

浙江大学 2014-2015 学年 秋冬 学期

《计算理论》课程期末考试试卷

课程号: 21120520 开课学院: 计算机学院

考试试卷: ☒ A卷 ☐ B卷

考试形式: ☒ 闭卷 ☐ 开卷, 允许带 _____ 入场

考试日期: 2015 年 1 月 27 日, 考试时间: 120 分钟

诚信考试, 沉着应考, 杜绝违纪

考生姓名 _____ 学号 _____ 所属院系 _____

题序	1	2	3	4	5	6	总分
得分							
评卷人							

Zhejiang University
Theory of Computation, Fall-Winter 2014
Final Exam

1. (24%) Determine whether the following statements are true or false. If it is true fill a \bigcirc otherwise a \times in the bracket before the statement.
- (a) () If L is any language, then the language LL^R must be equal to $\{ww^R \mid w \in L\}$.
 - (b) () Language $\{a^ib^jc^k \mid i, j, k \in \mathbb{N} \text{ and } i + j \not\equiv k \pmod{3}\}$ is not regular.
 - (c) () If A is non-regular and both of B and $A \cap B$ are regular, then $A \cup B$ is non-regular.
 - (d) () For all languages L_1, L_2 and L_3 , if $L_1 \subseteq L_2 \subseteq L_3$ and both L_1 and L_3 are regular, then L_2 is also regular.
 - (e) () Language $\{xycy \mid x, y \in \{a, b\}^* \text{ and } |x| \leq |y| \leq 2|x|\}$ is context-free.
 - (f) () Let L_1 be a regular language and L_2 be a context-free language, then $\{uv \mid u \in L_1, v \in L_2 \text{ and } |u| = |v|\}$ is also context-free.
 - (g) () Let $\mathbf{D}_{\text{DFA}} = \{\langle M \rangle \mid \text{DFA } M \text{ rejects } \langle M \rangle\}$, where " M " is the encoding of DFA M , just as Turing Machine, then \mathbf{D}_{DFA} is recursively enumerable but not regular.
 - (h) () Let L be a language and there is a Turing machine M halts on x for every $x \in L$, then L is decidable.
 - (i) () Every countably infinite language is recursively enumerable.
 - (j) () A language is recursively enumerable if and only if it is Turing enumerable.
 - (k) () Let A be a recursively enumerable language and $A \leq_{\tau} \bar{A}$, then A is recursive.
 - (l) () There are countably many Turing machines, and uncountably many languages, so most languages are not recursively enumerable.

2. (20%) Decide whether the following languages are regular or not and provide a formal proof for your answer. Let $\#_a(u)$ and $\#_a(v)$ be the number of a in string u and v , respectively.

(a) $L_1 = \{ucv \mid u, v \in \{a, b\}^*, \#_a(u) = 2 \cdot \#_a(v)\}$

Not Regular

$\exists ucv \in L$, , assume the pump length is $|u|$.

let $u = xa$ or $u = xb$.

for both case, $\forall i \neq 1, i \in \mathbb{N}$,

$xa^i cv \notin L$, $xb^i cv \notin L$.

(b) $L_2 = \{uv \mid u, v \in \{a, b\}^*, \#_a(u) = 2 \cdot \#_a(v)\}$

Regular. $uv = (b^* a b^* a b^* a b^*)^*$

3. (24%) On PDA and Context-Free Languages

Let $L_3 = \{xycy \mid x, y \in \{a, b\}^*, |x| = |y|, \text{ and } x \neq y^R\}$.

(a) Construct a context-free grammar that generates the language L_3 .

(b) Construct a pushdown automata that accepts L_3 .

Solution:

分两讨论: $\begin{cases} \text{外已不同: 随意, 可结束.} \\ \text{外元不同: 第一个不同.} \end{cases}$

(a)
$$\begin{aligned} S &\rightarrow aSa \mid bSb \mid aAb \mid bAa \\ A &\rightarrow aAa \mid bAb \mid aAb \mid bAb \mid \epsilon \end{aligned}$$

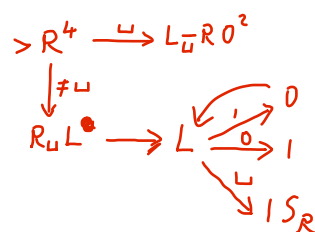
(b) The PDA $M = (K, \Sigma, \Gamma, \Delta, s, F)$ is defined below:

	(q, σ, β)	(p, γ)
$K = \underline{\hspace{2cm}}$		
$\Sigma = \{a, b, c\}$		
$\Gamma = \underline{\hspace{2cm}}$		
$s = \underline{\hspace{2cm}}$		
$F = \underline{\hspace{2cm}}$		

4. (20%) The function $\varphi : \mathbb{N} \rightarrow \mathbb{N}$ given by

$$\varphi(x) = \begin{cases} 4x, & \text{if } x < 8 \\ x + 2, & \text{if } x \geq 8 \end{cases}$$

(a) Try to construct a Turing Machine to compute the function $\varphi(x)$. When describing the Turing machines, you can use the elementary Turing machines described in textbook. Always assume that the Turing machines start computation from the configuration $\triangleright \sqcup x$ where x is represented by binary string, i.e. $x \in \{0, 1\}^*$.



(b) Show that the function $\varphi(x)$ is primitive recursive.

$$\left. \begin{array}{l} 4x = \text{mult}(x, 4) \\ x+2 = \text{plus}(x, 2) \\ \text{iszero}(x+1 \sim 8) \end{array} \right\} \text{p.r.}$$

5. (12%) Consider the language

$$\mathbf{Non-Empty} = \{ \langle M \rangle \mid \text{Turing machine } M \text{ halts on some strings} \}$$

Show that **Non-Empty** is recursively enumerable. Justify your answer, and an informal description suffices.