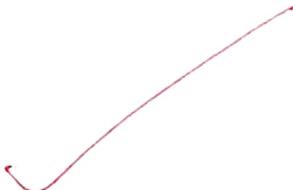


## Practical NO.-1

Aim :- Write a Program to implement caesar cipher

Objectives :-

1. To implement one of the encryption technique.
2. To use one of the Substitution technique.
3. To implement classical cipher



## Practical No. - 1

Page No. 2  
DATE / / 120

- Aim :- To demonstrate classical cipher.  
Inlab-aim :- Write a Program to implement Caesar cipher.

### Objectives :-

1. To implement one of the encryption technique.
2. To use one of the substitution technique.
3. To implement classical cipher.

### Theory :-

In cryptography, a Caesar Cipher, also known as Caesar's Cipher, the Shift cipher, Caesar's Code or Caesar Shift, is one of the simplest & most widely known encryption techniques. It is a type of Substitution cipher in which each letter in the Plaintext is replaced by a letter some fixed number of positions down the alphabet. For example, with a left shift of 3, D would be replaced by A, E would become B, f so on.

### Code :-

```
#include <stdio.h>
int main () {
    int k;
    char msg[100], chi;
    printf ("\n Enter message:");
    gets (msg);
    printf ("\n Enter key:");
    scanf ("%d", &k);
    int i=0;
```

Caesar cipher is a type of substitution cipher which replaces each letter by another letter shifted by a fixed number of positions in the alphabet. It is named after Julius Caesar who used it to protect his military communications.

The cipher works by shifting the letters of the alphabet by a fixed number of positions. For example, if we shift by 3, then A would be replaced by D, B by E, C by F, and so on. The resulting ciphertext is: D E F G H I J K L M N O P Q R S T U V W X Y Z A B C

The original message is: H E L L O W O R L D

After shifting by 3, the ciphertext is: K I M M U X R O M G

The cipher is considered secure because it requires knowledge of the shift value to decipher the message. However, it is vulnerable to frequency analysis attacks.

The cipher can be implemented using a simple algorithm. One way to do this is to use a mapping table where each letter is mapped to its corresponding shifted letter. For example, the mapping table for a shift of 3 would be:

Original Letter	Shifted Letter
A	D
B	E
C	F
D	G
E	H
F	I
G	K
H	L
I	M
J	N
K	O
L	P
M	R
N	S
O	T
P	U
Q	V
R	W
S	X
T	Y
U	Z
V	A
W	B
X	C
Y	D
Z	E

Conclusion 2— (In 200 words)  
Hence, we have successfully implemented Caesar cipher using classical cipher.

(Implementation of Caesar cipher)  
(Implementation of Caesar cipher)  
(Implementation of Caesar cipher)

```

for(i=0; msg[i]!='\0'; i++) {
    ch=msg[i];
    if(ch>='a' && ch<='z') {
        msg[i] = ch + K;
    }
    if(ch>'z')
        msg[i] = ch - 26;
}

```

```
{ printf("The Encrypted Message:");
```

```
    puts(msg);
```

```
-- for(i=0; msg[i]!='\0'; i++) {
```

```
    ch=msg[i];
```

```
    if(ch>='a' && ch<='z') {
```

```
        msg[i] = ch - K;
```

```
    if(ch<'a')
```

```
        msg[i] = ch + 26;
```

```
}
```

```
{ printf("the decrypted message:");
```

```
    puts(msg);
```

```
return 0;
```

```
}
```

Output :- Enter Message: attack is tomorrow

Enter Key : 3

The Encrypted Message: dwwdfn lv wprmuu

The decrypted Message: attack is tomorrow

(Q) Conclusion :-

Hence, we have successfully implemented  
Caesar cipher using classical cipher.

Done  
10/10

## Practical No. -2

Aim :- write a Program to implement Playfair cipher.

Objectives :-

1. To implement one of the encryption technique
2. To use one of the substitution technique.
3. To implement playfair cipher

## Practical No.-2

Page No. 4  
DATE / / 20

Aim :- To demonstrate Substitution ciphers.

Inlab Aim :- Write a program to implement Playfair cipher.

Objectives :-

1. To implement one of the encryption technique.
2. To use one of the Substitution technique.
3. To implement playfair cipher.

Theory :-

Playfair cipher is one of the substitution cipher technique. After that the blank cell is filled with letters A-Z which are not already there in the matrix. It works as follows:

1. Observe the given plaintext if it contains any consecutive repeated letters than insert a filler character 'X' in between the repeated letters.
2. observe the given plaintext if it contains odd numbers of letters after step-1 then add one filler letter 'X' at the end.
3. Make pair of characters to form a digram from given Plaintext.
4. ~~Encrypt the plaintext digrams as follows:~~  
  - a. If a diagram contains characters in same column then replace the character by character adjacent right cell of matrix.
  - b. If a diagram contains characters in same column then replace the characters by the character available in the adjacent cell below.

```

Code :- #include <stdio.h>
int main() {
    char arr[5][5] = { "MONAR", "CHYBD", "EFGIK",
                       "LPQST", "UVWxz" };
    char pt[10];
    int i, j, r1 = 0, r2 = 0, c1 = 0, c2 = 0;
    printf("Playfair keymatrix\n" == '\n');
    for (i = 0; i < 5; i++) {
        for (j = 0; j < 5; j++)
            printf("%c", arr[i][j]);
        printf("\n");
    }
    printf("Enter your Plaintext:");
    scanf("%s", pt);
    printf("Your plaintext = %s", pt);
    for (i = 0; i < 5; i++) {
        for (j = 0; j < 5; j++) {
            if (arr[i][j] == pt[0]) {
                r1 = i; c1 = j;
            }
        }
        if (arr[i][j] == pt[1]) {
            r2 = i; c2 = j;
        }
    }
    if (r1 == r2) {
        if (c2 == 4)
            printf("Ciphertext = %c%c\n",
                   arr[n][c1 + 1], arr[n][0]);
        else
            printf("Ciphertext = %c%c\n",
                   arr[r2][c2 + 1], arr[n][c1 + 1]));
    }
}

