uCOS-II Report (2)

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- Important Regs in Bf531
- Interrupt processing of Bf531
- VDSP 4.5 C Run-Time Mode for Bf531
- Porting uCOS-II to Bf531
- InterTask Communication and Synchronization
- Proposed Device Drivers Mode
- Q & A



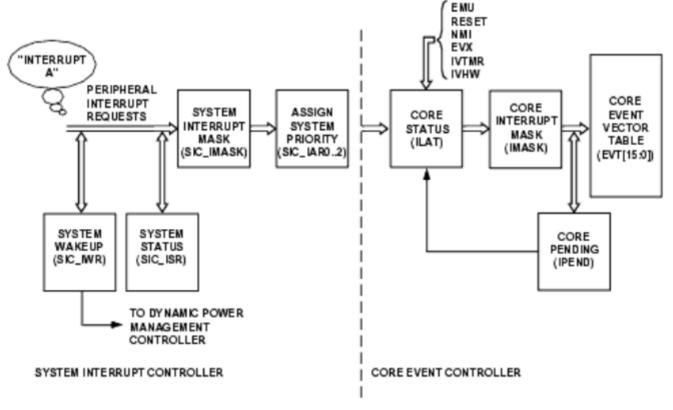
- SP & FP & PC (32位)
 - SP: 指向栈顶元素。堆栈从高到低生长。
 - FP: 指向最高Stack Frame 内存储有前一 Stack Frame 的FP寄存器内容的栈单元
 - PC: 指向下一条要执行的指令
- RETS (32(立)
 - 当用CALL指令调用函数的时候,处理器百动把返回地址保存到 RETS 中;函数执行完毕后的 RTS 指令将返回地址读到PC中



Important Regs in Bf531

- RETI (32(女)
 - 当发生中断时,在执行中断处理程序で前, 处理器将中断返回地址保存到 RETI中;在 中断处理程序完成时,执行 RTI 指令将中 断返回地址装入到 PC 中继续执行

■中断处理块图





- 全局中断的关闭和使能
 - CLI R0 ;
 - --将 IMASK 保存到 R0中,并清零 IMASK 寄存器,这会关闭所有通用目的的中断 (IVG5~IVG15)
 - STD R0
 - --将R0中的内容恢复到 IMASK 中,会使能原先允许的中断



- 事件向量表 (EVT)
 - 保存和名Core Event 所对应的中断向量;
 - ■对于一个外部中断a来说,假如在SIC_IAR 寄存器中将其映射到 Core Event IVG8上,则在中断发生的时候,处理器就从 EVT 中对应 IVG8 的位置处取出中断向量,并把某装入到 PC 寄存器中继续运行



- 软中断 (Software Interrupt)
 - 核心事件控制器(CEC)规定 将Core Event IVG14/IVG15 用做软中断
 - RAISE 指令用于发出软中断



- 万可嵌套的中断
 - 万要求将RETI保存在堆栈中,只需要保存 ISR中需要用到的那些寄存器即可;返回时 只需要用 RTI 指令

```
YoulSR()
{
//保存除 RETI 以外的其他要用到的寄存器

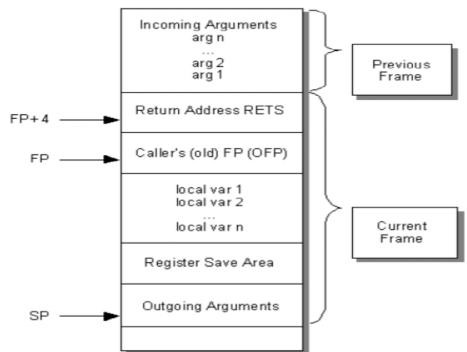
//其他处理

RTI
}
```

- ■可嵌套的中断
 - 其 ISR 要先负责将RETI 保存到堆栈中,等 到最后再从堆栈恢复RETI 寄存器

VDSP 4.5 C Run-Time Mode for Bf531

- ■堆栈的处理
 - Stack Frame: 一小段堆栈单元,用于保存 对应于当前正执行的C/C++函数的信息





VDSP 4.5 C Run-Time Mode for Bf531

- ■堆栈的处理
 - 进入被调用函数时:
 - a Linking Stack Frames ;
 - b Register Saving
 - 离开被调用函数时:
 - a · Restore Registers;
 - b Unlinking Stack Frame



- 被uCOS-II用到的处理器资源
 - Core Timer :
 - a,用于给uCOS-II的运行提供时钟节指;
 - b,使用IVG6 Core Event
 - Software Trap :
 - a,使用IVG14 Core Event 来进行任务级别的任务切换;



- 任务上下文(Task Context)
 - 寄存器上下文
 - a,当任务要放弃CPU时,需要把寄存器上下文保存到自己的堆栈中;
 - b,当任务要被处理器调度执行时,需要把寄存器上下文从自己堆栈中恢复出来。

