
COMP 472 Artificial Intelligence:

Adversarial Search *part 4*

Alpha-Beta Pruning *video 2*

- Russell & Norvig: Sections 5.1, 5.2, 5.3

Today

- Adversarial Search

1. Minimax

2. Alpha-Beta Pruning



Alpha-Beta Pruning

是对Mini Max的一个优化

Mini Max会计算n-ply那一层所有的 $e(n)$

■ Optimization over Minimax, that:

- ignores (cuts off/prunes) branches of the tree that cannot contribute to the solution

这里提前截断了那些不可能的subtree

因此那些subtree的 $e(n)$ 既不用算了，减少了 en 运算量

- reduces branching factor

~~$e(n)$~~

- allows deeper search with same effort

减少的运算量允许你进行更深的search



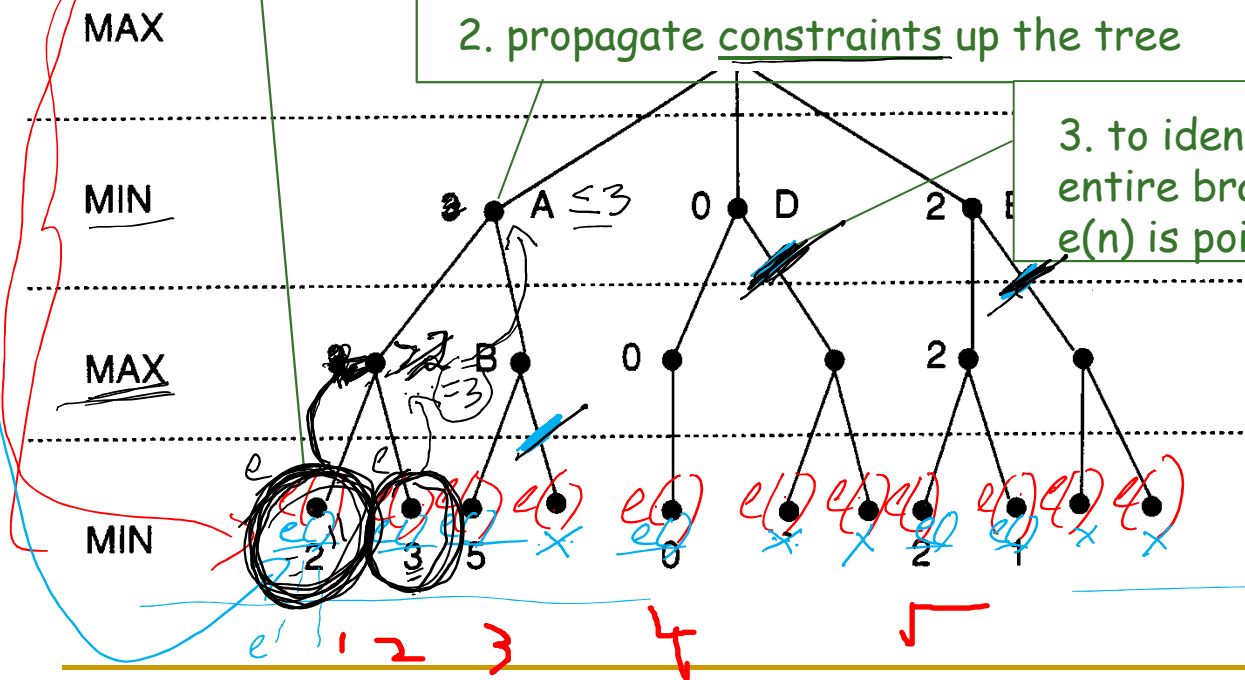
Alpha-Beta Pruning: Example 1

- With Minimax, we look at all nodes at depth n mini Max看全部
- With α - β pruning, we ignore branches that could not possibly contribute to the final decision

1. run $e(n)$ on leaf nodes in a depth first traversal manner on a necessity basis (not on all nodes)

2. propagate constraints up the tree

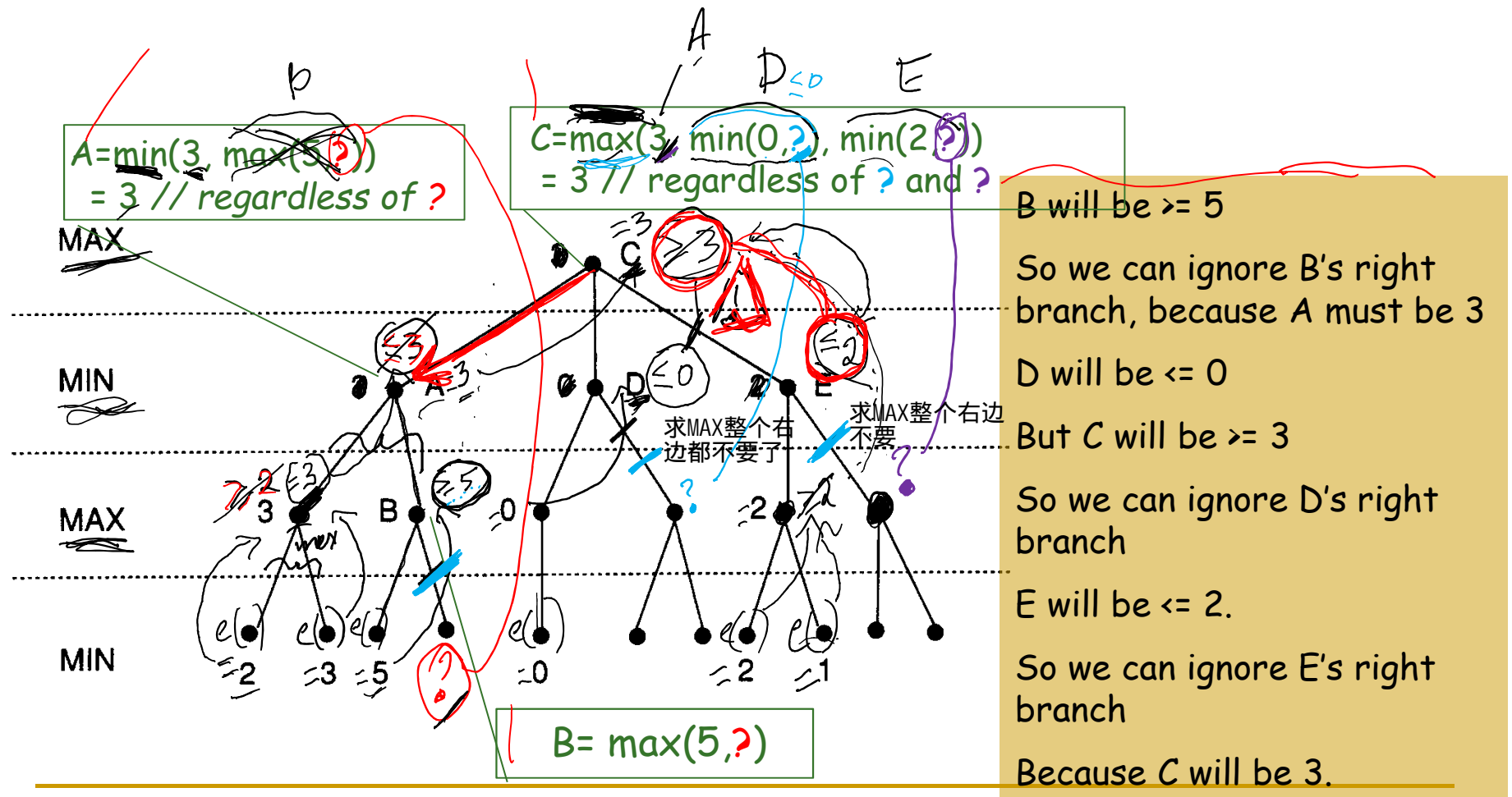
3. to identify leaf nodes (or entire branches) on which running $e(n)$ is pointless



用Depth First Traversal ,
所以第一个点是红色的1, 算出来是2,
Propagate给parent, parent因为是MAX,
constraint就是 ≥ 2 .
Depth first traversal所以下一个点是
2, e等于3, parent求 $\max=3$, propagate到
A点, 求MIN, 所以小于等于3, 然后点3,
发现是5, B至少大于等于5, 这个分支没
用, 舍去, 所以下一个点不用了

$$\begin{array}{l} 11 \times e(n) \\ 6 \times e(n) \end{array}$$

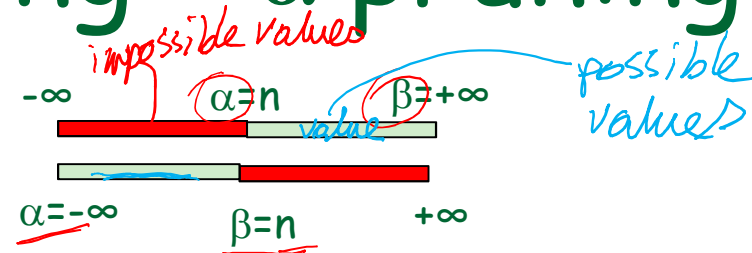
Alpha-Beta Pruning: Example 1



source: G. Luger (2005)

Alpha-Beta Pruning - α pruning

- 就是前面的大于等于, 小于等于
- α = minimum possible value $\text{value} \geq n$
- β = maximum possible value $\text{value} \leq n$

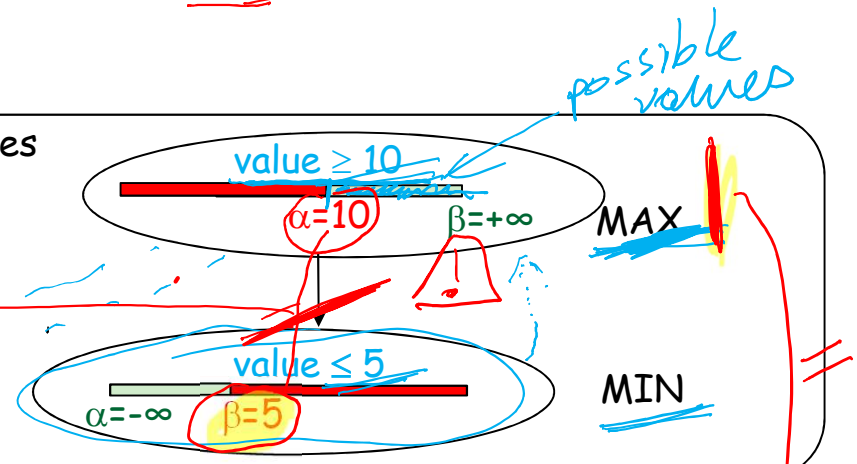


- Alpha pruning - parent is a MAX node

1. if MAX node's $\alpha = n$, then we can prune branches from a MIN descendant that has a $\beta \leq n$.

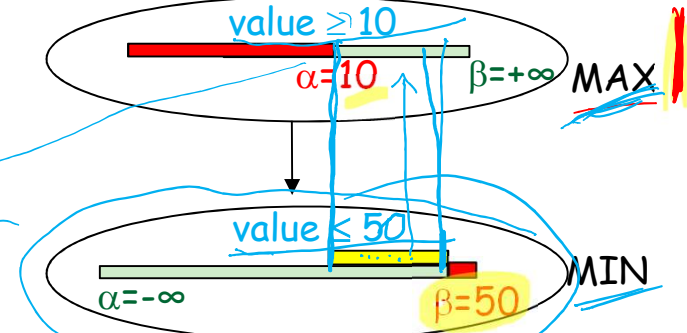
if $\beta_{\text{child}} \leq \alpha_{\text{ancestor}} \rightarrow \text{prune}$

incompatible...
so stop searching this branch;
the value cannot come from there!



2. if $\beta_{\text{child}} > \alpha_{\text{ancestor}} \rightarrow \text{cannot prune}$

compatible...
we need to search this branch;
the value could come from there!



MAX

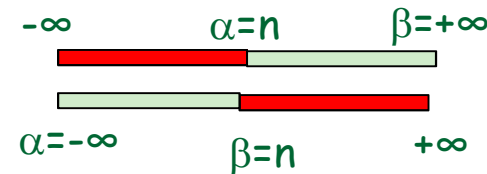
它本身: α 大于等于 n
要求 child: β 小于等于 n
这时可以舍去

MIN

它本身: β 小于等于 n
要求 child MAX: α 大于等于 n ,
这时可以舍去
如果 MAX 小于等于 n , 有希望, 我们还要算 n

Alpha-Beta Pruning - β pruning

- α = minimum possible value $\text{value} \geq n$
- β = maximum possible value $\text{value} \leq n$

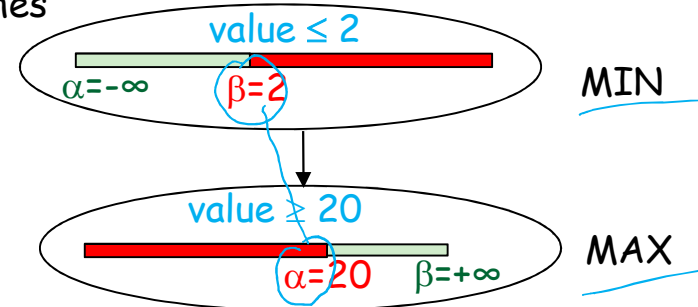


same

- Beta pruning - parent is a MIN node

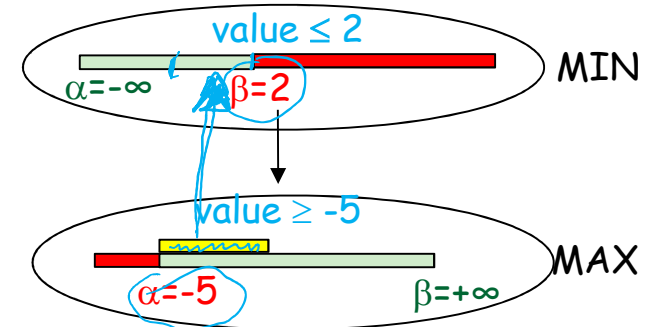
1. if a MIN node's $\beta = n$, then we can prune branches from a MAX descendant that has an $\alpha \geq n$.
if $(\alpha_{\text{child}} \geq \beta_{\text{ancestor}}) \rightarrow \text{prune}$

incompatible...
so stop searching this branch;
the value cannot come from there!



2. if $(\alpha_{\text{child}} < \beta_{\text{ancestor}}) \rightarrow \text{cannot prune}$

compatible...
we need to search this branch;
the value could come from there!



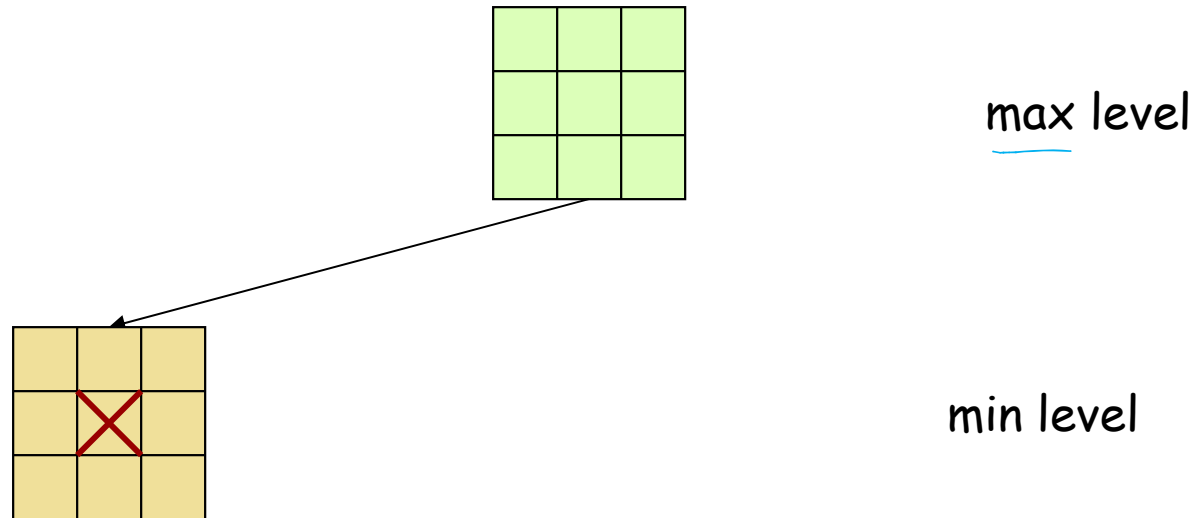
Alpha-Beta Pruning Algorithm

```
01 function alphabeta(node, depth,  $\alpha$ ,  $\beta$ , maximizingPlayer)
02     if depth = 0 or node is a terminal node
03         return the heuristic value of node
04     if maximizingPlayer
05         v :=  $-\infty$ 
06         for each child of node
07             v := max(v, alphabeta(child, depth - 1,  $\alpha$ ,  $\beta$ , FALSE))
08              $\alpha$  := max( $\alpha$ , v)
09             if  $\beta \leq \alpha$ 
10                 break (*  $\beta$  cut-off *)
11         return v
12     else
13         v :=  $\infty$ 
14         for each child of node
15             v := min(v, alphabeta(child, depth - 1,  $\alpha$ ,  $\beta$ , TRUE))
16              $\beta$  := min( $\beta$ , v)
17             if  $\beta \leq \alpha$ 
18                 break (*  $\alpha$  cut-off *)
19         return v
```

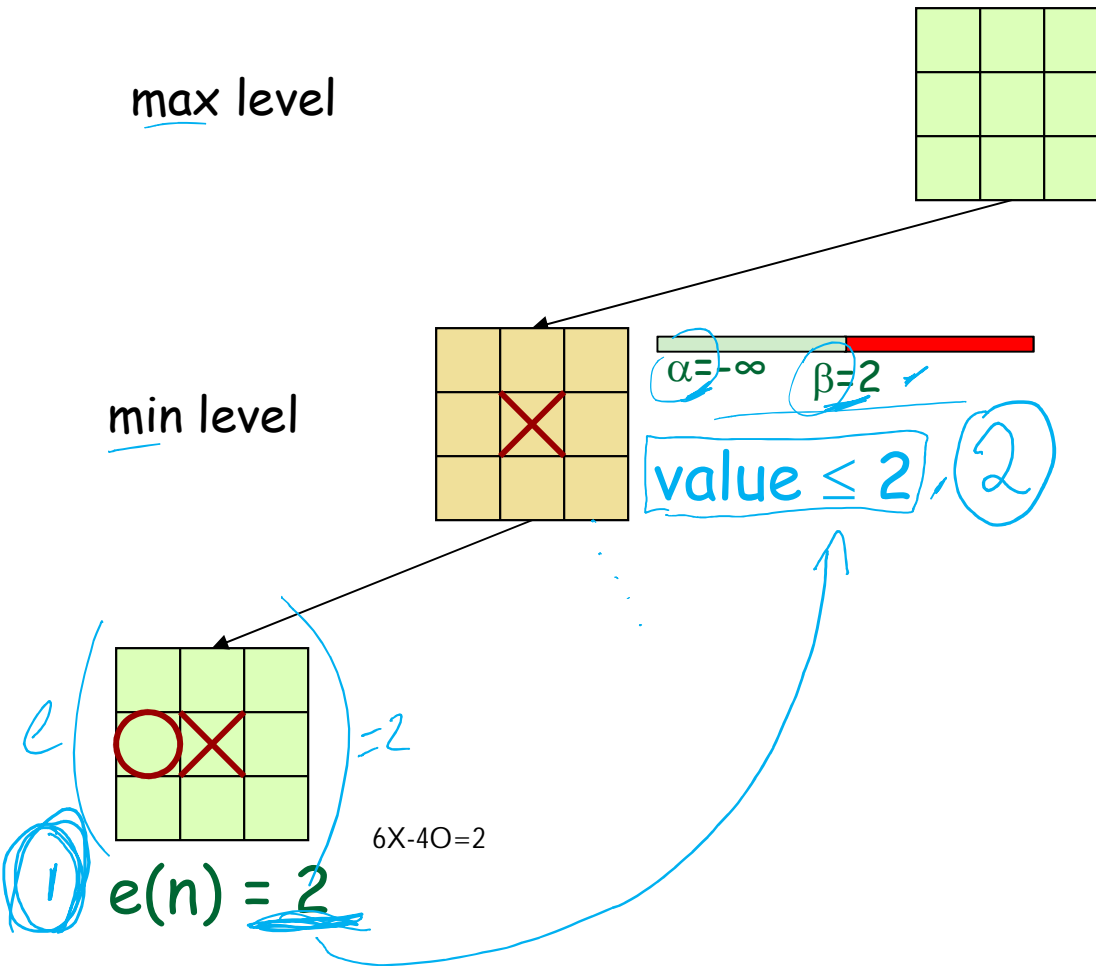
Initial call:

alphabeta(origin, depth, $-\infty$, $+\infty$, TRUE)

