COMP2610 / COMP6261 Information Theory
Lecture 11: Entropy and Coding

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Brief Recap of Course (Last 5 Weeks)

- How can we quantify information?
  - Basic Definitions and Key Concepts
  - Probability, Entropy & Information
- How can we make good guesses?
  - Probabilistic Inference
  - Bayes Theorem
- How much redundancy can we safely remove?
  - Compression
  - Source Coding Theorem, Kraft Inequality
  - Block, Huffman, and Lempel-Ziv Coding
- How much noise can we correct and how?
  - Noisy-Channel Coding
  - Repetition Codes, Hamming Codes
- What is randomness?
  - Kolmogorov Complexity
  - Algorithmic Information Theory

Brief Overview of Course (Next 6 Weeks)

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This time

Basic goal of compression

Key concepts: codes and their types, raw bit content, essential bit content

Informal statement of source coding theorem

What is Compression?

Cn y rd ths mssg wtht ny vwls?

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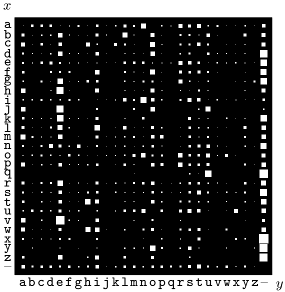
Cn y rd ths mssg wtht ny vwls?

It is not too difficult to read as there is redundancy in English text.
(Estimates of 1-1.5 bits per character, compared to $\log_2 26 \approx 4.7$)



(a) $P(y \mid x)$

- If you see a “q”, it is very likely to be followed with a “u”
- The letter “e” is much more common than “j”
- Compression exploits differences in relative probability of symbols or blocks of symbols

Compression in a Nutshell

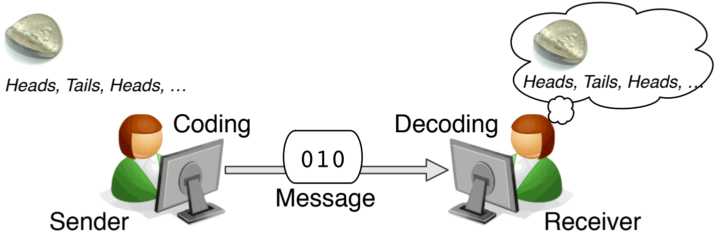
Compression

Data compression is the process of replacing a message with a smaller message which can be reliably converted back to the original.

A General Communication Game

Imagine the following game between Sender & Receiver:

- Sender & Receiver agree on code for each outcome ahead of time (e.g., `0` for *Heads*; `1` for *Tails*)
- Sender observes outcomes then codes and sends message
- Receiver decodes message and recovers outcome sequence



Goal: Want small messages on average when outcomes are from a fixed, known, but uncertain source (e.g., coin flips with known bias)

Sneak peek: source coding theorem

Consider a coin with $P(Heads) = 0.9$. If we want perfect transmission:

- Coding single outcomes requires 1 bit/outcome
- Coding 10 outcomes at a time needs 10 bits, or 1 bit/outcome

Not very interesting!

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- accept errors in transmission **(This Week)**
- allow variable length messages (Next week)

Sneak peek: source coding theorem

If we are happy to fail on up to 2% of the sequences we can ignore any sequence of 10 outcomes with more than 3 tails

Why? The number of tails follows a Binomial(10, 0.1) distribution

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Why? The number of tails follows a Binomial(10, 0.1) distribution

There are only $176 < 2^8$ sequences with 3 or fewer tails

So, we can just code those, and **ignore** the rest!

- Coding 10 outcomes with 2% failure doable with 8 bits, or 0.8 bits/outcome
- *Smallest bits/outcome needed for 10,000 outcome sequences?*

Generalisation: Source Coding Theorem

What happens when we generalise to arbitrary error probability, and sequence size?

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Source Coding Theorem (Informal Statement)

If: you want to uniformly code large sequences of outcomes with any degree of reliability from a random source

Then: the average number of bits per outcome you will **need** is roughly equal to the entropy of that source.

