COMP 3311: Database Management Systems

Lecture 18 Exercises Query Optimization

Exercise 1: Given relation R(A, B, C)								
Assume: R contains 10,000 tuples in 1,000 pages. A has 50 distinct values in the range 150. B has 100 distinct values in the range 0100.								
Estimate the size, SC (number of tuples), of each of the following operations assuming uniform distribution and attribute independence.								
a) $\sigma_{A=10}R$								
b) $\sigma_{A=10\land20< B}R$								
c) $\sigma_{C=1}R$								
d) $\sigma_{C=10 \land A=10}R$								

e) $\sigma_{C=10 \land A=10 \land 20 < B}R$

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Exercise 2: Consider the relation Sailor(sailorld, sName, rating, age) and the query:	
$n_{\text{Sailor}} = 10,000 \ B_{\text{Sailor}} = 1,000 \ \text{pages} \ bf_{\text{Sailor}} = \lceil 10,000 \ / \ 1,000 \rceil = 10$	select sName from Sailor
V(rating, Sailor) = 10 (10 distinct rating values)	where rating=7
V(age, Sailor) = 100 (100 distinct age values)	and age=40;
$SC(\text{rating=7, Sailor}) = n_{Sailor} / V(\text{rating=7, Sailor}) = 10,000 / 10 = 1,000 \text{ tuples}$	
SC(age=40, Sailor) = n _{Sailor} / V(age=40, Sailor) = 10,000 / 100 = 100 tuples	
Estimate the page I/O cost to process the query of the following alternative plans distribution and attribute independence. <i>Ignore the cost of searching any indexes</i> .	s assuming uniforr
a) file scan	
b) binary search cost to search on rating	
cost to search on age	
c) single B+-tree index (on either attribute) index on rating	
index on age	
index on age	
d) multiple B+-tree indexes (on both rating and age)	

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Exercise 3: Employee(empld: 4 bytes, name: 35 bytes, title: 2 bytes, salary: 5 bytes, deptld: 4 bytes)

Department(deptld: 4 bytes, name: 25 bytes, location: 7 bytes, projectld: 4 bytes)

DeptProject(deptId: 4 bytes, projectId: 4 bytes)

Project(projectId: 4 bytes, title: 20 bytes, budget: 6 bytes, report: 970 bytes)

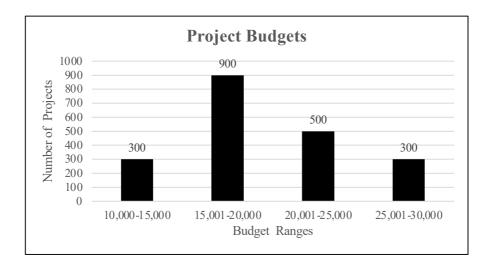
Employee:50 bytes/tuple;20,000 tuples250 pagesDepartment:40 bytes/tuple;100 tuples1 pageDeptProject8 bytes/tuple4,000 tuples8 pagesProject:1,000 bytes/tuple;2,000 tuples500 pages

Page size: 4,000 bytes Buffer pages: 12

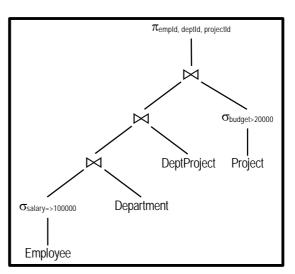
Employee salaries: uniformly distributed in the range 10,000 to 110,000.

Project budgets: distributed in the range 10,000 to 30,000 according to the histogram below.

- There is a clustering B⁺-tree index with 3 levels on salary for Employee.
- There is a hash index on deptld for Department; Department is ordered on deptld.
- There is a hash index on projectld for Project; Project is ordered on projectld.



select distinct empld, deptld, projectld from Employee natural join Department natural join DeptProject natural join Project where salary=>100000 and budget>20000;



Nam	e: (1)	Family/Given (PRINT)	/	Given/First (PRINT)	_ Student#: (1)	Date:
Nam	e : (2)	Family/Given (PRINT)	/	Given/First (PRINT)	_ Student#: (2)	
			You are I		ed to do this exercise with a	partner.
		COM	IP 331	1: Database	e Management Sys	stems
				Lecture 1	8 Exercises	
					ptimization	
				Exe	rcise 3	
a)	Use the	relational algebra	a tree to	estimate the o	utput size of the query in	tuples and in pages.
,		-				
						below. The goal is to minimize to reduce the page I/O cost.
	Where p	oossible, <u>use pip</u>	elining ra	ather than mat	terialization (i.e., keep ir	termediate results in memory
						above. For each step, give the Give the total page I/O cost to
	-	the query.				
		$\sigma_{\text{salary}=>100000}$ Employ	/ee ⇒ re	esult A		
	Strategy	/: lculation:				
	Cost cai	iculation.				
	Step 1 p	page I/O cost:				
	<u>Step 2</u> : ı	result A ⋈ Depar	tment ⇒	result B		
	Strategy	<i>/</i> :				
	Cost cal	lculation:				

Step 2 page I/O cost:

Name: (1) _ Name: (2) _	Family/Given (F	1	Given/First (PRINT)	_ Student#: (1)				
	Family/Given (F		Given/First (PRINT) highly encourag	ed to do this exercise with a part	tner.			
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Step 3	: result B ⋈	DeptProject =	⇒ result C					
Strate								
Cost c	alculation:							
Step 3	page I/O co	ost:						
Stop 4	· æ Di	roioct → rocult	+ D					
Strate(roject ⇒ result						
	alculation:							
Stop 4	naga 1/0 aa	a.t.						
Step 4	page I/O co	osi.						
Step 5	: result C 🖂	result D						
Strate								
Cost c	alculation:							
Step 5	page I/O co	ost:						

Query page I/O cost: