

# COMP 3311: Database Management Systems

## Lecture 20 Exercises Concurrency Control: Lock-based Protocols

**Exercise 1:** a) Is the schedule conflict serializable? ☐ Yes ☐ No

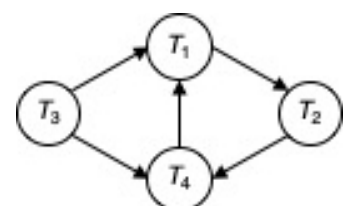
If yes, give the equivalent serial schedule: \_\_\_\_\_

b) Rewrite the schedule according to strict 2PL by adding a lock-s() before reading a data item, a lock-x() before writing a data item and an unlock() after all read/write operations have completed. Is the schedule deadlock free? ☐ Yes ☐ No → Show where the deadlock occurs.

$T_1$	$T_2$	$T_3$
read(X)		
	read(Y)	
	write(Y)	
		write(Z)
write(X)		
	read(X)	
	write(X)	
		read(Y)
		write(Y)
write(Z)		

**Exercise 2:** Which of the following statements is true about the wait-for graph (circle the correct answer)?

- a)  $T_4$  is waiting for  $T_3$  to release a data item.
- b) The system is in a deadlock state after removing the edge between  $T_2$  and  $T_4$ .
- c) The system is in a deadlock state after removing the edge between  $T_3$  and  $T_4$ .
- d) The system is in a deadlock state when  $T_1$  no longer holds a data item needed by  $T_4$ .



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**Exercise 3:** Rewrite the following schedule according to 2PL by adding a lock-s() before reading a data item, a lock-x() before writing a data item and unlock() as necessary.

Is the schedule serializable? ☐ Yes ☐ No → Show where the deadlock occurs.

$T_1$	$T_2$	$T_3$
read(X)		
	read(X)	
		read(Y)
read(Z)		
	read(Y)	
	write(X)	
		read(X)
		write(X)
write(Z)		

**Exercise 4:** In which positions, A to E, can an unlock(X) instruction be inserted if the schedule is according to:

a) strict 2PL (circle the correct answer)

- i. {A} {B} {C} {D}
- ii. {A} {B} {C} {D} {E}
- iii. {A} {C} {D}
- iv. {B} {E}
- v. {A} {C} {D} {E}

b) rigorous 2PL (circle the correct answer)

- i. {A} {B} {C} {D}
- ii. {A} {B} {C} {D} {E}
- iii. {A} {C} {D}
- iv. {B} {E}
- v. {A} {C} {D} {E}

$T_1$	$T_2$
lock-s(X)	
read(X)	
	lock-s(X)
lock-x(Y)	
<b>{A}</b>	
read(Y)	
write(Y)	
	read(X)
	<b>{C}</b>
commit	
unlock(Y)	
<b>{B}</b>	
	<b>{D}</b>
	commit
	<b>{E}</b>

**Exercise 5:** Consider the schedule shown below.

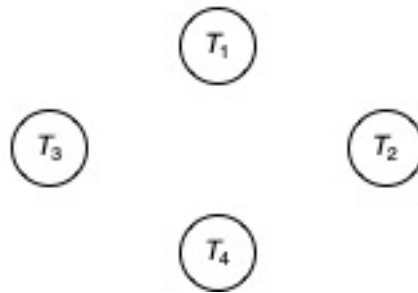
a) Is the schedule conflict serializable? ☐ Yes ☐ No

If yes, give the equivalent serial schedule \_\_\_\_\_

b) If  $T_3$  aborts after write(Y), which other transactions will be rolled back? \_\_\_\_\_

c) If  $T_1$  aborts after write(X), which other transactions will be rolled back? \_\_\_\_\_

d) Construct the wait-for graph that results from this schedule if all locks are only exclusive-locks (lock-x), no locks are released, and the execution process runs to the point of lock-x(Y) in  $T_1$ .



Wait-for Graph

e) Add lock-s(), lock-x() and unlock() instructions to the schedule according to strict 2PL.

$T_1$	$T_2$	$T_3$	$T_4$
read(X)			
write(X)			
	read(X)		
		read(Y)	
		write(Y)	
	write(X)		
			read(Y)
write(Y)			