System Calls

Invoking Kernel Operations

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System Calls

- · User programs can't directly
 - · Communicate with (most) hardware
 - Read/ write kernel data/ data structures

...unrestricted access would allow any program to read/ change anything

- System Calls
 - Offer standard interface to kernel/ OS functions
 - Provide access control mechanism
 - Single point of entry that can check all parameters

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Unix man pages

- 1. User Commands
- 2. System Calls
- 3. C Library Functions
- 4. Devices and Special Files
- 5. File Formats and Conventions
- 6. Games etc.
- 7. Miscellaneous
- 8. System Administration tools and Daemons

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Making a System Call on x86 Fill registers (EAX, EBX, ...) with Function code (for read, write, ...) Parameters Execute software interrupt: INT 80 (Unix) or INT 2E (WinNT) Function Code Parameter 1 Parameter 2 Parameter 3 Parameter 4 Parameter 5 eax ebx ecx edx esx edi OR Pointer to 6 (or more) parameters

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```
Making Unix System Calls
       Build with: as syscall.s -o syscall.o ld syscall.o -o syscall
.section
              .rodata
      .string
                     "Hello World\n"
      LEN, .-MSG
                             # calculate length of message string
.text
.global
              _start
start:
              $4.
                             # write: System call #4
      mov
                      %eax
              $1,
                      %ebx
                             # fd = 1 : stdout
      mov
              $MSG, %ecx
$LEN, %edx
                             # msg : char * msg
# number of characters to output
      mov
                              # syscall: int write(fd, msg, len)
                             # exit: System call #1
# code = 0 (success)
              $1,
      mov
                      %eax
      int
              $0x80
                              # syscall: exit(0)
```

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Direct System Call from C main () { char * msg = "Hello World\n"; Remember inline assembler allows us to directly specify the x86 registers variables should be passed in, e.g., to pass fd in reg. EBX we use "b" (fd) // stdout fd = 1; int len = (int) strlen(msg); -- this really simplifies the code retval; asm volatile ("mov \$4, "int \$0x80; "movl %%eax, %0 # get return value from function" : "=rm" (retval) // %0 refers to return variable retval : "b" (fd), "c" (msg), "d" (len) // Place in EBX, ECX, EDX : "%eax" // Clobbered/ overwritten printf ("write(fd, msg, len) reported %d characters written $\n"$ retval);

Kernel System Call Entry Point

- · Software Interrupts appear as any other interrupt
 - Normal ISR entry point
 - Simply check for Unix (0x80) or NT (0x2E) 'call' convention

```
    Note difference, one kernel could support both Unix and Windows

void
interrupt_handler(
        struct cpu registers regs, uint32_t irq, uint32_t code,
uint32_t eip, uint32_t cs, uint32_t flags,
uint32_t esp, uint32_t ss) {
    switch (irg) {
    case 0x80:
                                      // INT 0x80 (80 hex = 128 decimal)
        systemcall ( &regs );
```

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