    {
        return OpenFileTable[id]->Read(buffer, size);
    }
    return -1;
}

int FileSystem::CloseFile(OpenFileId id)
{
    if (isValidFileId(id))
    {
        delete OpenFileTable[id];
        OpenFileTable[id] = NULL;
        return 1;
    }
    return -1;
}

bool FileSystem::isValidFileId(OpenFileId id)
{
    return id >= 0 && id < FILE_OPEN_LIMIT && OpenFileTable[id] != NULL;
}

```

比照MP1, 加上System call的定義, 除了SC\_Create多帶了一個檔案大小的參數之外, 其他方法都與MP1的實作一模一樣, 為節省版面, 以下僅列出SC\_Create的switch case

d. code/userprog/exception.cc

```

case SC_Create:
    val = kernel->machine->ReadRegister(4);
    initialSize = kernel->machine->ReadRegister(5);
    {
        char *filename = &(kernel->machine->mainMemory[val]);
        // cout << filename << endl;
        status = SysCreate(filename, initialSize);
        kernel->machine->WriteRegister(2, (int)status);
    }
    // move PC and return...

```

## 2. File System

為了支援較大的檔案，必須先增加disk.h中Nachos模擬硬碟的容量，將SectorsPerTrack及NumTracks都從原先的32調整為1024，調整後，模擬硬碟大小從128 KB提升至128 MB。

$$1024 \text{ tracks} \times 1024 \text{ sectors/track} \times 128 \text{ bytes/sector} = 128 \text{ MB}$$

a. code/machine/disk.h

```
const int SectorsPerTrack = 1024; // number of sectors per disk track
const int NumTracks = 1024; // number of tracks per disk
```

為支援至多64 MB的大型檔案，我重新將FileHeader::dataSectors設計為30-way tree的結構，至多支援4層（檔案大小約98 MB），為增加可讀性並簡化實作，在filehdr.h定義LEVEL\_LIMIT表示此結構至多支援4層、定義MAX\_SIZE表示每層的單一檔案大小上限。

新增私有成員dataSectorMapping及children，前者能在ByteToSector方法中直接用索引算出offset位置在哪一個physical sector上，不用遍歷整個樹狀結構，後者則方便我們遍歷樹狀結構進行Allocate、Deallocate等操作，兩者都屬於in-core資料，不會寫回硬碟中，每次使用FetchFrom從硬碟讀取資料，必須重建這兩者，而為了在遍歷樹狀結構時重建dataSectorMapping，新增Allocate及FetchFrom的私有方法，幫助公開方法以遞迴方式重建dataSectorMapping。

除了上述變更外，也新增私有方法whichLv幫助計算該檔案屬於哪一階層，新增私有方法clear協助清理記憶體、回復私有成員的狀態，原有Print方法加上printContent的旗標，以滿足Bonus部分只印出file header的需求。

b. code/filesys/filehdr.h

```
#define INVALID_SECTOR -1
#define LEVEL_LIMIT 4
const int NUM_DIRECT = (SectorSize - 2 * sizeof(int)) / sizeof(int);
// 3840 bytes = 3.75 KB (30 sectors)
const int MAX_SIZE_L0 = NUM_DIRECT * SectorSize;
// 115200 bytes = 112.5 KB (900 sectors)
const int MAX_SIZE_L1 = NUM_DIRECT * NUM_DIRECT * SectorSize;
// 3456000 bytes = 3375 KB (27000 sectors)
const int MAX_SIZE_L2 = NUM_DIRECT * NUM_DIRECT * NUM_DIRECT * SectorSize;
```

```

// 103680000 bytes = 101250 KB = 98.876953125 MB (810000 sectors)
const int MAX_SIZE_L3 = NUM_DIRECT * NUM_DIRECT * NUM_DIRECT * NUM_DIRECT *
SectorSize;
const int MAX_SIZE[LEVEL_LIMIT] = {MAX_SIZE_L0, MAX_SIZE_L1, MAX_SIZE_L2,
MAX_SIZE_L3};

class FileHeader
{
public:
    // .....
    void Print(bool printContent = TRUE);

private:
    // =====disk part=====
    // Number of bytes in the file
    int numBytes;
    // Number of data sectors in the file
    int numDataSectors;
    // Disk sector numbers for each data block in the file
    int dataSectors[NUM_DIRECT];
    // =====disk part=====
    // =====in-core part=====
    // index: logical sector, value: physical sector
    vector<int> dataSectorMapping;
    FileHeader *children[NUM_DIRECT];
    // =====in-core part=====
    int whichLv(int fileSize);
    void clear();
    bool Allocate(PersistentBitmap *bitMap, int fileSize, vector<int>
&pSectors);
    void FetchFrom(int sectorNumber, vector<int> &pSectors);
};

```

FileHeader建構和解構時，都會呼叫clear()初始化成員的狀態，clear()將numBytes、numDataSectors設為-1，將dataSectors全部寫為INVALID\_SECTOR(-1)，清除dataSectorMapping，並將children中非NULL的指標釋放，此處會遞迴觸發child FileHeader的解構式，釋放整棵樹的heap空間。

whichLv方法簡單地遍歷MAX\_SIZE陣列，找到該檔案屬於哪個階層，此處假設不會有檔案超過第四層可容納的上限(約98 MB)。

ByteToSector方法直接使用dataSectorMapping轉換logical sector與physical sector，在轉換途中，也順便檢查私有成員numDataSectors和dataSectorMapping的大小是否一致，確保這兩個私有成員被正確地維護。

c. code/filesys/filehdr.cc 建構式、解構式、clear、whichLv及ByteToSector方法

