

# Graph Theory

## Trees

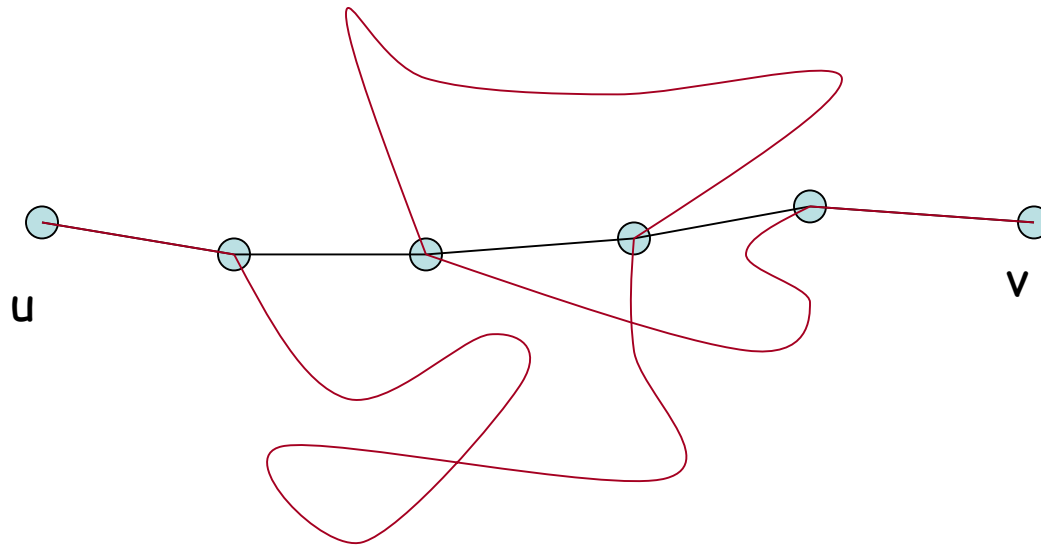
# Tree Characterization by Path

树是没有循环的连通图

**Definition.** A tree is a connected graph with no cycles.

Can there be no path between  $u$  and  $v$ ? NO

Can there be more than one simple path between  $u$  and  $v$ ? NO



This will create cycles.

一条路径上的节点除了  
起点和终点可以相同以外  
其他节点均不相同  
起点和终点相同的简单  
路径称为简单回路。

在树中，每对顶点之间都有一条唯一的简单路径

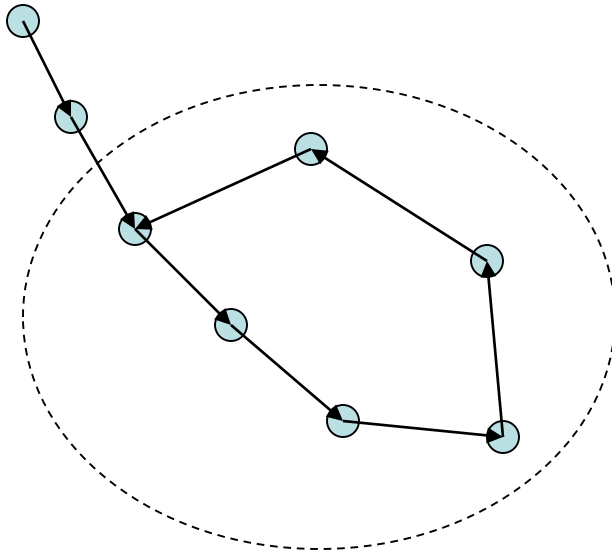
**Claim.** In a tree, there is a unique simple path between every pair of vertices.

# Tree Characterization by Number of Edges

**Definition.** A tree is a connected graph with no cycles.

Can a tree have no leaves? **NO**

Then every vertex has degree at least 2.



Go to unvisited edges as long as possible.

Cannot get stuck,  
unless there is a cycle.

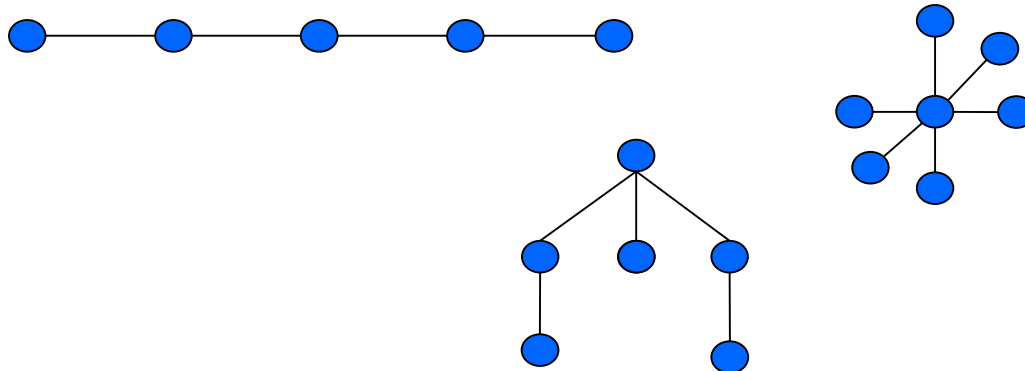
# Tree Characterization by Number of Edges

**Definition.** A tree is a connected graph with no cycles.

Can a tree have no leaves? NO

How many edges does a tree have?