```

### Conclusion :-

Hence, we successfully demonstrated  
Substitution cipher & implemented playfair cipher

{ if ( $c_1 == c_2$ ) {

    if ( $r_2 == 4$ )

        printf ("Cipher text = %c%c\n",  
            arr [ $r_1 + 1$ ] [ $c_1$ ], arr [ $r_2$ ] [ $c_2$ ]);

    else

        printf ("Cipher text = %c%c\n", arr [ $r_1 + 1$ ] [ $c_1$ ],  
            arr [ $r_2 + 1$ ] [ $c_2$ ]);

}

if ( $r_1 \neq r_2 \text{ & } c_1 \neq c_2$ ) {

    printf ("\n Cipher text = %c%c\n", arr [ $r_1$ ] [ $c_2$ ], arr [ $r_2$ ] [ $c_1$ ]);

}

return 0;

}

Output :- M O N A R

C H Y · B D

E F G I K

L P Q S T

U V W X Z

Do it for long string

Enter your plain text: IN Short str

Your Plain text = IN

Conclusion :-

Hence, we successfully demonstrated  
Substitution cipher & implemented playfair cipher

Do it  
long  
str

## Practical No. - 3

Page No. 7  
DATE / / 20

Aim :- To demonstrate Transposition cipher.  
Inlab-Aim :- Write a program to implement Railfence Cipher.

### Objectives :-

1. To implement one of the encryption technique.
2. To use Transposition technique.
3. To implement Railfence cipher.

### Theory :-

- In the rail fence cipher, the plain-text is written downwards & diagonally on successive coils of an imaginary fence.
- When we reach the bottom coil, we traverse upwards moving diagonally. After reaching the top coil, the direction is changed again. Thus the alphabets of the message are written in a zig-zag manner.
- After each alphabet has been written, the individual rows are combined to obtain the cipher-text.

Code :- #include <stdio.h>

#include <string.h>

#include <stdlib.h>

int main()

int i, j, len, rails, count, code[100][1000];

char str[1000];

printf("Enter Secret Message\n");

ACE gets(str);

```
len = strlen(str);
printf("Enter no. of rails\n");
scanf("%d", &rails);
for(i=0; i<rails; i++) {
    for(j=0; j<len; j++) {
        code[i][j] = 0;
    }
}
count = 0;
j = 0;
while(j < len) {
    if(count % 2 == 0)
        for(i=0; i<rails; j++) {
            code[i][j] = (int)str[i];
            j++;
        }
}
else {
    for(i=rails-2; i>0; i++) {
        code[i][j] = (int)str[j];
        j++;
    }
}
count++;
}
for(i=0; i<rails; i++) {
    for(j=0; j < len; j++) {
        if(code[i][j] != 0)
            printf("%c", code[i][j]);
    }
    printf("\n");
}
return 0;
```

Output :- Enter Secret Message

Hello World.

Enter No. of rails

2

H O R E L O L L W S D

Conclusion :-

Hence, we successfully demonstrated  
transposition cipher & implement Rail fence cipher

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# Practical No. - 4

Page No. 10  
DATE / / 120

Aim :- To demonstrate block cipher technique.

Inlab Aim :- Write a Program to implement AES

## Objectives :-

1. To implement block cipher.
2. To implement AES

## Theory :-

AES stands for Advance encryption algorithm designed for encrypting a fixed size block of data. AES was designed by Rijmen-Daemen in Belgium & it uses 128/192/256 bit keys, 128 bit data. It processes data as block of 4 columns of 4 bytes operates on entire data block in every round. It designed to have resistance against known attacks, speed & code compactness on many CPUs, design simplicity. The data block of 4 columns of 4 bytes is called state.

- byte substitution.
- shift rows
- mix columns
- add round key
- ✓ • view as alternating XOR key & Scramble data bytes.

Code :- # include <stdio.h>

# include <stdlib.h>

# include <string.h>

# include <mcrypt.h>

```
#include <stdint.h>
#include <stdlib.h>
int encrypt(
    void * buffer,
    int buffer-len,
    char * IV,
    char * key,
    int key-len
){
```

```
MCRYPT_td = mcrypt_module_open("rijndael-128",
    NULL, "cbc", NULL);
int blocksize = mcrypt_enc_get_block_size(td);
if (buffer-len % blocksize != 0) { return 1; }
mcrypt_generic_init(td, key, key-len, IV);
mcrypt_generic(td, buffer, buffer-len);
mcrypt_generic_deinit(td);
mcrypt_module_close(td);
return 0;
}
```

~~```
int decrypt(
    void * buffer,
    int buffer-len,
    char * IV,
    char * key,
    int key-len
){
```~~
~~```
MCRYPT_td = mcrypt_module_open("rijndael-128", NULL,
```~~

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```

    ("bc", NULL);
int blocksize = mCRYPT-enc-get-block-size(td);
if (buffer_len % blocksize != 0) { return 1; }
mCRYPT-generic-init(td, key, key_len, IV);
mCRYPT-generic(td, key, key_len, IV);
mCRYPT-generic-deinit(td);
mCRYPT-module-close(td);
return 0;
}

```

```

void display(char* ciphertext, int len) {
    int v;

```

```

    for (v = 0; v < len; v++) {
        printf("%d", ciphertext[v]);
    }
    printf("\n");
}

```

