

• If $f(n)$ is negligible, $\sqrt{f(n)}$ is neglig.

Let $p(n)$ be an arbitrary poly.

To prove: There exists a N s.t.

\exists for all $n > N$, $\sqrt{f(n)} < \frac{1}{p(n)}$

$f(n) < \frac{1}{p^2(n)}$ (use $f(n)$ is neglig.)
if $p(n)$ is poly, \leftarrow change of vars
so is $p^2(n)$

so $\exists N$, s.t. $n > N$ $f(n) < \frac{1}{p^2(n)}$ or $\sqrt{f(n)} < \frac{1}{p(n)}$

$f(n)$ is neglig.

$$G(s) = s \cdot \underbrace{1010 \dots}_{n \text{ bits}}$$

↓ concat

Is G a PRG?

$$l(n) = 2n > n$$

world 0

$$s \leftarrow \{0, 1\}^n$$

$$r = G(s)$$

give (r) to D

PRG

$$P[D(G(s)) = 1] = 1$$

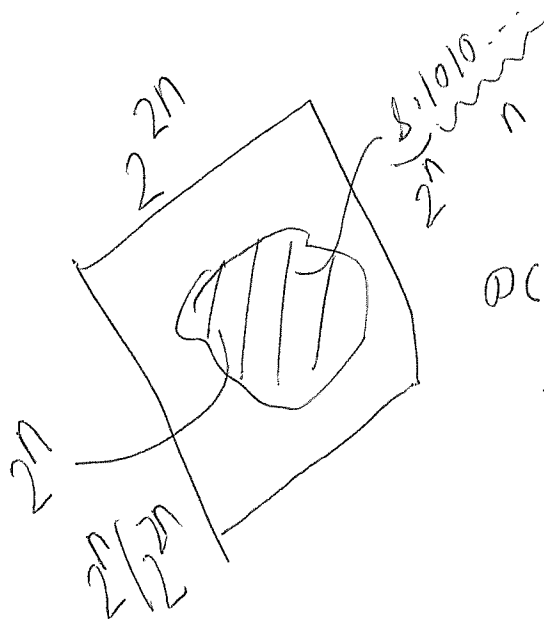
world 1

$$r \leftarrow \{0, 1\}^{2n}$$

give (r) to D

PRG
random

$$P[D(r) = 1] = \frac{2^n}{2^{2n}} = 2^{-n}$$

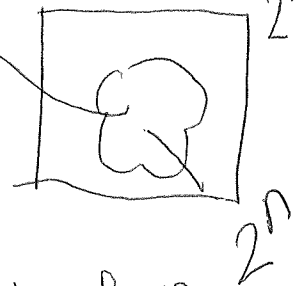


$$D(\omega) = \begin{cases} 1 \\ 0 \end{cases}$$

↑
2n bits

$$1 - 2^{-n} \neq \text{negl.}$$

$$s \cdot 1010 \dots$$



ω of the form
 $s \cdot 10 \dots$
otherwise

$$2^{-100 \log n} \quad \leftarrow \text{base 2}$$

$$(2^{-\log n})^{100} = \frac{1}{n^{100}} \neq \text{negl.}$$

$$\frac{1}{n^{100}} < \frac{1}{n^{200}}$$

$$\frac{1}{2\sqrt{n} + \underbrace{n^{1000}}_{\text{negl.}} + n^4} < \frac{1}{\underbrace{2\sqrt{n}}_{\text{negl.}}}$$

$$\underbrace{\frac{n^2 + n}{2^n}}_{\text{negl.}} + \frac{n^3 + n^4}{\underbrace{2\sqrt{n}}_{\text{negl.}}} = \text{negl.}$$

$\Pi = (\text{Gen}, \text{Enc}, \text{Dec})$

$k \leftarrow \text{Gen}(1^n)$

$b \leftarrow \{0, 1\}$

$c \leftarrow \text{Enc}_k(m_b)$

1^n

A

m_0, m_1

$|m_0| = |m_1|$

guess b'

output = $\begin{cases} 1 & b=b' \\ 0 & b \neq b' \end{cases}$

PPT

$\text{PrivK}_{A, \Pi}^{\text{eav}}(n)$

$\Pr[\text{PrivK}_{A, \Pi}^{\text{eav}}(n) = 1] \leq \frac{1}{2} + \text{negl}(n)$

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