$n-1$  ✓



我们通常使用 $n$ 表示顶点数，并使用 $m$ 表示图中的边数。

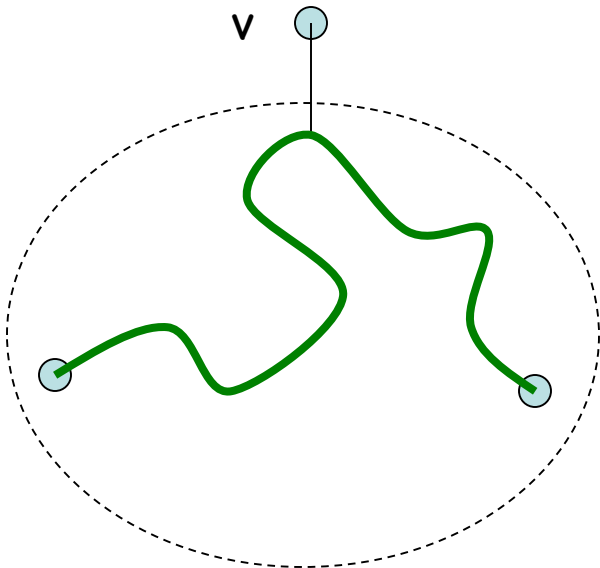
We usually use  $n$  to denote the number of vertices,  
and use  $m$  to denote the number of edges in a graph.

# Tree Characterization by Number of Edges

**Definition.** A tree is a connected graph with no cycles.

Can a tree have no leaves? **NO**

How many edges does a tree have?  $n-1$ ?



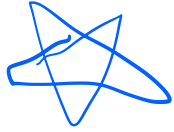
Look at a leaf  $v$ .

Is  $T-v$  a tree? **YES**

1. Can  $T-v$  has a cycle? **NO**
2. Is  $T-v$  connected? **YES**

By induction,  $T-v$  has  $(n-1)-1=n-2$  edges.

So  $T$  has  $n-1$  edges.



# Tree Characterizations

**Definition.** A tree is a connected graph with no cycles.

## Characterization by paths:

A graph is a tree if and only if

①任意两点之间<sup>且</sup>仅有一条 simple path  
则其一定是树

there is a unique simple path between every pair of vertices.

## Characterization by number of edges:

② A graph is a tree if and only if it is connected and has  $n-1$  edges.

(We have only proved one direction.

The other direction is similar and left as an exercise.)

① 树是 没有循环的连通图、无向图

# Trees

- **Definition:** A **tree** is a **connected undirected** graph with **no simple circuits**.

- Since a tree cannot have a simple circuit, a tree cannot contain multiple edges or loops.

- Therefore, any tree must be a **simple graph**.

任何树都是简单图

一棵树不能包含多个边或循环

无向的  
没有自循环  
没有多重边

- **Theorem:** An undirected graph is a tree if and only if there is a **unique simple path** between any of its vertices.

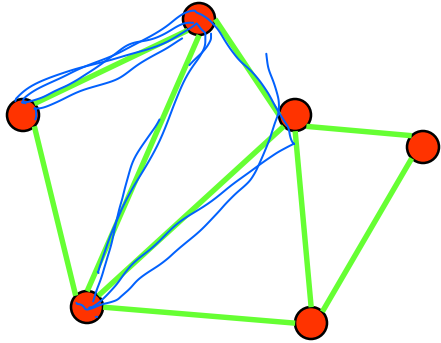
- **Definition:** An undirected graph that does not contain simple circuits and is not necessarily connected is called a **forest**.

- In general, we use trees to represent **hierarchical structures**.

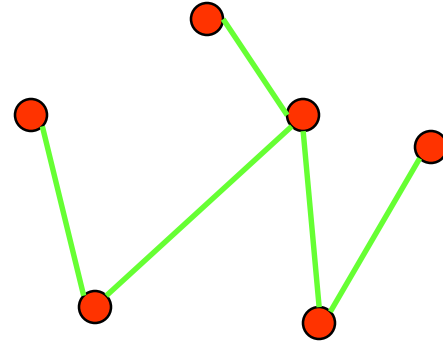
通常，我们使用树来表示层次结构

# Trees

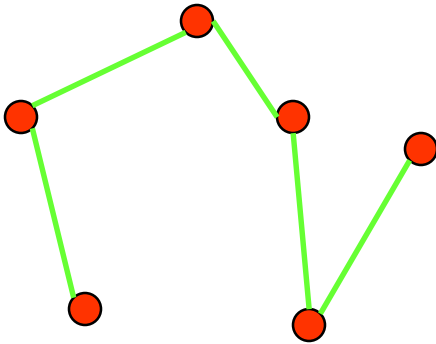
•Example: Are the following graphs trees?