```

int main() {

```

```

    MCRYPT td, td2;

```

```

    char* plaintext = "test text 123";

```

```

    char* IV = "AAAAAAAAAAAAAAA";

```

```

    char* key = "0123456789abcdef";

```

~~int keyszie = 16;~~

```

    char* buffer;

```

```

    int buffer_len = 16;

```

```

    buffer = (char*) malloc(1, buffer_len);

```

```

    strcpy(buffer, plaintext, buffer_len);

```

```

    printf("==c==\n");

```

```

    printf("plain: %s\n", plaintext);

```

```

    encrypt(buffer, buffer_len, IV, key, keyszie);

```

```

    printf("cipher: ");

```

```
    display(buffer, buffer_len);
    decrypt(buffer, buffer_len; IV, key, keysize);
    printf("decrypt : %s\n", buffer);
    return 0;
}
```

Output :-

Plain : Test Text 123

Cipher : 16-12441-33-16-12361-64-15-

Decrypt : Test Text 123

Conclusion :-

Hence, we successfully demonstrated  
block cipher by implementing AES algorithm.

Ques  
Ans  
20/07/20

# Practical No.-5

Page No. 14

Date: / /

Aim :- To demonstrate Symmetric secret key cipher.  
Inlab Aim :- Write a program to implement DES.

## Objectives :-

1. To implement Symmetric Secret key cipher
2. To implement DES
3. To use DES in real world examples.

## Theory :-

The DES encryption algorithm takes a 64-bit block of Plaintext & a 64-bit key as input & produces an 64-bit block of ciphertext as output. The DES decryption algorithm takes an 64-bit block of ciphertext & the same 64-bit key used to produce that ciphertext as input & produces the original 64-bit block of plaintext. The encryption algorithm involves function: an initial permutation (IP); a complex function labeled  $f_K$ , which involves expansions, S-Boxes, permutation, XOR & switching 32 bits halves of the data; the function  $f_K$  again; & finally a permutation function that is the inverse of the initial permutation (IP-1).

Code :-

```
#include <iostream>
#include <string.h>
#include <stdio.h>
#include <math.h>
using namespace std;
```

```
Void setB(int s; int in; int loc) {
```

```
    int h1=0;
```

```
    h1 = 1 << B1, -(loc-1);
```

```
    if(in);
```

```
        i = i & h1;
```

```
    else
```

```
        i = i & -h1;
```

```
} int getB(int i, int loc) {
```

```
    int sum=0;
```

```
    int t=0;
```

```
    for(int i=e; i>=b; i--) {
```

```
        sum += getB(in, i) << t;
```

```
        t += 1;
```

```
}
```

```
return sum;
```

```
} Void PKEY(int *k, int *r)
```

```
    int sum=0;
```

```
    int t=0;
```

```
    for(int i=e; i>=b; i--) {
```

```
        sum += getB(in, i) << t;
```

~~int map[56] = {57, 49, 41, 33, 25, 27, 9, 1, 58,~~

~~50, 42, 34, 26, 18, 10, 2, 59, 51, 43, 35, 27, 19, 11, 3, 60,~~

~~52, 44, 36, 63, 55, 47, 39, 31, 23, 46, 38, 30, 22};~~

```
    int hr=t;
```

```
    int hr=r;
```

```
    for(int t=0; t<56; t++) {
```

```
        if(map[t]>32) {
```

```
            if(t>27) {
```

```
                setB(hr, getB(*k, map[t]-32), t-27);
```

{else}

SetB(h1, getB(&.map[t]-32, ++));

}

{else}

if (t > 27) {

SetB(hc, getB(l, map[t]); t-27);

{else}

SetB(h1, getB(l, map[t], ++));

}

z = hr;

S = h1;

z = z >> 4;

z = z << 4;

z = z >> 8;

l = l << 8;

{void funct(int &l, int &r) {

int map[48] = {32, 1, 2, 3, 4, 5, 6, 7, 8, 9, 8, 9, 10, 11, 12,  
 13, 12, 13, 12, 13, 14, 15, 16, 17, 16, 17, 18, 19, 20, 21, 20,  
 21, 22, 23, 24, 24, 25, 26, 27, 28, 29, 30};

~~int h1 = l;~~

~~int h2 = l;~~

for (int t=0; t<48; ++) {

if (t > 23) {

setB(hc, getB(&.map[t], t-23))

{else}

SetB(h1, getB(&.map[t], t+1));

.

z = hr;

l = hl;

$x = x >> 8;$

$x = x << 8;$

$l = l >> 8;$

$l = l << 8;$

{ void final P(int &l, int &r) {

int map[64] = {40, 8, 48, 16, 56, 24, 64, 32, 39, 7, 47, 15, 35, 23, 31, 38, 6, 46, 14, 54, 22, 62, 30, 37, 5, 45, 13, 53, 21, 61, 29, 4, 44, 12, 52, 20, 60, 18, 33, 33}

int h1 = l;

int hr = r;

for (int t = 0; t < 64; t++) {

if (map[t] > 32) {

if (t > 31) {

SetB(hc, getB(r, map[t] - 32), t - 3));

} else {

SetB(hl, getB(r, map[t] - 32), t + 1));

}

else {

if (t > 31) {

SetB(hr, getB(l, map[t]), t - 31));

} else {

SetB(h, getB(r, map[t]), t + 1));

}

r = hc;

l = hl;

{ void Cipher(long, long int msg, long long int key) {

int f = 0;

int key\_r = (int)(key);

int key\_l = (int)(key >> 32);

Count << "key left int := " << key1 << "key right  
 int := " << keys << endl;

int l = 0, r = 0;  
 key (key l, key r);

l = key1;  
 r = key2;

Count << "msg left int := " << l << "msg  
 right int := " << r << endl;

p[i] = l, r;

for (int i = 1; i < 17; i++) {

int t = 2;

if (i == 11 || i == 2 || i == 9 || i == 16) {  
 t = 1;

} K1 = 10;

K2 = do;

PSUBKEY (K1, K2);

int plo = do;

lo = ro;

ro = plo funct (ro, K1, K2);

} int rl = 0;

rt = lo;

ro = lo;

lo = rt;

final P (lo, ro)

long long int final = lo;

final = final << 32;

final t = ro;

} int main () {

long long int keyIn [8] = {56, 18, 12, 11, 10, 107, 101,

1213}

