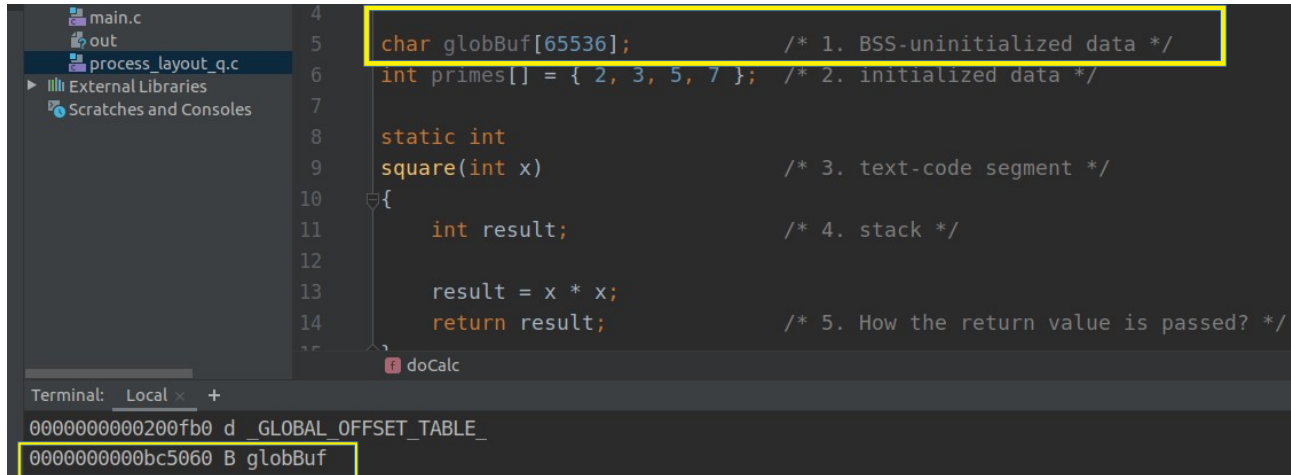


## Question 1:

1. **Question:** Where is allocated? `char globBuf[65536];`

**Answer:** BSS-uninitialized data

**proof:** we can see that after execute the “nm” command we got the type of “globBuf” to be B, that means according to “man” that “globBuf” is a global symbol and it’s found in BSS:



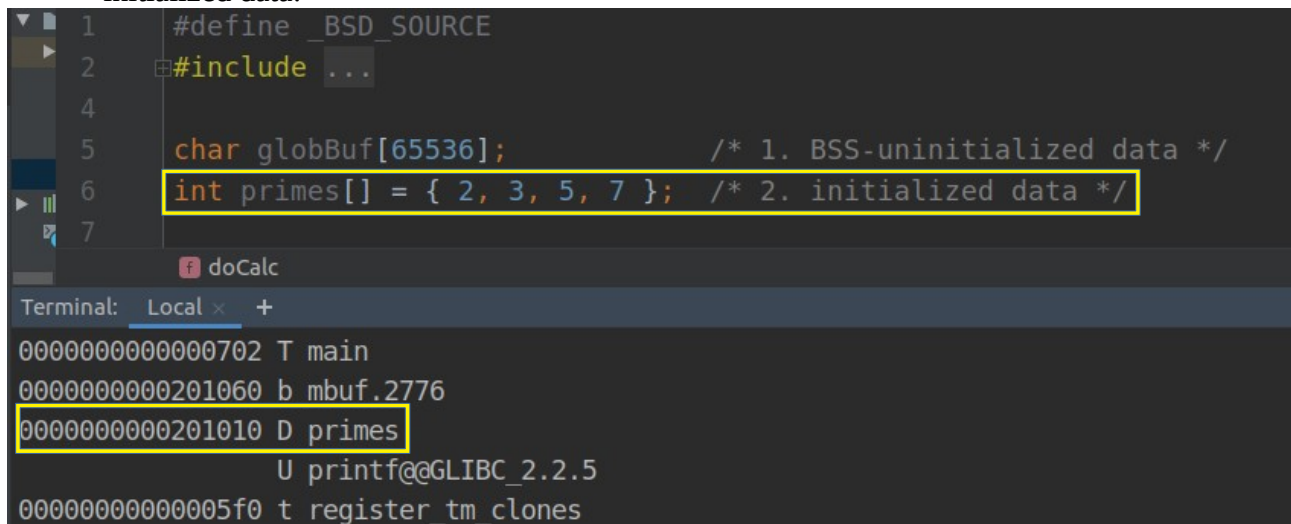
```
4 char globBuf[65536]; /* 1. BSS-uninitialized data */
5 int primes[] = { 2, 3, 5, 7 }; /* 2. initialized data */
6
7
8 static int
9 square(int x) /* 3. text-code segment */
10 {
11     int result; /* 4. stack */
12
13     result = x * x;
14     return result; /* 5. How the return value is passed? */
15 }
16
17 doCalc
```

```
Terminal: Local x +
0000000000200fb0 d _GLOBAL_OFFSET_TABLE_
0000000000bc5060 B globBuf
```