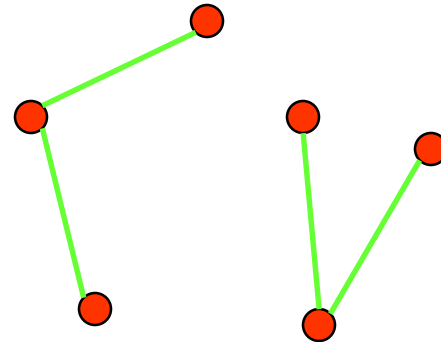
No.



Yes.



Yes.



No.



# Spanning Trees

令 $G$ 为简单图形。 $G$ 的生成树是 $G$ 的子图， $G$ 是包含 $G$ 的每个顶点的树。

• **Definition:** Let  $G$  be a simple graph. A spanning tree of  $G$  is a subgraph of  $G$  that is a tree containing every vertex of  $G$ .

$G = (V, E)$  的生成树是 $V$ 上具有最小边数的连接图

• **Note:** A spanning tree of  $G = (V, E)$  is a connected graph on  $V$  with a minimum number of edges  $(|V| - 1)$ .

• **Example:** Since winters in Boston can be very cold, six universities in the Boston area decide to build a tunnel system that connects their libraries.

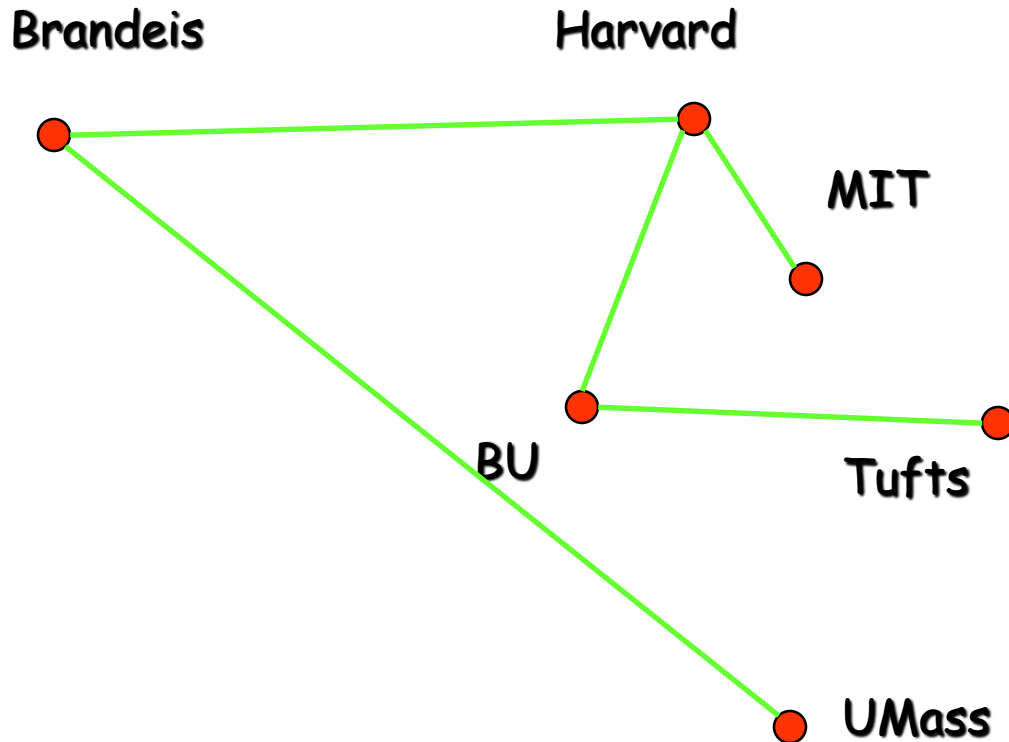
无向图的点不动，去掉边，直到边的个数为 $n-1$ 为止

## Spanning tree

- A spanning tree in an undirected graph  $G(V,E)$  is a subset of edges  $T \subseteq E$  that are acyclic and connect all the vertices in  $V$ .  
无向图
- A spanning tree must consist of exactly  $n-1$  edges.
- Suppose that each edge has a weight associated with it. Say that the weight of a tree  $T$  is the sum of the weights of its edges  $w(T) = \sum_{e \in T} w(e)$   
权重
- The minimum spanning tree in a weighted graph  $G(V,E)$  is one which has the smallest weight among all spanning trees in  $G(V,E)$   
最小生成树

