

## Problem Set 7, Part I

### Problem 1: Choosing an appropriate representation

#### 1-1) *ArrayList* or *LinkedList*? *ArrayList*

*Explanation:* As the number of events is roughly the same, we can predict the maximum length. So, the space efficiency of *ArrayList* always depends on the anticipated maximum length. It is also more time-efficient to use *ArrayList* because the events are added in order. The *ArrayList* does not need to shift, so it's the best case that only needs  $O(1)$  to add an event.

#### 1-2) *ArrayList* or *LinkedList*? *ArrayList*

*Explanation:* As we have to frequently access the runners' records through the id number as index, we need to use *ArrayList* which provides random access with time efficiency of  $O(1)$ . If we use *LinkedList*, we need to traverse the list to find the item, which is not time-efficient. Also, the length of the list is almost constant. It means that we can predict the maximum length of the *ArrayList*, and the space efficiency depends on the anticipated maximum length. There are very few changes to the list after the deadline, so the *ArrayList* does not need to spend much time on shifting.

#### 1-3) *ArrayList* or *LinkedList*? *LinkedList*

*Explanation:* The number of registrants varies significantly, so it is more space-efficient to use *LinkedList* which space efficiency depends on the number of added items in the list. *ArrayList* is not space-efficient because we have to preallocate the memory needed for the maximum sequence size. For time efficiency, we always add the recent one to the front, so it is more time-efficient to use *LinkedList* to add to the front of the list, which only needs  $O(1)$ .

### Problem 2: Scaling a list of integers

#### 2-1)

The for loop will execute for  $n$  times. For each for loop, the time efficiency of `getItem()` is  $O(1)$ , because *ArrayList* provides random access. And the time efficiency of `addItem` depends on the length  $n$ , because it needs to traverse the full list to add the item. Therefore, the running time of the algorithm is  $n*1 + (1 + 2 + \dots + n) = O(n^2)$ .

#### 2-2)

```
public static LinkedList scale(int factor, ArrayList vals) {
    LinkedList scaled = new LinkedList();

    for (int i = vals.length()-1; i >= 0; i--) {
        int val = (Integer)vals.getItem(i);
        scaled.addItem(val*factor, 0);
    }
    return scaled;
}
```

**2-3)**

The for loop will execute for  $n$  times. For each for loop, as we always add the item to the front of the LLList, the time efficiency of `addItem()` is  $O(1)$  because there is no need to walk down. Also, the time efficiency of `getItem()` is  $O(1)$  because `ArrayList` provides random access. Therefore, the running time of the algorithm is  $O(n)$ .

**Problem 3: Working with stacks and queues****3-1)**

```
public static void remAllStack(Stack<Object> stack, Object item) {
    Stack<Object> stack2 = new LLStack<Object>();
    while (!stack.isEmpty()) {
        if (stack.peek().equals(item)) {
            stack.pop();
        } else {
            Object O1 = stack.pop();
            stack2.push(O1);
        }
    }
    while (!stack2.isEmpty()) {
        Object O2 = stack2.pop();
        stack.push(O2);
    }
}
```

**3-2)**

```
public static void remAllQueue(Queue<Object> queue, Object item) {
    Queue<Object> queue2 = new LLQueue<Object>();
    while (!queue.isEmpty()) {
        if (queue.peek().equals(item)) {
            queue.remove();
        } else {
            Object O1 = queue.remove();
            queue2.insert(O1);
        }
    }
    while (!queue2.isEmpty()) {
        Object O2 = queue2.remove();
        queue.insert(O2);
    }
}
```

**Problem 4: Binary tree basics**

4-1)

The height is 3.

4-2)

It has 4 leaf nodes and 5 interior nodes.