```

long long int msga[8] = {109, 101, 115, 115, 97, 103, 101, 48};
long long int msg = 0;
long long int key = 0;
for (int i=0; i<8; i++) {
    msg += (msga[i] << ((7-i)*8));
}
for (int i=0; i<8; i++) {
    key += (keya[i] << ((7-i)*8));
}
cout << "[Your msg: " << msg << "] " << endl;
cout << "[Your key: " << key << "] " << endl;
Cipher (msg, key);
return 0;
}

```

Output:-

[Your msg : 7882833662174520622]

[Your key : 40629433546663103081]

Key left int := 945977716

Key right int := 1701537

Key left int := 4835365235

msg right int := 16341660

Conclusion :-

Hence, we successfully implemented DES.

(d) algorithm.

done  
Topic 10

# Practical No. - 6

Page No. 20

Date : / /

Aim :- To demonstrate asymmetric public key cipher  
Inlab Aim :- Write a program to implement RSA algorithm.

## Objectives :-

- To implement asymmetric public key cryptography
- To implement RSA algorithm.

## Theory :-

• It is an asymmetric cryptographic algorithm  
Asymmetric means that there are two different keys. This is also called public key cryptography, because one of them can be given to everyone. The other key must be kept private.

## Code :-

```
#include <iostream.h>
#include <math.h>
using namespace std;
int gcd(int a, int b)
{
    int temp;
    while(b != 0)
    {
        temp = a % b;
        if(temp == 0)
            return b;
        a = b;
        b = temp;
    }
}
```

```

int main() {
    double p=3;
    double q=4;
    double n=p*q;
    double count;
    double totient=(p-1)*(q-1);
    double e=2;
    while(e < totient) {
        count=gcd(e,totient);
        if(count==1)
            break;
        else
            e++;
    }
    double d;
    double k=2;
    d=(l+(k*totient))/e;
    double msg=12;
    double c=pow(msg,e);
    double m=pow(c,d);
    c=fmod(c,n);
    m=fmod(m,n);
}

cout << "message data = " << msg;
cout << "\n" << "p = " << p;
cout << "\n" << "q = " << q;
cout << "\n" << "n = p * q = " << n;
cout << "\n" << "totient = " << totient;
cout << "\n" << "e = " << e;
cout << "\n" << "encrypted data = " << c;
cout << "\n" << "original message sent = " << m;

```

Section 0;

{

Output :-

Message data = 12

$p = 3, q = 7$

$n = pq = 21$

totient  $\phi = 12$

$e = 5$

$d = 5$

encrypted data = 3

Original message sent = 12

Conclusion :-

Hence, we successfully demonstrated  
asymmetric public key cipher & implemented RSA.  
algorithm.

Copy  
Done  
Total 1/20

# Practical No.- f.

Aim :- To demonstrate secure key exchange.

Inlab Aim :- Write a program to implement Diffie-Hellman key exchange.

Objectives :-

- o To use secure key exchange
- o To encrypt & decrypt key Diffie-Hellman key exchange.

Theory :-

Diffie-Hellman key exchange, also called exponential key exchange, is a method of digital encryption that uses numbers raised to specific powers to produce decryption keys on the basis of components that are never directly transmitted, making the task of a would-be code breakers mathematically overwhelming.

To implement Diffie-Hellman the two end users Alice & Bob, while communicating over a channel they know to be private, mutually agree on positive whole members  $\cdot P$  &  $q$ , such that  $P$  is the prime number &  $q$  is a generator of  $P$ . The generator  $q$  is a number that, when raised to positive whole-number powers less than  $P$ , never produces the same result for any two such whole numbers. The values of  $P$  may be large but the value of  $q$  is usually small.

```

code :- #include <stdio.h>
int compute (int a, int m, int n) {
    int x;
    int y = 1;
    while (m > 0) {
        if (x == 1)
            y = (y * a) % n;
        a = a * a % n;
        m = m / 2;
    }
    return y;
}
int main () {
    int p = 23, g = 5, a, b, A, B, c = 6;
    A = compute (g, a, p);
    b = 15;
    B = compute (g, b, p);
    int key_A = compute (B, a, p);
    int key_B = compute (A, b, p);
    printf ("A's Secret Key is %d\nB's Secret Key is %d", key_A, key_B);
    return 0;
}

```

~~Output :-~~

A's Secret Key is 2.  
B's Secret Key is 4

Conclusion :-

Hence, we successfully demonstrated &  
implemented Diffie-Hellman key exchange algorithm

## Practical No. - 8

Aim :- To demonstrate hash technique for Security  
Inlab Aim :- Write a program to implement message digest.

### Objectives :-

1. To use message digest.
2. To produce message digest for given data.
3. To implement any message digest algorithm.

### Theory :-

Message digest is created by applying one of the message digest algorithm. The message digest can be calculated with the help of MDs, SHA, RIPEMD-10 etc algorithm.

Message digest algorithm takes a block of data of variable size or fixed size & produces a fixed size message digest output. Any message digest algorithm does not uses key but in place of that it uses many constant available in registers.

Code :- #include <stdlib.h>

#include <stdio.h>

#include <string.h>

#include <math.h>

typedef union NWB {

    unsigned w;

    unsigned char b[4];

{ mus Union;

unsigned func (unsigned abcd [ ]){

return (abcd[0] & abcd[2]) || (n abcd[1] & abcd[3]);

{ unsigned func1 (unsigned abcd [ ]){

return (abcd[3] & abcd[1]) || (abcd[3]& abcd[2]);

{ typedef unsigned (\* Digest funct) (unsigned a [ ]);

unsigned \* calculate (unsigned \* k){

double s, pwr;

int i;

pwr = Pow(2, 32);

for (i=0; i<64; i++) {

s = false (sin (i++));

k[i] = unsigned (s \* pwr);

} return k;

{ unsigned \* mod S (const char \* msg, int mlen){

static Digest Array h;

Digest Array abcd;

Digest fetc fetc;

short \* roth;

union {

unsigned w [16];

char b [64];

}; m m;

int as=0;

int grp, grps, g, P;

unsigned char p msg 2;

for (grp=0; grp < grps ; grp++) {

mmCpy ( mm, b, msg 2 + as :: 64 );