To define: "Uniformly code", "large sequences", "degree of reliability", "average number of bits per outcome", "roughly equal"

Entropy and Information: Recap

Ensemble

An **ensemble** X is a triple $(x, \mathcal{A}_X, \mathcal{P}_X)$; x is a **random variable** taking **values** in $\mathcal{A}_X = \{a_1, a_2, \ldots, a_I\}$ with **probabilities** $\mathcal{P}_X = \{p_1, p_2, \ldots, p_I\}$.

Information

The **information** in the observation that $x = a_i$ (in the ensemble X) is

$$h(a_i) = \log_2 \frac{1}{p_i} = -\log_2 p_i$$

Entropy

The **entropy** of an ensemble X is the average information

$$H(X) = \mathbb{E}[h(x)] = \sum_i p_i h(a_i) = \sum_i p_i \log_2 \frac{1}{p_i}$$

What is a Code?

A source code is a process for assigning names to outcomes. The names are typically expressed by strings of binary symbols.

We will denote the set of all finite binary strings by

$$\{0,1\}^+ \stackrel{\text{def}}{=} \{0,1,00,01,10,\ldots\}$$

Source Code

Given an ensemble X, the function $c : \mathcal{A}_X \to \{0,1\}^+$ is a **source code** for X. The number of symbols in $c(x)$ is the **length** $l(x)$ of the codeword for x. The **extension** of c is defined by $c(x_1 \ldots x_n) = c(x_1) \ldots c(x_n)$

Example:

- The code c names outcomes from $\mathcal{A}_X = \{\mathtt{r}, \mathtt{g}, \mathtt{b}\}$ by $c(\mathtt{r}) = \mathtt{00}$, $c(\mathtt{g}) = \mathtt{10}$, $c(\mathtt{b}) = \mathtt{11}$
- The length of the codeword for each outcome is 2.
- The extension of c gives $c(\mathtt{rgrb}) = \mathtt{00100011}$

Types of Codes

Let X be an ensemble and $c : \mathcal{A}_X \to \{0,1\}^+$ a code for X. We say c is a:

- Uniform Code if $l(x)$ is the same for all $x \in \mathcal{A}_X$ 长度相同.
- Variable-Length Code otherwise

Another important criteria for codes is whether the original symbol x can be unambiguously determined given $c(x)$. We say c is a:

- Lossless Code if for all $x_1, x_2 \in \mathcal{A}_X$ we have $x_1 \neq x_2$ implies $c(x_1) \neq c(x_2)$ 单一映射.
- Lossy Code otherwise

Types of Codes

Examples

Examples: Let $\mathcal{A}_X = \{\mathtt{a}, \mathtt{b}, \mathtt{c}, \mathtt{d}\}$

1. $c(\mathtt{a}) = 00$, $c(\mathtt{b}) = 01$, $c(\mathtt{c}) = 10$, $c(\mathtt{d}) = 11$ is uniform lossless
2. $c(\mathtt{a}) = 0$, $c(\mathtt{b}) = 10$, $c(\mathtt{c}) = 110$, $c(\mathtt{d}) = 111$ is variable-length lossless
3. $c(\mathtt{a}) = 0$, $c(\mathtt{b}) = 0$, $c(\mathtt{c}) = 110$, $c(\mathtt{d}) = 111$ is variable-length lossy
4. $c(\mathtt{a}) = 00$, $c(\mathtt{b}) = 00$, $c(\mathtt{c}) = 10$, $c(\mathtt{d}) = 11$ is uniform lossy
5. $\boxed{c(\mathtt{a}) = -,\ c(\mathtt{b}) = -,\ c(\mathtt{c}) = 10,\ c(\mathtt{d}) = 11 \text{ is uniform lossy}}$

undefined.

