

## 0.1 Exact functors

Let  $F : \mathcal{A} \rightarrow \mathcal{A}'$  be an additive functor.

**Definition 0.1.**  $F$  is called

- (a) *left exact*  $\iff F$  commutes with finite limits.
- (b) *right exact*  $\iff F$  commutes with finite colimits.
- (c) *exact*  $\iff F$  is left and right exact.

**Remark.** Since  $\mathcal{A}, \mathcal{A}'$  are additive categories, all finite limits and colimits exist in  $\mathcal{A}$  and  $\mathcal{A}'$ . So if  $D : J \rightarrow \mathcal{A}$  is a finite diagram, we have  $\lim_J D$  exists in  $\mathcal{A}$ ,  $\lim_J F \circ D$  exists in  $\mathcal{A}'$  and we have a natural morphism

$$F(\lim_J D) \rightarrow \lim_J F \circ D$$

in  $\mathcal{A}'$ .  $F$  is left exact if this morphism is an isomorphism  $\forall D$ .

**Proposition 0.2.** Let  $F : \mathcal{A} \rightarrow \mathcal{A}'$  be additive, then:

(a) The following are equivalent:

- (i)  $F$  is left exact.
- (ii)  $F$  commutes with the formation of kernels, i.e.  $\forall f : X \rightarrow Y$  in  $\mathcal{A}$ , the natural morphism
$$F(\ker f) \rightarrow \ker F(f)$$
is an isomorphism.
- (iii)  $\forall$  exact sequence  $0 \rightarrow X' \rightarrow X \rightarrow X''$  in  $\mathcal{A}$ , the sequence  $0 \rightarrow FX' \rightarrow FX \rightarrow FX''$  is exact in  $\mathcal{A}'$ .
- (iv)  $\forall$  exact sequence  $0 \rightarrow X' \rightarrow X \rightarrow X'' \rightarrow 0$  in  $\mathcal{A}$ , the sequence  $0 \rightarrow FX' \rightarrow FX \rightarrow FX''$  is exact in  $\mathcal{A}'$ .

(b) The following are equivalent:

- (i)  $F$  is right exact.
- (ii)  $F$  commutes with the formation of cokernels.
- (iii)  $\forall$  exact sequence  $X' \rightarrow X \rightarrow X'' \rightarrow 0$  in  $\mathcal{A}$ , the sequence  $FX' \rightarrow FX \rightarrow FX'' \rightarrow 0$  is exact in  $\mathcal{A}'$ .
- (iv)  $\forall$  exact sequence  $0 \rightarrow X' \rightarrow X \rightarrow X'' \rightarrow 0$  in  $\mathcal{A}$ , the sequence  $FX' \rightarrow FX \rightarrow FX'' \rightarrow 0$  is exact in  $\mathcal{A}'$ .

(c)  $F$  is exact  $\iff \forall$  exact sequences  $X' \xrightarrow{f} X \xrightarrow{g} X''$ , the sequence

$$FX' \xrightarrow{Ff} FX \xrightarrow{Fg} FX''$$

is exact.

*Proof.* TODO. □

**Proposition 0.3.**  $\forall X \in \mathcal{A}$  the co- and contravariant Hom functors  $\text{Hom}_{\mathcal{A}}(X, -), \text{Hom}_{\mathcal{A}}(-, X)$  are left exact.

*Proof.* TODO. □

**Proposition 0.4.** *Let  $F : \mathcal{A} \rightarrow \mathcal{A}'$  and  $G : \mathcal{A}' \rightarrow \mathcal{A}$  be additive functors with  $F \dashv G$ . Then  $F$  is right exact and  $G$  is left exact.*

*Proof.* TODO. □

**Example 0.5.** By Hom- $\otimes$  Adjunction (roughly  $\text{Hom}(M \otimes -, N) = \text{Hom}(M, \text{Hom}(-, N))$ )  $M \otimes -$  is left adjoint to  $\text{Hom}(-, ) \implies M \otimes_R - : {}_R\mathbf{Mod} \rightarrow \mathbf{Ab}$  is right exact, as is  $- \otimes_R M : \mathbf{Mod}_R \rightarrow \mathbf{Ab}$ .

$$(- \otimes_R M : \mathbf{Mod}_R \rightarrow \mathbf{Ab}) = (M \otimes_{R^{\text{op}}} - : {}_{R^{\text{op}}}\mathbf{Mod} \rightarrow \mathbf{Ab})$$

**Remark.** Next small goal:  $J$  any small index category. Are  $\lim_J : \mathcal{A}^J \rightarrow \mathcal{A}$ ,  $\text{colim}_J : \mathcal{A}^J \rightarrow \mathcal{A}$  exact?

**Proposition 0.6.**  $\mathcal{A}^J = \text{Fun}(J, \mathcal{A})$  is an abelian category.

**Proposition 0.7.** Suppose  $\mathcal{A}$  contains all limits (colimits) for a given small index category  $J$ , then  $\lim_J : \mathcal{A}^J \rightarrow \mathcal{A}$  is left exact. ( $\text{colim}_J : \mathcal{A}^J \rightarrow \mathcal{A}$  is right exact.)

**Corollary 0.8.** Let  $I$  be a set and  $\underline{I}$  the discrete category associated to  $I$ . Then:

- (a)  $\prod_{i \in I} : A^I \rightarrow A, (A_i)_{i \in I} \mapsto \prod_{i \in I} A_i$
- (b)  $\bigoplus_{i \in I} : A^I \rightarrow A, (A_i)_{i \in I} \mapsto \bigoplus_{i \in I} A_i$  assuming existence, are exact functors.

**Definition 0.9.** (a) A non-empty category  $J$  is called *filtered* iff

$$(i) \forall i, j \in J : \exists \text{ diagram } \begin{array}{ccc} i & & \\ & \searrow & \\ j & \longrightarrow & k \end{array} \in J.$$

(ii)  $\forall f, g : i \rightrightarrows j \exists h : j \rightarrow k$  such that

$$h \circ f = h \circ g : i \rightarrow k$$

- (b) A directed poset  $(I, \subseteq)$  is a poset such that  $\forall i, j \in I \exists h \in I$  such that  $h \geq i, h \geq j$  ( $\implies (I, \subseteq)$  is directed  $\implies \text{ord}(I, \subseteq)$  is filtered.)

**Definition 0.10.**  $\mathcal{A}$  has exact filtered colimits  $\iff$  for each filtered index category  $J$ , the functor  $\text{colim } \mathcal{A}^J \rightarrow \mathcal{A}$  is exact (and defined).

**Theorem 0.11.**  ${}_R\mathbf{Mod}$  has exact filtered colimits.

Fundamental question: how to investigate the non-exactness of functors (that are left or right exact). Answer of homological algebra: e.g.  $F : A \rightarrow B$  left exact, define “higher right derived functors”  $(R^i F)_{i \geq 0}$  from  $F$  such that  $\forall$  s.e.s.

$$0 \rightarrow X' \rightarrow X \rightarrow X'' \rightarrow 0$$

in  $\mathcal{A}$

$$\begin{array}{ccccccc}
0 & \longrightarrow & FX' & \longrightarrow & FX & \longrightarrow & FX'' \\
& & & & & \searrow & \\
& & R^1FX' & \longrightarrow & R^1FX & \longrightarrow & R^1FX'' \\
& & & & & \searrow & \\
& & R^2FX' & \longrightarrow & R^2FX & \longrightarrow & R^2FX'' \longrightarrow \dots
\end{array}$$

and  $R^0F = F$ . Study  $R^iF$  to understand the nonexactness of  $F$ , or to gain insight into some invariants of  $\mathcal{A}$ . Some  $R^iF$  (typically  $i \leq 3$ ) have concrete meanings.

### 0.1.1 To define $R^iF$ (or $L_iF$ )

One wants “enough” injectives (projectives) in  $\mathcal{A}$ .