```

for (q=p; q=0; q < 16; q++) {
    q = (m * q * 0) % 16;
    abcd[0] = abcd[3];
    abcd[3] = abcd[2];
    abcd[2] = abcd[1];
    abcd[1] = f;
}

```

}

[P] += abcd[P];

OS+=64;

return h;

int main () {

int j, k;

const char \*msg = message digest:";

printf ("Your hashcode generated by mds is: %s\n");

return 0;

}

Output:-

Your hashcode generated by mds is:

2 3 4 b g c 1 7 8 b 2 c g b e 4 e g g b l e i d  
5 b c b b f f 6

Conclusion :-

Hence, we successfully demonstrated hash technique for security & implemented message digest.

Ques 2/2

## Practical No. - 9

Page No. 28

Date: / /

Aim :- To demonstrate classical cryptography  
In lab Aim :- write a Program to implement Hill Cipher

### Objectives :-

1. To implement one of the encryption technique
2. To implement one of the classical Substitution cipher technique
3. To implement Hill cipher.

### Theory :-

Invented by Lester S. Hill in 1929, the Hill cipher is a Substitution cipher based on linear algebra. Hill used matrices & matrix multiplication to mix up the plaintext. To encrypt a message, each block of  $n$  letters is multiplied by an invertible  $n \times n$  matrix, against modulus 26. To decrypt the message, each block is multiplied by the inverse of the matrix used for encryption.

### Code :-

```
#include <iostream>
#include <math.h>
using namespace std;
float encrypt[3][1], decrypt[3][1],
a[3][3], b[3][3], mes[3][1], c[3][3];
void encryption();
```

```

Void decryption();
Void inverse();
int main() {
    get key message();
    Encryption();
    decryption();
}

```

```

Void encryption() {
    int i, j, k;
    for(i=0; i<3; i++)
        for(j=0; j<1; j++) {
            for(k=0; k<3; k++)
                encrypt[i][j] = encrypt[i][j] + a[i][k]
                * mes[k][j];
            cout << "The encrypted string is : ";
            for(i=0; i<3; i++)
                cout << (char) (fmod(encrypt[i][0], 16)+97);
}

```

```

Void decryption() {
    int i, j, k;
    inverse();
    for(i=0; i<3; i++)
        for(j=0; j<1; j++) {
            for(k=0; k<3; k++)
                decrypt[i][j] = decrypt[i][j] + b[i][k]
                * encrypt[k][j];
            cout << "The decrypted string is : ";
            for(i=0; i<3; i++)
                cout << (char) (fmod(decrypt[i][0], 26)+97);
}

```

```

cout << "\n";
} void getKeyMessage() {
    int i, j;
    char msg[3];
    cout << "Enter 3x3 matrix for key : \n";
    for (i=0; i<3; i++)
        for (j=0; j<3; j++)
            cin >> a[i][j];
    c[i][j] = a[i][j];
}
void inverse() {
    int i, j, k;
    float p, q;
    for (i=0; i<3; i++)
        for (j=0; j<3; j++) {
            if (i==j)
                b[i][j]=1;
            else
                b[i][j]=0;
        }
    for (k=0; k<3; k++) {
        for (i=0; i<3; i++) {
            p = c[i][k];
            q = c[k][k];
            for (j=0; j<3; j++) {
                if (i != k) {
                    c[i][j] = c[i][j] * q - p * c[k][j];
                    b[i][j] = b[i][j] * q - q * b[k][j];
                }
            }
        }
    }
}
}

```

```

cout << "\n\n Inverse matrix is :\n";
for(i=0; i<3; i++) {
    cout << b[i][j] << " ";
    cout << "\n";
}

```

Output :- Enter 3x 3 matrix for key

$$\begin{matrix} 6 & 24 & 1 \\ 13 & 16 & 10 \\ 20 & 17 & 15 \end{matrix}$$

Enter a 3 letter string : act

Encrypted String is : Poh

Inverse matrix is :

$$\begin{matrix} 0.15873 & -0.777228 \\ -0.0113379 & 0.15873 \\ -0.22449 & 0.057193 \end{matrix}$$

Decrypted string is : act.

Conclusion :-

Hence, we successfully demonstrated classical cryptography & implemented Hill cipher

(2)

Done  
06/03/20

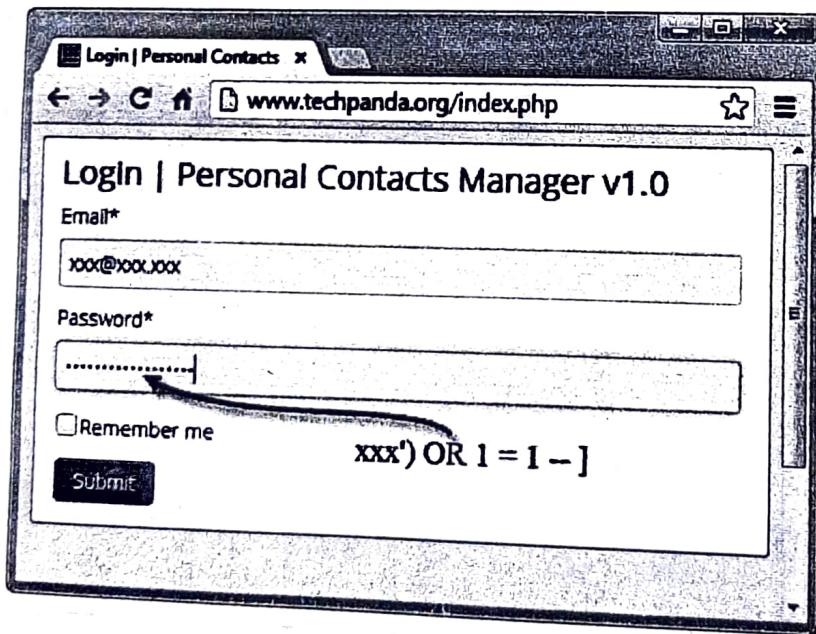
## Practical No. - 10

Aim :- Write a Program to implement SQL injection.

Objectives :-

- o To implement software vulnerability.
- o To implement SQL injection.

O/P :-



## Practical No. - 10

Aim :- To demonstrate Software Vulnerability  
Inlab Aim :- Write a program to implement SQL injection.

### Objectives :-

1. To implement Software Vulnerability
2. To implement SQL injection.

### Theory :-

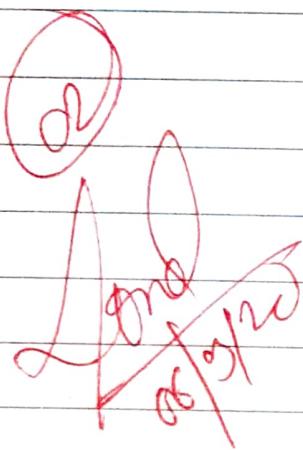
SQL injection is a technique of attacking on a web portal to gain access to the portal without knowing the credentials of the portal access. It uses the coding properties of SQL & finds a weak point to gain access to the portal without the knowledge of user id & user password. It also used to test the security of a website or database.

### Code :-

