

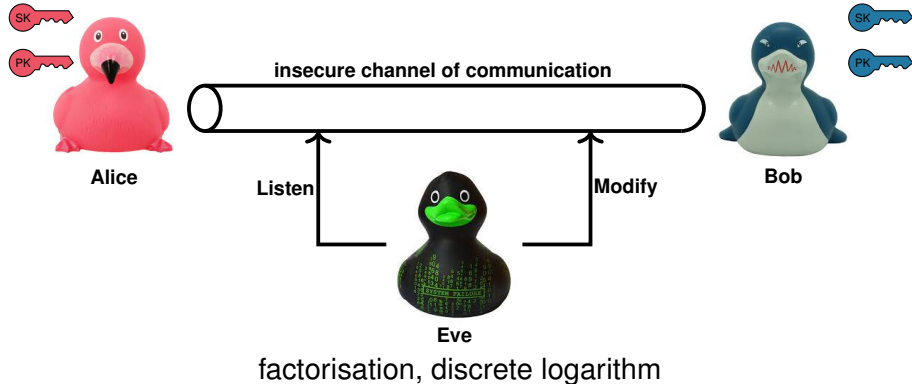
# Quantum Algorithms for Lattice-based Cryptography

Yixin Shen

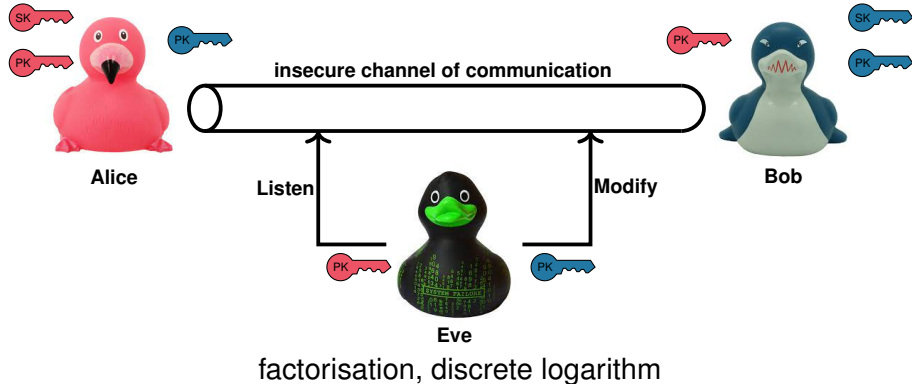
November 4, 2022



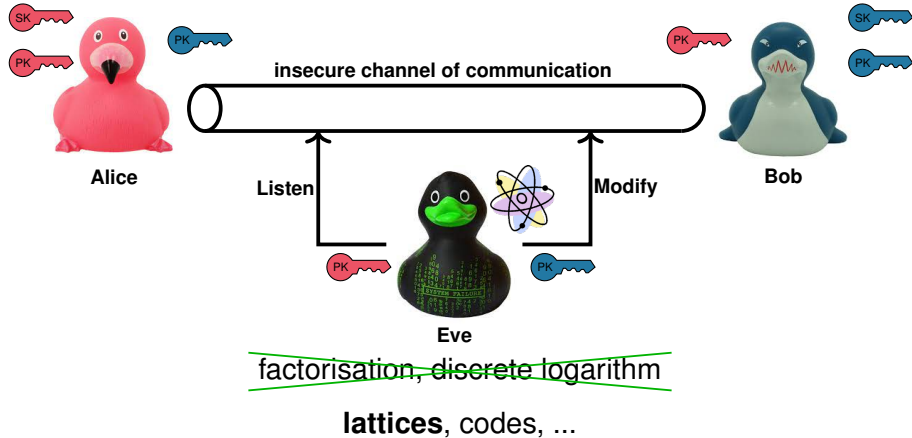
# Hard Problems in Public Key Cryptography



