Bidding Algorithm Primer

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Brand Auction

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Bidding Products

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- Auto-Bidding Algorithm

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- Cost Cap Bidding Algorithm

Three Major Bidding Products

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Formulation(use CPC ads as example)

$$\max_{x_t} \sum_{t=1}^{T} x_t \cdot r_t \quad \text{s.t.} \quad \sum_{t=1}^{T} x_t \cdot c_t \le B$$

where

t — t-th auction opportunity

T – total auction opportunity

 r_t — pctr at t

 c_t — cost at t

B − total budget

 x_t — indicator whether t-th auction win or lose, $x_t \in \{0,1\}$



Optimal Bidding Formula

Auction Mechanism: assume we use second price auction, then

$$\mathbf{x}_t = \mathbf{1}_{\{\mathsf{ecpm}_t > c_t\}} \quad \mathsf{where} \quad \mathsf{ecpm}_t = b_t * r_t$$

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Optimal Bidding Formula(cont'd)

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STEP 1: Write Lagrangian

$$\mathcal{L}(x_t, \lambda) = \sum_{t=1}^{T} x_t r_t + \lambda \left(B - \sum_{t=1}^{T} x_t c_t \right)$$
$$= \sum_{t=1}^{T} x_t (r_t - \lambda c_t) + \lambda B$$

Optimal Bidding Formula(cont'd)

STEP 2: Derive dual problem

$$\mathcal{L}^*(\lambda) = \max_{x_t} \mathcal{L}(x_t, \lambda)$$

$$=\sum_{t=1}^T (r_t - \lambda c_t)_+ + \lambda B$$

where $(z)_+ = \max(0, z)$ (a.k.a ReLU)

Optimal Bidding Formula(cont'd)

STEP 3: Solve dual problem

$$\lambda^* = \min_{\lambda \geq 0} \mathcal{L}^*(\lambda)$$

$$\Rightarrow \mathsf{ecpm}_t^* = \frac{r_t}{\lambda^*}$$

i.e.

$$b_t^* \cdot r_t = rac{r_t}{\lambda^*} \Rightarrow b_t^* = rac{1}{\lambda^*}$$

Note: this is true for all incentive compatible auctions

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Conclusion: optimal bid is constant bid

How To Find Optimal Bid b^*

(From KKT condition) b^* is the bid level that exactly depletes the budget, i.e.,

$$\sum_{t=1}^{T} x_t c_t \leq B \implies \sum_{t=1}^{T} x_t c_t = B$$

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Assume r_t , c_t are subject to i.i.d. at $(\tau, \tau + d\tau)$, the cost

$$\sum_{ au < t < au + d au} x_t c_t \sim \#$$
 of auction opportunities in $(au, au + d au)$

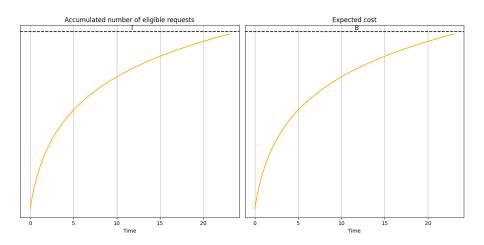


Figure: Left: Accumulated number of eligible requests. Right:Accumulated expected cost.

In practice, we construct an expected cost curve based on flow ratio map (e.g. 15 minute granularity):

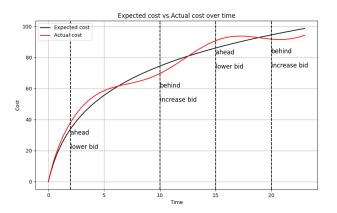


Figure: Expected cost vs Actual cost over time

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By following the curve: $b_t \to b^*$

Controller-based Algrorithm

Approach 1: PID Controller

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 $b(t_{k+1}) \leftarrow b(t_k) \cdot \exp(\phi(t_{k+1}))$

Controller-based Algrorithm(cont'd)

Approach 2: MPC Controller

MPC optimizes the entire future spend, at time t, remaining budget $B_{t,r}$

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- \Rightarrow expected bid price b_t

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Goal: find $cs_t = f(b_t)$

iMPC: fitting historical (bid, cost) pairs + interpolation

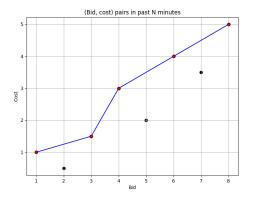


Figure: (Bid, cost) pair in past N minutes. Blue line represents the longest increasing subsequence used for interpolation.

iMPC Controller(cont'd)

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iMPC Controller(cont'd)

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- \odot Choose longest increasing subsequence, then interpolate to get f

iMPC Controller(cont'd)

- **1** At a given time, $cs_t = f(b_t)$ (non-decreasing)
- $b_t = f^{-1}(cs_t)$

MPC Controller

Isotonic Regression based MPC

Formulation(use CPC ads as example)

$$\max_{x_t} \sum_{t=1}^{T} x_t r_t$$

subject to

$$\sum_{t=1}^{T} x_t c_t \leq B$$

$$\sum_{t=1}^{T} x_t c_t \le C \left(\sum_{t=1}^{T} x_t r_t \right)$$

where C is cost cap (advertiser bid)

Optimal Bidding Formula

Use primal-dual method, similarly we get optimal bid:

$$b^* = \frac{1 + \mu^* C}{\lambda^* + \mu^*} = \frac{\lambda^*}{\lambda^* + \mu^*} \cdot \frac{1}{\lambda^*} + \frac{\mu^*}{\lambda^* + \mu^*} \cdot C$$

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$$b_t = \min(\text{flow_bid}, \text{risk_bid})$$

where

$$risk_bid = \frac{remaining \ budget}{remaining \ event \ goal}$$

flow_bid is the same as in auto-bidding.