所以,在a和b步两者它间,就应该有一个协议来规定寄存器上下文在堆栈中是如何保存的



- 任务上下文(Task Context)
 - 寄存器上下文在堆栈中的保存

| | High Memory |
|-------------|-------------|
| Ro | |
| - P1 | |
| RETS | |
| R1 | |
| R2 | |
| P0 | |
| P2 | |
| ASTAT | |
| RETI | |
| (R7:3,P5:3) | |
| FP | |
| 13:0 | |
| B3:0 | |
| L3:0 | |
| M3:0 | |
| A0.X | |
| AD.W | |
| A1.X | |
| A1.W | |
| LCO | |
| LC1 | |
| LTO | |
| LT1 | |
| LBO | |
| LB1 | |
| | High Memory |

注意:每个待运行的 任务(或者被剥夺CPU 使用权的任务),其 最后的堆栈顶端内容 都应该是寄存器上了 文是以后就可以正常 恢复执行。

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Porting uCOS-II to Bf531

os_cpu.h

os_cpu_c.c

os_cpu_c.c

```
void OS CPU RegisterHandler(INT8U ivg, FNCT PTR fn, BOOLEAN nesting)
 //...
 pEventVectorTable = (INT32U*)EVENT VECTOR TABLE ADDR;
 if (nesting == NESTED) {
      pEventVectorTable[ivg] = (INT32U)&OS CPU NESTING ISR;
 }else {
      pEventVectorTable[ivg] = (INT32U)&OS CPU NON NESTING ISR;
 OS_CPU_IntHanlderTab[ivg] = fn;
 //...
```

os_cpu_c.c

```
void OS CPU IntHandler (void)
    INT32U status;
    INT32U mask:
    INT8U
            i:
    mask
           = 1:
    status = *pIPEND & IPEND BIT 4 MASK;
    for (i = 0; i < IVG NUM; i++) {
        if ((1 << i) == (status & mask)) {
             if (OS CPU IntHanlderTab[i] != (void *)0) {
                 OS CPU IntHanlderTab[i]();
            break;
         mask <<=1;
```

-

Porting uCOS-II to Bf531

os_cpu_a.asm

```
_OSCtxSw:
    Save the CPU registers onto the old task's stack;

OSTCBCur->OSTCBStkPtr = SP;
OSTaskSwHook();

OSPrioCur = OSPrioHighRdy;
OSTCBCur = OSTCBHighRdy;
SP = OSTCBHighRdy->OSTCBStkPtr;

Restore the CPU registers from the new task's stack;
```

os_cpu_a.asm

```
OSIntCtxSw:
#if (OS TASK SW HOOK EN == 1)
   INIT C RUNTIME STACK (0x0)
   CALL OSTaskSwHook:
   DEL C RUNTIME STACK ()
#endif
   LOADA (P0, OSPrioCur);
   LOADA (P1, OSPrioHighRdy);
   R0 = B \overline{[P1](Z)}
   B[P0] = R0;
   LOADA (P0, OSTCBCur);
   LOADA (P1, OSTCBHighRdy);
   P2 = \overline{P1};
    [P0] = P2:
OSIntCtxSw modify SP:
   R0 = [P2];
   R0 += -8;
    [ FP ] = R0;
    SP += -4;
   RETI = [SP++];
 OSIntCtxSw.end:
  RTS;
```

os_cpu_a.asm

```
OS CPU NESTING ISR:
                                   OS CPU NON NESTING ISR:
   [ -- SP ] = R0;
                                       [ -- SP ] = R0;
                                      [ -- SP ] = P1;
   [ -- SP 1 = P1:
   [ -- SP ] = RETS;
                                      [ -- SP ] = RETS;
   R0 = NESTED;
                                       R0 = NOT NESTED;
   CALL.X OS CPU ISR Entry;
                                       CALL.X OS CPU ISR Entry;
   WORKAROUND 05000283()
                                       WORKAROUND 05000283()
   INIT C RUNTIME STACK (0x0)
                                       INIT C RUNTIME STACK (0x0)
   CALL.X OS CPU IntHandler;
                                       CALL.X OS CPU IntHandler
   SP
                                       CALL.X OSIntExit;
         = [ SP++ ];
                                       DEL C RUNTIME STACK()
   RETI
                                       JUMP.X OS CPU ISR Exit;
   CALL.X OSIntExit;
   DEL C RUNTIME STACK()
   JUMP.X OS CPU ISR Exit;
                                   OS CPU NON NESTING ISR.end:
OS CPU NESTING ISR.end:
  NOP:
                                       NOP;
                 #define DEL_C_RUNTIME_STACK()
```

UNLINK;

//ctong: 美中断

} ? end OSIntExit ?