# Spanning Trees

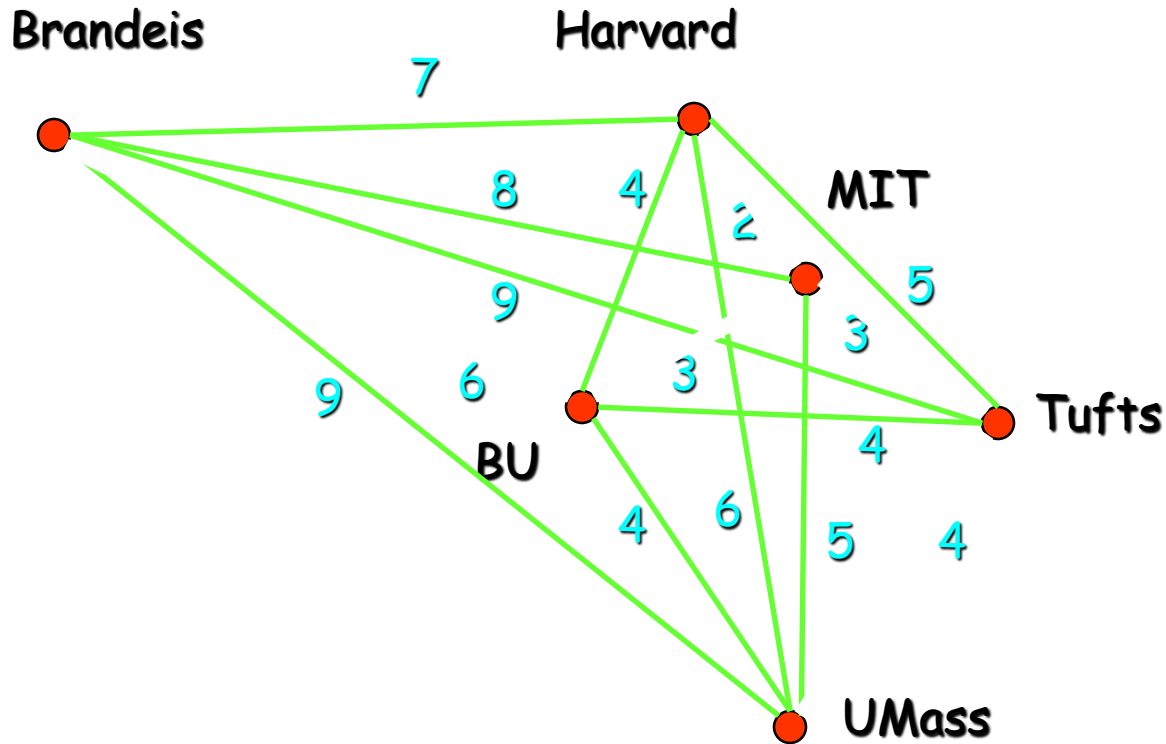
- Example for a spanning tree:



Since there are 6 libraries, 5 tunnels are sufficient to connect all of them.

# Spanning Trees

- The complete graph with cost labels (in billion \$):



The least expensive tunnel system costs \$20 billion.

# Spanning Trees

- Now imagine that you are in charge of the tunnel project. How can you determine a tunnel system of **minimal cost** that connects all libraries?
- **Definition:** A **minimum spanning tree** in a connected weighted graph is a spanning tree that has the smallest possible sum of weights of its edges.
- How can we find a minimum spanning tree?

# Spanning Trees

① 先把所有边排序

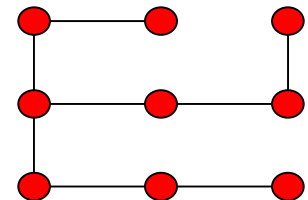
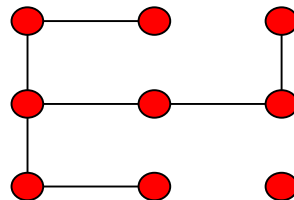
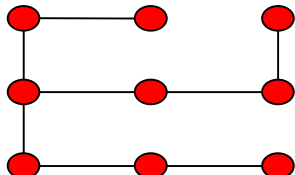
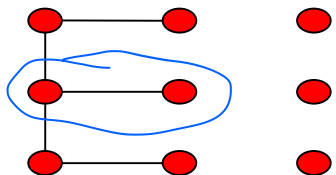
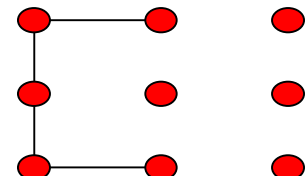
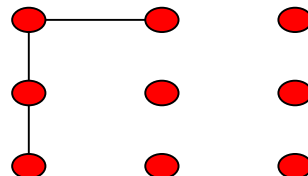
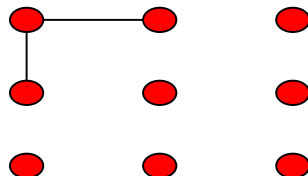
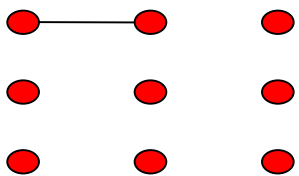
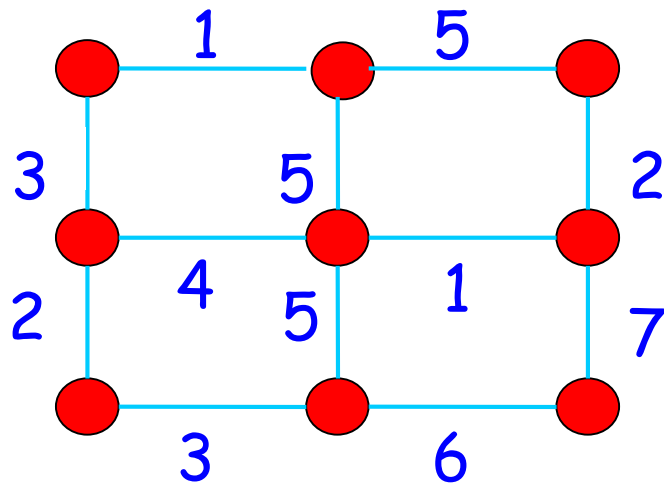
② 选最小的边

③ 选与这条边相邻的权重小的边

## • Prim's Algorithm:

- Begin by choosing any edge with **smallest weight** and putting it into the spanning tree,
- successively add to the tree edges of **minimum weight** that are incident to a vertex already in the tree and not forming a simple circuit with those edges already in the tree,
- stop when  $(n - 1)$  edges have been added.