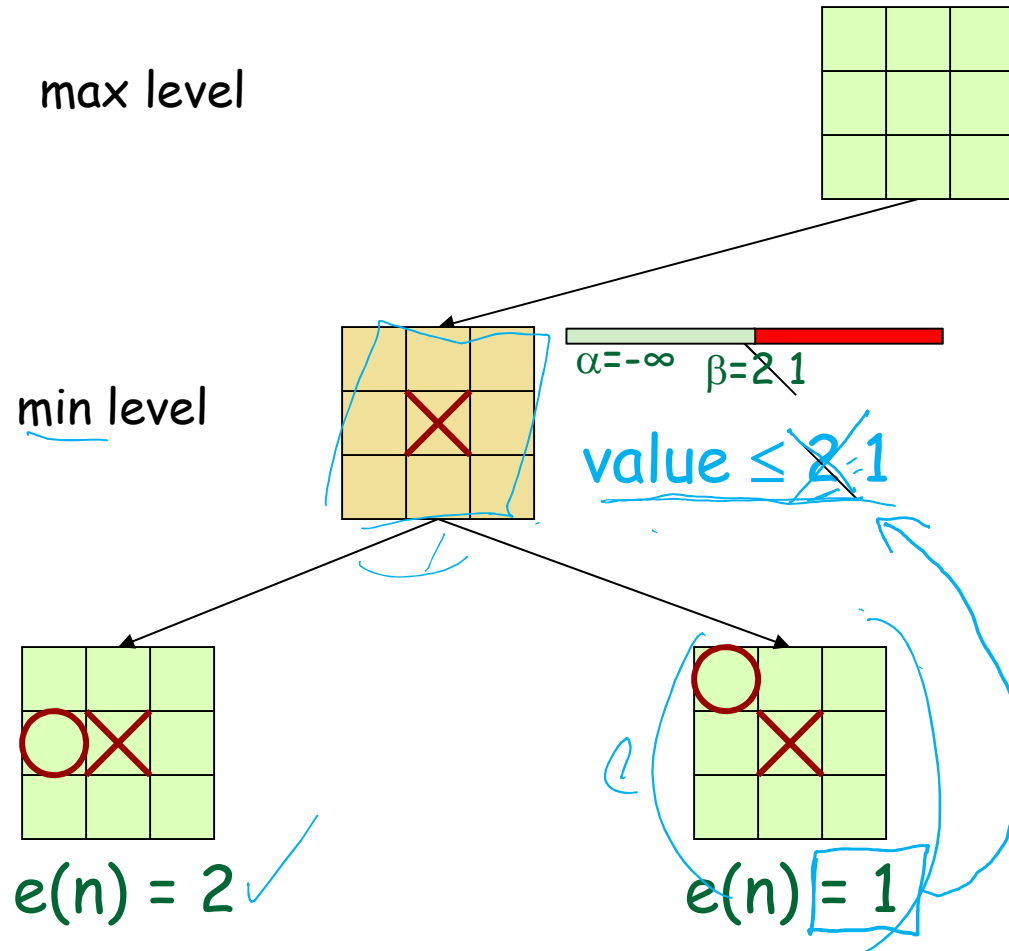
Example with tic-tac-toe



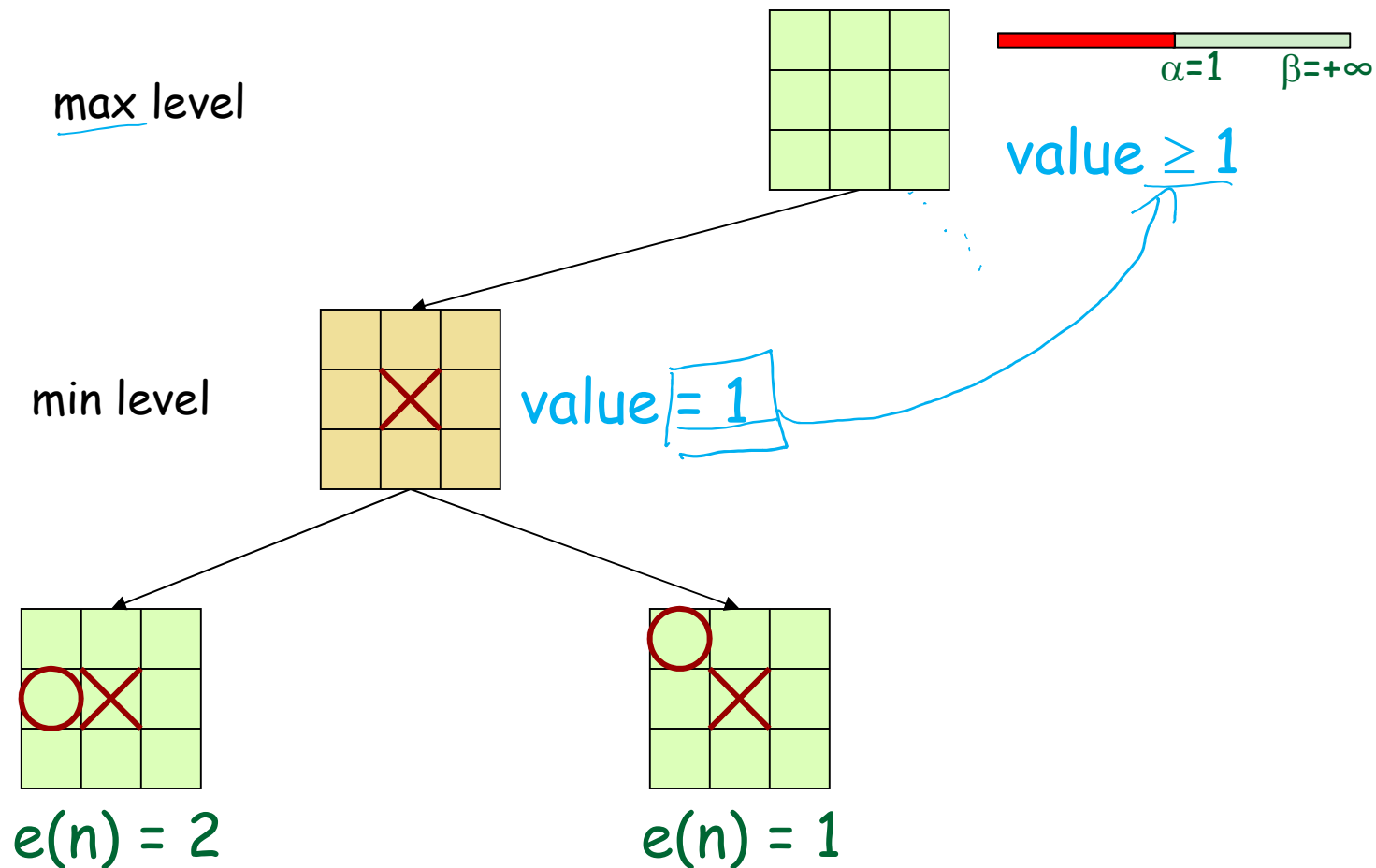
Example with tic-tac-toe



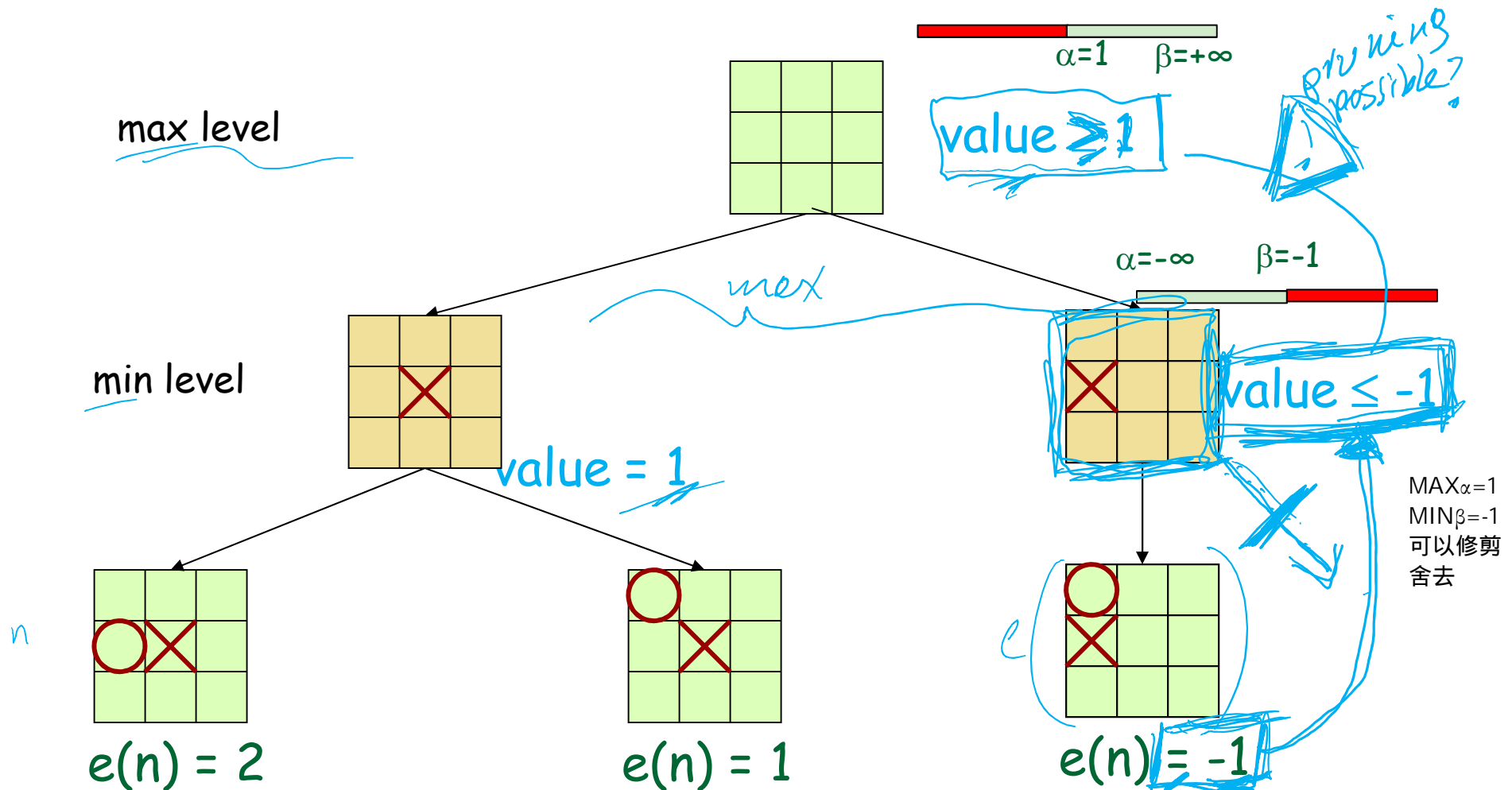
Example with tic-tac-toe



Example with tic-tac-toe

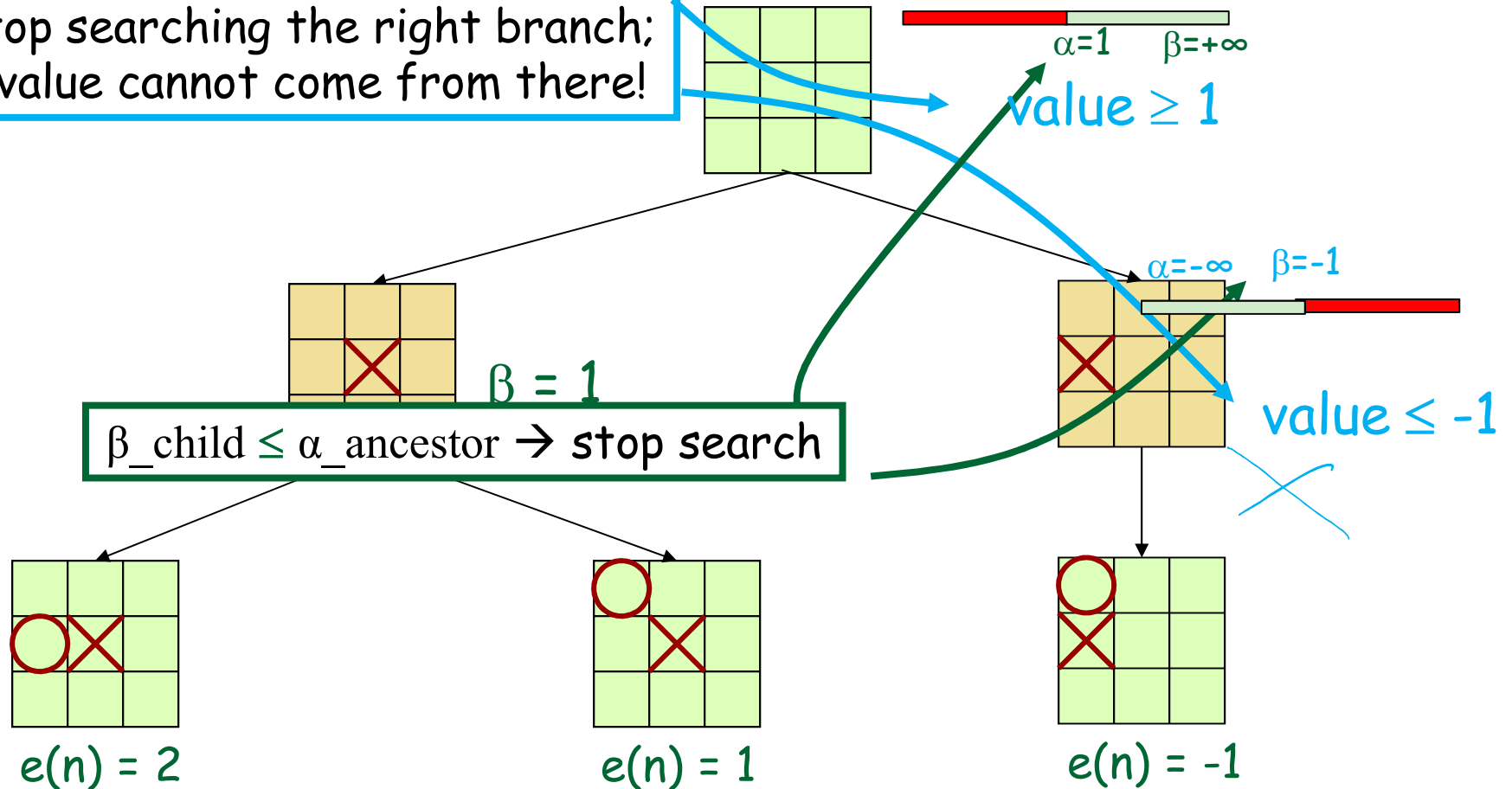


Example with tic-tac-toe

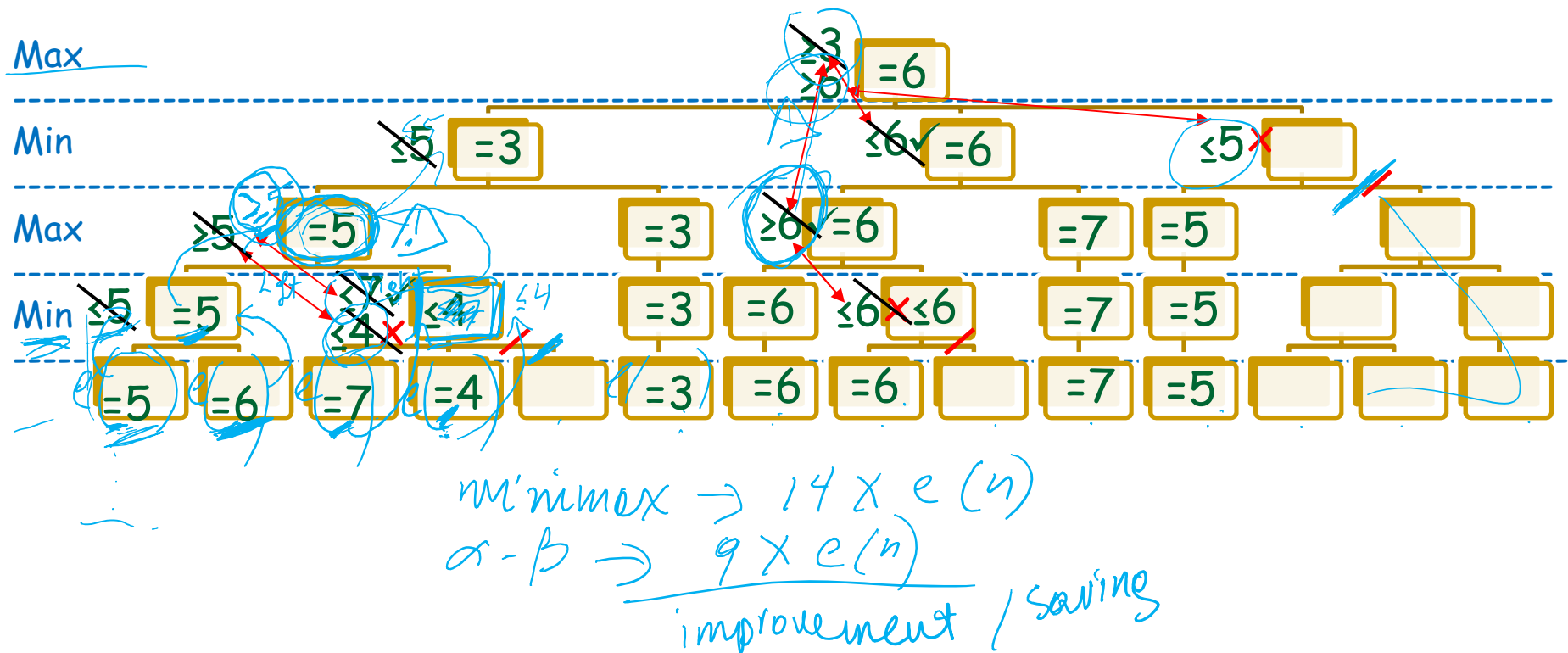


Example with tic-tac-toe

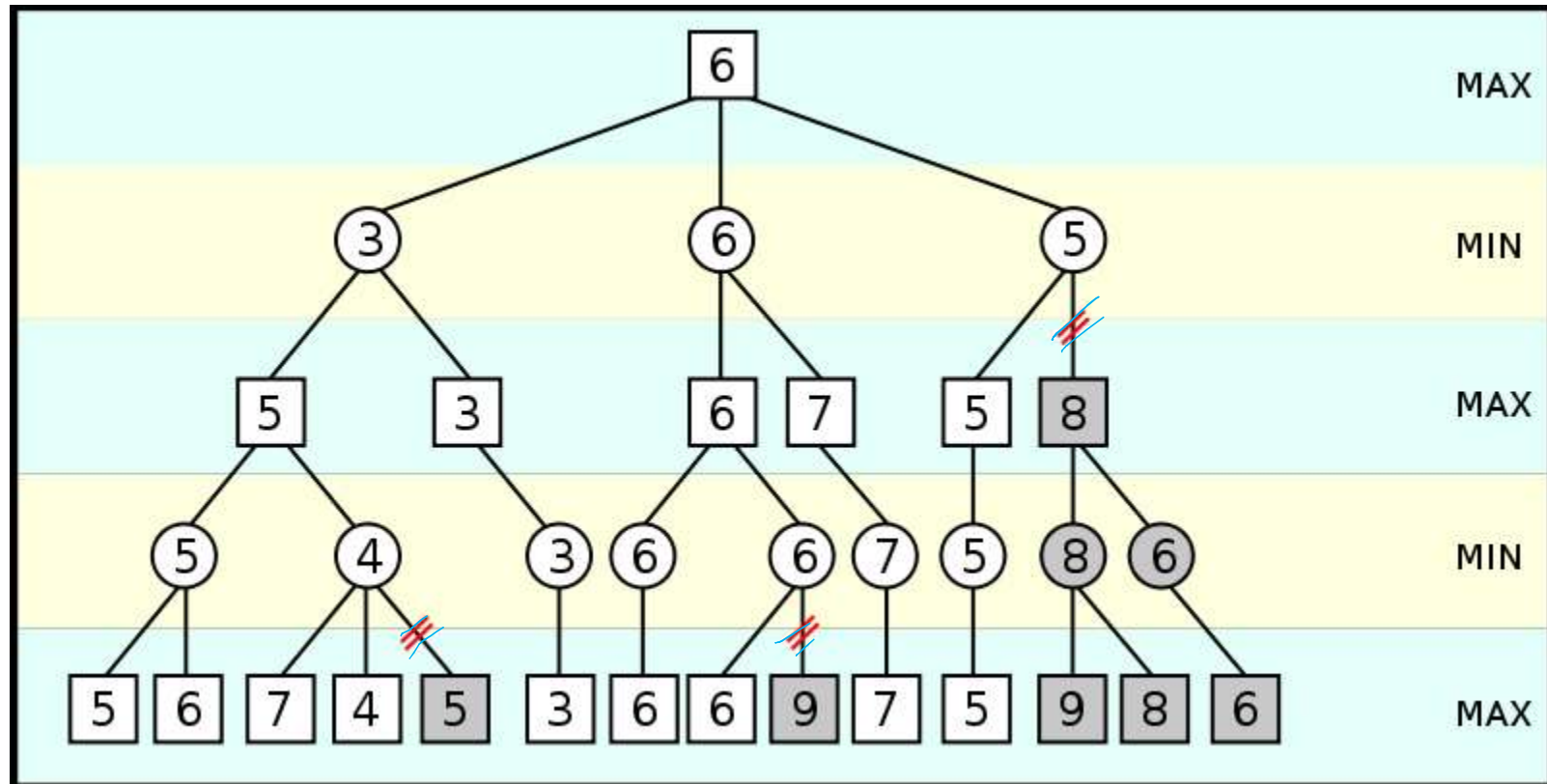
incompatible...
so stop searching the right branch;
the value cannot come from there!