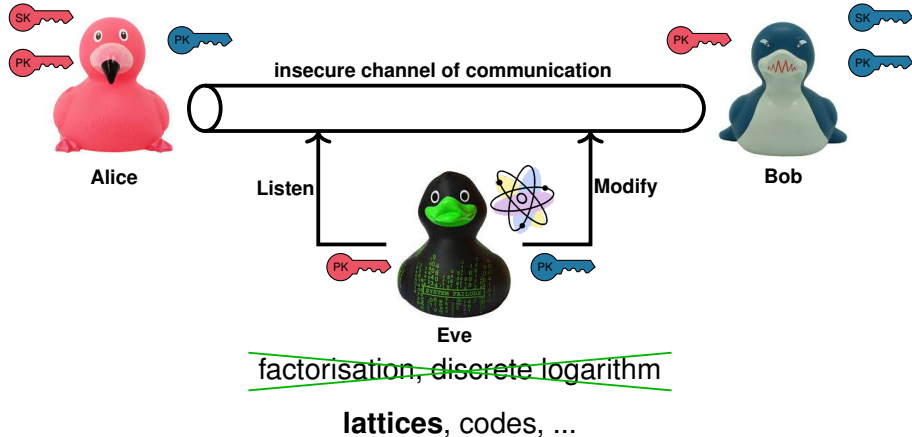
# Hard Problems in Public Key Cryptography



# Hard Problems in Public Key Cryptography



# Hard Problems in Public Key Cryptography



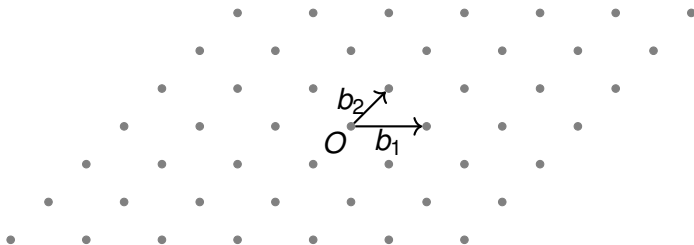
NIST selected algorithms:

- ▶ **encryption**: the only selected candidate is based on lattices
- ▶ **signatures**: 2 out of 3 based on lattices

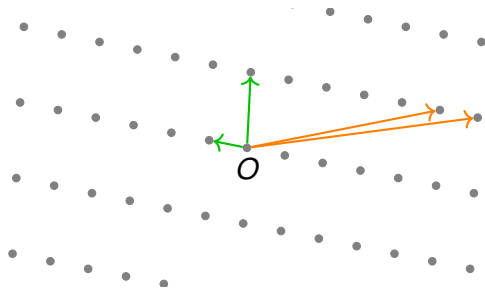
# What is a (Euclidean) lattice?

## Definition

$\mathcal{L}(\mathbf{b}_1, \dots, \mathbf{b}_n) = \left\{ \sum_{i=1}^n x_i \mathbf{b}_i : x_i \in \mathbb{Z} \right\}$  where  $\mathbf{b}_1, \dots, \mathbf{b}_n$  is a basis of  $\mathbb{R}^n$ .

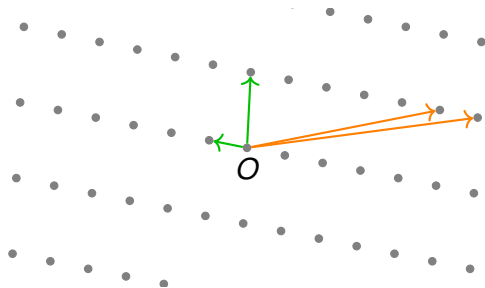


# Lattice-based cryptography: fundamental idea



- ▶ **good basis:** private information, makes problem easy
- ▶ **bad basis:** public information, makes problem hard

# Lattice-based cryptography: fundamental idea



- ▶ **good basis**: private information, makes problem easy
- ▶ **bad basis**: public information, makes problem hard

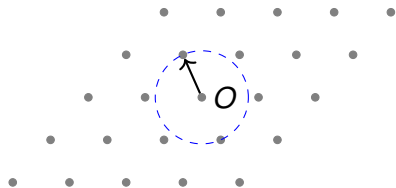
**Basis reduction**: transform a bad basis into a good one

**Main tool**: BKZ algorithm and its variants

Requires to solve the **(approx-)SVP problem** in smaller dimensions.