**Theorem-Definition 0.12.** For  $I \in \mathcal{A}$  the following are equivalent:

- (i)  $\forall$  diagrams with  $\iota$  monomorphism  $\exists$  extension  $g : B \rightarrow I$  such that the following commutes

$$\begin{array}{ccccc}
0 & \longrightarrow & A & \xrightarrow{\iota} & B \\
& & \downarrow f & \swarrow g & \\
& & I & & 
\end{array}$$

- (ii) The functor  $\text{Hom}_{\mathcal{A}}(-, I) : \mathcal{A}^{\text{op}} \rightarrow \mathbf{Ab}$  is exact.

- (iii) Every s.e.s.  $0 \rightarrow I \xrightarrow{h} C \xrightarrow{k} D \rightarrow 0$  in  $\mathcal{A}$  is split.

If any of these hold then  $I$  is called an injective object.

**Theorem-Definition 0.13.** For  $P \in \mathcal{A}$  the following are equivalent:

- (i)  $\forall$  diagrams with  $\pi$  epimorphism  $\exists$  lifting  $g : P \rightarrow B$  such that  $\pi \circ g = f$

$$\begin{array}{ccccc}
& & P & & \\
& \swarrow g & \downarrow f & & \\
B & \xrightarrow{\pi} & C & \longrightarrow & 0
\end{array}$$

- (ii) The functor  $\text{Hom}_{\mathcal{A}}(P, -) : \mathcal{A} \rightarrow \mathbf{Ab}$  is exact.

- (iii) Every s.e.s.  $0 \rightarrow A \rightarrow B \rightarrow P \rightarrow 0$  is split.

If these hold, then  $P$  is called a projective object in  $\mathcal{A}$ .

**Remark.**  $I$  injective in  $\mathcal{A} \iff I$  projective in  $\mathcal{A}^{\text{op}}$ .

**Proposition 0.14.** (a) If  $(I_j)_{j \in J}$  ( $J$  a set) is a family of injectives in  $\mathcal{A}$ , such that  $\prod_j I_j$  exist, then  $\prod_j I_j$  is injective.

(b) If  $(P_j)_{j \in J}$  ( $J$  a set) is a family of projectives in  $\mathcal{A}$ , such that  $\bigoplus_{j \in J} P_j$  exist, then  $\bigoplus_j P_j$  is projective.

**Example 0.15.** (a)  $R$  is a projective  $R$ -module (hence so is  $R^{(I)}$ )

(b)  $\mathbb{Q}/\mathbb{Z}$  is an injective in  ${}_{\mathbb{Z}}\mathbf{Mod}$ .

**Definition 0.16.** (a)  $\mathcal{A}$  has enough injective  $\iff \forall X \in \mathcal{A} \exists$  monomorphism  $X \rightarrow I$  with  $I \in \mathcal{A}$  injective.

(b)  $\mathcal{A}$  has enough projectives  $\iff \forall X \in \mathcal{A} \exists$  epimorphism  $P \rightarrow X$  with  $P \in \mathcal{A}$  projective.

**Definition 0.17.**  $Q \in \mathcal{A}$  is called a

(a) *generator*  $\iff \text{Hom}_{\mathcal{A}}(Q, -) : \mathcal{A} \rightarrow \mathbf{Ab}$  is faithful.

(b) *cogenerator*  $\iff \text{Hom}_{\mathcal{A}}(-, Q) : \mathcal{A}^{\text{op}} \rightarrow \mathbf{Ab}$  is faithful.

**Remark 0.18.** For  $Q \in \mathcal{A}$  the following are equivalent:

1.  $Q$  is a generator.
2.  $\forall X, Y \in \mathcal{A} \forall f, g \in \mathcal{A}(X, Y) :$

$$f \neq g \implies \exists h \in \mathcal{A}(Q, X) : f \circ h \neq g \circ h$$

3.  $\forall X, Y \in \mathcal{A} \forall f \in \mathcal{A}(X, Y) :$

$$f \neq 0 \implies \exists h \in \mathcal{A}(Q, X) : f \circ h \neq 0$$

**Examples 0.19.** 1.  $R$  is a generator for  ${}_R\mathbf{Mod}$ : Suppose  $f : M \rightarrow N$  in  ${}_R\mathbf{Mod}$  is nonzero. Then  $\exists m \in M : f(m) \neq 0$ . Define  $h : R \rightarrow M, r \mapsto rm \implies f \circ h : R \rightarrow N$  is nonzero.

2.  $\mathbb{Q}/\mathbb{Z}$  is a cogenerator in  $\mathbf{Ab} = {}_{\mathbb{Z}}\mathbf{Mod}$  let  $f : M \rightarrow N$  be non-zero in  $\mathbf{Ab}$ . Let  $X \in \text{im}(f) \setminus \{0\}$ . TODO

**Definition 0.20.**  $\mathcal{A}$  is called a *Grothendieck abelian category* (GAC) iff

1.  $\mathcal{A}$  is cocomplete.
2. For any filtered small category  $J : \text{colim} : A^J \rightarrow \mathcal{A}$  is exact.
3.  $\mathcal{A}$  possesses a cogenerator.

**Theorem 0.21.** For a GAC  $\mathcal{A}$  the following hold:

1. The subobjects and quotient objects of any  $A \in \mathcal{A}$  form a set.
2.  $\mathcal{A}$  has enough injectives.
3.  $\mathcal{A}$  has an injective cogenerator.
4.  $\mathcal{A}$  is complete.

**Example 0.22.** 1.  ${}_R\mathbf{Mod}$  is a GAC by examples 19(a) Thm 11 Cor II48

2. If  $\mathcal{A}$  is a GAC and  $J$  is a small category then  $\mathbf{PSh}(J, \mathcal{A})$  is a GAC

**Lemma 0.23.** Suppose  $F : \mathcal{A} \rightarrow \mathcal{A}'$  and  $G : \mathcal{A}' \rightarrow \mathcal{A}$  are functors such that  $F \dashv G$ , then:

1.  $F$  exact  $\implies G$  maps injectives in  $\mathcal{A}'$  to injectives in  $\mathcal{A}$ .
2.  $G$  exact  $\implies F$  maps projectives in  $\mathcal{A}$  to projectives in  $\mathcal{A}'$ .
3.  $F$  faithful  $\implies G$  maps cogenerators in  $\mathcal{A}'$  to cogenerators in  $\mathcal{A}$ .
4.  $G$  faithful  $\implies F$  maps generators in  $\mathcal{A}$  to generators in  $\mathcal{A}'$ .

**Theorem 0.24.** Let  $\mathcal{A}$  be a GAC with generator  $Q$ , then:

**Corollary 0.25.**  $\mathbb{Q}/\mathbb{Z}$  is injective in  $\mathbf{Ab}$ .

**Corollary 0.26.** For any ring  $R$ ,  $\text{Hom}(R, \mathbb{Q}/\mathbb{Z})$  is an injective  $R$ -module and a cogenerator.

**Proposition 0.27.** Let  $Q, J \in \mathcal{A}$ , then

1. Suppose  $\mathcal{A}$  contains all coproducts over index sets. Then the following are equivalent:
  - (a)  $Q$  is a generator.
  - (b)  $\forall X \in \mathcal{A} \exists$  set  $I, \exists$  epimorphism  $Q^{(I)} := \coprod_{i \in I} Q \rightarrow X$ . Moreover, if  $Q$  is a rojective generator then  $\mathcal{A}$  has enough projectievs of the form  $Q^{(I)}$ , where  $I$  is a set.
2. Suppose  $\mathcal{A}$  contains all products over index sets then the following are equivalent:
  - (a)  $J$  is a cogenerator.
  - (b)  $\forall Y \in \mathcal{A} \exists$  set  $I$ , monomorphism  $Y \rightarrow J^I := \prod_{i \in I} J$ . Moreover if  $J$  is an injective cogenerrator then  $\mathcal{A}$  has enough injectives of the form  $J^I$  where  $I$  is a set.

**Corollary 0.28.**  ${}_R\mathbf{Mod}$  has enough injectives of the form  $\text{Hom}_{\mathbb{Z}}(R, \mathbb{Q}/\mathbb{Z})^I$ ,  $I$  a set.

