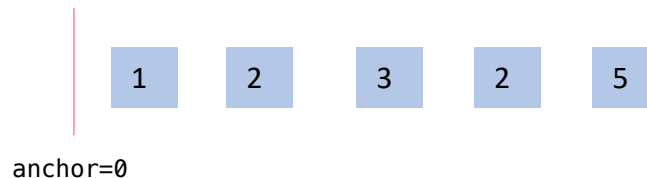


Sliding window:

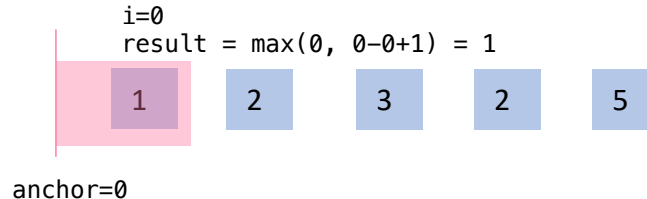
Let's say you are given the problem: for a given array  $L$  find the longest increasing subsequence, defined as a contiguous range made up of only increasing numbers. If  $L$  is  $[1, 2, 3, 2, 5]$ , the increasing subsequences of this array are  $[1, 2, 3]$  and  $[2, 5]$ . Of those,  $[1, 2, 3]$  is longer, so we return  $\text{len}([1, 2, 3]) \Rightarrow 3$ .

We instantiate our left-hand boundary, "anchor", and our result, "result", as 0. In our index-based loop, we check for if  $i > 0$  (as to prevent comparing the first and last indices) and if we are *not* still in an increasing subsequence. That's defined by the value at the previous index of  $L$  being  $\geq$  the value at our current index of  $L$ . If this is true, our left-hand boundary *anchor* must be dragged across to our current  $i$  as to reset the size of our window to 0.

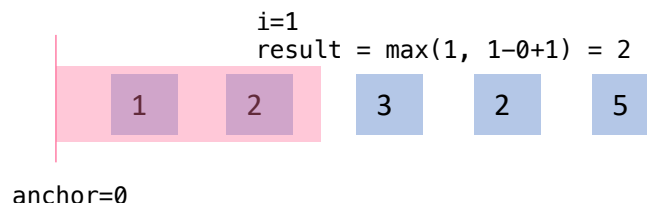
Let's follow along. We start with *anchor* being 0,



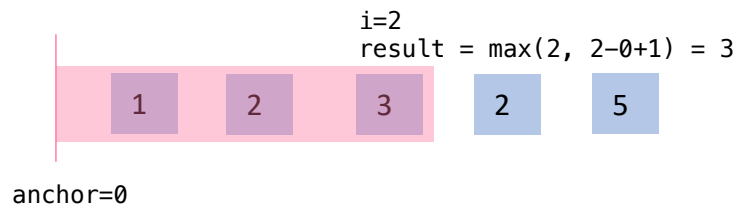
As we iterate over the indices of  $L$ , so long as  $i > 0$  we check to see if  $\text{val. @ index } i-1 \geq \text{index @ our current val.}$  Let's see how our window (in pink) changes:



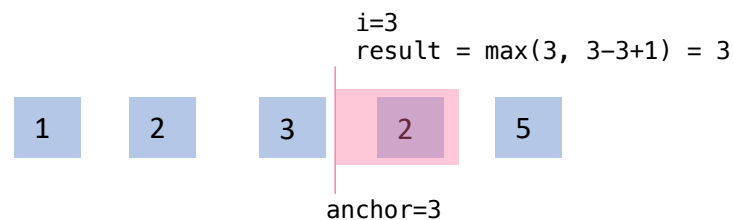
Now *result* has been set to 1, since our window covers 1 number. When  $i = 1$ , we evaluate the statement  $L[i-1] \geq L[i]$  as  $L[1-1] \geq L[1]$ , which equals  $1 \geq 2$ . This is false, indicating our subsequence is still increasing. That means our pink window will get 1 larger:



This continues for  $i=2$ , as  $L[i-1] \geq L[i] \rightarrow L[2-1] \geq L[2] \rightarrow L[1] \geq L[2] \rightarrow 2 \geq 3$ . The statement is False, meaning our anchor does *not* get moved, and our window keeps increasing.

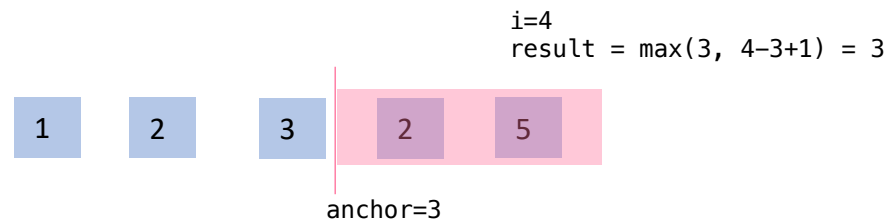


When we hit  $i=3$ , let's see what happens:



Since  $L[i-1] \geq L[i] \rightarrow L[3-1] \geq L[3] \rightarrow L[2] \geq L[3] \rightarrow 3 \geq 2$  evaluates to *True*, our previous val is bigger than our current, we drag our anchor our to  $i$ , which is 3, effectively resetting the size of our window to 1, only encompassing the current val.

Finally, we go over to the last index:



The code for this is shown below:

```
def longest_sub(L) -> int:
    anchor, result = 0, 0
    for i in range(len(L)):
        if i > 0 and L[i-1] >= L[i]: anchor = i
        result = max(result, i - anchor + 1)
    return result
```