# Full Specification of the Object Language

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# 1 Syntax

$$\begin{array}{lll} \Gamma \coloneqq \cdot \mid \Gamma, x : A & \text{Typing contexts} \\ A, B, C, \\ a, b, c, P \coloneqq & \text{Terms} \\ \mid x \mid \mathcal{U}_i & variables, \ type \ universes \\ \mid \Pi x : A . B \mid \lambda x . b \mid b \ a & function \ types, \ abstractions, \ applications \\ \mid \Sigma x : A . B \mid (a, b) \mid \pi_1 \ x \mid \pi_2 \ x & pair \ types, \ pairs, \ projections \\ \mid \mathbf{Nat} \mid \mathbf{zero} \mid \mathbf{succ} \ a \mid \mathbf{ind} \ P \ a \ b \ (\lambda x \ y . c) & naturals, \ zero, \ successor, \ induction \\ \end{array}$$

# 2 Typing

 $\Gamma \vdash \Sigma x : A.B : \mathcal{U}_i$ 

 $\overline{\Gamma \vdash (a,b) : \Sigma x : A.B}$ 

 $\vdash \Gamma$ 

(Context well-formedness)

 $\Gamma \vdash a : \Sigma x : A.B$ 

 $\Gamma \vdash \pi_2 \ a : B[\pi_1 \ a/x]$ 

WF-NIL 
$$\frac{\text{WF-Cons}}{\vdash \Gamma} \qquad \frac{\vdash \Gamma \quad \Gamma \vdash A : \mathcal{U}_i}{\vdash \Gamma, x : A}$$

 $\Gamma \vdash a : A$  (Typing)

 $\Gamma \vdash a : \Sigma x \colon\! A \ldotp B$ 

 $\Gamma \vdash \pi_1 \ a : A$ 

WT-ZERO

WT-Succ

WT-NAT

### 3 Reductions

Full  $\beta$ ,  $\eta$ , and  $\beta\eta$ -reductions  $(\leadsto_{\beta}, \leadsto_{\eta}, \leadsto_{\beta\eta})$  are the congruence closures of the respective primitive reductions and we omit their definitions because they are standard.

$$a \rhd_{\beta} b$$

### (Primitive $\beta$ -Reduction)

$$\frac{\text{PB-ProjPair1}}{\pi_1(a,b) \rhd_{\beta} a}$$

$$\frac{\text{PB-ProjPair2}}{\pi_2(a,b) \rhd_{\beta} b}$$

#### PB-IndZero

$$\overline{\operatorname{ind} P \operatorname{zero} b (\lambda x \ y. \ c) \rhd_{\beta} b}$$

PB-IndSuc

 $(\lambda x. \overline{a) b \rhd_{\beta} a [b/x]}$ 

 $\overline{\operatorname{ind} P(\operatorname{succ} a) b(\lambda x y. c)} \triangleright_{\beta} c[a/x, \operatorname{ind} P a b(\lambda x y. c)/y]$ 

## $a \rhd_{\eta} b$

(Primitive 
$$\eta$$
-Reduction)

$$\frac{x \notin \mathbf{freevar}(a)}{\lambda x. \ a \ x \rhd_{\eta} a}$$

$$\frac{\text{PE-PAIRETA}}{(\pi_1 \ a, \pi_2 \ a) \rhd_{\eta} \ a}$$

$$a \Rightarrow_{\eta} b$$

(Parallel  $\eta$ -Reduction)

$$\frac{a_0 \Rightarrow_{\eta} a_1}{x \notin \mathbf{freevar}(a_0)} \frac{x \notin \mathbf{freevar}(a_0)}{(\lambda x. a_0 x) \Rightarrow_{\eta} a_1}$$

$$\begin{array}{lll} a_0 \Rightarrow_{\eta} a_1 & & \text{P-PAIRETA} & & \text{P-ABSCONG} \\ \underline{x \notin \mathbf{freevar}(a_0)} & \underline{a_0 \Rightarrow_{\eta} a_1} & & \underline{a_0 \Rightarrow_{\eta} a_1} & & \underline{a_0 \Rightarrow_{\eta} a_1} \\ (\lambda x. \ a_0 \ x) \Rightarrow_{\eta} a_1 & & (\pi_1 \ a_0, \pi_2 \ a_0) \Rightarrow_{\eta} a_1 & & \lambda x. \ a_0 \Rightarrow_{\eta} \lambda x. \ a_1 \end{array}$$

P-AbsCong
$$a_0 \Rightarrow_{\eta} a_1$$

$$\lambda x. \ a_0 \Rightarrow_{\eta} \lambda x. \ a_1$$

P-APPCONG
$$a_0 \Rightarrow_{\eta} a_1 \qquad b_0 \Rightarrow_{\eta} b$$

$$a_1 b_2 \Rightarrow a_1 b_2$$

P-ProjCong2
$$a_0 \Rightarrow_{\eta} a_1$$

 $\pi_2 a_0 \Rightarrow_{\eta} \pi_2 a_1$ 

$$\frac{\text{P-VAR}}{x \Rightarrow_n x}$$

$$\frac{\text{P-VAR}}{x \Rightarrow_n x} \qquad \frac{\text{P-UNIV}}{\mathcal{U}_i \Rightarrow_n \mathcal{U}_i}$$

$$\frac{\text{P-VAR}}{x \Rightarrow_{\eta} x} \qquad \frac{\text{P-UNIV}}{\mathcal{U}_{i} \Rightarrow_{\eta} \mathcal{U}_{i}} \qquad \frac{A_{0} \Rightarrow_{\eta} A_{1}}{B_{0} \Rightarrow_{\eta} B_{1}} \frac{B_{0} \Rightarrow_{\eta} B_{1}}{\Pi x : A_{0} . B_{0} \Rightarrow_{\eta} \Pi x : A_{1} . B_{1}}$$

#### P-SigCong

$$\frac{A_0 \Rightarrow_{\eta} A_1}{B_0 \Rightarrow_{\eta} B_1}$$
$$\frac{\sum x : A_0. B_0 \Rightarrow_{\eta} \sum x : A_1. B_1}{\sum x : A_1. B_1}$$

$$rac{ ext{P-Nat}}{ ext{Nat} \Rightarrow_{\eta} ext{Nat}}$$

P-IndCong

$$P_0 \Rightarrow_{\eta} P_1$$

$$a_0 \Rightarrow_{\eta} a_1$$

$$b_0 \Rightarrow_{\eta} b_1 \qquad c_0 \Rightarrow_{\eta} c_1$$

$$\mathbf{ind} \ P_0 \ a_0 \ b_0 \ (\lambda x \ y. \ c_0) \Rightarrow_{\eta} \mathbf{ind} \ P_1 \ a_1 \ b_1 \ (\lambda x \ y. \ c_1)$$

$$\mathbf{zero} \Rightarrow_{\eta} \mathbf{zero}$$

P-SUCCONG
$$\frac{a_0 \Rightarrow_{\eta} a_1}{\mathbf{succ} \ a_0 \Rightarrow_{\eta} \mathbf{succ} \ a_1}$$

$$a \Rightarrow_r b$$

(Restrictive  $\eta$ -Reduction (not eliminated))

EP-APPETA
$$a_{0} \notin \mathbf{canf}$$

$$a_{0} \Rightarrow_{\overline{r}} a_{1}$$

$$x \notin \mathbf{freevar}(a_{0})$$

$$(\lambda x. a_{0} x) \Rightarrow_{r} a_{1}$$
EP-PAIRETA
$$a_{0} \notin \mathbf{canf}$$

$$a_{0} \neq \mathbf{canf}$$

$$a_{0} \Rightarrow_{\overline{r}} a_{1}$$

$$(\pi_{1} a_{0}, \pi_{2} a_{0}) \Rightarrow_{r} a_{1}$$

$$x \notin \mathbf{free}$$

$$(\pi_{1} a_{0}, \pi_{2} a_{0}) \Rightarrow_{r} a_{1}$$

$$(\pi_{1} a_{0}, \pi_{2} a_{0}) \Rightarrow_{r} a_{1}$$

 $a \Rightarrow_{\bar{r}} b$ 

(Restrictive  $\eta$ -Reduction (eliminated))

$$\begin{array}{lll} \text{NEP-PiCong} & \text{NEP-SigCong} \\ A_0 \Rightarrow_r A_1 & A_0 \Rightarrow_r A_1 \\ B_0 \Rightarrow_r B_1 & B_0 \Rightarrow_r B_1 \\ \hline \Pi x \colon A_0 \colon B_0 \Rightarrow_{\bar{r}} \Pi x \colon A_1 \colon B_1 & \overline{\Sigma} x \colon A_0 \colon B_0 \Rightarrow_{\bar{r}} \Sigma x \colon A_1 \colon B_1 & \overline{\textbf{Nat}} \Rightarrow_{\bar{r}} \mathbf{Nat} \end{array}$$

NEP-IndCong

$$\begin{array}{c} P_0 \Rightarrow_r P_1 \\ a_0 \Rightarrow_{\bar{r}} a_1 \\ b_0 \Rightarrow_r b_1 \qquad c_0 \Rightarrow_r c_1 \\ \hline \mathbf{ind} \ P_0 \ a_0 \ b_0 \ (\lambda x \ y. \ c_0) \Rightarrow_{\bar{r}} \mathbf{ind} \ P_1 \ a_1 \ b_1 \ (\lambda x \ y. \ c_1) \\ \hline \\ NEP\text{-SucCong} \\ \underline{a_0 \Rightarrow_r a_1} \\ \hline \\ \mathbf{succ} \ a_0 \Rightarrow_{\bar{r}} \mathbf{succ} \ a_1 \\ \end{array}$$