2. **Question:** Where is allocated? `int primes[] = { 2, 3, 5, 7 };`

**Answer:** initialized data

**proof:** we can see that after execute the “nm” command we got the type of “primes” to be D, that means according to “man” that “primes” is a global symbol and it’s found in initialized data:



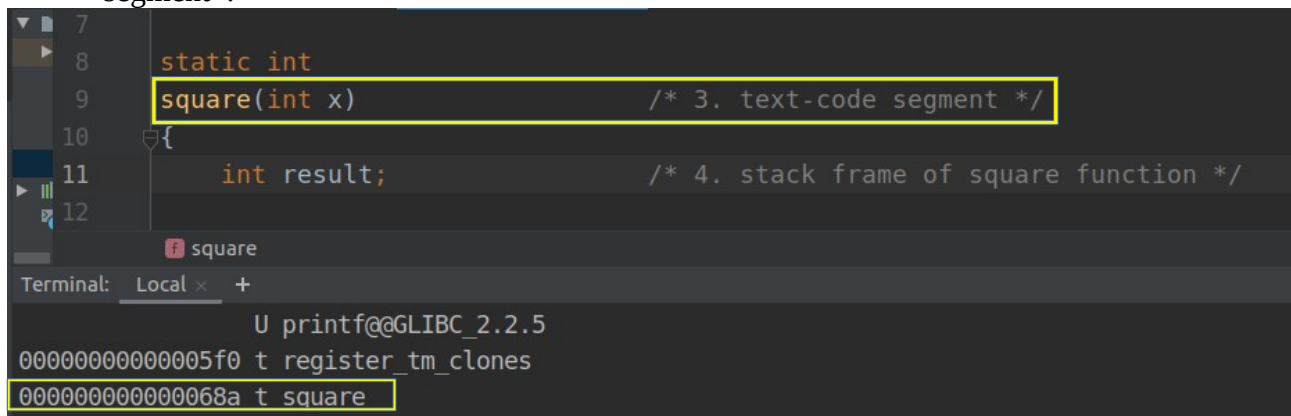
```
1 #define _BSD_SOURCE
2 #include ...
3
4
5 char globBuf[65536]; /* 1. BSS-uninitialized data */
6 int primes[] = { 2, 3, 5, 7 }; /* 2. initialized data */
7
```

```
Terminal: Local x +
0000000000000702 T main
0000000000201060 b mbuf.2776
0000000000201010 D primes
                U printf@@GLIBC_2.2.5
00000000000005f0 t register_tm_clones
```

3. **Question:** Where is allocated? `square(int x)`

**Answer:** text-code segment (The pointer of function).

**proof:** we can see that after execute the “nm” command we got the type of “square” to be t, that means according to “man” that “square” is a local symbol and it’s found in text-code segment :



```
7
8 static int
9 square(int x) /* 3. text-code segment */
10 {
11     int result; /* 4. stack frame of square function */
12 }

Terminal: Local x +
U printf@GLIBC_2.2.5
00000000000005f0 t register_tm_clones
000000000000068a t square
```