# Prim's algorithm



# Spanning Trees

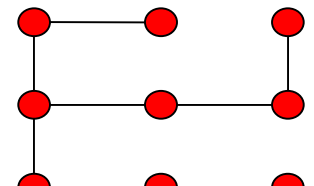
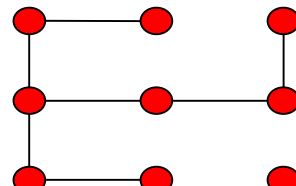
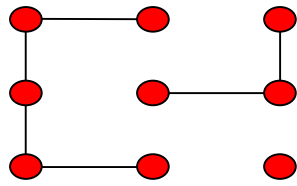
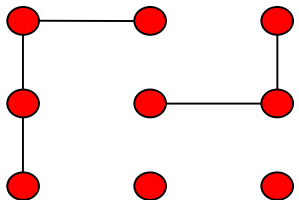
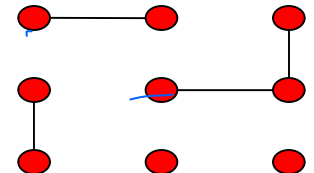
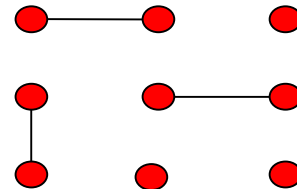
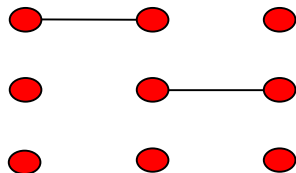
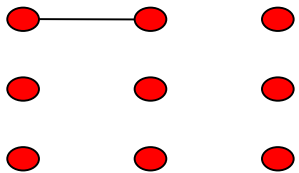
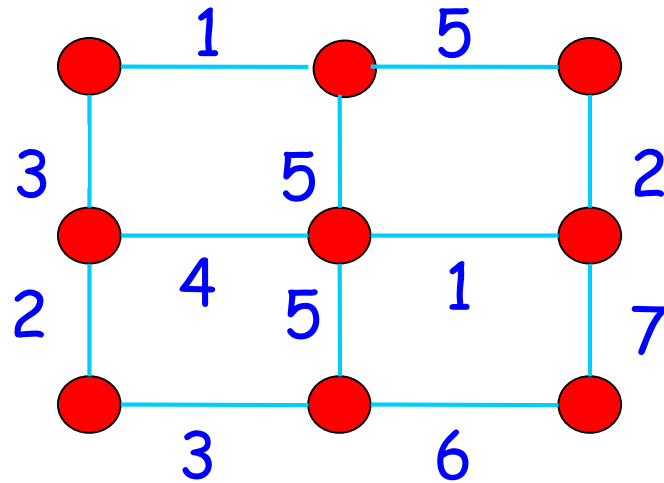
## •Kruskal's Algorithm:

- Kruskal's algorithm is identical to Prim's algorithm, except that it does not demand new edges to be incident to a vertex already in the tree.
- Both algorithms are **guaranteed** to produce a minimum spanning tree of a connected weighted graph.

先把各边权重从小到大排列。然后从权重最小的边开始加



# Kruskal's algorithm



# Rooted Trees 根树 (有向树)

- We often designate a particular vertex of a tree as the **root**. Since there is a unique path from the root to each vertex of the graph, we direct each edge away from the root.
- Thus, a tree together with its root produces a **directed graph** called a **rooted tree**.

根: 入度为0

其他入度均为1

叶: 出度为0

分枝点/内点 = 出度不为0

# Rooted Trees

- If  $v$  is a vertex in a rooted tree other than the root, the **parent** of  $v$  is the unique vertex  $u$  such that there is a directed edge from  $u$  to  $v$ .
- When  $u$  is the parent of  $v$ ,  $v$  is called the **child** of  $u$ .
- Vertices with the same parent are called **siblings**.
- The **ancestors** of a vertex other than the root are the vertices in the path from the root to this vertex, excluding the vertex itself and including the root.

# Rooted Trees

- The **descendants** of a vertex  $v$  are those vertices that have  $v$  as an ancestor.
- A vertex of a tree is called a **leaf** if it has no children.
- Vertices that have children are called **internal vertices**.
- If  $a$  is a vertex in a tree, then the **subtree** with  $a$  as its root is the subgraph of the tree consisting of  $a$  and its descendants and all edges incident to these descendants.

如果 $a$ 是树中的顶点，则以 $a$ 为根的子树是树的子图，该子图由 $a$ 及其后代以及入射到这些后代的所有边组成。

# Rooted Trees

**level**: 根树中顶点 $v$ 的级别是从根到此顶点的唯一路径的长度。

- The **level** of a vertex  $v$  in a rooted tree is the length of the unique path from the root to this vertex.

