

# Grammar of TeaPL

Each program is composed of variable declarations, function declarations, function definitions, and comments.

```
program := (varDeclStmt | structDef | fnDeclStmt | fnDef | comment | < ; >)*
```

## Basic Identifiers, Values, Expressions, and Assignments

Each identifier begins with an alphabet and contains only alphabets and digits, e.g., `alice`, `a0`.

```
id := [a-zA-Z][a-zA-Z0-9]*
```

TeaPL allows integers, e.g., `123`

```
num := [1-9][0-9]* | 0
```

### Arithmetic Expressions

An expression is composed of identifiers, values, and operators, e.g., `1+2`, `a*(b+c)`. For simplicity, we do not support unary operators, such as `++`, `+=`.

```
arithExpr := arithExpr arithBiOp arithExpr | exprUnit
exprUnit := num | id | < ( > arithExpr < ) > | fnCall | id < [ > id | num < ] > | id <
arithBiOp := < + > | < - > | < * > | < / >
arithUOp := < - >
```

主要可能是优先级的问题

### Condition Expressions

```
boolExpr := boolExpr boolBiOp boolUnit | boolUnit
boolUnit := < ( > exprUnit comOp exprUnit < ) > | < ( > boolExpr < ) > | boolUOp boolUn
boolBiOp := < && > | < || >
boolUOp := < ! >
comOp := < > > | < < > | < >= > | < <= > | < == > | < != >
```

## Assignment

We restrict neither the left value nor right value can be assignments.

```
assignStmt := leftVal < = > rightVal < ; >
leftVal   := id < [ > id | num < ] > | id < . > id
rightVal  := arithExpr | boolExpr
```

## Function Call

```
fnCall := id < ( > rightVal (< , > rightVal)* | ε < ) >
```

## Variable Declarations

TeaPL allows declaring one variable each time, which can be either a primitive or array type.

Developers can initialize the variable during declaration. For example, it supports the following variable declaration samples.

### Primitive Types

```
let a:int; // declare a variable of type int; the type field can be ignored.
let b:int = 0; // declare a variable of int and init it with value 0.
```

### One-level Array

```
let c[10]:int; // declare a variable of integer array.
let d[10]:int = {0}; // declare a variable of integer array and initialize it with zero
```

The grammar is defined as follows.

```
varDeclStmt := < let > (varDecl | varDef) < ; >
varDecl    := id < : > type | id < [ > num < ] >< : > type
varDef     := id < : > type < = > rightVal //primitive type
            | id < [ > num < ] >< : > type < = > < { > rightVal (< , > rightVal)* | ε < } >
type       := nativeType | structType | ε
nativeType := < int >
structType := id
```

## Define A New Structure

Developers can define new customized types with the preserved keyword struct, e.g.,

```
struct MyStruct {  
    node:int,  
    len:int  
}
```

The grammar is defined as follows.

```
structDef := < struct > id < { > (varDecl) (< , > varDecl)* < } >
```

## Function Declaration and Definition

Each function declaration starts with the keyword fn.

```
fn foo(a:int, b:int)->int;  
fn foo();
```

The grammar is defined as follows.

```
fnDeclStmt := fnDecl < ; >  
fnDecl := < fn > id < ( > paramDecl < ) > //without return value  
         | < fn > id < ( > paramDecl < ) > < -> > type //with return value  
paramDecl := varDecl (< , > varDecl)* | ε
```

### Function Definition

We can also define a function while declaring it.

```
fn foo(a:int, b:int)->int {  
    return a + b;  
}
```

The grammar is specified as follows.

```

fnDef := fnDecl codeBlock
codeBlock := < { > (varDeclStmt | assignStmt | callStmt | ifStmt | whileStmt | returnS
returnStmt := < ret > rightVal < ; >
continueStmt := < continue > < ; >
breakStmt := < break > < ; >

```

We have already defined the grammar of varDeclStmt and assignStmt. The callStmt is simply a function call terminated with an colon.

```

callStmt := fnCall < ; >

```

Next, we define the grammar of each rest statement type.

## Control Flows

### If-Else Statement

The condition should be surrounded with a paired parenthesis, and we further restrict the body should be within a paired bracket. The following shows an example.

```

if (x >0) {
    if (y >0) {
        x++;
    }
    else {
        x--;
    }
} else {

}

```

Besides, we restrict the condition expression to be explicit logical operations, e.g.,  $x > 0$ ; we donot allow implicit expressions like  $x$ , which means. We define the grammar as follows.

```

ifStmt := < if > < ( > boolExpr < ) > codeBlock ( < else > codeBlock | ε )

```

### While Statemet

Used for the representability of complicated loops.

Example:

```
while (x > 0) {  
    x--;  
}
```

Definition:

```
whileStmt := < while > < ( > boolExpr < ) > codeBlock
```

## Code Comments

Similar to most programming languages, TeaPL allows line comments with `//` and scope comments with `/* ... */`.

```
int a = 0; // this is a line comment.
```

```
/*  
    Feature: this is a scope comment  
*/  
fn foo(){  
    ...  
}
```

```
comment := < // > _* | < /* > _* < */ >
```