

COMP3600/6466 – Algorithms

Introduction

[CLRS ch. 1, sec. 2.1, 2.2]

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Topics

✓What is Algorithms?

- Setting up the stage
 - The Problems we'll focus on:
 - Sorting
 - Searching
 - Model of Computation
 - Math refresher
-

The Problems

- Recall: An algorithm is a well-defined procedure to solve well-specified problems
- In this class, we will focus on 2 classes of problems:
 - Sorting
 - Searching

These problems are building blocks of almost any program you see in practice today

Sorting – Problem Specification

- Input: A sequence of n numbers $[a_1, a_2, \dots, a_n]$
- Output: A reordering $[a_1', a_2', \dots, a_n']$ of the input sequence such that $a_1 \leq a_2 \leq \dots \leq a_n$ (for ascending order)

Where is sorting used?

- Often sorting is used as a building block of a bigger program. For instance:
 - Searching. Binary search can run faster once the keys are sorted. Internet search (e.g., Google) also uses sort to present the most relevant results first
 - Closest pair. Given n numbers, which two have the smallest difference between them? Once the numbers are sorted, the closest pair must be next to each other. Therefore, a linear-time scan will find the solution. Closest pair problem is often encountered in various applications, including in machine learning, computational geometry, etc.
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Searching – Problem Specification

- Input: A bounded set X with relation R among its elements, and a solution criteria C
 - Output: An element of X , $x \in X$, that satisfies the criteria C
 - Note: The above definition is very generic. In many problems, we'll be more specific. For instance:
 - What type of space is X ? (e.g., is it countable or uncountable?) What type of relation exists among the elements? Different algorithms are more suitable for different types of X and R .
 - The solution criteria will also influence the type of algorithms we should use to be efficient.
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Searching – A Problem Specification

- An example for a specific search problem that will be used often in this class
- Input: A sequence of n numbers $A = [a_1, a_2, \dots, a_n]$ and a value v
- Output: An index i such that $v = A[i]$ or the special value null if v does not appear in A

Where is searching used?

- More visible than sorting. For instance:
 - Internet search, search in your email, search in a document, etc.
 - Optimization. Computation wise, optimization problem is a search problem (i.e., finding a value that is the largest/smallest). Many are search in continuous space (e.g., in real number space). Since almost (if not all) problems can be framed as an optimization problem, they will eventually become a search problem 😊
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A bit more about problems

- Recall: An algorithm solves well-specified problems
 - In reality, problems do not come in well-specified form. Someone (usually the algorithm designer) needs to formulate the problem into a well-specified problem that an algorithm can solve
 - In fact, good problem formulation is usually half the solution
 - In this class, esp. from tutorial questions and assignments, we'll learn how to formulate problems too
 - Yes, you'll need to do this in your assignments 😊
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Analysing Algorithms

- Recall that analysing algorithms mean computing computational resources (typically, time and space) required to run the algorithm on a certain input
 - To do the above analysis, we need to first know the machine where the algorithms are executed
 - Model of computation
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Model of Computation

- In this class, we'll assume:
 - An abstract generic one-processor Random Access Machine (RAM)
 - Abstract: We don't consider memory swapping, garbage collection, etc.
 - One-processor RAM: Instructions are executed one after another (no concurrent operations)
 - To store numbers, RAM has integer and float data type. Each data type in RAM has a limit. We will introduce as necessary
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Model of Computation

- Have the following primitive instructions (each primitive instruction takes constant time)
 - Arithmetic: Add, subtract, multiply, divide, mod, floor, ceil
 - Data movement: Load, store, copy
 - Control: Conditional & unconditional branching, subroutine call and return
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Analysing Algorithms Executed in a RAM:

Example on Time Analysis

- Sorting problem:
 - Input: A sequence of n numbers $[a_1, a_2, \dots, a_n]$
 - Output: A reordering $[a_1', a_2', \dots, a_n']$ of the input sequence such that $a_1 \leq a_2 \leq \dots \leq a_n$