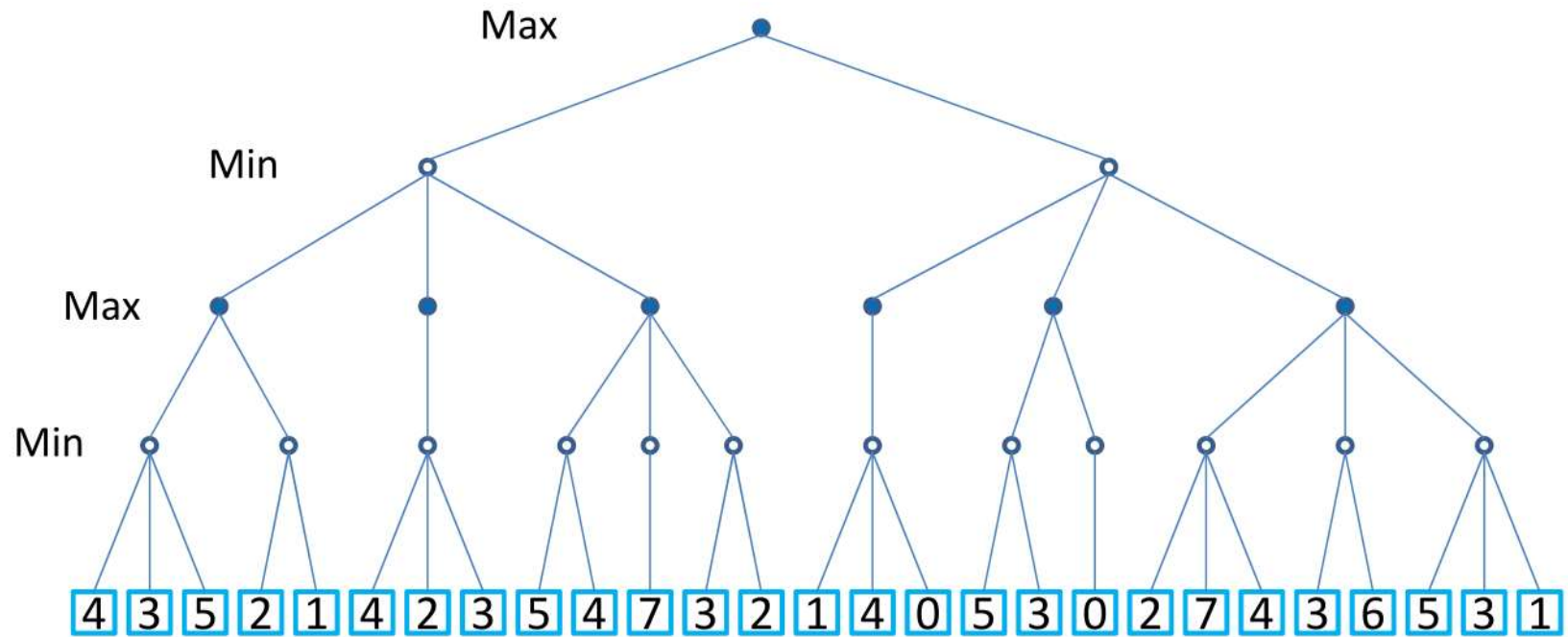
Alpha-Beta Pruning: Example 2



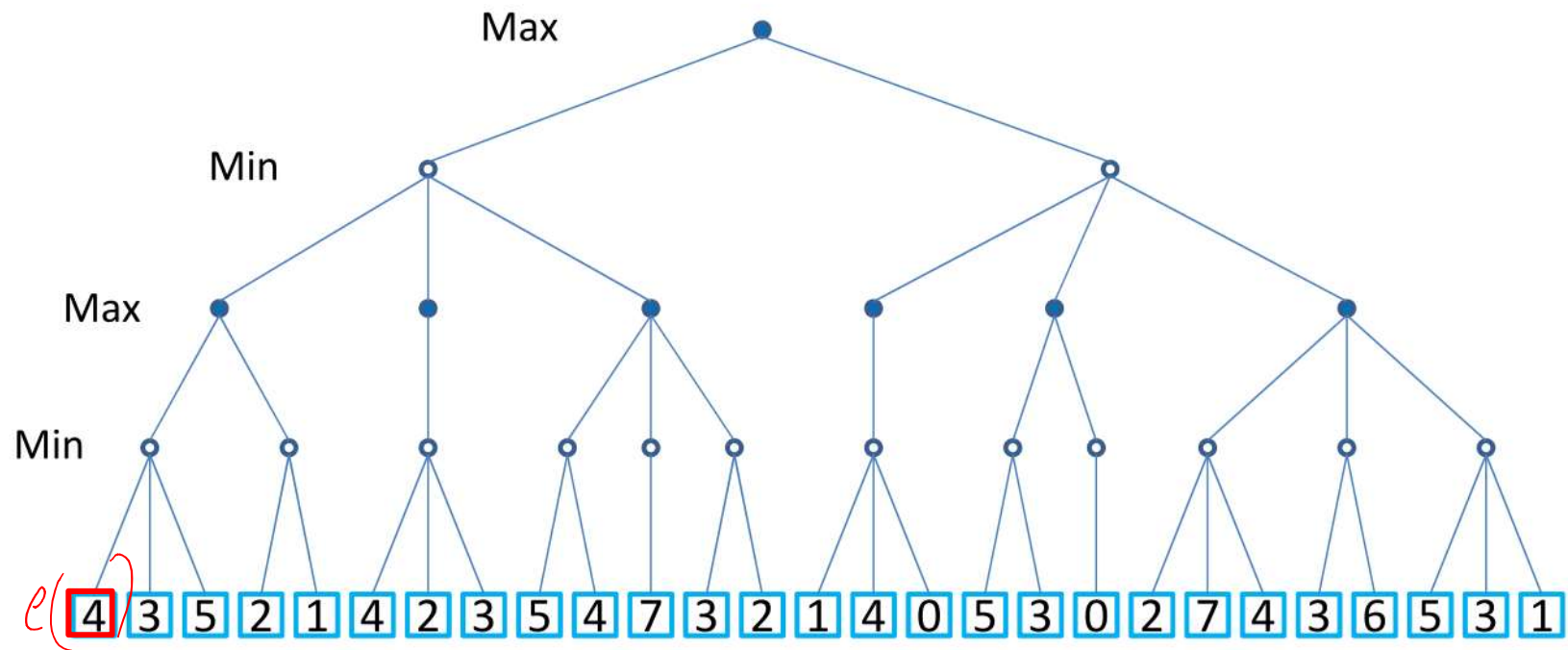
Alpha-Beta Pruning: Example 2



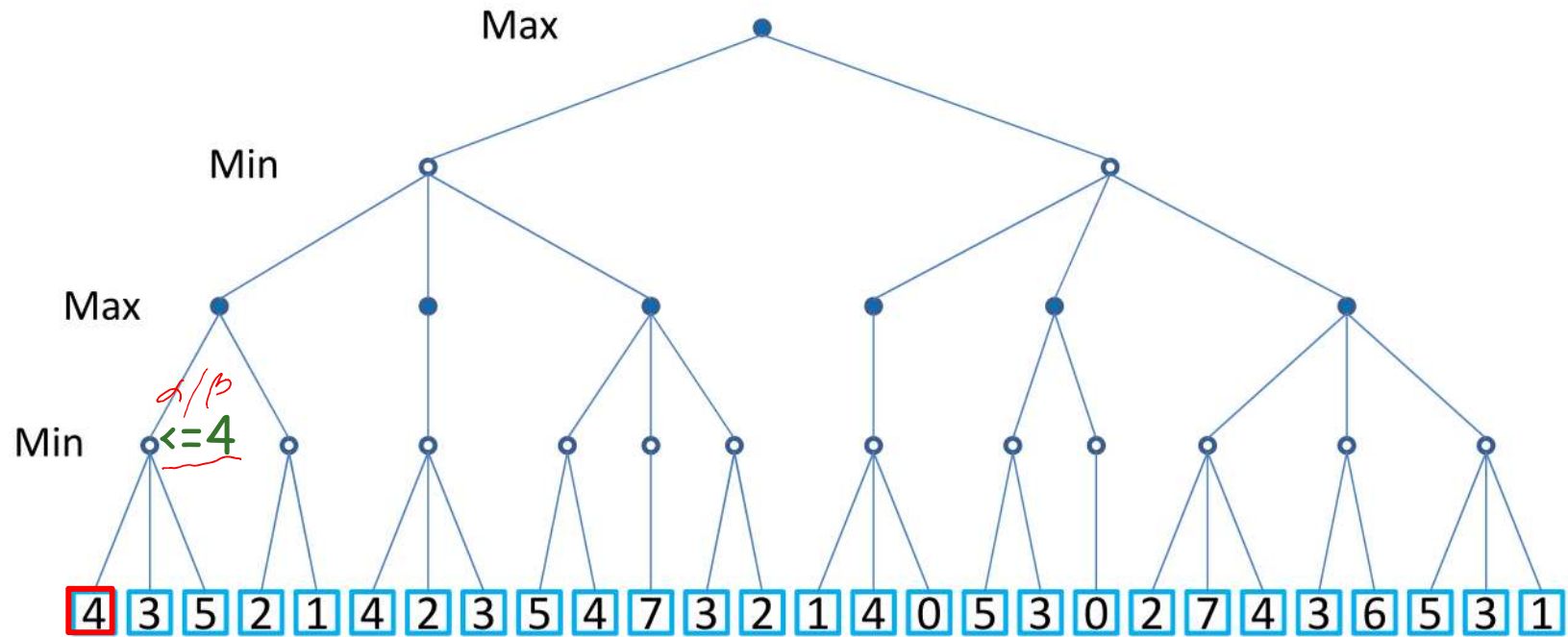
Alpha-Beta Pruning: Example 3



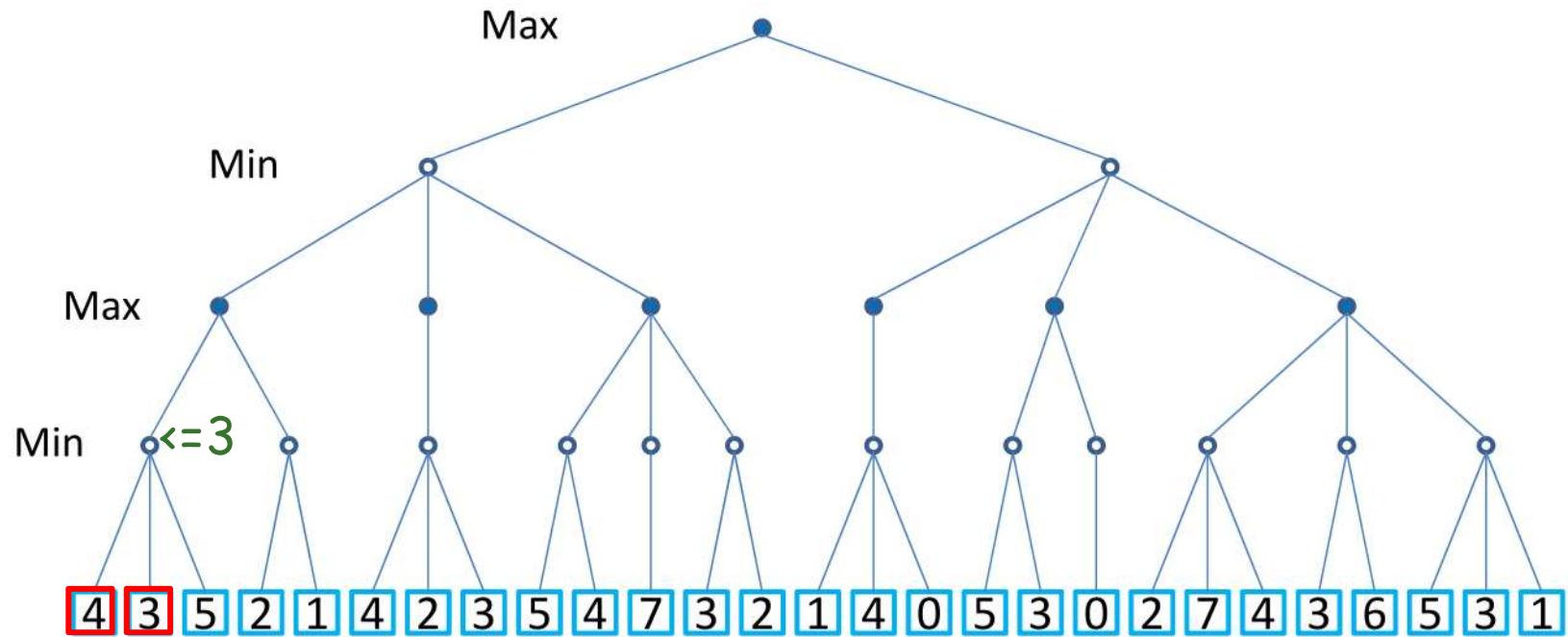
Alpha-Beta Pruning: Example 3



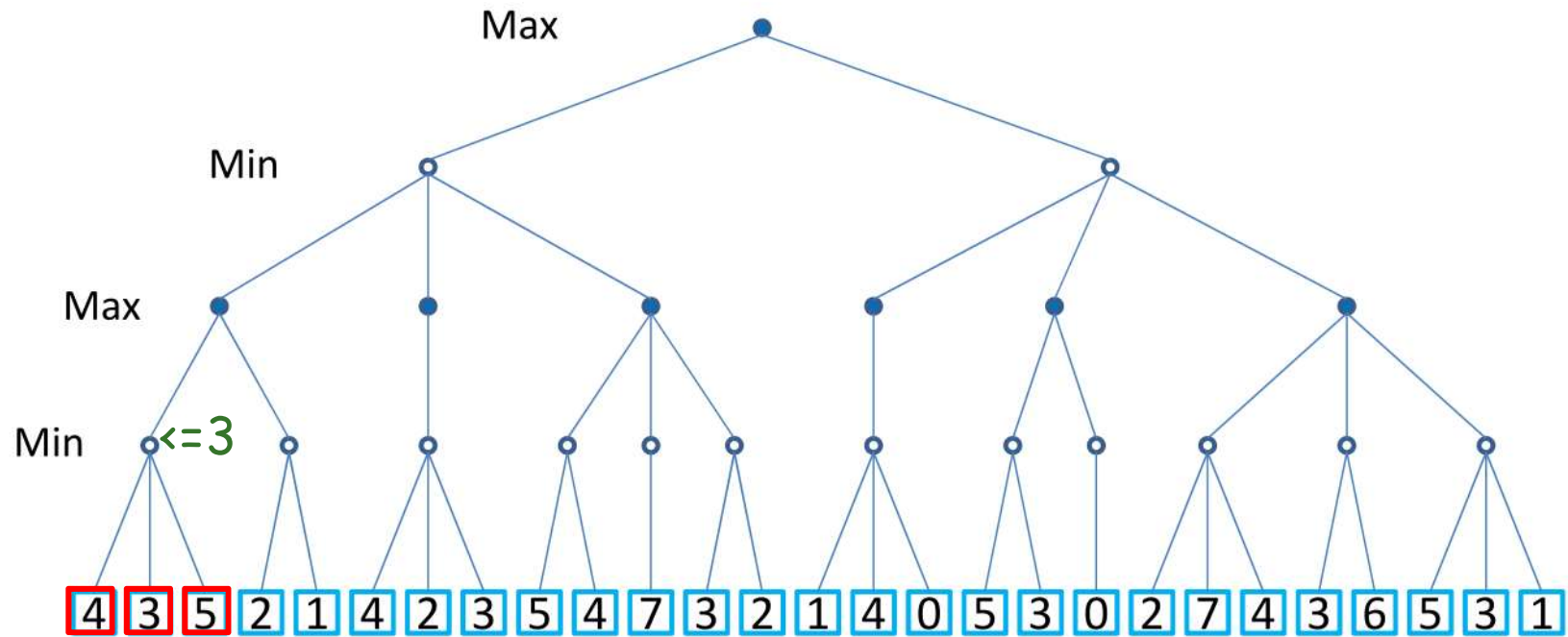
Alpha-Beta Pruning: Example 3



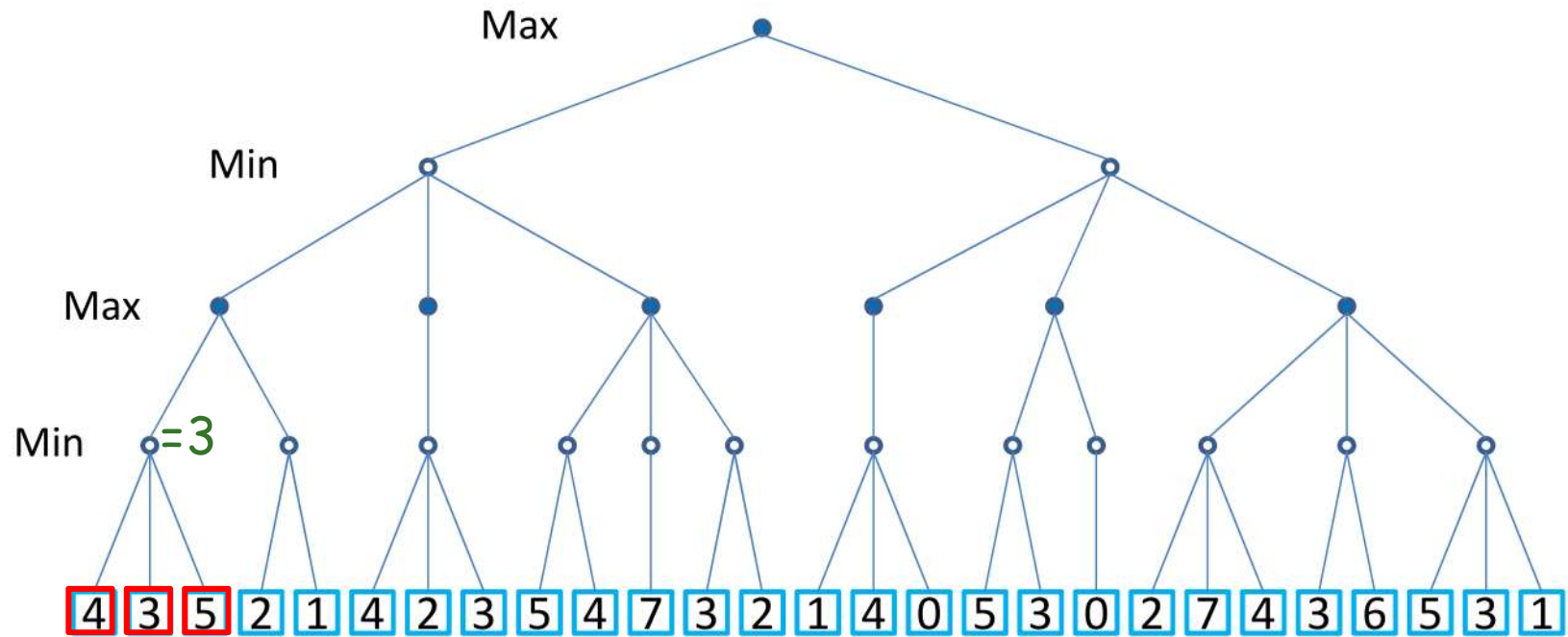
Alpha-Beta Pruning: Example 3



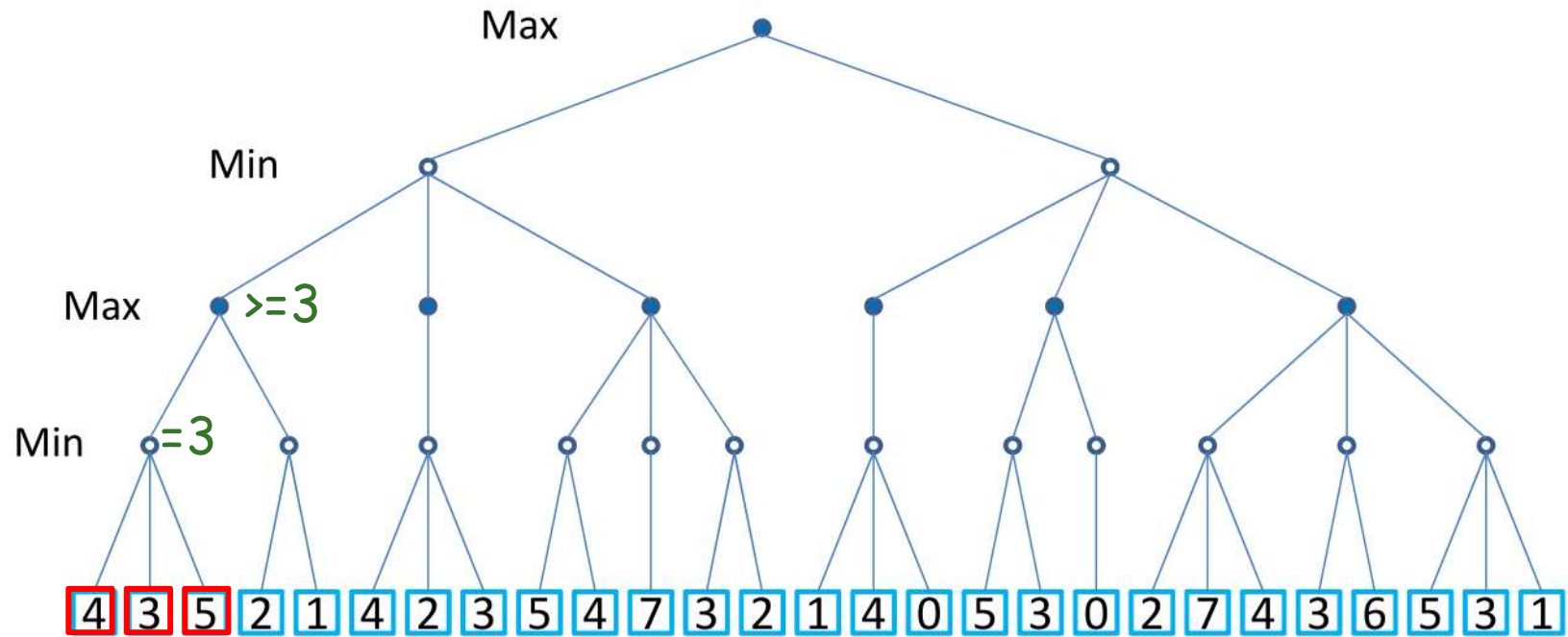
Alpha-Beta Pruning: Example 3



