## **COMP 3311: Database Management Systems**

# Lecture 16 Exercises Query Processing: Expression Evaluation

**Exercise 1:** The Student relation consists of 10,000 tuples sorted on student id. Each student has 5 attributes, each 20 bytes, so the tuple size is 100 bytes. The page size is 1,000 bytes, so  $bf_{Student}=[1000/100]=10$ . Therefore,  $B_{Student}=[1000/10]=1000$  pages. Assume that the buffer size M is 100 pages and that there are 5,000 different student names. There is no index. We want to evaluate the guery on the right.

select distinct name from Student;

| student names. There is no index. We want to evaluate the query on the right. |  |  |  |  |
|---|--|--|--|--|
| a)  | Projection using external sorting              |  |  |  |
|   | Pass 0 (sort) page I/O cost:                   |  |  |  |
|   | Pass 1 (merge) page I/O cost:                  |  |  |  |
|   | Projection page I/O cost:                      |  |  |  |
|   | Number of query result pages written:          |  |  |  |
| b)  | Projection using hashing (using 20 partitions) |  |  |  |
|   | Partitioning page I/O cost:                    |  |  |  |
|   |  |  |  |  |
|   | Duplicate elimination page I/O cost:           |  |  |  |
|   |  |  |  |  |
|   | Projection page I/O cost:                      |  |  |  |
|   | Number of query result pages written:          |  |  |  |
|   |  |  |  |  |

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Exercise 2: Sailor(sailorld, sName, rating, age) Reserves(sailorId, boatId, rDate)

For the Sailor relation, each tuple is 50 bytes, a page can hold 80 tuples and there are 500 full pages. For the Reserves relation, each tuple is 40 bytes, a page can hold 100 tuples and there are 1,000 full pages. There are 10 different sailor ratings and 100 different boats. Assume that sailors are distributed uniformly over the 10 ratings and reservations are distributed uniformly over the 100 boats.

Our goal is to process the query:

select sName from Sailor natural join Reserves where boatId=30 and rating>5;

#### Some useful statistics

- On average, each sailor has 2.5 reservations.
- On average, each boat has 1,000 reservations.
- On average, for each rating there are 4,000 sailors.

| not<br>As:<br>rela | timate the approximate cost of processing the query using a <u>fully pipelined execution method</u> (i.e., of materialize anything except the query result).  Sume that the Sailor relation contains a clustering B+-tree index with 3 levels on sailorld and the Reservation contains a hash index on boatld. Both the B+-tree and hash index can fit 400 index entries page. For non-clustering indexes, each pointer is assumed to lead to a different page. |
|--------------------|---|
|                    |   |
| a)                 | Cost to evaluate $\sigma_{\text{boatId}=30}$  |
|                    | Strategy 1:   |
|                    | Page I/O cost:  |
|                    | Strategy 2:   |
|                    | Page I/O cost:  |
| b)                 | Cost to evaluate σ <sub>rating-5</sub>  |
|                    | Strategy:   |
|                    | Page I/O cost:  |
| c)                 | Cost to evaluate Sailor ⋈ Reserves  |
|                    | Strategy:   |
|                    | Page I/O cost:  |
| d)                 | Cost to evaluate $\pi_{sName}$  |
|                    | Strategy:   |
|                    | Page I/O cost:  |
|                    | Query page I/O cost:  |
|                    |   |