A Note on Lossy Codes & Missing Codewords

When talking about a uniform lossy code c for $\mathcal{A}_X = \{\mathtt{a}, \mathtt{b}, \mathtt{c}\}$ we write

$$c(\mathtt{a}) = \mathtt{0} \quad c(\mathtt{b}) = \mathtt{1} \quad c(\mathtt{c}) = \texttt{-}$$

where the symbol - means "no codeword". This is shorthand for "the receiver will decode this codeword incorrectly"

For the purposes of these lectures, this is equivalent to the code

$$c(\mathtt{a}) = \mathtt{0} \quad c(\mathtt{b}) = \mathtt{1} \quad c(\mathtt{c}) = \mathtt{1}$$

and the sender and receiver agreeing that the codeword `1` should always be decoded as `b`

Assigning the outcome $\mathtt{a}_i$ the missing codeword "-" just means "it is not possible to send $\mathtt{a}_i$ reliably"

Lossless Coding

Example: Colours

Three colour ensemble with $\mathcal{A}_X = \{\mathtt{r}, \mathtt{g}, \mathtt{b}\}$ with $\mathtt{r}$ twice as likely as $\mathtt{b}$ or $\mathtt{g}$

- $p_{\mathtt{r}} = 0.5$ and $p_{\mathtt{g}} = p_{\mathtt{b}} = 0.25$.

Suppose we use the following **uniform lossless** code

$$c(\mathtt{r}) = 00;\ c(\mathtt{g}) = 10;\ \text{and}\ c(\mathtt{b}) = 11$$

For example $c(\mathtt{rrgbrbr}) = 00001011001100$ will have 14 bits.

On average, we will use $l(\mathtt{r})p_{\mathtt{r}} + l(\mathtt{g})p_{\mathtt{g}} + l(\mathtt{b})p_{\mathtt{b}} = 2$ bits per outcome

- $2N$ bits to code a sequence of N outcomes

Raw Bit Content

Uniform coding gives a crude measure of information: the number of bits required to assign equal length codes to each symbol

Raw Bit Content

If X is an ensemble with outcome set $\mathcal{A}_X$ then its **raw bit content** is

$$H_0(X) = \log_2 |\mathcal{A}_X|.$$

x	$c(x)$
a	000
b	001
c	010
d	011
e	100
f	101
g	110
h	111

Example:
This is a uniform encoding of outcomes in $\mathcal{A}_X = \{\mathtt{a}, \mathtt{b}, \mathtt{c}, \mathtt{d}, \mathtt{e}, \mathtt{f}, \mathtt{g}, \mathtt{h}\}$:

- Each outcome is encoded using $H_0(X) = 3$ bits
- The probabilities of the outcomes are ignored
- Same as assuming a uniform distribution

For the purposes of compression, the exact codes don't matter – just the number of bits used.

Lossy Coding

Example: Colours

Three colour ensemble with $\mathcal{A}_X = \{\mathtt{r}, \mathtt{g}, \mathtt{b}\}$

- $p_{\mathtt{r}} = 0.5$ and $p_{\mathtt{g}} = p_{\mathtt{b}} = 0.25$.

Using uniform lossy code:

- $c(\mathtt{r}) = 0$; $c(\mathtt{g}) = -$; and $c(\mathtt{b}) = 1$

ignored

Examples:

$c(\mathtt{rrrrrrr}) = 0000000$; $c(\mathtt{rrbbrbr}) = 0011010$; $c(\mathtt{rrgbrbr}) = -$

不可被解析

What is probability we can reliably code a sequence of N outcomes?

Given we can code a sequence of length N, how many bits are expected?

Lossy Coding

Example: Colours

What is probability we can reliably code a sequence of N outcomes?

$$P(x_1 \dots x_N \text{ has no } \mathtt{g}) = P(x_1 \neq \mathtt{g}) \dots P(x_N \neq \mathtt{g}) = (1 - p_{\mathtt{g}})^N$$

Given we can code a sequence of length N, how many bits are expected?

$$\mathbb{E}[I(X_1) + \cdots + I(X_N) | X_1 \neq \mathtt{g}, \dots, X_N \neq \mathtt{g}] = \sum_{n=1}^{N} \mathbb{E}[I(X_n) | X_n \neq \mathtt{g}]$$

$$= N \left(I(\mathtt{r}) p_{\mathtt{r}} + I(\mathtt{b}) p_{\mathtt{b}} \right) / (1 - p_{\mathtt{g}}) = N = N \log_2 |\{\mathtt{r}, \mathtt{b}\}|$$

since $I(p_{\mathtt{r}}) = I(p_{\mathtt{b}}) = 1$ and $p_{\mathtt{r}} + p_{\mathtt{b}} = 1 - p_{\mathtt{g}}$.