```
FileHeader::FileHeader()
{
    memset(children, NULL, sizeof(children));
    clear();
}
FileHeader::~FileHeader() { clear(); }
int FileHeader::whichLv(int fileSize)
{
    for (int i = 0; i < LEVEL_LIMIT; ++i)
    {
        if (fileSize <= MAX_SIZE[i])
        {
            return i;
        }
    }
    ASSERTNOTREACHED(); // don't support file larger than MAX_SIZE_L3
}
void FileHeader::clear()
{
    numBytes = -1;
    numDataSectors = -1;
    memset(dataSectors, INVALID_SECTOR, sizeof(dataSectors));
    dataSectorMapping.resize(0);
    for (int i = 0; i < NUM_DIRECT; ++i)
    {
        if (children[i])
        {
            delete children[i];
            children[i] = NULL;
        }
        else
    }
```

```

        {
            break;
        }
    }
}

int FileHeader::ByteToSector(int offset)
{
    int logicalSector = offset / SectorSize;
    int sectorsSz = static_cast<int>(dataSectorMapping.size());
    ASSERT(logicalSector >= 0 && logicalSector < sectorsSz && sectorsSz ==
numDataSectors);
    return dataSectorMapping[logicalSector];
}

```

公開的FileHeader::Allocate方法呼叫私有的Allocate方法，並傳入自己的dataSectorMapping以蒐集子FileHeader的data sector。Allocate先判斷freeMap是否有足夠的空間，之後根據FileHeader所屬層級分為兩種案例，第一種為葉節點，和Nachos原本的實作相同，不過要將取得的data sector寫入自己的dataSectorMapping，第二種為內部節點，內部節點在迴圈中將未配置的fileSize交給子節點配置，子節點配置的檔案大小不超過該層級所能容納的最大限制，值得注意的是，內部節點所取得的sector並非用於儲存data，故不加入dataSectorMapping之中。當此節點完成配置後，將此節點配置的data sectors寫入parent的dataSectorMapping之中，如此一來，根節點的dataSectorMapping就有整個檔案的data sectors資料。

#### d. code/filesys/filehdr.cc Allocate

```

bool FileHeader::Allocate(PersistentBitmap *freeMap, int fileSize)
{
    return Allocate(freeMap, fileSize, this->dataSectorMapping);
}

bool FileHeader::Allocate(PersistentBitmap *freeMap, int fileSize, vector<int>
&pSectors)
{
    numBytes = fileSize;
    numDataSectors = divRoundUp(fileSize, SectorSize);
    if (freeMap->NumClear() < numDataSectors)
    {
        return FALSE; // not enough space
    }
}

```

```

    }
    int lv = whichLv(fileSize);
    // DEBUG(dbgMp4, "allocate lv " << lv << " file header which requires " <<
fileSize << " bytes");
    if (!lv) // direct (original Nachos implementation)
    {
        for (int i = 0; i < numDataSectors; i++)
        {
            dataSectors[i] = freeMap->FindAndSet();
            // since we checked that there was enough free space, we expect
this to succeed
            ASSERT(dataSectors[i] >= 0);
            this->dataSectorMapping.push_back(dataSectors[i]);
        }
    }
    else
    {
        int i = 0;
        while (fileSize > 0)
        {
            dataSectors[i] = freeMap->FindAndSet();
            ASSERT(dataSectors[i] >= 0);
            this->children[i] = new FileHeader();
            int subHdrSize = min(fileSize, MAX_SIZE[lv - 1]);
            this->children[i]->Allocate(freeMap, subHdrSize,
this->dataSectorMapping); // recursive
            fileSize -= subHdrSize;
            ++i;
        }
    }
    if (this->dataSectorMapping != pSectors)
    {
        pSectors.insert(pSectors.end(), this->dataSectorMapping.begin(),
this->dataSectorMapping.end());
    }
    return TRUE;
}

```

FileHeader::Deallocate和Allocate方法十分類似，都是先判斷FileHeader所屬層級，再分為葉節點和內部節點兩個案例，葉節點Deallocate的方法和Nachos原本的實作相同，內部節

點則遍歷有配置的data sectors, 此時, 因為Allocate和FetchFrom在配置或讀取資料時, 維護了children中的指標, 故只要data sector並非INVALID\_SECTOR, 就能確定children[i]一定指向子節點, 只需要遞迴呼叫子節點的Deallocate方法, 就能釋放整個樹的sectors。最後當這個節點及其下的所有節點都釋放完sectors時, 呼叫clear()清理自身的私有成員並遞迴釋放非NULL的children指標。

e. code/filesys/filehdr.cc Deallocate

```
void FileHeader::Deallocate(PersistentBitmap *freeMap)
{
    int lv = whichLv(numBytes);
    if (!lv) // direct (original Nachos implementation)
    {
        for (int i = 0; i < numDataSectors; ++i)
        {
            ASSERT(freeMap->Test((int)dataSectors[i])); // ought to be marked!
            freeMap->Clear((int)dataSectors[i]);
        }
    }
    else
    {
        for (int i = 0; i < NUM_DIRECT; ++i)
        {
            if (dataSectors[i] == INVALID_SECTOR)
            {
                break;
            }
            ASSERT(children[i]);
            children[i]->Deallocate(freeMap);
        }
    }
    clear();
}
```

公開的FileHeader::FetchFrom方法呼叫呼叫私有的FetchFrom方法, 和Allocate方法類似, 都傳入自己的dataSectorMapping以蒐集子FileHeader的data sector。私有的FetchFrom使用大小為128 bytes的buffer讀取disk part的私有成員numBytes、numDataSectors、dataSectors資料, 和Allocate相似, 接下來根據FileHeader所屬層級分為葉節點和內部節點兩個案例, 葉節點直接將dataSectors放入自身的dataSectorMapping中, 內部節點則遍歷



有效的dataSectors, 這些dataSectors為子FileHeader所在的sector, 建立子節點指標, 並遞迴呼叫子節點的FetchFrom將整棵樹的資料從硬碟讀入kernel。當這個節點及其下的所有節點都讀取完成時, 將此節點配置的dataSectors寫入parent的dataSectorMapping之中。

f. code/filesys/filehdr.cc FetchFrom