**Remark 0.29.** TODO

## 0.2 (Co-)chain complexes, (co-)homology

Let  $\mathcal{A}$  be an additive category.

**Definition 0.30.** (a) A chain complex  $(C_*, \partial_*)$  over  $\mathcal{A}$  is a sequence  $(\partial_i : C_i \rightarrow C_{i-1})_{i \in \mathbb{Z}}$  of morphisms in  $\mathcal{A}$  such that  $\partial_i \circ \partial_{i+1} = 0, \forall i \in \mathbb{Z}$ . The map  $\partial_i$  is called the  $i$ -th differential or  $i$ -th boundary map of the complex. A morphism  $f_* : C_* \rightarrow D_*$  of chain complexes is a sequence of morphisms  $f_* = (f_i : C_i \rightarrow D_i)_{i \in \mathbb{Z}}$  such that the diagram commutes  $\forall i \in \mathbb{Z}$

$$\begin{array}{ccc} C_{i-1} & \xleftarrow{\partial_i^C} & C_i \\ f_{i-1} \downarrow & & \downarrow f_i \\ D_{i-1} & \xleftarrow{\partial_i^D} & D_i \end{array}$$

The category of chain complexes over  $\mathcal{A}$  is denoted  $\text{Ch}_*(\mathcal{A})$ .

- (b) A *cochain complex*  $(C^*, \delta^*)$  over  $\mathcal{A}$  is a sequence  $(\delta^i : C^i \rightarrow C^{i+1})_{i \in \mathbb{Z}}$  of morphisms in  $\mathcal{A}$ , such that  $\delta^{i+1} \circ \delta^i = 0, \forall i \in \mathbb{Z}$ . A morphism  $f^* : C^* \rightarrow D^*$  of cochain complexes is a sequence  $f^* = (f^i : C^i \rightarrow D^i)_{i \in \mathbb{Z}}$  of morphisms in  $\mathcal{A}$  such that the following diagram commutes  $\forall i \in \mathbb{Z}$

$$\begin{array}{ccc} D_i & \xrightarrow{\delta_D^i} & D_{i+1} \\ f_i \uparrow & & \uparrow f_{i+1} \\ C_i & \xrightarrow{\delta_C^i} & C_{i+1} \end{array}$$

The *category of cochain complexes* over  $\mathcal{A}$  is denoted  $\text{Ch}^*(\mathcal{A})$

**Exercise 0.31.** TODO

In the following we do most things only for cochain complexes.

- Definition 0.32.** (a) The *support* of a cochain complex  $C^* \in \text{Ch}^*(\mathcal{A})$  is  $\text{supp } C = \{i \in \mathbb{Z} \mid C^i \neq 0\}$
- (b) The full subcategory of  $\text{Ch}^*(\mathcal{A})$  on complexes supported on  $\mathbb{N}_0$  (or on  $-\mathbb{N}_0$ ) is denoted  $\text{Ch}_{\geq 0}^*(\mathcal{A})$  (or  $\text{Ch}_{\leq 0}^*(\mathcal{A})$ ).

**Proposition 0.33.** If  $\mathcal{A}$  is additive (or abelian), then so are  $\text{Ch}^*(\mathcal{A}), \text{Ch}_{\geq 0}^*(\mathcal{A}), \text{Ch}_{\leq 0}^*(\mathcal{A})$ .

**Definition 0.34.** For  $i \in \mathbb{Z}$  define left shift by  $i$  by

$$\begin{array}{ccc} \text{Ch}^*(\mathcal{A}) & \longrightarrow & \text{Ch}^*(\mathcal{A}) \\ C & \longmapsto & C[i] \\ f \downarrow & & \downarrow f[i] \\ D & \longmapsto & D[i] \end{array}$$

where  $C[i]^n = C^{n+i}, \delta_{C[i]}^n = \delta_C^{n+i}$  and  $f[i]^n = f^{n+1}$ .

**Convention 0.35.** We regard  $\mathcal{A}$  as a subcategory of  $\text{Ch}^*(\mathcal{A})$ , as the subcategory of complexes  $C^*$  with  $\text{supp } C^* \subseteq \{0\}$ . Identify  $X \in \mathcal{A}$  with complex

$$\cdots \rightarrow 0 \rightarrow 0 \rightarrow X \rightarrow 0 \rightarrow 0 \rightarrow \cdots$$

**Example 0.36.** TODO

### 0.3 Double complexes

Can iterate the formation of  $\mathcal{A} \rightarrow \text{Ch}^*(\mathcal{A})$ .

**Definition 0.37.** The category of double (cochain) complexes is  $\text{Ch}^{**}(\mathcal{A}) := \text{Ch}^*(\text{Ch}^*(\mathcal{A}))$ , so objects of  $\text{Ch}^{**}(\mathcal{A})$  are complexes of complexes

**Definition 0.38.**  $C \in \text{Ch}^{**}(\mathcal{A})$  is called *bounded*  $\iff \forall k \in \mathbb{Z} : \#\{(i, j) \in \mathbb{Z}^2 \mid i + j = k, C^{ij} \neq 0\} < \infty$ . Write  $\text{Ch}_b^{**}(\mathcal{A}) \subseteq \text{Ch}^{**}(\mathcal{A})$  for the full subcategory on bounded double complexes.

**Exercise 0.39.** If  $\mathcal{A}$  is additive or abelian, then so is  $\text{Ch}_b^{**}(\mathcal{A})$ .

**Definition 0.40.** The *total complex*  $\text{Tot}(C)$  of  $C \in \text{Ch}_b^{**}(\mathcal{A})$  is the complex  $\overline{C} \in \text{Ch}^*(\mathcal{A})$  defined as follows:

$$\overline{C}^k := \bigoplus_{\substack{(i,j) \in \mathbb{Z}^2 \\ i+j=k}} C^{ij}$$

$\delta_{\overline{C}}^k$  is constructed as follows:

$$\delta^{ij} : C^{ij} \xrightarrow{(\delta_1^{ij}, (-1)^i \delta_2^{ij})} C^{i+1,j} \oplus C^{i,j+1} \hookrightarrow \bigoplus_{i'+j'=k+1} C^{i'j'} = \overline{C}^{k+1}$$

Use the universal property of the direct sum to define

$$\delta_{\overline{C}}^k = \bigoplus_{i+j=k} \delta^{ij} : \overline{C}^k \rightarrow \overline{C}^{k+1}$$

TODO

**Exercise 0.41.** TODO

**Definition 0.42.** For  $C = (C^*, \delta^*) \in \text{Ch}^*(\mathcal{A})$  define:

- $Z^i(C) := \ker(\delta^i)$  as the  $i$ -th cocycle object.
- $B^i(C) := \text{im}(\delta^{i-1})$  as the  $i$ -th coboundary object.
- $u^i(C) : B^i(C) \rightarrow Z^i(C)$  the canonical monomorphism.
- $H^i(C) := \text{coker}(u^i(C)) = Z^i(C)/B^i(C)$  as the  $i$ -th cohomology object.

(Co-)homology measures the non-exactness of the complex  $C$ .

