

ECE108 Assignment2

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March 29, 2018

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1 Formalizing an Argument

1.1 Big Bang Theory

If the Big Bang Theory is correct then either there was a time before anything existed or the world will come to an end. The world will not come to an end. Therefore, if there was no time before anything existed, the Big Bang Theory is incorrect.

P : the Big Bang Theory is correct

Q : there was a time before anything existed

R : the world will come to an end

$$P \implies (Q \vee R) \quad (1)$$

$$\neg R \quad (2)$$

$$\text{as a conclusion : } \neg Q \implies \neg P \quad (3)$$

1. *Show* $\neg Q \rightarrow \neg P$

2.	$\neg Q$	1,CD
3.	<i>Show</i> $\neg P$	
4.	P	3,ID
5.	$P \rightarrow (Q \vee R)$	P2
6.	$Q \vee R$	4,5, MP
7.	$\neg R$	P1
8.	Q	5,6,MTP
9.	$\neg Q$	2

1.2 To Win a Gold Medal

To win a gold medal, an athlete must be very fit. If s/he does not win a gold medal, then either s/he arrived late for the competition or his/her training was interrupted. If s/he is not very fit, s/he will blame his/her coach. If s/he blames his/her coach, or his/her training is interrupted, then s/he will not get into the competition. Therefore, if s/he gets into the competition, s/he will not have arrived late.

P : win a gold medal

Q : an athlete must be very fit

R : s/he arrived late for the competition

S : his/her training was interrupted.

T : s/he will blame his/her coach

U : if s/he gets into the competition

$$P \implies Q \tag{4}$$

$$\neg P \implies (R \vee S) \tag{5}$$

$$\neg Q \implies T \tag{6}$$

$$(T \vee S) \implies \neg U \tag{7}$$

$$\text{as conclusion: } U \implies \neg R \tag{8}$$

1.3 Hector and the Battle of Priam

If Hector wins the battle, he will plunder the city. If he does not win the battle, he will either be killed or go into exile. If he plunders the city, then Priam will lose his kingdom. If Priam loses his kingdom or Hector goes into exile, then the war will end. Therefore, if the war does not end, Hector will be killed.

P : Hector wins the battle

Q : Hector plunders the city

R : Hector is killed

S : Hector goes into exile

T : Priam will lose his kingdom

U : war will end

$$P \implies Q \tag{9}$$

$$\neg P \implies (R \vee S) \tag{10}$$

$$(T \vee S) \implies U \tag{11}$$

$$\text{as conclusion: } \neg U \implies R \tag{12}$$

1.4 colonel and the murder

If the colonel was out of the room when the murder was committed then he couldn't have been right about the weapon used. Either the butler is lying or he knows who the murderer was. If Lady Barntree was not the murderer then either the colonel was in the room at the time or the butler is lying. Either the butler knows who the murderer was or the colonel was out of the room at the time of the murder. Therefore, if the colonel was right about the weapon then Lady Barntree was the murderer.

P : the colonel was out of the room when the murder was committed

Q : colonel couldn't have been right about the weapon used

R : the butler is lying

S : the butler knows who the murderer was

T : Lady Barntree was not the murderer

$$P \implies Q \tag{13}$$

$$R \vee S \tag{14}$$

$$T \implies (\neg P \vee R) \tag{15}$$

$$S \vee P \tag{16}$$

$$\text{as conclusion: } \neg Q \implies \neg T \tag{17}$$

2 Tautologies and Friends

$$(P \wedge Q) \rightarrow (P \rightarrow Q)$$

P	Q	$P \wedge Q$	$P \rightarrow Q$	$(P \wedge Q) \rightarrow (P \rightarrow Q)$
F	F	F	T	T
F	T	F	T	T
T	F	F	F	T
T	T	T	T	T

tautology

$$(P \wedge Q) \leftrightarrow (P \rightarrow Q)$$

P	Q	$P \wedge Q$	$P \rightarrow Q$	$(P \wedge Q) \iff (P \rightarrow Q)$
F	F	F	T	F
F	T	F	T	F
T	F	F	F	T
T	T	T	T	T