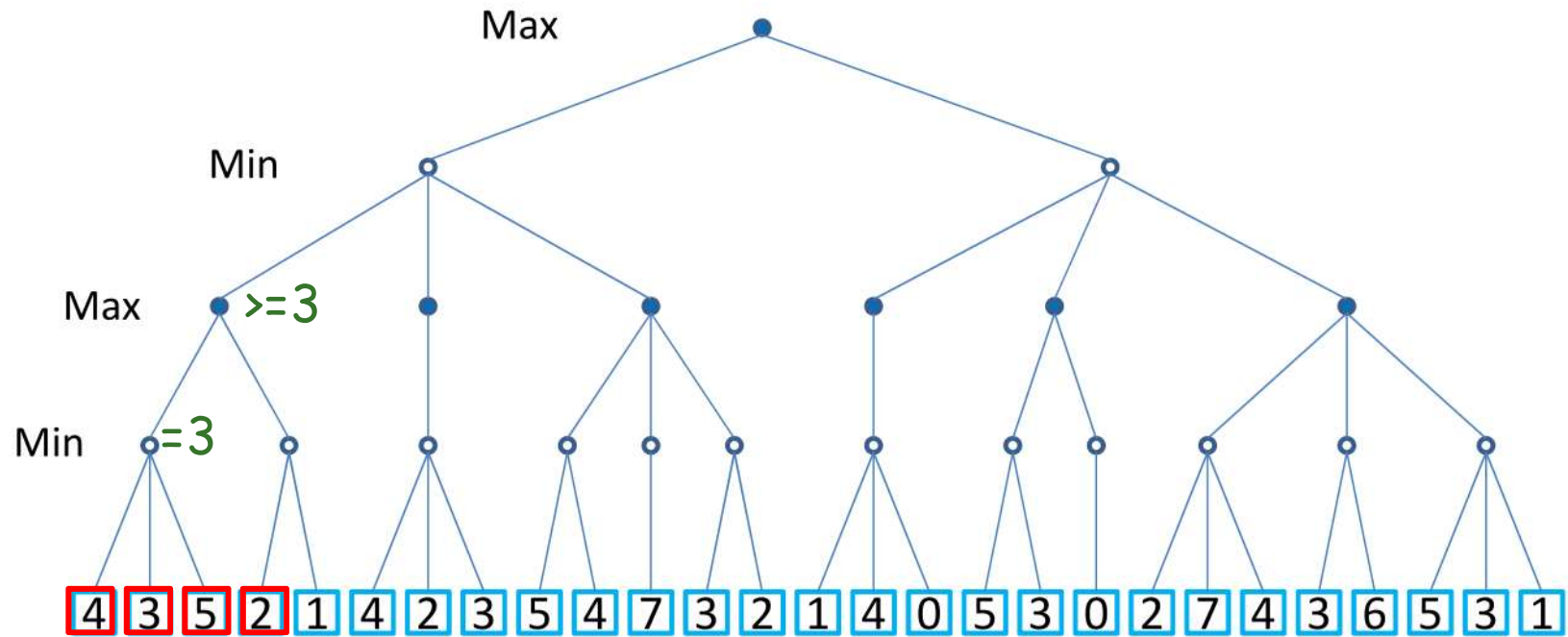
Alpha-Beta Pruning: Example 3



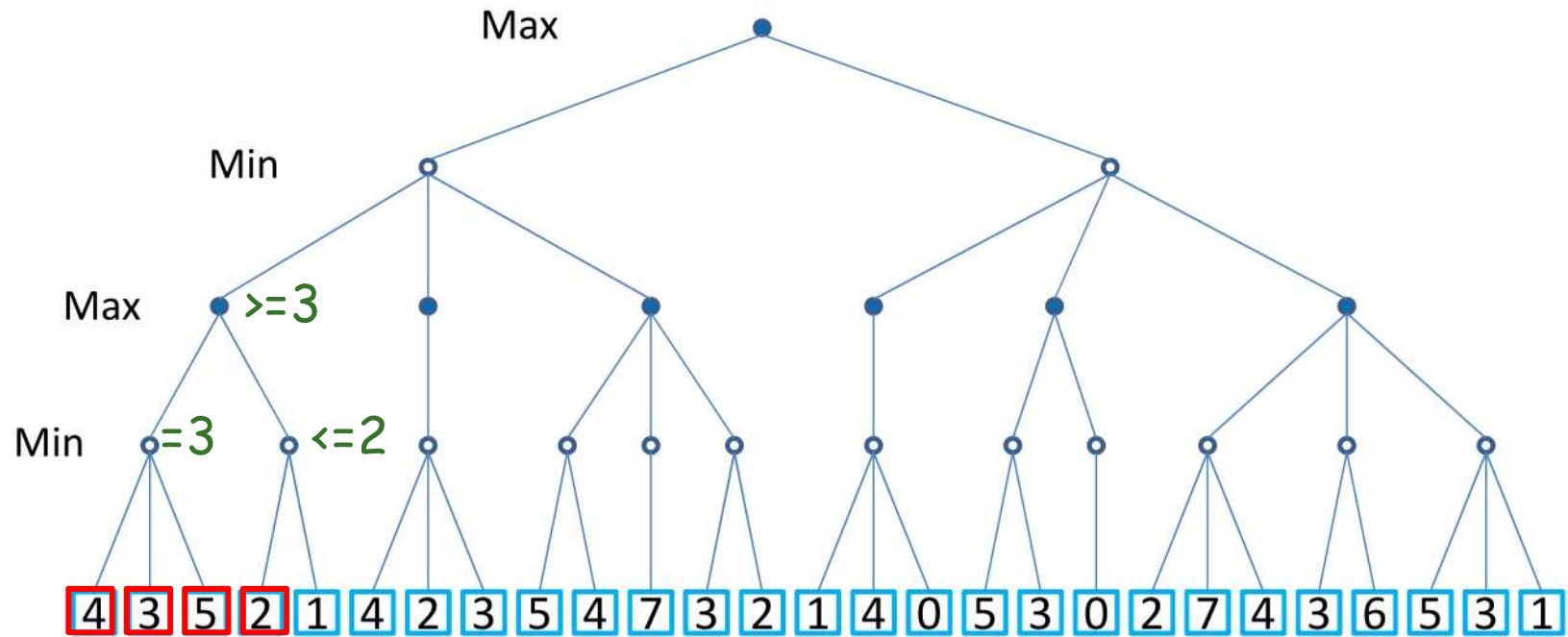
Alpha-Beta Pruning: Example 3



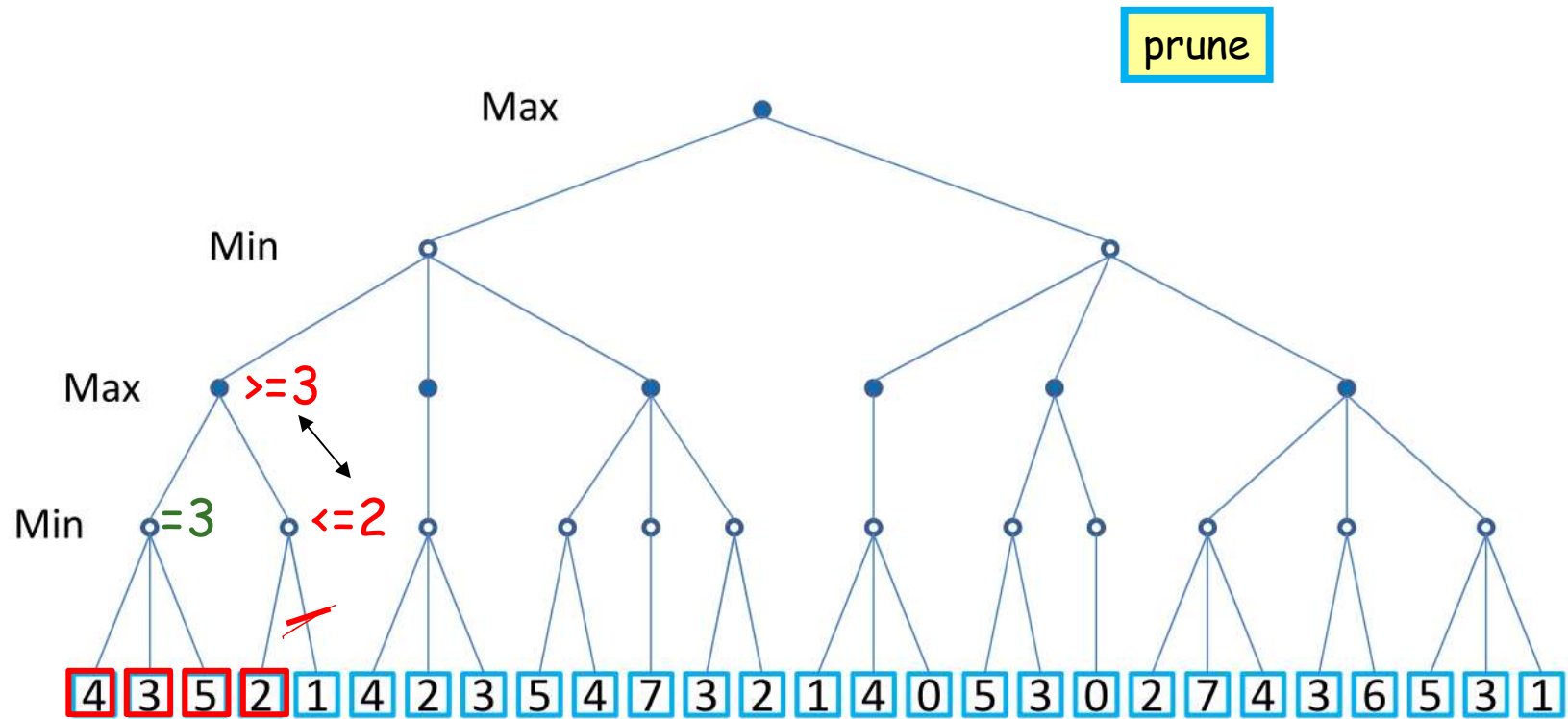
Alpha-Beta Pruning: Example 3



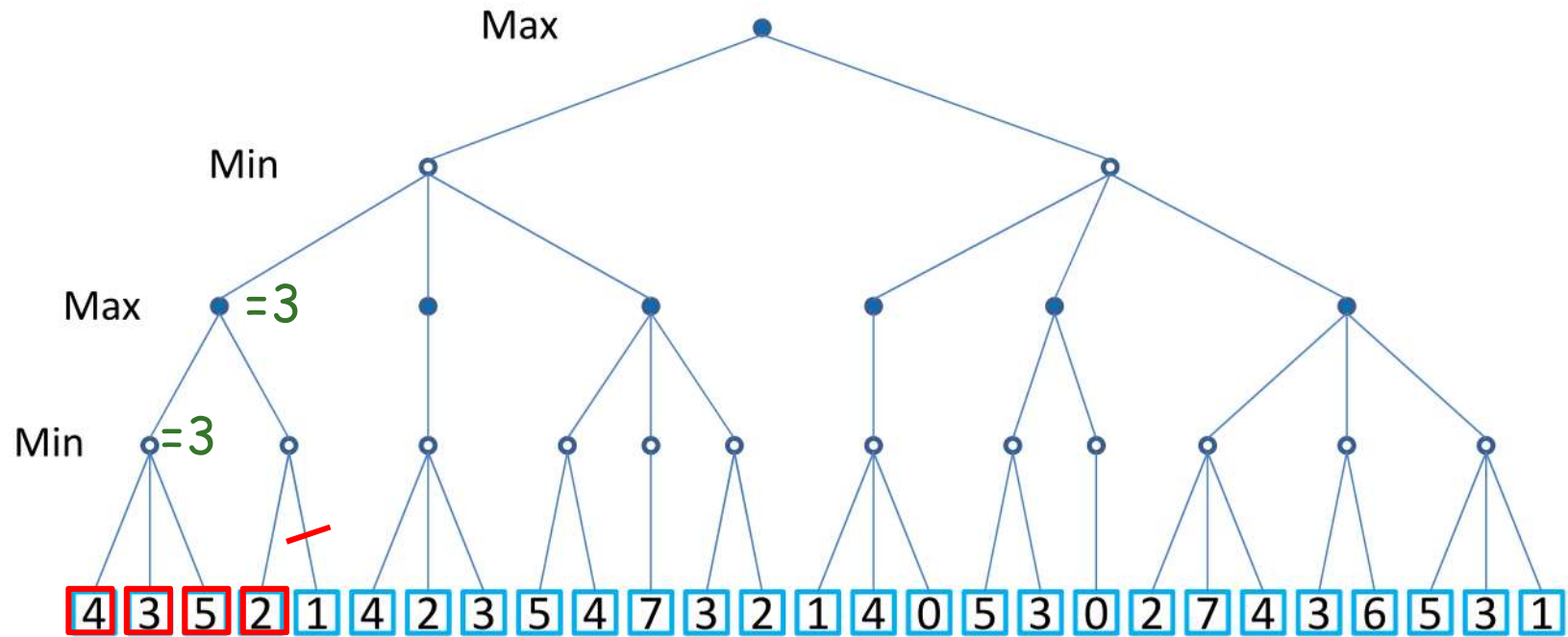
Alpha-Beta Pruning: Example 3



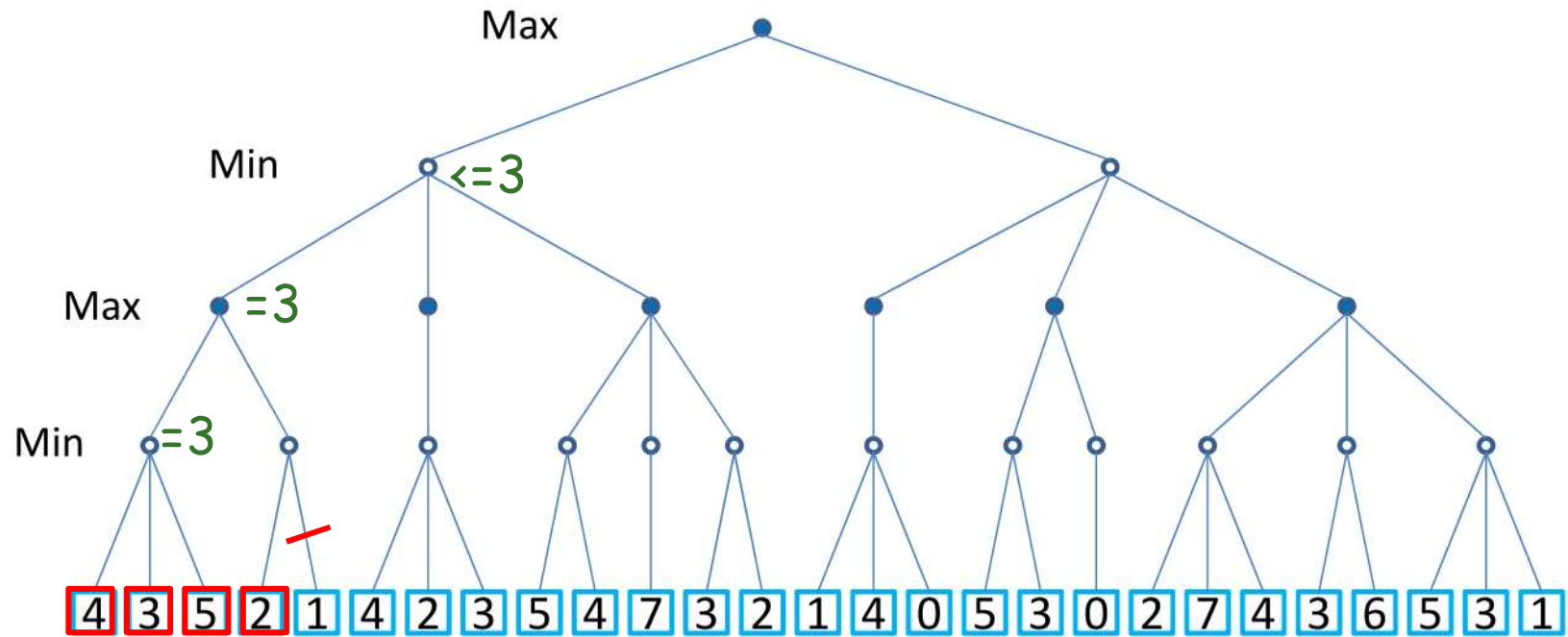
Alpha-Beta Pruning: Example 3



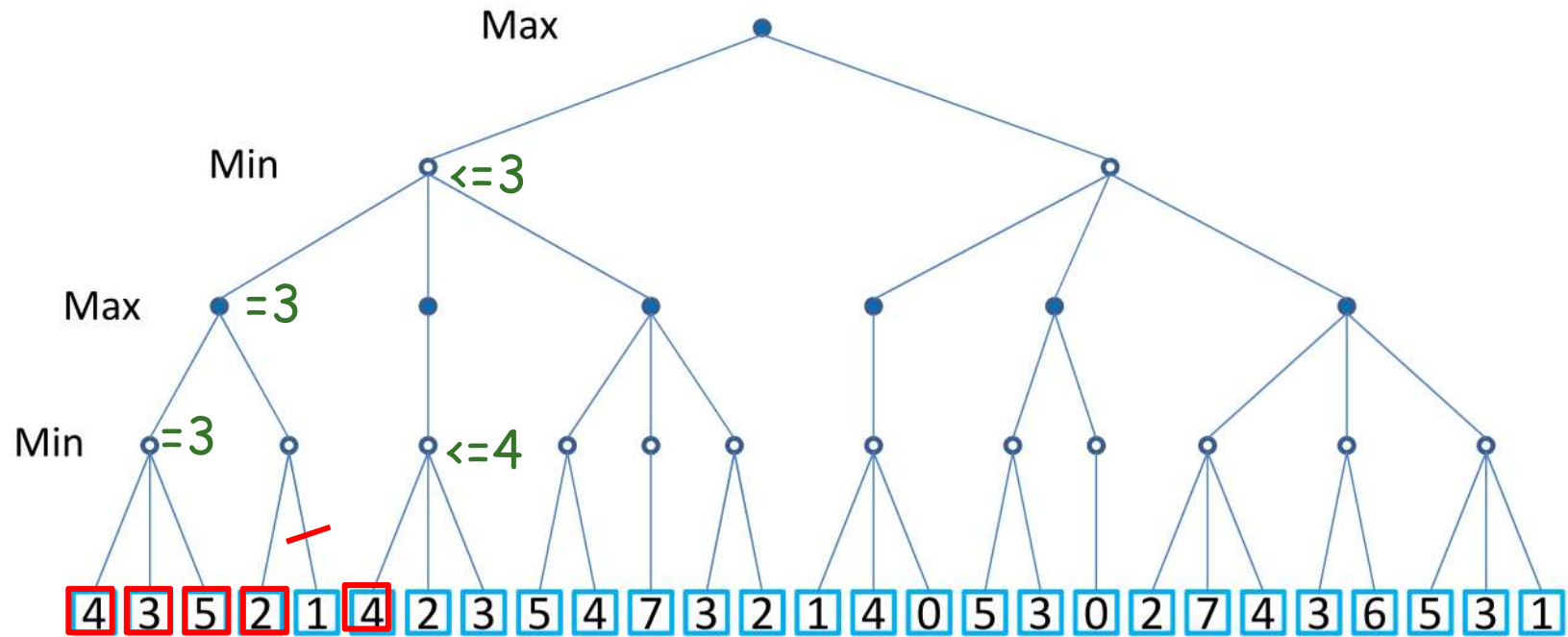
Alpha-Beta Pruning: Example 3



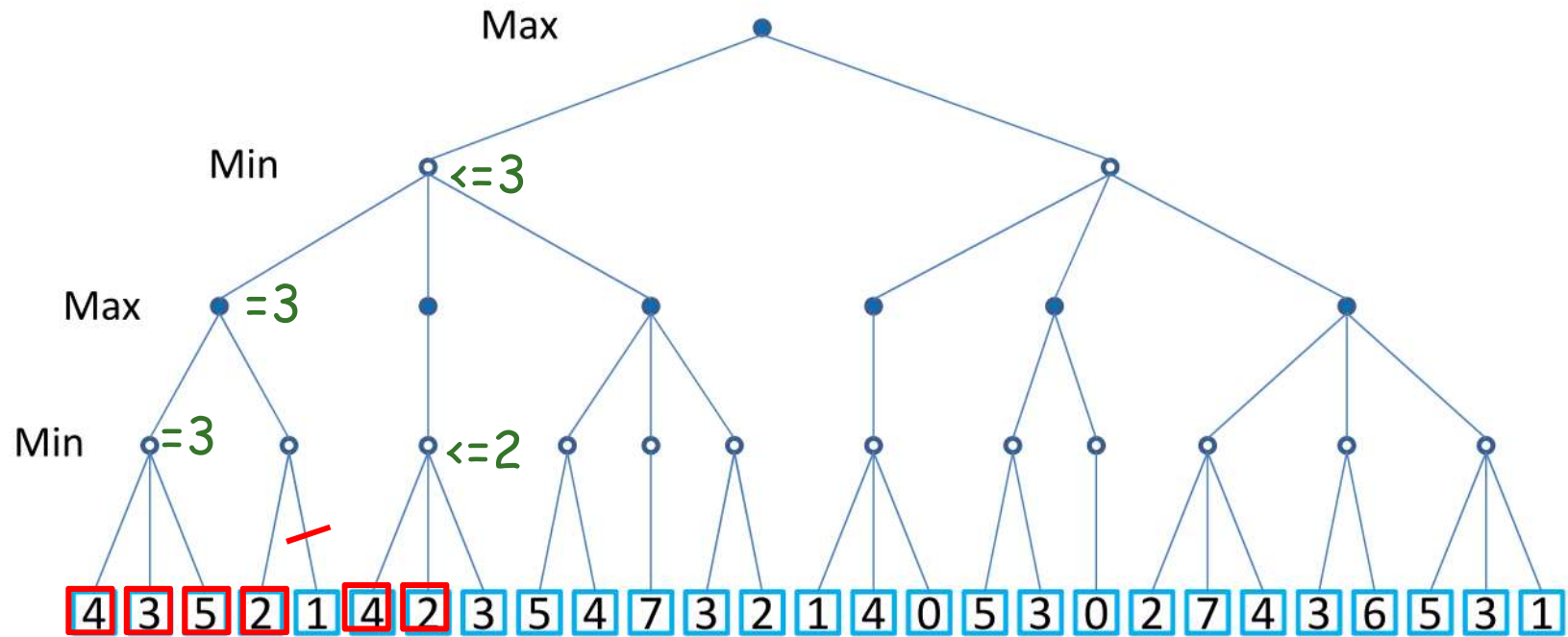
Alpha-Beta Pruning: Example 3



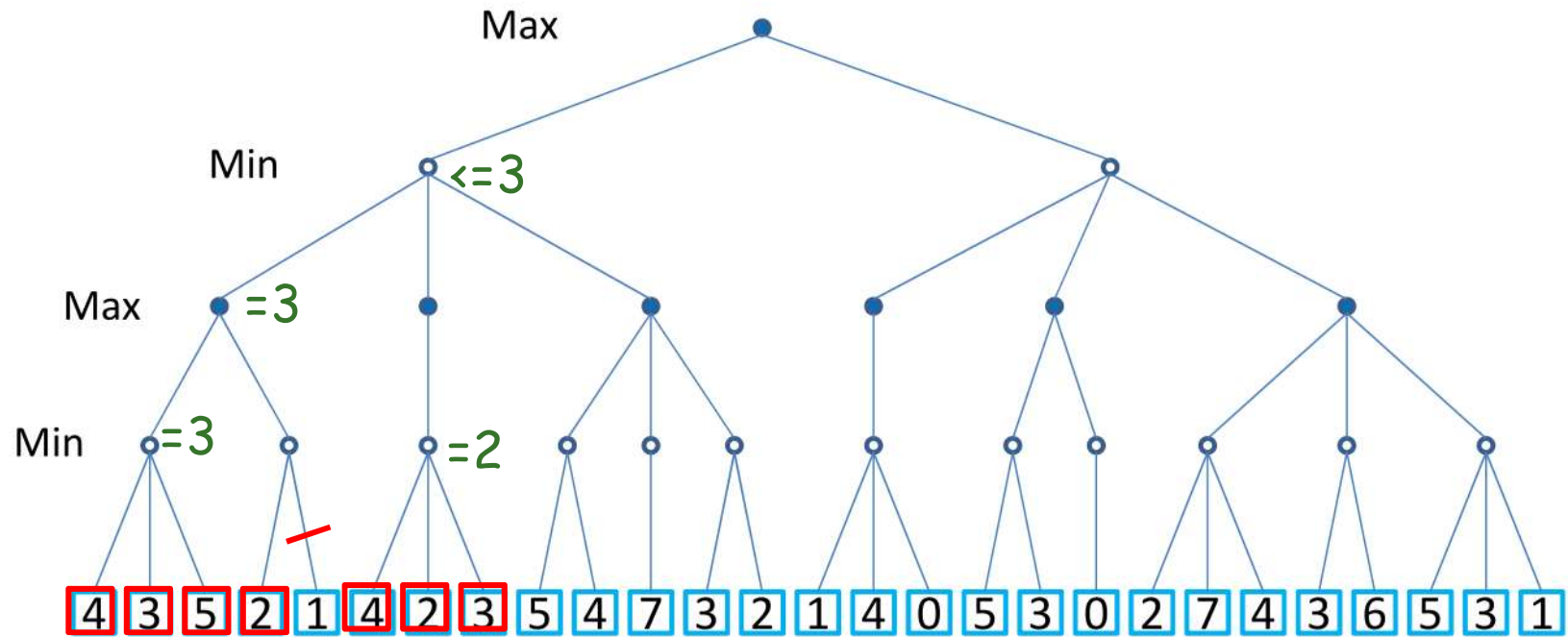