**Lemma 0.43** (Alternative description of cohomology object). *For  $C \in \text{Ch}^*(\mathcal{A})$  consider TODO*

**Lemma 0.44** (exer). *TODO*

**Theorem 0.45.** (a) *Given a s.e.s.*

$$\mathcal{E} : 0 \rightarrow C \xrightarrow{f} D \xrightarrow{g} E \rightarrow 0$$

*one obtains a long exact sequence*

$$\dots \rightarrow H^i(C) \xrightarrow{H^i(f)} H^i(D) \xrightarrow{H^i(g)} H^i(E) \xrightarrow{d_{\mathcal{E}}^i} H^{i+1}(C) \xrightarrow{H^{i+1}(f)} \dots$$

*for  $i \in \mathbb{Z}$ , where the connecting homomorphism  $d_{\mathcal{E}}^i$  is defined by the snake lemma and  $\mathcal{E} \mapsto d_{\mathcal{E}}^i$  is “functorial”.*

(b) *Given a morphism of short exact sequences  $\varphi : \mathcal{E} \rightarrow \mathcal{E}'$  in  $\text{Ch}^*(\mathcal{A})$ :*

$$\begin{array}{ccccccccc} \mathcal{E} : 0 & \longrightarrow & C & \xrightarrow{f} & D & \xrightarrow{g} & E & \longrightarrow & 0 \\ \varphi \downarrow & & \downarrow \alpha & & \downarrow \beta & & \downarrow \gamma & & \\ \mathcal{E}' : 0 & \longrightarrow & C' & \xrightarrow{f'} & D' & \xrightarrow{g'} & E' & \longrightarrow & 0 \end{array}$$

one obtains a commutative ladder of long exact sequences from (a)

$$\begin{array}{ccccccc}
\cdots & \longrightarrow & H^i(C) & \xrightarrow{H^i(f)} & H^i(D) & \xrightarrow{H^i(g)} & H^i(E) & \xrightarrow{d_{\mathcal{E}}^i} & H^{i+1}(C) & \longrightarrow & \cdots \\
& & \downarrow H^i(\alpha) & & \downarrow H^i(\beta) & & \downarrow H^i(\gamma) & & \downarrow H^{i+1}(\alpha) & & \\
\cdots & \longrightarrow & H^i(C') & \xrightarrow{H^i(f')} & H^i(D') & \xrightarrow{H^i(g')} & H^i(E') & \xrightarrow{d_{\mathcal{E}'}^i} & H^{i+1}(C') & \longrightarrow & \cdots
\end{array}$$

**Definition 0.46.**  $C \in \text{Ch}^*(\mathcal{A})$  is called *acyclic* if  $C$  is exact  $\iff H^i(C) = 0, \forall i \in \mathbb{Z}$ .

**Corollary 0.47** (To theorem 45). *Let  $\mathcal{E} : 0 \rightarrow C' \rightarrow C \rightarrow C'' \rightarrow 0$  be a s.e.s. in  $\text{Ch}^*(\mathcal{A})$ , then if any two of  $C', C, C''$  are acyclic, then so is the third.*

**Theorem 0.48** (Acyclicity criterion for total complexes). *Let  $C \in \text{Ch}_b^{**}(\mathcal{A})$  such that*

(a) *Each row  $(C^{ij}, \delta_1^{ij})_{i \in \mathbb{Z}}$  is acyclic  $\forall j \in \mathbb{Z}$ , or*

(b) *Each column  $(C^{ij}, \delta_1^{ij})_{j \in \mathbb{Z}}$  is acyclic  $\forall i \in \mathbb{Z}$*

*then  $\text{Tot}(C)$  is acyclic.*

**Definition 0.49.** An arrow  $f : C \rightarrow D$  in  $\text{Ch}^*(\mathcal{A})$  is called a *quasi-isomorphism* (qism)  $\iff \forall i \in \mathbb{Z} : H^i(f) : H^i(C) \rightarrow H^i(D)$  is an isomorphism.

**Lemma 0.50.** *TODO*

**Corollary 0.51.** *For  $f : C \rightarrow D$  in  $\text{Ch}^*(\mathcal{A})$  the following are equivalent:*

(a)  *$f$  is a quasi-isomorphism.*

(b) *Cone  $f$  is acyclic.*

**Definition 0.52.** Let  $f, g : C \rightarrow D$  in  $\text{Ch}^*(\mathcal{A})$ .

(a) A *homotopy* from  $f$  to  $g$  is a sequence of morphisms  $(s^i : C^i \rightarrow D^{i-1})_{i \in \mathbb{Z}}$  such that  $\forall i \in \mathbb{Z} : f^i - g^i = \delta_D^{i-1} \circ s^i + s^{i+1} \circ \delta_C^i$ , i.e.

$$\begin{array}{ccccc}
& & C^i & \xrightarrow{\delta_C^i} & C^{i+1} \\
& \swarrow s^i & \downarrow f^i - g^i & \swarrow s^{i+1} & \\
D^{i-1} & \xrightarrow{\delta_D^{i-1}} & D^i & & 
\end{array}$$

(b)  $f$  is called *homotopic* to  $g$  if  $\exists$  homotopy from  $f$  to  $g$ . Write  $f \sim g$ .

(c)  $f$  is called *nullhomotopic* if  $f \sim 0$ .

Note: This definition only requires that  $\mathcal{A}$  is additive.

**Proposition 0.53.** (a) *For  $C, D \in \text{Ch}^*(\mathcal{A})$ , homotopy defines an equivalence relation on  $\text{Hom}_{\text{Ch}^*(\mathcal{A})}(C, D)$ .*

(b) *For  $f, f' : C \rightarrow D$  and  $g, g' : D \rightarrow E$  in  $\text{Ch}^* \mathcal{A}$  one has:*

$$f \sim f', g \sim g' \implies g \circ f \sim g' \circ f'$$



(c) Suppose  $F : \mathcal{A} \rightarrow \mathcal{B}$  is an additive functor, then one has a functor  $F : \text{Ch}^* \mathcal{A} \rightarrow \text{Ch}^* \mathcal{B}$  by

$$F(\delta^n : C^n \rightarrow C^{n+1})_{n \in \mathbb{Z}} = (F\delta^n : FC^n \rightarrow FC^{n+1})_{n \in \mathbb{Z}}$$

This functor preserves homotopy, i.e.  $f \sim g \implies Ff \sim Fg$ .

**Proposition 0.54.** Suppose  $f, g : C \rightarrow D$  in  $\text{Ch}^* \mathcal{A}$  are homotopic, then

$$H^i(f) = H^i(g) : H^i(C) \rightarrow H^i(D), \forall i \in \mathbb{Z}$$

**Definition 0.55.** A morphism  $f : C \rightarrow D$  in  $\text{Ch}^* \mathcal{A}$  is called a *homotopy equivalence* from  $C$  to  $D$  if  $\exists g : D \rightarrow C$  in  $\text{Ch}^* \mathcal{A}$  such that

$$g \circ f \sim 1_C, \quad f \circ g \sim 1_D$$

and in this case  $C$  and  $D$  are called homotopy equivalent.

**Proposition 0.56.** Suppose  $f : C \rightarrow D$  is a homotopy equivalence, then  $f$  is a quasi-isomorphism.

**Example.** TODO

## 0.4 Injective and projective resolutions

**Notation.**  $\text{Inj}$  or  $\text{Inj}_{\mathcal{A}}$  and  $\text{Proj}$  or  $\text{Proj}_{\mathcal{A}}$  are the full subcategories of  $\mathcal{A}$  on injective or projective objects respectively. Note that  $\text{Inj}_{\mathcal{A}}^{\text{op}} = \text{Proj}_{\mathcal{A}^{\text{op}}}$ .

**Definition 0.57.** (a) An *injective resolution* of  $A \in \mathcal{A}$  is a quism  $f : \underline{A} \rightarrow I$  in  $\text{Ch}_{\geq 0}^* \mathcal{A}$  with  $I \in \text{Ch}_{\geq 0}^*(\text{Inj}_{\mathcal{A}})$  i.e.  $\text{Cone } f$ , which is

$$0 \rightarrow A \rightarrow I^0 \rightarrow I^1 \rightarrow I^2 \rightarrow \dots$$

is acyclic and all the  $I^j$  are injective.

(b) A *projective resolution* of  $A \in \mathcal{A}$  is a quism  $g : P \rightarrow \underline{A}$  in  $\text{Ch}_{\leq 0}^* \mathcal{A}$  with  $P \in \text{Ch}_{\leq 0}^*(\text{Proj}_{\mathcal{A}})$ .

