

# Bidding Algorithm Overview

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Brand Auction

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# Agenda

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## ① Bidding Products

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- 1 **Bidding Products**
- 2 **Auto-Bidding Algorithm**

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- ② **Auto-Bidding Algorithm**
- ③ **Cost Cap Bidding Algorithm**

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# Auto Bidding Algorithm

## Formulation(use CPC ads as example)

$$\max_{x_t} \sum_{t=1}^T x_t \cdot r_t \quad \text{s.t.} \quad \sum_{t=1}^T x_t \cdot c_t \leq B$$

where

$t$  – t-th auction opportunity

$T$  – total auction opportunity

$r_t$  – pctr at  $t$

$c_t$  – cost at  $t$

$B$  – total budget

$x_t$  – indicator whether t-th auction win or lose,  $x_t \in \{0, 1\}$

# Auto Bidding Algorithm

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**Idea:** use primal-dual + online gradient descent to derive an optimal bidding formula for auto-bidding campaigns

## Optimal Bidding Formula(cont'd)



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### STEP 1: Write Lagrangian

$$\begin{aligned}\mathcal{L}(x_t, \lambda) &= \sum_{t=1}^T x_t r_t + \lambda \left( B - \sum_{t=1}^T x_t c_t \right) \\ &= \sum_{t=1}^T x_t (r_t - \lambda c_t) + \lambda B\end{aligned}$$

## Optimal Bidding Formula(cont'd)

### STEP 2: Derive dual problem

$$\begin{aligned}\mathcal{L}^*(\lambda) &= \max_{x_t} \mathcal{L}(x_t, \lambda) \\ &= \sum_{t=1}^T (r_t - \lambda c_t)_+ + \lambda B\end{aligned}$$

where  $(z)_+ = \max(0, z)$  (a.k.a ReLU)

## Optimal Bidding Formula(cont'd)

### STEP 3: Solve dual problem

$$\lambda^* = \min_{\lambda \geq 0} \mathcal{L}^*(\lambda)$$

$$\Rightarrow \text{ecpm}_t^* = \frac{r_t}{\lambda^*}$$

i.e.

$$b_t^* \cdot r_t = \frac{r_t}{\lambda^*} \Rightarrow b_t^* = \frac{1}{\lambda^*}$$

*Note: this is true for all incentive compatible auctions*

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**Conclusion:** optimal bid is constant bid

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## How To Find Optimal Bid $b^*$

(From KKT condition)  $b^*$  is the bid level that exactly depletes the budget, i.e.,

$$\sum_{t=1}^T x_t c_t \leq B \implies \sum_{t=1}^T x_t c_t = B$$

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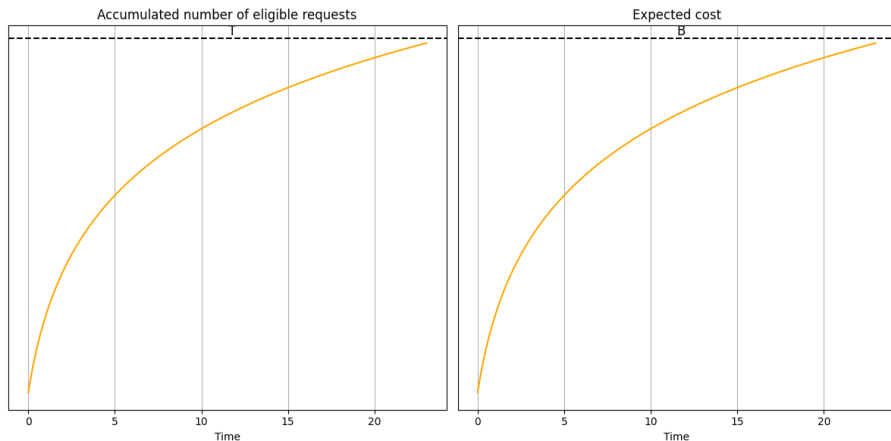
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Assume  $r_t, c_t$  are subject to i.i.d. at  $(\tau, \tau + d\tau)$ , the cost

$$\sum_{\tau \leq t \leq \tau + d\tau} x_t c_t \sim \# \text{ of auction opportunities in } (\tau, \tau + d\tau)$$

# Auto Bidding Algorithm



**Figure:** Left: Accumulated number of eligible requests. Right: Accumulated expected cost.

# Auto Bidding Algorithm

In practice, we construct an expected cost curve based on flow ratio map (e.g. 15 minute granularity):

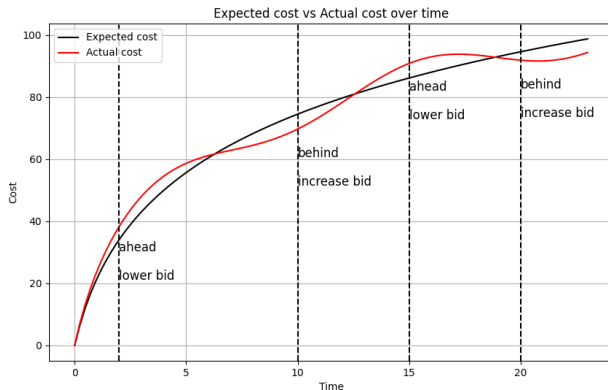


Figure: Expected cost vs Actual cost over time



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By following the curve:  $b_t \rightarrow b^*$

