

Design and Analysis of Algorithms I

Introduction Merge Sort (Pseudocode)

Merge Sort: Pseudocode

- -- recursively sort 1st half of the input array
- -- recursively sort 2nd half of the input array
- -- merge two sorted sublists into one

[ignores base cases]

Pseudocode for Merge:

```
C = output [length = n]

A = 1^{st} sorted array [n/2]

B = 2^{nd} sorted array [n/2]

i = 1

j = 1
```

for
$$k = 1$$
 to n
if $A(i) < B(j)$
 $C(k) = A(i)$
 $i++$
else $[B(j) < A(i)]$
 $C(k) = B(j)$
 $j++$
end
(ignores end cases)

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Merge Sort Running Time?

Key Question: running time of Merge Sort on array of n numbers?

[running time ~ # of lines of code executed]

Pseudocode for Merge:

```
C = output [length = n]

A = 1<sup>st</sup> sorted array [n/2]

B = 2<sup>nd</sup> sorted array [n/2]

i = 1

j = 1
```

for
$$k = 1$$
 to n

if $A(i) < B(j)$

$$C(k) = A(i) - i + + - i + - i$$

else $[B(j) < A(i)]$

$$C(k) = B(j) - i + - i + - i$$

end

(ignores end cases)

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Running Time of Merge

Upshot: running time of Merge on array of m numbers is $\leq 4m+2$

$$\leq 6m$$
 (Since $m \geq 1$)

一共有logn个C array 所以最后的 merge sort operations = 6n * logn + 6n

Running Time of Merge Sort

Claim: Merge Sort requires

 $\leq 6n\log_2 n + 6n$ operations

to sort n numbers.

Recall : $= \log_2 n$ is the # of times you divide by 2 until you get down to 1