**Proposition 0.58.** Consider the functor

$$\begin{array}{ccc} \hat{\cdot} : \text{Ch}^*(\mathcal{A})^{\text{op}} & \longrightarrow & \text{Ch}^*(\mathcal{A}^{\text{op}}) \\ (C^n, \delta_C^n) & \xrightarrow{\hat{\cdot}} & (\hat{C}^n, \hat{\delta}_C^n) \\ f \downarrow & & \uparrow \hat{f} \\ (D^n, \delta_D^n) & \xrightarrow{\hat{\cdot}} & (\hat{D}^n, \hat{\delta}_D^n) \end{array}$$

where  $\hat{C}^n = C^{-n}$ ,  $\hat{\delta}_C^n = \delta^{-n-1} \in \mathcal{A}(C^{-n-1}, C^{-n}) = \mathcal{A}^{\text{op}}(\hat{C}^n, \hat{C}^{n+1})$  and  $\hat{f}^n := f^{-n} \in \mathcal{A}(C^{-n}, D^{-n}) = \mathcal{A}^{\text{op}}(\hat{D}^n, \hat{C}^n)$ . Then

(a)  $\hat{\cdot}$  is well defined, satisfies  $\hat{\cdot} \circ \hat{\cdot} = \text{id}$  and  $\hat{\cdot}$  is an isomorphism of categories.

(b)  $\hat{\cdot}(\text{Ch}_{\geq 0/\leq 0}^*(\mathcal{A})^{\text{op}}) = \text{Ch}_{\leq 0/\geq 0}^*(\mathcal{A}^{\text{op}})$ .

$$(c) \sim (\text{Ch}_{\geq 0/\leq 0}^*(\text{Inj}_{\mathcal{A}} / \text{Proj}_{\mathcal{A}}))^{\text{op}} = \text{Ch}_{\leq 0/\geq 0}^*(\text{Proj}_{\mathcal{A}^{\text{op}}} / \text{Inj}_{\mathcal{A}^{\text{op}}})$$

(d) If  $\underline{A} \rightarrow I$  is an injective resolution in  $\text{Ch}^* \geq 0(\mathcal{A})$ , then  $\widehat{I} \rightarrow \underline{A}$  is a projective resolution in  $\text{Ch}_{\leq 0}^*(\mathcal{A}^{\text{op}})$

**Theorem 0.59.** Let  $\mathcal{A}$  be an abelian category with enough injectives, then

(a) Each  $A \in \mathcal{A}$  possesses an injective resolution.

(b) Let  $h : A \rightarrow B$  be a morphism, let  $f : \underline{B} \rightarrow I^\bullet$  be an injective resolution and  $g : \underline{A} \rightarrow C^\bullet$  be a quism in  $\text{Ch}_{\geq 0}^*(\mathcal{A})$  ( $C^\bullet$  is a resolution of  $A$ ), then there is a commutative diagram in  $\text{Ch}_{\geq 0}^*\mathcal{A}$

$$\begin{array}{ccc} \underline{A} & \xrightarrow{g} & C^\bullet \\ h \downarrow & & \downarrow H \\ \underline{B} & \xrightarrow{f} & I^\bullet \end{array}$$

(c) If diagram in (b) commutes with  $H, H' : C^\bullet \rightarrow I^\bullet$ , then  $H' \sim H$ .

**Corollary 0.60.** Suppose  $\underline{A} \xrightarrow{g} I^\bullet$  and  $\underline{A} \xrightarrow{g} J^\bullet$  are inj resolutions of  $A$ . Then:

(a)  $\exists H : I^\bullet \rightarrow J^\bullet$  such that

$$\begin{array}{ccc} & \underline{A} & \\ g \swarrow & & \searrow f \\ J^\bullet & \xleftarrow{\quad H \quad} & I^\bullet \end{array}$$

commutes.

(b)  $H$  in (a) is always a homotopy equivalence.

**Lemma 0.61** (Horseshoe 2). Let  $0 \rightarrow A' \rightarrow A \rightarrow A'' \rightarrow 0$  be a s.e.s. in  $\mathcal{A}$ , let  $\underline{A}' \xrightarrow{f'} I^\bullet$  and  $\underline{A}'' \xrightarrow{f''} J^\bullet$  be injective resolutions. Then  $\exists$  commutative diagram:

$$\begin{array}{ccccccc} 0 & \longrightarrow & \underline{A}' & \longrightarrow & \underline{A} & \longrightarrow & \underline{A}'' \longrightarrow 0 \\ & & \downarrow f' & & \downarrow f & & \downarrow f'' \\ 0 & \longrightarrow & I^\bullet & \longrightarrow & K^\bullet & \longrightarrow & J^\bullet \longrightarrow 0 \end{array}$$

in  $\text{Ch}_{\geq 0}^*\mathcal{A}$  with exact rows and an injective resolution  $f : \underline{A} \rightarrow K^\bullet$ . Moreover  $\forall n \geq 0$  the following commutative diagram has exact rows and columns:

$$\begin{array}{ccccccc} & & 0 & & 0 & & 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & \text{im}(\delta_I^{n-1}) & \longrightarrow & \text{im}(\delta_K^{n-1}) & \longrightarrow & \text{im}(\delta_J^{n-1}) \longrightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & I^n & \longrightarrow & K^n & \longrightarrow & J^n \longrightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & \text{im}(\delta_I^n) & \longrightarrow & \text{im}(\delta_K^n) & \longrightarrow & \text{im}(\delta_J^n) \longrightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ & & 0 & & 0 & & 0 \end{array}$$

So  $f' = \delta_I^{-1}$ ,  $f = \delta_K^{-1}$  and  $f'' = \delta_J^{-1}$ .

**Definition 0.62.** Define  $\text{Ex}_{\mathcal{A}}$  as the category of s.e.s. in  $\mathcal{A}$  with objects:

$$\mathcal{E}_{\mathcal{A}} : 0 \rightarrow A' \rightarrow A \rightarrow A'' \rightarrow 0$$

in  $\mathcal{A}$ , and morphisms are commutative diagrams

$$\begin{array}{ccccccccc} 0 & \longrightarrow & A' & \longrightarrow & A & \longrightarrow & A'' & \longrightarrow & 0 \\ & & \downarrow f' & & \downarrow f & & \downarrow f'' & & \\ 0 & \longrightarrow & B' & \longrightarrow & B & \longrightarrow & B'' & \longrightarrow & 0 \end{array}$$

in  $\mathcal{A}$  with exact rows represented by  $\underline{f} = (f', f, f'')$  in  $\text{Ex}_{\mathcal{A}}(\mathcal{E}_A, \mathcal{E}_B)$ . Composition of arrows is componentwise. TODO

**Proposition 0.63.** (a)  $\text{Ex}_{\mathcal{A}}$  is an additive category.

(b) We have additive functors  $\text{pr}_i : \text{Ex}_{\mathcal{A}} \rightarrow \mathcal{A}$ , mapping  $0 \rightarrow A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow 0$  to  $A_i$ .

(c)  $\text{Ex}_{\mathcal{A}}$  is not abelian.

**Definition 0.64.** An arrow  $\underline{f} = (f', f, f'')$  in  $\text{Ex}_{\mathcal{A}}$  is called

(a) a *strict monomorphism (strict epimorphism)*  $\iff f', f, f''$  are monics (epics) in  $\mathcal{A}$ .

(b) *strict*  $\iff$

$$0 \rightarrow \ker f' \rightarrow \ker f \rightarrow \ker f'' \rightarrow 0$$

is exact which is (by the snake lemma) equivalent to exactness of

$$0 \rightarrow \text{coker } f' \rightarrow \text{coker } f \rightarrow \text{coker } f'' \rightarrow 0$$

**Proposition 0.65.** If  $\underline{f}$  is strict in  $\text{Mor}(\text{Ex}_{\mathcal{A}})$  then  $\ker \underline{f}$ ,  $\text{coker } \underline{f}$ ,  $\text{im } \underline{f}$ ,  $\text{coim } \underline{f}$  exist in  $\text{Ex}_{\mathcal{A}}$  and the canonical map  $\text{coim } \underline{f} \rightarrow \text{im } \underline{f}$  is an isomorphism.

