

Collaborated with yc2454 (Yalu Cai)

- a. Assume for the sake of contradiction that the halting problem is computable, so that given a machine  $M$  and its input  $x$ , we know whether  $M$  halts on  $x$  (1 for yes, 0 for no). If  $M$  does halt on  $x$ , then let  $S$  be the set of all trivial functions  $f$  where  $f(x) = 1$ ; similarly, if  $M$  does not halt on  $x$ , then let  $S$  be the set of all trivial functions  $f$  where  $f(x) = 0$ . Then in either these cases,  $f_S$  would be computable, and the answer would be yes, the given machine  $M$  computes one of the functions in  $S$ . But Rice's theorem says  $f_S$  is uncomputable, reaching a contradiction. Thus, the halting problem is uncomputable.
- b. We assume that  $f_S$  is computable and reach the contradiction that this would allow the halting problem to be computable. Given a Turing machine  $M$  and its input  $x$ , the halting problem is concerned with whether  $M$  halts on  $x$ . Let a set  $S$  be a non-trivial set of partial functions. We make  $S$  only have one partial function  $f_\emptyset$ , which is undefined for every input. We make a new Turing machine  $M_b$  that loops infinitely for every input other than  $x$  and just simulates  $M$  on  $x$  for  $x$ , i.e.  $M_b$ 's pseudocode looks like this:

```
M_b(input){  
    if (input!=x){  
        loop infinitely;  
    }  
    else{  
        return M(x);  
    }  
}
```

If there were a Turing machine that computes  $f_S$ , then we would know whether  $M_b$  halts on  $x$ , and thus whether  $M$  halts on  $x$ . Because if  $f_S$  outputs 1 on input  $M_b$ , then we know  $M_b$  doesn't halt on  $x$  because that means  $M_b$  computes  $f_\emptyset$ , which is undefined for  $x$ . The complement of  $S$  (let's call it  $S'$ ) is also a non-trivial set of partial functions. If there were a Turing machine that computes  $f_{S'}$  and outputs 1 on input  $M_b$ , then we know  $M_b$  (and thus,  $M$ ) halts on  $x$ , because then that means there's a function in  $S'$  that is undefined for every input other than  $x$  and is defined for  $x$ , and has output  $M_b(x)$ . Either of these Turing machines that decide  $f_S$  or  $f_{S'}$  is a decider for the halting problem, which we know is uncomputable. Rice's theorem is thus proven.