```

UName = getRequestString (" Username ");
UPass = getRequestString (" Userpassword ");
Sql = "SELECT * from USER WHERE Name= " "
Sql or=" = " AND Pass= " " OR " = "
Sql = "SELECT * from USER WHERE
      Name= " " + UName + " " AND
      Pass = " " + UPass + " "
  
```

Conclusion :- Hence, we successfully demonstrated the program for SQL injection.



# Practical No. - 11

Page No. 34

Date : / /

Aim :- To demonstrate Various encryption Techniques using GUI

Inlab Aim :- write a program to implement key agreement using Elliptic Curve based Algorithm.

Objectives :-

- To understand the concept of Elliptic curve based algorithm.
- To demonstrate Various encryption Techniques.

Theory :-

Elliptic-curve cryptography is an approach to Public-key Cryptography based on the algebraic structure of elliptic curves over finite fields.

ECC requires smaller keys compared to non-ECC cryptography to provide equivalent security.

Elliptic curves are applicable for key agreement, digital signatures, pseudo random generation & other tasks. elliptic curves that have application in cryptography.

Code :-

```
#include <stdio.h>
```

```
#include <math.h>
```

```
#define PI
```

```
typedef struct {double x, y } pt;
```

```
pt zero(void) { return (pt){INFINITY,
```

```
INFINITY};}
```

```
int is_zero(pt p) {return !(p.x < -1e20);}
```

```

Pt msg (Pt p) {
    return (Pt) { p.x - p.y };
}
    db (Pt p) {
        if (iszero (p))
            return p;
        Pt x;
    }

```

$$\text{double } l = (3 \cdot p.x \cdot p.x) / (2 \cdot p.y);$$

$$x.x = l \cdot l - p.x \cdot p.x;$$

$$x.y = l \cdot (p.x - x.x) - p.y;$$

$$\text{return } x;$$

3

```

Pt add (Pt p, Pt q) {
    if (p.x == q.x && p.y == q.y) return db(p);
    if (iszero(p)) return q;
    if (iszero(q)) return p;
    Pt x;
    double l = (q.y - p.y) / (q.x - p.x);
    x.x = l * l - p.x * q.x;
    x.y = l * (p.x - x.x) - p.y;
    return x;
}

```

3

```

Pt mul (Pt p, int n) {
    int i;

```

$$\text{Pt } r = \text{zero}();$$

$$\text{for } (i=1; i \leq n; i++)$$

$$\text{if } (i \neq 1) \quad r = \text{add} (r, p);$$

$$p = \text{db} (p);$$

3 return r;

3 void show (const char \*s, Pt p) {
 printf ("%s", s);
 if (iszero(p)) "zero" n": "(%f, %f)\n", px, py);

$\{ \text{pt from-y (double y)} \}$

$\text{Pt } z;$

$z \cdot n = \text{pow}(y * y \cdot c, 10/3);$

$z \cdot y = y;$

$\text{return } z;$

$\}$

$\text{int main(Void)}$

$\text{Pt } a, b, c, d;$

$a = \text{from-y}(1);$

$b = \text{from-y}(2);$

$\text{show}( "a = ", a);$

$\text{show}( "b = ", b);$

$\text{show}( "c = a + b = ", c = \text{add}(a, b));$

$\text{show}( "d = -c = ", d = \text{neg}(c));$

$\text{show}( "c + d = ", \text{add}(c, d));$

$\text{show}( "a + b + d = ", \text{add}(a, \text{add}(b, d)));$

$\text{show}( "a + 12345 = ", \text{mod}(a, 12345));$

$\text{return } 0;$

$\}$

Output :-

$a = (-1.817, 1.000)$

$b = (-1.442, 2.000)$

$c = a + b = (10.375, -33.525)$

$d = -c = (10.375, 33.525)$

$c + d = \text{zero}$

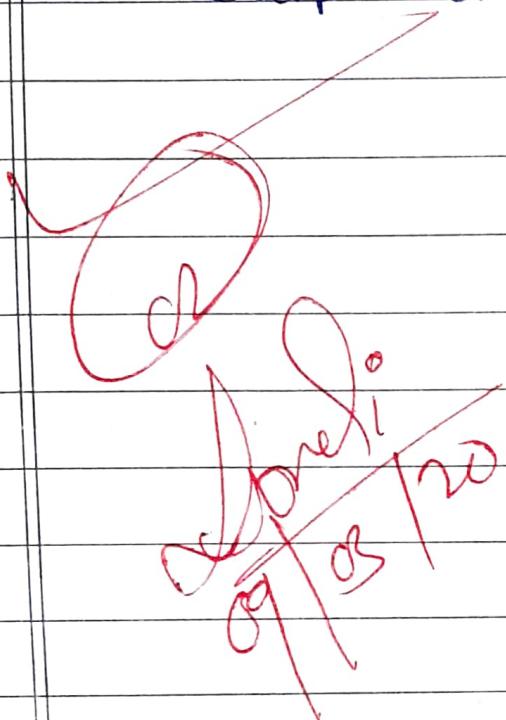
$a + b + d = \text{zero}$

$a + 12345 = (10.789, 33.387)$

3 pt from y (double y) {

Conclusion :-

Thus, we have implemented  
Elliptical curve cryptography.



## Practical No. - 12

Aim :- To demonstrate user authentication

In Lab Aim :- WAP to implement user authentication based on user id & password.

Theory :-

Authentication is the act of confirming the truth of a single piece of data claimed true by an entity in context with identification which refers to the act of stating or otherwise indicating a claim professed by at testing to a person on things identity authentication is the process of confirming that identity. It might involve verifying the authenticating of a website with digital certificate data mining the age of an user's artifact by carbon dating or ensuring that a product as what its packaging of labeling claim to be in other words , authenticating after involves verifying the validity of atleast one of identification

Code :-