- c.f. $2N$ bits with lossless code

Essential Bit Content

There is an inherent trade off between the number of bits required in a uniform lossy code and the probability of being able to code an outcome

Smallest δ-sufficient subset

Let X be an ensemble and for $0 \le \delta \le 1$, define S_δ to be the smallest subset of $\mathcal{A}_X$ such that

$$P(x \in S_\delta) \ge 1 - \delta$$

For small δ, **smallest** collection of **most likely** outcomes

If we uniformly code elements in S_δ, and ignore all others:

- We can code a sequence of length N with probability $(1-\delta)^N$
- If we can code a sequence, its expected length is $N \log_2 |S_\delta|$

Essential Bit Content

Example

Intuitively, construct S_δ by removing elements of X in ascending order of probability, till we have reached the $1-\delta$ threshold

x	$P(\mathbf{x})$
a	$1/4$
b	$1/4$
c	$1/4$
d	$3/16$
e	$1/64$
f	$1/64$
g	$1/64$
h	$1/64$

- Outcomes ranked (high - low) by $P(x = a_i)$ removed to make set S_δ with $P(x \in S_\delta) \geq 1 - \delta$

$$\delta = 0 \ : S_\delta = \{a\ ,b,\ c,\ d,\ e,\ f,\ g,\ h\}$$

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- Outcomes ranked (high - low) by $P(x = a_i)$ removed to make set S_δ with $P(x \in \mathrm{S}_\delta) \geq 1 - \delta$

$$\delta = 0 \ : \mathrm{S}_\delta = \{\text{a ,b, c, d, e, f, g, h}\}$$
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$$\delta = 1/64 \ : S_\delta = \{a, b, c, d, e, f, g\}$$
$$\delta = 1/16 \ : S_\delta = \{a, b, c, d\}$$
$$\delta = 3/4 \ : S_\delta = \{a\}$$

Essential Bit Content

Trade off between a probability of δ of not coding an outcome and size of uniform code is captured by the essential bit content

Essential Bit Content

For an ensemble X and $0 \leq \delta \leq 1$, the **essential bit content** of X is

$$H_\delta(X) \stackrel{\text{def}}{=} \log_2 |S_\delta|$$

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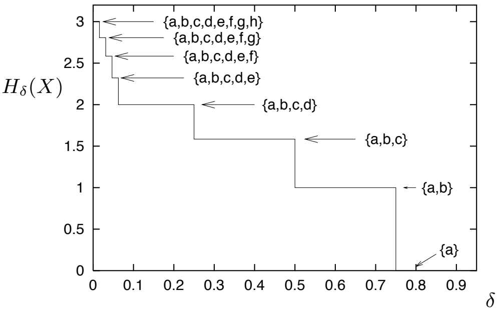

The Source Coding Theorem for Uniform Codes

(Theorem 4.1 in MacKay)

Our aim next time is to understand this:

The Source Coding Theorem for Uniform Codes

Let X be an ensemble with entropy $H = H(X)$ bits. Given $\epsilon > 0$ and $0 < \delta < 1$, there exists a positive integer N_0 such that for all $N > N_0$

$$\left| \frac{1}{N} H_\delta\left(X^N\right) - H \right| < \epsilon.$$

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What?

- The term $\frac{1}{N} H_\delta(X^N)$ is the average number of bits required to uniformly code all but a proportion δ of the symbols.
- Given a tiny probability of error δ, the average bits per symbol can be made as close to H as required.
- Even if we allow a large probability of error we cannot compress more than H bits ber symbol.

Some Intuition for the SCT

- Don't code individual symbols in an ensemble; rather, consider sequences of length N.
- As length of sequence increases, the probability of seeing a "typical" sequence becomes much larger than "atypical" sequences.
- Thus, we can get by with essentially assigning a unique codeword to each typical sequence

Next time

Recap: typical sets

Formal statement of source coding theorem

Proof of source coding theorem