

Lab #3 Applications of Vitis HLS & PYNQ

ECE4810J System-on-Chip Design



Yihua Liu
UM-SJTU Joint Institute

ayka_tsuzuki@sjtu.edu.cn
Oct. 10, 2022

Overview



- 1 Overview
- 2 High Level Synthesis
- 3 Reference



Algorithmic HLS does several things automatically that an RTL designer does manually:

- HLS analyzes and exploits the concurrency in an algorithm.
- HLS inserts registers as necessary to limit critical paths and achieve a desired clock frequency.
- HLS generates control logic that directs the data path.
- HLS implements interfaces to connect to the rest of the system.
- HLS maps data onto storage elements to balance resource usage and bandwidth.
- HLS maps computation onto logic elements performing user specified and automatic optimizations to achieve the most efficient implementation.

We make the following assumptions about the input function specification:

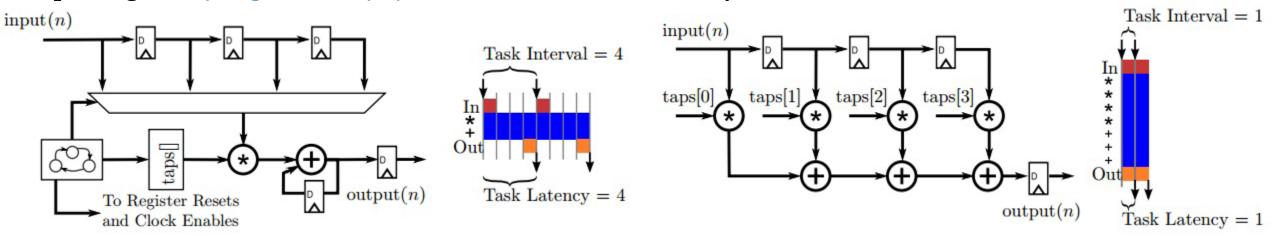
- No dynamic memory allocation (no operators like malloc(), free(), new, and delete())
- Limited use of pointers-to-pointers (e.g., may not appear at the interface)
- System calls are not supported (e.g., abort(), exit(), printf(), etc. They can be used in the code, e.g., in the testbench, but they are ignored (removed) during synthesis.
- Limited use of other standard libraries (e.g., only common math.h functions are supported)
- Limited use of function pointers and virtual functions in C++ classes (function calls must be compile-time determined by the compiler).
- No recursive function calls.
- The interface must be precisely defined.



```
fir:
                                                                     A key question to understand is: "What
        .frame
                r1,0,r15
                                  # vars= 0, regs= 0, args= 0
                0x00000000
        .mask
                                                                     circuit is generated from this code?".
        addik
                r3,r0,delay line.1450
        lwi
                                  # Unrolled loop to shift the delay line
                r4, r3, 8
        swi
                r4, r3, 12
        lwi
               r4, r3,4
                r4, r3,8
        swi
        lwi
                r4, r3,0
                                                         branch
               r4, r3,4
        swi
                r5, r3,0
                                  # Store the new input sample into the delay line
        swi
        addik
                r5, r0,4
                                  # Initialize the loop counter
        addk
                r8, r0, r0
                                  # Initialize accumulator to zero
        addk
                r4, r8, r0
                                  # Initialize index expression to zero
$L2:
        muli
                r3, r4, 4
                                  # Compute a byte offset into the delay line array
        addik
                r9, r3, delay_line.1450
                r3,r3,r7
                                  # Load filter tap
        1w
                                                                         One possible circuit would execute
                                  # Load value from delay line
        lwi
                r9, r9, 0
                                                                         the code sequentially, as would a
                                  # Filter Multiply
        mul
                r3, r3, r9
        addk
                r8, r8, r3
                                  # Filter Accumulate
                                                                         simple RISC microprocessor.
               r5,r5,-1
        addik
                                  # update the loop counter
        bneid
                r5,$L2
                                                                                 microblazeel-xilinx-linux-gnu-gcc
        addik
                r4,r4,1
                                  # branch delay slot, update index expression
                                                                                 -01 -mno -xl -soft -mul S fir.c
                r15, 8
        rtsd
                r8, r6, 0
                                  # branch delay slot, store the output
        swi
        end
                fir
```



However, the Vitis HLS tool can also generate higher performance pipelined and parallel architectures. One important class of architectures is called a *function pipeline*. A function pipeline architecture is derived by considering the code within the function to be entirely part of a computational data path, with little control logic. Loops and branches in the code are converted into unconditional constructs. The Vitis HLS tool can be directed to generate a function pipeline by placing the #pragma HLS pipeline directive in the body of a function.