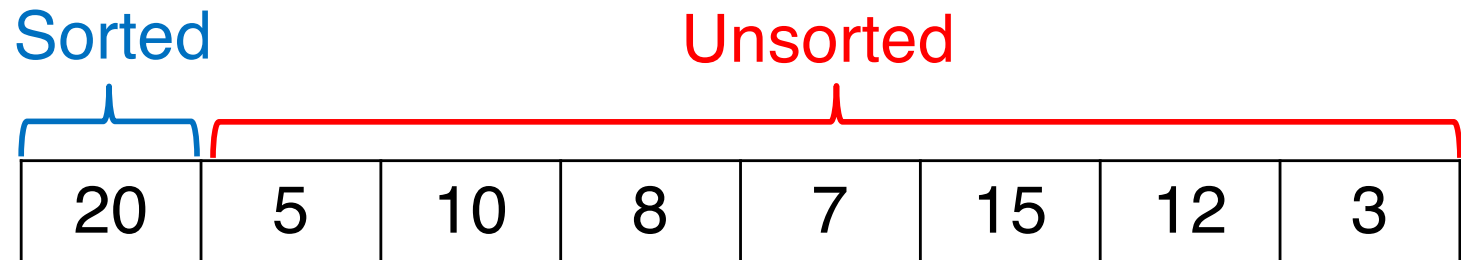
- InsertionSort(A)

1. for $j = 2$ to $A.length$
2. $Key = A[j]$
3. $i = j - 1$
4. While $i > 0$ and $A[i] > key$
5. $A[i + 1] = A[i]$
6. $i = i - 1$
7. $A[i + 1] = key$

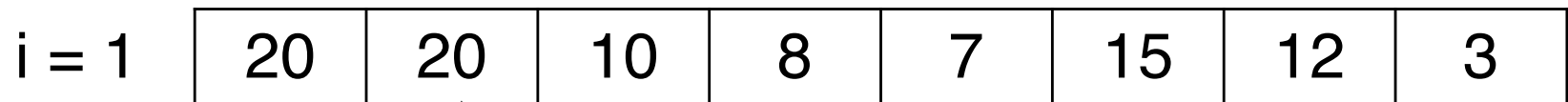
Intuitively:

- Separate the sequence of numbers / array into: Sorted and unsorted.
- At each iteration take an element from the unsorted part and put it in the right place in the sorted part of the sequence

Illustration of Insertion Sort



Key = 5, j = 2



Key = 5, j = 2

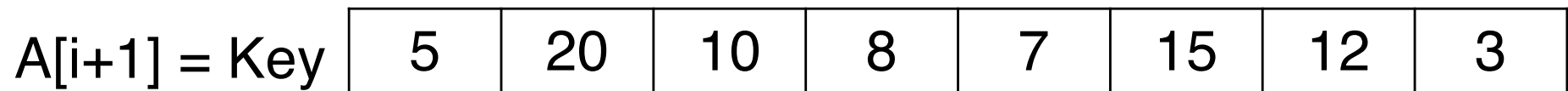
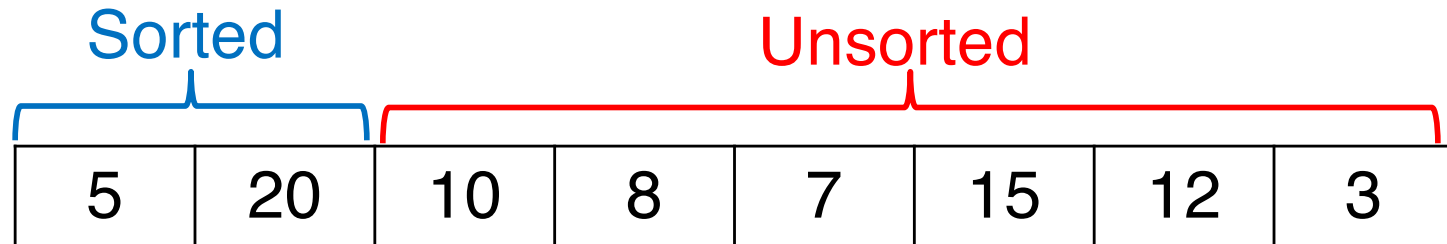