```
void FileHeader::FetchFrom(int sector)
{
    ASSERT(this->dataSectorMapping.empty());
    FetchFrom(sector, this->dataSectorMapping);
}

void FileHeader::FetchFrom(int sector, vector<int> &pSectors)
{
    int buf[SectorSize / sizeof(int)];
    kernel->synchDisk->ReadSector(sector, (char *)buf);
    numBytes = buf[0];
    numDataSectors = buf[1];
    memcpy(dataSectors, buf + 2, sizeof(dataSectors));
    // rebuild in-core part
    int lv = whichLv(numBytes);
    if (!lv) // leaf
    {
        for (int i = 0; i < numDataSectors; i++)
        {
            this->dataSectorMapping.push_back(dataSectors[i]);
        }
    }
    else
    {
        for (int i = 0; i < NUM_DIRECT; ++i)
        {
            if (dataSectors[i] == INVALID_SECTOR)
            {
                break;
            }
            this->children[i] = new FileHeader();
            this->children[i]->FetchFrom(dataSectors[i],
this->dataSectorMapping);
        }
    }
    if (this->dataSectorMapping != pSectors)
```

```

{
    pSectors.insert(pSectors.end(), this->dataSectorMapping.begin(),
this->dataSectorMapping.end());
}
}

```

FileHeader::WriteBack方法因不需要建立、維護in-core的私有成員，較FetchFrom簡單。WriteBack使用大小為128 bytes的buffer將disk part的私有成員numBytes、numDataSectors、dataSectors寫入傳入的sector中，如果這個FileHeader有子節點，則遍歷有效的dataSectors，遞迴呼叫子節點的WriteBack方法，將整棵FileHeader樹的disk part都寫入對應的sector中。

g. code/filesys/filehdr.cc WriteBack

```

void FileHeader::WriteBack(int sector)
{
    int buf[SectorSize / sizeof(int)];
    buf[0] = numBytes;
    buf[1] = numDataSectors;
    memcpy(buf + 2, dataSectors, sizeof(dataSectors));
    kernel->synchDisk->WriteSector(sector, (char *)buf);
    // write disk part
    int lv = whichLv(numBytes);
    if (lv)
    {
        for (int i = 0; i < NUM_DIRECT; ++i)
        {
            if (dataSectors[i] == INVALID_SECTOR)
            {
                break;
            }
            ASSERT(children[i]);
            children[i]->WriteBack(dataSectors[i]);
        }
    }
}

```

FileHeader::Print印出該FileHeader的disk part和in-core part的大小，如果需要印出檔案內容，則使用dataSectorMapping印出該節點及其下的所有data sectors，接下來根據

FileHeader的層級分為葉節點和內部節點兩個案例，葉節點的部分和Nachos的原始實作相同，但會先判斷是否需要印出檔案內容，內部節點則遍歷有效的dataSectors，遞迴呼叫子節點的Print方法。

h. code/filesys/filehdr.cc Print

```
void FileHeader::Print(bool printContent)
{
    cout << "FileHeader contents:" << endl
        << "1. File size: " << numBytes << " bytes (" << numDataSectors << "
sectors)" << endl
        << "2. FileHeader size in disk: " << (sizeof(FileHeader) -
sizeof(dataSectorMapping) - sizeof(children)) << " bytes" << endl
        << "3. FileHeader size in memory: " << (sizeof(FileHeader) +
dataSectorMapping.size() * sizeof(int) + sizeof(children)) << " bytes" <<
endl;
    if (printContent)
    {
        cout << "4. Data blocks: " << endl;
        for (int i = 0; i < numDataSectors; ++i)
        {
            ASSERT(dataSectorMapping[i] != INVALID_SECTOR);
            cout << dataSectorMapping[i] << " ";
        }
        printf("\nFile contents:\n");
    }

    int lv = whichLv(numBytes);
    if (!lv) // leaf
    {
        if (!printContent)
        {
            return;
        }
        char *data = new char[SectorSize];
        for (int i = 0, k = 0; i < numDataSectors; ++i)
        {
            kernel->synchDisk->ReadSector(dataSectors[i], data);
            for (int j = 0; (j < SectorSize) && (k < numBytes); j++, k++)
            {
                if ('\040' <= data[j] && data[j] <= '\176') //
```

```

isprint(data[j])
    {
        printf("%c", data[j]);
    }
    else
    {
        printf("\\%x", (unsigned char)data[j]);
    }
}
printf("\n");
}
delete[] data;
}
else
{
    for (int i = 0; i < NUM_DIRECT; ++i)
    {
        if (dataSectors[i] == INVALID_SECTOR)
        {
            break;
        }
        ASSERT(children[i]);
        children[i]->Print(printContent);
    }
}
}
}

```

將NumDirEntries調整為64，使一個目錄能容納64個檔案和目錄。為支援子目錄，DirectoryEntry新增isDir屬性，以區分該DirectoryEntry為檔案還是目錄，同時調整Directory原先Find、FindIndex、Add、Remove的方法宣告，加入isDir參數以區分要操作、尋找的DirectoryEntry。Directory加上RemoveAll及RecursivelyList的宣告以支援-rr和-lr指令。

#### i. code/filesys/directory.h