A "one tap per clock" architecture for an FIR filter. A "one sample per clock" architecture for an FIR filter.



Arbitrary Precision Data Types Library

Language		Integer Data Type		Required Header	
C++		<pre>ap_[u]int<w> (Can be extende</w></pre>	•	<pre>#include <ap_int.h></ap_int.h></pre>	
C++		ap_[u]fixed <w,< th=""><th>I,Q,O,N></th><th><pre>#include <ap_fixed.h></ap_fixed.h></pre></th></w,<>	I,Q,O,N>	<pre>#include <ap_fixed.h></ap_fixed.h></pre>	
W I Q	The number	Word length in bits The number of bits used to represent the integer value Quantization mode			
•	AP_RND AP_RND_ AP_RND_ AP_RND_	ZERO MIN_INF INF	Rounds to plus Rounds to zero Rounds to min Rounds to infin	us infinity nity	
O	AP_RND_ AP_TRN AP_TRN_2 Overflow n	ZERO	Convergent rou Truncation to a	minus infinity (default)	
	AP_SAT AP_SAT_Z AP_SAT_S AP_WRAP AP_WRAP	SYM	Saturation Saturation to z Symmetrical sa Wrap around (Sign magnitud	aturation default)	
N		This defines the number of saturation bits in overflow wrap mode			



```
#include "cpp_ap_int_arith.h"
void cpp ap int arith (din A inA, din B inB, din C
inC, din D inD, dout 1 *out1, dout 2 *out2, dout 3
*out3, dout 4 *out4) {
    *out1 = inA * inB;
    *out2 = inB + inA;
    *out3 = inC / inA;
    *out4 = inD % inA;
#ifndef CPP_AP_INT_ARITH_H_
#define CPP AP INT ARITH H
#include <stdio.h>
#include "ap int.h"
typedef ap int<6> dinA t;
typedef ap int<12> dinB t;
typedef ap int<22> dinC t;
typedef ap int<33> dinD t;
typedef ap int<18> dout1 t;
typedef ap uint<13> dout2 t;
typedef ap int<22> dout3 t;
typedef ap_int<6> dout4_t;
void cpp ap int arith (dinA t inA, dinB t inB,
dinC t inC, dinD t inD, dout1 t *out1, dout2 t *out2,
dout3 t *out3, dout4 t *out4);
#endif
```

The default maximum width allowed is 1024 bits. You can override this default by defining the macro AP_INT_MAX_W with a positive integer value less than or equal to 4096 before inclusion of the ap_int.h header file.

```
#define AP_INT_MAX_W 4096 // Must be defined before next line
#include "ap_int.h"
ap_int<4096> very_wide_var;
```

One disadvantage of AP data types is that arrays are not automatically initialized with a value of 0. You must manually initialize the array if desired.

If C++ Arbitrary Precision Integer Types are synthesized, it results in a design that is functionally identical to Standard Types. To allow assignment of values wider than 64-bits, the ap_[u]int<> classes provide constructors that allow initialization from a string of arbitrary length. To avoid unexpected behavior during cosimulation, do not initialize ap_uint<N> a ={0}.

```
ap_uint<96> wide_var("76543210fedcba9876543210", 16);
wide_var = ap_int<96>("0123456789abcdef01234567", 16);
```



Base FIR Architecture

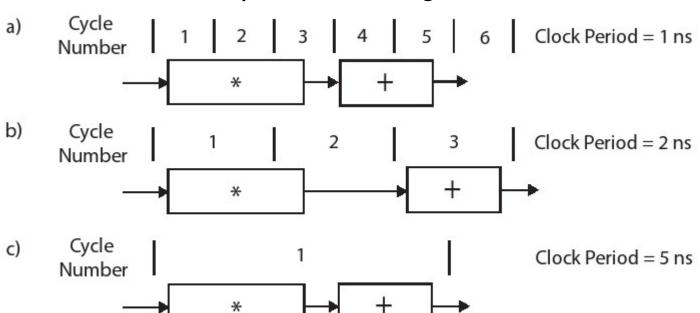
```
#define N 11
#include "ap int.h"
typedef int coef t;
typedef int data_t;
typedef int acc t;
void fir(data_t *y, data_t x) {
    coef_t c[N] = \{53, 0, -91, 0, 313, 500, 313, 0, -91, 0, 53\};
    static data_t shift_reg[N]; The label Shift_Accum_Loop: is not
    acc t acc;
                                 necessary. However it can be useful
    int i;
                                 for debugging. The Vitis HLS tool adds
    acc = 0;
                                 these labels into the views of the code.
Shift Accum Loop:
    for (i = N - 1; i >= 0; i--) {
        if (i == 0) {
            acc += x * c[0];
            shift reg[0] = x;
        } else {
            shift reg[i] = shift reg[i - 1];
            acc += shift_reg[i] * c[i];
```