# The Shortest Vector Problem

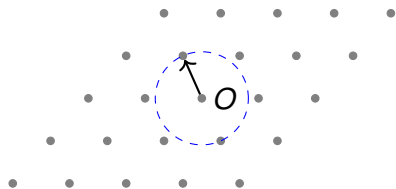


Shortest Vector Problem (SVP):

Given a basis for the lattice  $\mathcal{L}$ , find a shortest nonzero lattice vector.

$\lambda_1(\mathcal{L}) =$  length of such a vector.

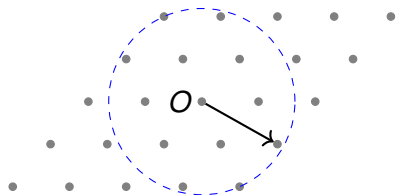
# The Shortest Vector Problem



**Shortest Vector Problem (SVP):**

Given a basis for the lattice  $\mathcal{L}$ , find a shortest nonzero lattice vector.

$\lambda_1(\mathcal{L}) = \text{length of such a vector.}$



**$\gamma$ -approx-SVP ( $\gamma > 1$ ):**

Given a basis of  $\mathcal{L}$ , find a nonzero lattice vector of length at most  $\gamma \cdot \lambda_1(\mathcal{L})$ .

$\gamma$  is approximation factor.

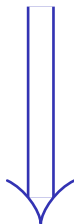
# The Shortest Vector Problem

Depending on the dimension  $n$ :

- ▶ NP-Hardness (randomized reduction)
- ▶  $\text{NP} \cap \text{co-NP}$
- ▶ Subexponential-time algorithms
- ▶ Poly-time algorithms

Approx factor:

- ▶  $O(1)$
- ▶  $\sqrt{n}$
- ▶  $2^{\sqrt{n}}$
- ▶  $2^{\frac{n \log \log n}{\log n}}$



# The Shortest Vector Problem

Depending on the dimension  $n$ :

- ▶ NP-Hardness (randomized reduction)
- ▶  $\text{NP} \cap \text{co-NP}$
- ▶ Subexponential-time algorithms
- ▶ Poly-time algorithms

Approx factor:

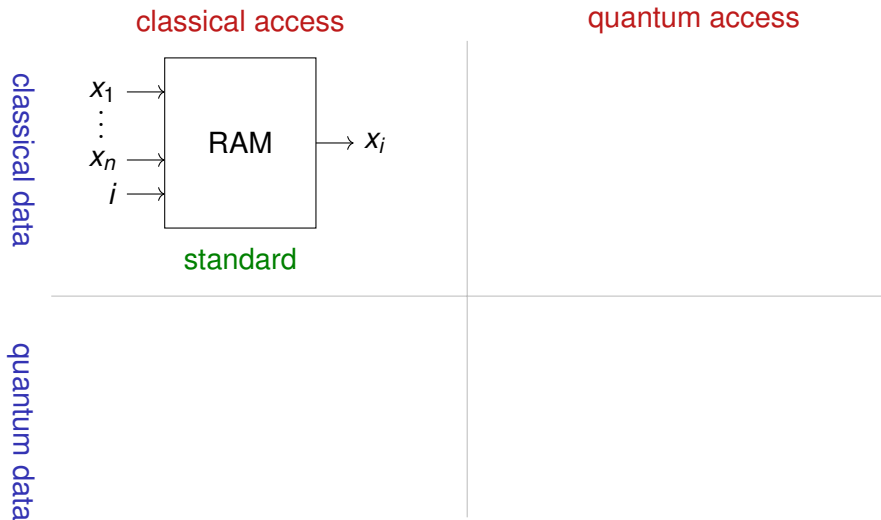
- ▶  $O(1)$
- ▶  $\sqrt{n}$
- ▶  $2^{\sqrt{n}}$
- ▶  $2^{\frac{n \log \log n}{\log n}}$



Main approaches for SVP:

- ▶ Enumeration:  $2^{O(n \log(n))}$  time and  $\text{poly}(n)$  space
- ▶ Sieving:  $2^{O(n)}$  time and  $2^{O(n)}$  space