```

#define NumDirEntries 64
#define FILE false
#define DIR true

class DirectoryEntry
{
public:

```

```

    // Is this a dir entry?
    bool isDir;
    // Is this directory entry in use?
    bool inUse;
    // Location on disk to find the FileHeader for this file
    int sector;
    // Text name for file, with +1 for the trailing '\0'
    char name[FileNameMaxLen + 1];
};

class Directory
{
public:
    // .....
    // Find the sector number of the FileHeader for file: "name"
    int Find(const char *name, bool isDir);
    // Add a file name into the directory
    bool Add(const char *name, int newSector, bool isDir);
    // Remove a file from the directory
    bool Remove(const char *name, bool isDir);
    // Remove all files and directories under this directory
    bool RemoveAll(PersistentBitmap *freeMap);
    // Print the names of all the files in the directory (command -l)
    void List();
    // command -lr
    void RecursivelyList(int depth);
private:
    //---disk part---//
    // Table of pairs: <file name, file header location>
    DirectoryEntry *table;
    //---disk part---//
    //---in-core part---//
    // Number of directory entries
    int tableSize;
    //---in-core part---//
    int FindIndex(const char *name, bool isDir);
};

```

Directory建構式在初始化迴圈中設定每個table entry的isDir屬性。

j. code/filesys/directory.cc 建構式

```
Directory::Directory(int size)
{
    table = new DirectoryEntry[size];
    memset(table, 0, sizeof(DirectoryEntry) * size); // dummy operation to
keep valgrind happy
    tableSize = size;
    for (int i = 0; i < tableSize; i++)
    {
        table[i].isDir = FALSE;
        table[i].inUse = FALSE;
    }
}
```

Directory::FindIndex方法增加判斷isDir屬性是否和要找的目標相容, Find、Add和Remove則是將外部傳入的isDir參數直接傳給FindIndex, 其中Add要記得設定新table entry的isDir屬性, 其餘都和原實作相同

k. code/filesys/directory.cc FindIndex、Find、Add及Remove

```
int Directory::FindIndex(const char *name, bool isDir)
{
    for (int i = 0; i < tableSize; i++)
    {
        if (table[i].inUse && table[i].isDir == isDir &&
!strncmp(table[i].name, name, FileNameMaxLen))
        {
            return i;
        }
    }
    return -1; // name not in directory
}

int Directory::Find(const char *name, bool isDir)
{
    int i = FindIndex(name, isDir);
    if (i != -1)
    {
        return table[i].sector;
    }
}
```

```

        return INVALID_SECTOR;
    }

bool Directory::Add(const char *name, int newSector, bool isDir)
{
    if (FindIndex(name, isDir) != -1)
    {
        return FALSE;
    }

    for (int i = 0; i < tableSize; i++)
    {
        if (!table[i].inUse)
        {
            table[i].isDir = isDir;
            table[i].inUse = TRUE;
            strncpy(table[i].name, name, FileNameMaxLen);
            table[i].sector = newSector;
            return TRUE;
        }
    }
    return FALSE; // no space.  Fix when we have extensible files.
}

bool Directory::Remove(const char *name, bool isDir)
{
    int i = FindIndex(name, isDir);
    if (i == -1)
    {
        return FALSE; // name not in directory
    }
    table[i].inUse = FALSE;
    DEBUG(dbgMp4, "remove " << name << " (dir)");
    return TRUE;
}

```

Directory::RemoveAll遍歷每個inUse的table entry, 如果該entry是子目錄, 則使用entry中紀錄的sector開啟子目錄, 並遞迴呼叫子目錄的RemoveAll, 清除子目錄下的所有檔案和目錄(不含清除子目錄), 完成後, 或者該entry為檔案, 將該entry的FileHeader及data占用的

sectors歸還freeMap。因為Directory物件並未記錄自身的FileHeader sector資訊，此方法的呼叫者必須負責將此Directory的變更寫回硬碟中。

#### l. code/filesys/directory.cc RemoveAll

```
bool Directory::RemoveAll(PersistentBitmap *freeMap)
{
    for (int i = 0; i < tableSize; i++)
    {
        if (!table[i].inUse)
        {
            continue;
        }
        if (table[i].isDir)
        {
            OpenFile *openFh = new OpenFile(table[i].sector);
            Directory *dir = new Directory(NumDirEntries);
            dir->FetchFrom(openFh);
            dir->RemoveAll(freeMap);
            delete openFh;
            delete dir;
        }
        DEBUG(dbgMp4, "remove " << table[i].name << " (dir or file)");
        ASSERT(freeMap->Test(table[i].sector));
        freeMap->Clear(table[i].sector); // return header sector
        FileHeader *fh = new FileHeader();
        fh->FetchFrom(table[i].sector);
        fh->Deallocate(freeMap); // return data sectors
        delete fh;

        table[i].inUse = FALSE;
    }
    return TRUE;
}
```

Directory::RecursivelyList遍歷每個inUse的table entry，根據深度階層印出空格縮排後，印出自身為檔案或目錄以及自身的名字，如果該table entry為子目錄，開啟子目錄並遞迴呼叫子目錄的RecursivelyList方法，傳入的深度階層參數往上+1，以正確印出縮排。

#### m. code/filesys/directory.cc RecursivelyList