4. **Question:** Where is allocated? `int result;`

**Answer:** stack frame of square function

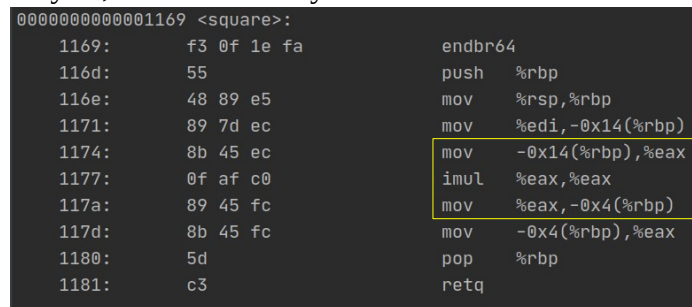
**Proof:** we can see that after execute the “objdump -d” command, we received the code of our program in Assembly Language.

We can be sure that **-0x14(%rbp)** hold the value of x variable because the “**mov %edi, -0x14(%rbp)**” command.

The “**result=x\*x**” command can be seen in the yellow square.

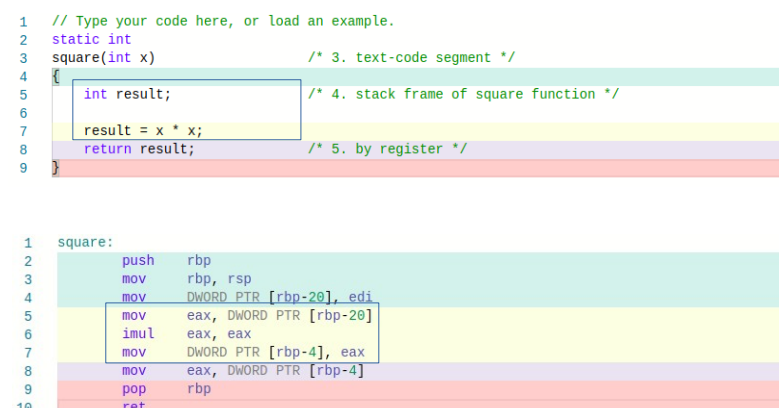
It can be seen that after the calculation of  $x * x$  is performed, the value passed by the rbp-base pointer register, which points to the base of the square stack frame.

The **-0x4(%rbp)** signals us that space has been allocated in a square stack frame the size of 4 bytes, which is exactly the size of the result variable.



```
0000000000001169 <square>:
1169: f3 0f 1e fa    endbr64
116d: 55             push %rbp
116e: 48 89 e5       mov %rsp,%rbp
1171: 89 7d ec       mov %edi,-0x14(%rbp)
1174: 8b 45 ec       mov -0x14(%rbp),%eax
1177: 0f af c0       imul %eax,%eax
117a: 89 45 fc       mov %eax,-0x4(%rbp)
117d: 8b 45 fc       mov -0x4(%rbp),%eax
1180: 5d             pop %rbp
1181: c3             retq
```

In addition, a screenshot is attached that proves to us that the “**result=x\*x**” command is indeed translated into the yellow square command (by the “godbolt” website).



```
1 // Type your code here, or load an example.
2 static int
3 square(int x) /* 3. text-code segment */
4 {
5     int result; /* 4. stack frame of square function */
6     result = x * x;
7     return result; /* 5. by register */
8 }

1 square:
2     push    rbp
3     mov     rbp, rsp
4     mov     DWORD PTR [rbp-20], edi
5     mov     eax, DWORD PTR [rbp-20]
6     imul    eax, eax
7     mov     DWORD PTR [rbp-4], eax
8     mov     eax, DWORD PTR [rbp-4]
9     pop     rbp
10    ret
```

5. **Question:** How the return value is passed? return result;

**Answer:** by register

**Proof:** we can see that after execute the “objdump -d” command, we received the code of our program in Assembly Language.

The “return result” command can be seen in the yellow square.