# Auto Bidding Algorithm

## Controller-based Algorithm

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- 1  $e(t_k) = C_e(t_k) - C_a(t_k)$
- 2  $\phi(t_{k+1}) \leftarrow \lambda_P e(t_k) + \lambda_I \sum_{j=1}^k e(t_j) \Delta t_j + \lambda_D \frac{\Delta e(t_k)}{\Delta t_k}$

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### Approach 1: PID Controller

- ①  $e(t_k) = C_e(t_k) - C_a(t_k)$
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- ③  $b(t_{k+1}) \leftarrow b(t_k) \cdot \exp(\phi(t_{k+1}))$

## Controller-based Algorithm(cont'd)

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MPC optimizes the entire future spend, at time  $t$ , remaining budget  $B_{t,r}$

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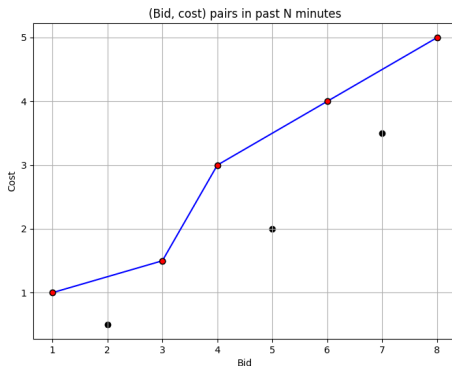
⇒ expected cost speed  $cs_t$

⇒ expected bid price  $b_t$

**Goal:** find  $cs_t = f(b_t)$  where  $f$  is monotonically non-decreasing

# Auto Bidding Algorithm

iMPC: Longest Increasing Subsequence of historical (bid, cost) pairs  
+ interpolation



**Figure:** (Bid, cost) pair in past N minutes. Blue line represents the longest increasing subsequence used for interpolation.

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## iMPC Controller(cont'd)

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### iMPC algorithm:

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- 2 Choose longest increasing subsequence(LIS), then interpolate to get  $f$

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### iMPC algorithm:

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- 3  $b_t = f^{-1}(cs_t)$

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**Idea:** Use Isotonic Regression

# Auto Bidding Algorithm

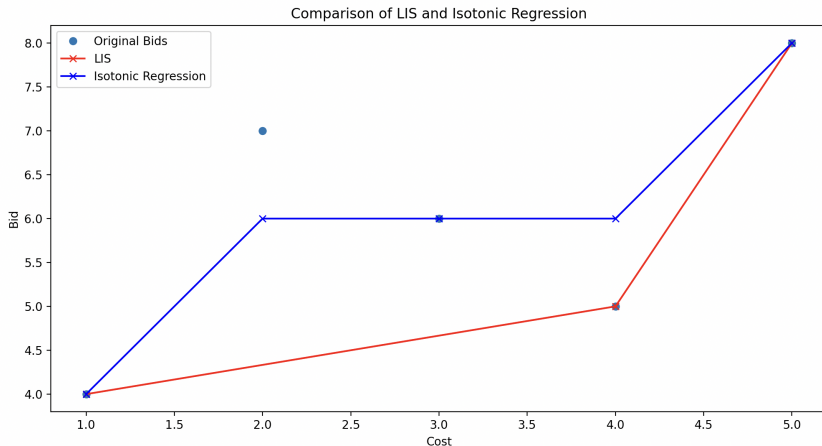


Figure: Isotonic Regression vs LIS.



# Cost Cap Algorithm

Formulation(use CPC ads as example)

$$\max_{x_t} \sum_{t=1}^T x_t r_t$$

subject to

$$\sum_{t=1}^T x_t c_t \leq B$$

$$\sum_{t=1}^T x_t c_t \leq C \left( \sum_{t=1}^T x_t r_t \right)$$

where  $C$  is cost cap (advertiser bid)

# Cost Cap Algorithm

## Optimal Bidding Formula

Use primal-dual method, similarly we get optimal bid:

$$b^* = \frac{1 + \mu^* C}{\lambda^* + \mu^*} = \frac{\lambda^*}{\lambda^* + \mu^*} \cdot \frac{1}{\lambda^*} + \frac{\mu^*}{\lambda^* + \mu^*} \cdot C$$

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In practice, we use the "cost-min" algorithm:

$$b_t = \min(\text{flow\_bid}, \text{risk\_bid})$$

where

$$\text{risk\_bid} = \frac{\text{remaining budget}}{\text{remaining event goal}}$$

flow\_bid is the same as in auto-bidding.