# Interlude: quantum memory models



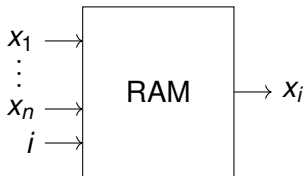
Assumption:  $O(1)$  time cost

# Interlude: quantum memory models

classical access

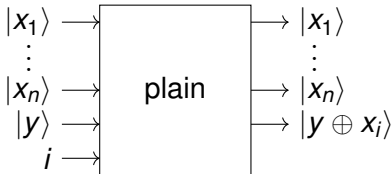
quantum access

classical data



standard

quantum data

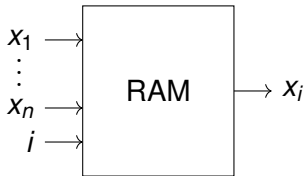


standard

Assumption:  $O(1)$  time cost

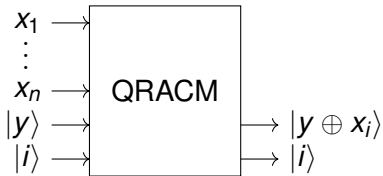
# Interlude: quantum memory models

classical access



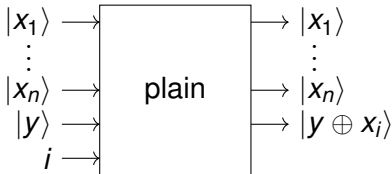
standard

quantum access



potentially strong assumption

quantum data

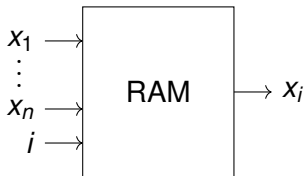


standard

Assumption:  $O(1)$  time cost

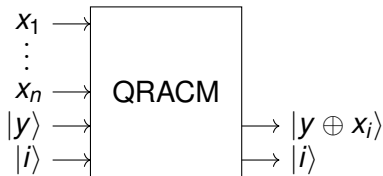
