

Assignment 4

1. In Sets^2 , the subobject classifier is the functor

$$\Delta_2 : \text{Sets} \rightarrow \text{Sets}, \quad 0 \mapsto \{\{0, 1\}, \{1\}\}, \quad 1 \mapsto \{\{1\}\}$$

with natural transformation

$t : \Delta_2 \rightarrow \Delta_2$ such that

$$t_0(x) = \{0, 1\}, \quad t_1(x) = \{1\}$$

where $* = 1$ is the only element in the singleton set 1 .

and Δ_1 denotes the constant functor that maps to $1\!\!1$.

Now prove that it indeed forms a subobject classifier.

Proof: For any Sets^2 -object $F : 2 \rightarrow \text{Sets}$ and its subobject U with

$m : U \rightarrow F$, let $\alpha : U \rightarrow \Delta_1$ be the unique morphism towards the terminal object.

Define $w : F \rightarrow \Delta_2$ such that

$$w_0 : F(0) \rightarrow \Delta_2(0), \quad w_0(a) = \{0, 1\} \text{ if } a \in m_0(U(0)), \quad \{1\} \text{ otherwise}$$

$$w_1 : F(1) \rightarrow \Delta_2(1), \quad w_1(a) = \{1\} \text{ for all } a \in F(1)$$

Then for some $a \in U(0)$, $b \in U(1)$

$$(t \circ w)_0(a) = t_0(w_0(a)) = t_0(*) = \{0, 1\} = m_0(m_0(a))$$

$$(t \circ w)_1(b) = t_1(w_1(b)) = t_1(*) = \{1\} = m_1(m_1(b))$$

Thus $t \circ w = u \circ m$.

$$\begin{array}{ccccc} U & \xrightarrow{\alpha} & \Delta_1 & & \\ \downarrow r & \nearrow s & \downarrow t & & \\ m & \downarrow \phi & G & \xrightarrow{u} & \Delta_2 \\ F & \xrightarrow{w} & \Delta_2 & \xrightarrow{v} & \Delta_2 \end{array}$$

Then for all Sets^2 -object G with $\phi : G \rightarrow F$ and $\psi : G \rightarrow \Delta_2$, such that

$$u \circ \phi = t \circ \psi$$

It follows that $\phi(G(0)) \subseteq m_0(U(0))$

Therefore for all $x \in G(0)$, $\phi(x) \in \phi(G(0)) \subseteq m_0(U(0))$. This shows $x \in U(0)$ that

$$\phi(x) = m_0(y)$$

Thus define $v : G \rightarrow U$ such that $v_0(x) = y$ with $\phi(x) = m_0(y)$.

$$\text{Then } (m \circ v)_0(x) = m_0(v_0(x)) = m_0(y) = \phi(x)$$

$$\Rightarrow (m \circ v)_0 = \phi$$

Suppose $! : 0 \rightarrow 1$ in part 2. By naturality,

$$v_1 \circ G(1) = U(1) \circ v_0, \quad m_1 \circ U(1) = F(1) \circ m_0, \quad \phi_1 \circ G(1) = F(1) \circ \phi_0$$

It follows that $m_1 \circ v_1 \circ G(1) = F(1) \circ m_0 \circ v_0$

$$\Rightarrow m_1 \circ v_1 \circ G(1) = F(1) \circ \phi_0$$

$$\Rightarrow m_1 \circ v_1 \circ G(1) = \phi_1 \circ G(1)$$

As functors on Sets preserve epimorphisms, we have

$$m_1 \circ v_1 = \phi_1$$

Therefore $m \circ v = \phi$. And clearly $\alpha \circ v = \tau$ as both sides map to $*$.

Additionally, the uniqueness of v follows from m is monic. Hence such structure actually forms a pullback.

Now suppose $w : F \rightarrow \Delta_2$ also makes it a pullback, and $w \neq u$. Then

$$w \circ m = t \circ \alpha$$

Hence $w_0(x) = \{1\}$ for all $x \in F(0)$,

$$w_1(x) = \{0, 1\} \text{ if } x \in R, \quad \{1\} \text{ otherwise, where } m_0(U(0)) \not\subseteq R \subseteq F(0)$$

Since $m_0(U(0)) \not\subseteq R$, pick $r \in R, r \notin m_0(U(0))$, $w_1(r) = \{1\}$.

For some Sets^2 -object G with $\phi : G \rightarrow F$ and $\psi : G \rightarrow \Delta_2$, such that

$$\phi_0(a) = r \quad \text{and} \quad w \circ \phi = t \circ \psi$$

Then there is unique $v : G \rightarrow U$ that $m \circ v = \phi$, $\alpha \circ v = \tau$.

However $\phi_0(a) = r \notin m_0(U(0))$, yielding a contradiction.

Thus $w = u$, proved uniqueness.