It can be seen that after the calculation of  $x * x$  is performed, it is passed to the `eax` register which is responsible among other things for the return values from the function, so we can be sure that the result value is returned by register.

```
0000000000001169 <square>:
1169: f3 0f 1e fa    endbr64
116d: 55             push  %rbp
116e: 48 89 e5       mov   %rsp,%rbp
1171: 89 7d ec       mov   %edi,-0x14(%rbp)
1174: 8b 45 ec       mov   -0x14(%rbp),%eax
1177: 0f af c0       imul  %eax,%eax
117a: 89 45 fc       mov   %eax,-0x4(%rbp)
117d: 8b 45 fc       mov   -0x4(%rbp),%eax
1180: 5d             pop   %rbp
1181: c3             retq
```

In addition, a screenshot is attached that proves to us that the “return result” command is indeed translated into the yellow square command (by the “godbolt” website).

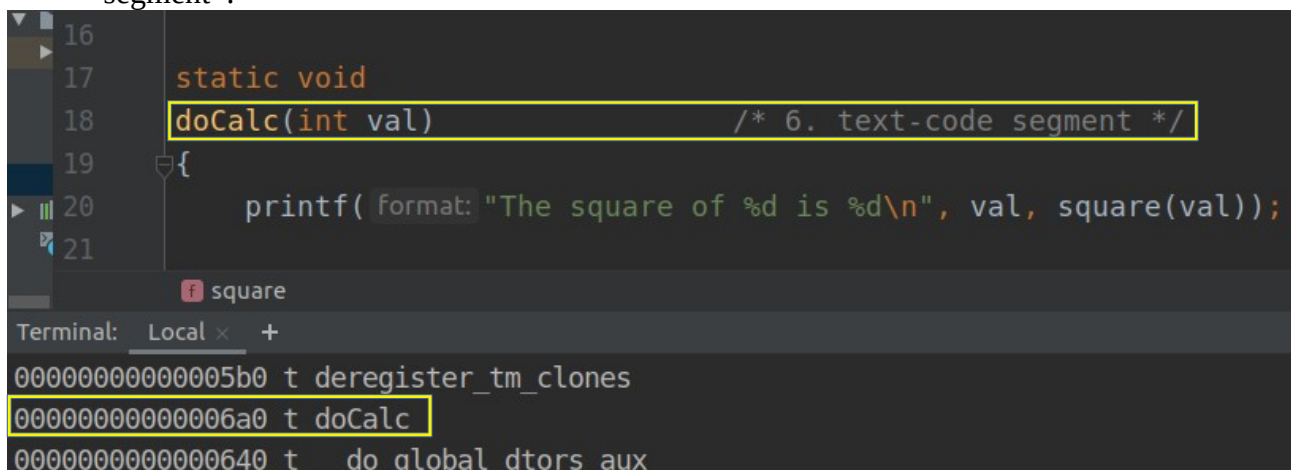
```
1 // Type your code here, or load an example.
2 static int
3 square(int x) /* 3. text-code segment */
4 {
5     int result; /* 4. stack frame of square function */
6
7     result = x * x;
8     return result; /* 5. by register */
9 }
```

```
1 square:
2 push    rbp
3 mov     rbp, rsp
4 mov     DWORD PTR [rbp-20], edi
5 mov     eax, DWORD PTR [rbp-20]
6 imul    eax, eax
7 mov     DWORD PTR [rbp-4], eax
8 mov     eax, DWORD PTR [rbp-4]
9 pop     rbp
10 ret
```

6. **Question:** Where is allocated? `doCalc(int val)`

**Answer:** text-code segment (The pointer of function).

**proof:** we can see that after execute the “nm” command we got the type of “doCalc” to be `t`, that means according to “man” that “doCalc” is a local symbol and it’s found in text-code segment :



```
16
17 static void
18 doCalc(int val) /* 6. text-code segment */
19 {
20     printf(format: "The square of %d is %d\n", val, square(val));
21 }

f square

Terminal: Local x +
00000000000005b0 t deregister_tm_clones
00000000000006a0 t doCalc
0000000000000640 t __do_global_dtors_aux
```

7. **Question:** Where is allocated? int t;

**Answer:** stack frame of doCalc function

**Proof:** we can see that after execute the “objdump -d” command, we received the code of our program in Assembly Language.

The “t=val\*val\*val” command can be seen in the yellow square.