not tautology

$$(\neg P \vee Q) \rightarrow (P \rightarrow \neg Q)$$

P	Q	$\neg P \vee Q$	$P \rightarrow \neg Q$	$(\neg P \vee Q) \rightarrow (P \rightarrow \neg Q)$
F	F	T	T	T
F	T	T	T	T
T	F	F	T	T
T	T	T	F	F

not tautology

$$(((P \rightarrow Q) \rightarrow P) \rightarrow Q)$$

P	Q	$(P \rightarrow Q)$	$((P \rightarrow Q) \rightarrow P)$	$((P \rightarrow Q) \rightarrow P) \rightarrow Q$
F	F	T	F	T
F	T	T	F	T
T	F	F	T	F
T	T	T	T	T

not tautology

$$(P \rightarrow (Q \rightarrow (P \rightarrow Q)))$$

P	Q	$(P \rightarrow Q)$	$(Q \rightarrow (P \rightarrow Q))$	$(P \rightarrow (Q \rightarrow (P \rightarrow Q)))$
F	F	T	T	T
F	T	T	T	T
T	F	F	T	T
T	T	T	T	T

tautology

$$((P \wedge \neg Q) \rightarrow \neg R) \leftrightarrow ((P \wedge R) \rightarrow Q)$$

P	Q	R	$(P \wedge \neg Q)$	$((P \wedge \neg Q) \rightarrow \neg R)$	$P \wedge R$	$(P \wedge R) \rightarrow Q$
F	F	F	F	T	F	T
F	F	T	F	T	F	T
F	T	F	F	T	F	T
F	T	T	F	T	F	T
T	F	F	T	T	F	T
T	F	T	T	F	T	F
T	T	F	F	T	F	T
T	T	T	F	T	T	T

we see that column 3 = column 5

the if and only if is true, so tautology

$$(((P \vee Q) \vee R) \vee S) \leftrightarrow (P \vee (Q \vee (R \vee S)))$$

P	Q	R	S	$P \vee Q$	$P \vee Q) \vee R$	$((P \vee Q) \vee R) \vee S$
F	F	F	F	F	F	F
F	F	F	T	F	F	T
F	F	T	F	F	T	T
F	F	T	T	F	T	T
F	T	F	F	T	T	T
F	T	F	T	T	T	T
F	T	T	F	T	T	T
F	T	T	T	T	T	T
T	F	F	F	T	T	T
T	F	F	T	T	T	T
T	F	T	F	T	T	T
T	F	T	T	T	T	T
T	T	F	F	T	T	T
T	T	F	T	T	T	T
T	T	T	F	T	T	T
T	T	T	T	T	T	T

$R \vee S$	$Q \vee (R \vee S)$	$P \vee (Q \vee (R \vee S))$
F	F	F
T	T	T
T	T	T
T	T	T
T	T	T
F	T	T
T	T	T
T	T	T
T	T	T
F	F	T
T	T	T
T	T	T
T	T	T
F	T	T
T	T	T
T	T	T
T	T	T

as anyone can clearly see from commutativity of OR, this is a tautology
 $((P \rightarrow Q) \rightarrow R) \rightarrow S \leftrightarrow (P \rightarrow (Q \rightarrow (R \rightarrow S)))$

P	Q	R	S	$P \rightarrow Q$	$(P \rightarrow Q) \rightarrow R$	$((P \rightarrow Q) \rightarrow R) \rightarrow S$
F	F	F	F	T	F	T
F	F	F	T	T	F	T
F	F	T	F	T	T	F
F	F	T	T	T	T	T
F	T	F	F	T	F	F
F	T	F	T	T	F	T
F	T	T	F	T	T	F
F	T	T	T	T	T	T
T	F	F	F	F	T	F
T	F	F	T	F	T	T
T	F	T	F	F	T	F
T	F	T	T	F	T	T
T	T	F	F	T	F	F
T	T	F	T	T	F	T
T	T	T	F	T	T	F
T	T	T	T	T	T	T