**Remark.**  $\text{Ex}_{\mathcal{A}}$  is an exact category.

**Definition 0.66.** (a) A complex  $\mathcal{E}^\bullet = (\underline{\delta}^i :)$  TODO

**Theorem 0.67.** TODO

## 0.5 Derived Functors

Let  $\mathcal{A}, \mathcal{B}$  be abelian categories.

**Definition 0.68.** (a) A homological (resp. cohomological)  $\delta$ -functor  $(T_n, \delta_n)_{n \geq 0}$  (resp.  $(T^n, \delta^n)_{n \geq 0}$ ) from  $\mathcal{A}$  to  $\mathcal{B}$  consists of

(i) a sequence of additive functors  $(T_n : \mathcal{A} \rightarrow \mathcal{B})_{n \geq 0}$  resp.  $(T^n : \mathcal{A} \rightarrow \mathcal{B})_{n \geq 0}$ .

(ii) A sequence of natural transformations

$$\begin{array}{ccc} \text{Ex}_{\mathcal{A}} & \begin{array}{c} \xrightarrow{T_{n+1} \circ \text{pr}_3} \\ \Downarrow \delta_n \\ \xrightarrow{T_n \circ \text{pr}_1} \end{array} & \mathcal{B} \\ \text{resp.} & & \\ \text{Ex}_{\mathcal{A}} & \begin{array}{c} \xrightarrow{T^n \circ \text{pr}_3} \\ \Downarrow \delta^n \\ \xrightarrow{T^{n+1} \circ \text{pr}_1} \end{array} & \mathcal{B} \end{array}$$

such that this data assigns to each  $\underline{\mathcal{E}} : 0 \rightarrow A_1 \xrightarrow{f} A_2 \xrightarrow{g} A_3 \rightarrow 0$  in  $\text{Ex}_{\mathcal{A}}$  a long exact sequence in  $\mathcal{B}$

$$\cdots \rightarrow T_{n+1} \xrightarrow{\delta_n^{\underline{\mathcal{E}}}} T_n A_1 \xrightarrow{T_n f} T_n A_2 \xrightarrow{T_n g} T_n A_3 \xrightarrow{\delta_{n-1}^{\underline{\mathcal{E}}}} T_{n-1} A_1 \rightarrow \cdots$$

resp.

$$\cdots \rightarrow T^{n-1} A_3 \xrightarrow{\delta_{\mathcal{E}}^{n-1}} T^n A_1 \xrightarrow{T^n f} T^n A_2 \xrightarrow{T^n g} T^n A_3 \xrightarrow{\delta_{\mathcal{E}}^n} T^{n+1} A_1 \rightarrow \cdots$$

Moreover the assignment  $\mathcal{E} \rightarrow \text{l.e.s.}$  is functorial in  $\text{Ex}_{\mathcal{A}}$

(iii) TODO WTF IS THIS?

**Example 0.69.** TODO

**Construction 0.70.** Suppose  $\mathcal{A}$  has enough injectives and let  $F : \mathcal{A} \rightarrow \mathcal{B}$  be a left exact additive functor.

1.  $\forall A \in \mathcal{A}$  choose an injective resolution  $\iota_A : \underline{A} \rightarrow I_A^\bullet$  in  $\text{Ch}_{\geq 0}^*(\mathcal{A})$  and define  $R^n F(A) := H^n(FI_A^\bullet)$  choice in  $\mathcal{B}$  of  $n$ -th cohomology object.
2.  $\forall f : A \rightarrow A'$  morphism in  $\mathcal{A}$  choose an arrow  $\iota_f$  such that

$$\begin{array}{ccc} \underline{A}' & \xrightarrow{\iota_{A'}} & I_{A'}^\bullet \\ f \uparrow & & \uparrow \iota_f \\ \underline{A} & \xrightarrow{\iota_A} & I_A^\bullet \end{array}$$

commutes in  $\text{Ch}_{\geq 0}^*(\mathcal{A})$  which implies that  $F(\iota_f) : FI_A^\bullet \rightarrow FI_{A'}^\bullet$  morphism in  $\text{Ch}_{\geq 0}^*(\mathcal{B})$ . Define:  $R^n F(f) := H^n(F(\iota_f)) : R^n F(A) \rightarrow R^n F(A')$  ...

**Lemma 0.71.** (a)  $R^n F$  is a functor  $\mathcal{A} \rightarrow \mathcal{B}$  (additive.)

(b) If one makes other choices of injective resolutions  $\tilde{\iota}_A : A \rightarrow \tilde{I}_A^\bullet$  and  $\tilde{\iota}_f : \tilde{I}_A^\bullet \rightarrow \tilde{I}_{A'}^\bullet$ , then we get a natural isomorphism  $R^n F \cong \tilde{R}^n F$

3. Given  $\mathcal{E} : 0 \rightarrow A_1 \xrightarrow{f} A_2 \xrightarrow{g} A_3 \rightarrow 0$  in  $\text{Ex}_{\mathcal{A}}$  and an injective resolution  $\iota_{\mathcal{E}} : \mathcal{E} \rightarrow J_{\mathcal{E}}^\bullet$  in  $\text{Ch}_{\geq 0}^*(\text{Ex}_{\mathcal{A}})$

**Lemma 0.72.** (a)  $FJ_{\mathcal{E}}^\bullet$  lies in  $\text{Ch}_{\geq 0}^*(\text{Ex}_{\mathcal{B}})$ , in particular, it is a s.e.s. of complexes in  $\text{Ch}_{\geq 0}^*(\mathcal{B})$ .

- (b) Write  $J_{\mathcal{E}}^{\bullet} = 0 \rightarrow I_1^{\bullet} \rightarrow I_2^{\bullet} \rightarrow I_3^{\bullet} \rightarrow 0$  as a s.e.s. in  $\text{Ch}_{\geq 0}^*(\mathcal{A})$ , then the  $I_i^n$  are injective and by (a) we have a s.e.s.

$$0 \rightarrow FI_1^{\bullet} \rightarrow FI_2^{\bullet} \rightarrow FI_3^{\bullet} \rightarrow 0 \quad (*)$$

in  $\text{Ch}_{\geq 0}^*(\mathcal{B})$ , then the following diagram commutes and the top row is a l.e.s. in  $\mathcal{B}$ , and the vertical maps are isomorphisms.

$$\begin{array}{ccccccc} \cdots & \longrightarrow & H^{n-1}(FI_j^{\bullet}) & \xrightarrow{\delta_{FJ_{\mathcal{E}}}^{n-1}} & H^n(FI_1^{\bullet}) & \longrightarrow & H^n(FI_2^{\bullet}) \longrightarrow H^n(FI_3^{\bullet}) \xrightarrow{\delta_{FJ_{\mathcal{E}}}^n} \cdots \\ & & u_{A_3}^{n-1} \uparrow & & u_{A_1}^n \uparrow & & u_{A_2}^n \uparrow & & u_{A_3}^n \uparrow \\ \cdots & \longrightarrow & R^{n-1}F(A_3) & \longrightarrow & R^nF(A_1) & \xrightarrow{R^nF(f)} & R^nF(A_2) & \xrightarrow{R^nF(g)} & R^nF(A_3) \longrightarrow \cdots \end{array}$$

This is not even funny anymore. Define  $\delta_{RF}^n : R^nF \circ \text{pr}_3 \Rightarrow R^nF \circ \text{pr}_1$  for  $\mathcal{E}$  as  $(u_{A_1}^{n+1})^{-1} \circ \delta_{FJ_{\mathcal{E}}}^n \circ u_{A_3}^n$ .

4. Show that  $\delta_{RF}^n$  is well defined.

**Lemma 0.73.** *TODO*

**Theorem 0.74.** *Suppose  $\mathcal{A}$  has enough injectives and  $F$  is an additive left exact functor, then  $RF := (R^nF, \delta_{RF}^n)_{n \geq 0}$  is a universal cohomological  $\delta$ -functor and  $R^nF$  is called the  $n$ -th right derived functor of  $F$ . It satisfies:*

- (a)  $R^0F = F$ .  
(b)  $I \in \text{Inj}_{\mathcal{A}} \implies \forall n \geq 1, R^nF(I) = 0$ .