2. Proof: Suppose $\phi: \text{Hom}_D(F(-), *) \rightarrow \text{Hom}_C(-, U(*))$ is a natural isomorphism, and define

$$\eta_C = \phi_{C, F(C)}(1_{F(C)})$$

For all $f: X \rightarrow Y$ in C , by naturality of ϕ , we have

$$\eta_Y \circ 1_C(f) = \phi_{Y, F(Y)}(1_{F(Y)}) \circ f = UF(f) \circ \phi_{X, F(X)}(1_{F(X)}) = UF(f) \circ \eta_X$$

Thus $\eta: I_C \Rightarrow UF$ indeed forms a natural transformation.

Consider arbitrary C -object C and D -object D with $f: C \rightarrow U(D)$.

Since ϕ is a natural isomorphism, $\phi_{C, D}$ is a bijection, thus there exists

unique $g: F(C) \rightarrow D$ such that

$$\phi_{C, D}(g) = f$$

By naturality, we have

$$\phi_{C, D}(g) = U(g) \circ \phi_{C, F(C)}(1_{F(C)}) = U(g) \circ \eta_C$$

It follows that $f = U(g) \circ \eta_C$. the required UMP holds.

3. Proof: Consider the natural transformation

$$g: 1_{P(A)} \rightarrow f^* \circ f$$

Clearly this is well-defined as $X \subseteq (f^* \circ f)(x)$ for all $x \in P(A)$.

Then for all $X \in P(A)$, $Y \in P(B)$ with $h: X \rightarrow f^*(Y)$, by definition of poset category, it is equivalent to

$$X \subseteq f^*(Y)$$

$$\Leftrightarrow \forall x \in X, x \in f^*(Y)$$

$$\Leftrightarrow \forall x \in X, f(x) \in Y$$

$$\Leftrightarrow \forall f(x) \in f(X), f(x) \in Y$$

$$\Leftrightarrow f(X) \subseteq Y$$

That is there uniquely exists $g: f(X) \rightarrow Y$ by definition of poset.

Thus $f^*g: f^*f(X) \rightarrow f^*(Y)$ exists uniquely, hence the following diagram commutes:

$$\begin{array}{ccc} f^*f(X) & \xleftarrow{\quad h \quad} & X \\ & \searrow f^*g & \downarrow h \\ & f^*(Y) & \end{array}$$

It follows that $f \dashv f^*$.

4. Proof: Consider $P: \text{Sets}^{\text{op}} \rightarrow \text{Sets}$ and $P': \text{Sets}^{\text{op}} \rightarrow \text{Sets}$.

Then for all sets A, B , since $P(X) \cong 2^X$ for all set X , we have

$$\text{Hom}_{\text{Sets}}(A, P(B)) \cong \text{Hom}_{\text{Sets}}(A, 2^B)$$

Since in Sets , $I \times A \cong A$, by proposition 8.13 in Awodey's, we have

$$\text{Hom}_{\text{Sets}}(A, 2^B) \cong \text{Hom}_{\text{Sets}}(I \times A, 2^B)$$

$$\cong \text{Hom}_{\text{Sets}}(I \times B, 2^A)$$

$$\cong \text{Hom}_{\text{Sets}}(B, 2^A)$$

$$\cong \text{Hom}_{\text{Sets}}(B, P(A))$$

By definition of opposite category, we have $\text{Hom}_{\text{Sets}}(B, P(A)) \cong \text{Hom}_{\text{Sets}}(P(A), B)$ naturally.

Therefore,

$$\text{Hom}_{\text{Sets}}(A, P(B)) \cong \text{Hom}_{\text{Sets}}(P(A), B)$$

yielding that $P \dashv P'$

5. Proof: Suppose $i: P(I) \rightarrow \text{Sets}/I$, $U \mapsto i(U): U \mapsto I$

First show that i indeed forms a functor.

For arbitrary functions $f: U \rightarrow V$, $g: V \rightarrow W$ in $P(I)$,

by definition of slice category, $i(g): i(V) \rightarrow i(W)$ satisfies

$$i(W) \circ i(g) = i(V)$$

$$\text{Similarly } i(V) \circ i(f) = i(U)$$

$$\text{And } i(W) \circ i(g \circ f) = i(U)$$

$$\text{Hence } i(W) \circ i(g \circ f) = i(W) = i(V) \circ i(f) = i(W) \circ i(g) \circ i(f)$$

Since W is a $P(I)$ -object, $W \in I$. Thus $i(W): W \rightarrow I$ is an injection,

i.e. monomorphism in the category Sets .

Therefore $i(W) \circ i(g \circ f) = i(W) = i(g) \circ i(f)$ implies $i(g \circ f) = i(g) \circ i(f)$.

Additionally, $i(1_U)$ satisfies $i(U) \circ i(1_U) = i(U)$. That is $i(1_U) = 1_{i(U)}$.

Hence i is a functor.

Now let $\sigma: \text{Sets}/I \rightarrow P(I)$, $(f: U \rightarrow I) \mapsto \text{im}_f$.

and for any Sets/I -morphism $g: (f: U \rightarrow I) \rightarrow (f': V \rightarrow I)$, define

$\sigma(g): \text{im}_f \rightarrow \text{im}_{f'}$ be the inclusion map.

