

# Classical IPC Problems

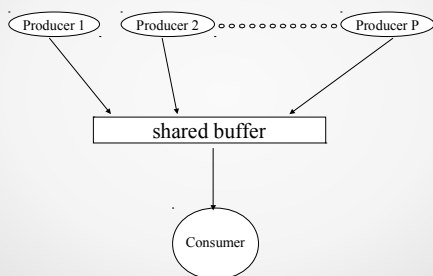
Computer Operating Systems  
BLG 312E

2017-2018 Spring

## Problems

- producer – consumer
- readers – writers
- dining philosophers
- sleeping barber

## Producer - Consumer



## Producer - Consumer

- access to shared buffer through mutual exclusion
- circular buffer
- if buffer empty  $\rightarrow$  consumer waits (synchronization)

## Producer – Consumer

- use counting semaphores
  - takes on  $\geq 0$  integers
  - used when resource capacity  $> 1$
  - initial value = initial free resource capacity
  - P: one more unit of capacity in use
  - V: one unit of capacity freed

## Producer – Consumer

- shared buffer implemented through a shared array of size N
  - array[N]
- binary semaphore:  
mutex  $\leftarrow 1$
- counting semaphores:  
full  $\leftarrow 0$  : number of full buffer locations  
empty  $\leftarrow N$  : number of free **buffer locations**

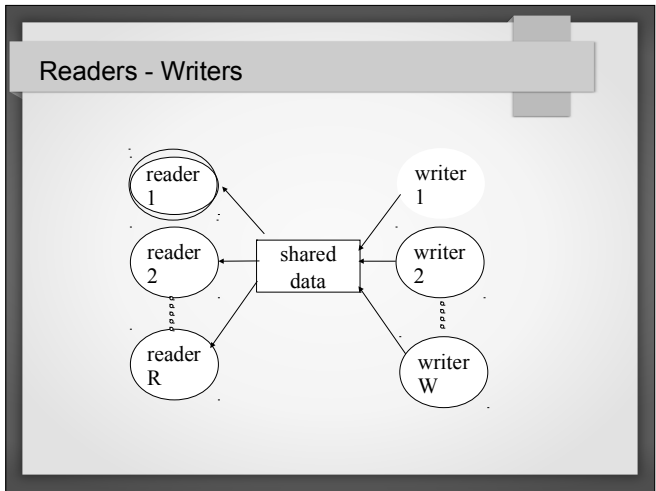
```

constant N=100;
semaphore full=0, empty=N, mutex=1;
item array[N];
int in=0, out=0;
item data;

process producer(){
  while (true) {
    -- produce data --
    p(empty);
    p(mutex);
    array[in]=data;
    in=(in+1)%N;
    v(mutex);
    v(full);
  }
}

process consumer(){
  while (true) {
    p(full);
    p(mutex);
    data=array[out];
    out=(out+1)%N;
    v(mutex);
    v(empty);
    -- use data --
  }
}

```



### Readers - Writers

- more than one reader may read shared data (no writers)
- when a writer uses shared data, all other writers and readers must be excluded

```

int reader=0;
semaphore read_mutex=1, data_mutex=1;

process reader() {
  while (true) {
    p(read_mutex);
    reader=reader+1;
    if (reader==1)
      p(data_mutex);
    v(read_mutex);
    -- read data --
    p(read_mutex);
    reader=reader-1;
    if (reader==0)
      v(data_mutex);
    v(read_mutex);
  }
}

process writer() {
  while (true) {
    -- execute --
    p(data_mutex);
    -- write data --
    v(data_mutex);
  }
}

```

readers have priority over writers!  
Possible indefinite postponement !

### Readers - Writers

- must find a fair solution
- apply rules for access order:
  - if a writer is waiting for readers to be finished, do not allow any more readers
  - if a reader is waiting for a writer to finish, give reader priority

### Dining Philosophers

The diagram illustrates the Dining Philosophers problem. Five philosophers, represented by circles and numbered 1 through 5, are seated around a circular table. Between each pair of adjacent philosophers is a fork, represented by a small line segment. The philosophers are numbered 1, 2, 3, 4, and 5 in clockwise order starting from the top-left. The forks are numbered 1 through 4, with the fifth fork being the one between philosopher 5 and philosopher 1.

