

## Homework 5

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Due: April 21

## 1. Exercise 4.8

Let  $F$  be a pseudorandom function. Show that the following MAC for messages of length  $2n$  is insecure: **Gen** outputs a uniform  $k \in \{0, 1\}^n$ . To authenticate a message  $m_1 \| m_2$  with  $|m_1| = |m_2| = n$ , compute the tag  $F_k(m_1) \| F_k(m_2)$ .

## 2. Exercise 4.1

Say  $\Pi = (\text{Gen}, \text{Mac}, \text{Vrfy})$  is a secure MAC, and for  $k \in \{0, 1\}^n$  the tag-generation algorithm  $\text{Mac}_k$  always outputs tags of length  $t(n)$ . Prove that  $t$  must be super-logarithmic or, equivalently, that if  $t(n) = \mathcal{O}(\log n)$  then  $\Pi$  cannot be a secure MAC.

**Note.** Super-logarithmic means not of the form  $\mathcal{O}(\log n)$ .

**Note.** Super-logarithmic means not of the form  $\mathcal{O}(\log n)$ , i.e. it is not the case that  $t(n) \leq C \log n$  for large  $n$  and some fixed  $C > 0$ .

**Hint:** Consider the probability of randomly guessing a valid tag.

3. Alice has five files  $F_1, F_2, F_3, F_4, F_5$  that she wants to store on Bob's computer (Bob just purchased a new server that has a gigantic hard disk). However, Alice is worried that Bob might corrupt or modify the files. Answer the following:

- (a) Show the Merkle hash tree for  $F_1, F_2, F_3, F_4, F_5$ .
- (b) What is stored on Alice's computer?

4. Now Alice wants to retrieve file  $F_3$  from Bob's computer.

- (a) What does Bob send to Alice? Recall that Bob needs to "prove" to Alice that the file has not been modified.
- (b) Show that it is "hard" for Bob to generate a "proof" for Alice for a file  $F'_3$  different from  $F_3$ . We of course assume that hash functions that the Merkle hash tree is constructed from is *collision resistant*.