We can be sure that -0x14(%rbp) hold the value of val variable because the “mov %edi, - 0x14(%rbp)” command.

This value passed into eax register.

After that, the value of eax multiplied by itself and saved into eax register.

Now, the value of -0x14(%rbp) passed into edx register.

After that, the value of eax multiplied by edx and saved into eax register.

At last, eax register hold the value of val\*val\*val, and this value copied into -0x4(%rbp).

That prove that t variable is allocated on the doCalc stack frame.

```
0000000000001182 <doCalc>:
1182: f3 0f 1e fa      endbr64
1186: 55              push %rbp
1187: 48 89 e5         mov %rsp,%rbp
118a: 48 83 ec 20      sub $0x20,%rsp
118e: 89 7d ec         mov %edi,-0x14(%rbp)
1191: 8b 45 ec         mov -0x14(%rbp),%eax
1194: 89 c7           mov %eax,%edi
1196: e8 ce ff ff ff   callq 1169 <square>
119b: 89 c2           mov %eax,%edx
119d: 8b 45 ec         mov -0x14(%rbp),%eax
11a0: 89 c6           mov %eax,%esi
11a2: 48 8d 3d 5b 0e 00 00 lea 0xe5b(%rip),%rdi # 2004 <_IO_stdin_used+0x4>
11a9: b8 00 00 00 00   mov $0x0,%eax
11ae: e8 ad fe ff ff   callq 1060 <printf@plt>
11b3: 81 7d ec e7 03 00 00 cmpl $0x3e7,-0x14(%rbp)
11ba: 7f 28           jg 11e4 <doCalc+0x62>
11bc: 8b 45 ec         mov -0x14(%rbp),%eax
11bf: 0f af c0        imul %eax,%eax
11c2: 8b 55 ec         mov -0x14(%rbp),%edx
11c5: 0f af c2        imul %edx,%eax
11c8: 89 45 fc         mov %eax,-0x4(%rbp)
11cb: 8b 55 fc         mov -0x4(%rbp),%edx
11ce: 8b 45 ec         mov -0x14(%rbp),%eax
11d1: 89 c6           mov %eax,%esi
11d3: 48 8d 3d 42 0e 00 00 lea 0xe42(%rip),%rdi # 201c <_IO_stdin_used+0x1c>
11da: b8 00 00 00 00   mov $0x0,%eax
11df: e8 7c fe ff ff   callq 1060 <printf@plt>
11e4: 90              nop
11e5: c9              leaveq
11e6: c3              retq
```

In addition, a screenshot is attached that proves to us that the “t=val\*val\*val” command is indeed translated into the yellow square command (by the “godbolt” website).

