

## Homework 5

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Due: Nov 14

## 1. Exercise 4.8

Let  $F$  be a pseudorandom function. Show that the following MAC for messages of length  $2n$  is insecure: Gen outputs a uniform  $k \in \{0, 1\}^n$ . To authenticate a message  $m_1 \| m_2$  with  $|m_1| = |m_2| = n$ , compute the tag  $F_k(m_1) \| F_k(m_2)$ .

2. In Lecture 21, we did the Merkle-Damgrad construction which took a hash function  $h$  that compresses by a factor of  $1/2$  and constructs a hash function  $H$  by a factor of  $1/3$ . Repeat the construction to construct a hash function  $H$  that compresses by a factor of  $1/5$ . Also redo the proof of collision resistance that was done in the Lecture.
3. Alice has six files  $F_1, F_2, F_3, F_4, F_5, F_6$  that she wants to store on Bob's computer (Bob just purchased a new server that has a gigantic hard disk). However, Alice is worried that Bob might corrupt or modify the files. Answer the following:
  - (a) Show a Merkle hash tree for  $F_1, F_2, F_3, F_4, F_5, F_6$  where the root is binary and the internal nodes are ternary. This shows that Merkle hash tree doesn't necessarily have to be binary.
  - (b) What is stored on Alice's computer?
4. Now Alice wants to retrieve file  $F_3$  from Bob's computer.
  - (a) What does Bob send to Alice? Recall that Bob needs to "prove" to Alice that the file has not been modified.
  - (b) Show that it is "hard" for Bob to generate a "proof" for Alice for a file  $F'_3$  different from  $F_3$ . We of course assume that hash functions that the Merkle hash tree is constructed from is *collision resistant*.