4-3)

21 18 7 25 19 27 30 26 35

4-4)

7 19 25 18 26 35 30 27 21

4-5)

21 18 27 7 25 30 19 26 35

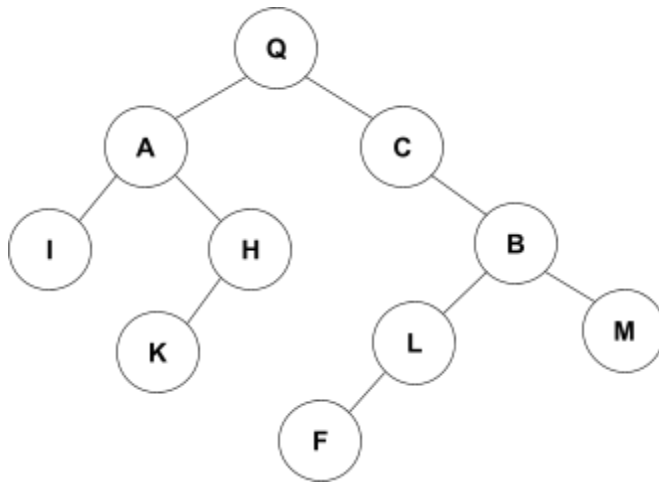
4-6)

No. For Node 25, 25 is greater than 21, so it should be on the right side of Node 21. For Node 26, 26 is smaller than 27, so it should be on the left side of Node 27.

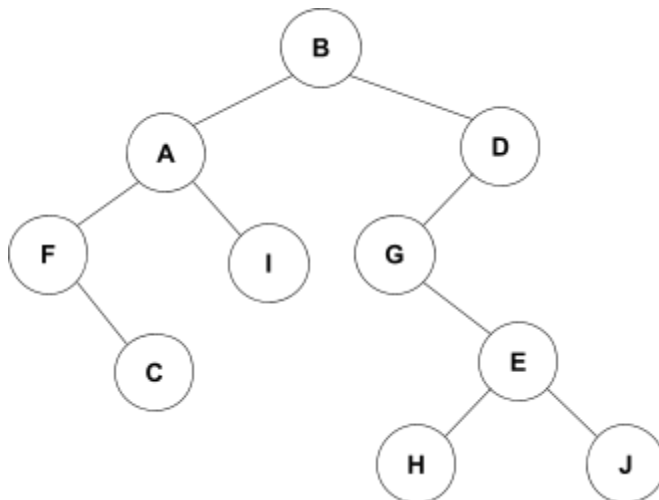
4-7)

Yes. The tree is balanced because the node's subtrees have the same height or have heights that differ by 1.

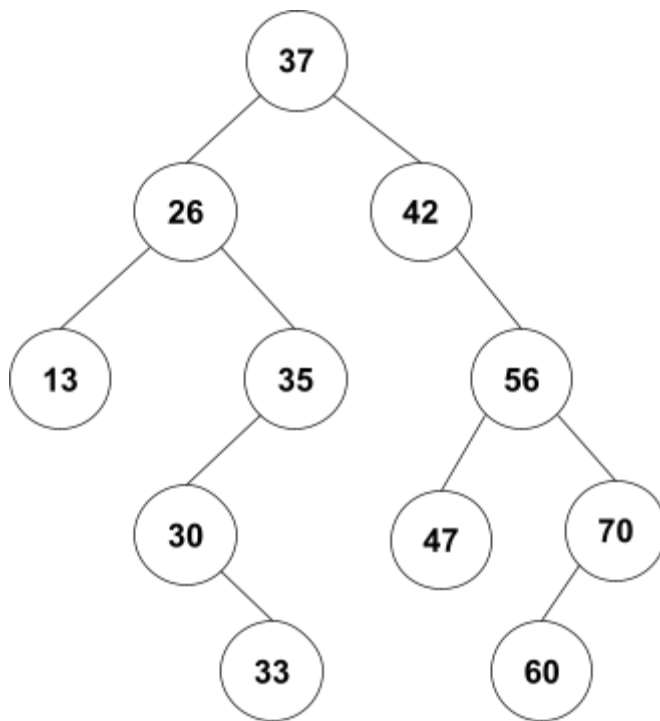
**Problem 5: Tree traversal puzzles**  
**5-1)**



**5-2)**



**Problem 6: Binary search trees**  
6-1)



6-2)