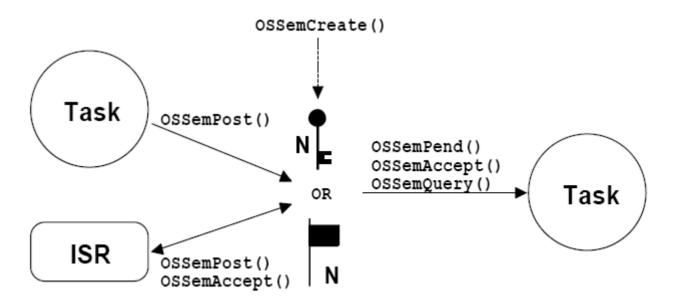
CT Close Interrupt();

Os core.c void OSIntExit (void) #if OS CRITICAL METHOD == 3 OS CPU SR cpu sr = 0; #endif if (OSRunning == OS TRUE) { OS ENTER CRITICAL(); if (OSIntNesting > 0) { OSIntNesting--; if (OSIntNesting == 0) { if (OSLockNesting == 0) { OS SchedNew(); if (OSPrioHighRdy != OSPrioCur) { OSTCBHighRdv = OSTCBPrioTb1[OSPrioHighRdv]; #if OS TASK PROFILE EN > 0 OSTCBHighRdy->OSTCBCtxSwCtr++; #endif OSCtxSwCtr++; OSIntCtxSw(); OS EXIT CRITICAL(); } ? end if OSRunning==OS_TRUE ?



InterTask Communication and Synchronization

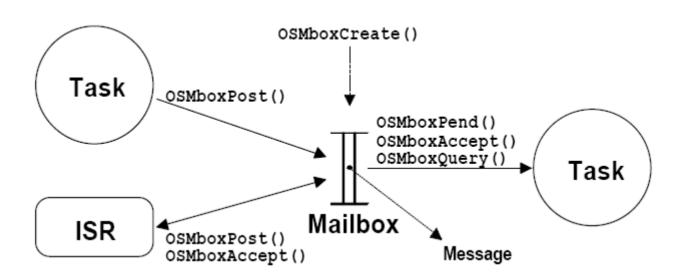
Semaphore





InterTask Communication and Synchronization

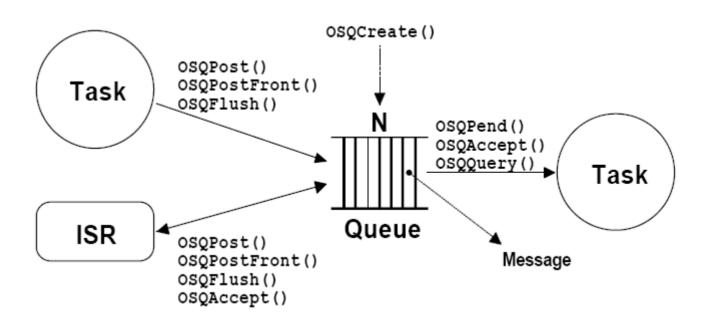
Message Box





InterTask Communication and Synchronization

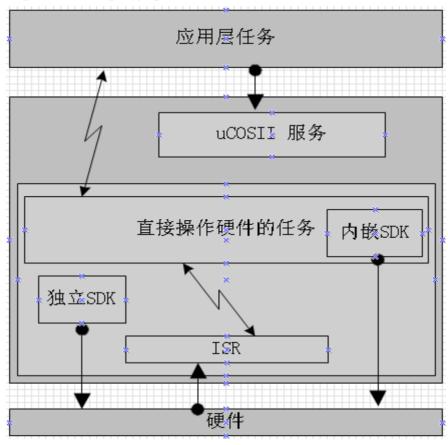
Message Queue



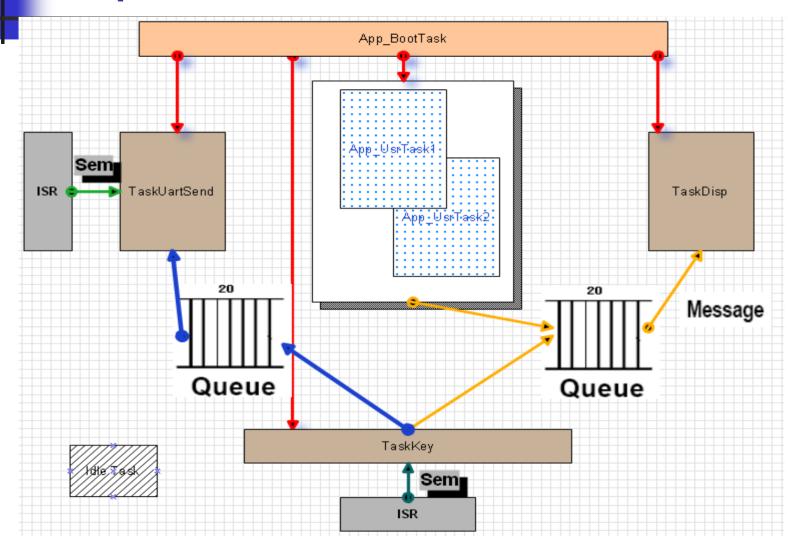


Proposed Device Drivers Mode

Driver Mode



Proposed Device Drivers Mode





THE END

Thanks a lot, any suggestion?