Clearly it is well defined because $f \circ g = f$ implies that $\text{im}_f \subseteq \text{im}_{f'}$, and

$$\sigma(1_f) = 1_{\text{im}_f} = 1_{\sigma(f)}.$$

For all $h: (f_0: U \rightarrow I) \rightarrow (f_1: V \rightarrow I)$, $g: (f_1: V \rightarrow I) \rightarrow (f_2: W \rightarrow I)$,

$$\sigma(g \circ h) = \sigma(g) \circ \sigma(h)$$

as $\sigma(g \circ h)$, $\sigma(g)$, $\sigma(h)$ are all inclusion maps.

Hence σ is a functor.

Now let $j: \text{Sets}/I \rightarrow i \circ \sigma$ be a natural transformation.

For all Sets/I -object $f: U \rightarrow I$ and $V \in P(I)$ with $k: f \rightarrow i(V)$ in Sets/I ,

$$i(V) \circ k = f \quad \text{in } \text{Sets}$$

$$\text{Thus } \text{im}_f = \text{im}(i(V) \circ k) \subseteq \text{im}(i(V)) = V$$

Hence there is a unique $P(I)$ -morphism $g: \text{im}_f \rightarrow V$, meeting following

diagram commutes:

$$\begin{array}{ccc} \text{im}_f & \xrightarrow{i(g)} & i(V) \\ \downarrow j_f & \swarrow & \downarrow k \\ f & & \end{array}$$

Therefore $\sigma \dashv i$.

6. Proof: Let $U: \text{Cat} \rightarrow \text{Graphs}$ be the forgetful functor

Fix some graph G . Let $\{\circ\}$ be the solution set, where \circ is the empty category.

For all small category C with graph homomorphism $f: G \rightarrow UC$, we have the following diagram commutes:

$$\begin{array}{ccc} G & \xrightarrow{\varphi} & \mathcal{U}(\circ) \\ f \searrow & & \downarrow U(!) \\ & & U(C) \end{array}$$

where $!: \circ \rightarrow C$ is the unique morphism, φ is the trivial graph homomorphism from G to the empty graph.

Additionally as U is clearly continuous, by adjoint functor theorem, we have U admits a left adjoint functor F .

By UMP definition of adjunction, $F: \text{Graphs} \rightarrow \text{Cat}$ exactly produces free categories by UMP of free category.

That is, free category exists.

7. Claim that $\Delta \dashv \text{ver}\lim$ where 'ver' extracts the vertex of some cone.

Proof: Consider the natural transformation

$$\phi: \text{Hom}(\Delta(-), *) \rightarrow \text{Hom}(-, \text{ver}\lim(*))$$

For all C -object X and C^J -object D , we have

$$\phi_{X,D}: \text{Hom}(\Delta(X), D) \rightarrow \text{Hom}(X, \text{ver}\lim(D))$$

By definition of diagonal functor Δ , $\Delta(X)$ is the constant functor $\Delta(X): J \rightarrow C$ such that $\Delta(X)(j) = X$ for all indices j .

Therefore, the only natural transformation $\alpha \in \text{Hom}(\Delta(X), D)$ is just the family of morphisms of the cone of type J at object X .

By definition of limit, $\lim D$ is the terminal object in $\text{Cone}(D)$.

Thus there is only one morphism from the cone at X to $\lim D$,

by definition of cone morphisms, correspondingly there's only one morphism from X to $\text{ver}\lim(D)$.

Hence $\phi_{X,D}$ is a one-to-one (injective) map for all X and D .

Therefore ϕ forms a natural isomorphism and $\Delta \dashv \text{ver}\lim$. \square

Thus, we can easily determine the unit and counit by:

$$\eta_C = \phi_{C, \Delta(C)}(1_{\Delta(C)}) = 1_C: C \rightarrow \text{ver}\lim \Delta(C) = C \rightarrow C$$

$$\epsilon_D = \phi_{\text{ver}\lim(D), D}^{-1}(1_{\text{ver}\lim(D)}) : \Delta \text{ver}\lim(D) \rightarrow D \text{ is the natural transformation such that}$$

$$\epsilon_D: \text{ver}\lim(D) \rightarrow D \quad \text{for all } i \in \text{obj } J.$$

By duality, the left adjoint to Δ is the counit:

$$\text{ver}\lim + \Delta$$

And if $\text{Hom}(\text{ver}\lim(*), -) \rightarrow \text{Hom}(*, \Delta(-))$ is a natural isomorphism.

Therefore, we can determine the units and counits:

$$\eta_D = \phi_{D, \text{ver}\lim(D)}(1_{\text{ver}\lim(D)}) : D \rightarrow \Delta \text{ver}\lim(D) \text{ is a natural transformation that}$$

$$\eta_D: D \rightarrow \Delta \text{ver}\lim(D) \text{ for all } i \in \text{obj } J.$$

$$\epsilon_C = \phi_{\Delta(C), C}^{-1}(1_{\Delta(C)}) = 1_C: C \rightarrow C$$