- The level of the root is defined to be zero.

高是最大的 level

- The **height** of a rooted tree is the maximum of the levels of vertices.

节点深度: 从根节点到该节点最长简单路径边的条数

高度: 从某点到叶子节点的最长简单路径边的条数

深度为  $n$ , 最多有  $2^n - 1$  个节点 (根节点深度按 1 算)

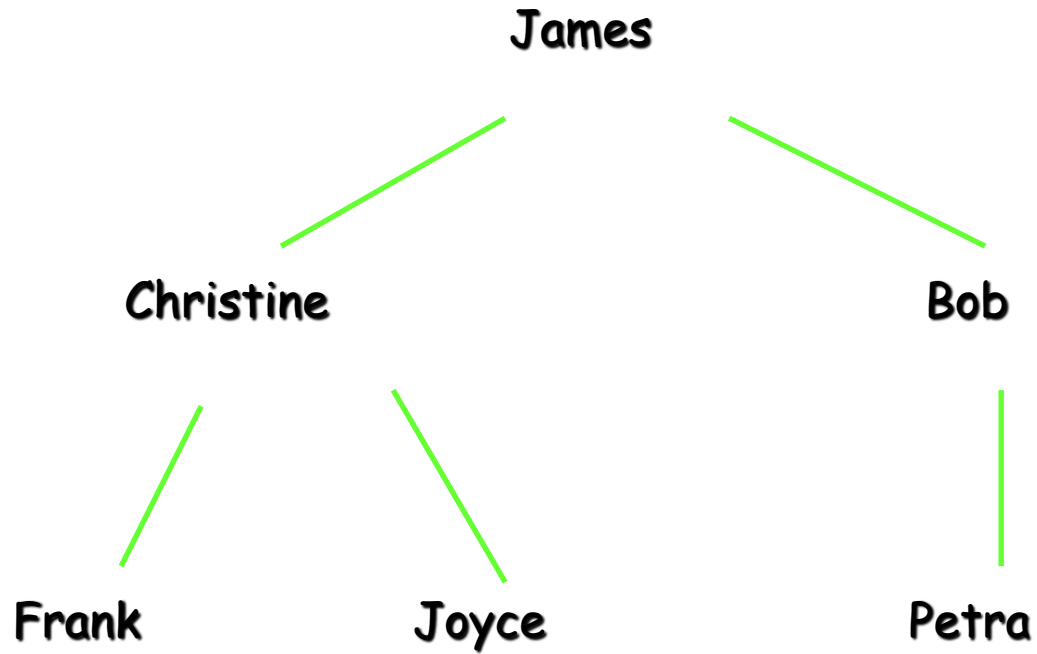
在第  $i$  层, 最多有  $2^{i-1}$  个节点

具有  $n$  个结点的二叉树最小深度为  $\lg_2(n+1)$

最大深度为  $n-1$

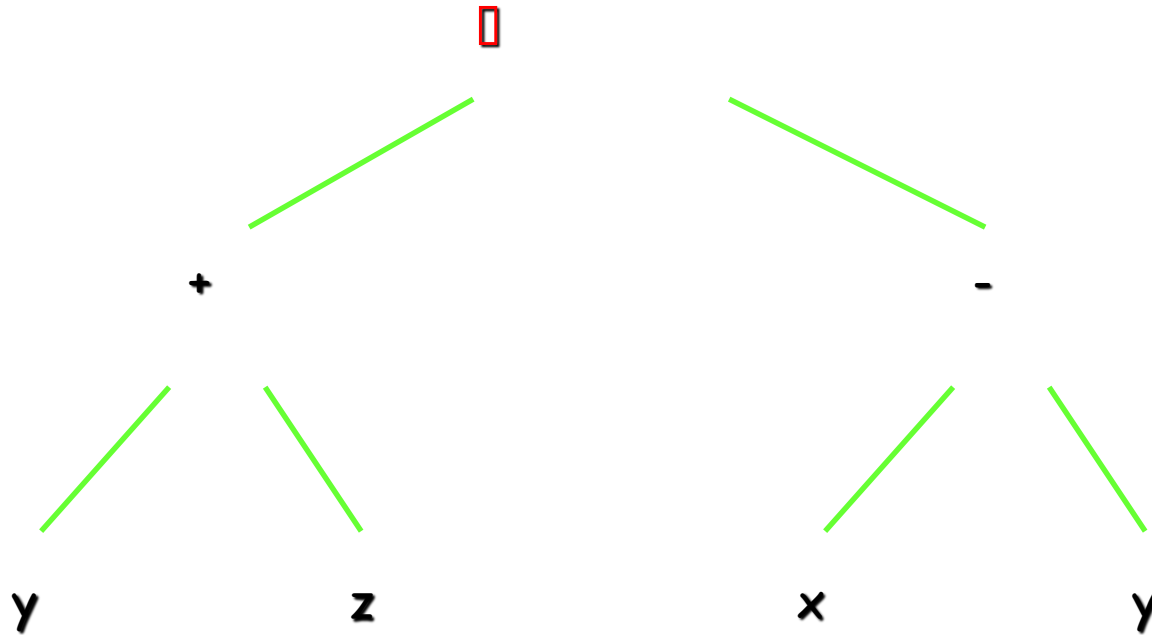
# Trees

- **Example I:** Family tree



# Trees

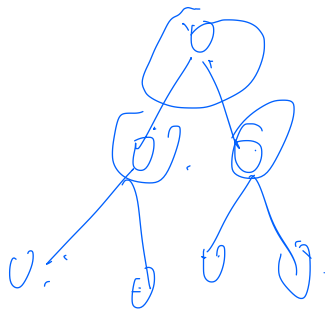
- **Example III:** Arithmetic expressions



This tree represents the expression  $(y + z) * (x - y)$ .

# Trees m 叉树

- **Definition:** A rooted tree is called an **m-ary tree** if every internal vertex has no more than m children.