Alpha-Beta Pruning: Example 3



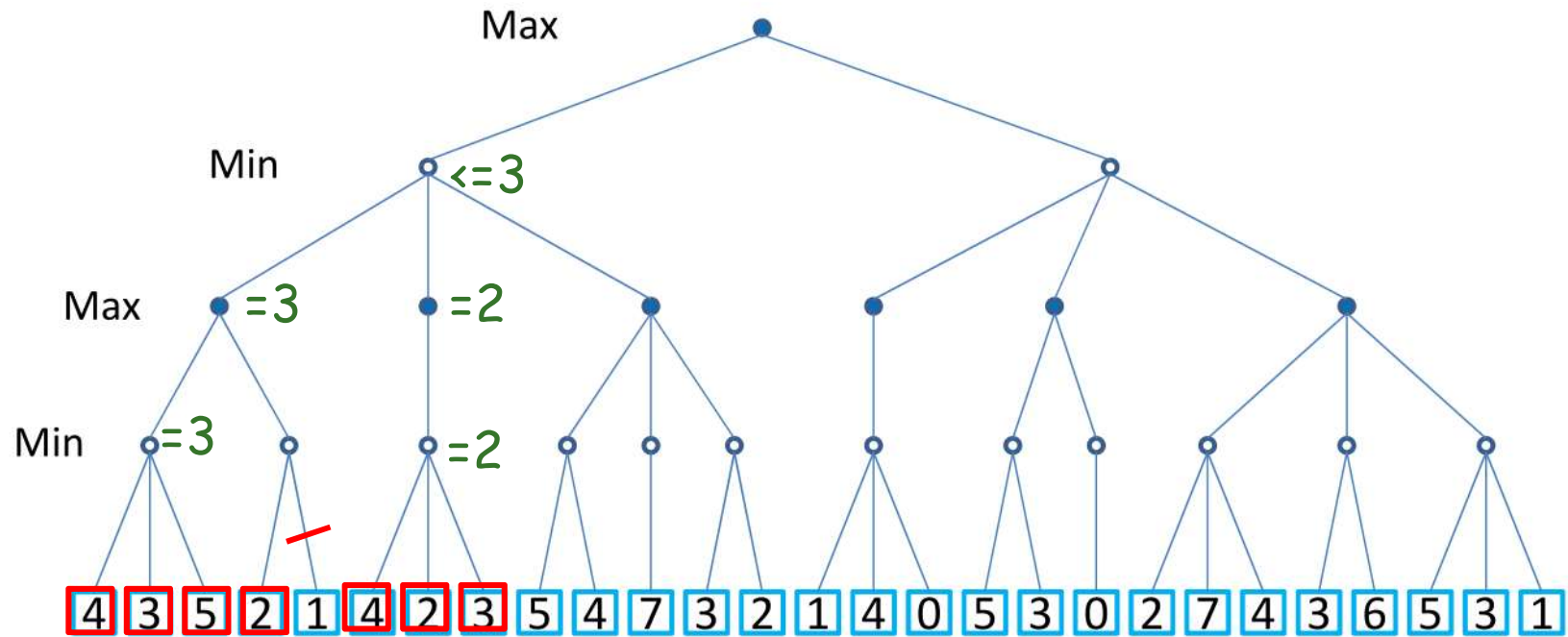
Alpha-Beta Pruning: Example 3



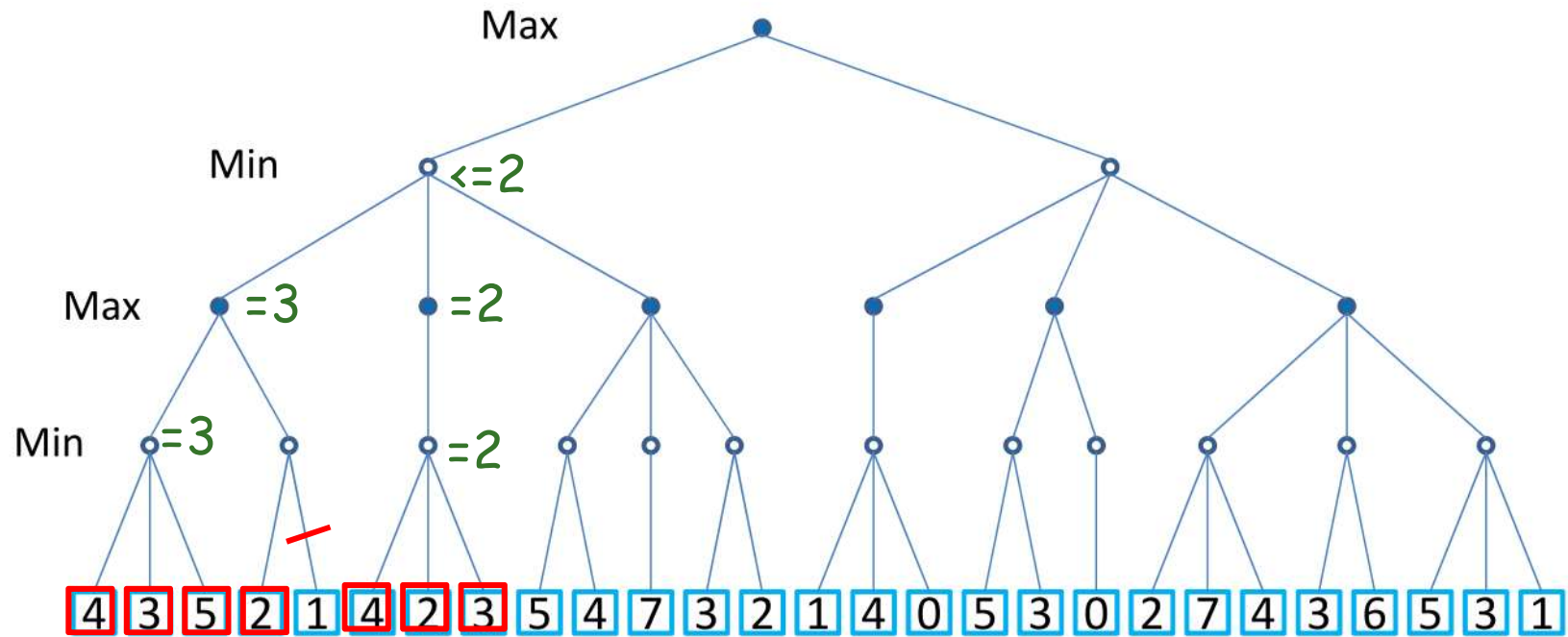
Alpha-Beta Pruning: Example 3



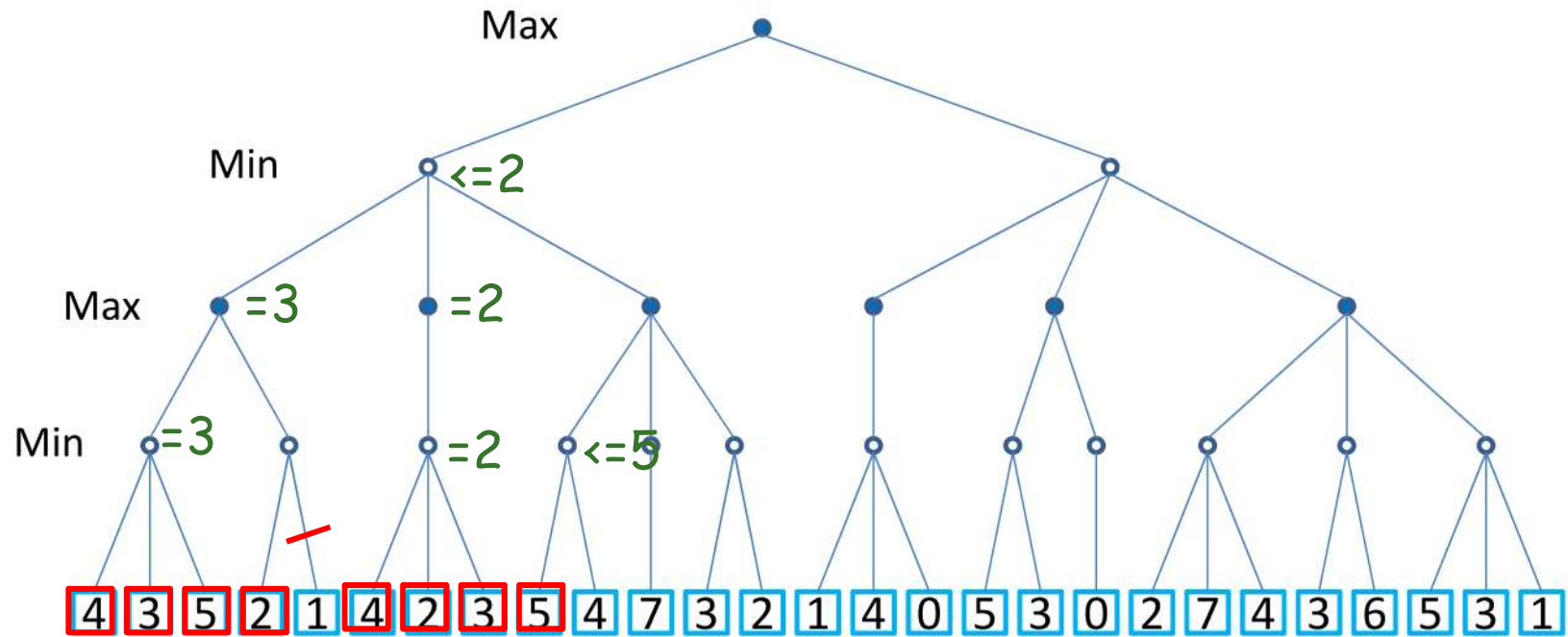
Alpha-Beta Pruning: Example 3



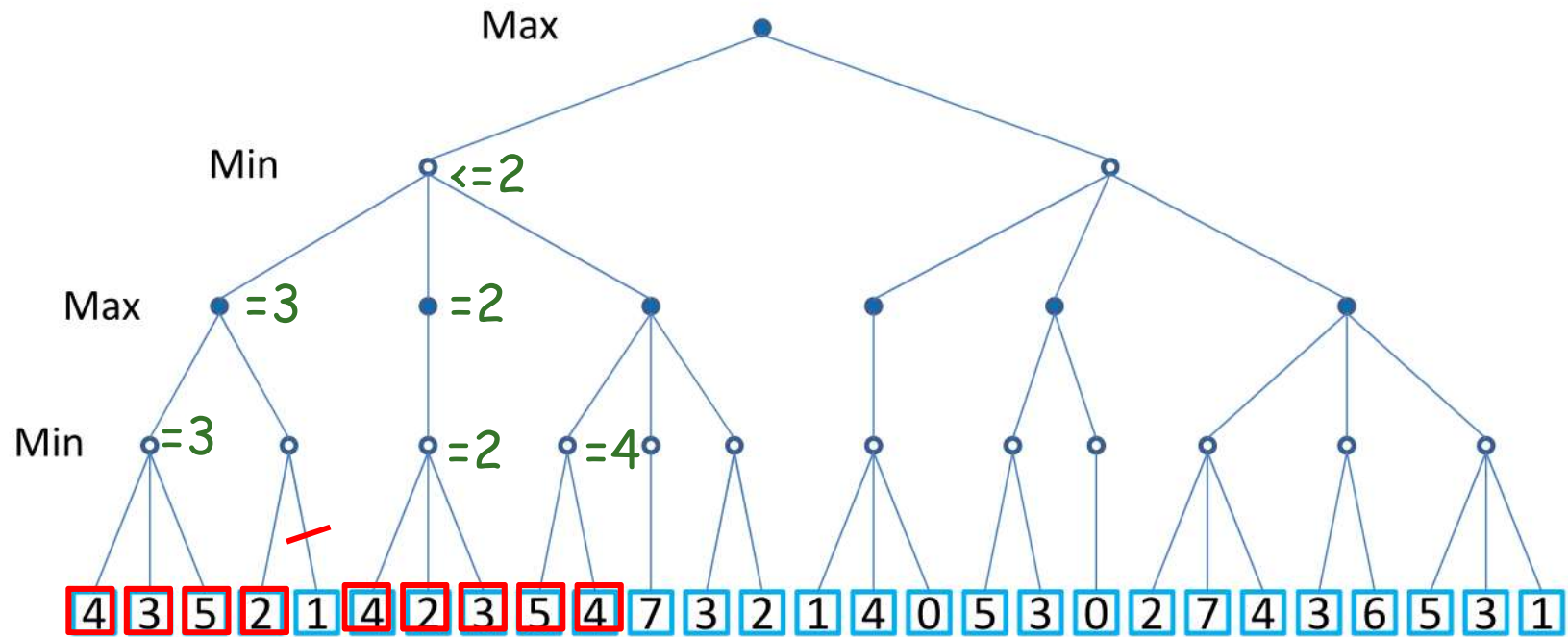
Alpha-Beta Pruning: Example 3



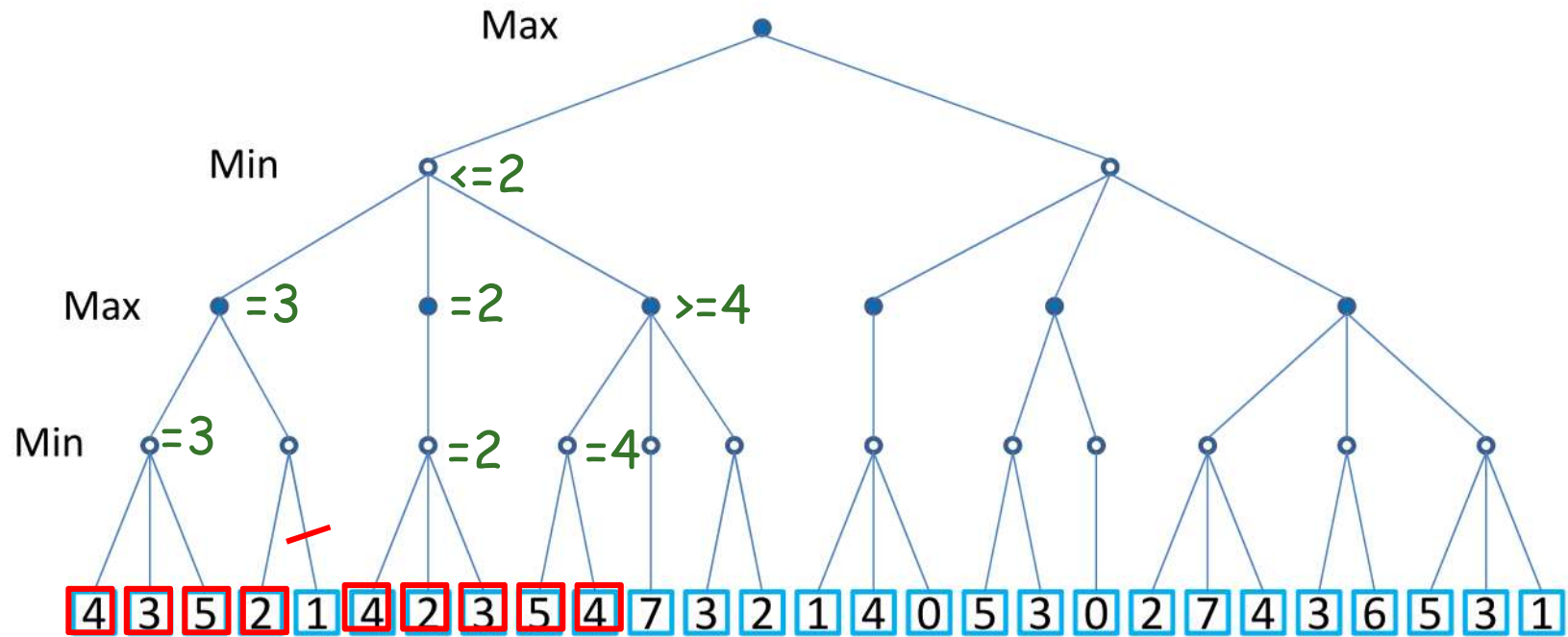
Alpha-Beta Pruning: Example 3



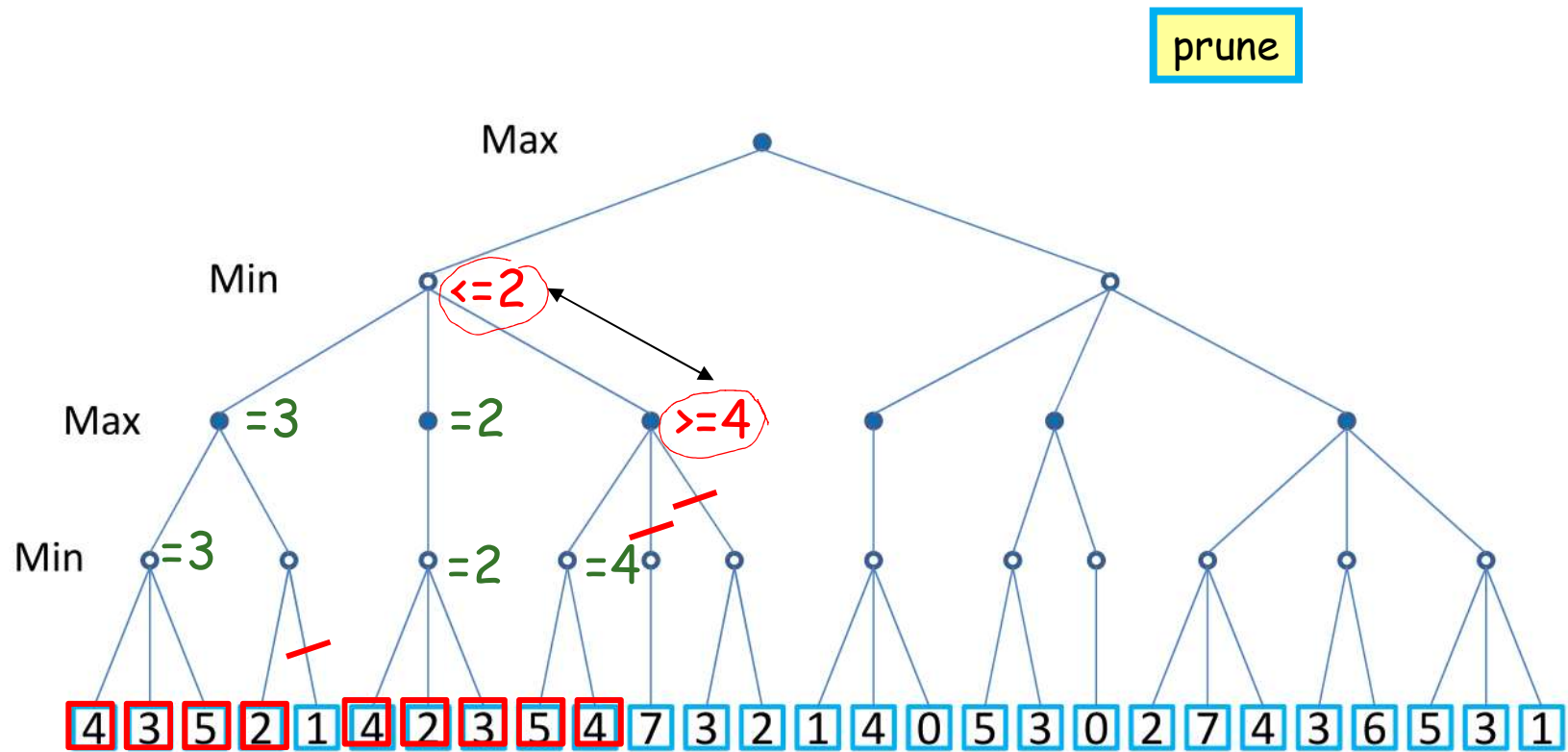
Alpha-Beta Pruning: Example 3



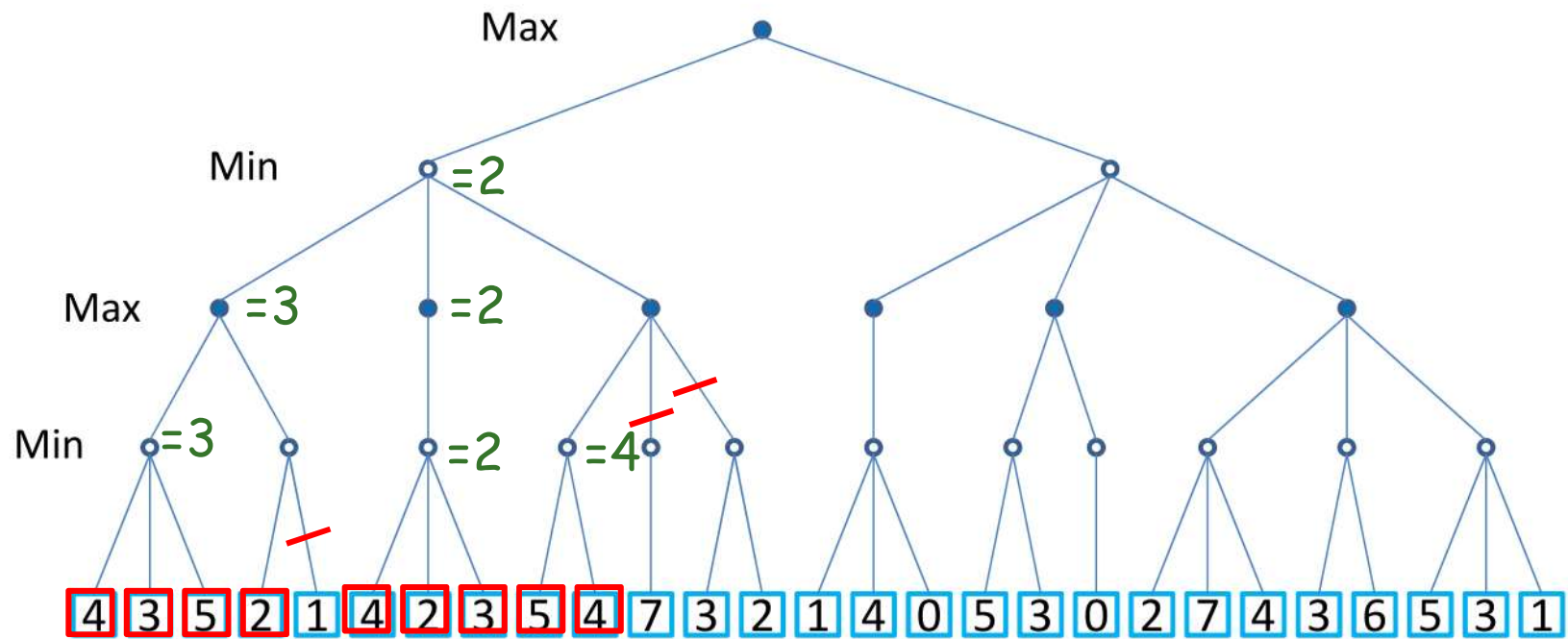