Calculating Performance

It is possible to specify a target frequency to the Vitis HLS tool using the create_clock tcl command. For example, the command create_clock -period 5 directs the tool to target a clock period of 5 ns and equivalently a clock frequency of 200 MHz. Vitis HLS deals with clock frequency estimates including some margin to account for the fact that there is some error in the estimate. The goal of this margin is to ensure enough timing slack in the design that the generated RTL can be successfully placed and routed. This margin can be directly controlled using the set_clock_uncertainty TCL command.







Code Hoisting

```
Shift_Accum_Loop:
for (i = N - 1; i > 0; i--) {
    shift_reg[i] = shift_reg[i - 1];
    acc += shift_reg[i] * c[i];
}
Removing the conditional statement from the for loop
acc += x * c[0];
shift_reg[0] = x;

creates a more efficient hardware implementation.
```

Loop Fission

Loop fission alone often does not provide a more efficient hardware implementation. However, it allows each of the loops to be optimized independently, which could lead to better results than optimizing the single, original **for** loop.

The reverse is also true; merging two (or more)

for loops into one for loop may yield the best
results. This is highly dependent upon the
application, which is true for most optimizations.



Loop Unrolling

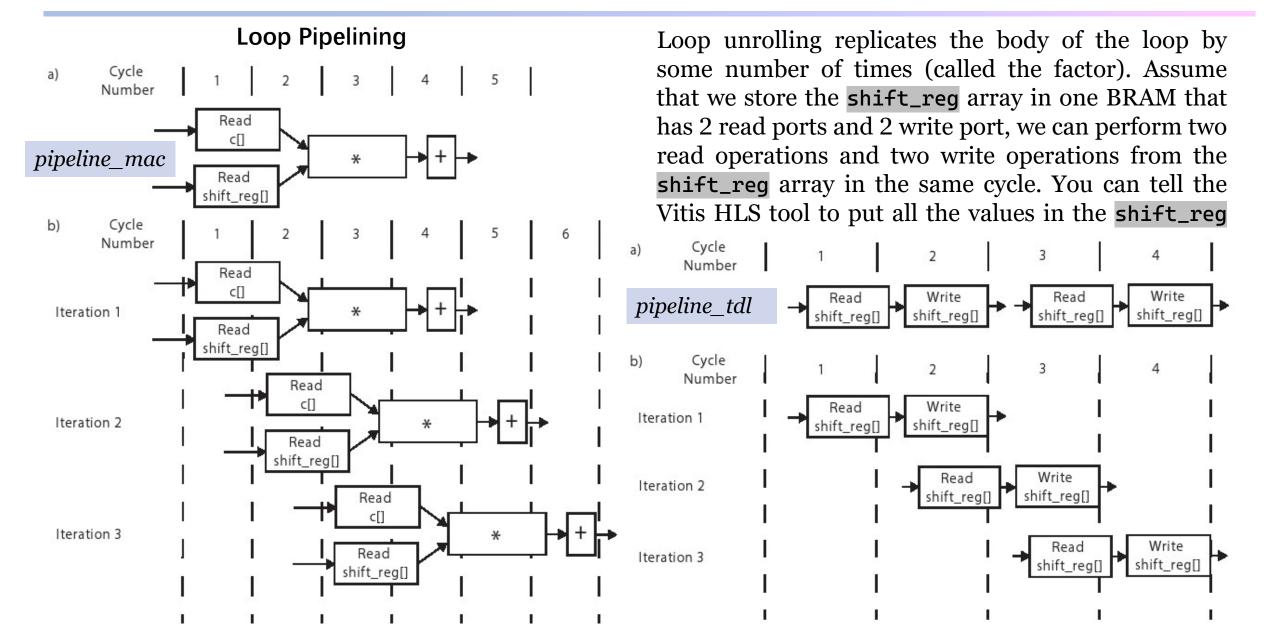
```
TDL:
for (i = N - 1; i > 1; i = i - 2) {
    shift_reg[i] = shift_reg[i - 1];
    shift_reg[i - 1] = shift_reg[i - 2];
}
if (i == 1) {
    shift_reg[1] = shift_reg[0];
}
shift_reg[0] = x;
```

Loop unrolling replicates the body of the loop by some number of times (called the factor). Assume that we store the <code>shift_reg</code> array in one BRAM that has 2 read ports and 2 write port, we can perform two read operations and two write operations from the <code>shift_reg</code> array in the same cycle. You can tell the Vitis HLS tool to put all the values in the <code>shift_reg</code> array into registers using the directive <code>#pragma HLS array_partition variable=shift_reg complete</code>.