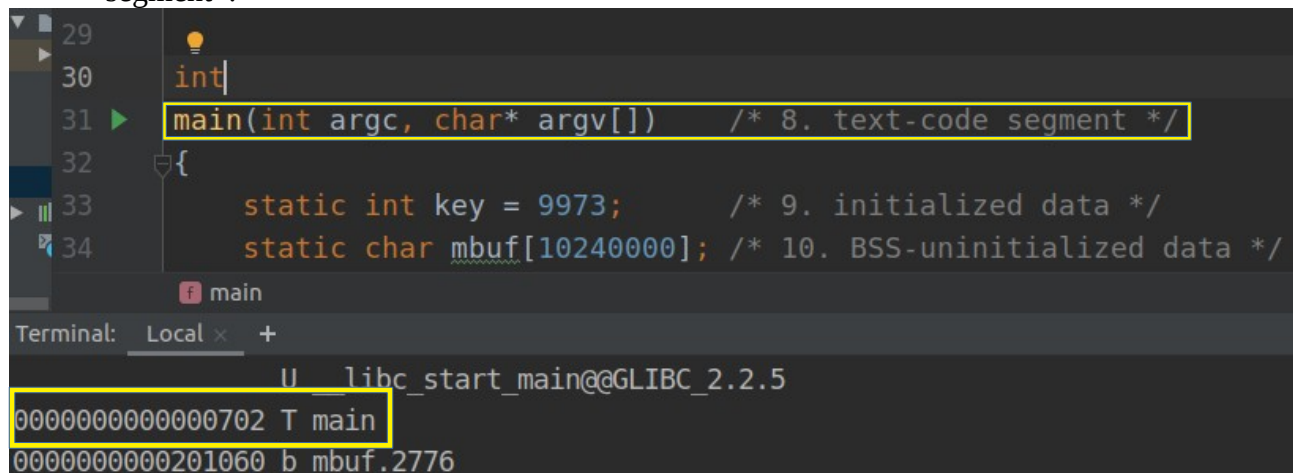
```
22 doCalc:
23     push    rbp
24     mov     rbp, rsp
25     sub     rsp, 32
26     mov     DWORD PTR [rbp-20], edi
27     mov     eax, DWORD PTR [rbp-20]
28     mov     edi, eax
29     call    square
30     mov     edx, eax
31     mov     eax, DWORD PTR [rbp-20]
32     mov     esi, eax
33     mov     edi, OFFSET FLAT:_L0
34     mov     eax, 0
35     call    printf
36     cmp     DWORD PTR [rbp-20], 999
37     jg      .L5
38     mov     eax, DWORD PTR [rbp-20]
39     imul    eax, eax
40     mov     edx, DWORD PTR [rbp-20]
41     imul    eax, edx
42     mov     DWORD PTR [rbp-4], eax
43     mov     edx, DWORD PTR [rbp-4]
44     mov     eax, DWORD PTR [rbp-20]
45     mov     esi, eax
46     mov     edi, OFFSET FLAT:_L1
47     mov     eax, 0
48     call    printf
49 .L5:
50     nop
51     leaveq
52     ret
```

```
18 static void
19 doCalc(int val) /* 6. text-code segment */
20 {
21     printf("The square of %d is %d\n", val, square(val));
22
23     if (val < 1000) {
24         int t; /* 7. stack frame of doCalc function */
25
26         t = val * val * val;
27         printf("The cube of %d is %d\n", val, t);
28     }
29 }
30
```

8. **Question:** Where is allocated? `main(int argc, char*\ argv[])`

**Answer:** text-code segment (The pointer of function).

**Proof:** we can see that after execute the “nm” command we got the type of “main” to be T, that means according to “man” that “main” is a global symbol and it’s found in text-code segment :

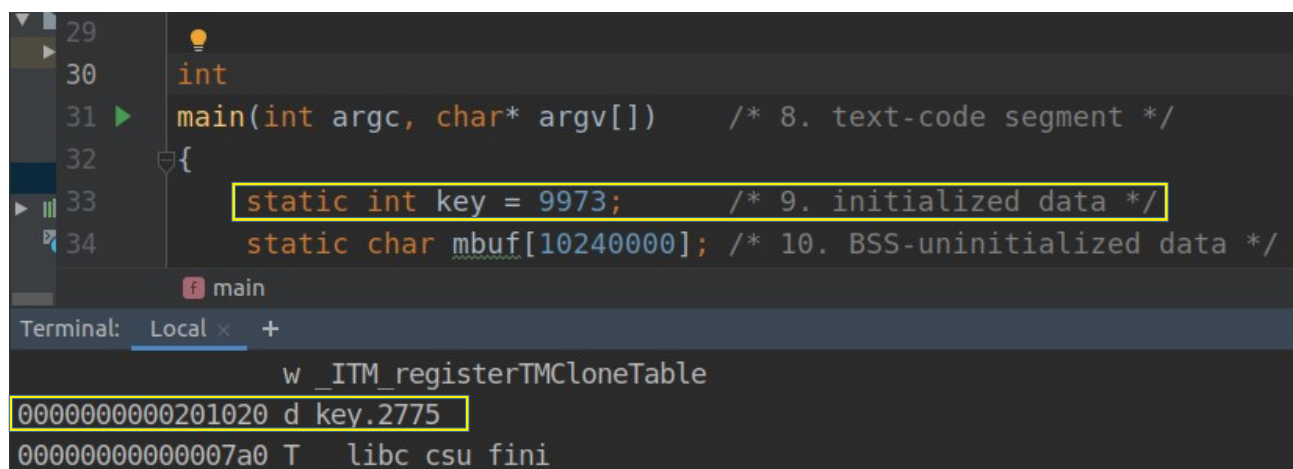


```
29
30 int
31 main(int argc, char* argv[]) /* 8. text-code segment */
32 {
33     static int key = 9973; /* 9. initialized data */
34     static char mbuf[10240000]; /* 10. BSS-uninitialized data */
35 }
main
Terminal: Local x +
U _libc_start_main@@GLIBC_2.2.5
00000000000000702 T main
00000000000201060 b mbuf.2776
```