Alpha-Beta Pruning: Example 3



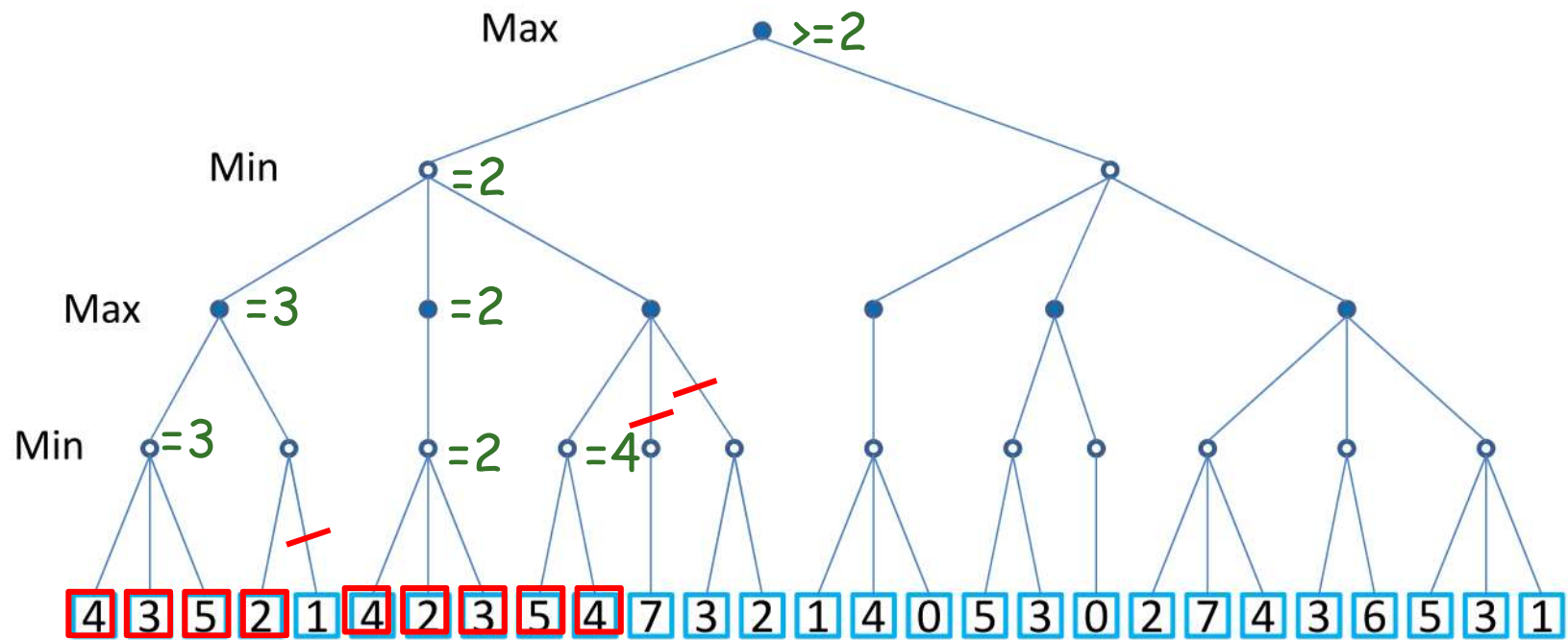
Alpha-Beta Pruning: Example 3



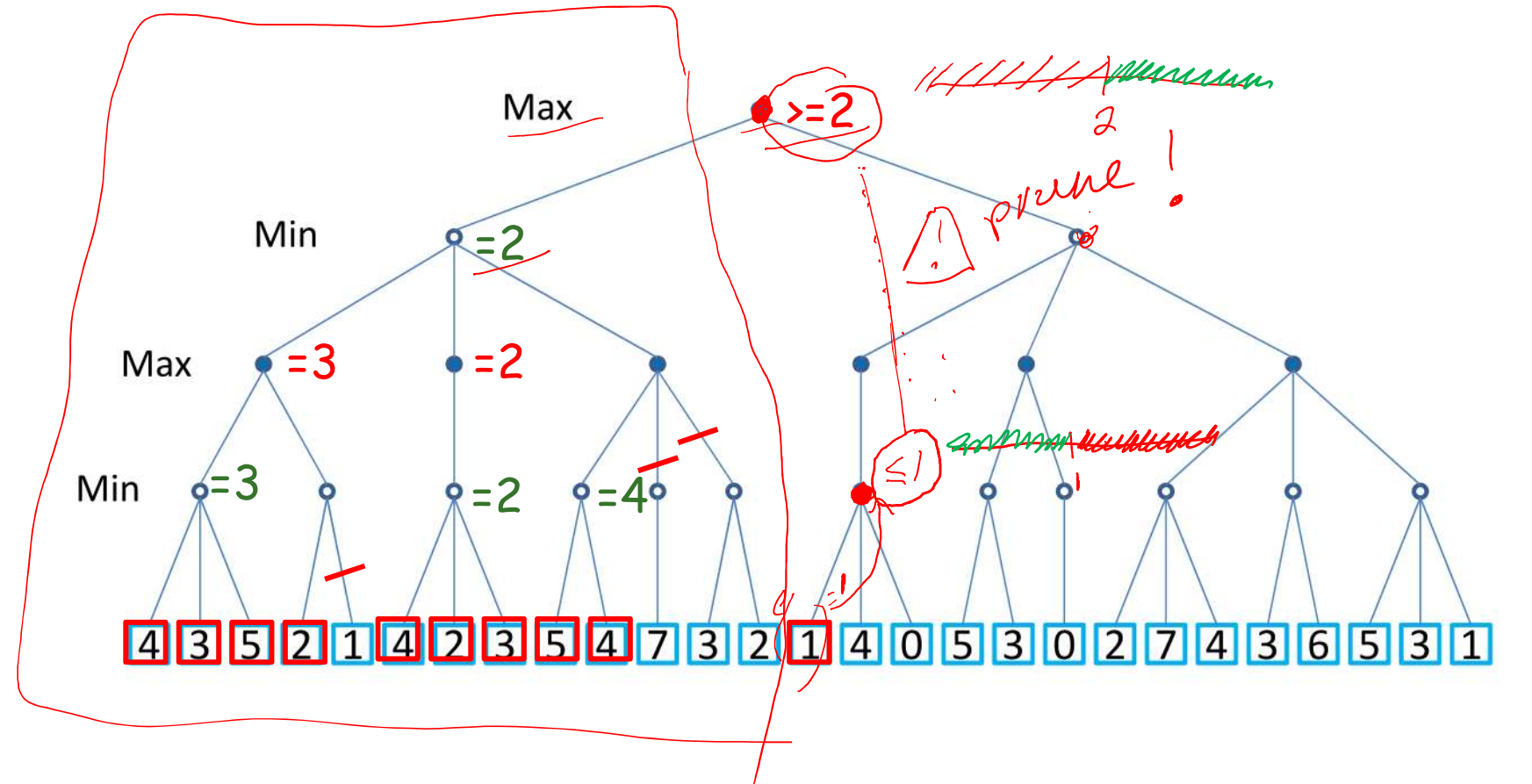
Alpha-Beta Pruning: Example 3



Alpha-Beta Pruning: Example 3



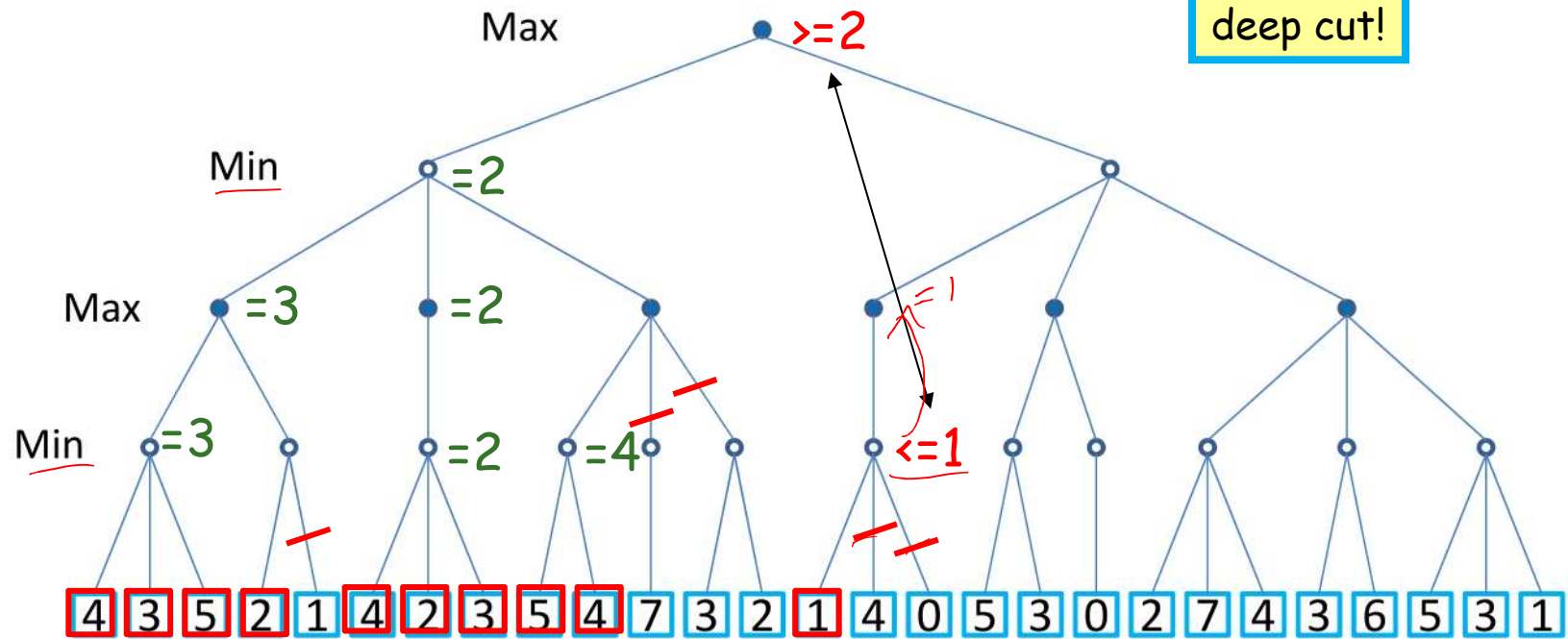
Alpha-Beta Pruning: Example 3



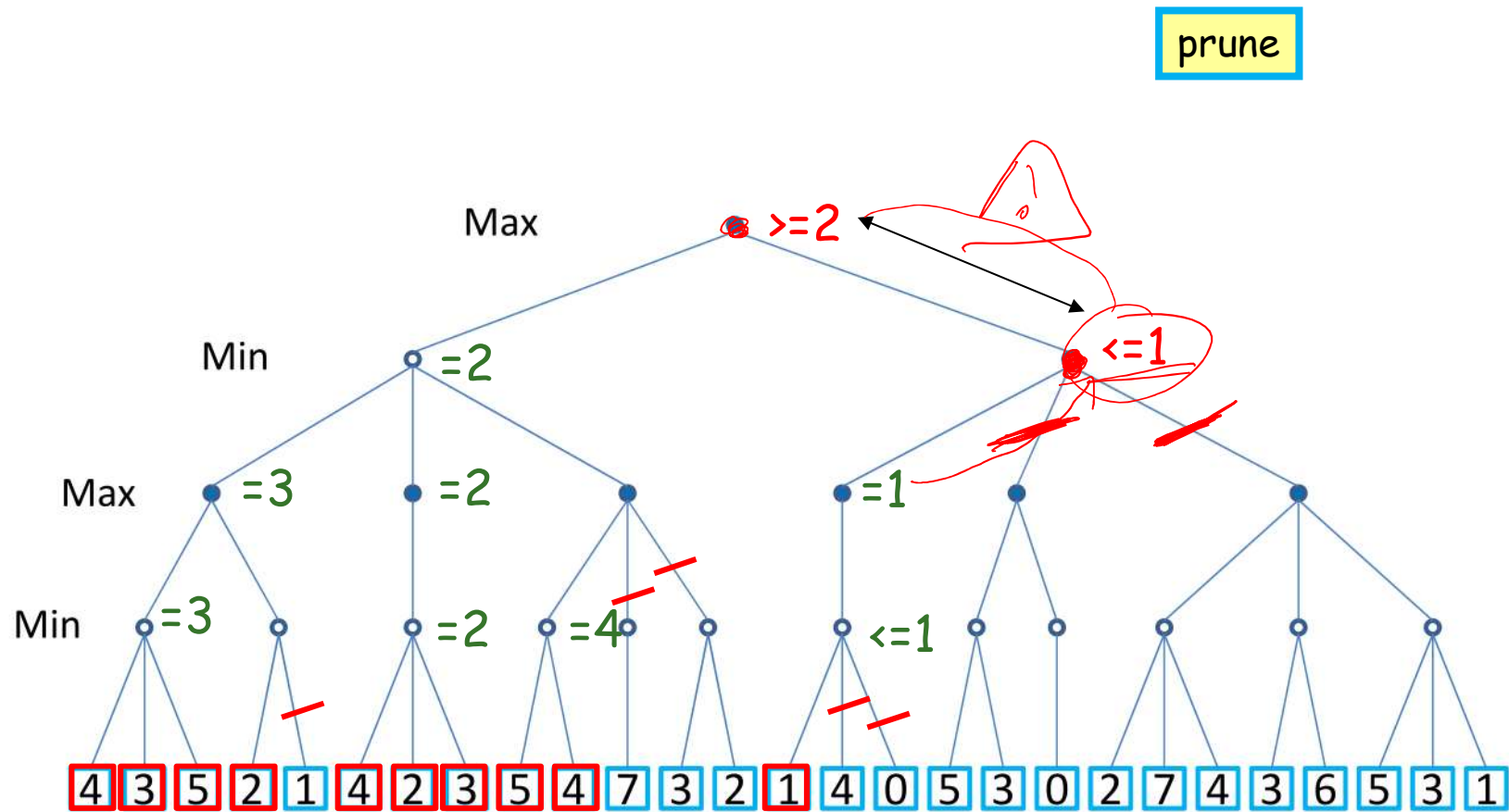
Alpha-Beta Pruning: Example 3

prune

deep cut!



Alpha-Beta Pruning: Example 3



10 nodes explored out of 27

Efficiency of Alpha-Beta Pruning

效率取决于sibling的排列顺序

- Depends on the order the siblings