P	Q	R	S	$R \rightarrow S$	$Q \rightarrow (R \rightarrow S)$	$P \rightarrow (Q \rightarrow (R \rightarrow S))$
F	F	F	F	T	T	T
F	F	F	T	T	T	T
F	F	T	F	F	T	T
F	F	T	T	T	T	T
F	T	F	F	T	T	T
F	T	F	T	T	T	T
F	T	T	F	F	F	T
F	T	T	T	T	T	T
T	F	F	F	T	T	T
T	F	F	T	T	T	T
T	F	T	F	F	T	T
T	F	T	T	T	T	T
T	T	F	F	T	T	T
T	T	F	T	T	T	T
T	T	T	F	F	F	F
T	T	T	T	T	T	T

after brainlessly bruteforce everything in a truth table, we see that this is not a tautology.

now... if I was smart I could have seen that the the only way for $P \rightarrow (Q \rightarrow (R \rightarrow S))$ to be true is S:False, P: True

$$(P \rightarrow (\neg R \rightarrow \neg S)) \vee ((S \rightarrow (P \vee \neg T)) \vee (\neg Q \rightarrow R))$$

P	Q	R	S	T	$\neg R \rightarrow \neg S$	$(P \rightarrow (\neg R \rightarrow \neg S))$
F	F	F	F	F	T	T
F	F	F	F	T	T	T
F	F	F	T	F	F	T
F	F	F	T	T	F	T
F	F	T	F	F	T	T
F	F	T	F	T	T	T
F	F	T	T	F	T	T
F	F	T	T	T	T	T
F	T	F	F	F	T	T
F	T	F	F	T	T	T
F	T	F	T	F	F	T
F	T	F	T	T	F	T
F	T	T	F	F	T	T
F	T	T	F	T	T	T
F	T	T	T	F	T	T
F	T	T	T	T	T	T
T	F	F	F	F	T	T
T	F	F	F	T	T	T
T	F	F	T	F	F	F
T	F	F	T	T	F	F
T	F	T	F	F	T	T
T	F	T	F	T	T	T
T	F	T	T	F	T	T
T	F	T	T	T	T	T
T	T	F	F	F	T	T
T	T	F	F	T	T	T
T	T	F	T	F	F	F
T	T	F	T	T	F	F
T	T	T	F	F	T	T
T	T	T	F	T	T	T
T	T	T	T	F	T	T
T	T	T	T	T	T	T

ok now, since the rest is linked by OR \vee , we only need to find 4 cases where the first statement doesn't evaluate to true.

P	Q	R	S	T	$(P \rightarrow (\neg R \rightarrow \neg S))$	$P \vee T$
T	F	F	T	F	F	T
T	F	F	T	T	F	T
T	T	F	T	F	F	T
T	T	F	T	T	F	T

well we see in $S \rightarrow (P \vee T)$ that if P is True, then the whole thing is true no matter T or S