```
void Directory::RecursivelyList(int depth)
```



```

{
    for (int i = 0; i < tableSize; i++)
    {
        if (!table[i].inUse)
        {
            continue;
        }
        // indentation
        for (int j = 0; j < depth * 4; ++j)
        {
            cout << " ";
        }
        cout << "[" << (table[i].isDir ? "D" : "F") << "]" " " << table[i].name
<< endl;
        if (table[i].isDir)
        {
            OpenFile *f = new OpenFile(table[i].sector);
            Directory *dir = new Directory(NumDirEntries);
            dir->FetchFrom(f);
            delete f;
            dir->RecursivelyList(depth + 1);
            delete dir;
        }
    }
}

```

為了使尋找檔案和目錄的操作變得輕鬆，在filesys.h新增class FileFinder，負責解析路徑、尋找檔案，尋找途中會將parent FileHeader sector寫入pFhSector屬性、將自身FileHeader sector寫入fhSector屬性、將檔名（路徑字串中，最後一個 '/' 後的部分）寫入filename屬性，並將搜尋結果寫入exist屬性。

n. code/filesys/filesys.h FileFinder

```

#define PATH_NAME_MAX_LEN 256
#define FILE_OPEN_LIMIT 20

class FileFinder
{
    friend class FileSystem;
    // is the file exist?
    bool exist;

```

```

    // sector number of the parent dir
    int pFhSector;
    // sector number of the file header
    int fhSector;
    string filename;

    FileFinder();
    static vector<string> splitPath(const char *name, bool
&exceedPathLenLimit);
    // find the file by path (name) and set the result to private fields
    void find(const char *name, bool isDir, OpenFile *root);
    void deleteOpenFile(OpenFile *openfile, OpenFile *root);
};

```

FileFinder建構式初始化表示搜尋結果的屬性，避免使用時誤解搜尋結果。splitPath將傳入的路徑字串依'/'符號進行分割，並回傳結果，分割途中順便檢查路徑字串是否超過PATH\_NAME\_MAX\_LEN(256)的限制。deleteOpenFile則是為了簡化FileFinder::find中重複的刪除操作而存在的方法。

#### o. code/filesys/filesys.cc FileFinder建構式、splitPath及deleteOpenFile

```

FileFinder::FileFinder() : exist(FALSE), pFhSector(INVALID_SECTOR),
fhSector(INVALID_SECTOR) {}

vector<string> FileFinder::splitPath(const char *str, bool
&exceedPathLenLimit)
{
    vector<string> result;
    string token;
    int i = 0;
    while (*str)
    {
        if (*str == '/')
        {
            result.push_back(token);
            token.clear();
        }
        else
        {
            token += *str;
        }
    }
}

```

```

        ++str;
        ++i;
    }
    exceedPathLenLimit = i >= PATH_NAME_MAX_LEN;
    result.push_back(token);
    return result;
}

void FileFinder::deleteOpenFile(OpenFile *openfile, OpenFile *root)
{
    if (openfile != NULL && openfile != root)
    {
        delete openfile;
    }
}

```

FileFinder::find方法先調用splitPath處理路徑字串，並檢查路徑長度是否超過上限，如果路徑字串為根目錄，將fhSector設定為DirectorySector、exist設定為true後返回表示尋找成功，否則遍歷分割的路徑字串，從根目錄開始尋找檔案，尋找的途中維護pFhSector及fhSector，除了最後一層尋找的目標型態由傳入的isDir參數決定，途中都是尋找目錄，如果中途找不到對應的目錄，或最後一層不存在目標型態、名稱相同的檔案或目錄，清除指標後返回，表示尋找失敗，否則遍歷到最後，表示找到目標，設定exist為true後返回。

p. code/filesys/filesys.cc FileFinder find

```

void FileFinder::find(const char *name, bool isDir, OpenFile *root)
{
    bool exceedPathLenLimit = FALSE;
    vector<string> path = splitPath(name, exceedPathLenLimit);
    filename = path.back();
    if (exceedPathLenLimit)
    {
        return;
    }
    // the file is the root dir
    if (!strcmp(name, "/"))
    {
        ASSERT(isDir);
        exist = true;
        fhSector = DirectorySector;
    }
}

```

```

        return;
    }
    OpenFile *openPfh = root; // open parent file header
    Directory *pDir = new Directory(NumDirEntries);
    int pathSz = static_cast<int>(path.size());
    for (int i = 1; i < pathSz; ++i)
    {
        pDir->FetchFrom(openPfh);
        // set pfh sector to previous fh sector
        if (i == 1)
        {
            pFhSector = DirectorySector;
        }
        else
        {
            pFhSector = fhSector;
        }
        // find dir until the leaf
        fhSector = pDir->Find(path[i].c_str(), (i == pathSz - 1 ? isDir :
DIR));
        if (fhSector == INVALID_SECTOR)
        {
            deleteOpenFile(openPfh, root);
            delete pDir;
            return;
        }
        deleteOpenFile(openPfh, root);
        openPfh = new OpenFile(fhSector);
    }
    exist = true;
    deleteOpenFile(openPfh, root);
    delete pDir;
}

```

FileSystem除了新增System call的宣告外，既有的Remove及List方法加上recursive選項，以支援-rs和-lr指令，新增PrintHeader公開方法對應Bonus印出不同大小的FileHeader指令（新的自訂指令-h），新增Mkdir公開方法以支援新增目錄的指令(-mkdir)，新增私有方法createFileOrDir以整合Create和Mkdir兩個公開方法的共通邏輯，新增私有方法

recursivelyRemove以實作-r指令，新增私有方法returnSectorsToFreeMap以簡化將sectors歸還freeMap的操作。

q. code/filesys/filesys.h FileSystem