Suppose  $\mathcal{A}$  has enough projectives, let  $G : \mathcal{A} \rightarrow \mathcal{B}$  be right exact, then

- (a)  $\exists$  homological  $\delta$ -functor  $LG = (L_iG, \delta_i)_{i \geq 0}$  such that  $(*)$

- (i)  $L_0G = G$   
(ii)  $P \in \text{Proj}_{\mathcal{A}} \implies \forall n \geq 1, L_nG(P) = 0$ .

- (b)  $LG$  with  $(*)$  is universal.

- (c)  $LG$  with  $(*)$  is unique up to unique isomorphism.

- (d)  $LG$  can be computed via projective resolutions, i.e.  $\forall A \in \mathcal{A}$  with projective resolution  $P^{\bullet} \rightarrow A$  in  $\text{Ch}_{\leq 0}^*(\mathcal{A})$  we have

$$L_iG(A) = H^{-i}(GP^{\bullet})$$

**Lemma.** *If  $T = (T_n, \delta_n)_{n \geq 0}$  is a homological  $\delta$ -functor from  $\mathcal{A}$  to  $\mathcal{B}$  and if  $\mathcal{A}$  has enough projectives and  $T_iP = 0, \forall P \in \text{Proj}_{\mathcal{A}}, i \geq 1$ , then  $T$  is universal.*

## 0.6 The Ext Functor

Let  $\mathcal{A}$  be abelian,  $M, N \in \mathcal{A}$  and suppose  $\mathcal{A}$  has enough projectives/injectives if needed.

**Definition 0.75.** Define

$$\mathrm{Ext}_{\mathcal{A}}^i(-, N) := R^i \mathrm{Hom}_{\mathcal{A}}(-, N)$$

with the natural transformations  $\delta^i$ .

**Example.** To compute: if

$$\cdots \rightarrow P^{-2} \rightarrow P^{-1} \rightarrow P^0 \rightarrow M \quad (*)$$

is a projective resolution, then we have the complex  $P^\bullet$ , apply  $\mathrm{Hom}_{\mathcal{A}}(-, N)$  and get

$$0 \rightarrow \mathrm{Hom}_{\mathcal{A}}(P^0, N) \rightarrow \mathrm{Hom}_{\mathcal{A}}(P^{-1}, N) \rightarrow \mathrm{Hom}_{\mathcal{A}}(P^{-2}, N) \rightarrow \cdots$$

and  $\mathrm{Ext}_{\mathcal{A}}^i(M, N) = H^i(\mathrm{Hom}_{\mathcal{A}}(P^\bullet, N))$

**Proposition 0.76.** (a)  $M$  projective  $\implies \mathrm{Ext}_{\mathcal{A}}^i(M, N) = 0, \forall i \geq 1$ .

(b)  $N$  injective  $\implies \mathrm{Ext}_{\mathcal{A}}^i(M, N) = 0, \forall i \geq 1$ .

**Definition 0.77.** Define

$$\overline{\mathrm{Ext}}_{\mathcal{A}}^i(M, -) := R^i(\mathrm{Hom}_{\mathcal{A}}(M, -))$$

with the natural transformations  $\delta^i$ .

**Example.** Compute  $\overline{\mathrm{Ext}}_{\mathcal{A}}^i(M, -)$  via injective resolutions of the 2nd argument:

$$0 \rightarrow N \rightarrow I^0 \rightarrow I^1 \rightarrow I^2 \rightarrow \cdots$$

and we get

$$\mathrm{Hom}_{\mathcal{A}}(M, I^0) \rightarrow \mathrm{Hom}_{\mathcal{A}}(M, I^1) \rightarrow \mathrm{Hom}_{\mathcal{A}}(M, I^2) \rightarrow \cdots$$

so  $\overline{\mathrm{Ext}}_{\mathcal{A}}^i(M, -) = H^i(\mathcal{A}(M, I^\bullet))$ .

**Proposition 0.78.** For  $M$  projective or  $N$  injective we have  $\overline{\mathrm{Ext}}_{\mathcal{A}}^i(M, N) = 0, \forall i \geq 1$ .

**Remark.** If  $\mathcal{A}$  has enough projectives and injectives,  $\mathrm{Ext}_{\mathcal{A}}^i(-, -)$  and  $\overline{\mathrm{Ext}}_{\mathcal{A}}^i(-, -)$  turn out to be isomorphic!

**Remark.**  $\mathrm{Ext}_{\mathcal{A}}^i(-, -)$  and  $\overline{\mathrm{Ext}}_{\mathcal{A}}^i(-, -)$  are bifunctors.

**Theorem 0.79.** Suppose  $\mathcal{A}$  has enough injectives and projectives, then  $\exists$  natural isomorphisms as bifunctors

$$\begin{aligned} \mathcal{A}^{\mathrm{op}} \times \mathcal{A} &\rightarrow \mathbf{Ab} \\ u_{M,N}^i : \mathrm{Ext}_{\mathcal{A}}^i(M, N) &\rightarrow \overline{\mathrm{Ext}}_{\mathcal{A}}^i(M, N) \end{aligned}$$

### 0.6.1 Classical interpretation of $\text{Ext}^1$ as extension objects

**Definition 0.80.** (a) An extension of  $M$  by  $N$  is a s.e.s. in  $\mathcal{A}$ :

$$\mathcal{E} : 0 \rightarrow N \rightarrow E \rightarrow M \rightarrow 0$$

(b) Extensions  $\mathcal{E}$  and

$$\mathcal{E}' : 0 \rightarrow N \rightarrow E' \rightarrow M \rightarrow 0$$

are called *equivalent* (write  $\sim$ ), iff  $\exists$  commutative diagram:

$$\begin{array}{ccccccccc} \mathcal{E} : 0 & \longrightarrow & N & \longrightarrow & E & \longrightarrow & M & \longrightarrow & 0 \\ & & \downarrow \text{id}_N & & \downarrow \varphi & & \downarrow \text{id}_M & & \\ \mathcal{E}' : 0 & \longrightarrow & N & \longrightarrow & E' & \longrightarrow & M & \longrightarrow & 0 \end{array}$$

(c)  $\text{Ext}_{\mathcal{A}}^1(M, N)$  denotes the set of equivalence classes of extensions of  $M$  by  $N$ .

**Remark.** (a) By the snake lemma:  $\mathcal{E} \sim \mathcal{E}' \implies E \cong E'$ .

(b)  $\sim$  is an equivalence relation.

**Proposition 0.81.** (a)  $\text{Ext}_{\mathcal{A}}(M, N)$  is an abelian group for addition  $\mathcal{E} + \mathcal{E}'$  defined by the Baer sum:

$$\begin{array}{ccccccccc} 0 & \longrightarrow & N \oplus N & \longrightarrow & E \oplus E' & \longrightarrow & M \oplus M & \longrightarrow & 0 \\ & & \downarrow \Sigma & & \downarrow & & \downarrow \text{id}_{M \oplus M} & & \\ 0 & \longrightarrow & N & \longrightarrow & E \oplus E' \amalg_{N \oplus N} N & \longrightarrow & M \oplus M & \longrightarrow & 0 \\ & & \uparrow \text{id}_N & & \uparrow & & \uparrow \Delta & & \\ 0 & \longrightarrow & N & \longrightarrow & E + E' & \longrightarrow & M & \longrightarrow & 0 \end{array}$$

where  $E \oplus E' \amalg_{N \oplus N} N$  is the pushout by the sum map  $+: N \oplus N \rightarrow N, (n, m) \mapsto n + m$  and  $E + E'$  is the pullback by the diagonal map  $\Delta : M \rightarrow M \oplus M, m \mapsto (m, m)$ . And the zero object is given by the split s.e.s

$$0 \rightarrow N \rightarrow N \oplus M \rightarrow M \rightarrow 0$$

(b) One has a natural isomorphism as bifunctors

$$u : \text{Ext}_{\mathcal{A}}(-, -) \Rightarrow \text{Ext}_{\mathcal{A}}^1(-, -)$$

if  $\mathcal{A}$  has enough projectives (or to  $\overline{\text{Ext}}_{\mathcal{A}}^1(-, -)$  if  $\mathcal{A}$  has enough injectives).

