

Problem Set 5, Part I

Problem 1: Using recursion to print an array

1-1)

```
public static void print(int[] arr, int start) {  
    if (arr == null || start < 0) {  
        throw new IllegalArgumentException();  
    }  
    if (start >= arr.length) {  
        return;  
    } else {  
        System.out.println(arr[start]);  
        print(arr, start+1);  
    }  
}
```

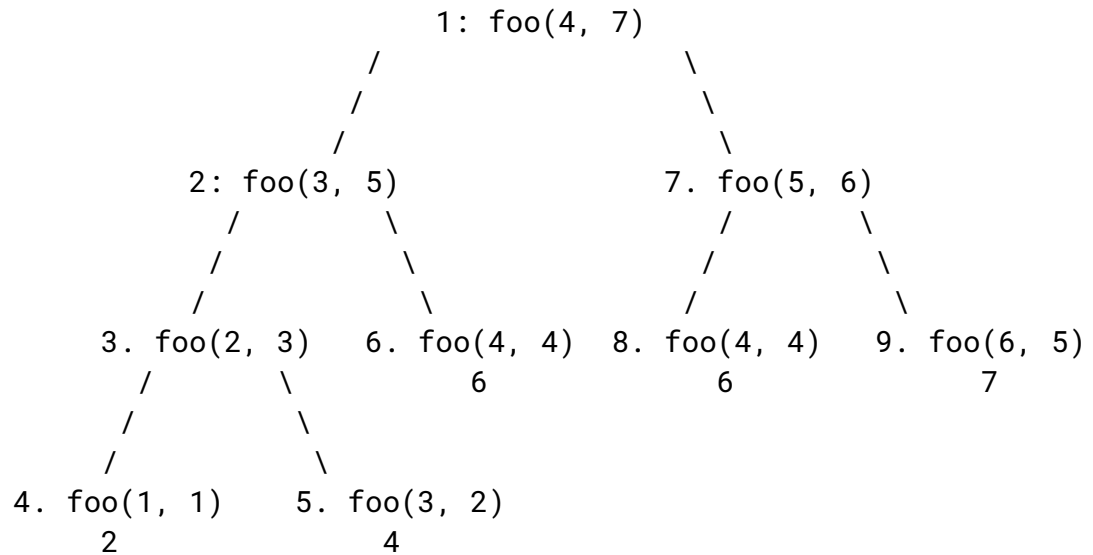
1-2)

```
public static void printReverse(int[] arr, int i) {  
    if (arr == null || i < 0 || i > arr.length) {  
        throw new IllegalArgumentException();  
    }  
    if (i == arr.length) {  
        return;  
    } else {  
        printReverse(arr, i+1);  
        System.out.println(arr[i]);  
    }  
}
```

1-3) **initial call:** printReverse(arr, 0);

Problem 2: A method that makes multiple recursive calls

2-1)



2-2)

call 4 (foo(1, 1)) returns 2
call 5 (foo(3, 2)) returns 4
call 3 (foo(2, 3)) returns 6
call 6 (foo(4, 4)) returns 6
call 2 (foo(3, 5)) returns 12
call 8 (foo(4, 4)) returns 6
call 9 (foo(6, 5)) returns 7
call 7 (foo(5, 6)) returns 13
call 1 (foo(4, 7)) returns 25

Problem 3: Sorting practice

3-1)

{7, 10, 13, 27, 24, 20, 14, 33}

3-2)

{7, 13, 14, 24, 27, 20, 10, 33}

3-3)

{7, 13, 14, 20, 10, 24, 27, 33}

3-4)

{10, 7, 13, 27, 24, 20, 14, 33}

3-5)

{7, 10, 13, 27, 24, 20, 14, 33}

3-6)

{7, 13, 14, 27, 24, 20, 10, 33}

Problem 4: Practice with big-O

4-1)

function	big-O expression
$a(n) = 5n + 1$	$a(n) = O(n)$
$b(n) = 2n^3 + 3n^2 + n\log(n)$	$b(n) = O(n^3)$
$c(n) = 10 + 5n\log(n) + 10n$	$c(n) = O(n\log n)$
$d(n) = 4\log(n) + 7$	$d(n) = O(\log n)$
$e(n) = 8 + n + 3n^2$	$e(n) = O(n^2)$

4-2)

$O(n^2)$

The outer loop will be executed for $2n = O(n)$ times. As the outer loop executes once, the inner loop will be executed for $n-1 = O(n)$ times. Therefore, the count() method is called for $2n(n-1) = O(n^2)$ times.

4-3)

$O(n\log n)$

The i loop will be executed for 5 times. As the i loop executes once, the j loop will be executed for $n = O(n)$ times. As the j loop executes once, the k loop will be executed for $\log(n) = O(\log n)$ times. Therefore, the count() method is called for $5 \cdot n \cdot \log n = O(n\log n)$ times.

Problem 5: Comparing two algorithms

worst-case time efficiency of algorithm A: $O(n \log n)$

Explanation: The worst case is the array in order from the largest to the smallest. The algorithm uses Mergesort, and there are $\log(n)$ levels in the call tree. Merging two halves of an array of size n requires $2n$ moves, so the number of moves in each level is $2n$. The total number of moves is $2n * \log(n) = O(n \log n)$. The number of comparisons is also $O(n \log n)$. Thus, the running time is $O(n \log n)$.

worst-case time efficiency of algorithm B: $O(n)$

Explanation: The worst case is the array in order from the smallest to the largest. The i in the for loop will increment from 0 to $n-1$. The number of comparison is $O(n)$. The number of moves is also $O(n)$. Thus, the running time is $O(n)$.