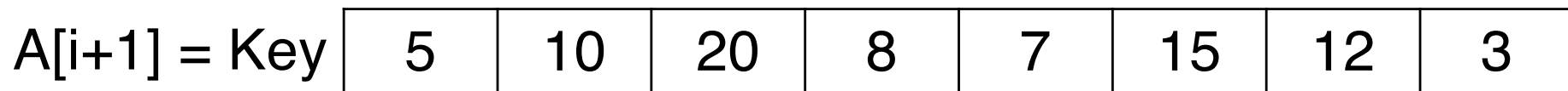
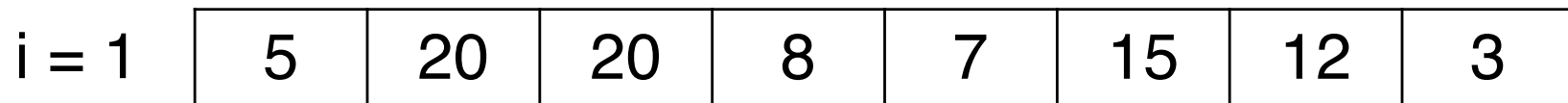
Illustration of Insertion Sort



Key = 10, j = 3



Key = 10, j = 3



Correctness? We'll discuss after Asymptotic Analysis

Back to time analysis

InsertionSort		Cost	Times
1	for j = 2 to A.length		
2	Key = A[j]		
3	i = j-1		
4	While i > 0 and A[i] > key		
5	A[i+1] = A[i]		
6	i = i-1		
7	A[i+1] = key		

Total time $T(n)$: sum of cost X times

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Important Summation Formulas

1. $\sum_{i=l}^u 1 = \underbrace{1 + 1 + \cdots + 1}_{u-l+1 \text{ times}} = u - l + 1$ (l, u are integer limits, $l \leq u$); $\sum_{i=1}^n 1 = n$ _____

2. $\sum_{i=1}^n i = 1 + 2 + \cdots + n = \frac{n(n+1)}{2} \approx \frac{1}{2}n^2$

3. $\sum_{i=1}^n i^2 = 1^2 + 2^2 + \cdots + n^2 = \frac{n(n+1)(2n+1)}{6} \approx \frac{1}{3}n^3$

4. $\sum_{i=1}^n i^k = 1^k + 2^k + \cdots + n^k \approx \frac{1}{k+1}n^{k+1}$

5. $\sum_{i=0}^n a^i = 1 + a + \cdots + a^n = \frac{a^{n+1} - 1}{a - 1}$ ($a \neq 1$); $\sum_{i=0}^n 2^i = 2^{n+1} - 1$

6. $\sum_{i=1}^n i2^i = 1 \cdot 2 + 2 \cdot 2^2 + \cdots + n2^n = (n-1)2^{n+1} + 2$

7. $\sum_{i=1}^n \frac{1}{i} = 1 + \frac{1}{2} + \cdots + \frac{1}{n} \approx \ln n + \gamma$, where $\gamma \approx 0.5772 \dots$ (Euler's constant)

8. $\sum_{i=1}^n \lg i \approx n \lg n$ _____

Source: [Lev] Appendix A

Properties of logarithms

1. $\log_a 1 = 0$

2. $\log_a a = 1$

3. $\log_a x^y = y \log_a x$

4. $\log_a xy = \log_a x + \log_a y$

5. $\log_a \frac{x}{y} = \log_a x - \log_a y$

6. $a^{\log_b x} = x^{\log_b a}$

A function where the exponent is log can be transformed into a polynomial function

7. $\log_a x = \frac{\log_b x}{\log_b a} = \log_a b \log_b x$

In this class, unless otherwise stated, the base of the log is 2, i.e.: $\log x = \log_2 x$

Examples

- Can we transform $2^{3 \log \sqrt{n}}$ into a polynomial form?
- Can we transform $n^{\sqrt{n}}$ into the form of a power of 2?

You need to be comfortable with the basic math (basic transformations as in previous slides, basic arithmetic and algebra + basic calculus + basic probability).

If not, please catch up **NOW!!!**

To catch up, see the notes of Week-1 in Wattle

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Next: Asymptotic Analysis