**Remark.** In fact, Yoneda also considered higher Ext-groups (in the absence of injectives/projectives), e.g.  $\text{Ext}^2(M, N)$  is the “group” of exact sequences

$$0 \rightarrow N \rightarrow E_1 \rightarrow E_2 \rightarrow M \rightarrow 0$$

modulo a suitable equivalence relation.

**Notation.** For  ${}_R\text{Mod}$  one usually abbreviates

$$\text{Ext}_R^i(-, -) := \text{Ext}_{{}_R\text{Mod}}^i(-, -)$$

For  $R = \mathbb{Z}[G]$  and a group  $G$ , one also considers the group cohomology for  $M \in {}_{\mathbb{Z}[G]}\mathbf{Mod}$  as

**Definition 0.82.** Define

$$H^i(G, M) := \mathrm{Ext}_{\mathbb{Z}[G]}^i(\mathbb{Z}, M)$$

such that  $\mathbb{Z}$  is the trivial  $G$ -action.

**Proposition.** In face  $H^i(G, M)$  is the  $i$ -th right derived functor of

$$\begin{aligned} M \rightarrow \mathrm{Hom}_{\mathbb{Z}[G]}(\mathbb{Z}, M) &\stackrel{\text{exer.}}{=} M^G := \{m \in M \mid \forall g \in G : gm = m\} \\ \varphi_m : n \mapsto nm &\leftarrow m \in M^G \end{aligned}$$

So far we had  $\mathrm{Ext}_{\mathcal{A}}^i$  valued in  $\mathbf{Ab}$ . But recall that  $\mathrm{Hom}_R(-, -)$  is also a bifunctor

$$\begin{aligned} \mathrm{Hom}_R(-, -) : {}_R\mathbf{Mod}_{R'} \times {}_R\mathbf{Mod}_{R''} &\rightarrow {}_{R'}\mathbf{Mod}_{R''} \\ (M, N) &\mapsto \mathrm{Hom}_R(M, N) \end{aligned}$$

Fact:  ${}_R\mathbf{Mod}_{R'}$  and  ${}_R\mathbf{Mod}_{R''}$  have enough injectives and projectives because  ${}_R\mathbf{Mod}_{R'}$  is isomorphic to  ${}_{R \otimes (R')^{\mathrm{op}}}\mathbf{Mod}$ .

**Proposition 0.83** (Exer). *This Bifunctor induces*

$$\mathrm{Ext}_R^i(-, -) : {}_R\mathbf{Mod}_{R'} \times {}_R\mathbf{Mod}_{R''} \rightarrow {}_{R'}\mathbf{Mod}_{R''}$$

## 0.7 The Tor Functor

Let  $R$  be a ring,  $M \in \mathbf{Mod}_R$  and  $N \in {}_R\mathbf{Mod}$ . We had the bifunctor

$$- \otimes - : \mathbf{Mod}_R \times {}_R\mathbf{Mod} \rightarrow \mathbf{Ab}$$

and it is right exact, and also  $\mathbf{Mod}_R$  and  ${}_R\mathbf{Mod}$  have enough projectives.

**Definition 0.84.** For  $i \geq 0$  define:

$$\begin{aligned} \mathrm{Tor}_i^R(M, -) &:= L_i(M \otimes_R -) \\ \overline{\mathrm{Tor}}_i^R(-, N) &:= L_i(- \otimes_R N) \end{aligned}$$

**Proposition 0.85.** *If  $P^\bullet \rightarrow M$  (or  $Q^\bullet \rightarrow N$ ) are projective resolutions, then one has:*

$$\begin{aligned} \overline{\mathrm{Tor}}_i^R(M, N) &= H^{-i}(P^\bullet \otimes_R N) \\ \mathrm{Tor}_i^R(M, N) &= H^{-i}(M \otimes_R Q^\bullet) \end{aligned}$$

Moreover, if  $M$  or  $N$  are projective, then

$$\mathrm{Tor}_i^R(M, N) \cong 0 \cong \overline{\mathrm{Tor}}_i^R(M, N)$$

**Theorem 0.86.** *One has natural isomorphism of bifunctors*

$$\mathrm{Tor}_i^R(-, -) \cong \overline{\mathrm{Tor}}_i^R(-, -)$$



Other ways to compute  $Tor_i$ , not using projective resolutions.

**Definition 0.87.** Let  $T = (T_n, \delta_n)$  be a homological  $\delta$ -functor from  $\mathcal{A}$  to  $\mathcal{B}$ . Call  $A \in \mathcal{A}$   $T$ -acyclic if  $T_i A = 0, \forall i \geq 1$  (with  $T_0$  right exact) For example  $M$  projective  $\implies M$  is  $L(- \otimes_R N)$ -acyclic

**Facts 0.88.** Let  $0 \rightarrow A' \rightarrow A \rightarrow A'' \rightarrow 0$  be a s.e.s. in  $\mathcal{A}$ , then

(a)  $A''$  is  $T$ -acyclic then

$$0 \rightarrow T_0 A' \rightarrow T_0 A \rightarrow T_0 A'' \rightarrow 0$$

is a s.e.s. (because  $T_1 A'' = 0$  and l.e.s. from  $\delta$ -functor).

(b) If  $A$  and  $A''$  are  $T$ -acyclic, then so is  $A'$  (because of the l.e.s. from being a  $\delta$ -functor).

$$\rightarrow T_2 A' \rightarrow \underbrace{T_2 A}_{=0} \rightarrow \underbrace{T_2 A''}_{=0} \rightarrow T_1 A' \rightarrow \underbrace{T_1 A}_{=0} \rightarrow \underbrace{T_1 A''}_{=0} \rightarrow T_0 A' \rightarrow T_0 A \rightarrow T_0 A''$$

(c) If  $0 \rightarrow A_0 \rightarrow A_1 \rightarrow \dots \rightarrow A_n \rightarrow 0$  is exact in  $\mathcal{A}$  and if  $A_1, \dots, A_n$  are  $T$ -acyclic then  $A_0$  is also  $T$ -acyclic. This follows from (b) by ind. using the exact sequences  $0 \rightarrow A_0 \rightarrow A_1 \rightarrow X \rightarrow 0$  and

$$0 \rightarrow X \rightarrow A_2 \rightarrow \dots \rightarrow A_n \rightarrow 0$$

**Lemma 0.89.** If  $C^\bullet \in \text{Ch}_{\leq 0}^*(\mathcal{A})$  is acyclic with all  $C^i$   $T$ -acyclic, then  $T_0 C^\bullet$  is acyclic (we assume  $T_0$  is right exact)

**Corollary 0.90.** Suppose  $\mathcal{A}$  has enough projectives,  $G : \mathcal{A} \rightarrow \mathcal{B}$  is right exact and  $T = LG$ . If  $Q^\bullet \xrightarrow{g} \underline{A}$  is a resolution by  $T$ -acyclic objects. Then:

$$L_i G(A) \cong H^{-i}(GQ^\bullet)$$

**Definition 0.91.**  $M \in \text{Mod}_R$  is called *flat* if  $M \otimes_R - : {}_R\text{Mod} \rightarrow \text{Ab}$  is exact.

**Proposition 0.92.** For  $M \in \text{Mod}_R$  the following are equivalent:

- (a)  $M$  is flat.
- (b)  $\text{Tor}_1^R(M, -) = 0$ .
- (c)  $\text{Tor}_i^R(M, -) = 0 \forall i \geq 1$ .

**Theorem 0.93.** For  $M \in \text{Mod}_R$  the following are equivalent:

- (a)  $M$  is flat.
- (b)  $\forall N \in {}_R\text{Mod} : M$  is  $L(- \otimes_R N)$ -acyclic.

And thus  $\text{Tor}_i^R$  can be computed via flat resolutions.