A user can tell the Vitis HLS tool to automatically unroll the loop using the unroll directive. To automatically perform the unrolling, we should put the directive **#pragma** HLS unroll factor=2 into the body of the code, right after the **for**-loop header. By specifying the optional argument **skip_exit_check** in that directive, the Vitis HLS tool will not add the final **for** loop to check for partial iterations.

```
acc = 0;
MAC:
for (i = N - 1; i >= 3; i -= 4) {
    acc += shift_reg[i] * c[i] +
    shift_reg[i - 1] * c[i - 1] +
    shift_reg[i - 2] * c[i - 2] +
    shift_reg[i - 3] * c[i - 3];
}
for (; i >= 0; i--) {
    acc += shift_reg[i] * c[i];
}
```







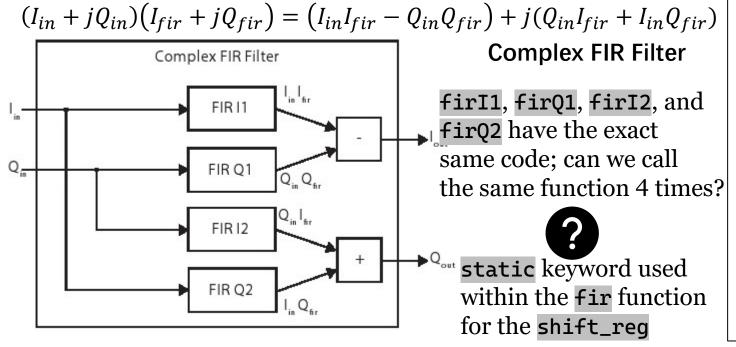
Bitwidth Optimization

coef t $c[N] = \{53, 0, -91, 0, 313, 500, 313, 0, -91, 0, 53\};$

The maximum absolute value for any of these 11 entries is 500, which requires $\lceil \log 2 \ 500 \rceil = 9$ bits.

Since we need negative numbers, we add an additional bit. Thus, $coef_t$ can be declared as $ap_int<10>$. Consider the operation a = b + c where $ap_uint<10> b$ and $ap_uint<10> c$. The data type for a is $ap_uint<2>$ where z = max(x, y) + 1.

The value the bitwidth z given the bitwidths x and y (i.e., $ap_int < z > a$, $ap_int < x > b$, $ap_int < y > c$) for the operation a = b * c is z = x + y.



```
typedef int data_t;
void firI1(data_t *y, data_t x);
void firQ1(data_t *y, data_t x);
void firI2(data_t *y, data_t x);
void firQ2(data_t *y, data_t x);
void complexFIR(data_t Iin, data_t Qin,
data_t *Iout, data_t *Qout) {
    data_t IinIfir,QinQfir,QinIfir,IinQfir;
    firI1(&IinIfir,Iin);
    firQ1(&QinQfir,Qin);
    firQ2(&QinIfir,Qin);
    firQ2(&IinQfir,Iin);
    *Iout = IinIfir + QinQfir;
    *Qout = QinIfir - IinQfir;
}
```



Bitwidth Optimization

The Vitis HLS tool does not optimize across function boundaries, i.e., each fir function synthesized independently, and treated as a black box in the **complexFIR** function. You can use the **inline** directive if you want the Vitis HLS tool to co-optimize a particular function within its parent function.

- 1. can increase the potential for benefits in performance and area
- 2. creates a large amount of code that the tool must synthesize, which may take a long time, even fail to synthesize, or may result in a non-optimal design
- 3. eliminates any overhead associated with performing the function call.

Therefore, use the inline directive carefully. Also note that the Vitis HLS tool may choose to inline functions on its own. These are typically functions with a small amount of code. You can force the tool to keep the function hierarchy by **inline off**. The **inline** directive also has a **recursive** argument that recursively adds the code into the parent functions from every child function.

```
float mul(int x, int y) {
    return x * y;
}
float top_function(float a, float b, float c, float d) {
    return mul(a, b) + mul(c, d) + mul(b, c) + mul(a, d);
}
float inlined_top_function(float a, float b, float c, float d) {
    return a * b + c * d + b * c + a * d;
}
```

Reference



1 Ryan Kastner, Janarbek Matai, and Stephen Neuendorffer. "Parallel Programming for FPGAs." Sept. 6, 2022.