```
class FileSystem
{
    friend class FileFinder;

public:
    // .....
    // Create a file (UNIX creat)
    bool Create(char *name, int initialSize);
    // Open a file (UNIX open)
    OpenFile *Open(char *name);
    // This function is used for kernel open system call
    OpenFileId OpenAFile(char *name);
    int WriteFile_(char *buffer, int size, OpenFileId id);
    int ReadFile(char *buffer, int size, OpenFileId id);
    int CloseFile(OpenFileId id);
    // Delete a file (UNIX unlink)
    bool Remove(const char *name, bool recursive);
    // List all the files in the file system
    void List(char *name, bool recursive);
    void PrintHeader(char *name);
    bool Mkdir(char *name);
    // .....

private:
    // .....
    OpenFile *OpenFileTable[FILE_OPEN_LIMIT];
    bool isValidFileId(OpenFileId id);
    /**
     * @brief Create a File Or Dir
     *
     * @param name absolute path
     * @param isDir is this a dir or a file
     * @param initialSize file size (will be ignored if this is a dir)
     * @return true success
     * @return false fail
     */
    bool createFileOrDir(char *name, bool isDir, int initialSize);
```

```

bool recursivelyRemove(const char *name);
// Return data and header sectors to freeMap
void returnSectorsToFreeMap(int fhSector, PersistentBitmap *freeMap);
};

```

Create和Mkdir兩者都是直接呼叫私有方法createFileOrDir完成實作，其中Create對應的是檔案的新增，Mkdir是目錄的新增。createFileOrDir的實作可分為4個部分，第一，根據路徑和屬性尋找檔案，如果檔案已存在，返回無法新增檔案，否則parent目錄必定存在，進入第二部分，將檔名寫入parent目錄的其中一個entry，我們假設測試案例總是有足夠的空間可以寫入，第三部分，配置data sectors，如果要新增的是目錄，則大小為NumDirEntries \* sizeof(DirectoryEntry)，否則為傳入的initialSize，建立FileHeader配置data sectors後，第四部份，將FileHeader、parent目錄、freeMap的變更寫回硬碟，如果要新增的是目錄，也將目錄的table寫回硬碟，清除指標後返回新增成功。

r. code/filesys/filesys.cc Create、Mkdir及createFileOrDir

```

bool FileSystem::Create(char *name, int initialSize)
{
    return createFileOrDir(name, FILE, initialSize);
}
bool FileSystem::Mkdir(char *name)
{
    return createFileOrDir(name, DIR, -1);
}
bool FileSystem::createFileOrDir(char *name, bool isDir, int initialSize)
{
    // 1. find the parent dir
    FileFinder finder = FileFinder();
    finder.find(name, isDir, directoryFile);
    if (finder.exist)
    {
        DEBUG(dbgMp4, "File or dir exist, cannot create: " << name);
        return FALSE;
    }
    ASSERT(finder.pFhSector != INVALID_SECTOR); // parent dir should exist

    // 2. add file header to the parent dir
    PersistentBitmap *freeMap = new PersistentBitmap(freeMapFile, NumSectors);
    int sector = freeMap->FindAndSet();
}

```

```

ASSERT(sector >= 0);
OpenFile *openPfh = new OpenFile(finder.pFhSector);
Directory *pDir = new Directory(NumDirEntries);
pDir->FetchFrom(openPfh);

ASSERT(pDir->Add(finder.filename.c_str(), sector, isDir)) // assume always
have space

// 3. allocate data blocks
int size = isDir ? NumDirEntries * sizeof(DirectoryEntry) : initialSize;
ASSERT(size >= 0);
FileHeader *fh = new FileHeader();
ASSERT(fh->Allocate(freeMap, size));

// 4. write back
fh->WriteBack(sector);
pDir->WriteBack(openPfh);
freeMap->WriteBack(freeMapFile);
if (isDir)
{
    Directory *dir = new Directory(NumDirEntries);
    OpenFile *openFh = new OpenFile(sector);
    dir->WriteBack(openFh);
    delete dir;
    delete openFh;
}
delete freeMap;
delete openPfh;
delete pDir;
delete fh;
return TRUE;
}

```

FileSystem::Open使用FileFinder嘗試在傳入的路徑下尋找檔案或目錄，測試案例假設要開啟的檔案必定存在，故finder.exist必定為true，且FileHeader sector必定為非負整數，回傳開啟的檔案指標。

s. code/filesys/filesys.cc Open

```

OpenFile *FileSystem::Open(char *name)
{

```

```

FileFinder finder = FileFinder();
finder.find(name, FILE, directoryFile);
if (!finder.exist)
{
    finder.find(name, DIR, directoryFile);
}
ASSERT(finder.exist && finder.fhSector >= 0);
return new OpenFile(finder.fhSector);
}

```

FileSystem::Remove增加遞迴刪除的選項，如果要遞迴刪除，直接交給recursivelyRemove方法，否則為單一檔案的刪除，使用FileFinder尋找檔案，不存在則回傳刪除失敗，否則將該檔案data和header占用的sectors歸還freeMap，開啟parent目錄，將檔名從parent目錄中移除，最後將freeMap和parent目錄的變更寫回硬碟。

FileSystem::recursivelyRemove先使用FileFinder尋找目錄，如果找不到，表示為檔案，呼叫Remove進行單檔刪除，否則呼叫Directory::RemoveAll方法，將該目錄下所有檔案和目錄移除（不包含該目錄），成功之後將該目錄從其parent目錄刪除（如果為根目錄則不需要進行這一步），並將freeMap和目錄的變更寫回硬碟。

t. code/filesys/filesys.cc Remove和recursivelyRemove