```
#include <stdio.h>
int main () {
    char Password [10], Username [10], ch;
    int i;
    printf ("Enter User Name:");
    gets ( Username);
```

```

printf("Enter the Password:");
for(i=0; i<8; i++){
    ch = getchar();
    password[i] = ch;
    ch = '*';
    printf("%c", ch);
}
password[i] = '\0';
printf("\n Your Password is:");
for (i=0; i<8; i++){
    printf("%c", password[i]);
}
return 0;
}

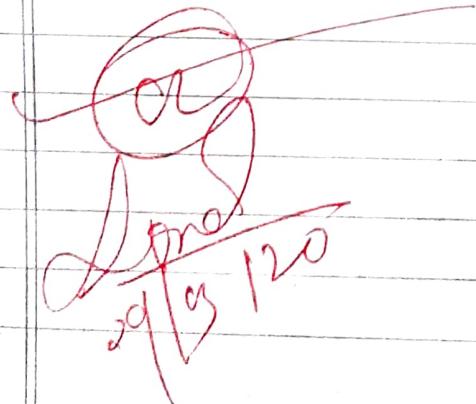
```

Output :-

Enter User Name: Damini @123  
 Enter the Password [any 8 characters]: - \* \* \* \* \* \* \* \*  
 Your Password is: damini @124 \* \* \* \* \*

Conclusion :-

Hence, we successfully implemented user authentication based on user id & Password.



## Practical No. -13

Aim :- To implement wireshark Open Source Tool

Objectives :-

- To implement wireshark open source tool.
- To know the use of wireshark tool.

Theory :-

Wireshark is a free & open-source packet analysis tool. It is used for network troubleshooting, analysis, s/w & communicating Protocol.

Wireshark intercepts traffic & converts that binary traffic into human-readable format. This makes it easy to identify what traffic is crossing your network, how much of it? how frequently, how much latency there is between certain hops, so forth.

While Wireshark supports more than two thousand network protocols, many of them are uncommon, or old, the modern security professional will find analyzing IP packets to be of most immediate usefulness. The majority of the packets on your network are likely to be TCP, UDP, & ICMP.

Given the large volume of traffic are what make it easily useful.

Name of Practical

## Features :-

Wireshark is a data capturing program that understands the structure of different networking protocols. Wireshark uses Pcap to capture packets. So it can only capture packets on the types of network that Pcap supports.

- 1) Data can be captured from the wire from a live network connection or read from a file or already-captured packets.
- 2) Captured N/W later can be treated via GUI
- 3) Data display can be refined using a display filter.
- 4) Plug-in can be created for displaying new protocols.
- 5) VoIP calls in the captured traffic can be deleted.
- 6) Wireless connection can be also seen as long as they traverse the ethernet.

Output :- On the blank page

Conclusion :- Hence, we have successfully implemented Wireshark tool.

Apply a display filter - <Ctrl-f>

| No. | Time                  | Source          | Destination     | Protocol | Length                                                                          | Info                                                                                  |
|-----|-----------------------|-----------------|-----------------|----------|---------------------------------------------------------------------------------|---------------------------------------------------------------------------------------|
| 343 | 65.142415             | 192.168.0.21    | 174.129.249.228 | TCP      | 66                                                                              | 405555 → 80 [ACK] Seq=1 Ack=1 Win=5888 Len=0 Tsvval=551811827                         |
| 344 | 65.142215192.168.0.21 | 174.129.249.228 | HTTP            | 253      | GET /clients/netflix/flash/application_swf>flash_version=flash_lite_2.1&v=1.58n |                                                                                       |
| 345 | 65.230738             | 174.129.249.228 | 192.168.0.21    | TCP      | 66                                                                              | 80 → 405555 [ACK] Seq=1 Ack=188 Win=6864 Len=0 Tsvval=551811850 Tscr=491519347        |
| 346 | 65.249742             | 174.129.249.228 | 192.168.0.21    | TCP      | 828                                                                             | HTTP/2.1 302 Moved Temporarily                                                        |
| 347 | 65.241592             | 192.168.0.21    | 174.129.249.228 | TCP      | 66                                                                              | 405555 → 80 [ACK] Seq=188 Ack=63 Win=7424 Len=0 Tsvval=491519446 Tscr=551811852       |
| 348 | 65.242532             | 192.168.0.21    | 192.168.0.1     | DNS      | 77                                                                              | Standard query 0x2188 A cdn-0.netflix.com                                             |
| 349 | 65.276870             | 192.168.0.1     | 192.168.0.21    | DNS      | 489                                                                             | Standard query response 0x2188 A cdn-0.netflix.com CNAME images.netflix.com edge      |
| 350 | 65.277992             | 192.168.0.21    | 63.80.242.48    | TCP      | 74                                                                              | 37063 → 80 [SFIN] Seq=0 Win=5840 Len=0 MSS=1460 SACK_PERM=1 Tsvval=491519482 Tscr=351 |