- The tree is called a **full m-ary tree** if every internal vertex has exactly m children. 每个点出度均为 m
- An m-ary tree with  $m = 2$  is called a **binary tree**.
- **Theorem:** A tree with n vertices has  $(n - 1)$  edges.
- **Theorem:** A full m-ary tree with i internal vertices contains  $n = mi + 1$  vertices.



$$3 \times 2 + 1$$

2

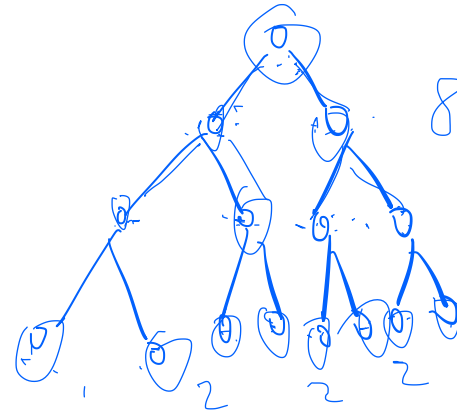
4 +

13

6

13

$$6 \times 2 + 1 = 13$$



8 +

$$7 \times 3 + 1 = 22$$

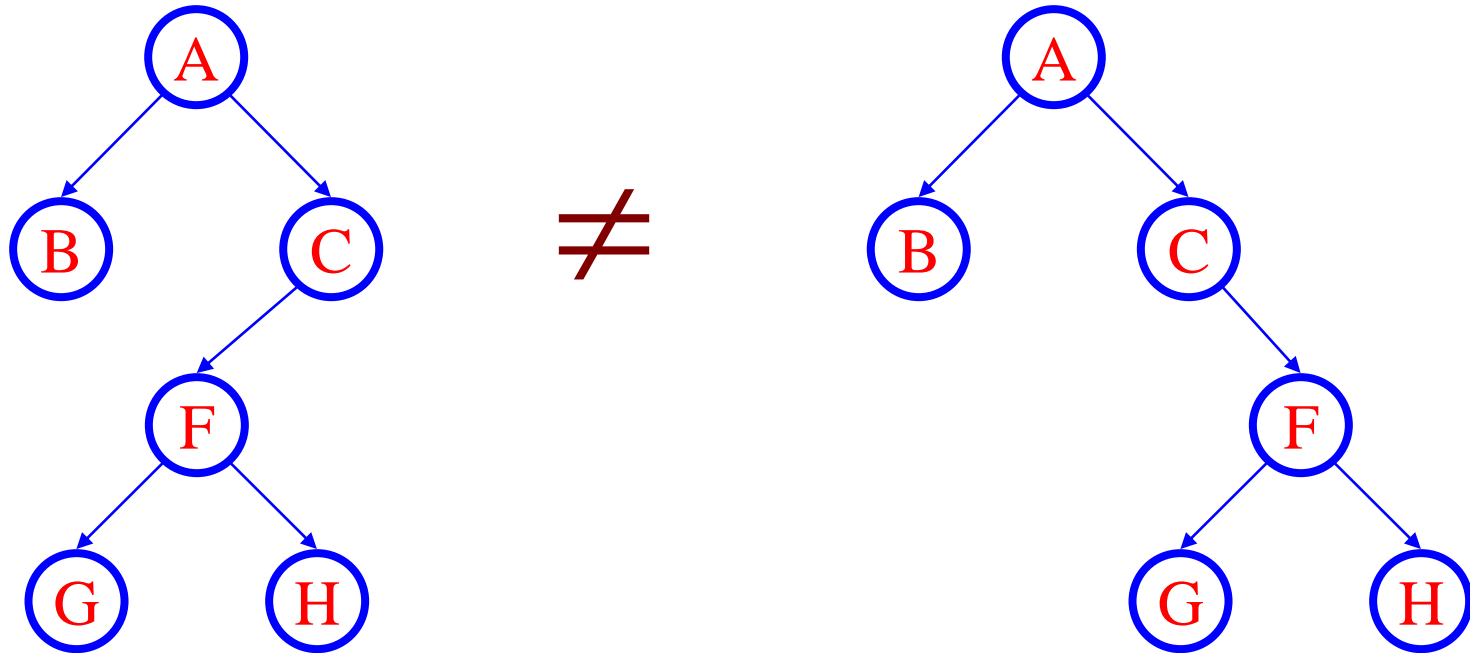


# Binary Trees

- Every node has at most two children 每个节点最多有2个孩子。
- Most popular tree in computer science
- Given  $N$  nodes, what is the minimum depth of a binary tree?  $\log_2 (N+1)$
- What is the <sup>最大的深度</sup> maximum depth of a binary tree with  $N$  nodes? <sup>二叉树有  $n$  个节点</sup>  $N-1$ .