therefore this is a tautology

3 Semantic Tableaux

$$(P \wedge Q) \rightarrow (P \rightarrow Q) : F$$

|

$$(P \wedge Q) : T, (P \rightarrow Q) : F$$

|

$$P : T, Q : F$$

|

$$P : T, Q : T$$

|

X

$$(P \wedge Q) \leftrightarrow (P \rightarrow Q)$$

$$(P \wedge Q) \leftrightarrow (P \rightarrow Q) : F$$

└──┬──┘

$$(P \wedge Q) : T, (P \rightarrow Q) : F$$

$$(P \wedge Q) : F, (P \rightarrow Q) : T$$

|

$$P : T, Q : F$$

$$P : F \quad Q : F$$

|

$$P : T, Q : T$$

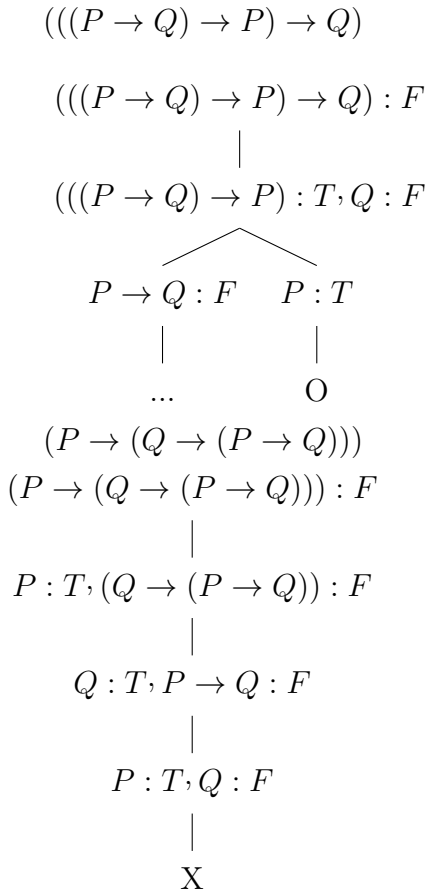
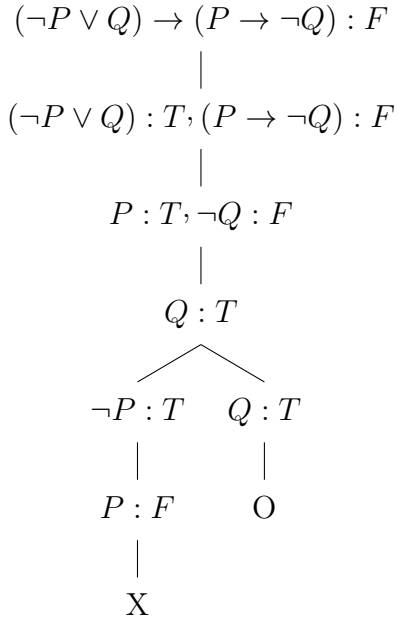
$$P : F, Q : T$$

|

X

O

$$(\neg P \vee Q) \rightarrow (P \rightarrow \neg Q)$$



$$((P \wedge \neg Q) \rightarrow \neg R) \leftrightarrow ((P \wedge R) \rightarrow Q)$$

$$((P \wedge \neg Q) \rightarrow \neg R) \leftrightarrow ((P \wedge R) \rightarrow Q)) : F$$

$$((P \wedge \neg Q) \rightarrow \neg R) : T, ((P \wedge R) \rightarrow Q)) : F \quad ((P \wedge \neg Q) \rightarrow \neg R) : F, ((P \wedge R) \rightarrow Q)) : T$$

$$P \wedge R : T, Q : F$$

$$(P \wedge \neg Q) : T, \neg R : F$$

$$P : T, R : T$$

$$R : T$$

$$P \wedge \neg Q : F$$

$$\neg R : T$$

$$P : T, \neg Q : T$$

$$P : F$$

$$\neg Q : F$$

$$R : F$$

$$Q : F$$

$$X$$

$$Q : T$$

$$X$$

$$(P \wedge R) : F$$

$$Q : T$$

$$X$$

$$P : F$$

$$R : F$$

$$X$$

$$X$$

$$X$$

$$(((P \vee Q) \vee R) \vee S) \leftrightarrow (P \vee (Q \vee (R \vee S)))$$

$$\begin{array}{c}
(((P \vee Q) \vee R) \vee S) \leftrightarrow (P \vee (Q \vee (R \vee S))) : F \\
\swarrow \quad \searrow \\
((P \vee Q) \vee R) \vee S : T, (P \vee (Q \vee (R \vee S))) : F \quad ((P \vee Q) \vee R) \vee S : F \text{ (AND)} \\
\qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \downarrow \\
\qquad \qquad \qquad \qquad \qquad \qquad \qquad \dots(P \vee (Q \vee (R \vee S))) : T
\end{array}$$