| 351 | 65.297757             | 63.80.242.48    | 192.168.0.21    | TCP      | 74                                                                              | 80 → 37063 [SFIN] ACK Seq=0 ACK=1 Win=5792 Len=0 MSS=1460 SACK_PERN=1 Tsvval=329      |
| 352 | 65.298396             | 192.168.0.21    | 63.80.242.48    | TCP      | 96                                                                              | 37063 → 80 [ACK] Seq=1 Ack=1 Win=5888 Len=0 Tsvval=491519502 Tscr=3295534130          |
| 353 | 65.298687             | 192.168.0.21    | 63.80.242.48    | HTTP     | 153                                                                             | GET /us/nord/clients/flash/814540.html HTTP/1.1                                       |
| 354 | 65.318730             | 63.80.242.48    | 192.168.0.21    | TCP      | 66                                                                              | 80 → 37063 [ACK] Seq=1 Ack=88 Win=5792 Len=0 Tsvval=3295534151 Tscr=491519503         |
| 355 | 65.321733             | 63.80.242.48    | 192.168.0.21    | TCP      | 154                                                                             | [TCP segment of a reassembled PDU]                                                    |

> frame 349: 489 bytes on wire (3912 bits), 489 bytes captured (3912 bits)  
 > Ethernet II, Src: Globals<00:3b:0a (fe:00:3b:0a), Dst: Vito\_14:3a:e1 (00:19:9d:14:3a:e1)  
 > Internet Protocol Version 4, Src: 192.168.0.1, Dst: 192.168.0.21  
 > User Datagram Protocol, Src Port: 53 (53), Dst Port: 34936 (34936)  
 > Domain Name System (response)  
 [Request In: 348]  
 [Time: 0.034338200 seconds]  
 Transaction ID: 0x2188  
 > Flags: 0x8100 Standard query response, No error  
 Questions: 1  
 Answer RRs: 4  
 Authority RRs: 9  
 Additional RRs: 9  
 ✓ Queries  
 > cdn-0.netflix.com: type A, class IN  
 > Answers  
 > Authoritative nameservers

|      |                                                 |             |
|------|-------------------------------------------------|-------------|
| 0026 | 80 15 00 35 84 f4 01 c7 83 3f f1 3e 81 80 00 01 | ... 5 ... ? |
| 0038 | 00 04 00 00 00 00 00 05 63 64 6c 2d 30 07 62    | ... 7 ... ? |
| 0048 | 78 69 6d 67 03 63 6f 6d 00 00 01 00 01 00 00    | ... 7 ... ? |
| 0050 | 05 09 91 00 00 05 29 00 22 00 6d 61 67 65 73    | ... 7 ... ? |
| 0058 | 07 6c 65 74 66 6c 69 78 03 63 6f 6d 09 65 64    | ... 7 ... ? |
| 0070 | 65 73 75 69 74 65 03 6e 65 74 00 50 2f 00 05 00 | ... 7 ... ? |

File Edit View Go Capture Analyze Status Help

| No. | Time     | Source      | Destination     | Protocol | Info                                     |
|-----|----------|-------------|-----------------|----------|------------------------------------------|
| 4   | 1.025659 | 192.168.0.2 | igrp.mcast.net  | IGMP     | V3 Membership Report                     |
| 5   | 1.044366 | 192.168.0.2 | 192.168.0.1     | DNS      | Standard query SRV _Iab_Support_         |
| 6   | 1.048652 | 192.168.0.2 | 239.255.255.250 | UDP      | Source port: 3193 Destination 00         |
| 7   | 1.050784 | 192.168.0.2 | 192.168.0.1     | DNS      | Standard query SOA ns1.oldid.m004        |
| 8   | 1.055053 | 192.168.0.1 | 192.168.0.2     | UDP      | Source port: 1900 Destination 00         |
| 9   | 1.082038 | 192.168.0.2 | 192.168.0.255   | NBNS     | Registration NB.NB10061D.002             |
| 10  | 1.14211  | 192.168.0.2 | 192.168.0.1     | DNS      | Standard query A proxyconf.m004          |
| 11  | 1.14211  | 192.168.0.2 | 192.168.0.1     | TCP      | 1196 > 1196 [SYN] SEQ=0 LEN=0 MSS        |
| 12  | 0.115337 | 192.168.0.1 | 192.168.0.2     | TCP      | 1196 > 3196 [SYN] ACK=0 SEQ=0 ACK=1      |
| 13  | 0.115380 | 192.168.0.2 | 192.168.0.1     | TCP      | 3196 > 1196 [ACK] ACK=1 ACK=1            |
| 14  | 0.115506 | 192.168.0.2 | 192.168.0.1     | TCP      | 3196 > 3196 [PSH, ACK] ACK=1 SEQ=1 ACK=1 |
| 15  | 0.117354 | 192.168.0.1 | 192.168.0.2     | TCP      | 3196 > 3196 [ACK] SEQ=1 ACK=256          |
| 16  | 0.136410 | 192.168.0.1 | 192.168.0.2     | TCP      | 1025 > 5000 [SYN] SEQ=0 LEN=0 MSS        |
| 17  | 0.136410 | 192.168.0.1 | 192.168.0.2     | TCP      | 1025 > 5000 [SYN] SEQ=0 LEN=0 MSS        |

Identification: 0x1847 (6215)

Flags: 0x00

Fragment offset: 0

Time to live: 128

Protocol: UDP (0x11)

Header checksum: 0xa109 [correct]

Source: 192.168.0.2 (192.168.0.2)

Destination: 192.168.0.1 (192.168.0.1)

0000 00 09 5b 2d 75 9a 00 0b 5d 20 cd 02 08 00 45 00 :: [-U... ] ... E.  
0010 00 49 18 47 00 00 80 11 a1 09 cd a8 00 02 d0 a8 :: I.G.::  
0020 00 01 0b d2 00 35 00 35 46 69 00 21 01 00 00 01 :: ... 5.5 Fi. ! ..  
0030 00 00 00 00 00 00 09 70 72 6f 78 79 63 6f 6e 66 :: ... P roxyconf  
0040 05 77 77 30 34 07 73 69 65 6d 65 6e 73 03 6e .m004.s ientens.n