# Binary Trees

- Notice:
- we distinguish between left child and right child



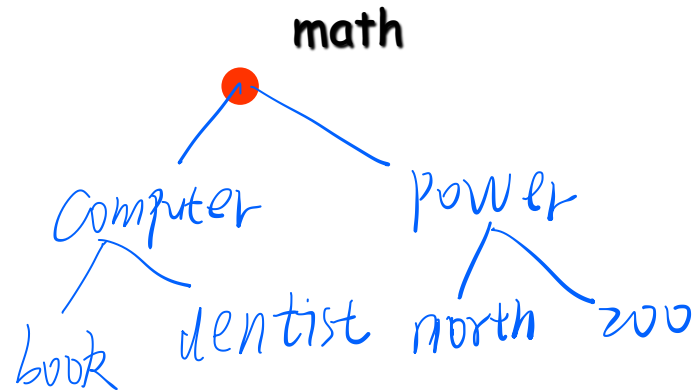
# Binary Search Trees = 二叉搜索树

- If we want to perform a large number of searches in a particular list of items, it can be worthwhile to arrange these items in a **binary search tree** to facilitate the subsequent searches.
- A binary search tree is a binary tree in which each child of a vertex is designated as a **right or left child**, and each vertex is labeled with a **key**, which is one of the items.
- When we construct the tree, vertices are assigned keys so that the key of a vertex is both larger than the keys of all vertices in its left subtree and smaller than the keys of all vertices in its right subtree.

小放左  
大放右

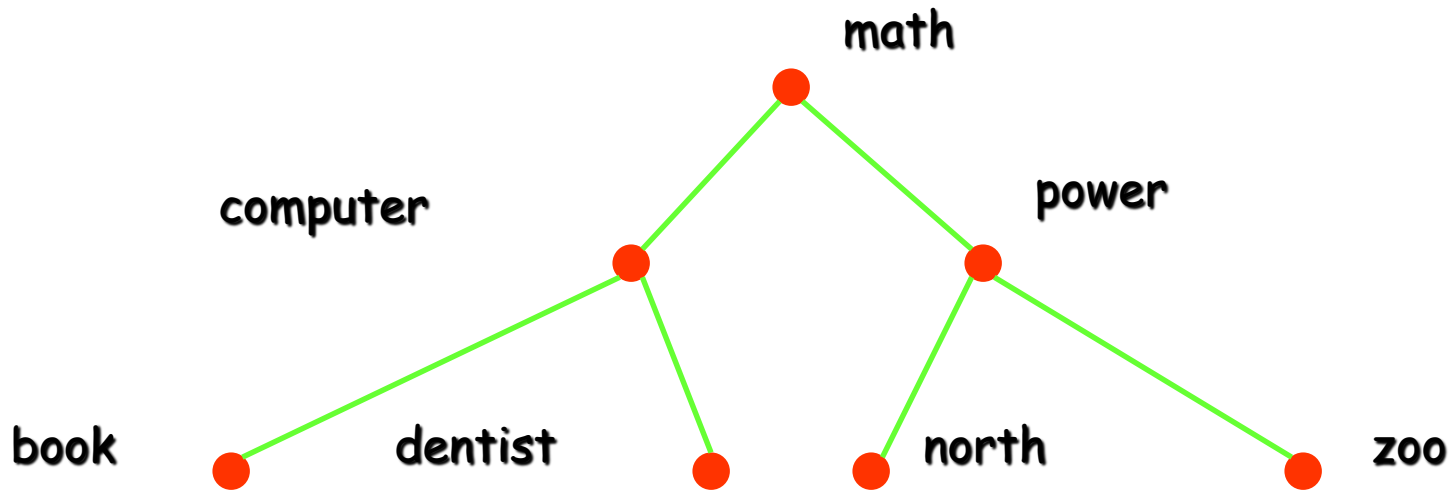
# Binary Search Trees

•**Example:** Construct a binary search tree for the strings **math**, **computer**, **power**, **north**, **zoo**, **dentist**, **book**.



# Binary Search Trees

•**Example:** Construct a binary search tree for the strings **math**, **computer**, **power**, **north**, **zoo**, **dentist**, **book**.



# Binary Search Trees

- To perform a search in such a tree for an item  $x$ , we can start at the root and compare its key to  $x$ . If  $x$  is **less** than the key, we proceed to the **left** child of the current vertex, and if  $x$  is **greater** than the key, we proceed to the **right** one.  
小放左  
大放右
- This procedure is repeated until we either found the item we were looking for, or we cannot proceed any further.

The End