4 Kalish-Montague Derivations

$$(P \wedge Q) \rightarrow (P \rightarrow Q)$$

- | | | |
|----|---|--------|
| 1. | $Show (P \wedge Q) \rightarrow (P \rightarrow Q)$ | |
| 2. | $P \wedge Q$ | 1,CD |
| 3. | $Show P \rightarrow Q$ | |
| 4. | P | 3,CD |
| 5. | P | 2,SIMP |
| 6. | Q | 2,SIMP |

$$(P \rightarrow (Q \rightarrow (P \rightarrow Q)))$$

1. *Show* ($P \rightarrow (Q \rightarrow (P \rightarrow Q))$)

2.	P	1,CD
3.	<i>Show</i> $Q \rightarrow (P \rightarrow Q)$	1,subDer
4.	Q	3,CD
5.	<i>Show</i> $P \rightarrow Q$	3,subDerv
6.	P	5,CD
7.	Q	4

$$((P \wedge \neg Q) \rightarrow \neg R) \leftrightarrow ((P \wedge R) \rightarrow Q)$$

1.	<i>Show</i> $((P \wedge \neg Q) \rightarrow \neg R) \rightarrow ((P \wedge R) \rightarrow Q)$	
2.	$(P \wedge \neg Q) \rightarrow \neg R$	1,CD
3.	<i>Show</i> $(P \wedge R) \rightarrow Q$	2,subDer
4.	$P \wedge R$	3,CD
5.	P	4,SIMP
6.	R	4,SIMP
7.	$\neg\neg R$	6,DN
8.	$\neg(P \wedge \neg Q)$	2,7,MT
9.	<i>Show</i> Q	8,subDerv
10.	$\neg Q$	9,ID
11.	$P \wedge \neg Q$	5,10,ADJ
12.	$\neg(P \wedge \neg Q)$	8

1.	$\textit{Show} ((P \wedge \neg Q) \rightarrow \neg R) \leftarrow ((P \wedge R) \rightarrow Q)$	
2.	$(P \wedge R) \rightarrow Q$	1,CD
3.	$\textit{Show} (P \wedge \neg Q) \rightarrow \neg R$	subDer
4.	$P \wedge \neg Q$	3,CD
5.	P	4,SIMP
6.	$\neg Q$	4,SIMP
7.	$\neg(P \wedge R)$	2,6,MT
8.	$\textit{Show} \neg R$	subDer
9.	R	8,ID
10.	$P \wedge R$	5,9,ADJ
11.	$\neg(P \wedge R)$	7

$$(((P \vee Q) \vee R) \vee S) \leftrightarrow (P \vee (Q \vee (R \vee S)))$$

1. *Show* $((P \vee Q) \vee R) \vee S \rightarrow (P \vee (Q \vee (R \vee S)))$

2.	$((P \vee Q) \vee R) \vee S$	1,CD
3.	<i>Show</i> $P \vee (Q \vee (R \vee S))$	2,subDer
4.	$\neg(P \vee (Q \vee (R \vee S)))$	3,ID
5.	P	4,SIMP
6.	R	4,SIMP
7.	$\neg\neg R$	6,DN
8.	$\neg(P \wedge \neg Q)$	2,7,MT
9.	<i>Show</i> Q	8,subDerv
10.	$\neg Q$	9,ID
11.	$P \wedge \neg Q$	5,10,ADJ
12.	$\neg(P \wedge \neg Q)$	8

5 Premises

if our premises form a contradiction, then

$$\neg(P_1 \wedge P_2 \cdots \wedge P_N) \rightarrow T$$

we turn it into a tautology, then we do what we are used to do:

we would then use a semantic tableau to show contradiction will always occur if we start with:

$$\neg(P_1 \wedge P_2 \cdots \wedge P_N) : F$$

this will end up being:

$$(P_1 \wedge P_2 \cdots \wedge P_N) : T$$

we would go through and split all the elements inside the ANDs showing that our starting logic equation is indeed a tautology

we would also be able to use KM derivation indeed we would start with

$$\neg(P_1 \wedge P_2 \cdots \wedge P_N)$$

then use the simplification inference rule on the ANDs to simplify the equation. We would end up

6 Normal Form

7 Generalized DeMorgan's Laws