```

bool FileSystem::Remove(const char *name, bool recursive)
{
    if (recursive)
    {
        return recursivelyRemove(name);
    }
    FileFinder finder = FileFinder();
    finder.find(name, FILE, directoryFile);
    if (!finder.exist)
    {
        return FALSE;
    }
    PersistentBitmap *freeMap = new PersistentBitmap(freeMapFile, NumSectors);
    return SectorsToFreeMap(finder.fhSector, freeMap);
    OpenFile *openPfh = new OpenFile(finder.pFhSector);
    Directory *pDir = new Directory(NumDirEntries);
    pDir->FetchFrom(openPfh);
}

```



```

    ASSERT(pDir->Remove(finder.filename.c_str(), FILE));

    freeMap->WriteBack(freeMapFile);
    pDir->WriteBack(openPfh);
    DEBUG(dbgMp4, "remove " << name << " (single file)");
    delete freeMap;
    delete openPfh;
    delete pDir;
    return TRUE;
}

bool FileSystem::recursivelyRemove(const char *name)
{
    FileFinder finder = FileFinder();
    finder.find(name, DIR, directoryFile);
    if (!finder.exist)
    {
        return Remove(name, FALSE); // remove single file
    };
    // remove all files/dirs in this dir
    PersistentBitmap *freeMap = new PersistentBitmap(freeMapFile, NumSectors);
    OpenFile *openfh = new OpenFile(finder.fhSector);
    Directory *dir = new Directory(NumDirEntries);
    dir->FetchFrom(openfh);
    ASSERT(dir->RemoveAll(freeMap));
    dir->WriteBack(openfh);
    freeMap->WriteBack(freeMapFile);
    // remove the dir
    if (finder.pFhSector != INVALID_SECTOR)
    {
        returnSectorsToFreeMap(finder.fhSector, freeMap);
        OpenFile *openPfh = new OpenFile(finder.pFhSector);
        Directory *pDir = new Directory(NumDirEntries);
        pDir->FetchFrom(openPfh);
        pDir->Remove(finder.filename.c_str(), DIR);
        pDir->WriteBack(openPfh);
        delete openPfh;
        delete pDir;
    }
    delete freeMap;
    delete openfh;
    delete dir;
}

```

```
    return TRUE;
}
```

FileSystem::List使用FileFinder尋找目錄，如果不存在則直接回傳，否則開啟目錄，並根據傳入的recursive選項決定呼叫Directory::RecursivelyList或Directory::List。

u. code/filesys/filesys.cc List

```
void FileSystem::List(char *name, bool recursive)
{
    FileFinder finder = FileFinder();
    finder.find(name, DIR, directoryFile);
    // file not exist
    if (!finder.exist)
    {
        return;
    }
    ASSERT(finder.fhSector != INVALID_SECTOR);
    OpenFile *f = new OpenFile(finder.fhSector);
    Directory *dir = new Directory(NumDirEntries);
    dir->FetchFrom(f);
    delete f;
    if (recursive)
    {
        dir->RecursivelyList(0);
    }
    else
    {
        dir->List();
    }
    delete dir;
}
```

FileSystem::PrintHeader對應新增的自訂指令-h，找到對應的檔案或目錄後，呼叫FileHeader::Print並傳入FALSE，表示只需印出header，不需印出檔案內容。

v. code/filesys/filesys.cc PrintHeader

```
void FileSystem::PrintHeader(char *name)
{
    FileFinder finder = FileFinder();
```

```

finder.find(name, DIR, directoryFile);
if (!finder.exist)
{
    finder.find(name, FILE, directoryFile);
}
ASSERT(finder.fhSector != INVALID_SECTOR);
FileHeader *hdr = new FileHeader();
hdr->FetchFrom(finder.fhSector);
hdr->Print(FALSE);
delete hdr;
}

```

FileSystem::returnSectorsToFreeMap簡化將header和data sectors歸還freeMap的操作，呼叫者必須自行將freeMap的變更寫回硬碟。

w. code/filesys/filesys.cc returnSectorsToFreeMap

```

void FileSystem::returnSectorsToFreeMap(int fhSector, PersistentBitmap
*freeMap)
{
    ASSERT(freeMap->Test(fhSector));
    freeMap->Clear(fhSector); // return header sector
    FileHeader *fh = new FileHeader();
    fh->FetchFrom(fhSector);
    fh->Deallocate(freeMap); // return data sectors
    delete fh;
}

```

main方法新增-h指令以支援Bonus僅印出FileHeader的部分，對應-rr和-lr指令，將原本的Remove和List方法加上recursive選項，新增目錄則改為直接呼叫FileSystem::Mkdir方法。

x. code/threads/main.cc

```

// .....
int main(int argc, char **argv)
{
    // .....
    char *printHeaderName = NULL;
    bool mkdirFlag = false;
    bool recursiveListFlag = false;
    bool recursiveRemoveFlag = false;

```

