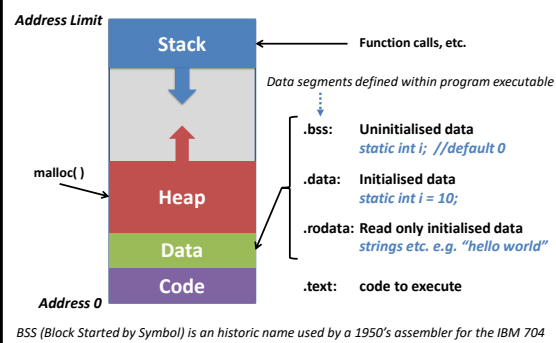


Processes: Memory Allocation

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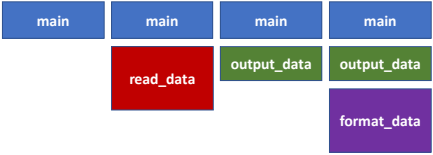
Program Address Space



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Process Stack

- Common way of implementing function calls
 - Functions parameters + return address placed on stack before call
 - There is no standard, could use mix of registers and stack, ...
- Grows and shrinks in blocks (frames) with function calls
 - On x68 stack grows down



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Stack Allocation Test

- What would this code output?

```
#include <stdio.h>
#include <string.h>

char *
first() {
    char first_buff[ 7 ];

    return strcpy( first_buff, "Hello" );
}

char *
second() {
    char second_buff[ 7 ];

    return strcpy( second_buff, "World!" );
}

int
main() {
    char * str1 = first( );
    char * str2 = second( );

    printf( "%s %s\n", str1, str2 );
}
```

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Stack Allocation Test

- What would this code output?

```
#include <stdio.h>
#include <string.h>

char *
first() {
    char first_buff[ 7 ];

    return strcpy( first_buff, "Hello" );
}

char *
second() {
    char second_buff[ 7 ];

    return strcpy( second_buff, "World!" );
}

int
main() {
    char * str1 = first( );
    char * str2 = second( );

    printf( "%s %s\n", str1, str2 );
}
```

- World! World!

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Stack Allocation Test Output

- Let's output some addresses...

```
str1 = 0x8BD1
str2 = 0x8BD1
World! World!
```

...they're the same address!

```
#include <stdio.h>
#include <string.h>

char *
first() {
    char first_buff[ 7 ];

    return strcpy( first_buff, "Hello" );
}

char *
second() {
    char second_buff[ 7 ];

    return strcpy( second_buff, "World!" );
}

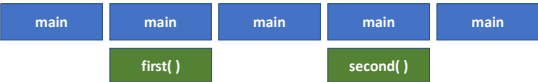
int
main() {
    char * str1 = first( );
    char * str2 = second( );

    → printf( "str1 = %p\nstr2 = %p\n", str1, str2 );
    printf( "%s %s\n", str1, str2 );
}
```

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Think About the Process Stack

- To understand what’s happening...
...think how the stack grows and shrinks



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Let’s watch the stack pointer...

- Notice
 - The stack grows down as we enter first function, F0→C0
 - Then shrinks back to same address it started at, C0→F0
- Same pattern repeated for second function, F0→C0→F0
 - Addresses are the same, so same address (8BD1) used for first_buff and second_buff
- Therefore, when we fill second_buff, we overwrite what was there before

```
unsigned long long stackVal;  
asm( "mov %%rsp, %0" : "=rm" ( stackVal ) )
```

```
Stack before first( ): 0x8BF0  
Stack in first( ): 0x8BC0  
Stack after first( ): 0x8BF0  
  
Stack in second( ): 0x8BC0  
Stack after second( ): 0x8BF0  
  
str1 = 0x8BD1, str2 = 0x8BD1  
World! World!
```

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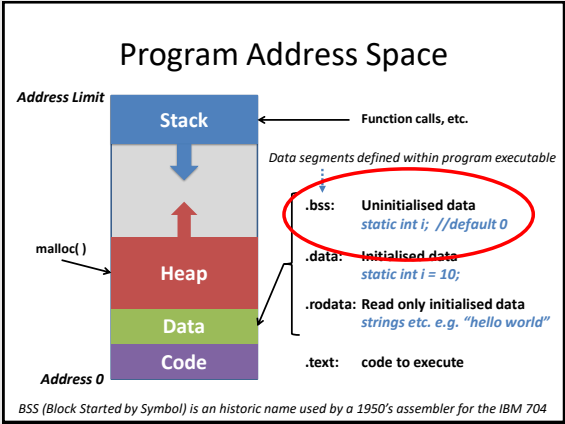
Stack Allocation Test: *a simple fix?*

- Buffers not on stack
 - 2x7 bytes now safe in BSS segment (*next slide*)
...and never freed, but...
- Wasteful if much bigger
 - Need better way...

```
Stack before first( ): 0x8BF0  
Stack in first( ): 0x8BD0  
Stack after first( ): 0x8BF0  
  
Stack in second( ): 0x8BD0  
Stack after second( ): 0x8BF0  
  
str1 = 0x5011, str2 = 0x5018  
Hello World!
```

```
#include <stdio.h>  
#include <string.h>  
  
char *  
first( ) {  
    static char first_buff[ 7 ];  
    return strcpy( first_buff, "Hello" );  
}  
  
char *  
second( ) {  
    static char second_buff[ 7 ];  
    return strcpy( second_buff, "World!" );  
}  
  
int  
main( ) {  
    char * str1 = first( );  
    char * str2 = second( );  
  
    printf( "%s %s\n", str1, str2 );  
}
```

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Process Heap

- Predefined memory areas often too restrictive
 - For example, arrays can't expand if more data than expected
 - End up over sizing to be safe
- Heap allows dynamic allocation of memory as needed
 - Elements in linked-list, queue, graph, ...
 - New object instances, etc.
- Use library calls to expand and (possibly) contract heap
 - `memory = malloc()`
 - `free(memory)`

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