# 4 Strong Normalization

 $a \in \mathbf{SN}$ 

(Strong Normal Forms)

$$\begin{array}{ll} \text{N-Abs} & & \text{N-Pair} \\ \underline{b \in \mathbf{SN}} & \underline{a \in \mathbf{SN}} & \underline{b \in \mathbf{SN}} & & \underline{A \in \mathbf{SN}} \\ \lambda x. \ \underline{b \in \mathbf{SN}} & & \underline{(a,b) \in \mathbf{SN}} & & \underline{\Pi x : A. \ B \in \mathbf{SN}} \end{array}$$

# 5 Coquand's Algorithm

 $a \leftrightarrow b$ 

(Algorithmic equality)

CE-RED
$$a \leadsto_h^* f_0$$

$$b \leadsto_h^* f_1 \qquad f_0 \sim f_1$$

$$a \leftrightarrow b$$

 $f_0 \sim f_1$ 

(Algorithmic equality for head normal forms)

 $A \ll B$ 

(Algorithmic subtyping)

$$\begin{array}{c} \text{CLE-Red} \\ A \leadsto_h^* f_0 \\ \underline{B \leadsto_h^* f_1} \quad f_0 \lesssim f_1 \\ \overline{A \ll B} \end{array}$$

 $\overline{\operatorname{ind} P_0 e_0 b_0 (\lambda x \ y. \ c_0) \sim \operatorname{ind} P_1 e_1 b_1 (\lambda x \ y. \ c_1)}$ 

$$f_0 \lesssim f_1$$

(Algorithmic subtyping for head normal forms)

### 5.1 Domains

The diamond operator (\$\diamond\$) in rules SA-Conf and A-Conf is meant for the catch-all cases where none of the other cases match (i.e. the input terms have conflicting weak-head normal forms that can never be convertible).

(Domain for Algorithmic Equality)

 $(a,b) \in \mathcal{A}$  (Domain for Algorithmic Equality (normal forms))

$$\frac{\text{A--Univ}}{(\mathcal{U}_{i}, \mathcal{U}_{j}) \in \mathcal{A}} \qquad \frac{\text{A--AbsAbs}}{(\lambda x. \ a, \lambda x. \ b) \in \mathcal{A}} \qquad \frac{\text{A--AbsNeu}}{(a, e \ x) \in \mathcal{A}} \qquad \frac{\text{A--NeuAbs}}{(e \ x, a) \in \mathcal{A}^{*}}$$

$$\begin{array}{lll} \text{A-PairPair} & \text{A-PairNeu} & \text{A-NeuPair} \\ (a_0, a_1) \in \mathcal{A}^* & (a_0, \pi_1 \, e) \in \mathcal{A}^* & (\pi_1 \, e, a_0) \in \mathcal{A}^* \\ (b_0, b_1) \in \mathcal{A}^* & (a_1, \pi_2 \, e) \in \mathcal{A}^* & (\pi_2 \, e, a_1) \in \mathcal{A}^* \\ \hline ((a_0, b_0), (a_1, b_1)) \in \mathcal{A} & ((a_0, a_1), e) \in \mathcal{A} & (e, (a_0, a_1)) \in \mathcal{A} \end{array}$$

$$\begin{array}{ll} \text{A-Zero} & \begin{array}{c} \text{A-Succ} \\ \hline (zero, zero) \in \mathcal{A} \end{array} & \begin{array}{c} \text{A-Succ} \\ \hline (succ \ a, succ \ b) \in \mathcal{A} \end{array} & \begin{array}{c} \text{A-Univ} \\ \hline (\mathcal{U}_i, \mathcal{U}_j) \in \mathcal{A} \end{array} \end{array}$$

A-P<sub>1</sub>

$$(A_0, A_1) \in \mathcal{A}^*$$

$$(B_0, B_1) \in \mathcal{A}^*$$

$$(\Pi x : A_0. B_0, \Pi x : A_1. B_1) \in \mathcal{A}$$
A-S<sub>1G</sub>

$$(A_0, A_1) \in \mathcal{A}^*$$

$$(B_0, B_1) \in \mathcal{A}^*$$

$$(\Sigma x : A_0. B_0, \Sigma x : A_1. B_1) \in \mathcal{A}$$

$$\begin{array}{lll} & \begin{array}{lll} \text{A-NAT} & \begin{array}{lll} & \text{A-PPCONG} \\ (e_0,e_1) \in \mathcal{A} \\ (a_0,a_1) \in \mathcal{A}^* \end{array} & \begin{array}{lll} & \text{A-PROJCONG1} \\ (e_0,e_1) \in \mathcal{A} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG1} \\ (e_0,e_1) \in \mathcal{A} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG1} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG1} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \\ (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \\ & (e_0,e_1) \in \mathcal{A} \end{array} & \begin{array}{lll} & \text{A-PROJCONG2} \end{array} & \begin{array}{lll} & \text{A-PROJCONG3} \end{array} & \begin{array}{lll} & \text{A-P$$

 $(h_0,h_1)\in\mathcal{S}$