```

bool printHeaderFlag = false;
for (i = 1; i < argc; i++)
{
    // .....
    else if (strcmp(argv[i], "-h") == 0)
    {
        ASSERT(i + 1 < argc);
        printHeaderFlag = true;
        printHeaderName = argv[i + 1];
        i++;
    }
    // .....
}

if (removeFileName != NULL)
{
    kernel->fileSystem->Remove(removeFileName, recursiveRemoveFlag);
}
// .....
if (dirListFlag)
{
    kernel->fileSystem->List(listDirectoryName, recursiveListFlag);
}
if (mkdirFlag)
{
    kernel->fileSystem->Mkdir(createDirectoryName);
}
// .....
if (printHeaderFlag)
{
    kernel->fileSystem->PrintHeader(printHeaderName);
}
}

```

因模擬硬碟大小提升至128 MB (1048576 sectors), freeMap大小增加為1048576 bits, 為加速 Bitmap::FindAndSet和Bitmap::NumClear方法, 加入兩個新的私有成員numClear和cur, 前者表示目前剩下多少仍是0的bits, 後者表示下次從此處開始尋找可用的bit。

y. code/lib/bitmap.h

```
class Bitmap
```

```

{
// .....
protected:
    // Number of clear bits
    int numClear;
    // Start looking for clear bit here. Assuming only the first clear bit
    will be marked.
    int cur;
// .....
};

```

Bitmap建構時，將numClear初始化為numItems，cur初始化為0。

Mark時，如果這個bit原本為0，將cur設到此處，並減少numClear。Clear時，如果這個bit位置比cur小，則將cur設到此處，表示下次從此處開始尋找可用的bit，同時，如果這個bit原先為1，則增加numClear。

FindAndSet直接使用numClear判斷目前是否有可用的bit，如果有，則從cur位置開始尋找可用的bit，最後，因為外部都使用FindAndSet和Clear來操作BitMap，不可能發生有可用的0 bit小於cur的狀況，故斷言不會到達最後一行（有剩餘的0 bit但在大於等於cur的部分找不到）。

NumClear改為直接回傳numClear。修改SelfTest，將原先SelfTest中隨機Mark bit的測試案例改掉，因為除了FileSystem建構式中使用Mark(FreeMapSector)和Mark(DirectorySector)，其他都使用FindAndSet和Clear來操作BitMap，而且FreeMapSector為0、DirectorySector為1，並不存在隨機Mark bit的狀況，故以上調整能大幅增加BitMap的效能，且不違反目前的使用案例。

z. code/lib/bitmap.cc

```

Bitmap::Bitmap(int numItems)
{
    // .....
    numClear = numItems;
    cur = 0;
    // .....
}

```

```

void Bitmap::Mark(int which)
{
    ASSERT(which >= 0 && which < numBits);
    if (!Test(which))
    {
        cur = which;
        --numClear;
    }
    map[which / BitsInWord] |= 1 << (which % BitsInWord);
    ASSERT(Test(which));
}

void Bitmap::Clear(int which)
{
    ASSERT(which >= 0 && which < numBits);
    cur = min(cur, which);
    if (Test(which))
    {
        ++numClear;
    }
    map[which / BitsInWord] &= ~(1 << (which % BitsInWord));
    ASSERT(!Test(which));
}

int Bitmap::FindAndSet()
{
    if (!numClear)
    {
        return -1;
    }
    for (int i = cur; i < numBits; i++)
    {
        if (!Test(i))
        {
            Mark(i);
            return i;
        }
    }
    ASSERTNOTREACHED();
}

int Bitmap::NumClear() const

```

```

{
    return numClear;
}

void Bitmap::SelfTest()
{
    int i;

    ASSERT(numBits >= BitsInWord); // bitmap must be big enough

    ASSERT(NumClear() == numBits); // bitmap must be empty
    ASSERT(FindAndSet() == 0);
    Mark(1);
    ASSERT(Test(0) && Test(1));

    ASSERT(FindAndSet() == 2);
    Clear(0);
    Clear(1);
    Clear(2);

    for (i = 0; i < numBits; i++)
    {
        Mark(i);
    }
    ASSERT(FindAndSet() == -1); // bitmap should be full!
    for (i = 0; i < numBits; i++)
    {
        Clear(i);
    }
}

```

使用以下指令執行測試腳本，能在FS\_fileheader\_cout0.log中看到不同file header在硬碟和在記憶體所佔用的大小。

```
bash FS_fileheader.sh > FS_fileheader_cout0.log 2> FS_fileheader_cerr0.log
```

aa. Bonus測試腳本(位於~/NachOS-4.0\_MP4/code/test/FS\_fileheader.sh)

```

../build.linux/nachos -f
../build.linux/nachos -mkdir /t0
../build.linux/nachos -mkdir /t1
../build.linux/nachos -cp num_100.txt /t0/f1

```

```
../build.linux/nachos -cp num_1000.txt /t0/f2
../build.linux/nachos -mkdir /t0/aa
../build.linux/nachos -cp num_100.txt /t0/aa/f1
../build.linux/nachos -cp 64MB.txt /t0/aa/64MB
../build.linux/nachos -l /
echo =====
../build.linux/nachos -lr /
../build.linux/nachos -h /t0/aa/64MB
```



## Part III - Feedback

本次的實作和除錯的時間都遠超前面3個報告，確實學到很多東西，如果有更多的時間，應該還有很多可以優化的部分（註解文件、變數和方法命名、報告用字），但很可惜，將Bonus全做完後，目前剩下44個小時要完成Pthread實作和報告，明天還要出門投票，十分極限。