**Problem:** share resources (forks) among philosophers without causing deadlock or starvation

## Dining Philosophers

- philosophers
  - eat pasta
  - think
- philosophers need two forks to eat

## Dining Philosophers

- fact: two philosophers sitting side by side cannot eat at the same time
  - e.g. for  $N=5$ , at most 2 philosophers can eat at the same time
- solution must provide maximum amount of parallelism

## Dining Philosophers

```
philosopher(i) {  
  while (true) {  
    think();  
    take_fork(i); //left fork  
    take_fork((i+4)%5); //right fork  
    --- eat ----  
    leave_fork(i);  
    leave_fork((i+4)%5);  
  }  
}
```

what happens if  
all philosophers  
take their left  
forks?

## Dining Philosophers

```
philosopher(i) {  
  while (true) {  
    think();  
    take_fork(i); //left fork  
    if (fork_free((i+4)%5)==FALSE)  
      leave_fork(i);  
    else {  
      take_fork((i+4)%5); //right fork  
      --- eat ----  
      leave_fork(i);  
      leave_fork((i+4)%5);  
    }  
  }  
}
```

is it possible  
that all  
philosophers  
starve?

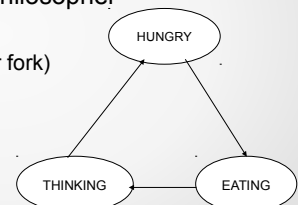
## Dining Philosophers

```
philosopher(i) {  
  while (true) {  
    P(mutex); //binary semaphore  
    think();  
    take_fork(i); //left fork  
    take_fork((i+4)%5); //right fork  
    --- eat ----  
    leave_fork(i);  
    leave_fork((i+4)%5);  
    V(mutex);  
  }  
}
```

at most how  
many  
philosophers can  
eat together?

## Dining Philosophers (Correct Solution)

- $state[i]$  : state of  $i$ th philosopher
  - 0 : THINKING
  - 1 : HUNGRY (wait for fork)
  - 2 : EATING



## Dining Philosophers (Correct Solution)

- a philosopher can be “EATING” only if both neighbors are not “EATING”
- use a binary semaphore per philosopher
  - blocks on semaphore if a fork is not available when requested

## Variables:

- N=5 philosophers
- states:
  - THINKING = 0
  - HUNGRY = 1
  - EATING = 2
- state[5]: array of size 5
- semaphores:
  - mutex ← 1
  - s[5] ← 0 array of size 5

```
process philosopher(i) {
    while (true) {
        think();
        take_fork(i);
        --- eat ---
        leave_fork(i);
    }
}

take_fork(i) {
    P(mutex);
    state[i]=HUNGRY; //request to eat
    try[i]; //try to take forks
    V(mutex);
    P(s[i]); //blocks if can't take forks
}
```

```
leave_fork(i) {
    left=(i+1)%5;
    right=(i+4)%5;
    P(mutex);
    state[i]=THINK;
    try(left);
    try(right);
    V(mutex);
}

try(i) {
    left=(i+1)%5;
    right=(i+4)%5;
    if ((state[i]=HUNGRY) ^
        (state[left]≠EATING) ^
        (state[right]≠EATING))
    {
        state[i]=EATING;
        v(s[i]);
    }
}
```

## Sleeping Barber

- in a barber shop
  - 1 barber
  - 1 customer seat
  - N waiting seats
- barber sleeps if there are no customers
- arriving customer wakes barber up
- if barber is busy when customer arrives
  - waits if waiting seats available
  - leaves if no waiting seats available

## Sleeping Barber

- 3 semaphores needed for the solution
  - customers : number of customers waiting (excluding the one in the customer seat)
  - barbers : number of available barbers (0/1 in this problem)
  - mutex : for mutual exclusion

```
constant CHAIRS=5;
int waiting=0;
semaphore customers=0, barber=0, mutex=1;
```

```
process barber() {
    while(true) {
        P(customers); //sleep if no customers
        P(mutex);
        waiting--; //remove customer
        V(barber); //barber ready to cut hair
        V(mutex);
        -- cut hair --
    }
}
```

```
process customer() {
    P(mutex);
    if (waiting<CHAIRS) { //shop full?
        waiting++; //admit customer
        V(customers); //wake-up barber (possibly)
        V(mutex);
        P(barber); //sleep if barber busy
        -- cut hair --
    }
    else
        V(mutex); //shop is full, so leave
}
```