9. **Question:** Where is allocated? `static int key = 9973;`

**Answer:** initialized data

**proof:** we can see that after execute the “nm” command we got the type of “key” to be d, that means according to “man” that “key” is a local symbol and it’s found in initialized data:



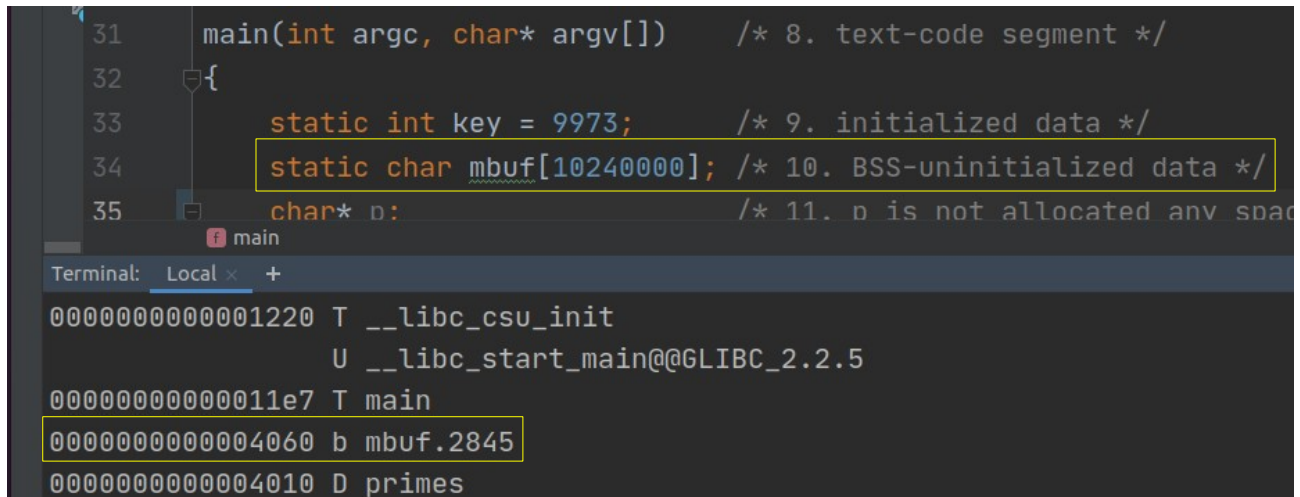
```
29
30 int
31 main(int argc, char* argv[]) /* 8. text-code segment */
32 {
33     static int key = 9973; /* 9. initialized data */
34     static char mbuf[10240000]; /* 10. BSS-uninitialized data */
35 }
main
Terminal: Local x +
w _ITM_registerTMCloneTable
00000000000201020 d key.2775
00000000000007a0 T __libc_csu_fini
```



10. **Question:** Where is allocated? static char mbuf[10240000];

**Answer:** BSS-uninitialized data

**proof:** we can see that after execute the “nm” command we got the type of “mbuf” to be b, that means according to “man” that “mbuf” is a local symbol and it’s found in BSS :



The image shows a code editor with a C program snippet and a terminal window below it. The code defines a static integer 'key' and a static character array 'mbuf' of size 10240000. The terminal shows the output of the 'nm' command, listing symbols in memory. The entry for 'mbuf' is highlighted with a yellow box, showing it is of type 'b' (BSS-uninitialized data) and located at address 00000000000004060.

```
31 main(int argc, char* argv[]) /* 8. text-code segment */
32 {
33     static int key = 9973; /* 9. initialized data */
34     static char mbuf[10240000]; /* 10. BSS-uninitialized data */
35     char* p; /* 11. p is not allocated any space */
}

Terminal: Local x +
00000000000001220 T __libc_csu_init
U __libc_start_main@@GLIBC_2.2.5
000000000000011e7 T main
00000000000004060 b mbuf.2845
00000000000004010 D primes
```