- which is an arbitrary choice ;-(

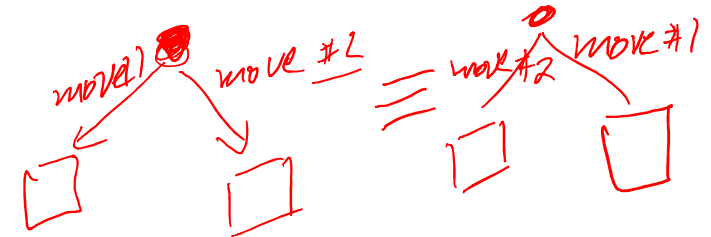
- In worst case: 最差没用

- alpha-beta provides no pruning
- plus extra overhead cost ;-(

最好开根

- In best case:

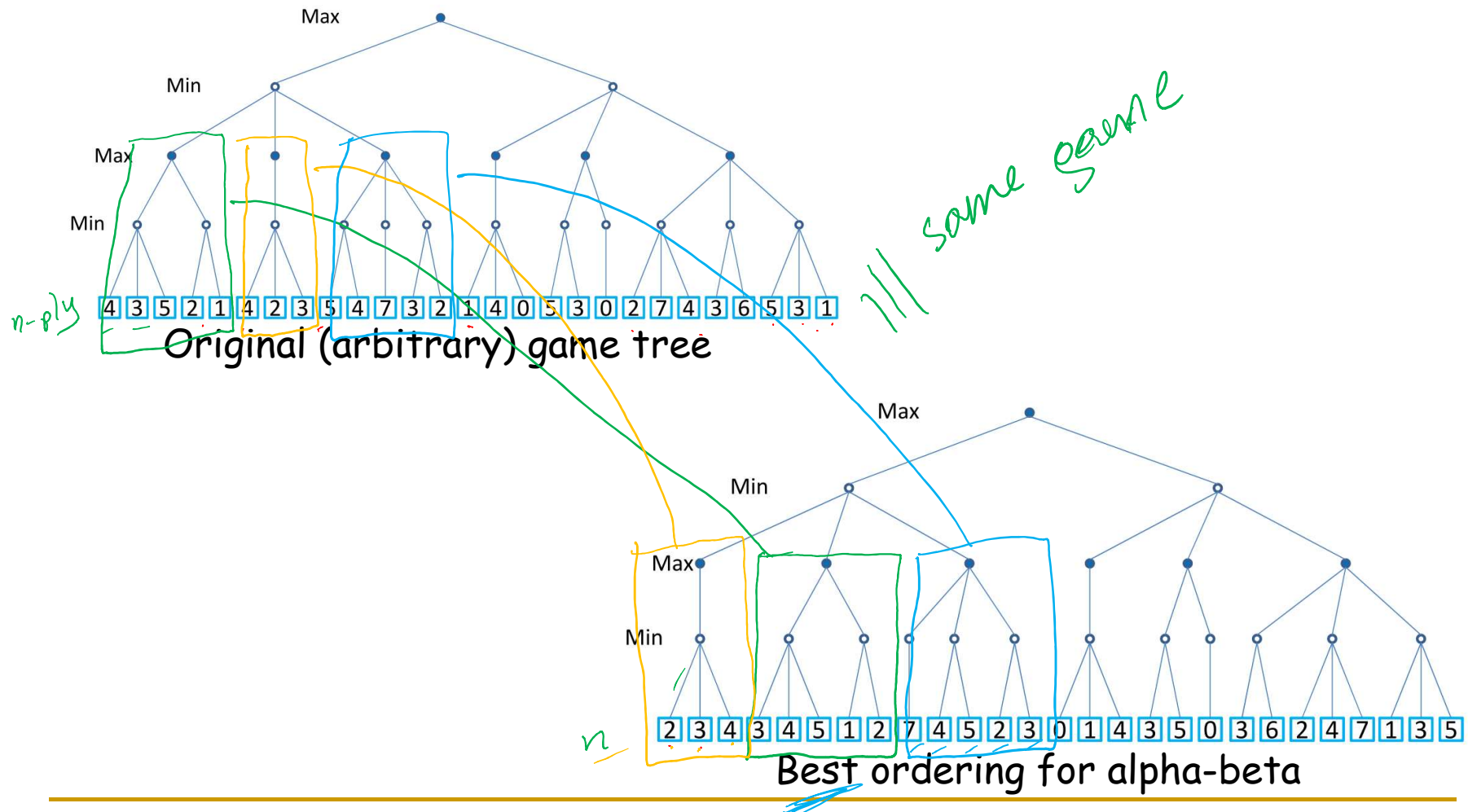
- branching factor is reduced to its square root



试图generate最优化的tree

$$\begin{aligned} & 36 e(n) \\ \Rightarrow & 6 e(n) \end{aligned}$$

Alpha-Beta: Best ordering



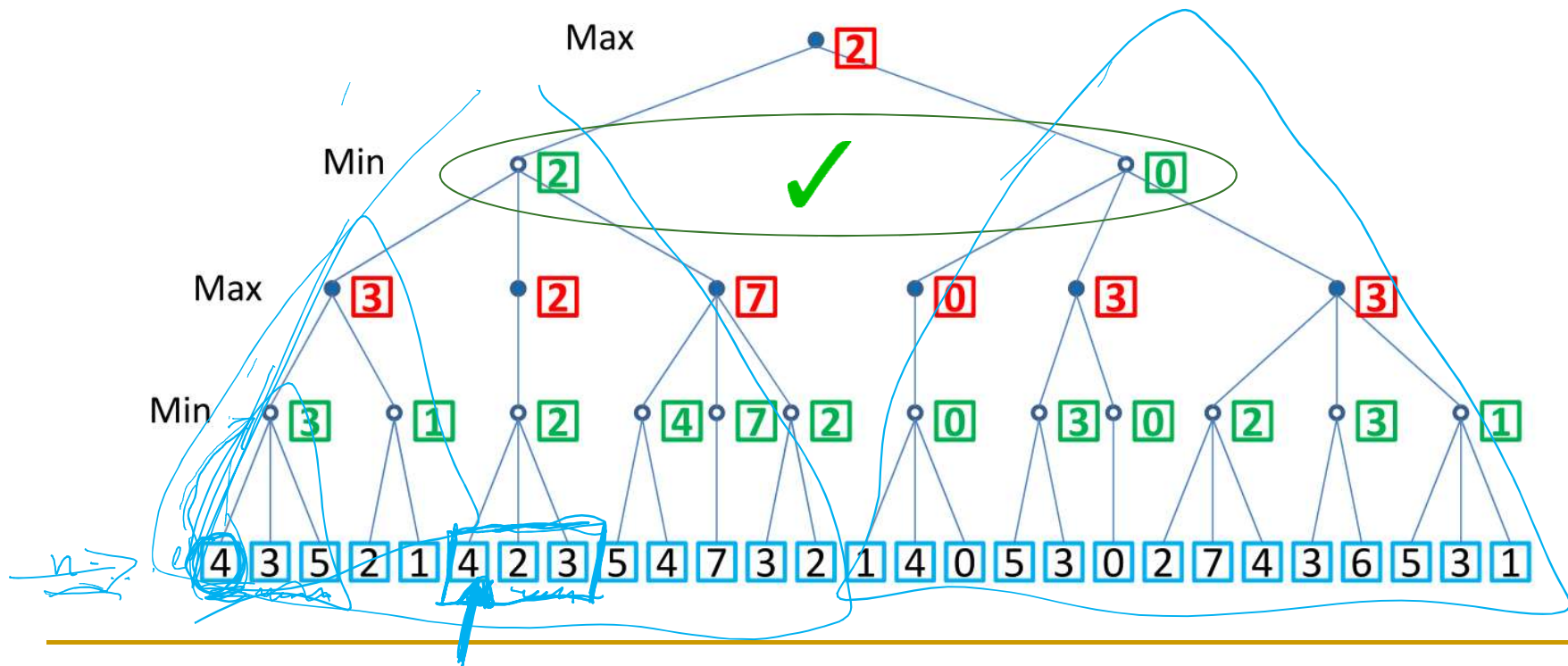
Alpha-Beta: Best ordering

- best ordering:

- strongest constraint placed first, ie:

- children of MIN: smallest node first

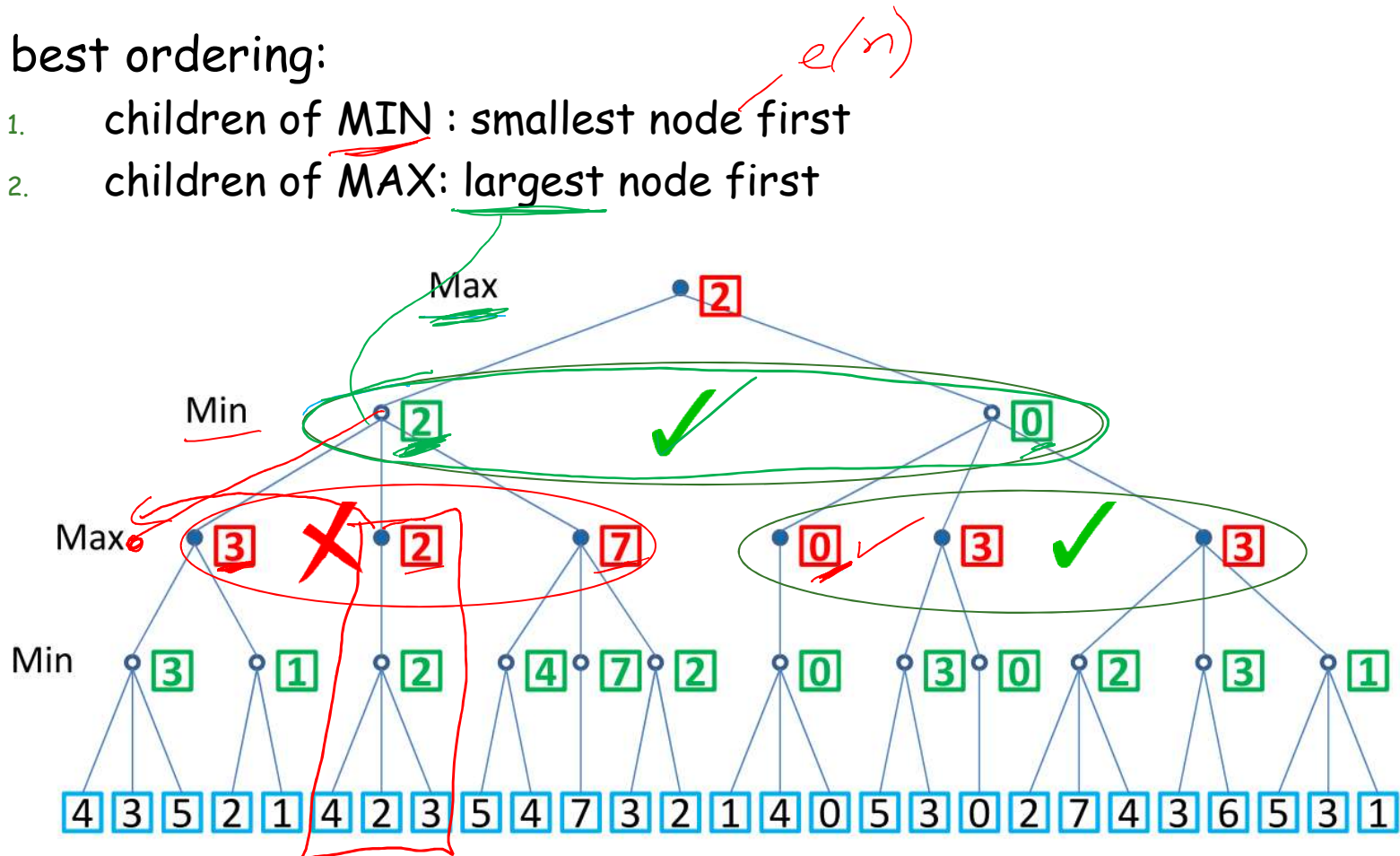
- children of MAX: largest node first



Alpha-Beta: Best ordering

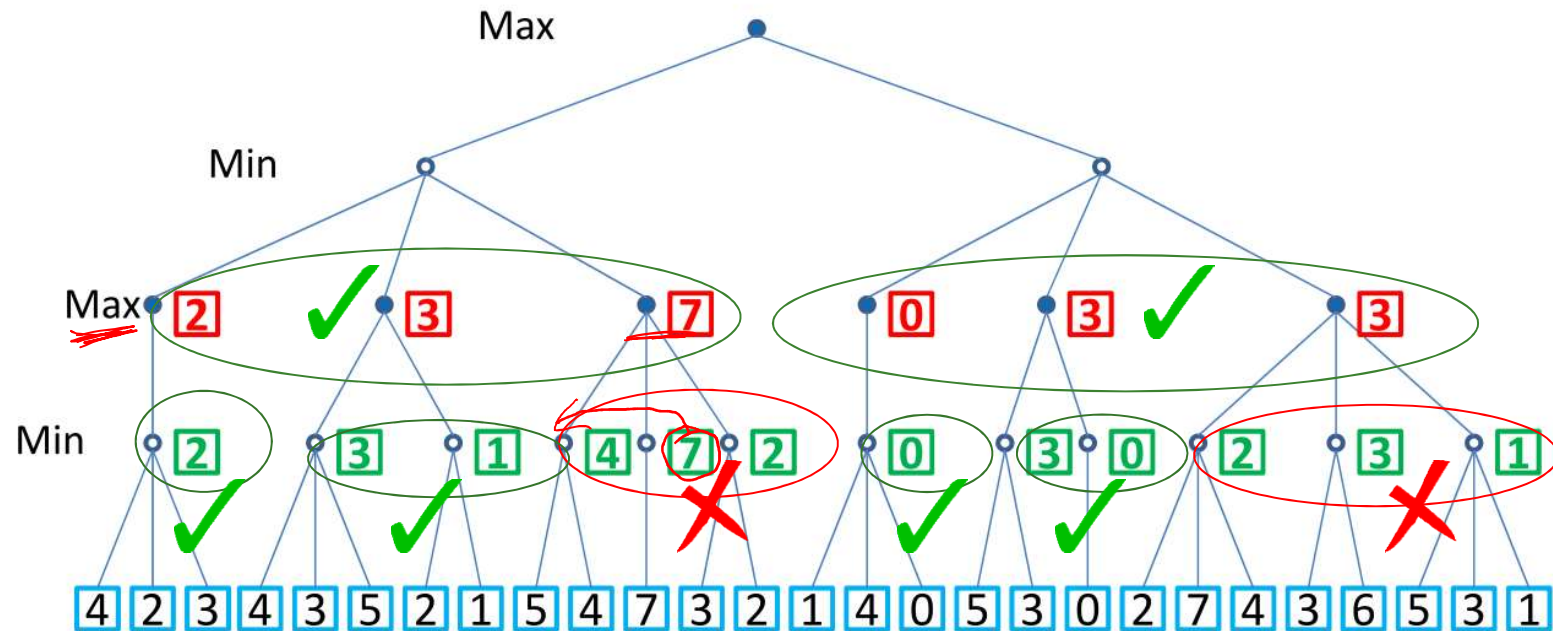
- best ordering:

1. children of MIN: smallest node first ^{$e(n)$}
2. children of MAX: largest node first

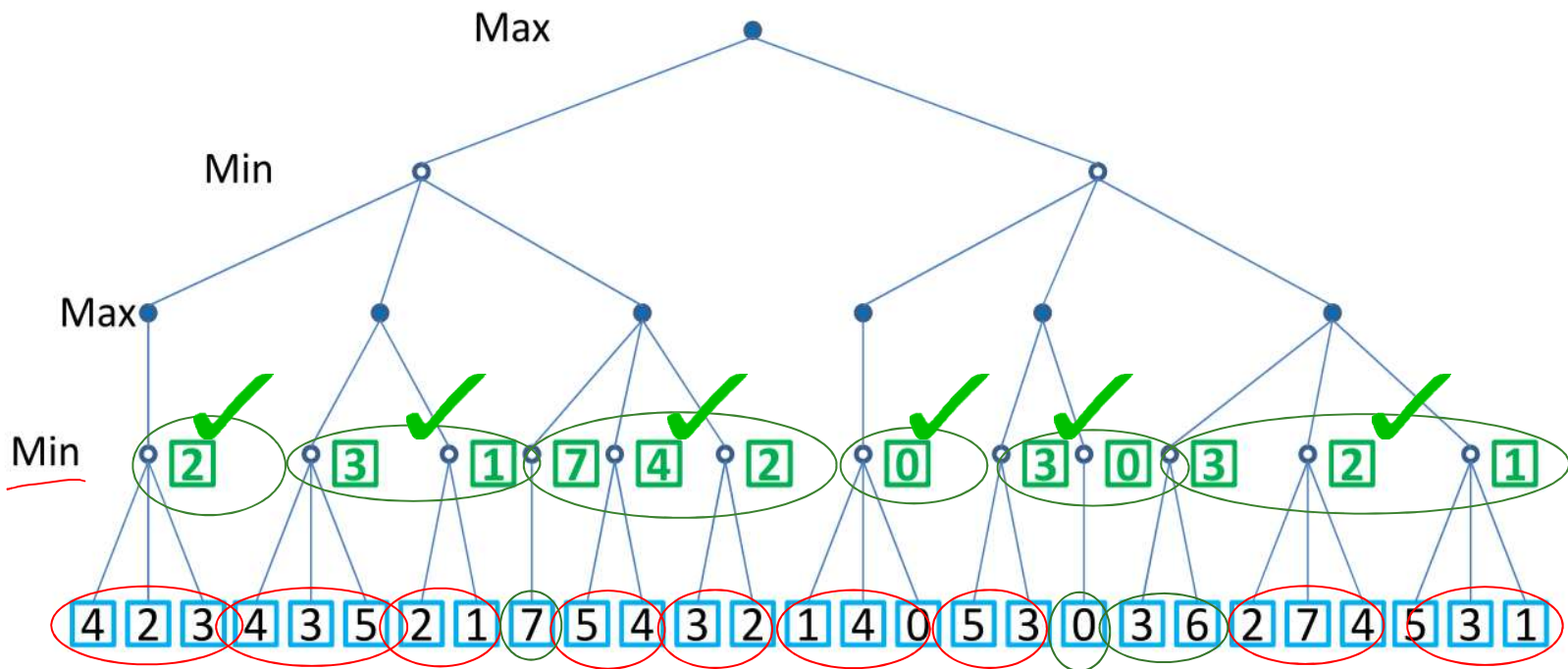


Alpha-Beta: Best ordering

- best ordering:
 1. children of MIN : smallest node first
 2. children of MAX: largest node first



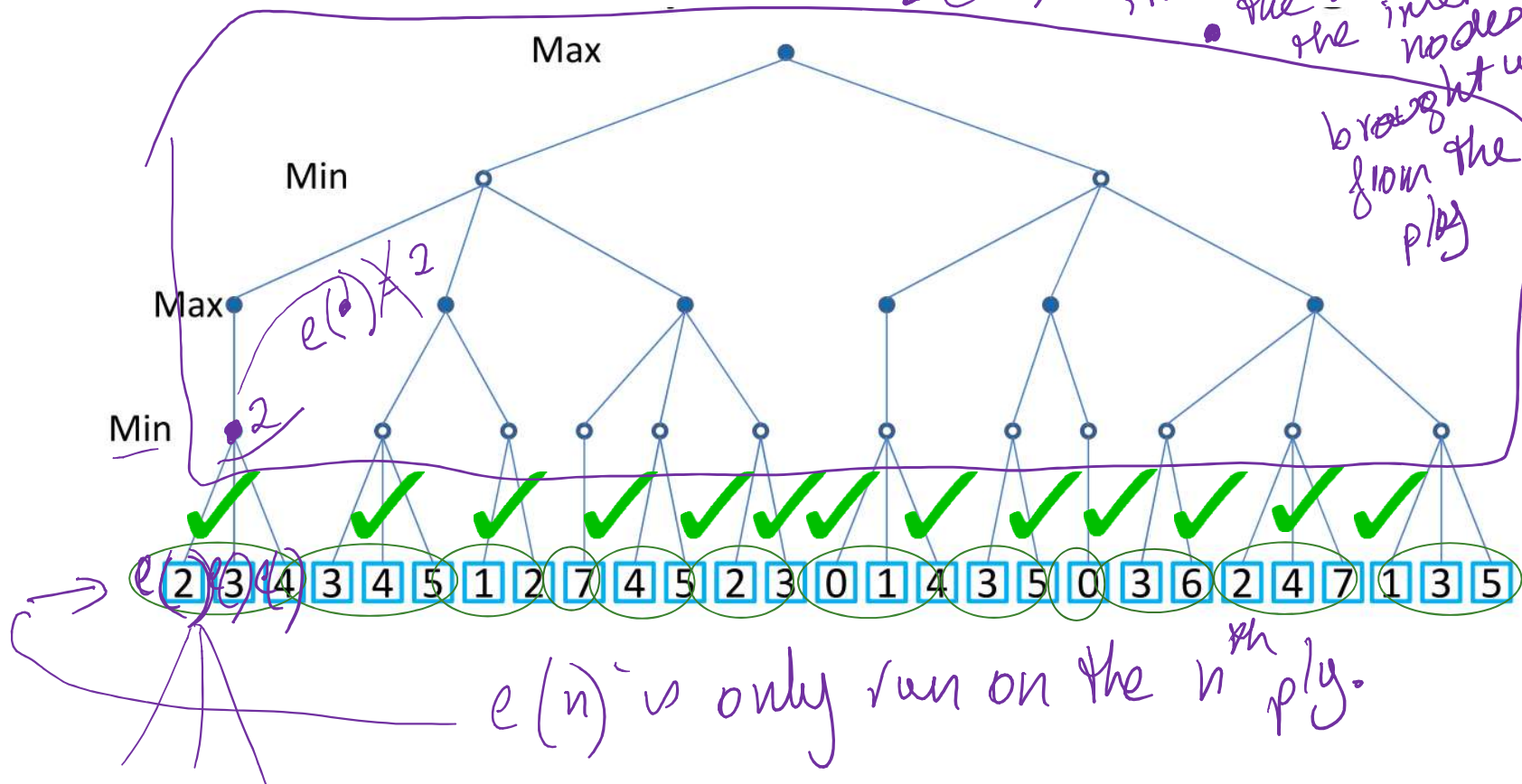
Alpha-Beta: Best ordering



Alpha-Beta: Best ordering

en 永远不是 internal node 算出来的
他是 nply 经过 ab 算法带上的

$e(n)$ is not run on internal nodes
the value of the internal nodes is brought up from the n^{th} ply



Alpha-Beta: Best ordering

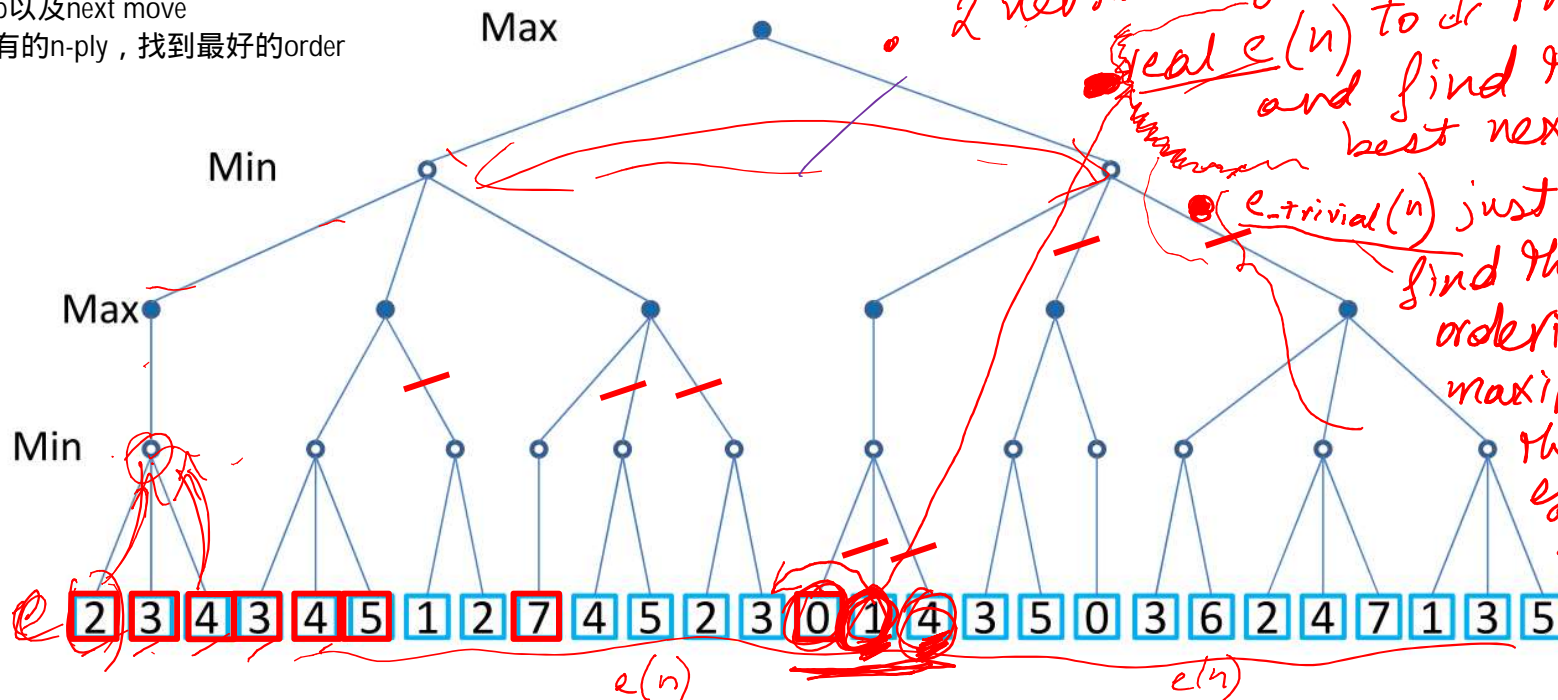


问题在于，想要进行排列，我们先要把所有 $e(n)$ 算出来，这样是矛盾的

因此我们设计两个版本的 $e(n)$

real $e(n)$: 用来进行 α - β 以及next move

e-trivial n ，遍历所有的 n -ply，找到最好的order



8 nodes explored out of 27

best ordering \Rightarrow run $e(n)$ less times

2 versions of $e(n)$

real $e(n)$ to do the α - β and find the best next move

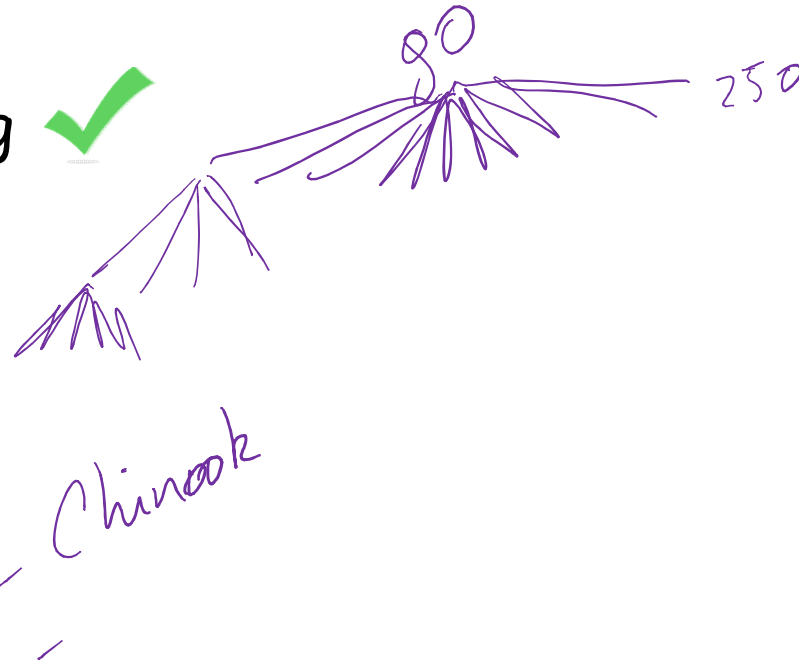
e-trivial n just to find the best ordering to maximize the effect of α - β

Today

■ Adversarial Search

1. Minimax ✓

2. Alpha-beta pruning ✓



Up Next

- Part 5: Natural Language Processing