

## Problem Set 6, Part I

### Problem 1: Counting unique values

1-1)

The worst case occurs when every element in the array is unique.

1-2)

The line comparing `arr[j]` to `arr[i]` will execute for  $(n-1)+(n-2)+\dots+2+1 = n(n-1)/2$  times.

1-3)

The overall time efficiency is  $O(n^2)$  times. In the worst case, every element needs to compare with the rest of the elements in the array. The outer loop will execute for  $n$  times. The inner loop will execute for  $n-1$  times. The comparison of `arr[j]` and `arr[i]` will execute for  $(n-1)+(n-2)+\dots+2+1 = n(n-1)/2 = O(n^2)$  times.

1-4)

The best case occurs when all the elements in the array are the same.

1-5)

The overall time efficiency is  $O(n)$  times. The outer loop will execute for  $n$  times. For each outer loop, the inner loop will execute for 1 time except in the case  $i = \text{arr.length}-1$ . Therefore, the overall time efficiency is  $n-1 = O(n)$  times.

## Problem 2: Improving the efficiency of an algorithm

2-1)

```
public static int numUnique(int[] arr) {
    Sort.mergeSort(arr);
    int count = 1;
    if (arr[0] == arr[arr.length-1]) {
        return count;
    }
    for (int i = 1; i < arr.length; i++) {
        if (arr[i] != arr[i-1]) {
            count++;
        }
    }
    return count;
}
```

2-2)

The worst-case time efficiency is  $O(n \log n)$ , which is more efficient than  $O(n^2)$ . The comparison and move in mergesort are both  $O(n \log n)$ , so the running time is also  $O(n \log n)$ . In the worst case, every element in the array is unique, and the if statement in the for loop will execute for  $n-1$  times. Therefore, the overall time efficiency is  $O(n \log n) + O(n) = O(n \log n)$ .

2-3)

The best-case time efficiency is also  $O(n \log n)$ , which is not more efficient than  $O(n)$ . The running time of mergesort is still  $O(n \log n)$  in the best case. All the elements in the array are the same, so the if statement will execute for 1 times and the for loop will not execute. Therefore, the overall efficiency is  $O(n \log n)$ .

### Problem 3: Practice with references

3-1)

Expression	Address	Value
n	0x128	0x800
n.ch	0x800	'e'
n.next	0x802	0x240
n.prev.next	0x182	0x800
n.next.prev	0x246	0x800
n.next.prev.prev	0x806	0x180

3-2)

```
m.next = n.next;  
n.next.prev = m;  
n.next = m;  
m.prev = n;
```

3-3)

```
public static void addNexts(DNode last) {  
    DNode trav = last;  
    trav.next = null;  
    while (trav.prev != null) {  
        trav.prev.next = trav;  
        trav = trav.prev;  
    }  
}
```

#### **Problem 4: Printing the odd values in a list of integers**

**4-1)**

```
public static void printOddsRecur(IntNode first) {  
    if (first == null) {  
        return;  
    } else {  
        printOddsRecur(first.next);  
        if (first.val % 2 != 0 ) {  
            System.out.println(first.val);  
        }  
    }  
}
```

**4-2)**

```
public static void printOddsIter(IntNode first) {  
    IntNode trav = first;  
    while (trav != null) {  
        if (trav.val % 2 != 0) {  
            System.out.println(trav.val);  
        }  
        trav = trav.next;  
    }  
}
```