# Interlude: quantum memory models

classical access



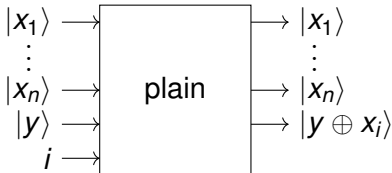
standard

quantum access

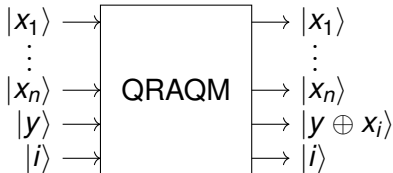


potentially strong assumption

classical data



standard



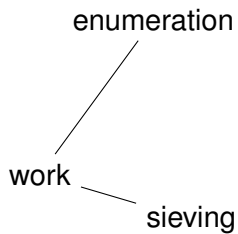
strong assumption

quantum data

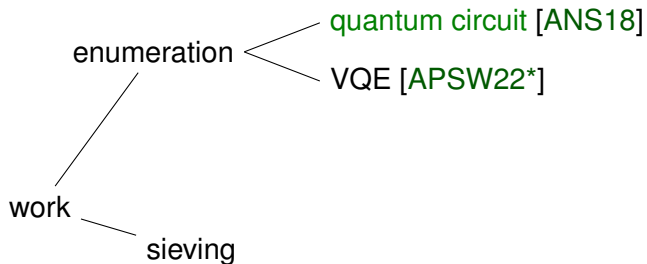
Assumption:  $O(1)$  time cost



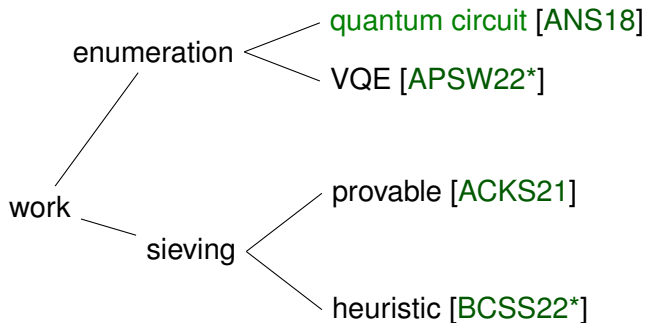
# Overview of my work



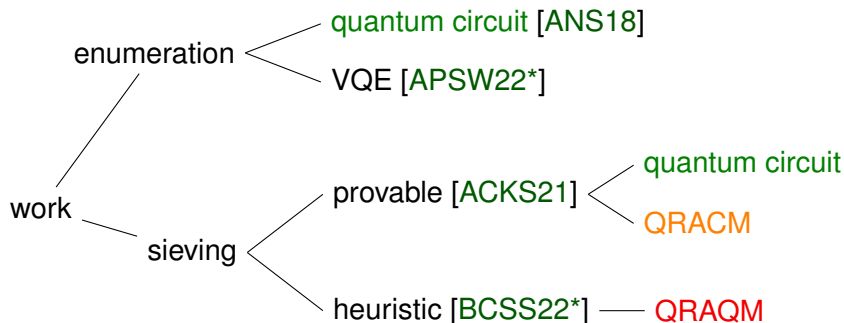
# Overview of my work



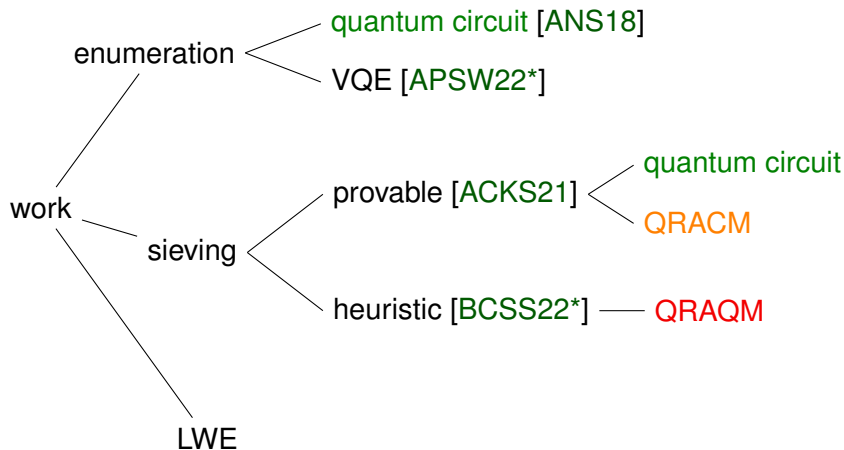
# Overview of my work



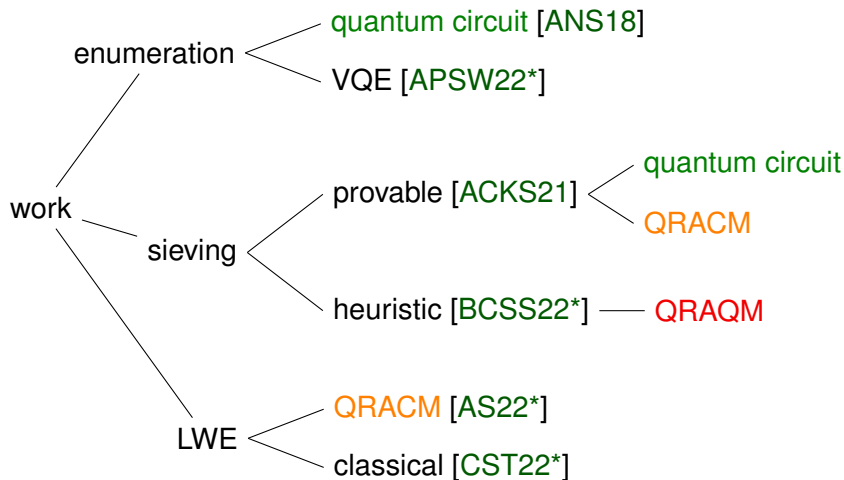
# Overview of my work



# Overview of my work



# Overview of my work



# Overview of my work