11. **Question:** Where is allocated? char\* p;

**Answer:** p is not allocated any space in memory because it is not initialized during the program.

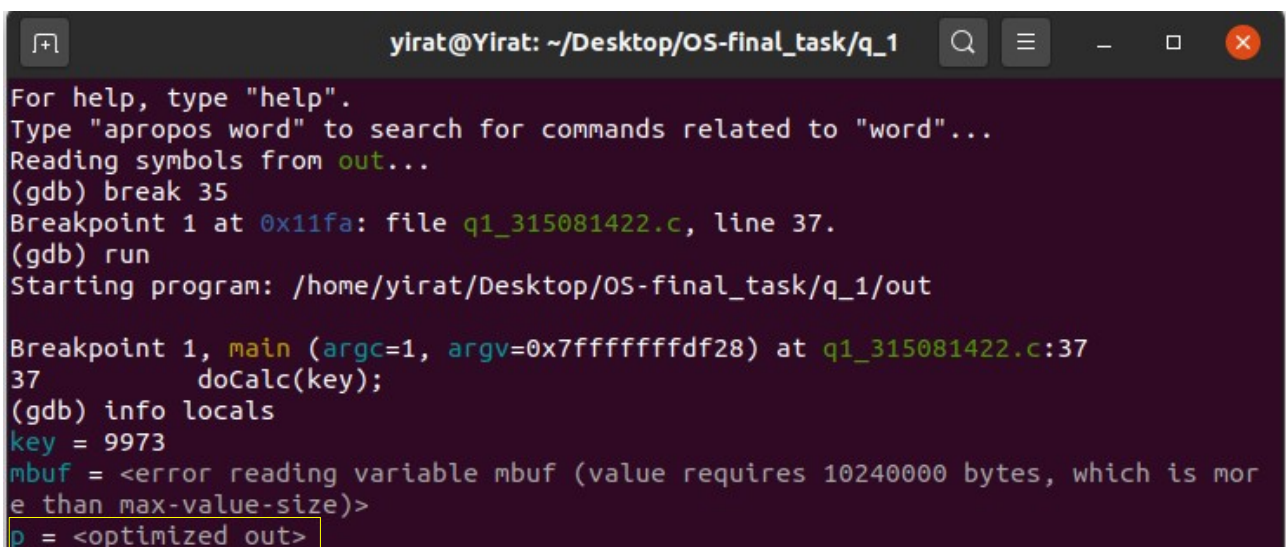
If it was initialized, it would be allocated space in the stack frame of main function.(Because the compiler has certain optimizations).

**NOTE:** There are compilers with certain optimizations that will not allocate space in the memory for a variable that will not be initialized during the program.

In our program we are referring to a compiler without these optimizations.

**Proof:** We can see that we ran with gdb the info locals command that shows the variables that are on the stack.

It can be seen in the yellow square that the variable p was optimized and therefore it is not on the stack, but without the optimization it would have been there.



The image shows a GDB terminal window. It displays the execution of a program with a breakpoint at line 37. The 'info locals' command is used to show the state of local variables. The output shows 'key' as 9973, 'mbuf' as an error (due to its large size), and 'p' as '<optimized out>', which is highlighted with a yellow box.

```
yirat@Yirat: ~/Desktop/OS-final_task/q_1
For help, type "help".
Type "apropos word" to search for commands related to "word"...
Reading symbols from out...
(gdb) break 35
Breakpoint 1 at 0x11fa: file q1_315081422.c, line 37.
(gdb) run
Starting program: /home/yirat/Desktop/OS-final_task/q_1/out

Breakpoint 1, main (argc=1, argv=0x7fffffffdf28) at q1_315081422.c:37
37      doCalc(key);
(gdb) info locals
key = 9973
mbuf = <error reading variable mbuf (value requires 10240000 bytes, which is more than max-value-size)>
p = <optimized out>
```