| Name  | P: (1) /  | _ Student#: (1)                        | _ Date: |  |  |  |  |  |  |  |
|---|---|--|---------|--|--|--|--|--|--|--|
|   | e: (2)  | _ Student#: (2)                        |         |  |  |  |  |  |  |  |
|   | Family/Given (PRINT) Given/First (PRINT)          |  |         |  |  |  |  |  |  |  |
|   | NOTE: You are highly encourage                    | ed to do this exercise with a partner. |         |  |  |  |  |  |  |  |
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|   | Lecture 16 Exercises                              |  |         |  |  |  |  |  |  |  |
|   | Query Processing: I                               | Expression Evaluation                  |         |  |  |  |  |  |  |  |
| <b>Exercise 3:</b> The Sailor relation consists of 40,000 tuples sorted on sailorld. Each sailor has 4 attributes, each 10 bytes, so the tuple size is 40 bytes. The page size is 800 bytes so $bf_{Sailor} = [800/40] = 20$ . Therefore, $B_{Sailor} = [40000/20] = 2000$ pages. Assume that the buffer size $M$ is 100 pages and that 5% of sailor names are the same. There is no index. We want to evaluate the query on the right. |   |  |         |  |  |  |  |  |  |  |
| a)  | Projection using external sorting                 |  |         |  |  |  |  |  |  |  |
|   | Pass 0 (sort) page I/O cost:                      |  |         |  |  |  |  |  |  |  |
|   | Pass 1 (merge) page I/O cost:                     |  |         |  |  |  |  |  |  |  |
|   | Projection page I/O cost:                         |  |         |  |  |  |  |  |  |  |
|   | Number of query result pages written:             |  |         |  |  |  |  |  |  |  |
| b)  | Projection using hashing (using 40 partitions; no | o optimization)                        |         |  |  |  |  |  |  |  |
|   | Partitioning page I/O cost:                       |  |         |  |  |  |  |  |  |  |
|   | Duplicate elimination page I/O cost:              |  |         |  |  |  |  |  |  |  |
|   | Projection page I/O cost:                         |  |         |  |  |  |  |  |  |  |

Number of query result pages written:

| Nan   | ne: (1)/  | St        | udent#: (1)                                    | Date:  |  |  |  |  |  |  |
|---|---|-----------|--|--|--|--|--|--|--|--|
|   | Family/Given (PRINT) Given/First (PRINE: (2)  | RINT)     | udent#: (2)                                    |  |  |  |  |  |  |  |
|   | Family/Given (PRINT) Given/First (PRINT)  | RINT)     | do this exercise with a par                    |  |  |  |  |  |  |  |
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|   |   | ture 16 E |  |  |  |  |  |  |  |  |
|   | Query Processing: Expression Evaluation   |           |  |  |  |  |  |  |  |  |
| Ex  | Exercise 4: Student(sld, name, deptId, address) EnrollsIn(courseld, sld, semester, grade) |           |  |  |  |  |  |  |  |  |
| The Student relation contains 10,000 tuples in 1,000 pages and the EnrollsIn relation contains 50,000 tuple in 5,000 pages. There are 25 different departments and 1,000 different courses. All attributes have the same length. Each available index is a B*-tree with 3 levels. For non-clustering indexes, each pointer assumed to lead to a different page. |   |           |  |  |  |  |  |  |  |  |
| Ou  |   | Some use  | eful statistics                                |  |  |  |  |  |  |  |
|   | select name<br>from Student natural join EnrollsIn  |           | rage, a student enrolls i                      |  |  |  |  |  |  |  |
|   | where courseld='COMP3311' and deptld='COMP';  |           | rage, a department has rage, each course has a | 400 students.<br>an enrollment of 50 students. |  |  |  |  |  |  |
| Estimate the approximate cost of processing the query using a fully pipelined execution method (i.e., do not materialize anything except the query result).  Assume that the Student relation contains a clustering index on deptld and the EnrollsIn relation contains a non-clustering index on sld.  |   |           |  |  |  |  |  |  |  |  |
| a)  | Cost to evaluate $\sigma_{\text{courseld='COMP3311'}}$                                    |           |  |  |  |  |  |  |  |  |
|   | Strategy:   |           |  |  |  |  |  |  |  |  |
|   | Page I/O cost:  |           |  |  |  |  |  |  |  |  |
| b)  | Cost to evaluate $\sigma_{\text{deptId='COMP'}}$  |           |  |  |  |  |  |  |  |  |
|   | Strategy:   |           |  |  |  |  |  |  |  |  |
|   | Page I/O cost:  |           |  |  |  |  |  |  |  |  |
| c)  | Cost to evaluate Student ⋈ EnrollsIn  |           |  |  |  |  |  |  |  |  |
|   | Strategy:   |           |  |  |  |  |  |  |  |  |
|   | Page I/O cost:  |           |  |  |  |  |  |  |  |  |
| d)  | Cost to evaluate $\pi_{\text{name}}$  |           |  |  |  |  |  |  |  |  |

Strategy:

Page I/O cost:

Query page I/O cost: