



Technische
Universität
Braunschweig

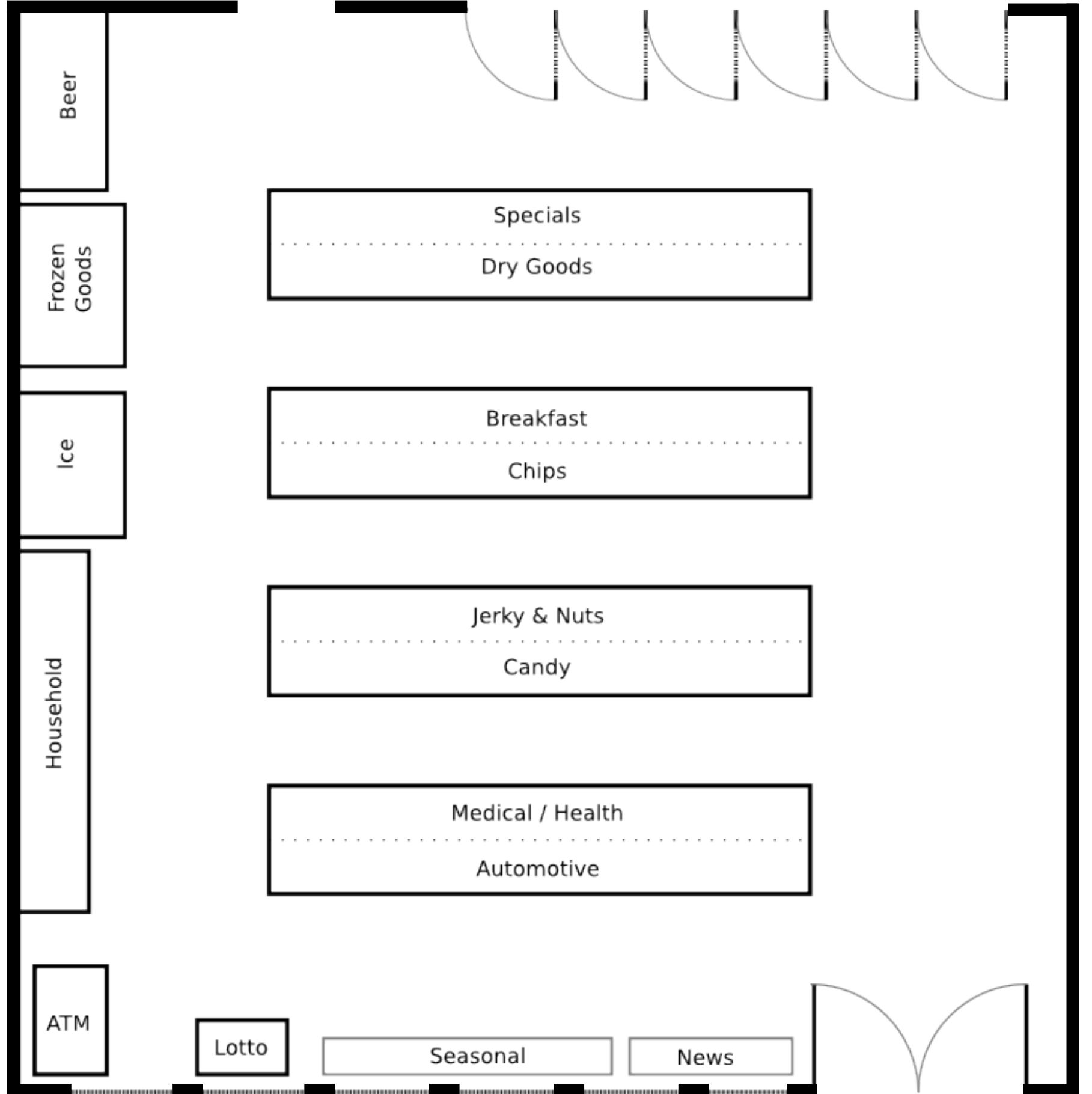


New Variants of the Floodlight Problem

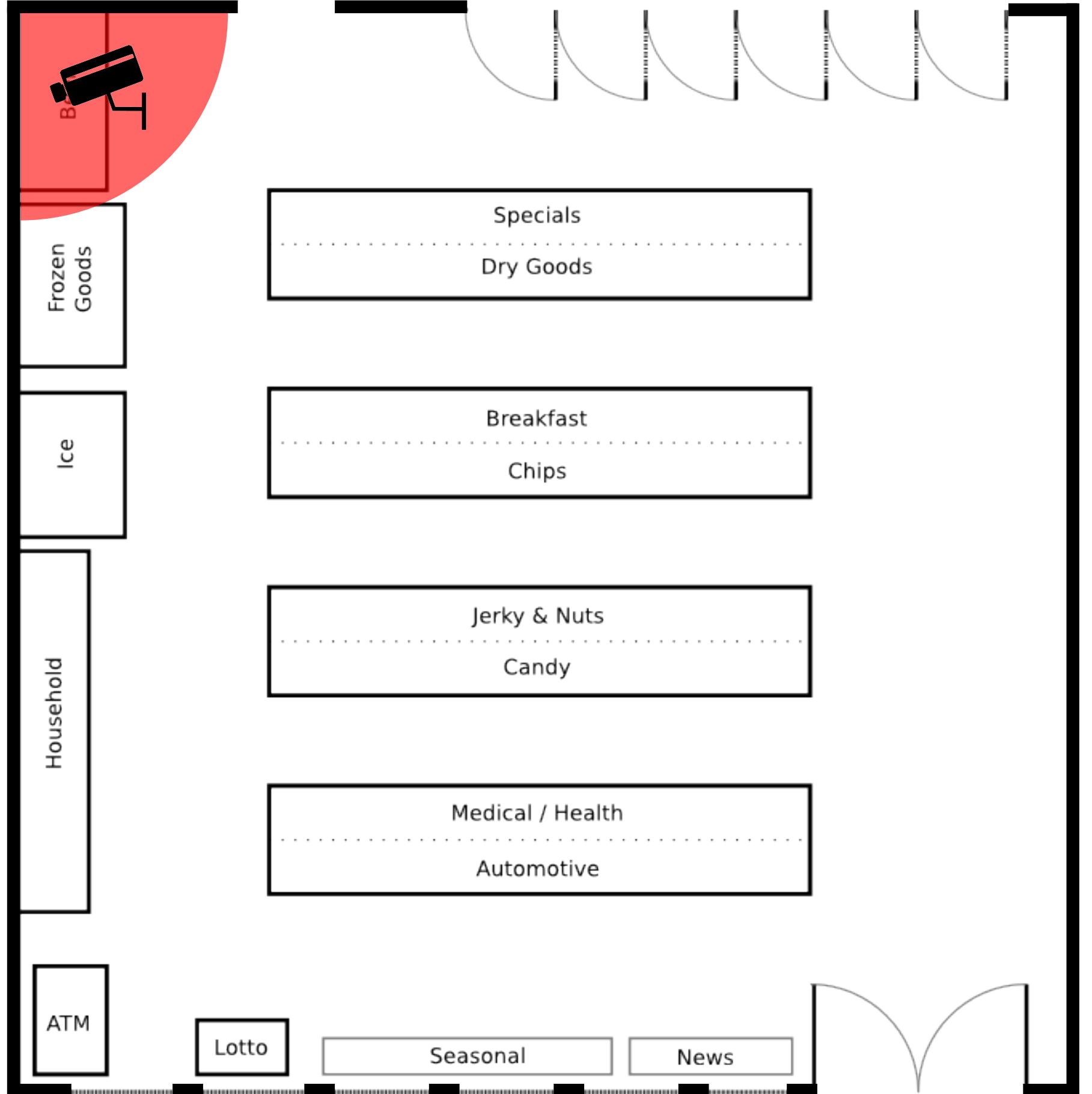
Bachelor's Thesis

Yannic Lieder, September 5, 2018

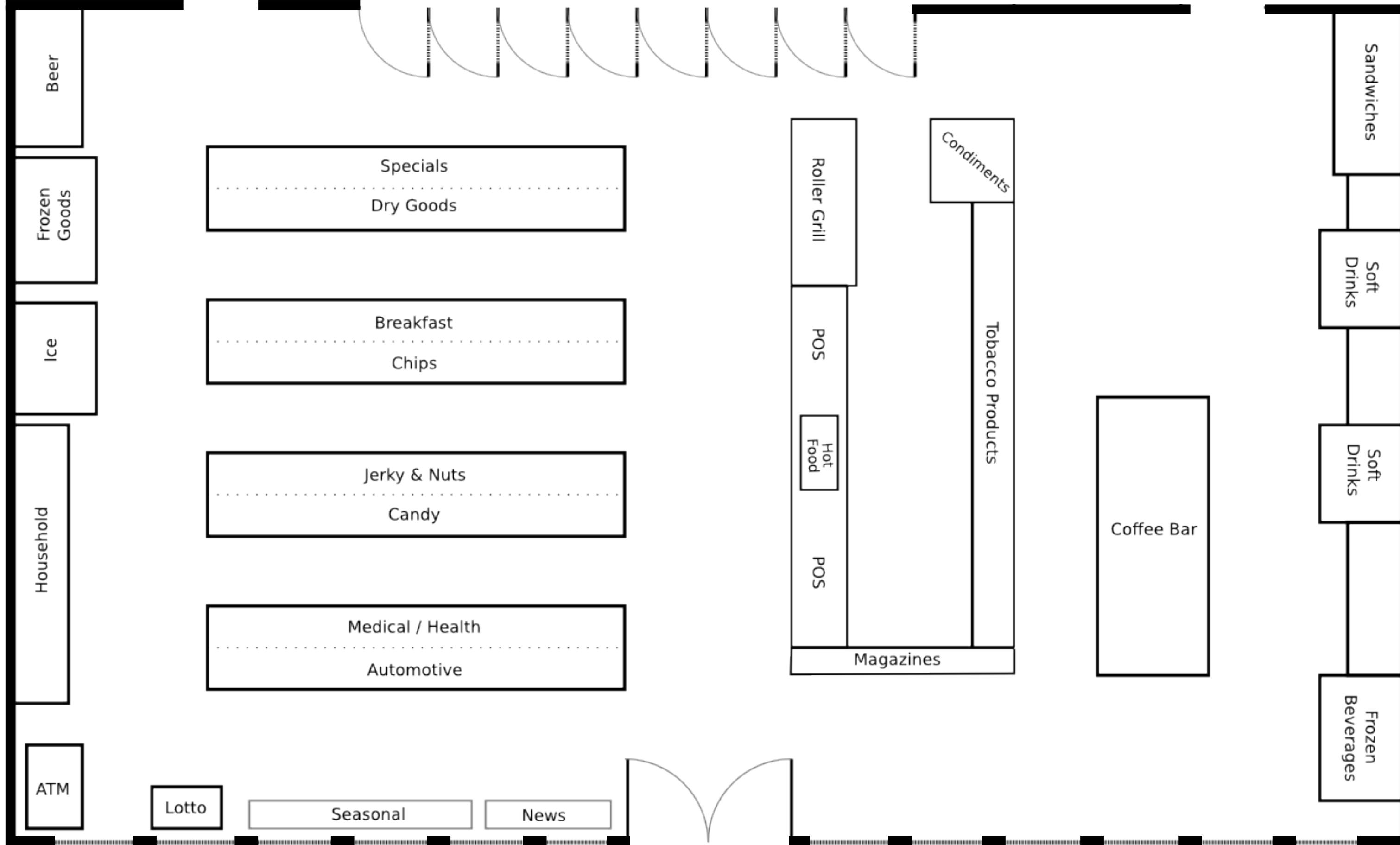
Introduction



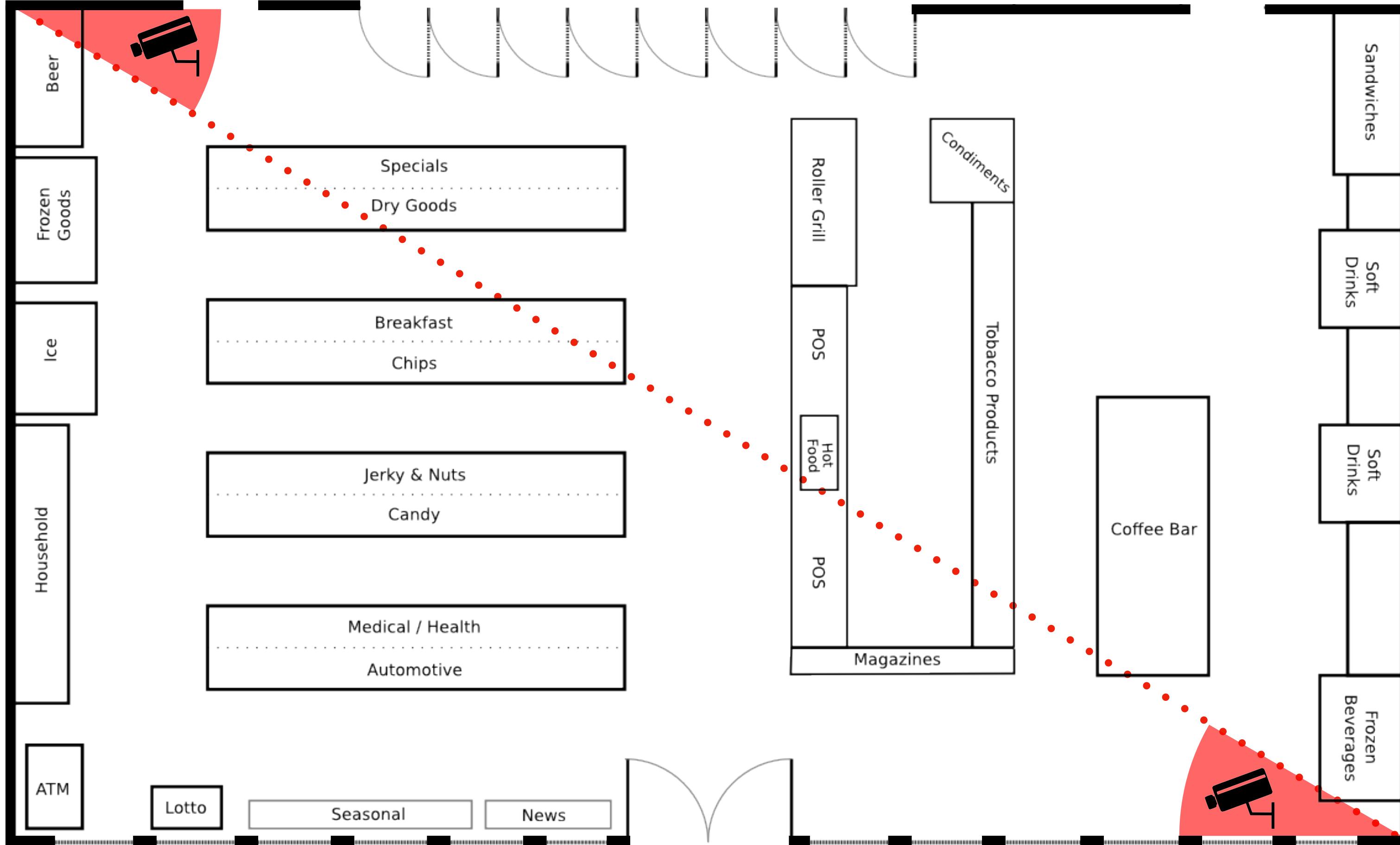
Introduction



Introduction



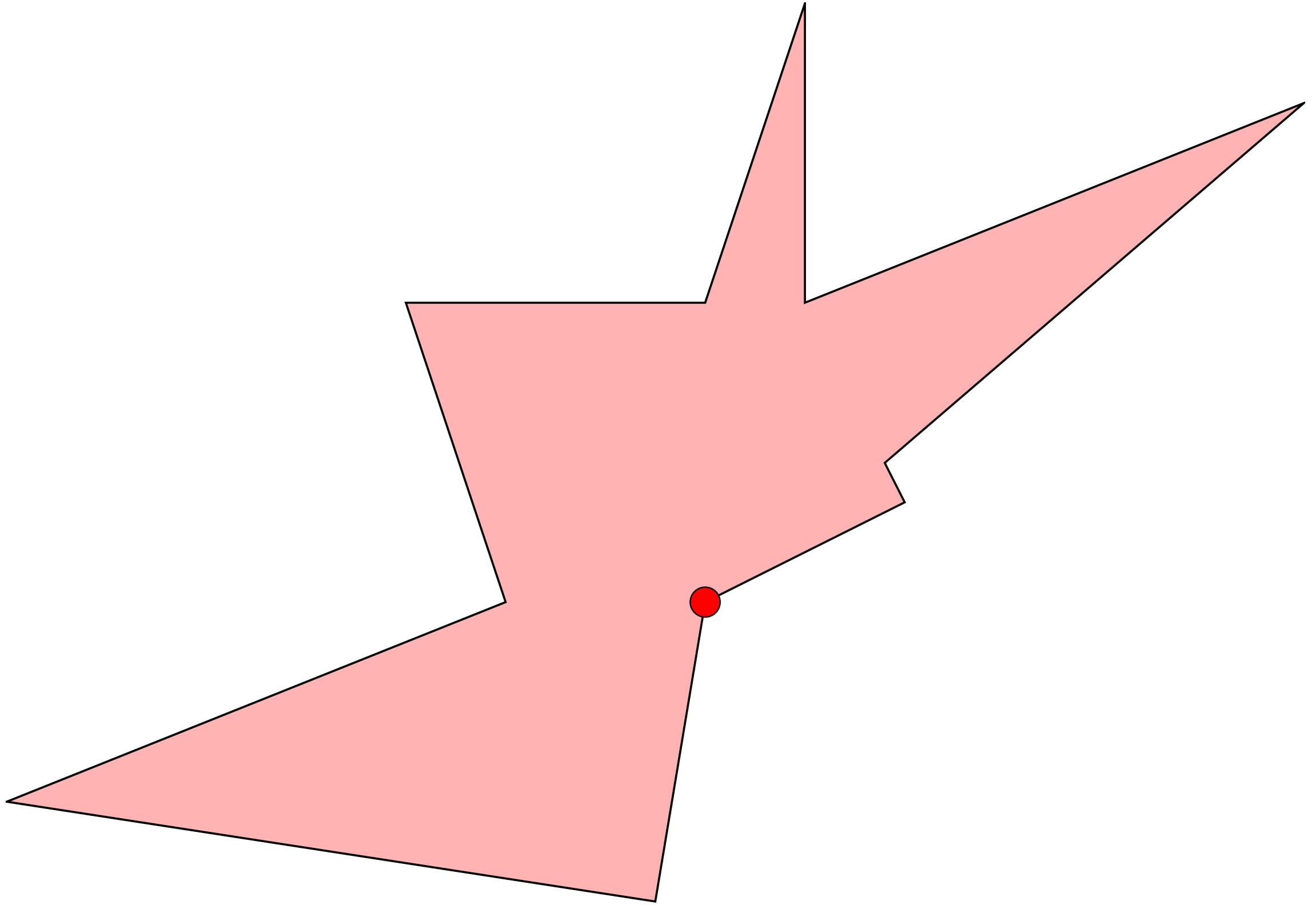
Introduction



Background

1973

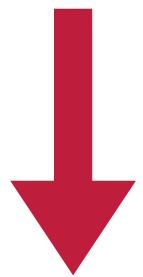
The Art Gallery Problem



Background

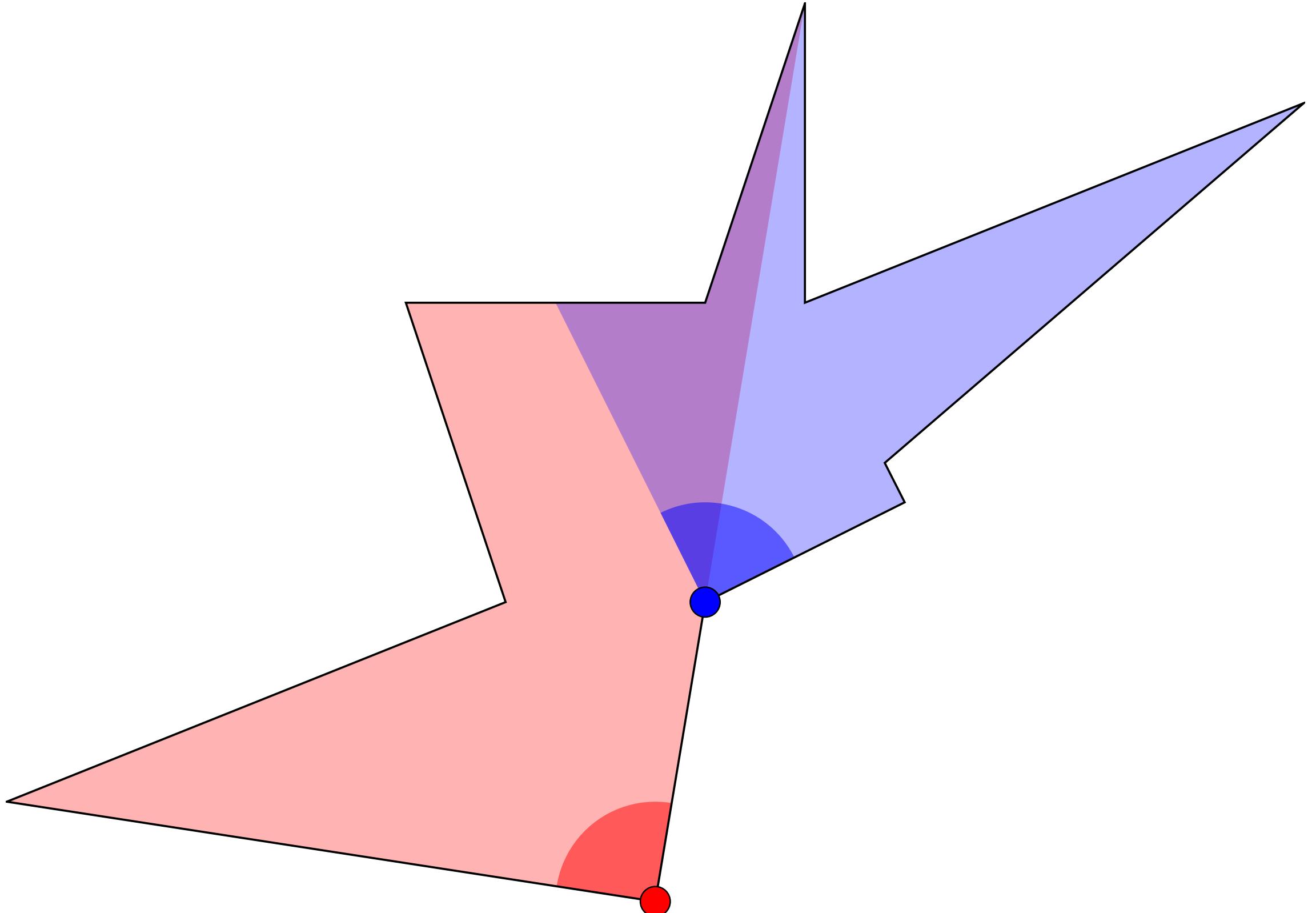
1973

The Art Gallery Problem



1992

The α - Floodlight Problem



Background

1973

The Art Gallery Problem



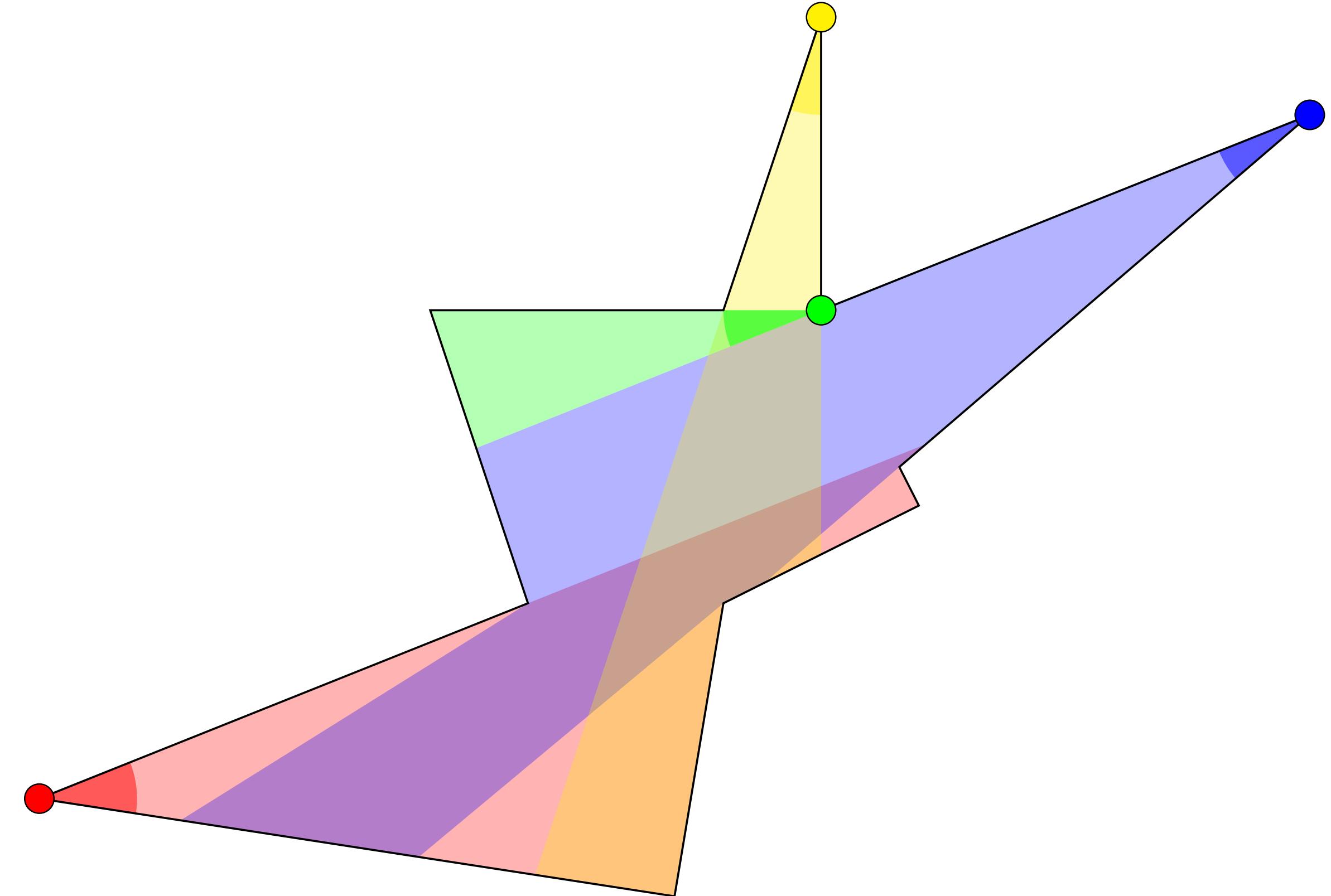
1992

The α - Floodlight Problem



This
thesis

The Angular Art Gallery Problem



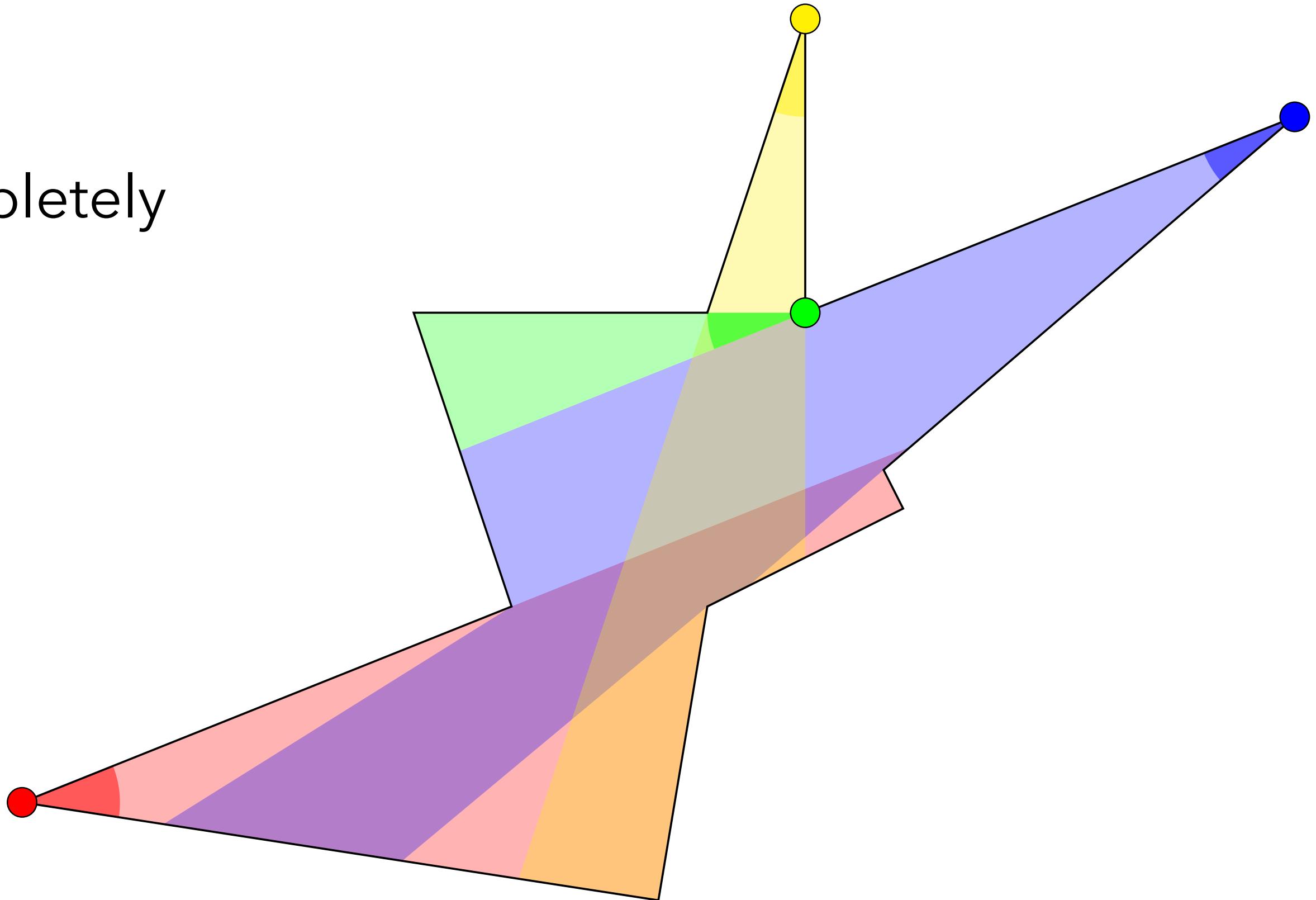
Problem Definition

Angular Art Gallery Problem

Instance: A simple polygon P

Wanted: A set of floodlights, covering P completely

Minimize: The total angle of all floodlights



Results

Lower Bound

Upper Bound

Results

	Lower Bound	Upper Bound
Equilateral Triangles	$\frac{\pi}{3}$	$\frac{\pi}{3}$

Results

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Equilateral Triangles	$\frac{\pi}{3}$	$\frac{\pi}{3}$
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Equilateral Triangles	$\frac{\pi}{3}$	$\frac{\pi}{3}$
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Simple Polygons	$(n - 1)\frac{\pi}{6} - \epsilon$	$(n - 2)\frac{\pi}{4}$

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Duality to independent circle packing

Outline

Introduction

Equilateral Triangles

Histograms

Simple Polygons

Duality

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Introduction

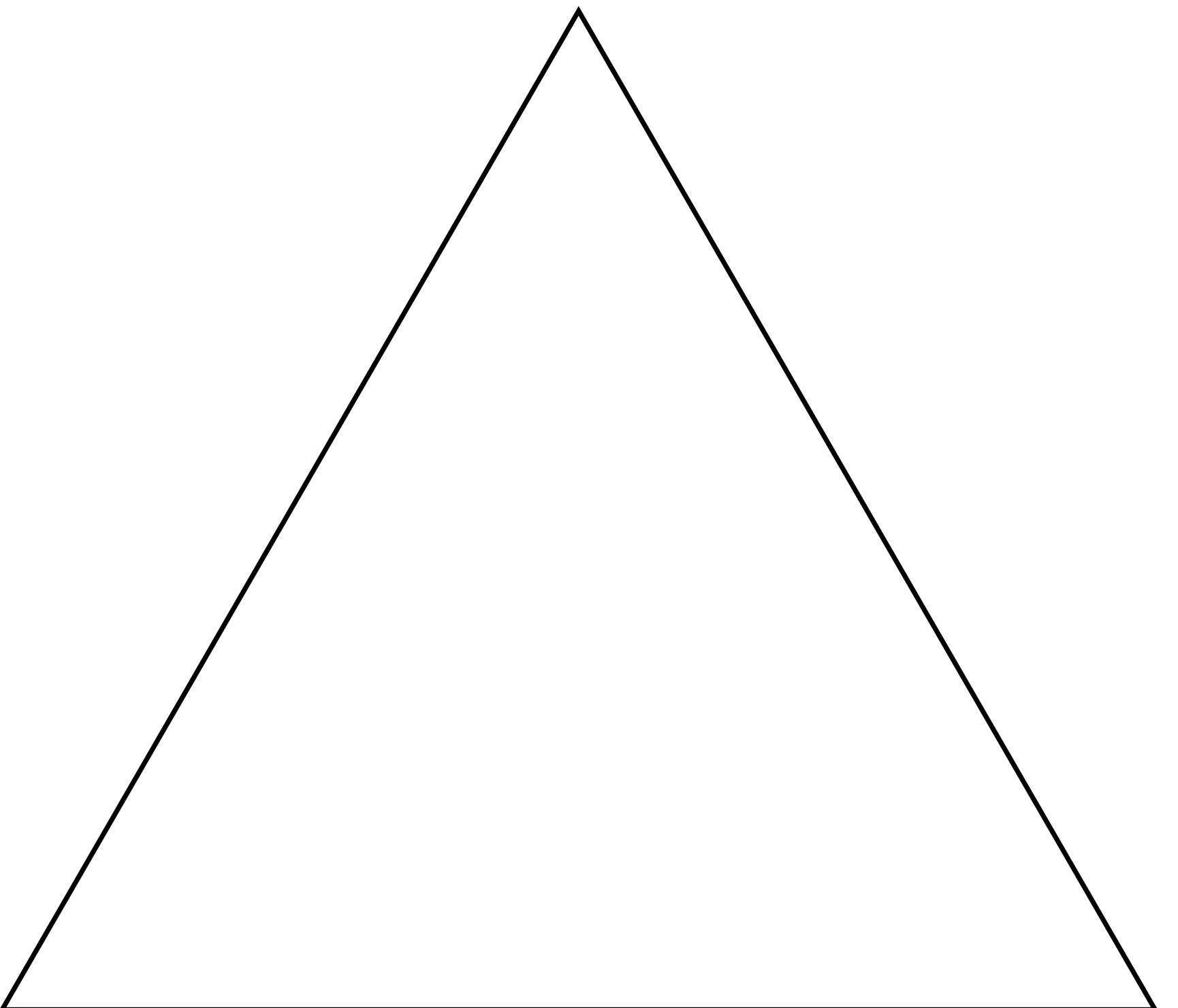
Optimal Covering of **Equilateral Triangles**

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Introduction

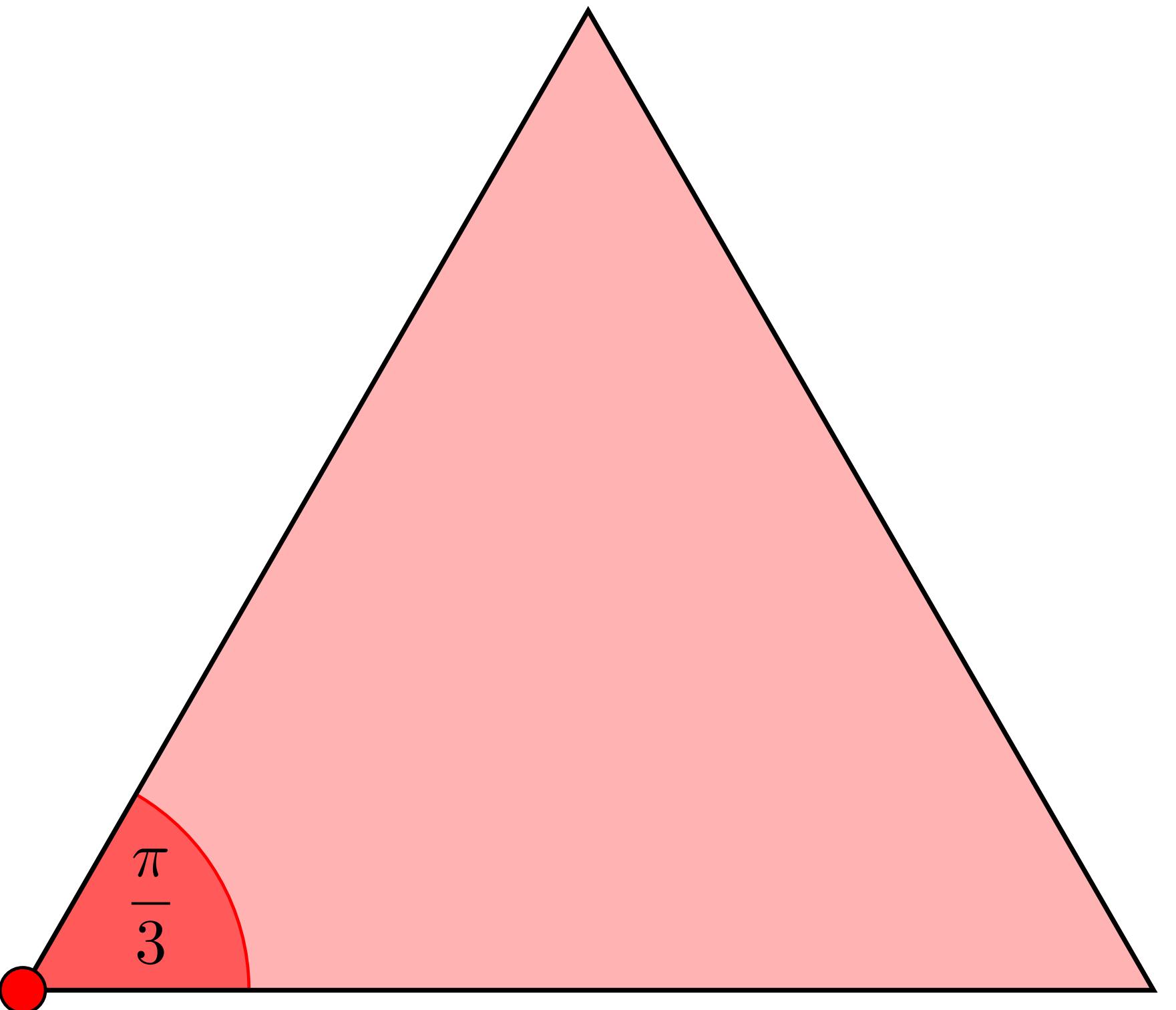
Optimal Covering of **Equilateral Triangles**

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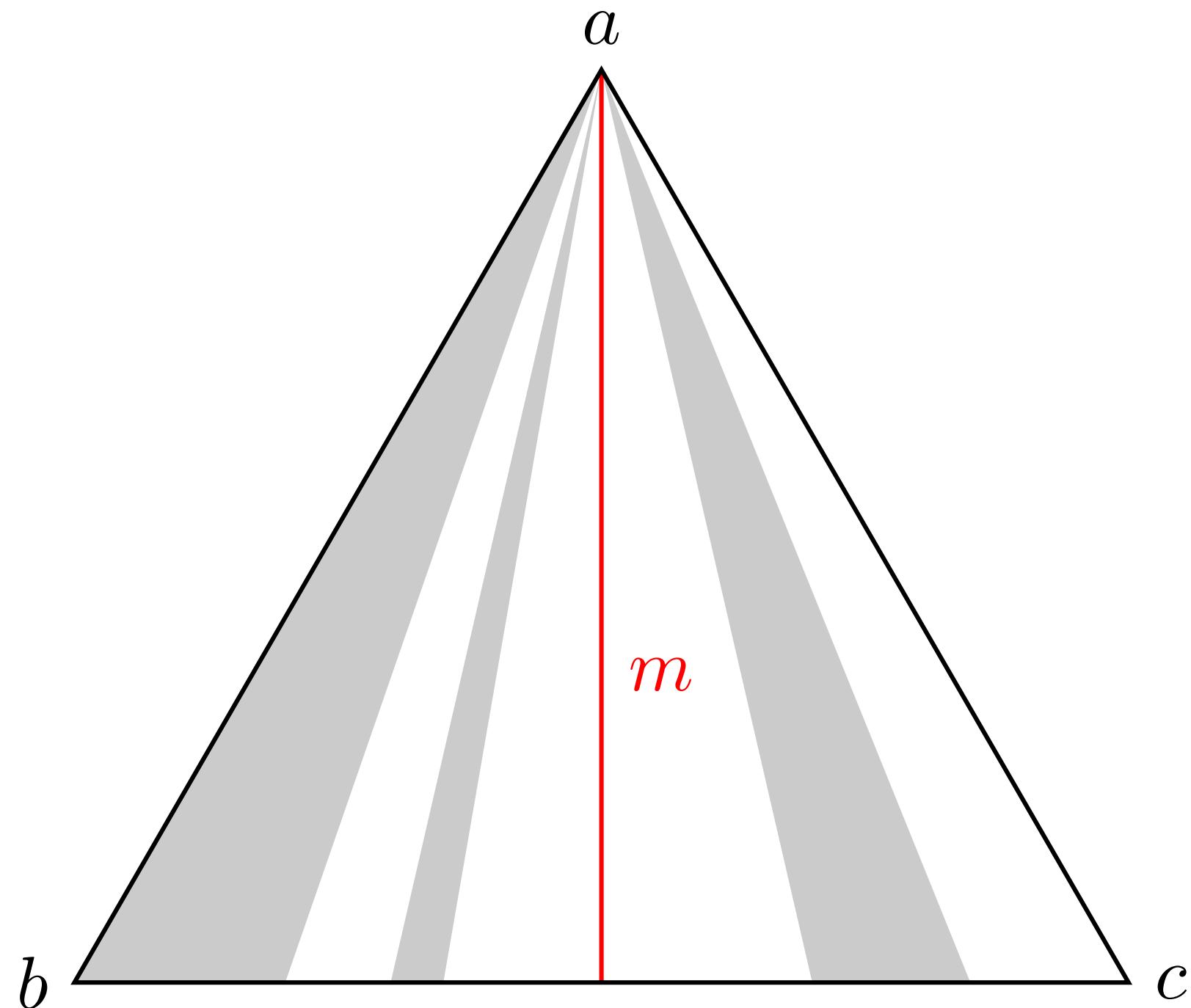
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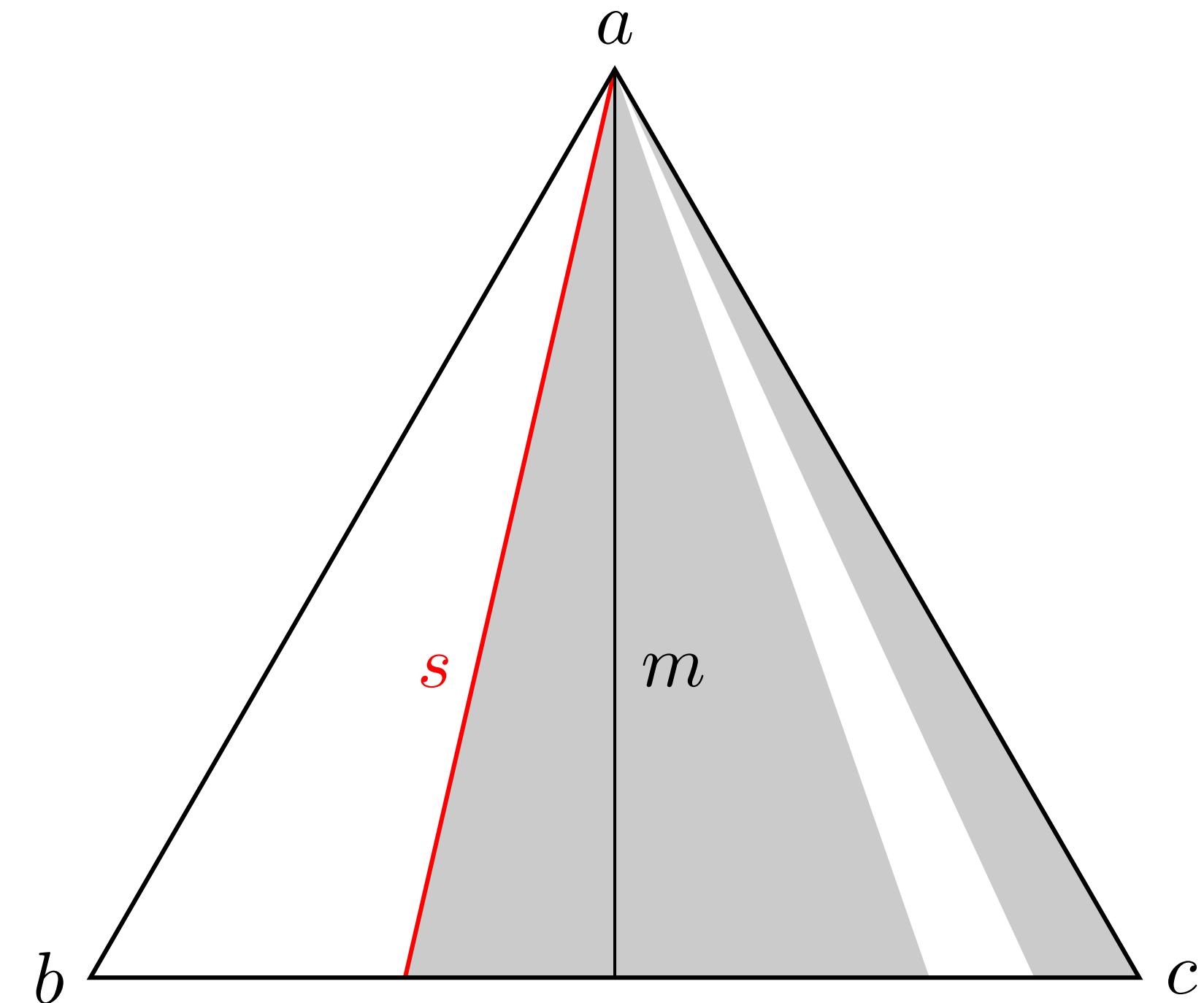


Lower Bound for Equilateral Triangles

Case 1: m is not covered by a

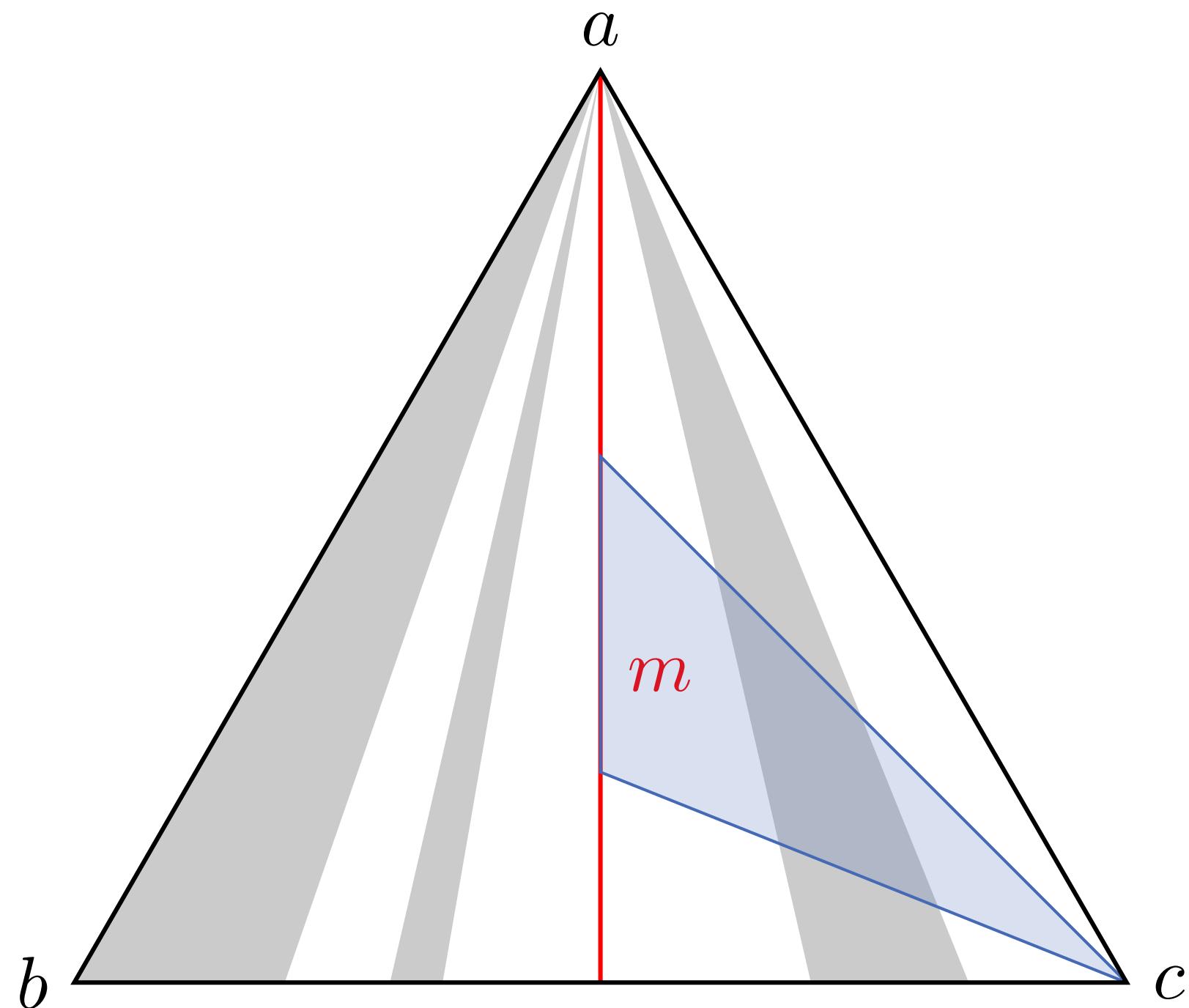


Case 2: m is covered by a

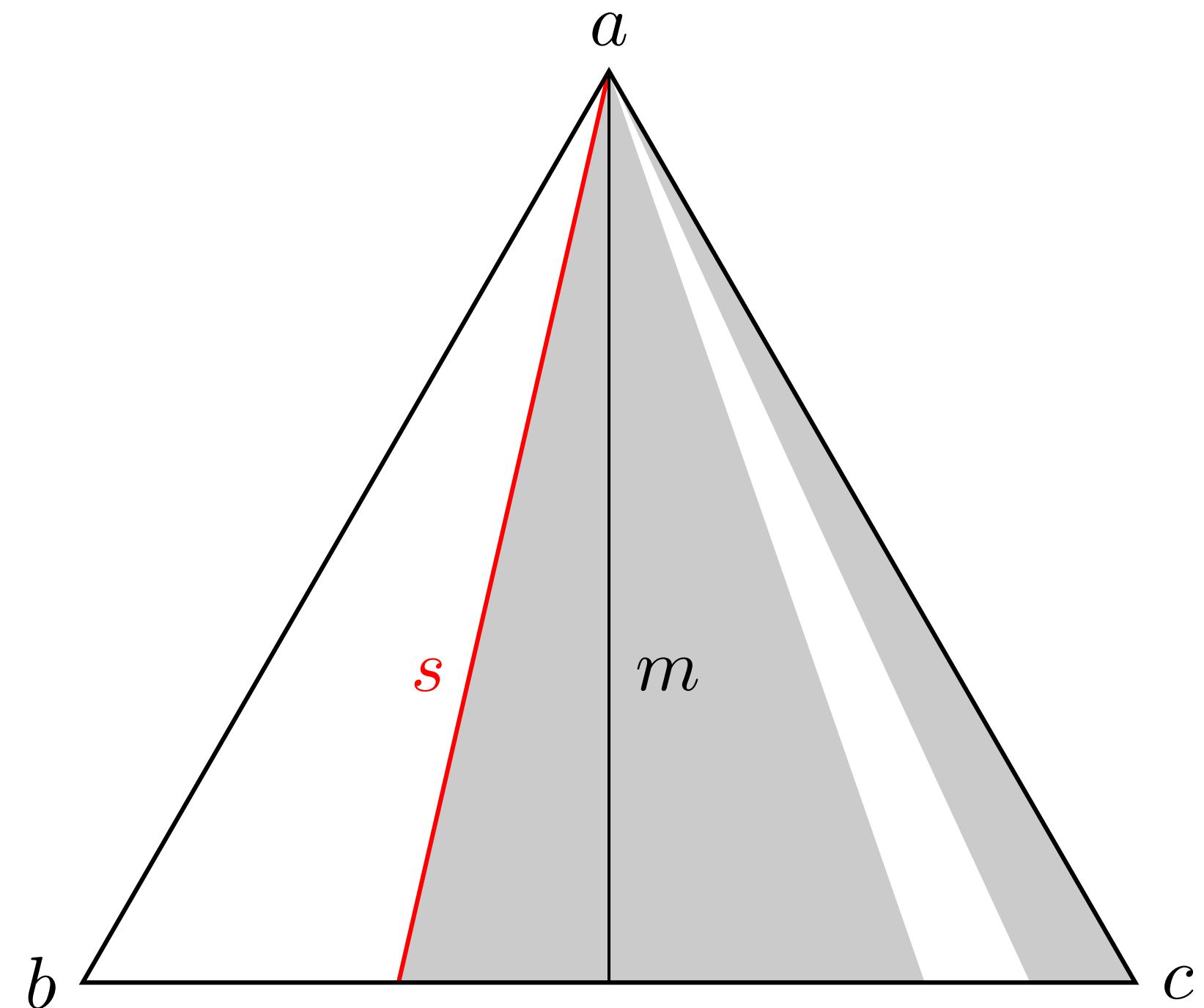


Lower Bound for Equilateral Triangles

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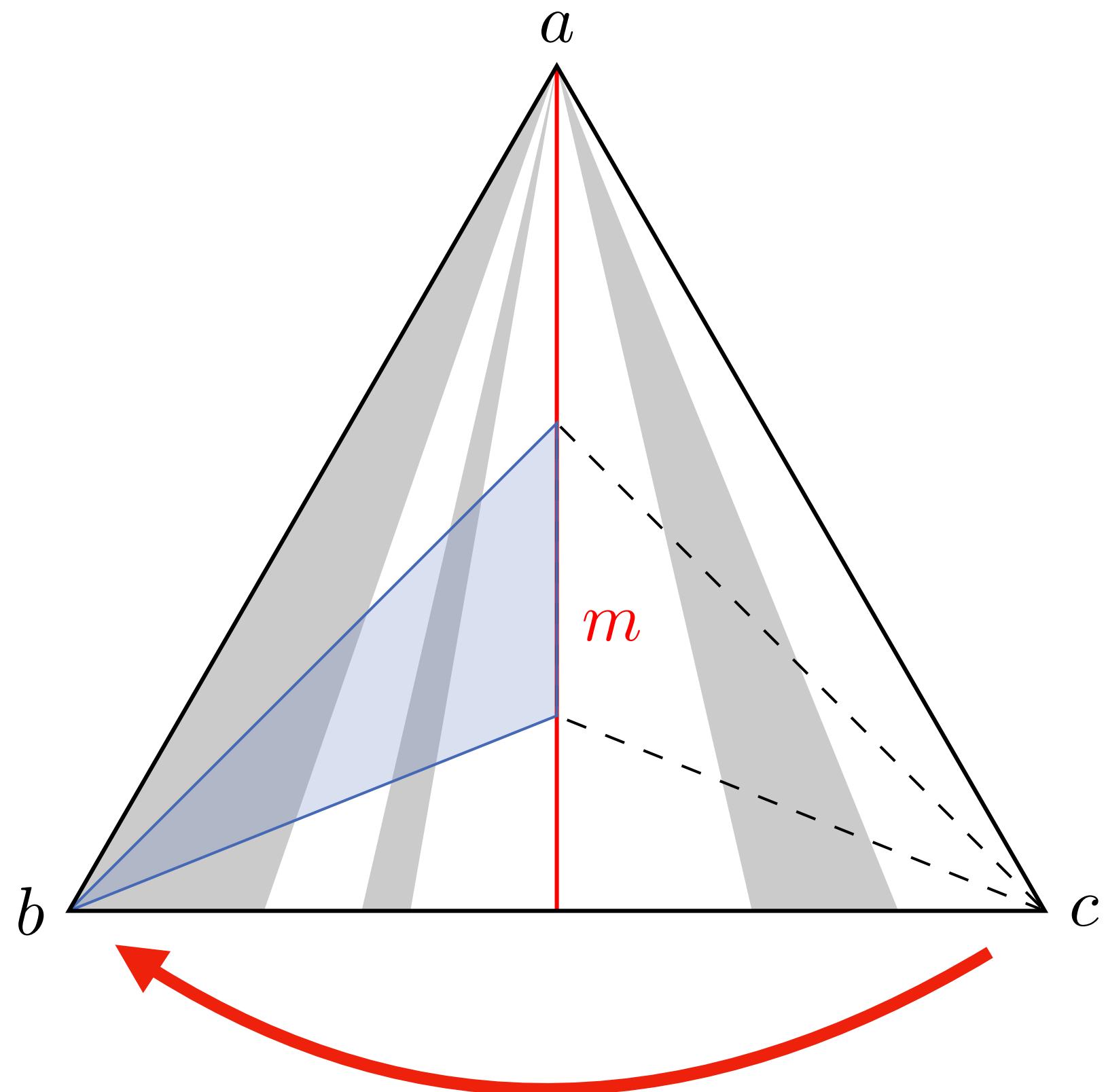


Case 2: m is covered by a

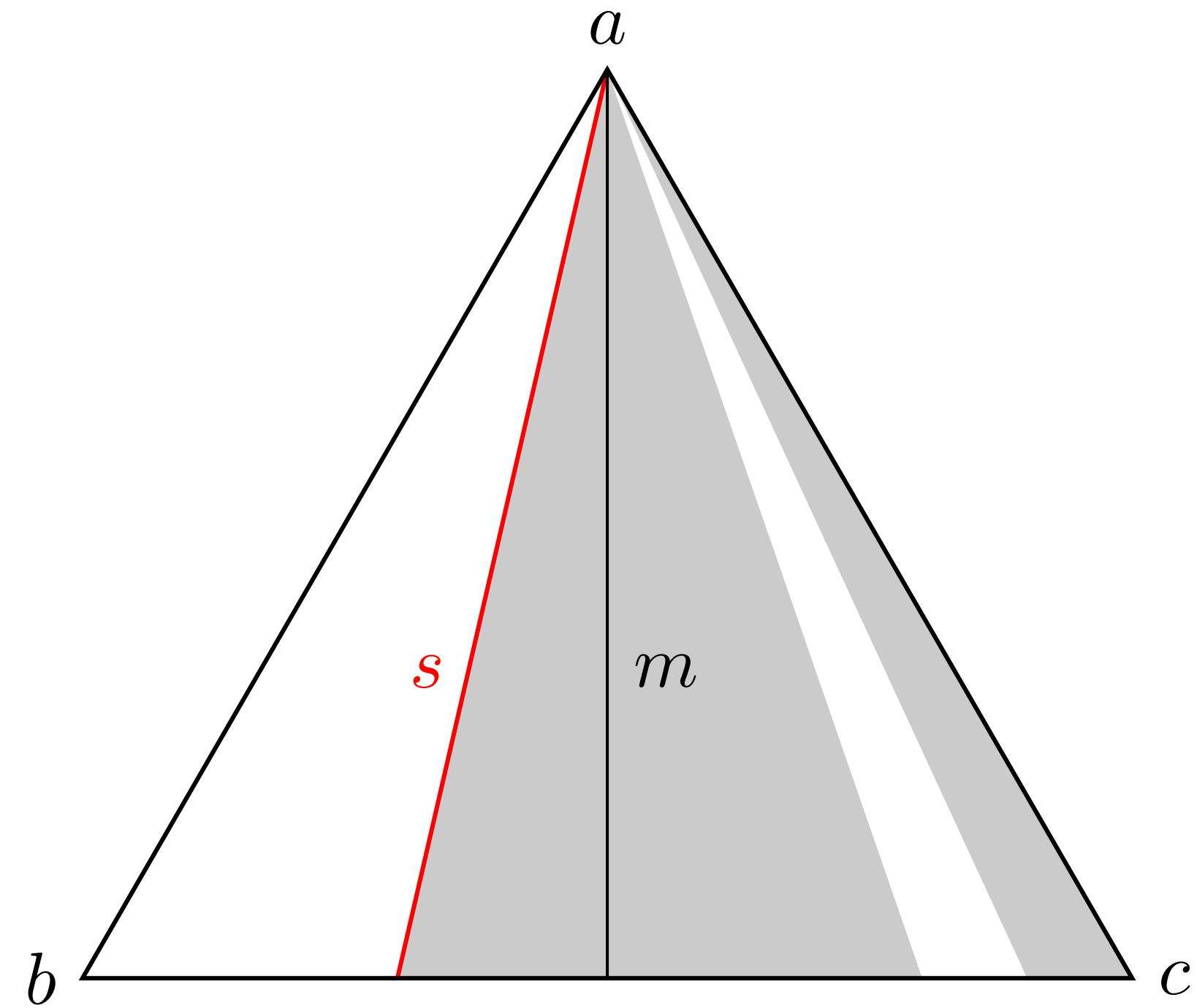


Lower Bound for Equilateral Triangles

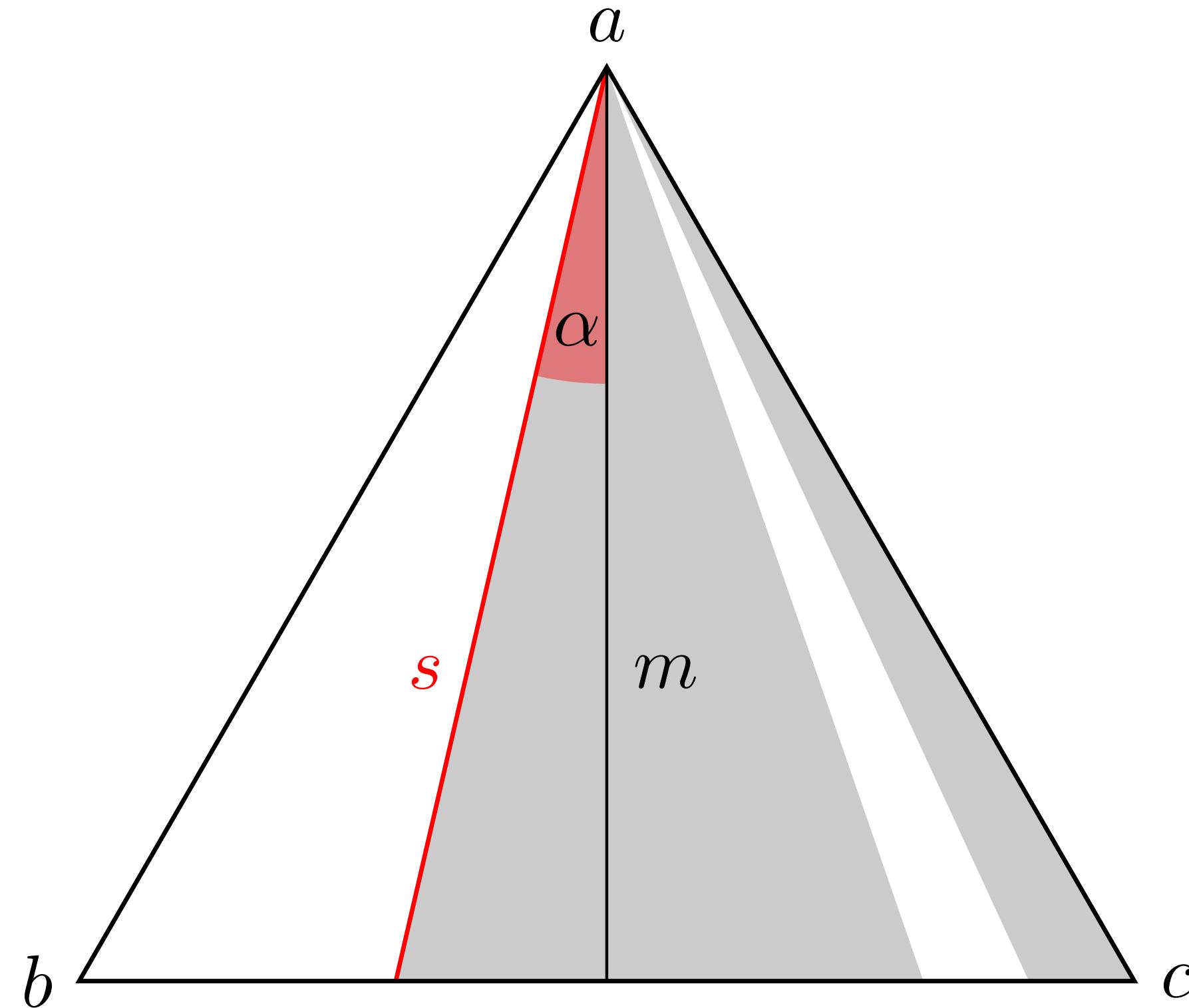
Case 1: m is not covered by a



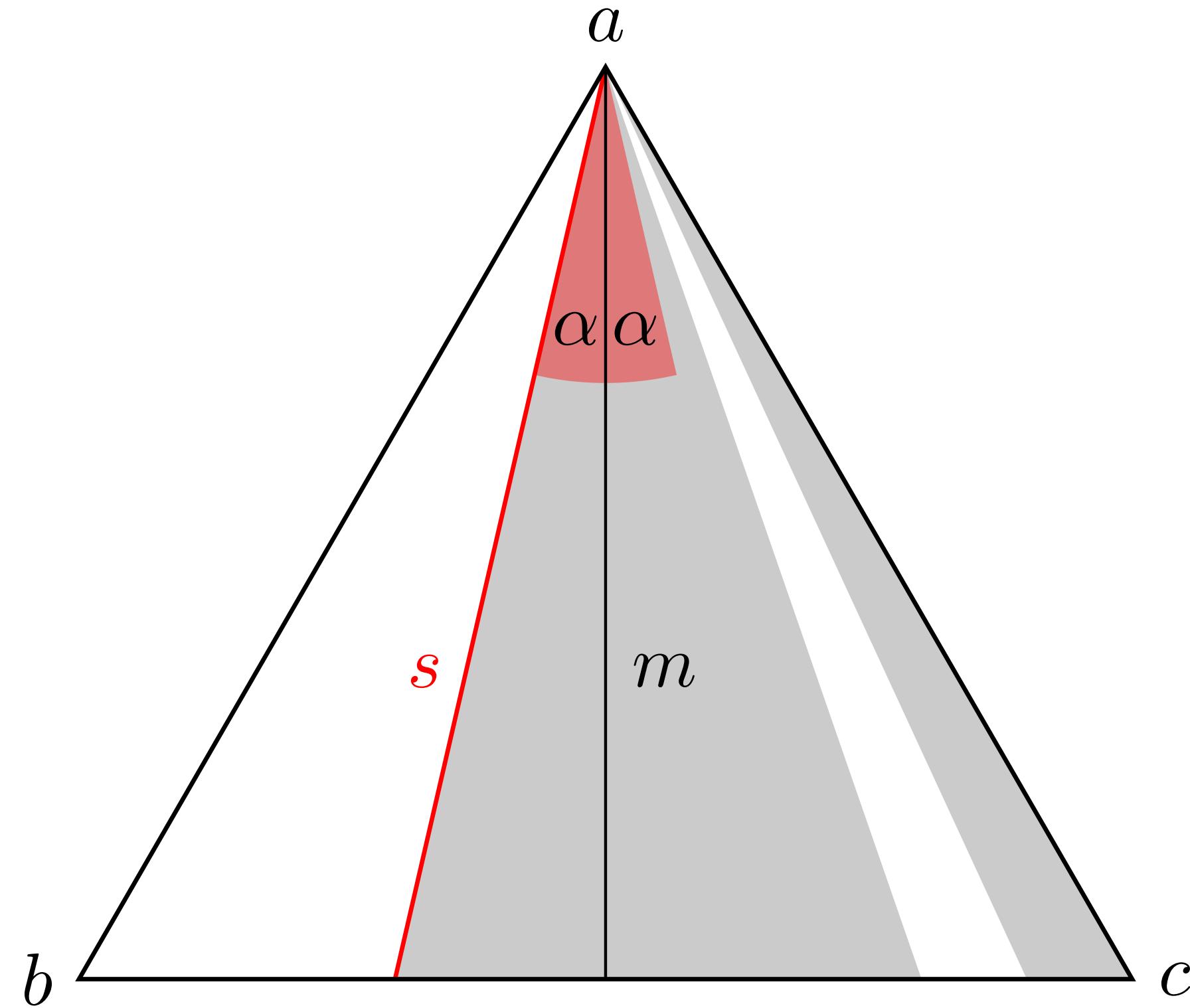
Case 2: m is covered by a



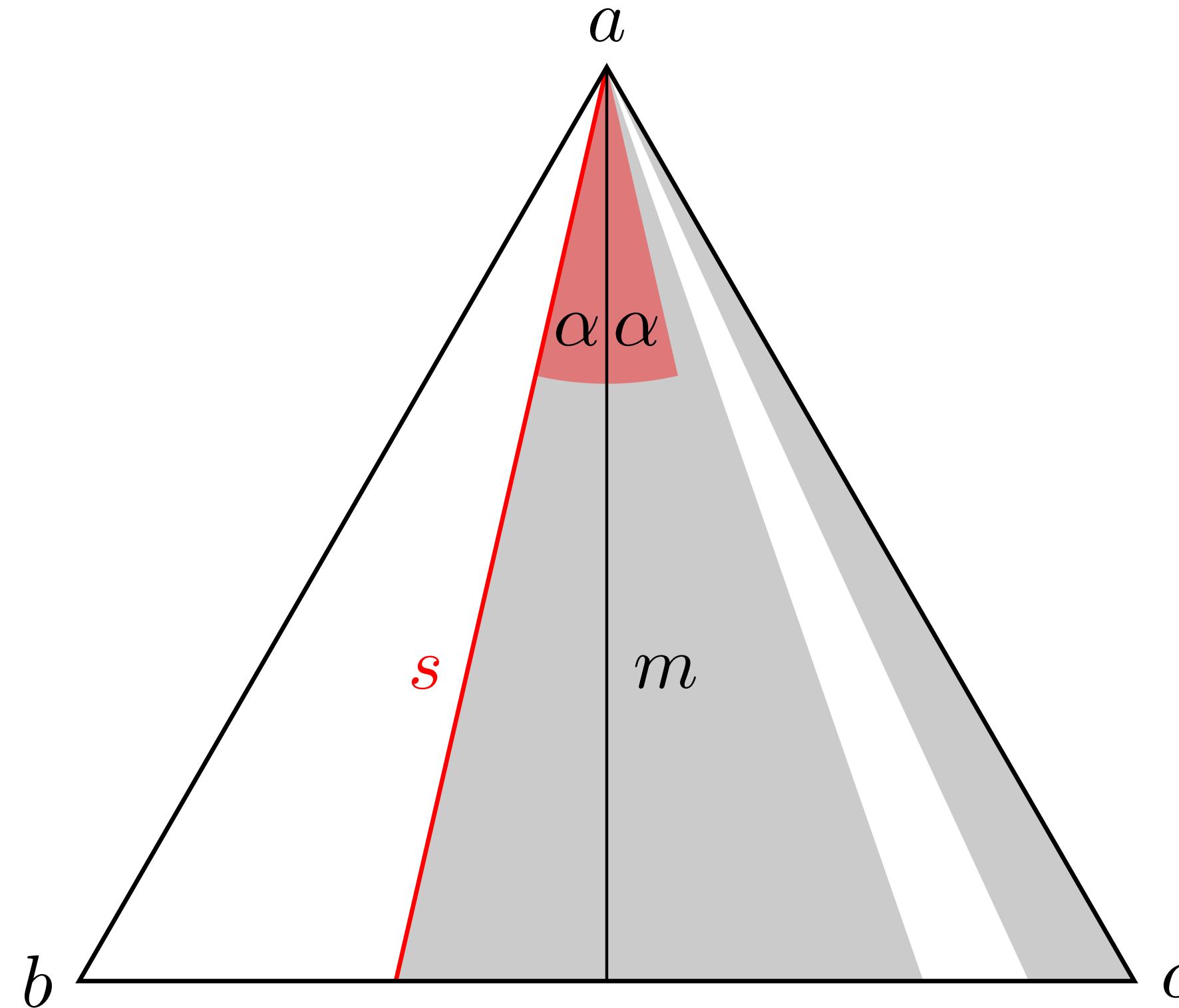
Lower Bound for Equilateral Triangles, case 2



Lower Bound for Equilateral Triangles, case 2



Lower Bound for Equilateral Triangles, case 2

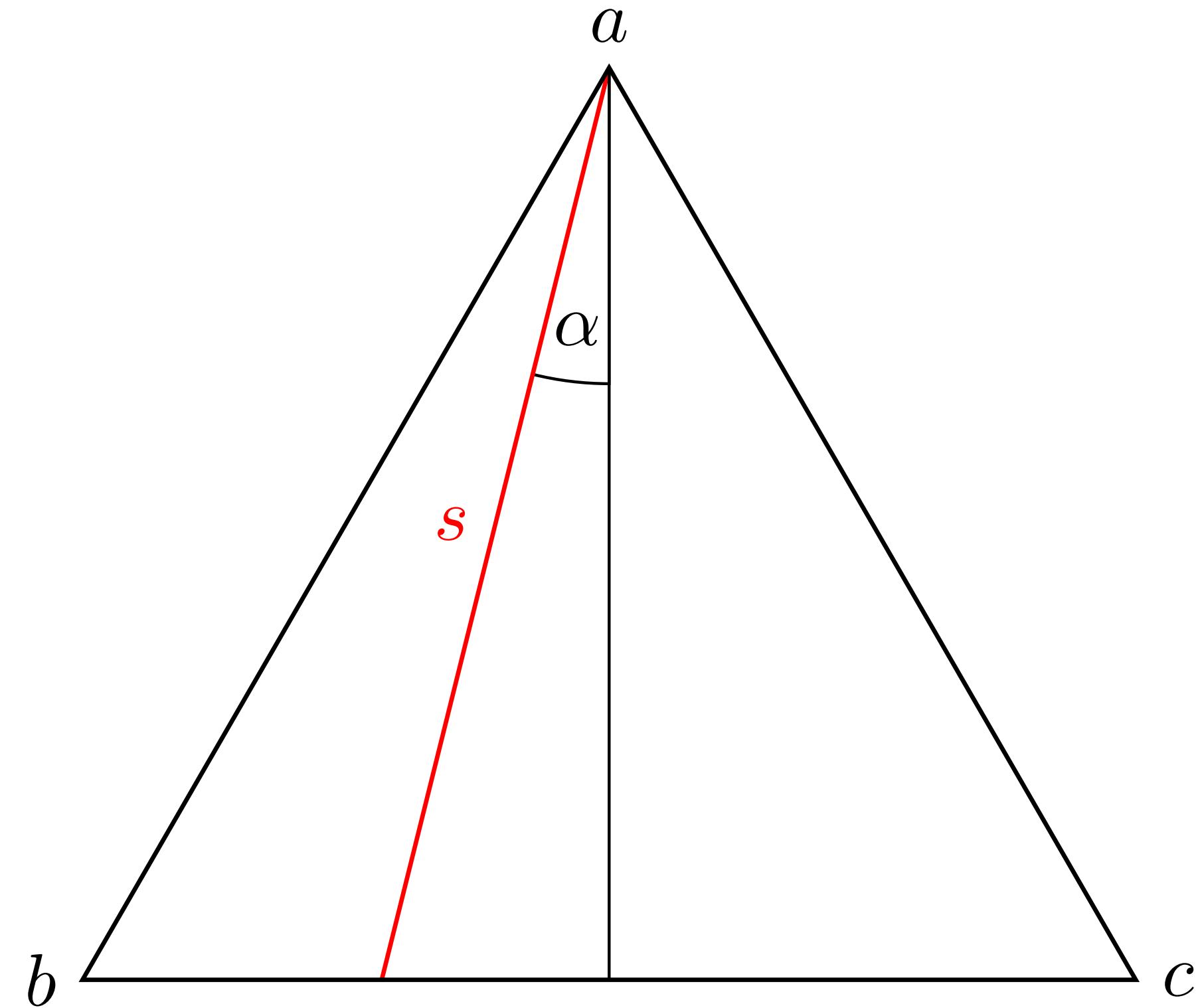


Idea: Showing that an angle of at least

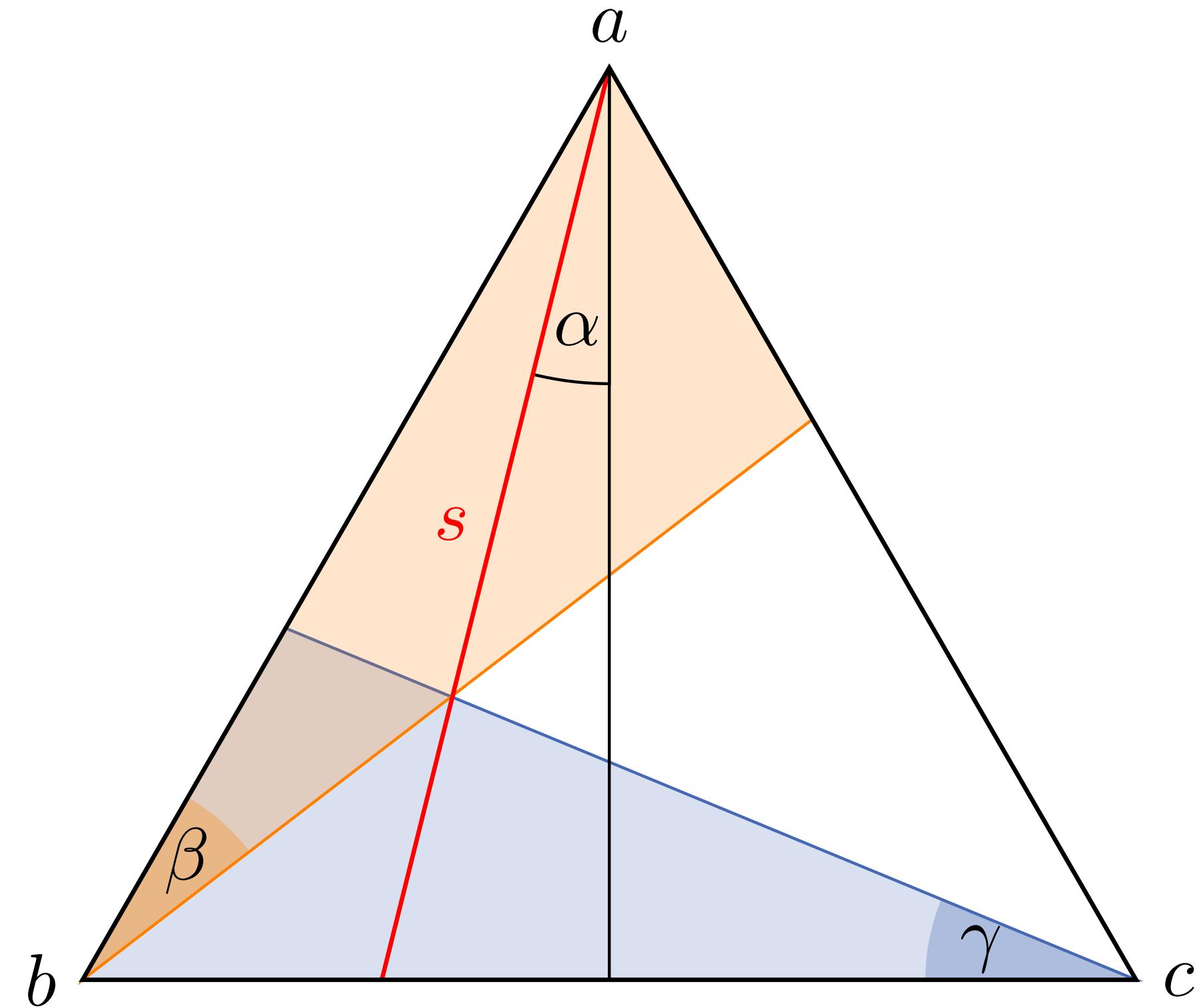
$$\frac{\pi}{3} - 2\alpha$$

is required to cover s .

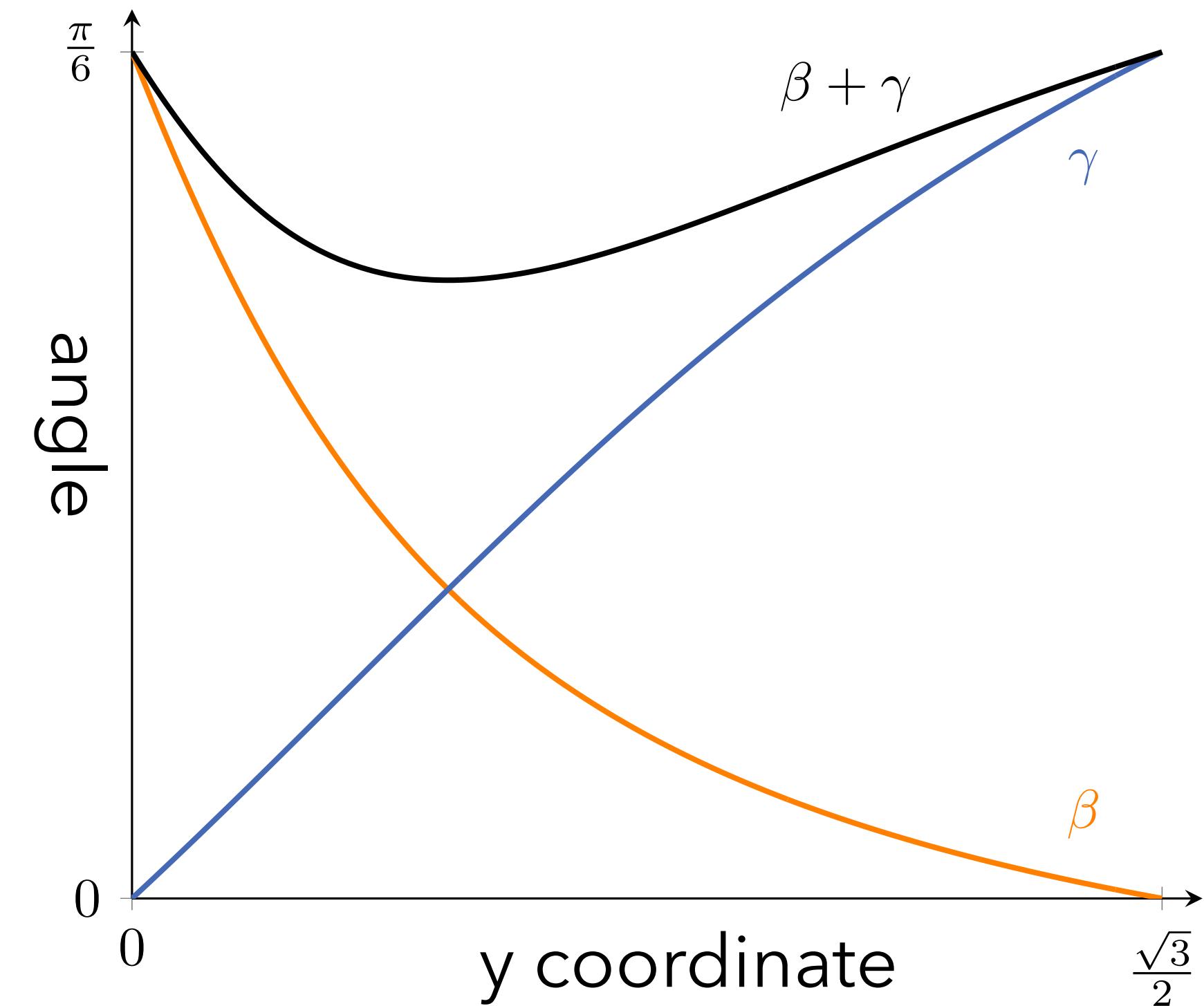
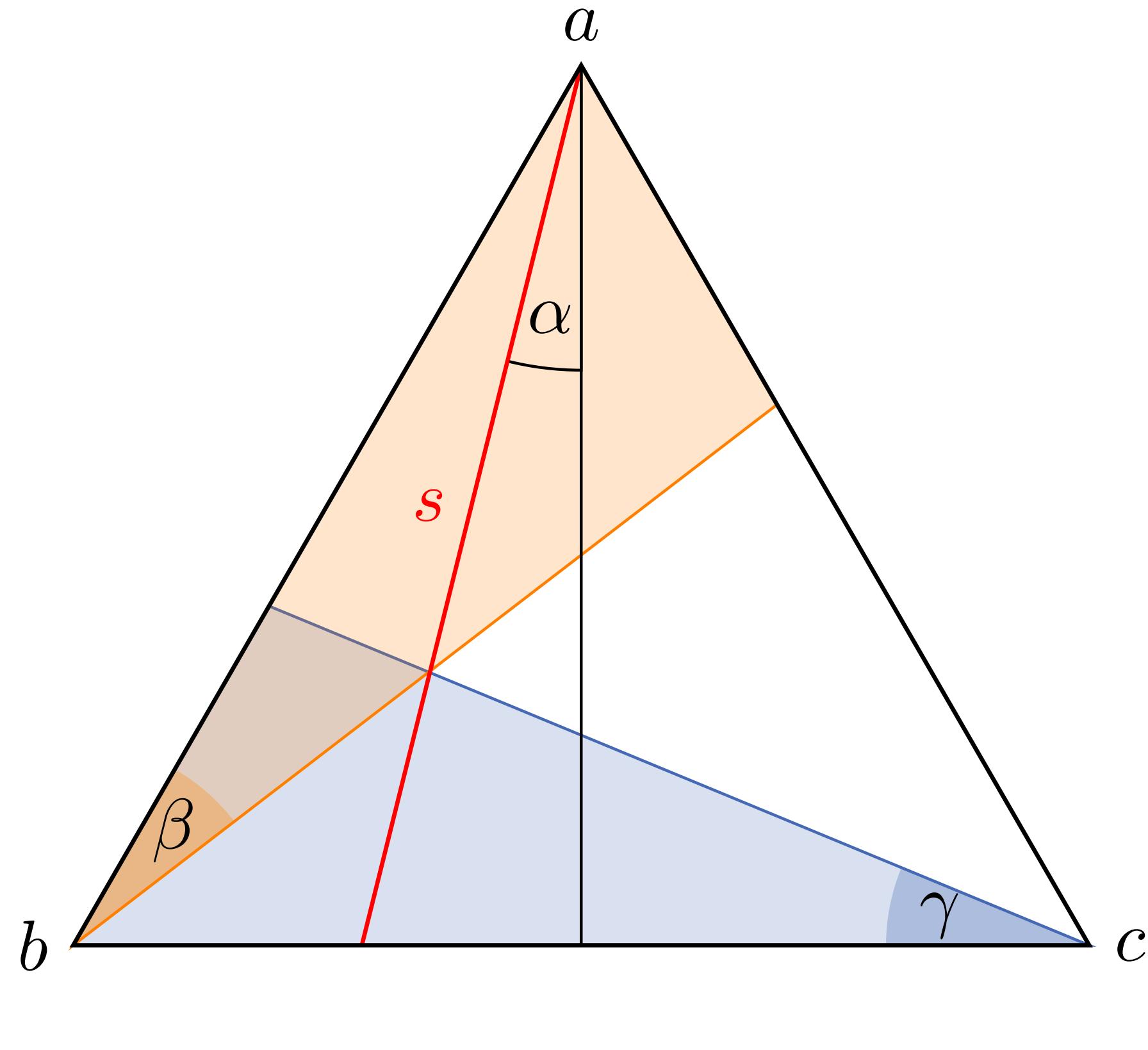
Minimum Covering of s



Minimum Covering of s



Minimum Covering of s



Minimum Covering of s

$$(\beta + \gamma)(y) = \arctan \left(\frac{\frac{1}{2} - \tan(\alpha) \left(\frac{\sqrt{3}}{2} - y \right)}{y} \right) - \frac{\pi}{6} + \arctan \left(\frac{y}{\frac{1}{2} + \tan(\alpha) \left(\frac{\sqrt{3}}{2} - y \right)} \right)$$

Minimum Covering of s

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$$y_{\min} = -\frac{1 - 3 \tan^2(\alpha)}{2\sqrt{3} (\tan^2(\alpha) + 1)} + \sqrt{\left(\frac{1 - 3 \tan^2(\alpha)}{2\sqrt{3} \tan^2(\alpha) + 2} \right)^2 - \frac{3 \tan^2(\alpha) - 1}{4 (\tan^2(\alpha) + 1)}}$$

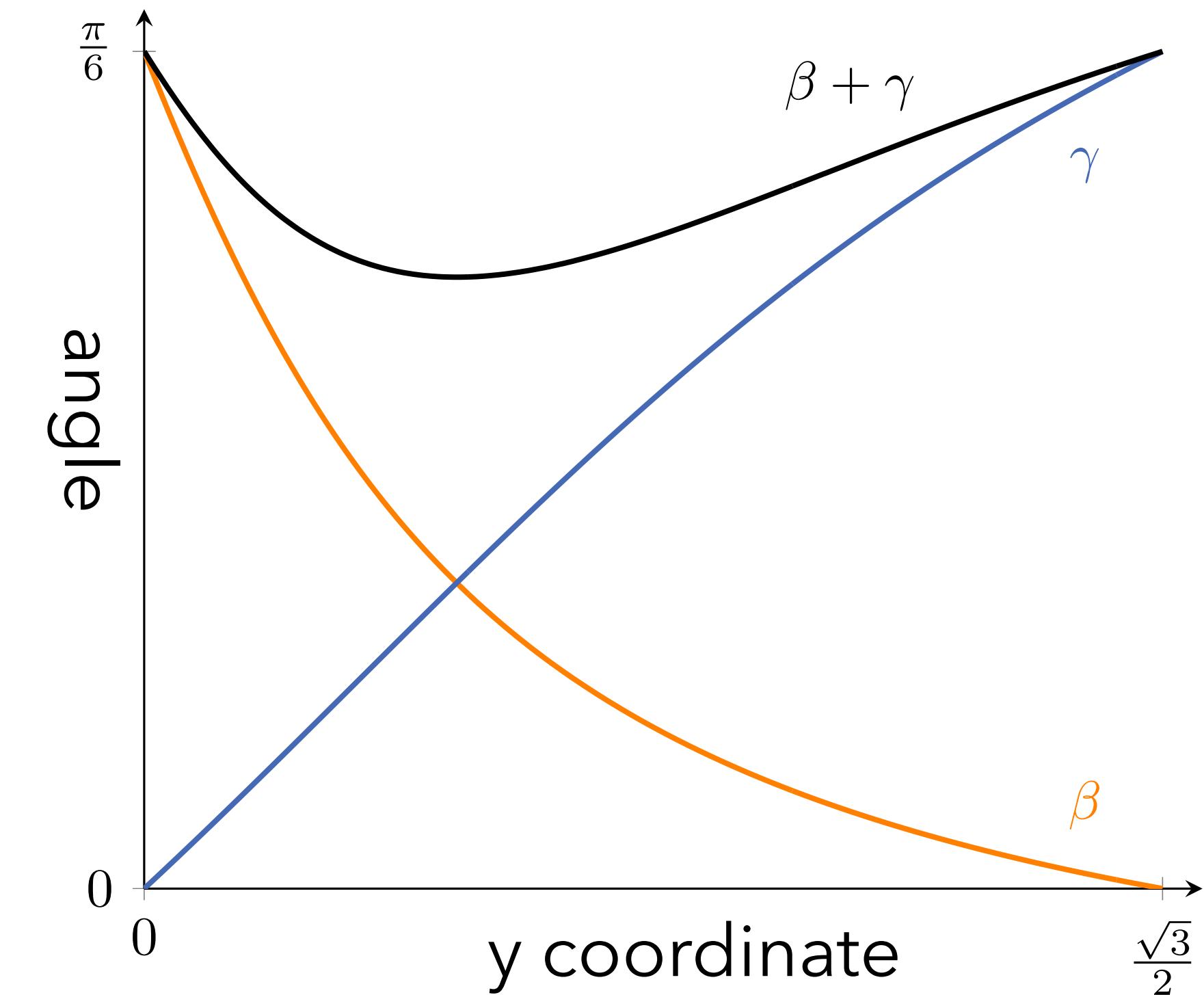
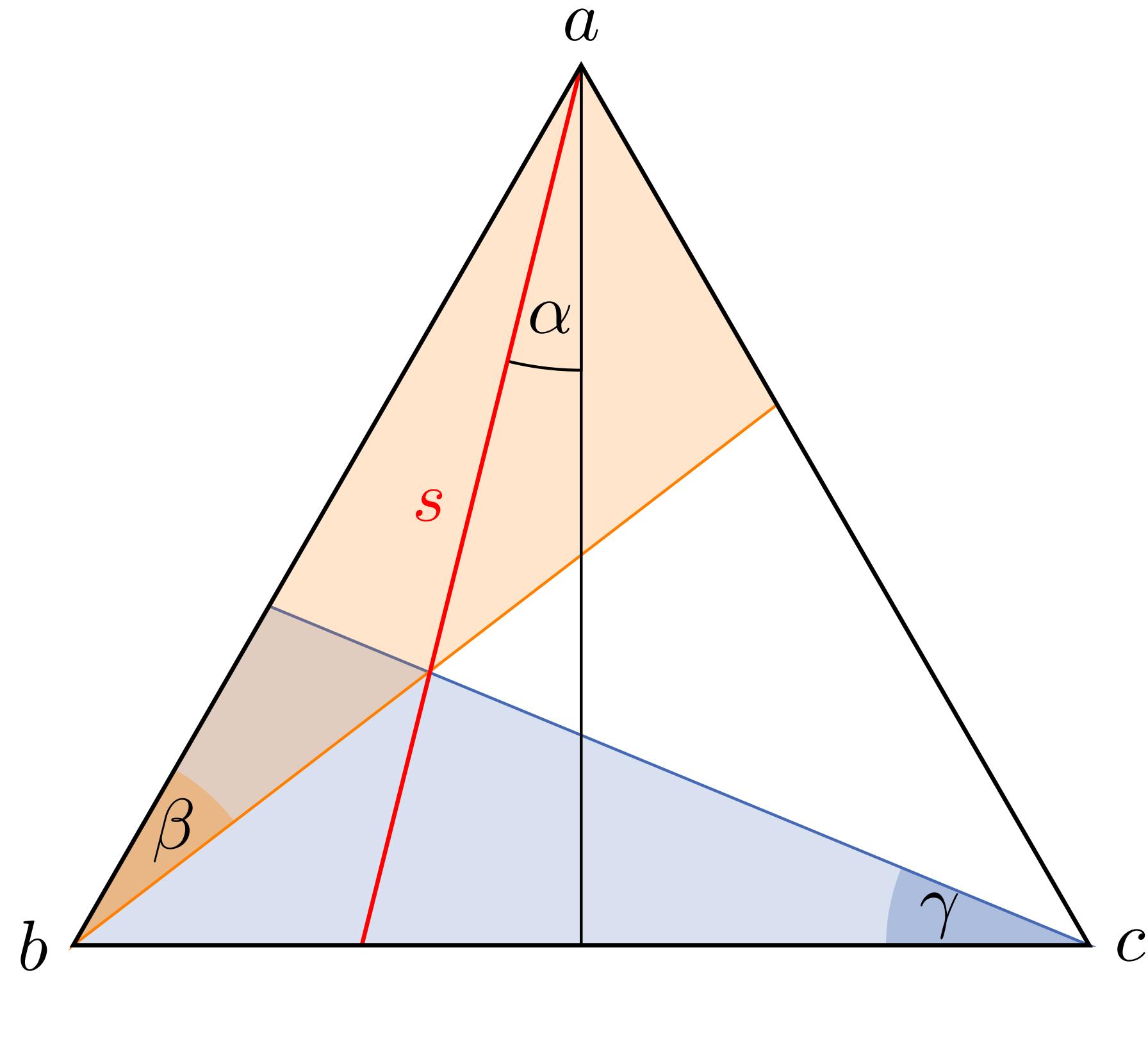
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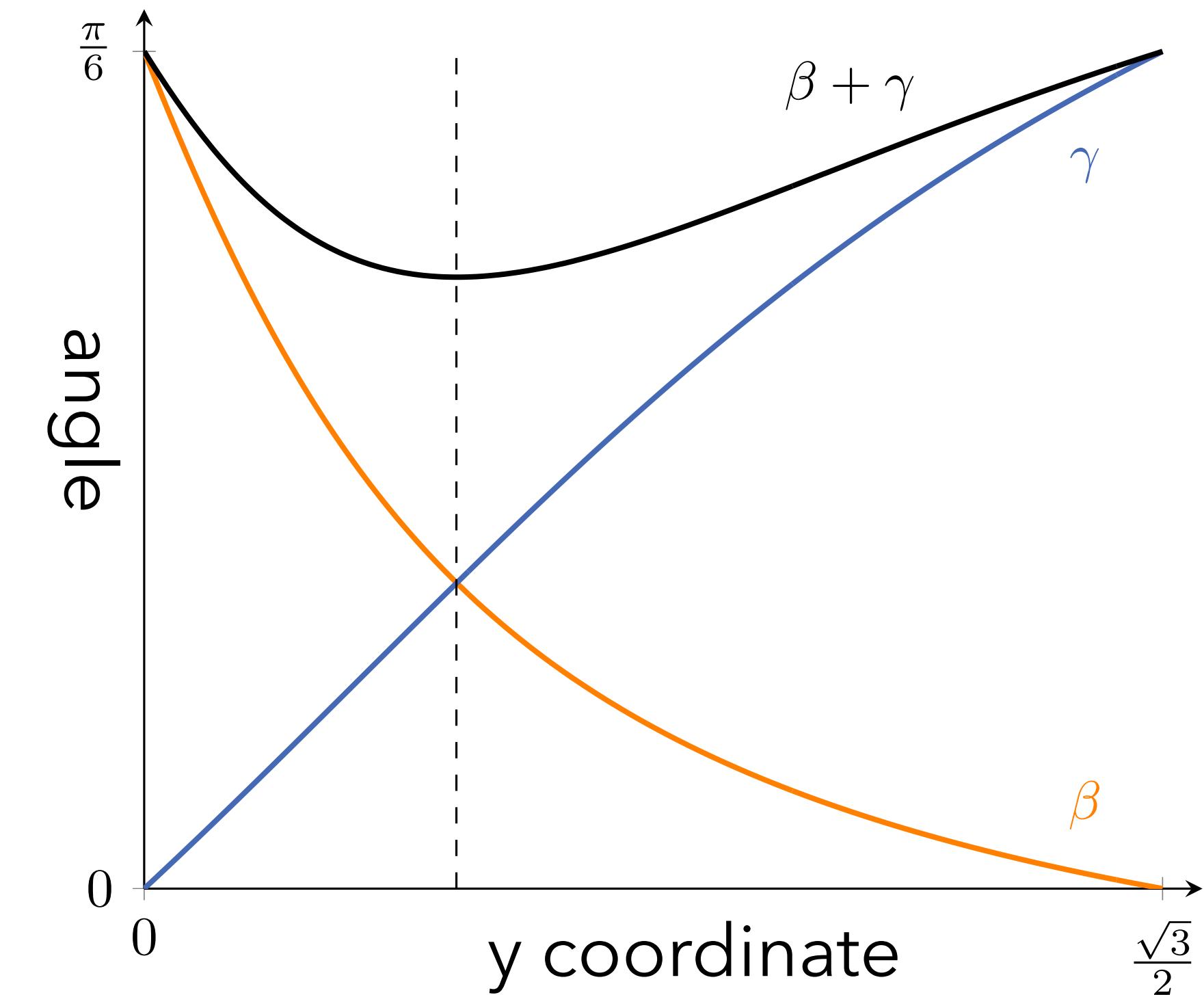
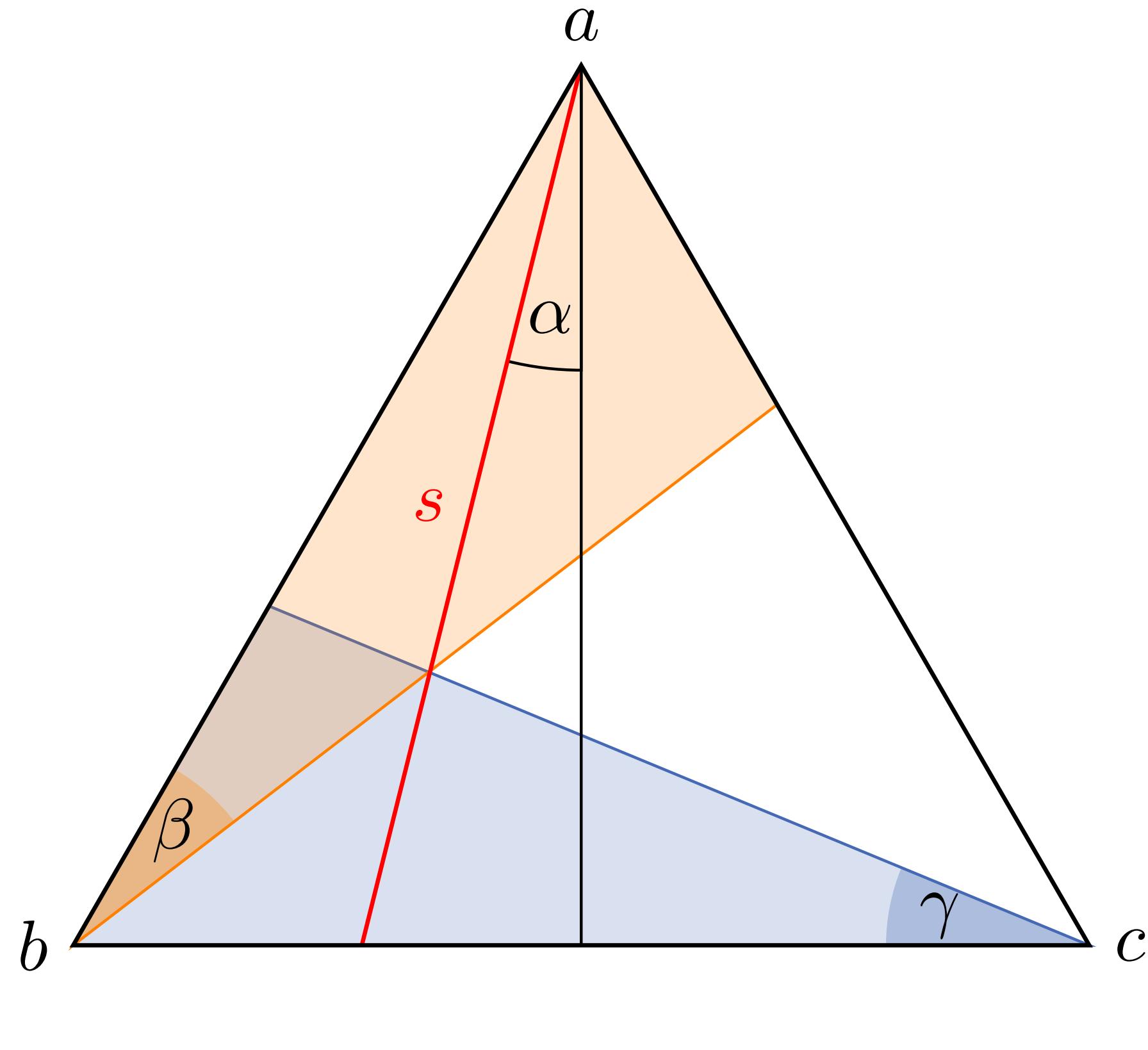
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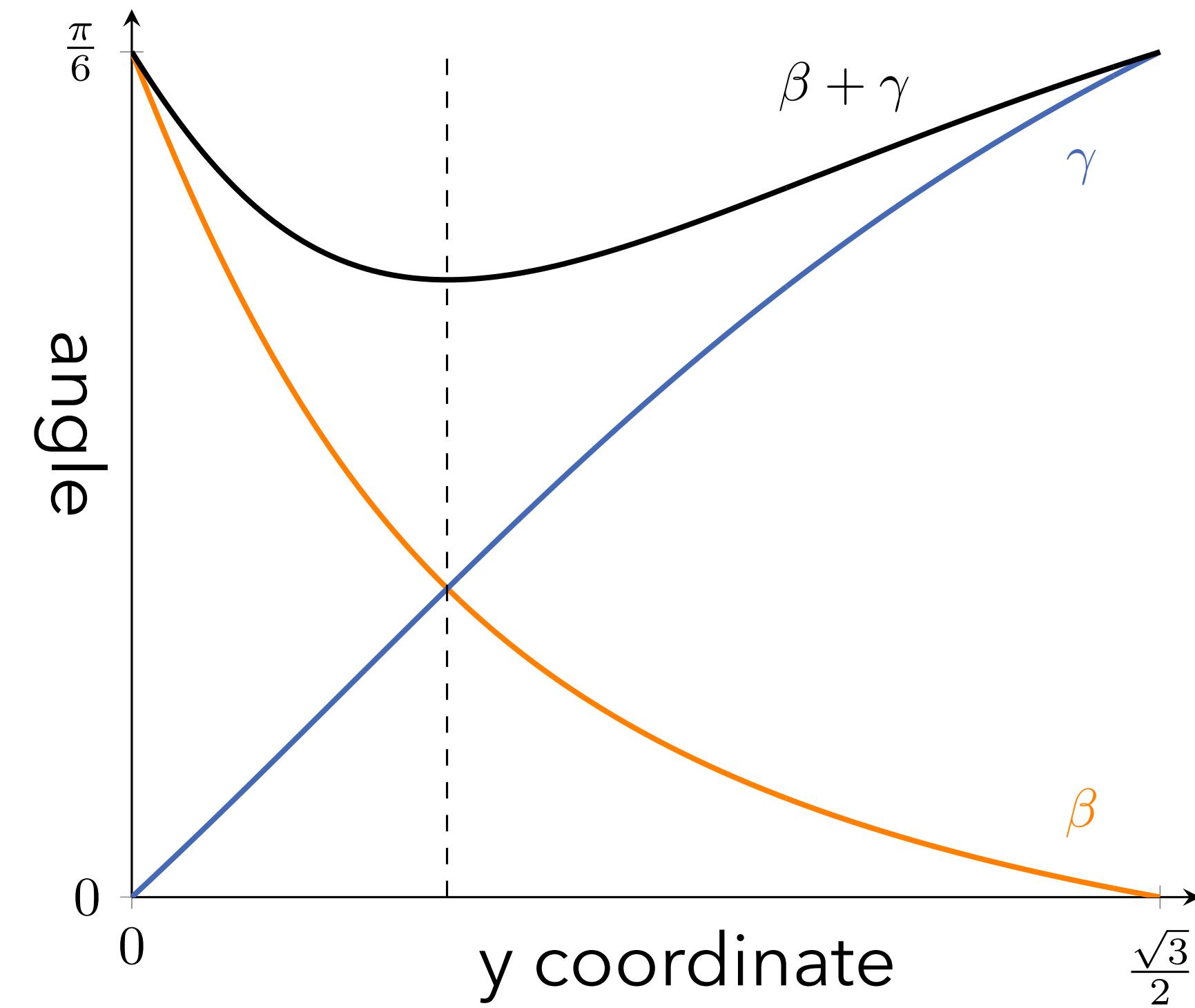
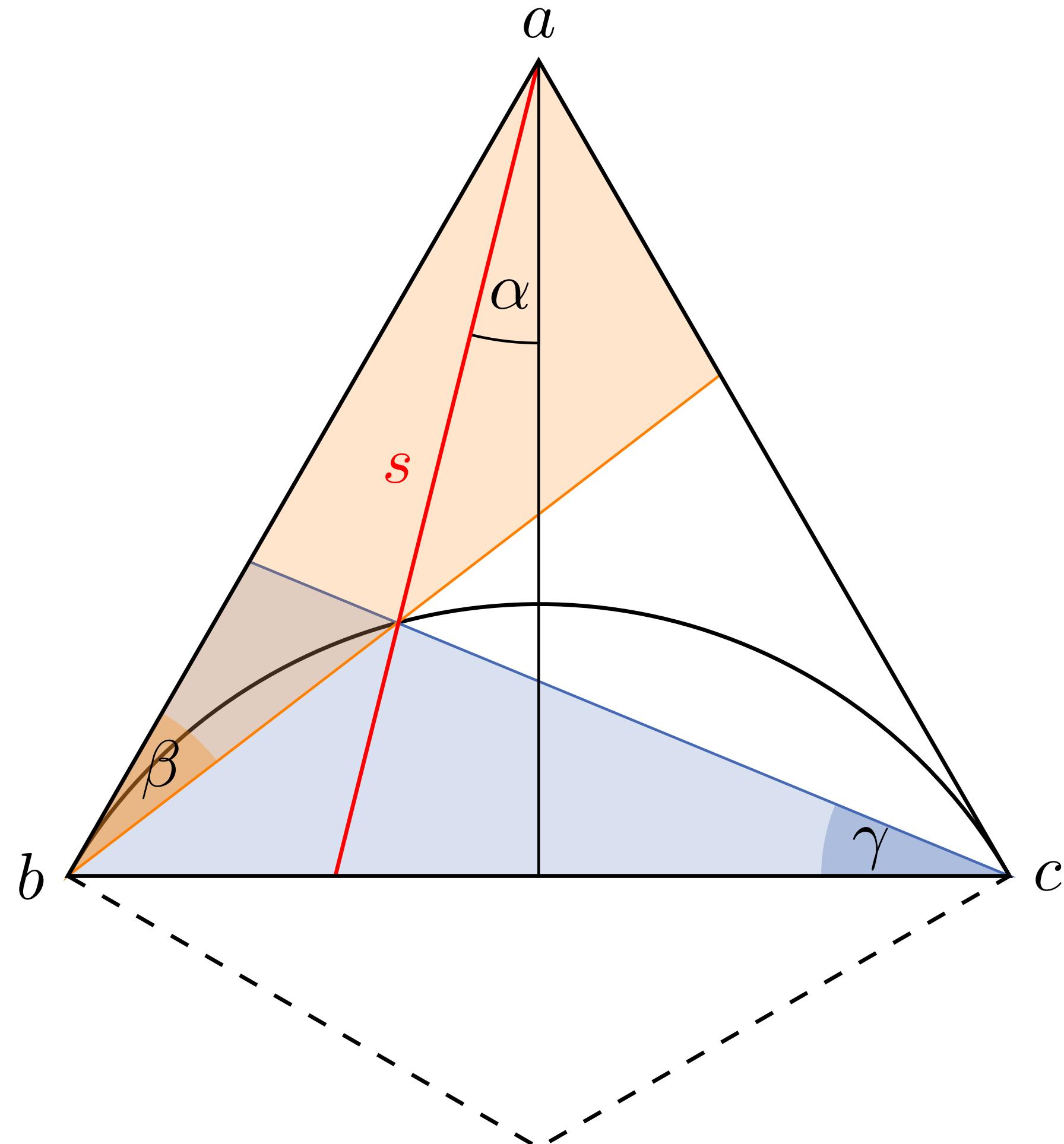
Minimum Covering of s



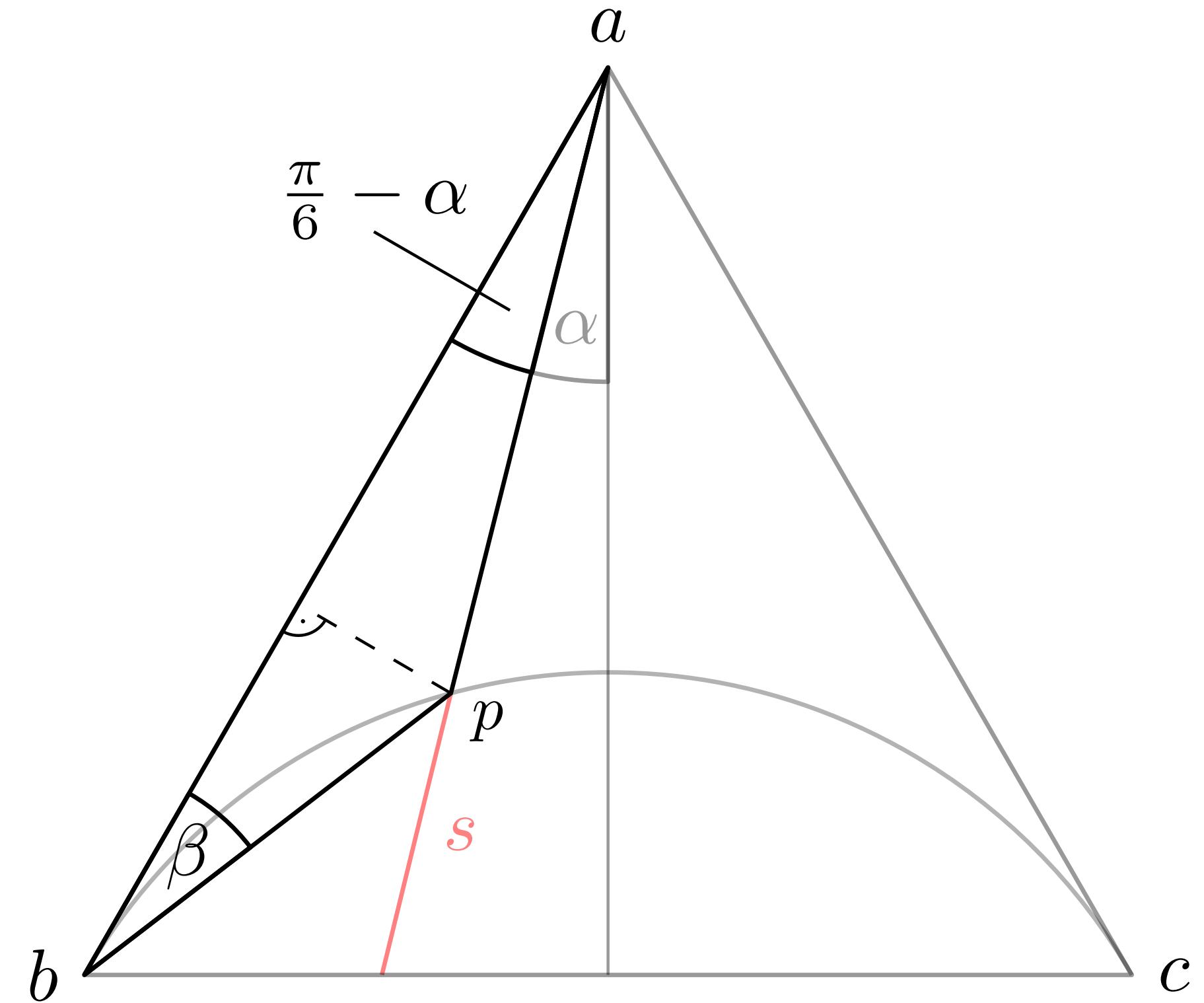
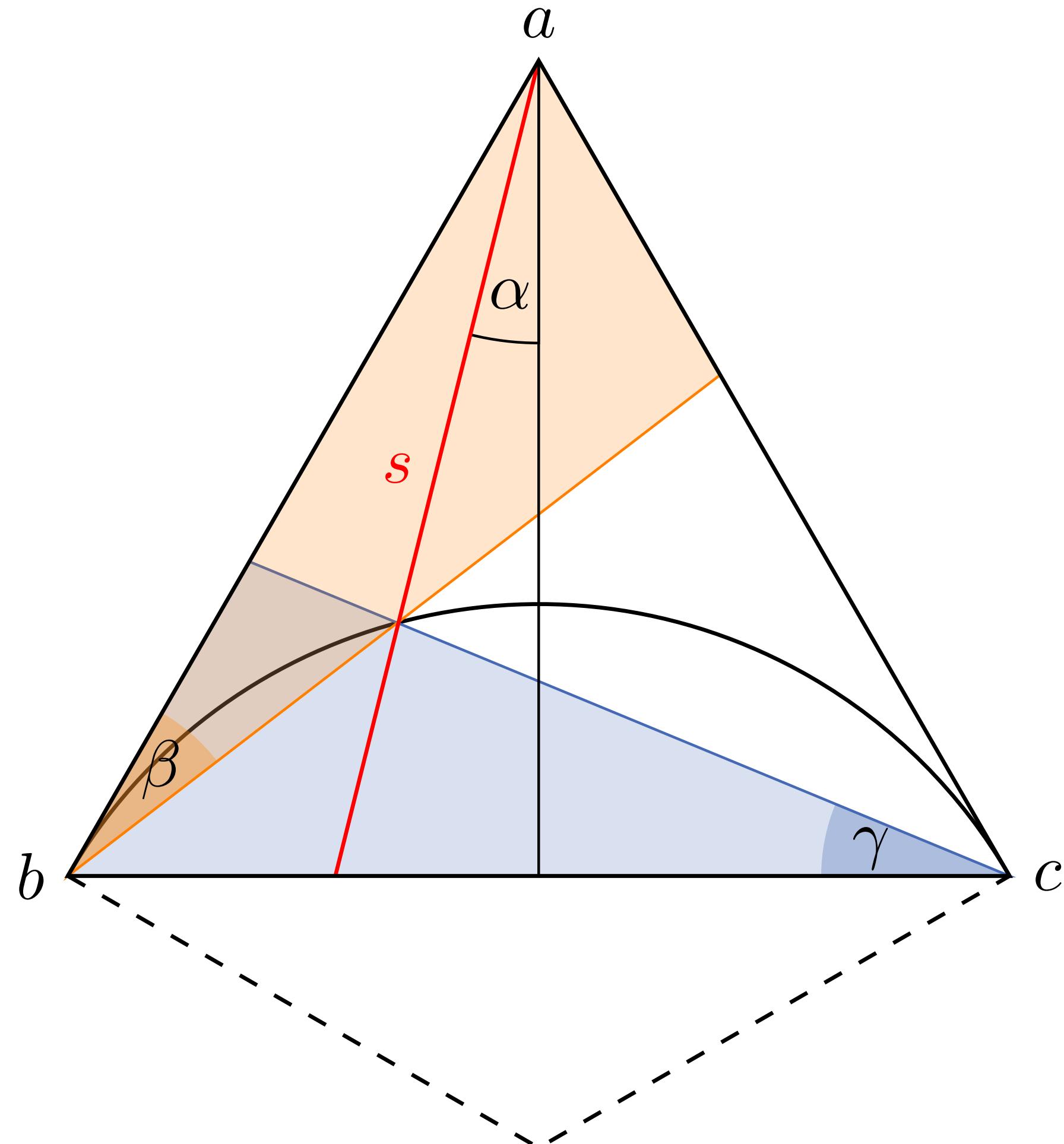
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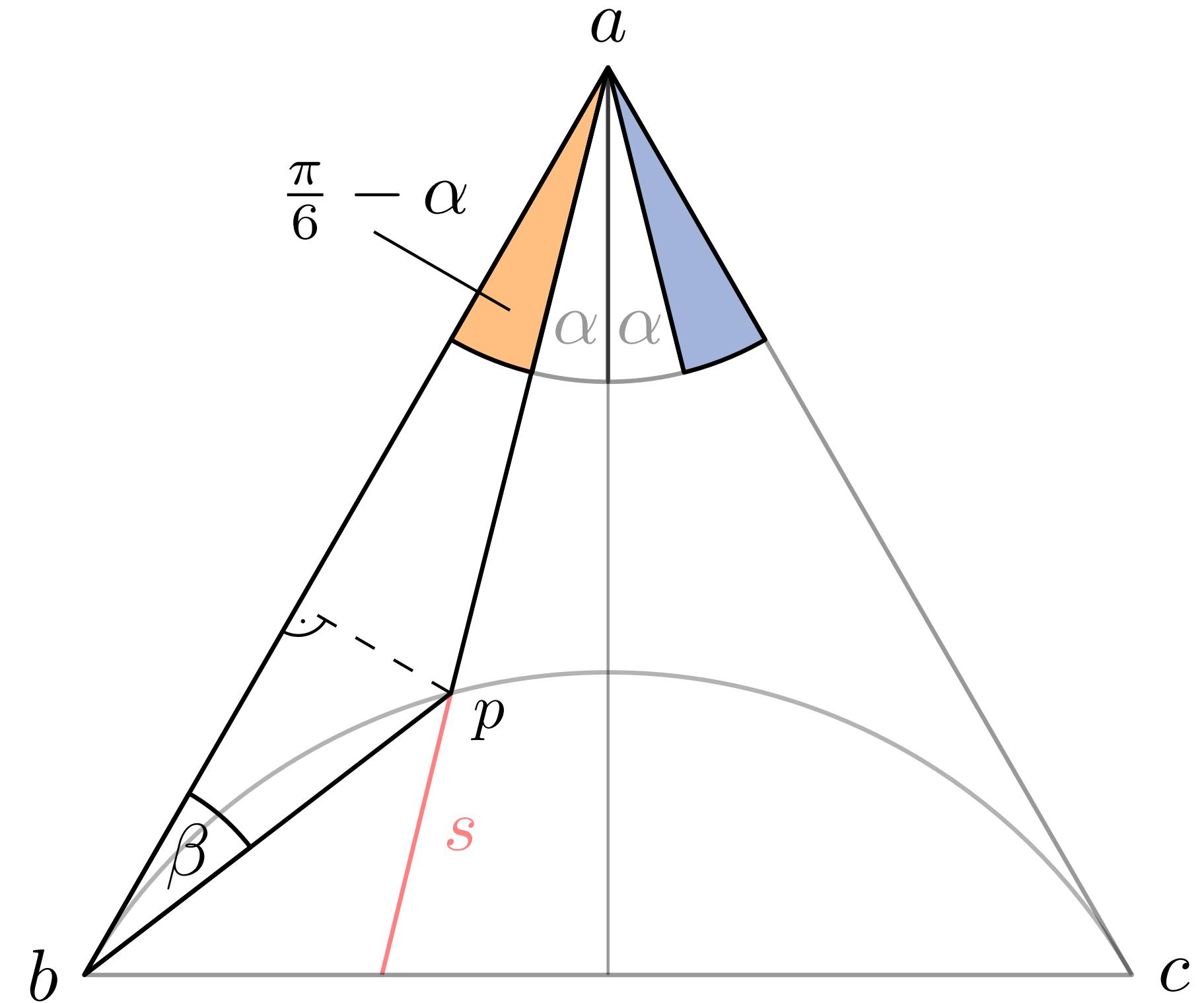
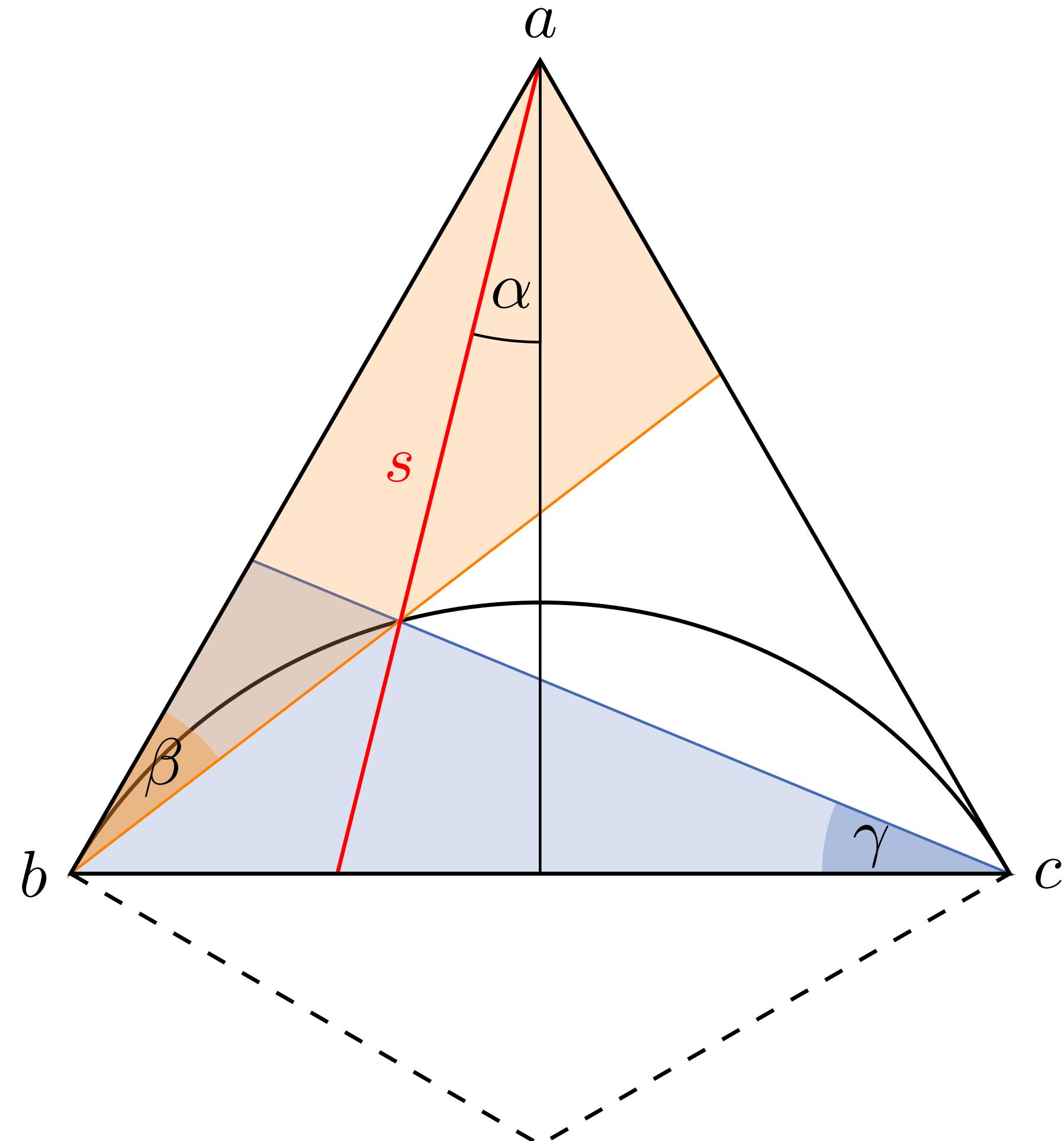
Minimum Covering of s



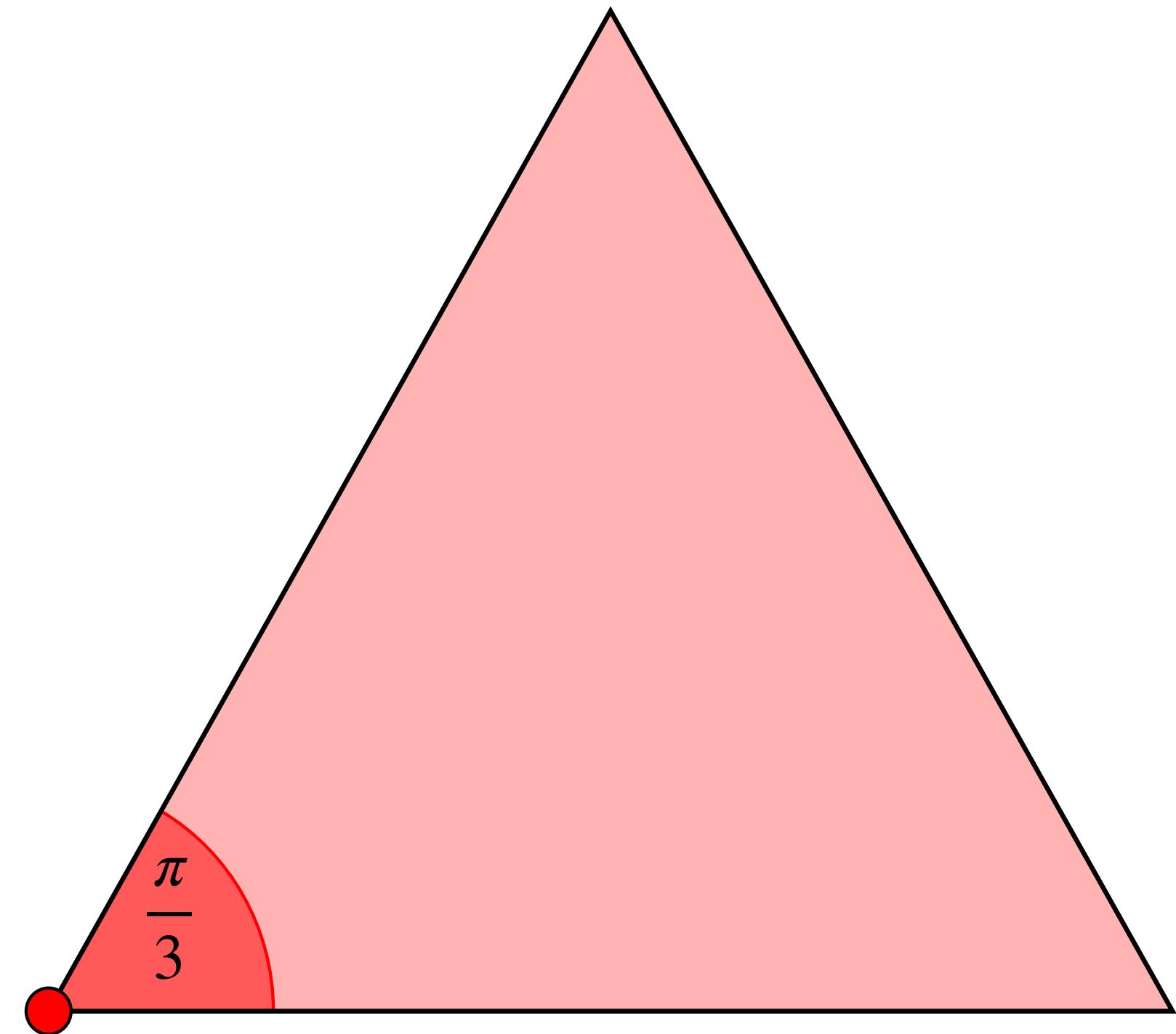
Minimum Covering of s



Minimum Covering of s



Optimal Covering of Equilateral Triangles



Introduction

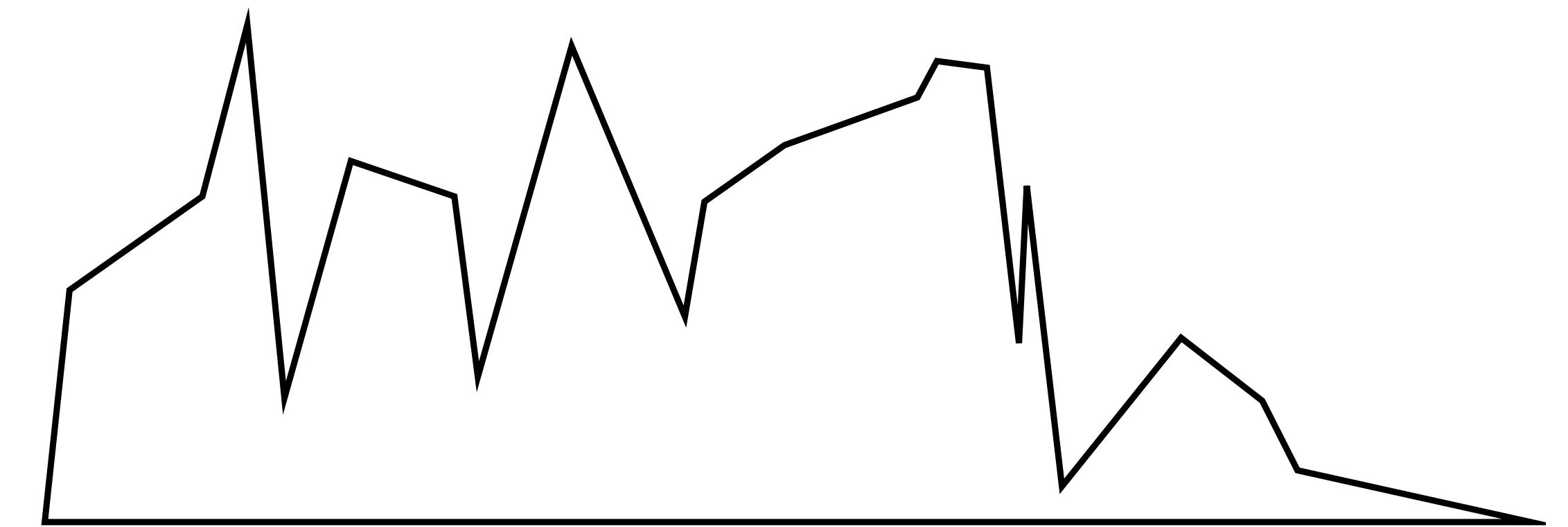
Equilateral Triangles

Tight Upper Bound for **Histograms**

Simple Polygons

Duality

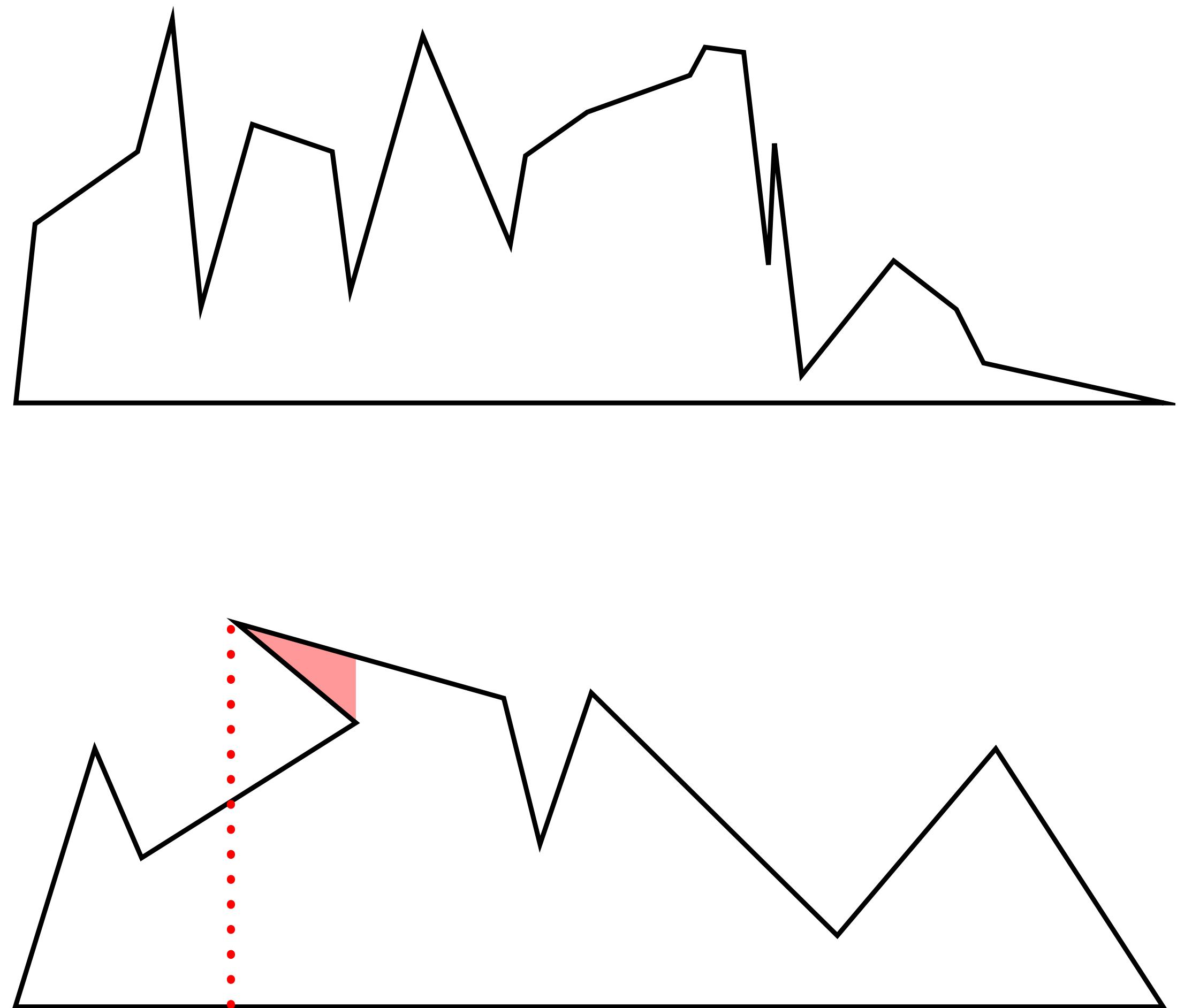
Conclusion



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Equilateral Triangles

Tight Upper Bound for **Histograms**

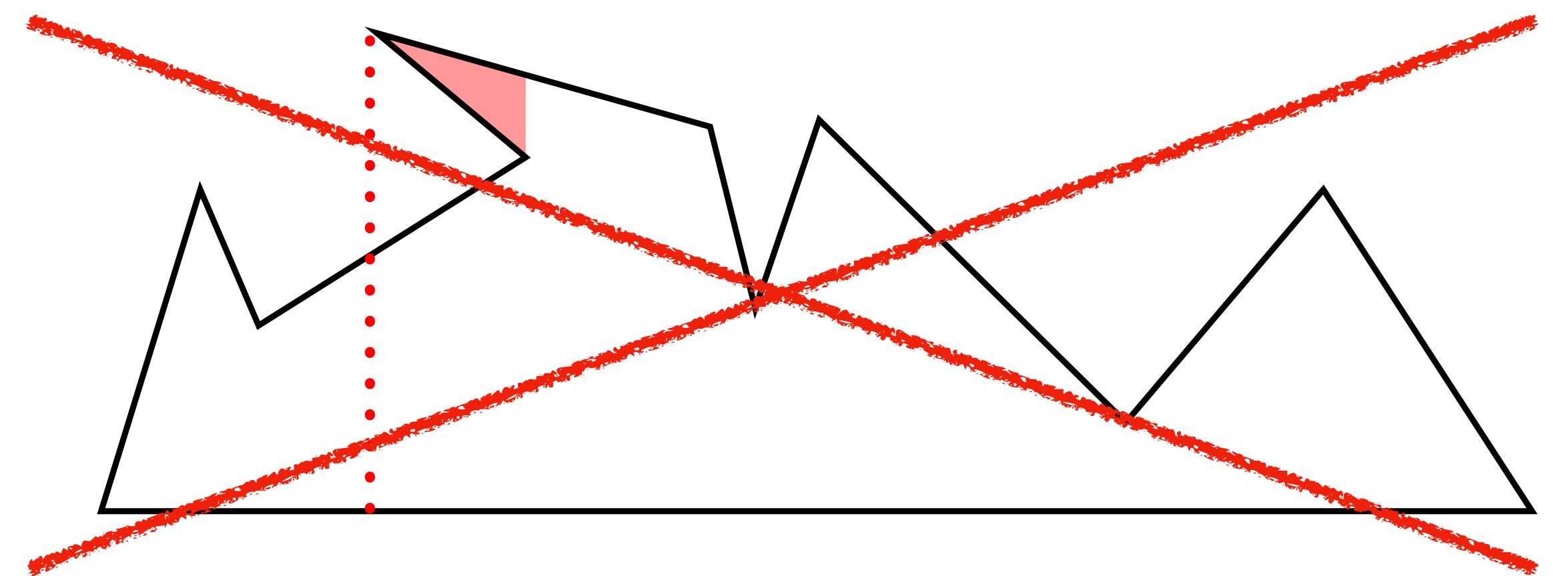
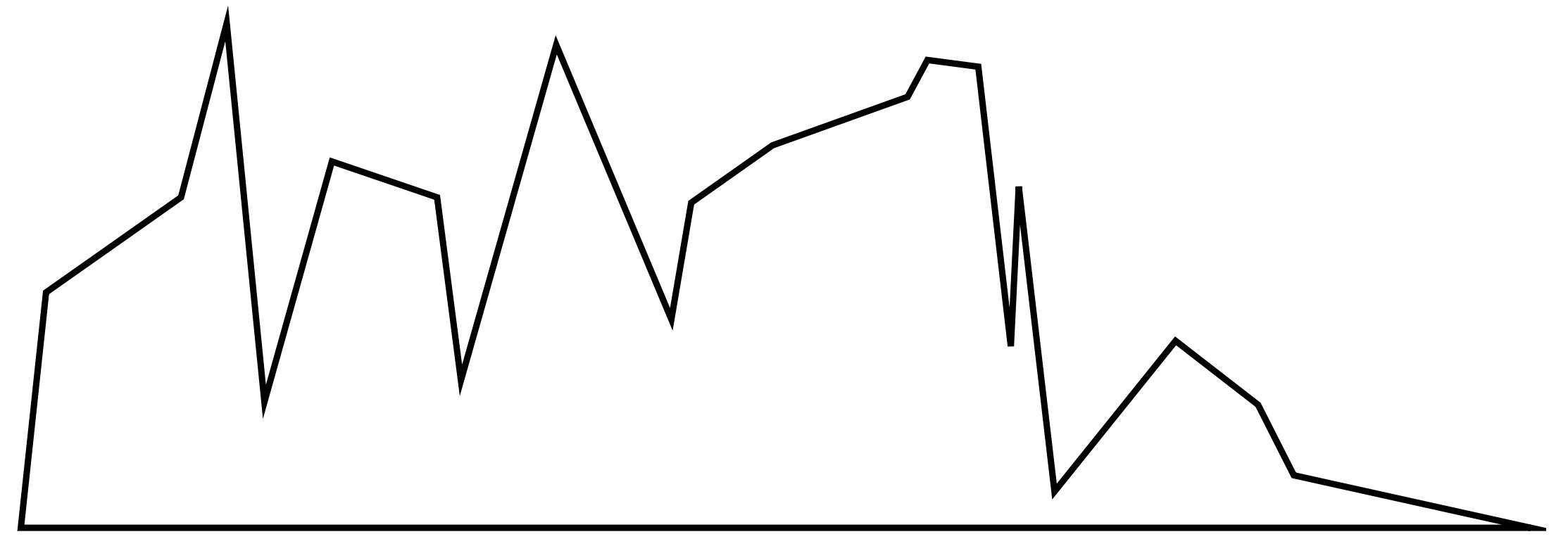
Simple Polygons
Duality
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Equilateral Triangles

Tight Upper Bound for **Histograms**

Simple Polygons
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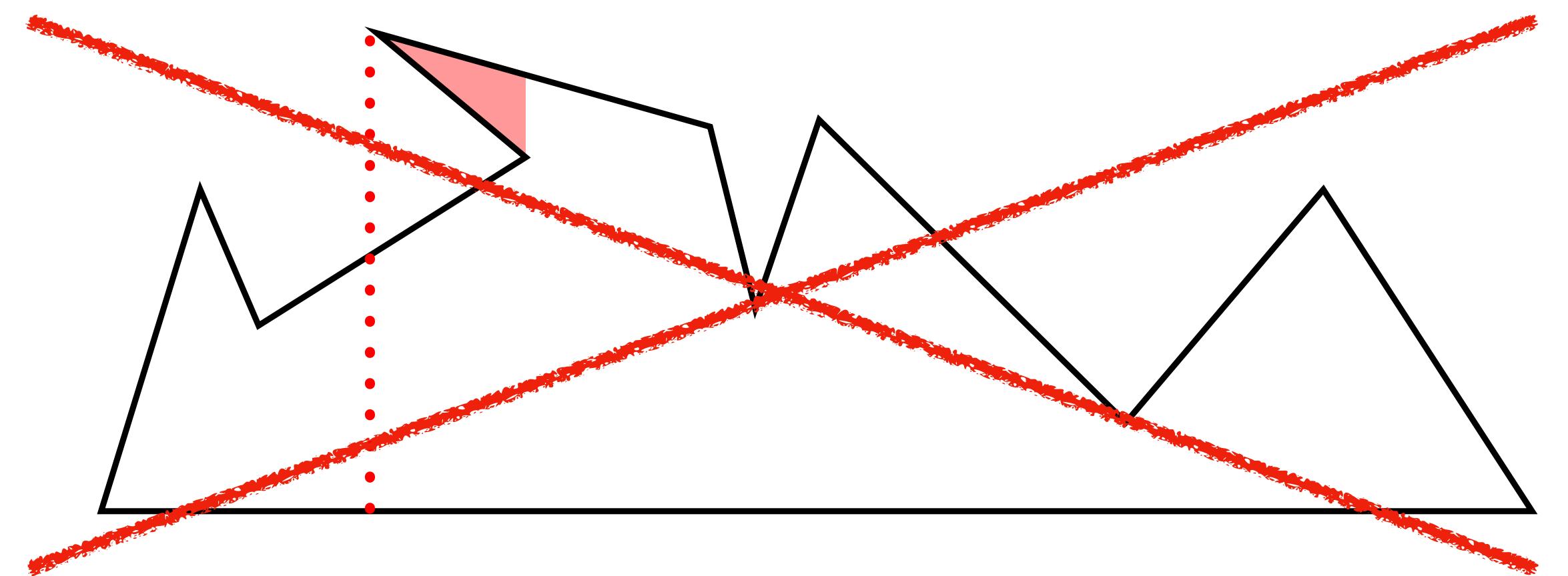
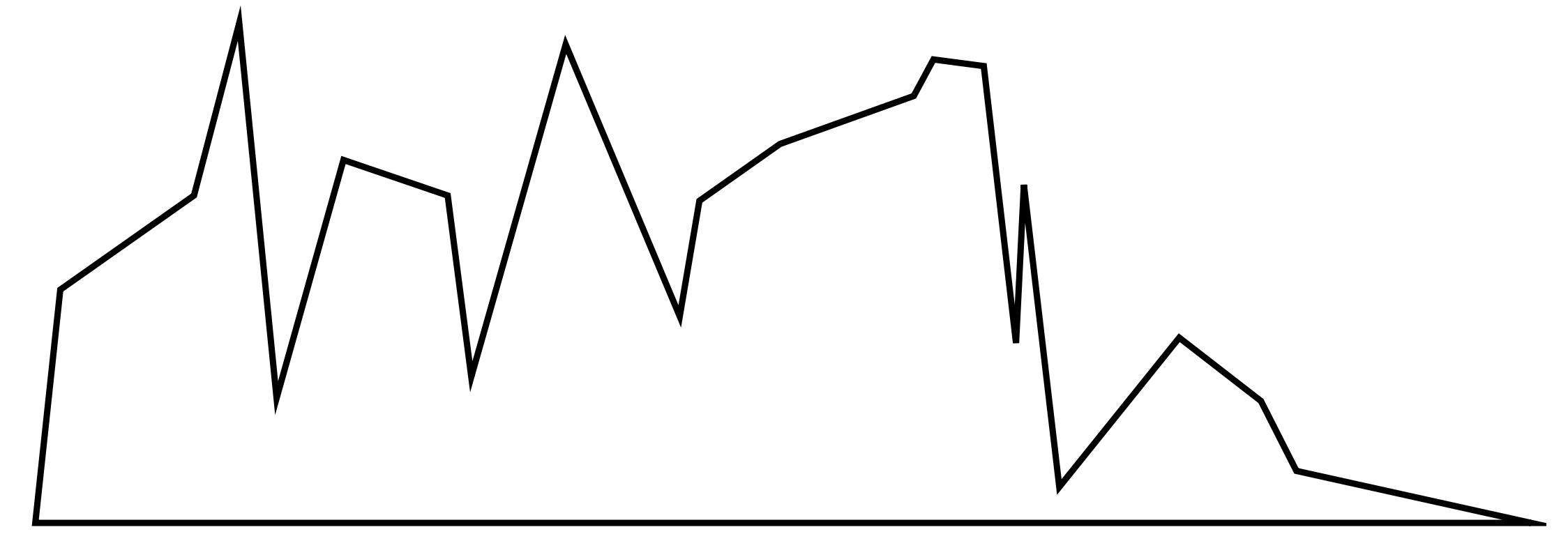


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Equilateral Triangles

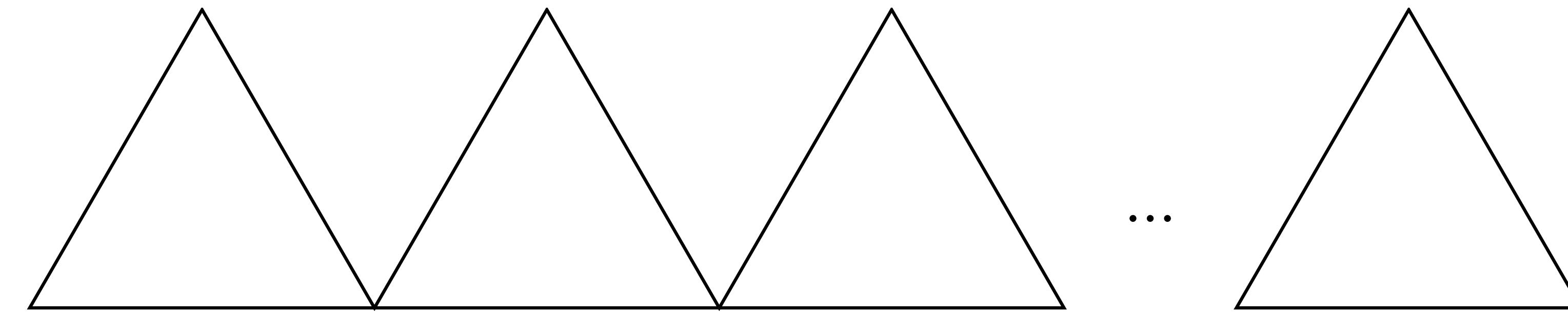
Tight Upper Bound for **Histograms**

Simple Polygons
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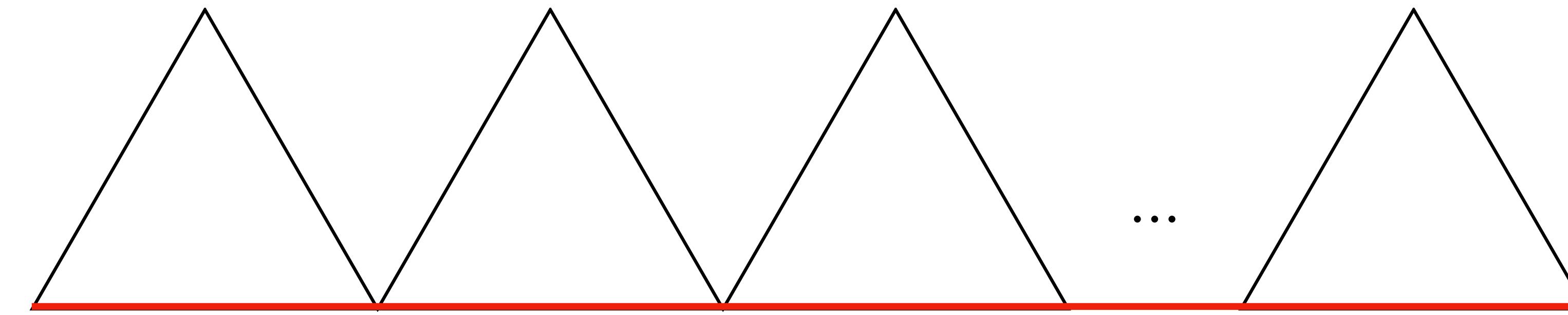
$$(n - 1) \frac{\pi}{6}$$



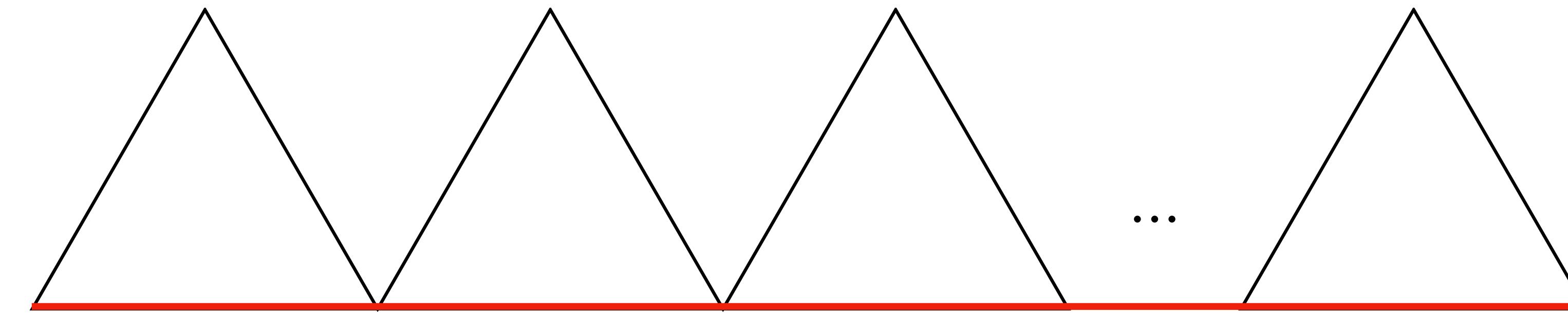
Lower Bound for Histograms



Lower Bound for Histograms

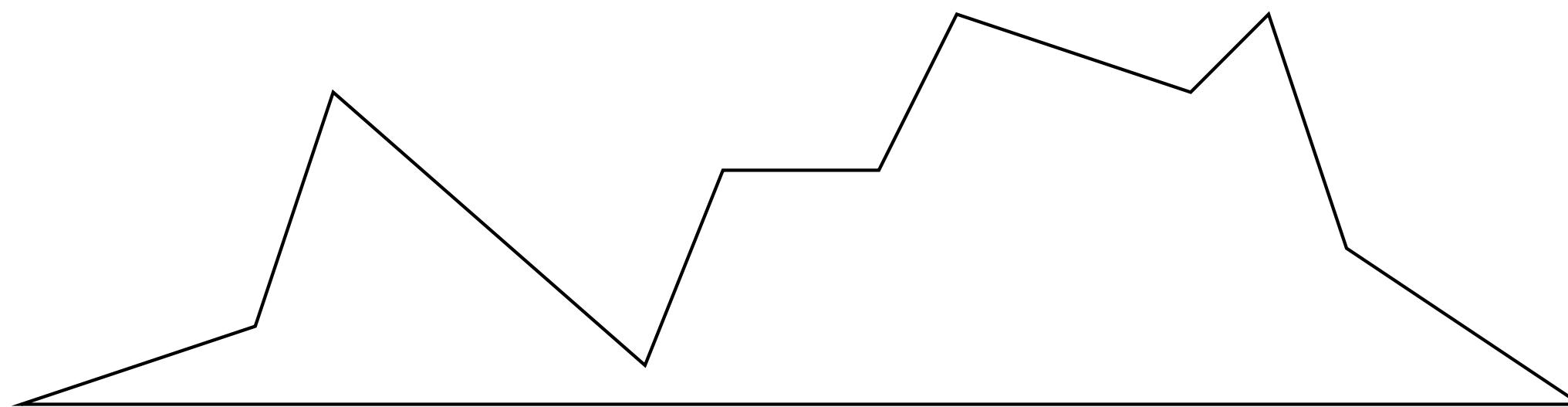


Lower Bound for Histograms

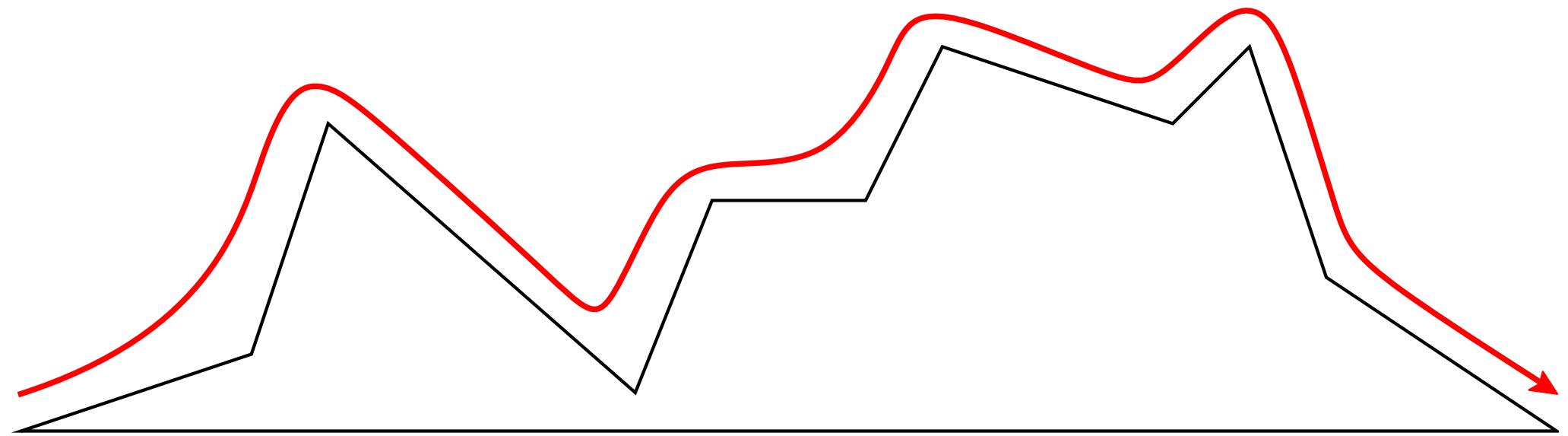


$$\text{Lower Bound: } \frac{n-1}{2} \frac{\pi}{3} = (n-1) \frac{\pi}{6}$$

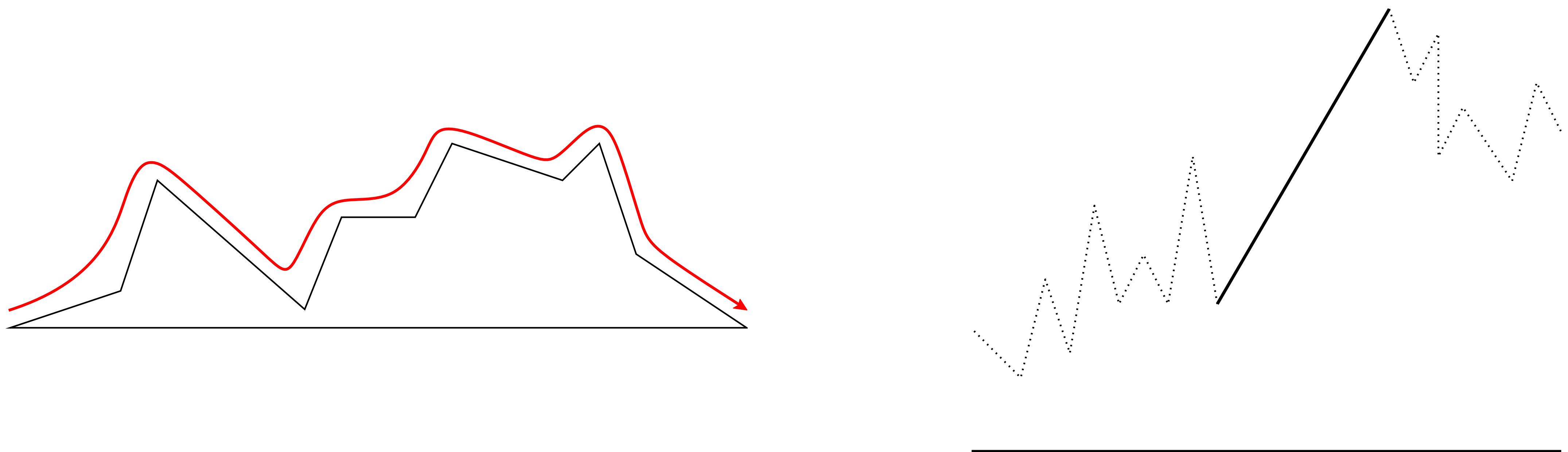
Upper Bound for Histograms



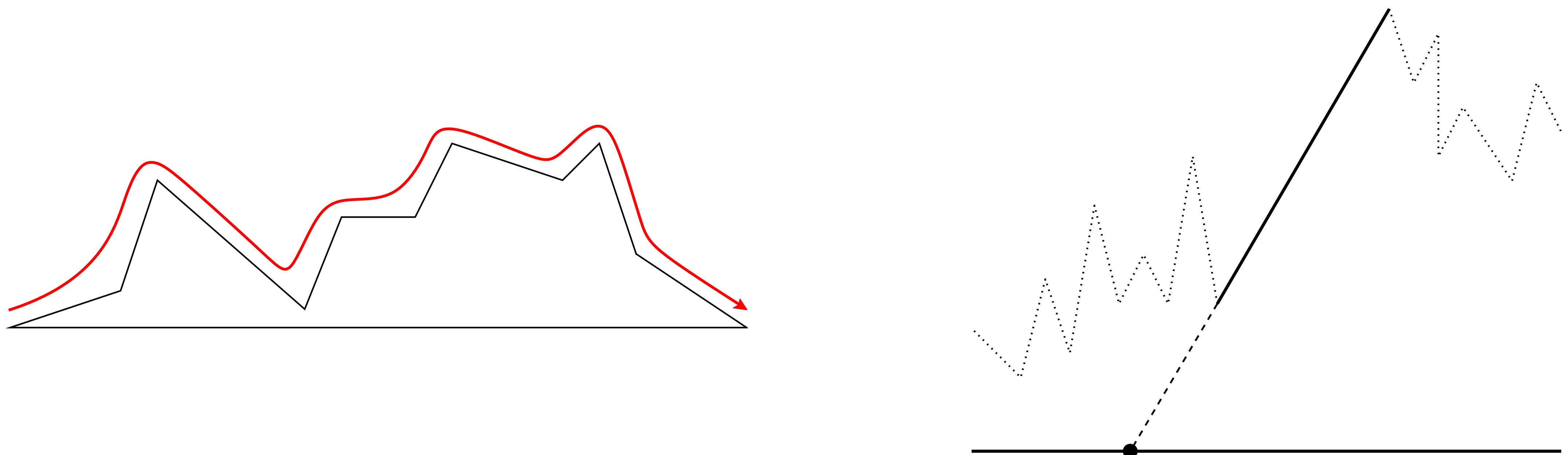
Upper Bound for Histograms



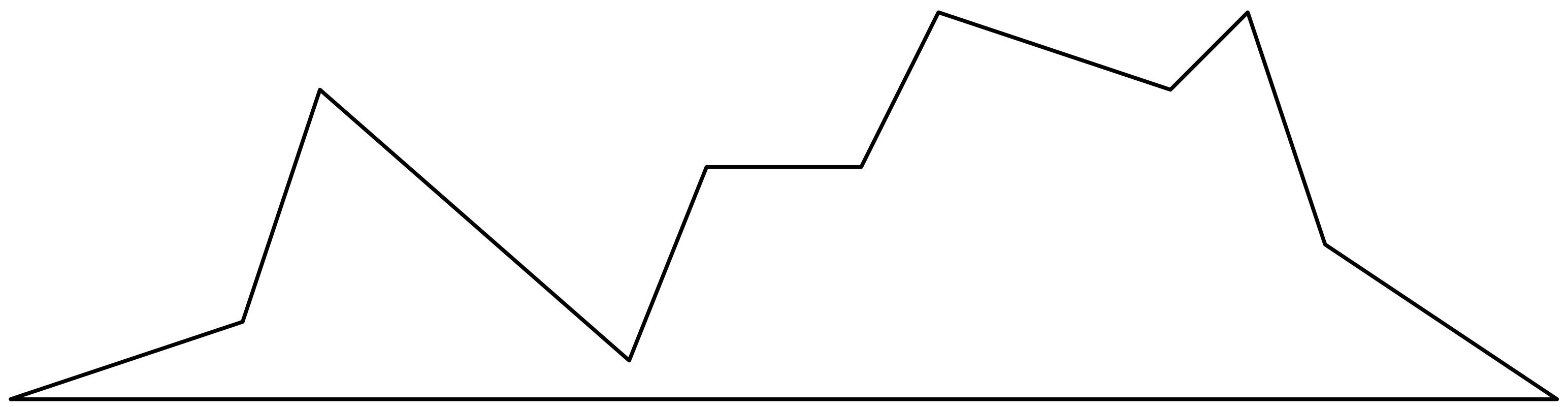
Upper Bound for Histograms



Upper Bound for Histograms

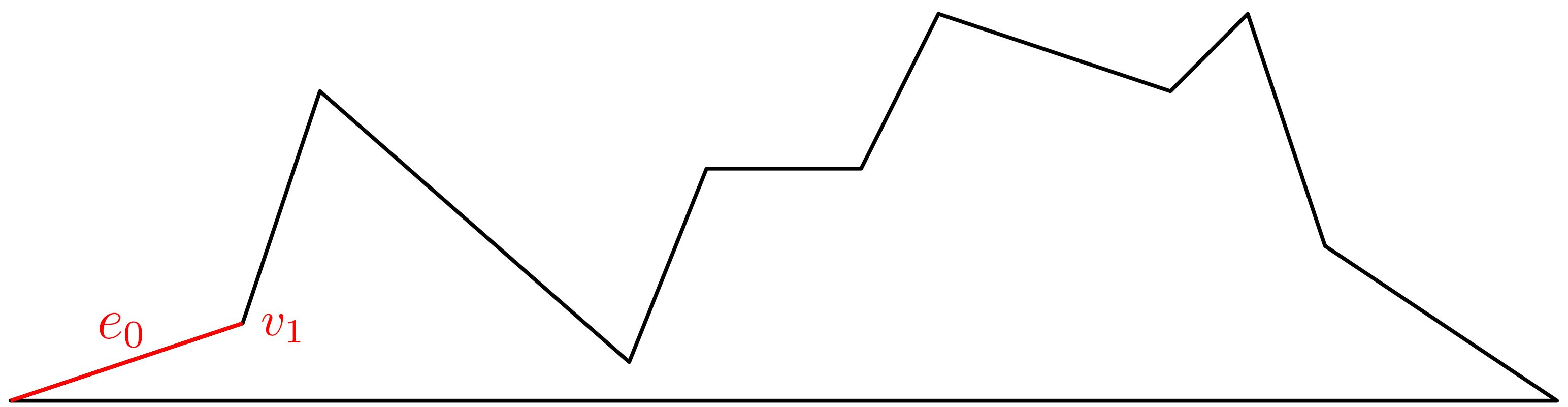


Histogram Covering Algorithm – Example



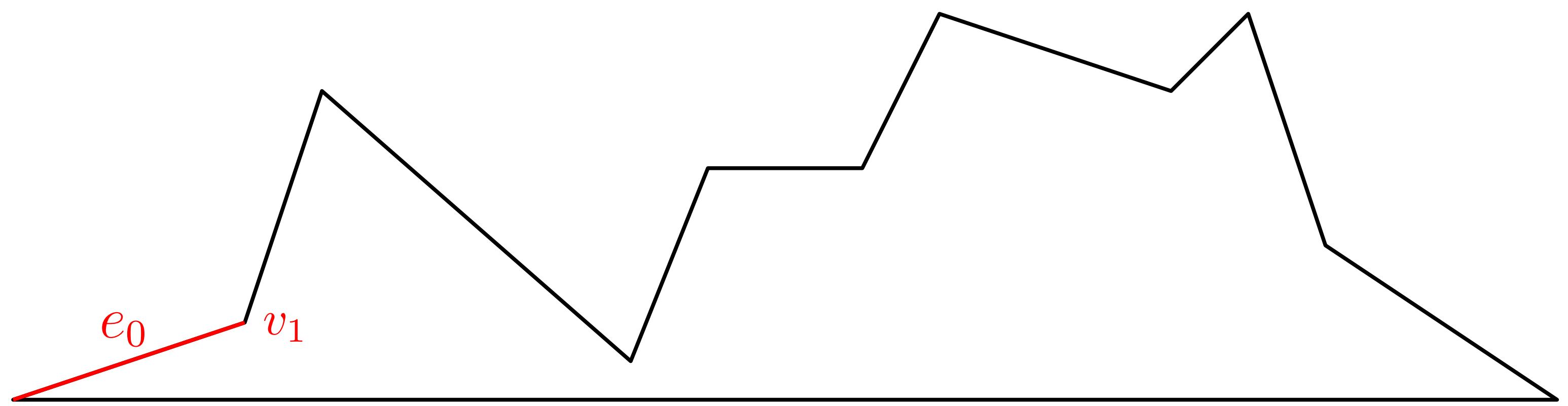
Case	Condition
1	e_i is covered
2	v_{i+1} is reflex
3	v_{i+2} is convex
4	there is another uncovered, completely visible edge
5	visibility extension of e_{i+1} intersects with baseline
6	none of the above applies

Histogram Covering Algorithm – Example



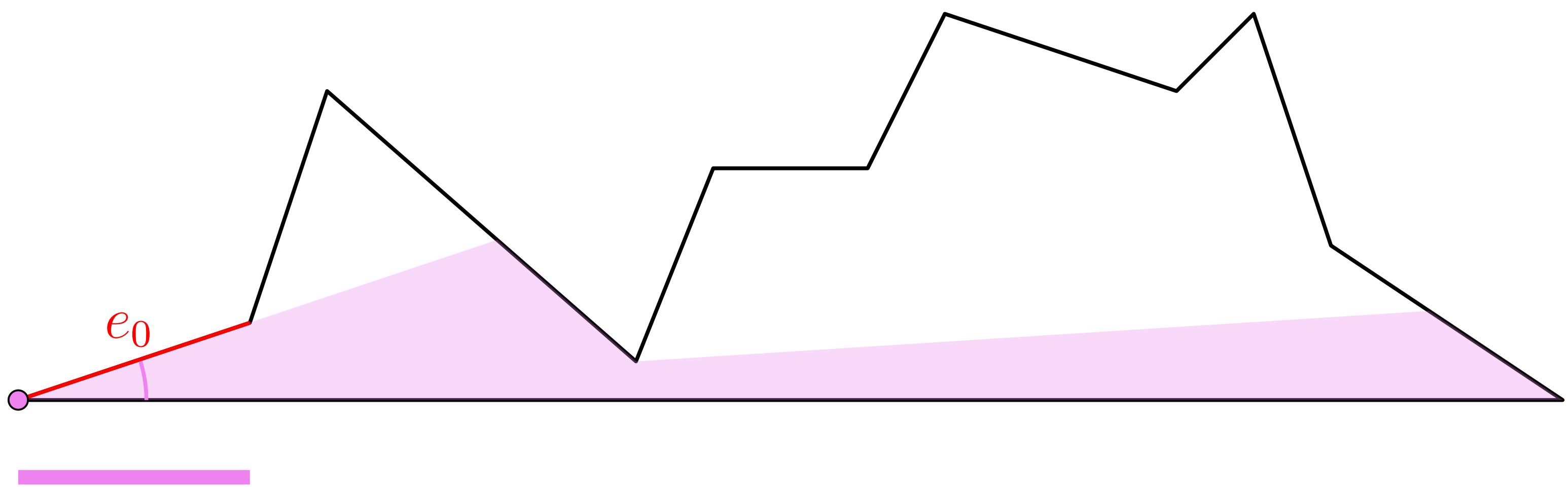
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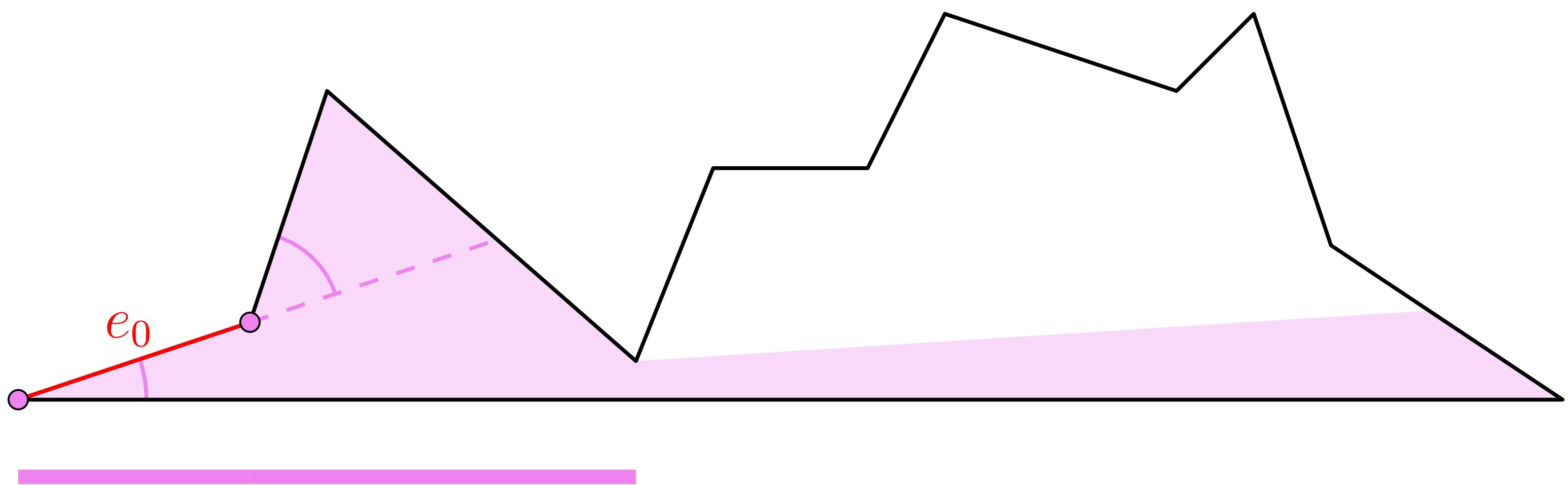
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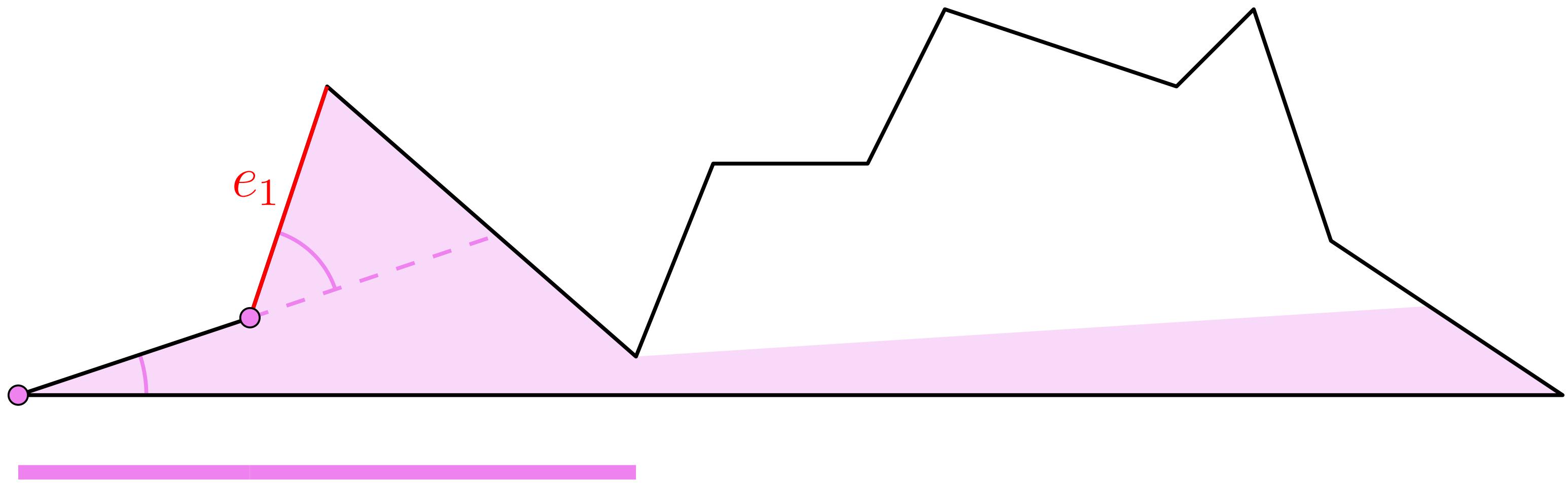
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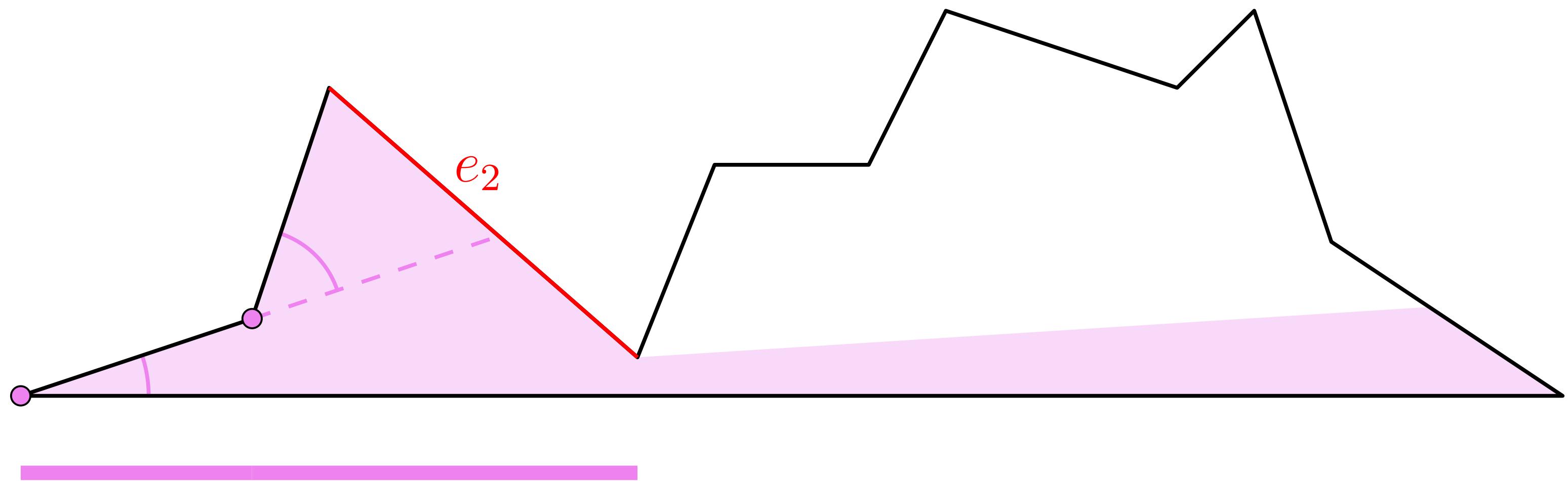
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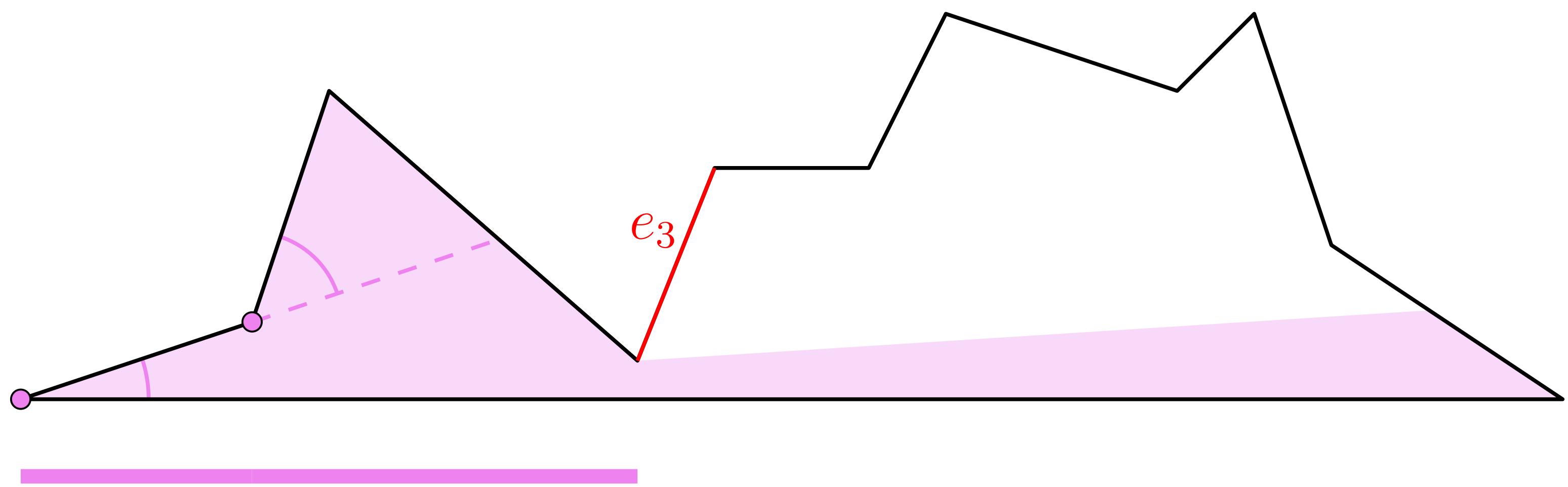
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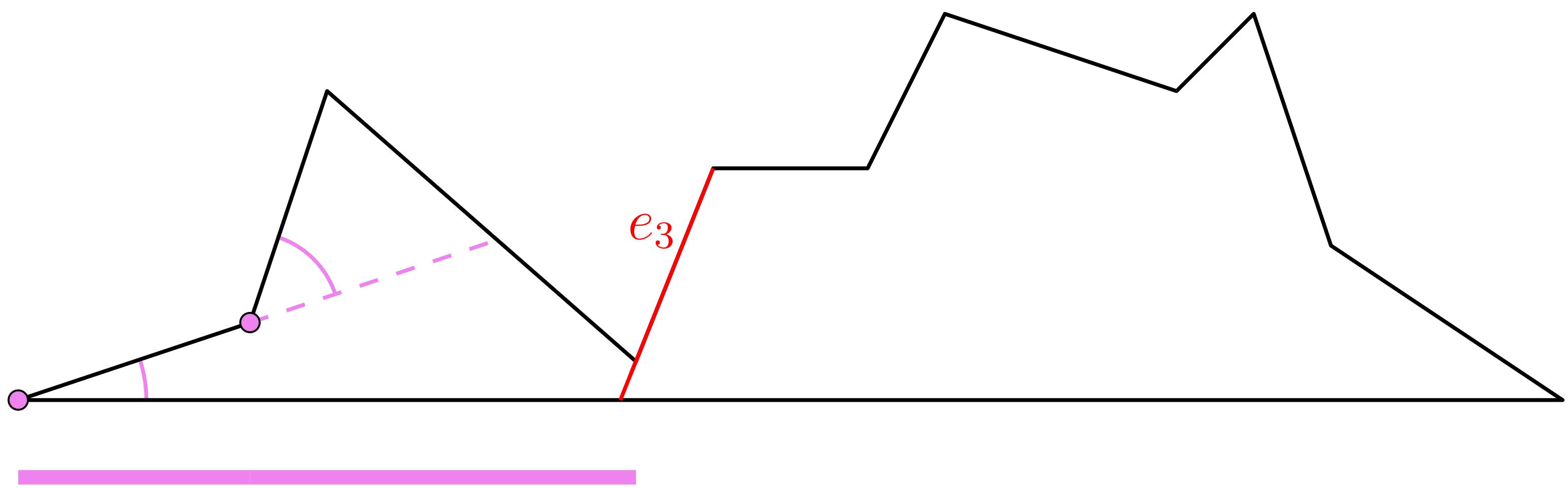
Case	Condition
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Histogram Covering Algorithm – Example



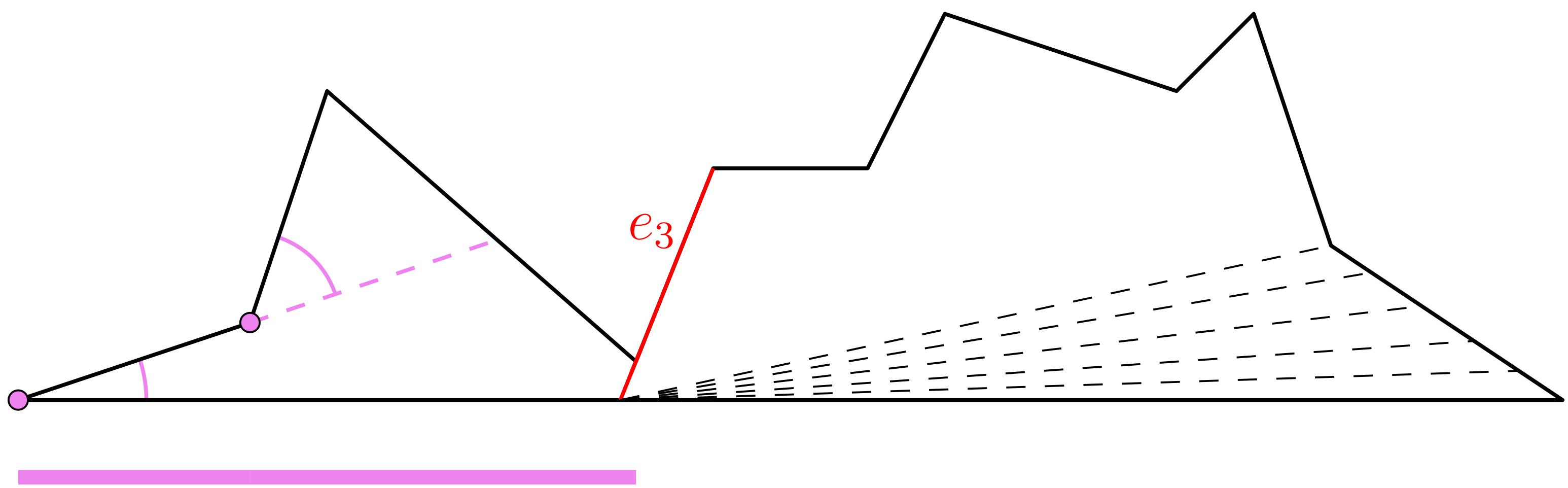
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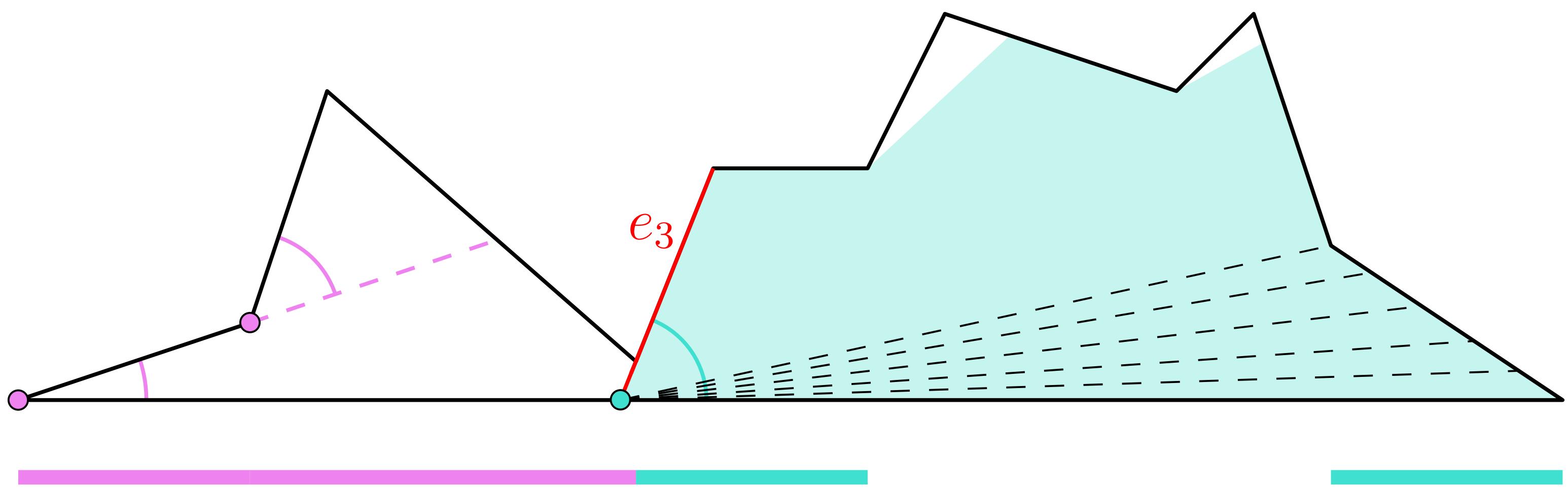
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Histogram Covering Algorithm – Example



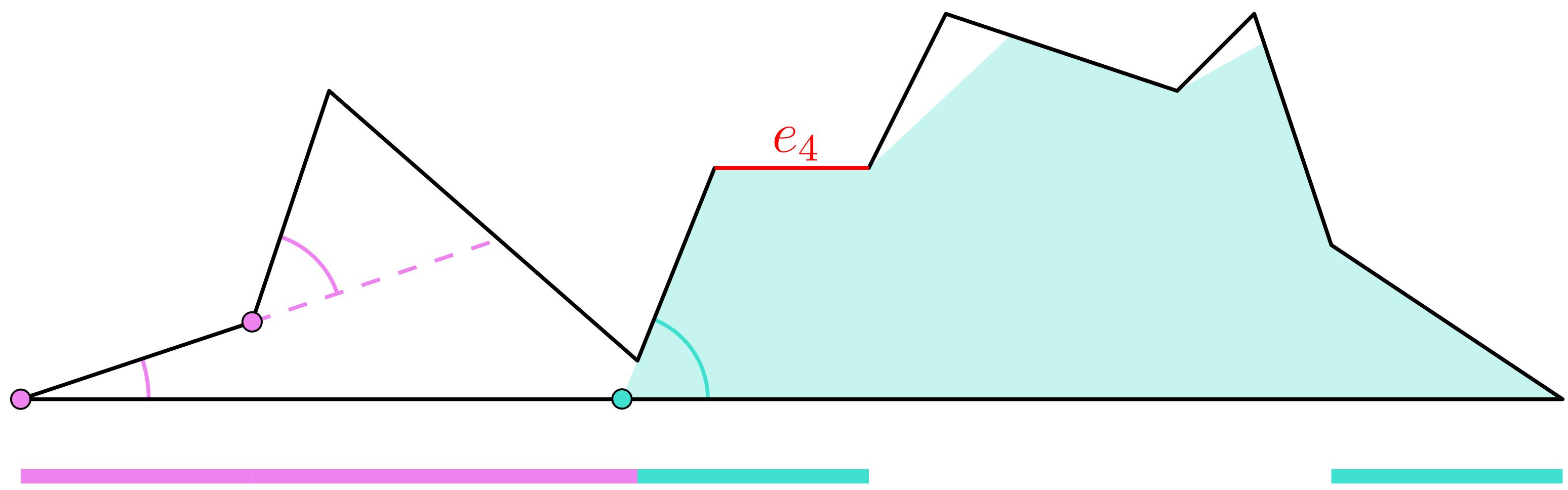
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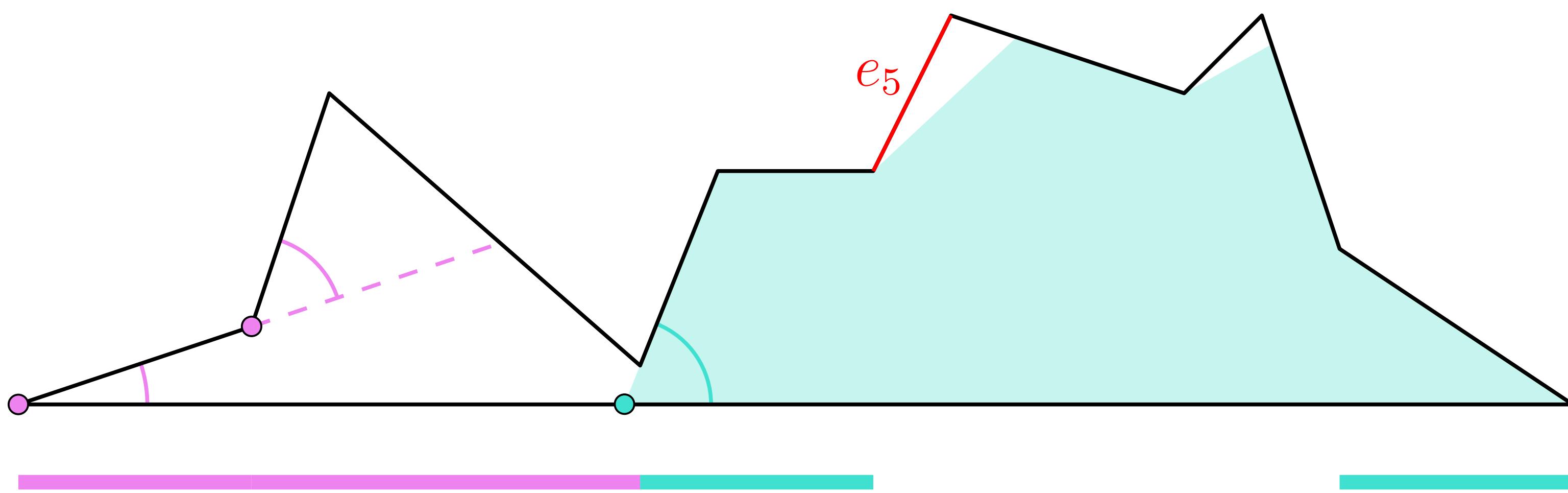
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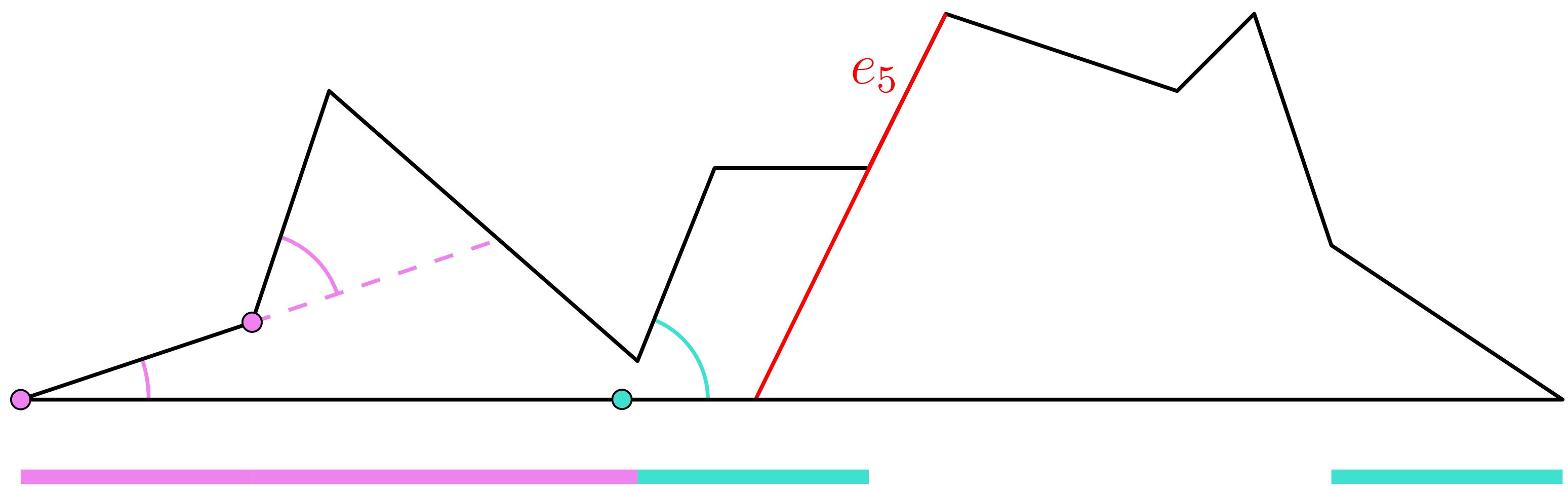
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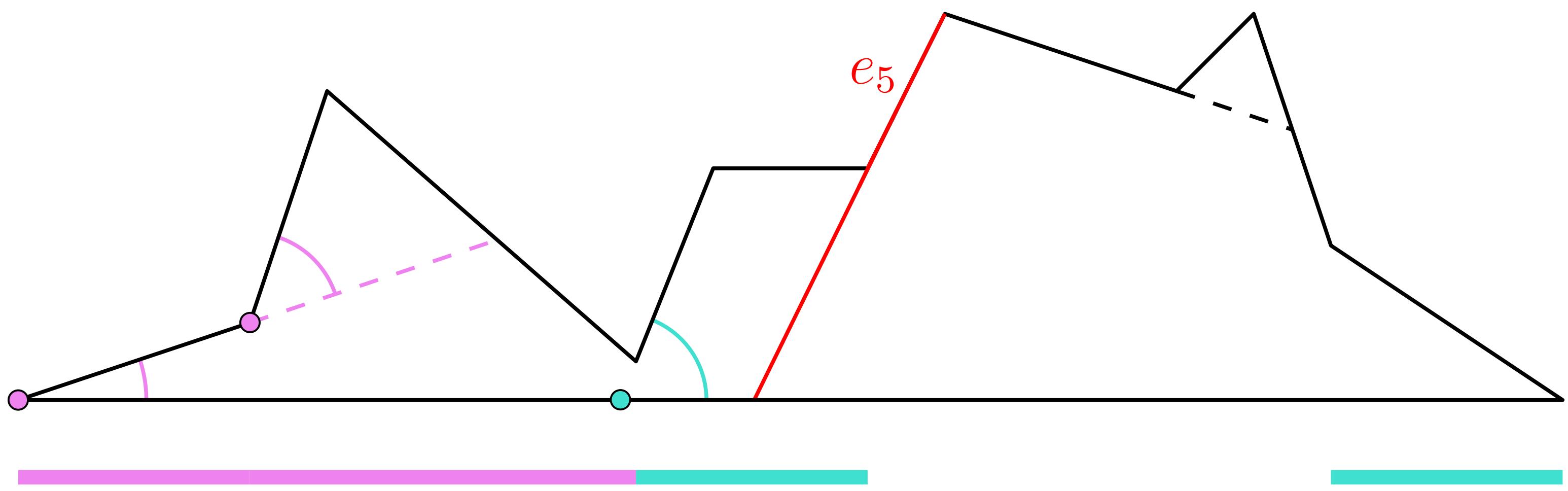
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Histogram Covering Algorithm – Example



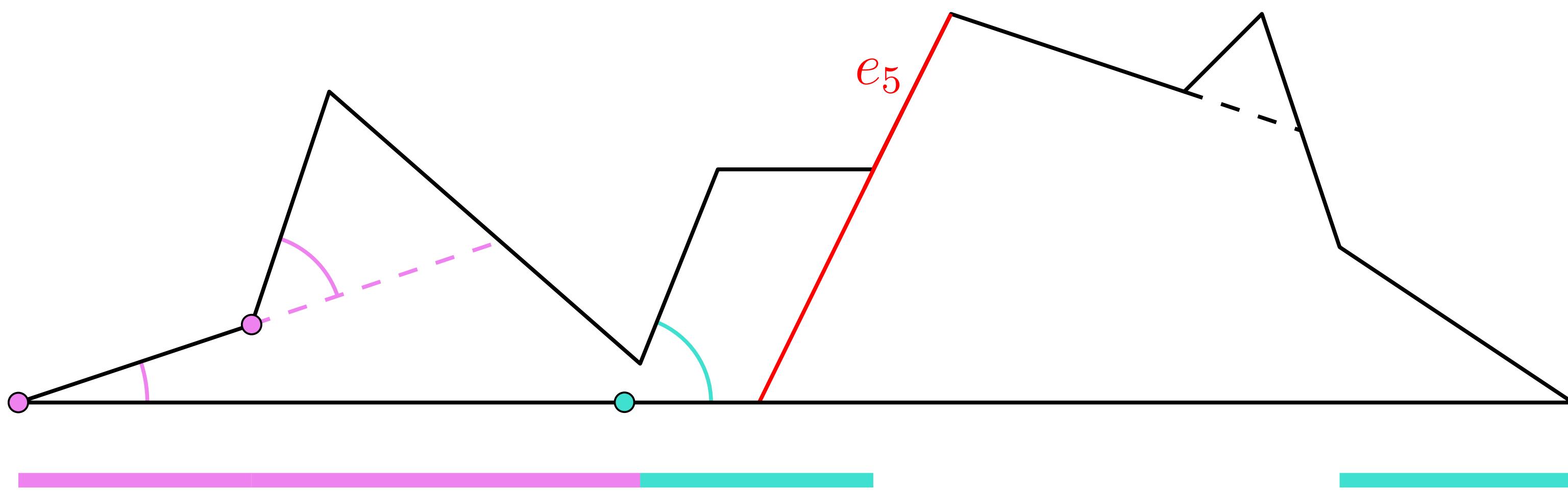
Case	Condition
1	e_i is covered
2	v_{i+1} is reflex
3	v_{i+2} is convex
4	there is another uncovered, completely visible edge
5	visibility extension of e_{i+1} intersects with baseline
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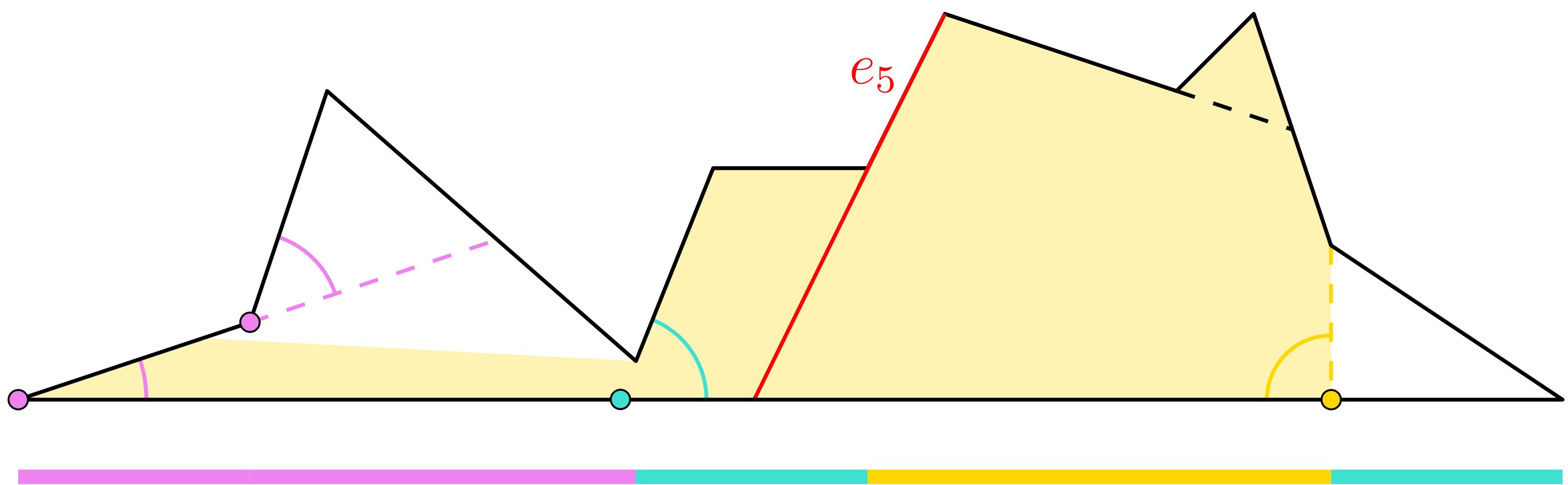
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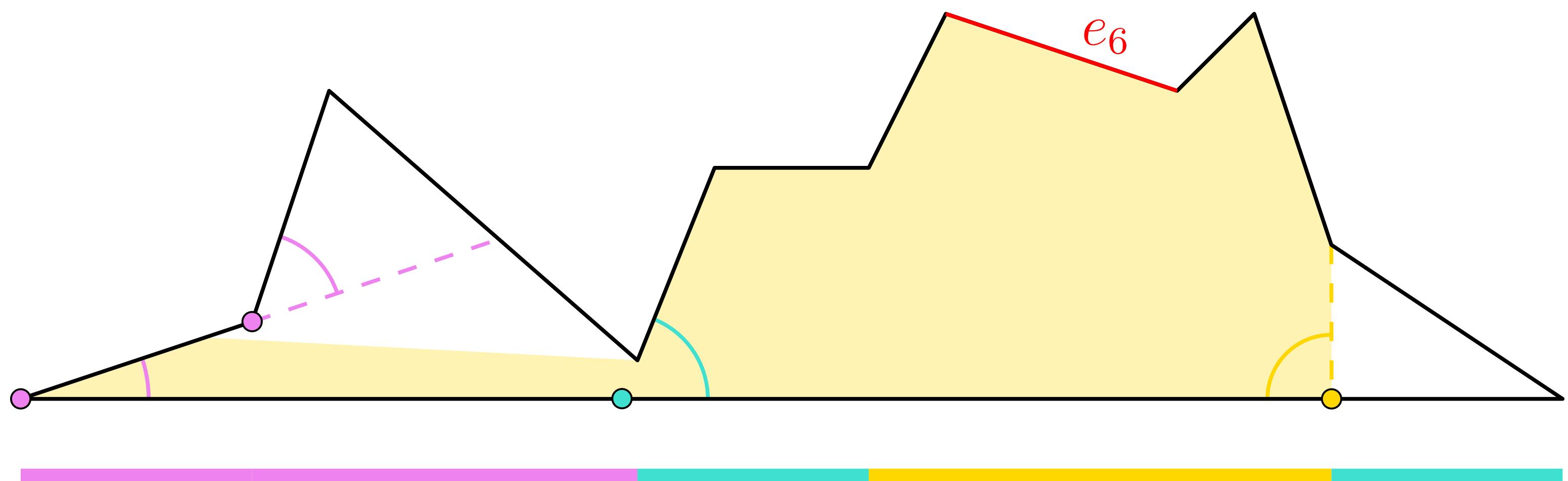
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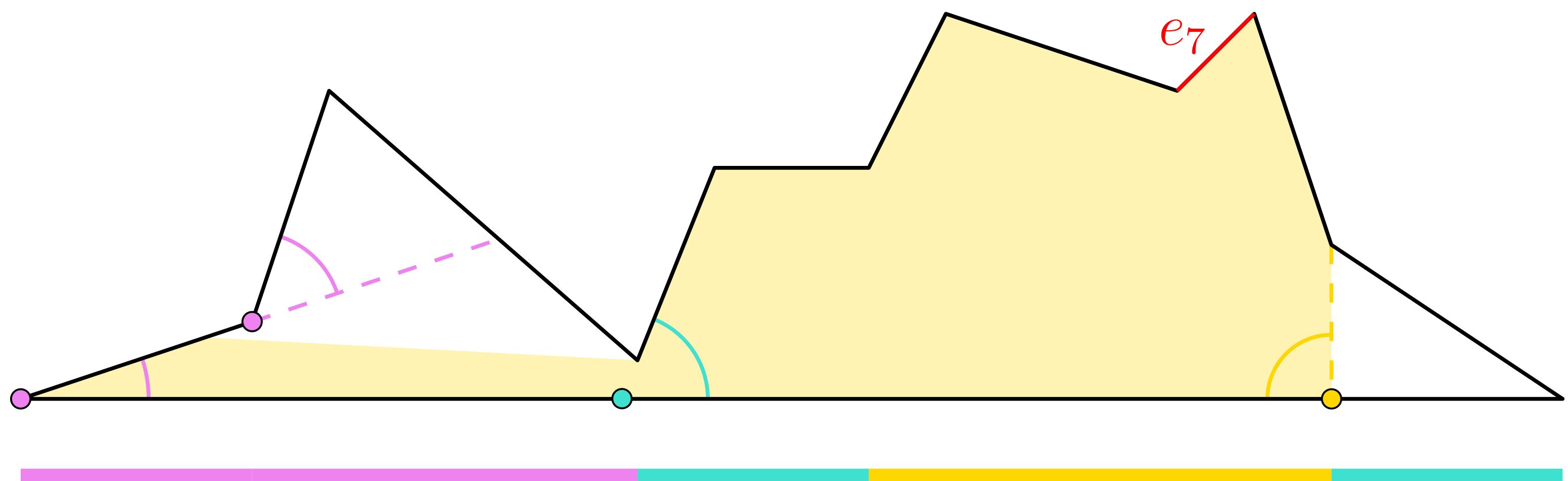
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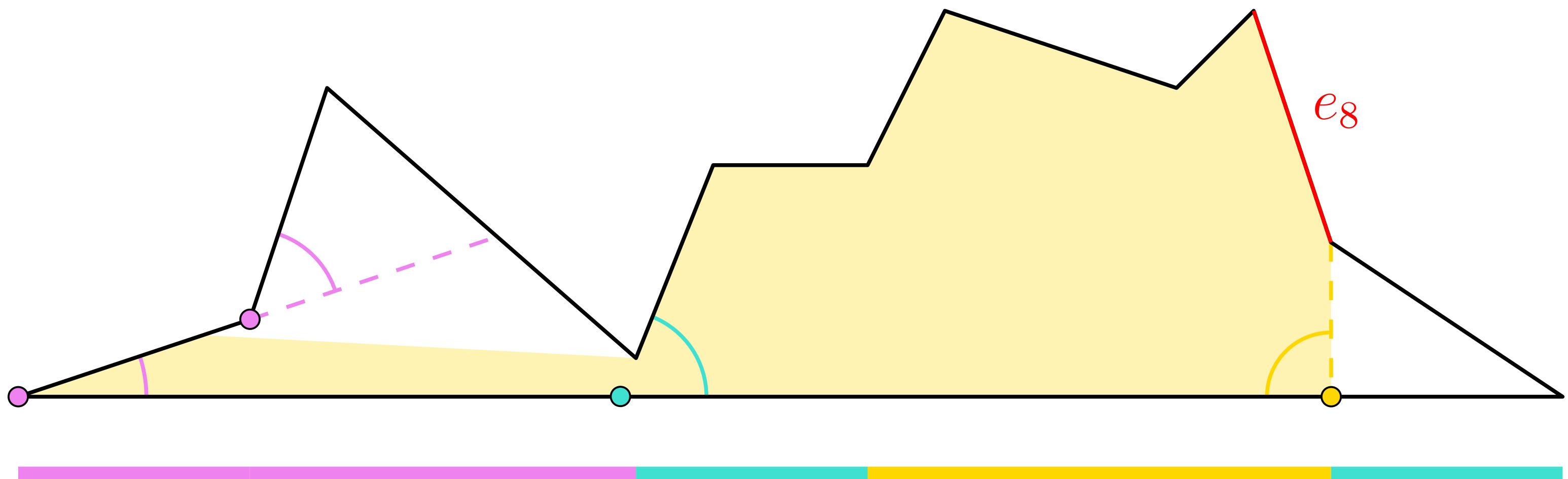
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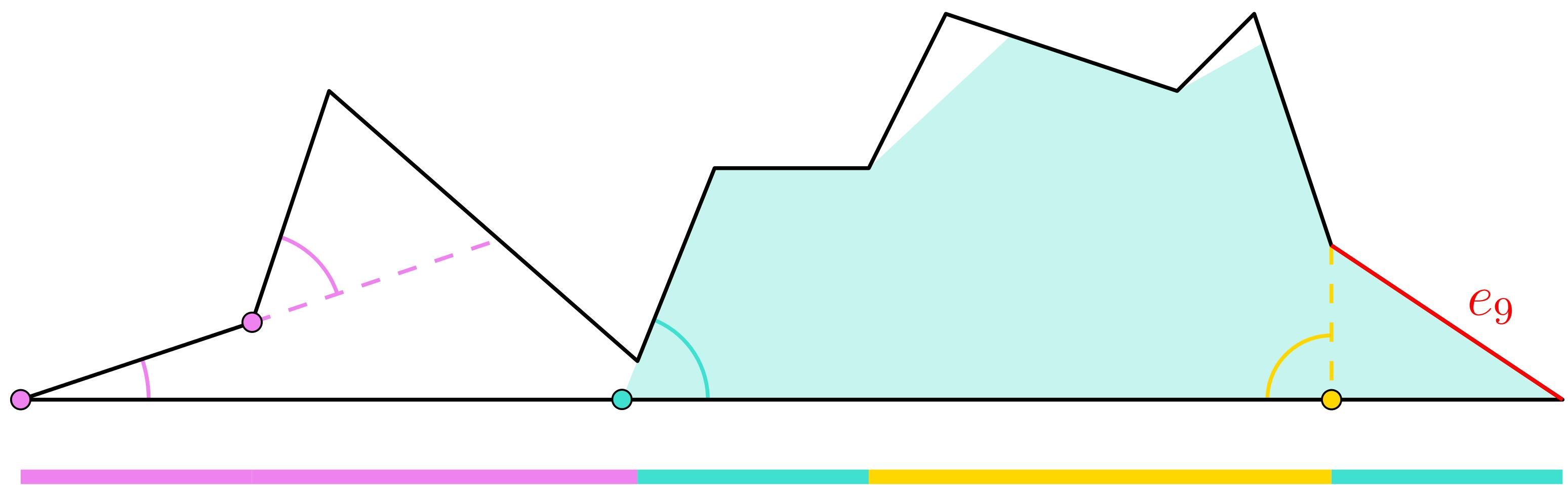
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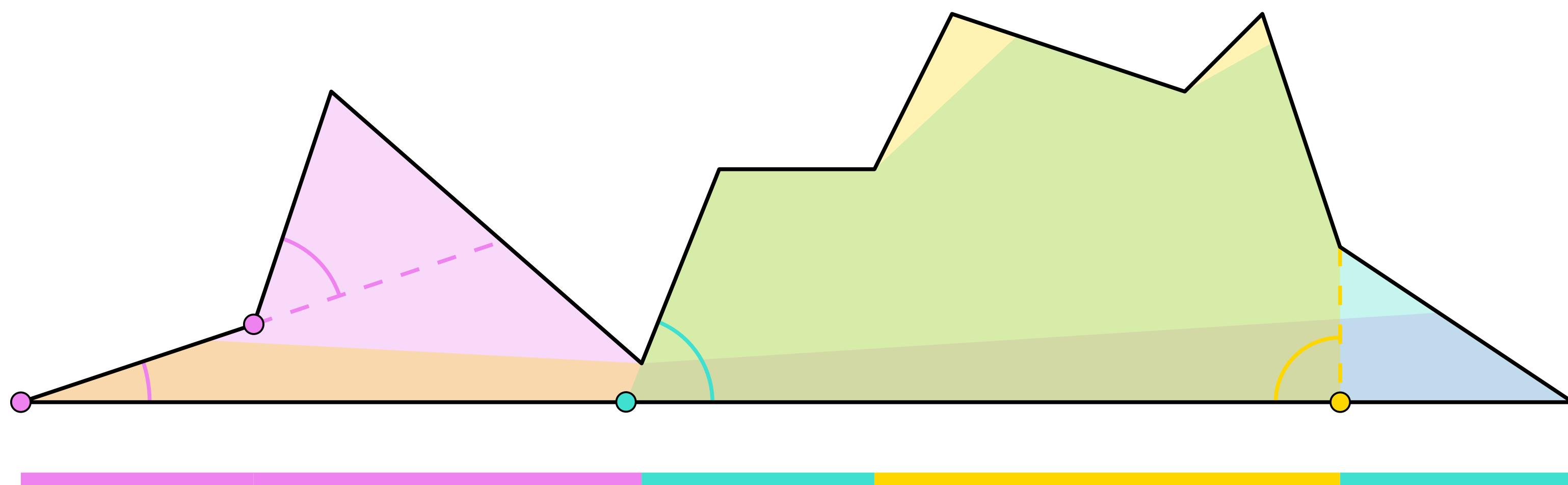
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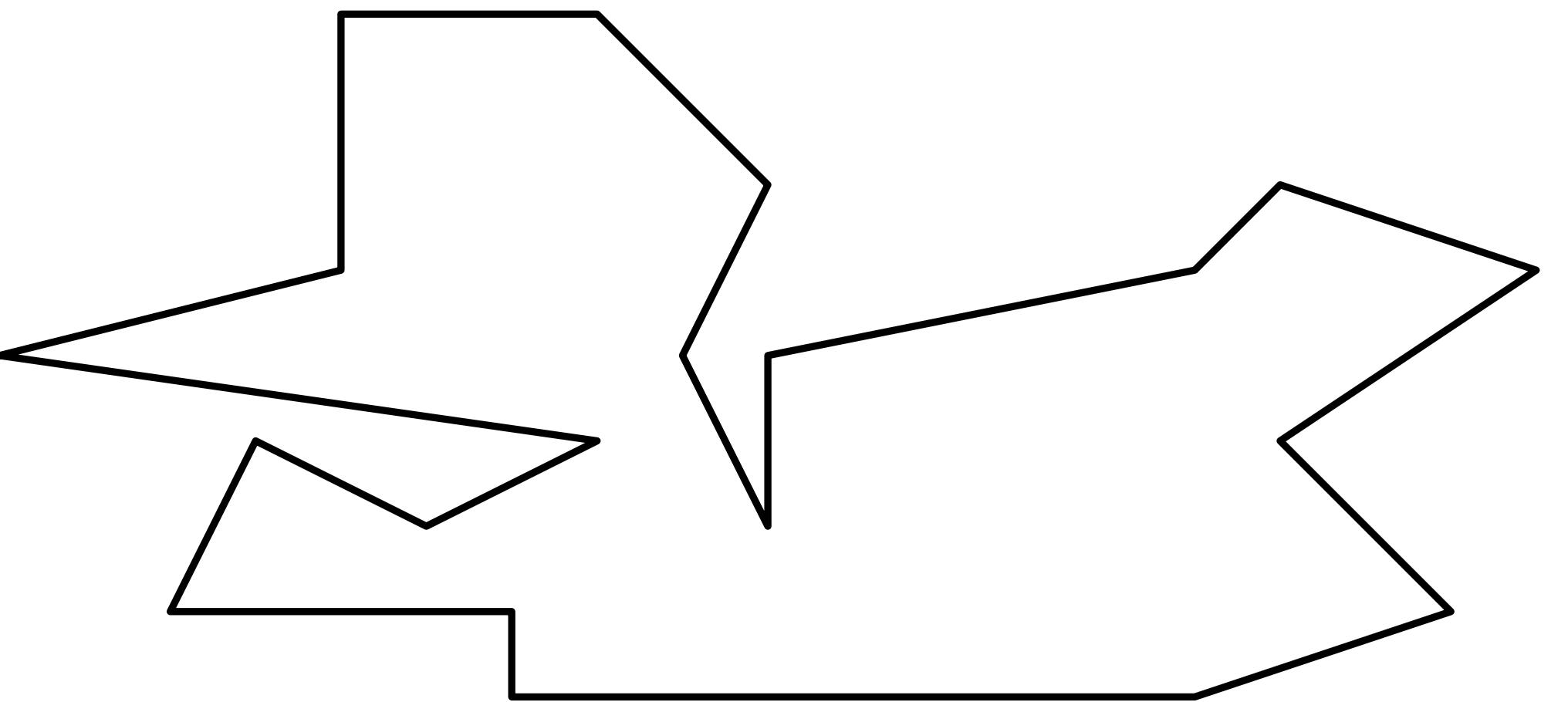
Equilateral Triangles

Histograms

Upper Bound for **Simple Polygons**

Duality

Conclusion



Introduction

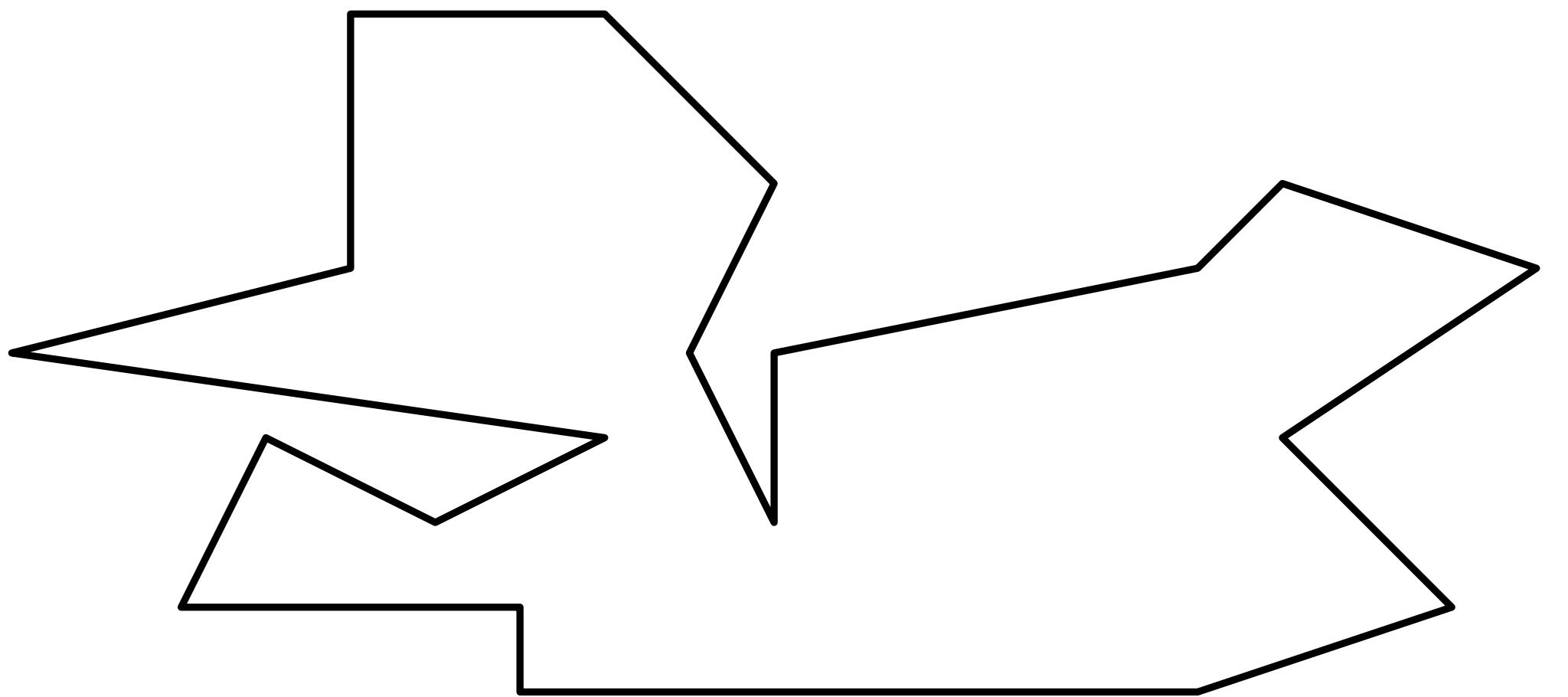
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$$(n - 2) \frac{\pi}{4}$$

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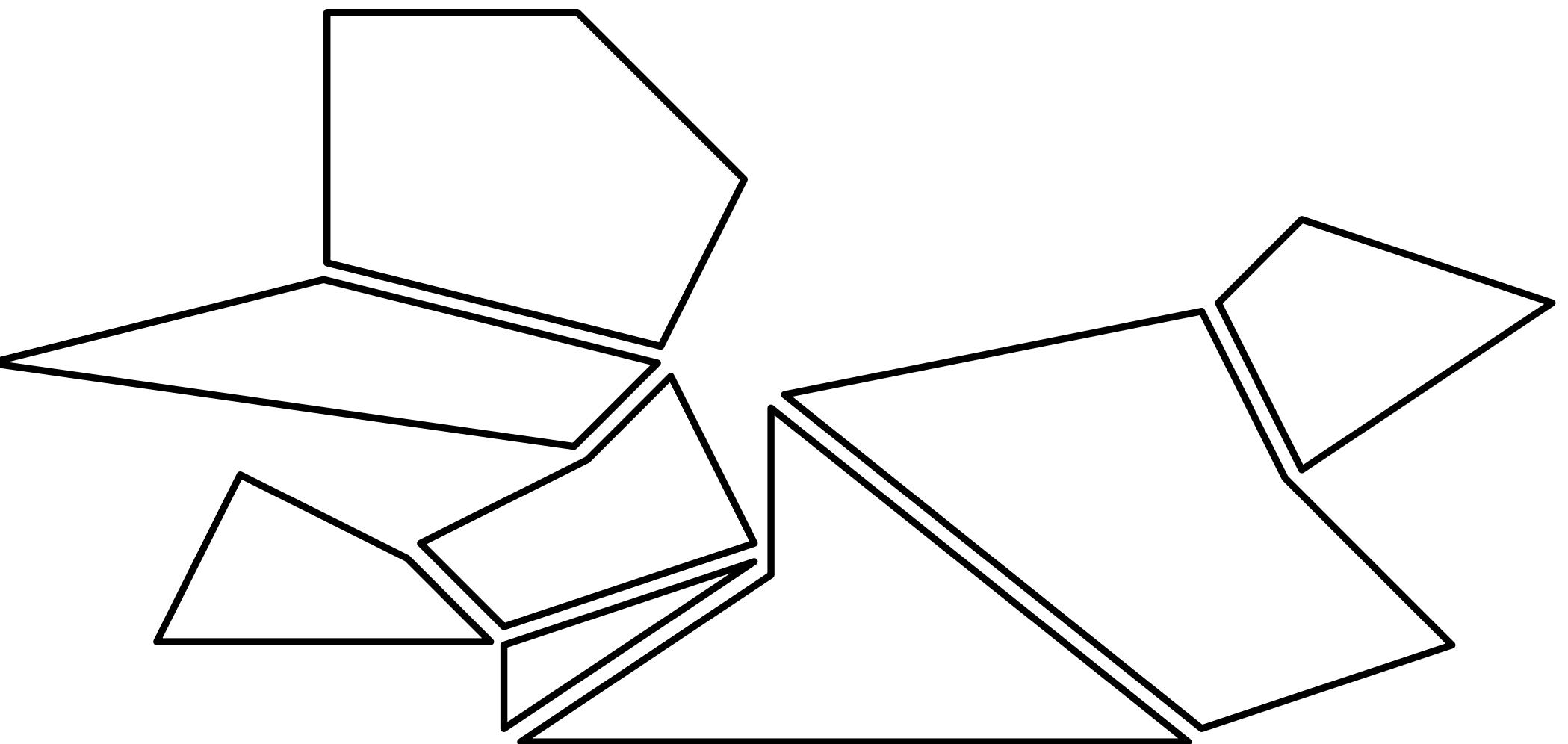
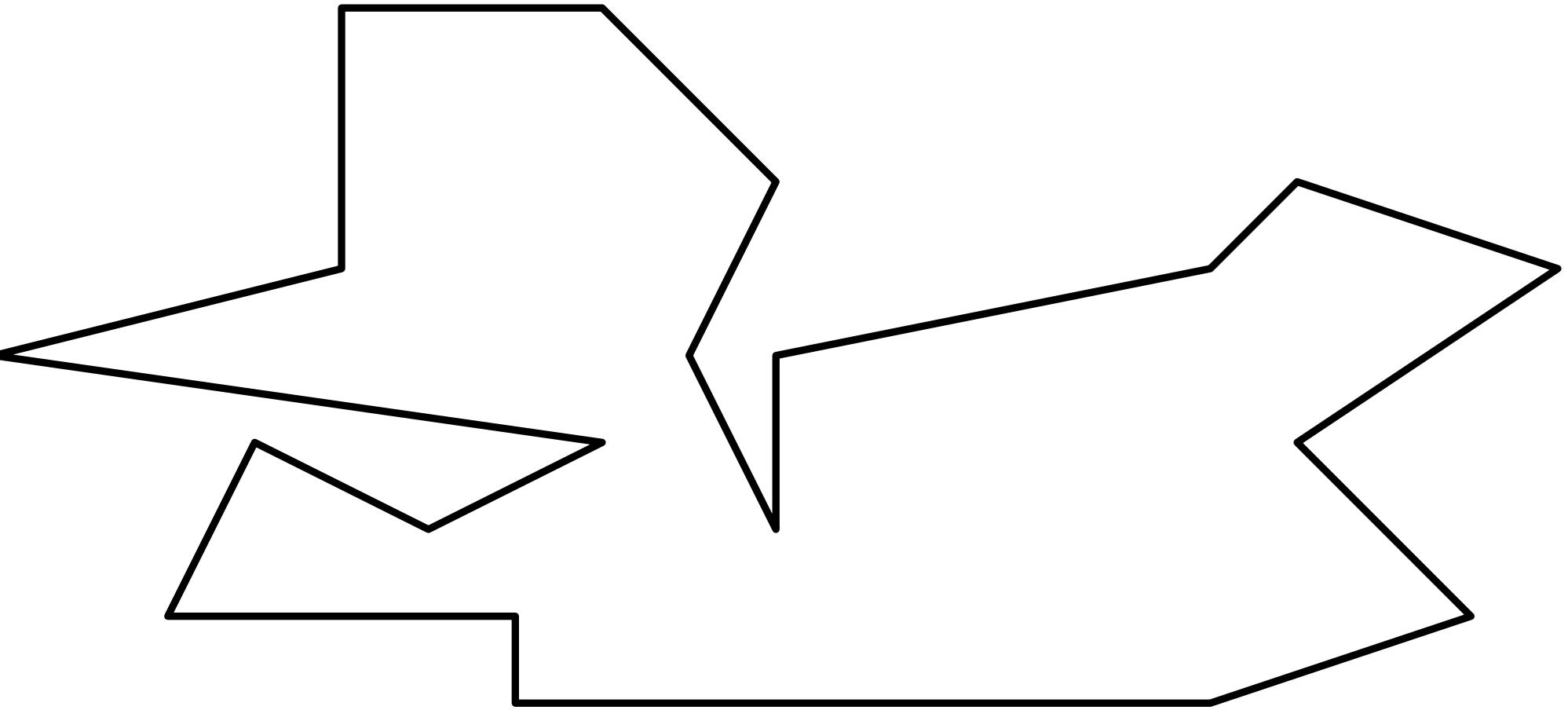
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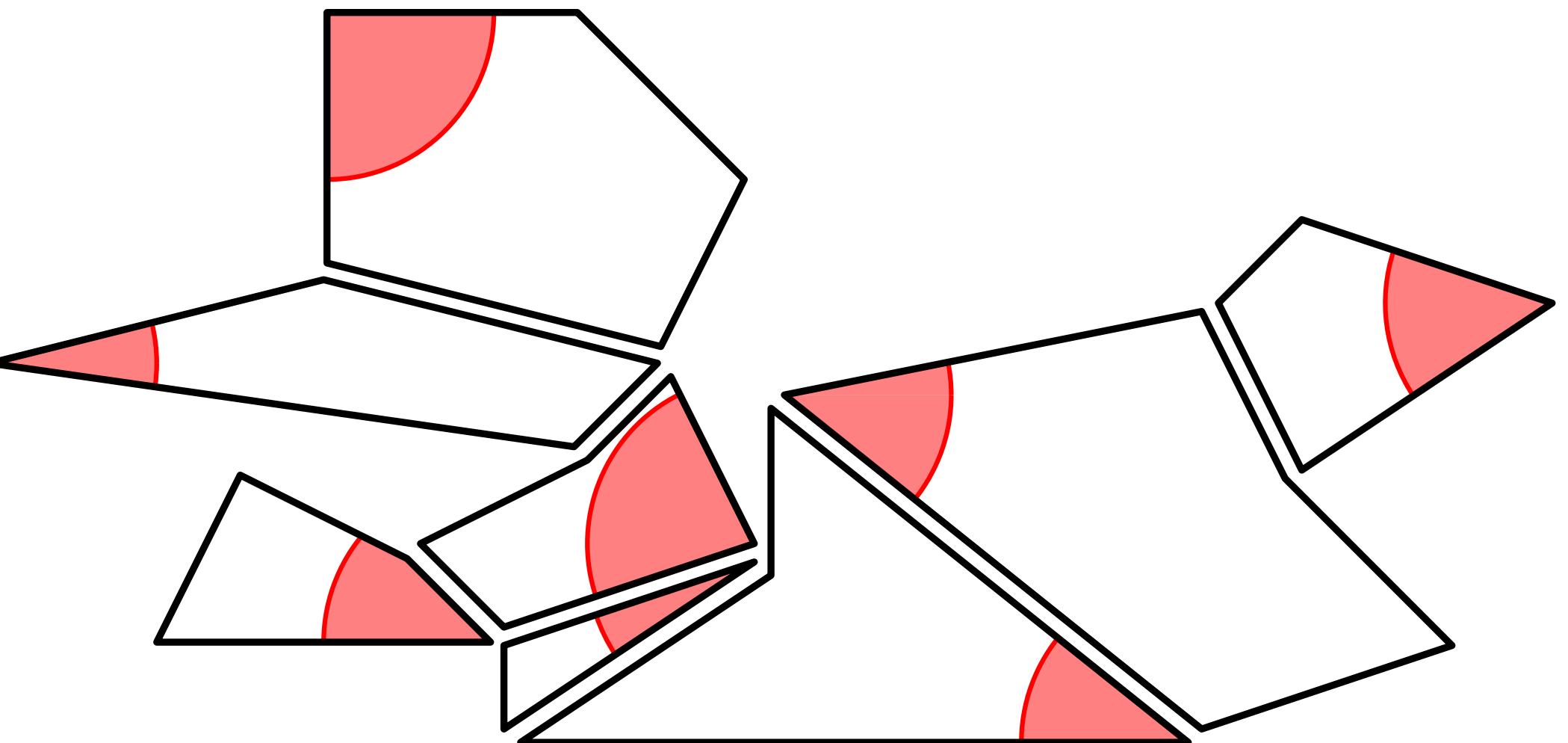
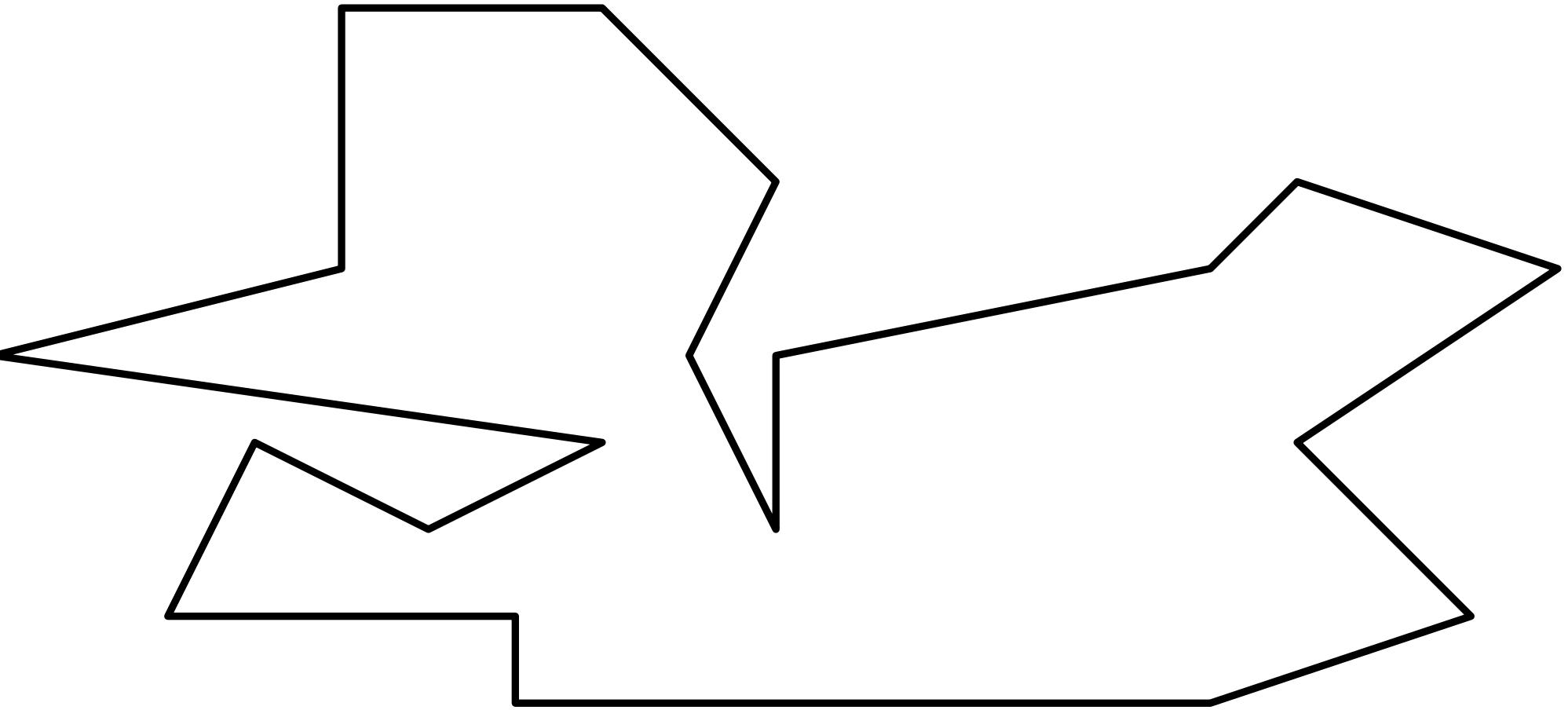
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Upper bound for simple polygons

Upper bound for $n \geq k + 2$ of

$$\left(1 + \frac{1}{k}\right) (n - 2) \frac{\pi}{6}$$

- Proven for $k = 2$
- Prove for $k = 10 \rightarrow UB \leq 1.1 LB$

Introduction

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Duality to independent circle packing

Conclusion

Duality to independent circle packing

Independent Circle Packing Problem

Instance: A polygon P

Wanted: A set of independent circles in P

Maximize: The minimum required angle to cover all circles

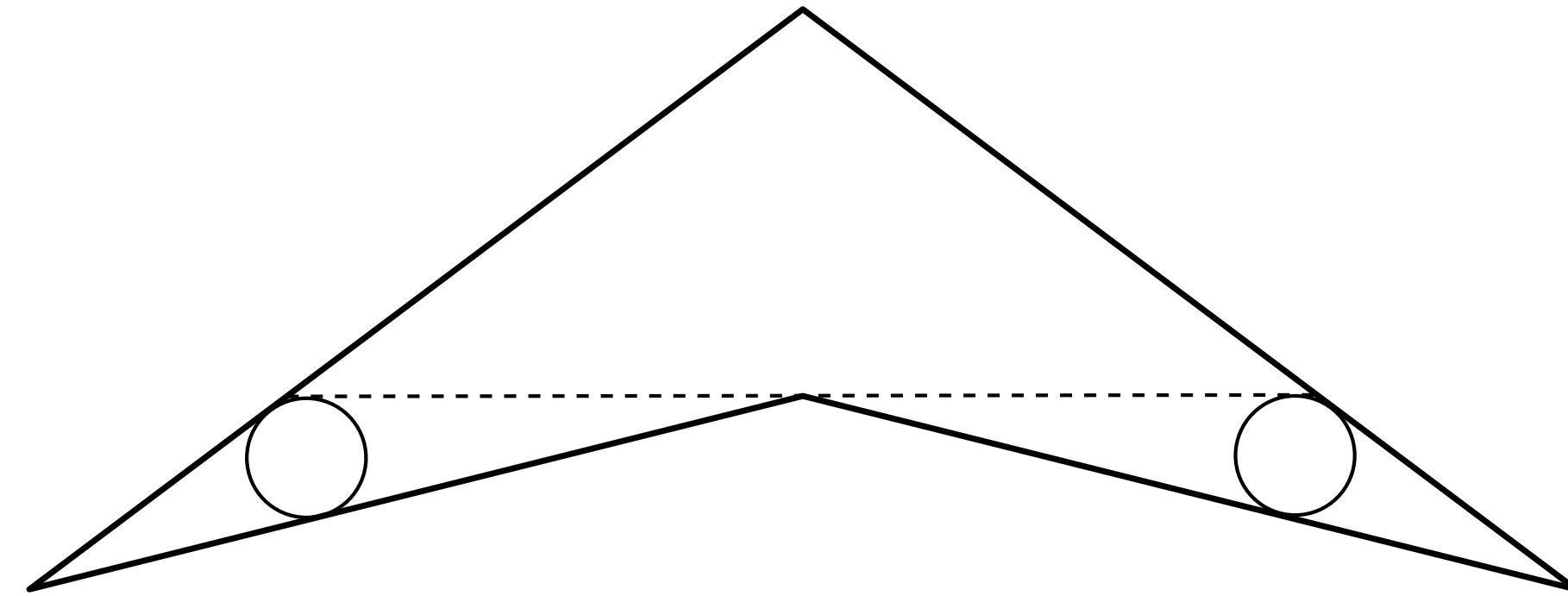
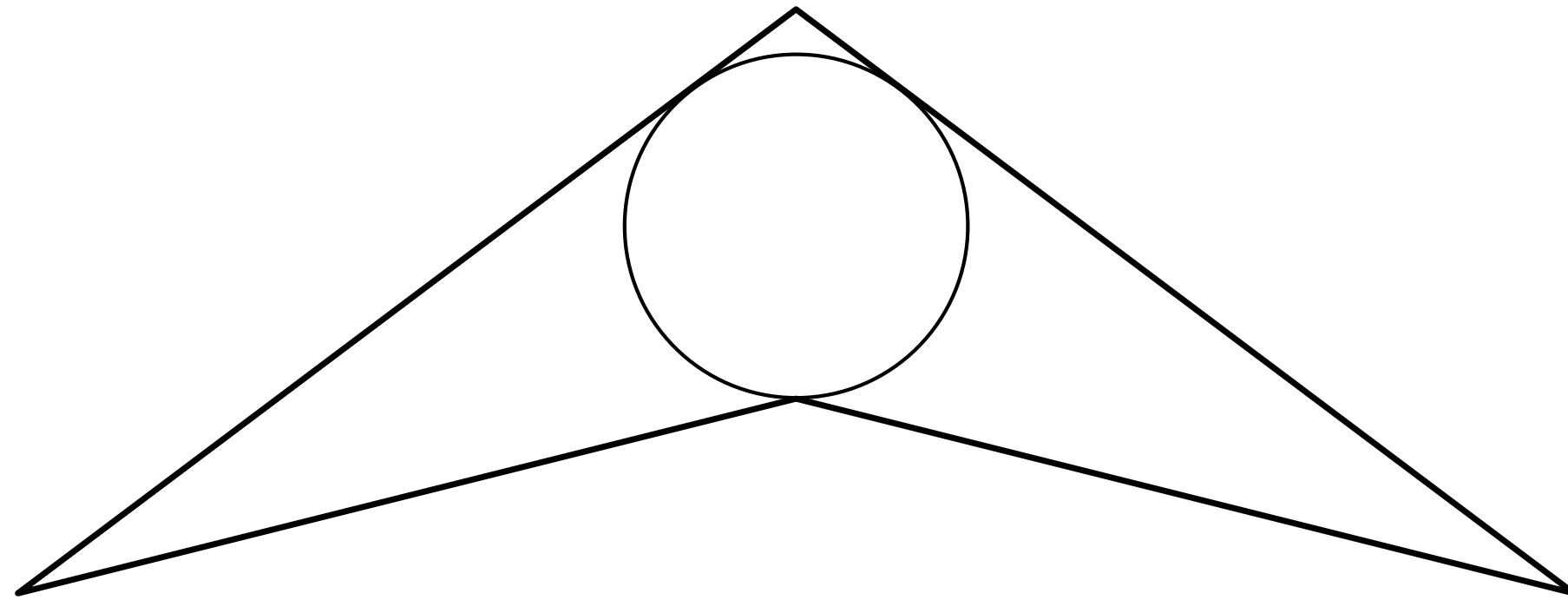
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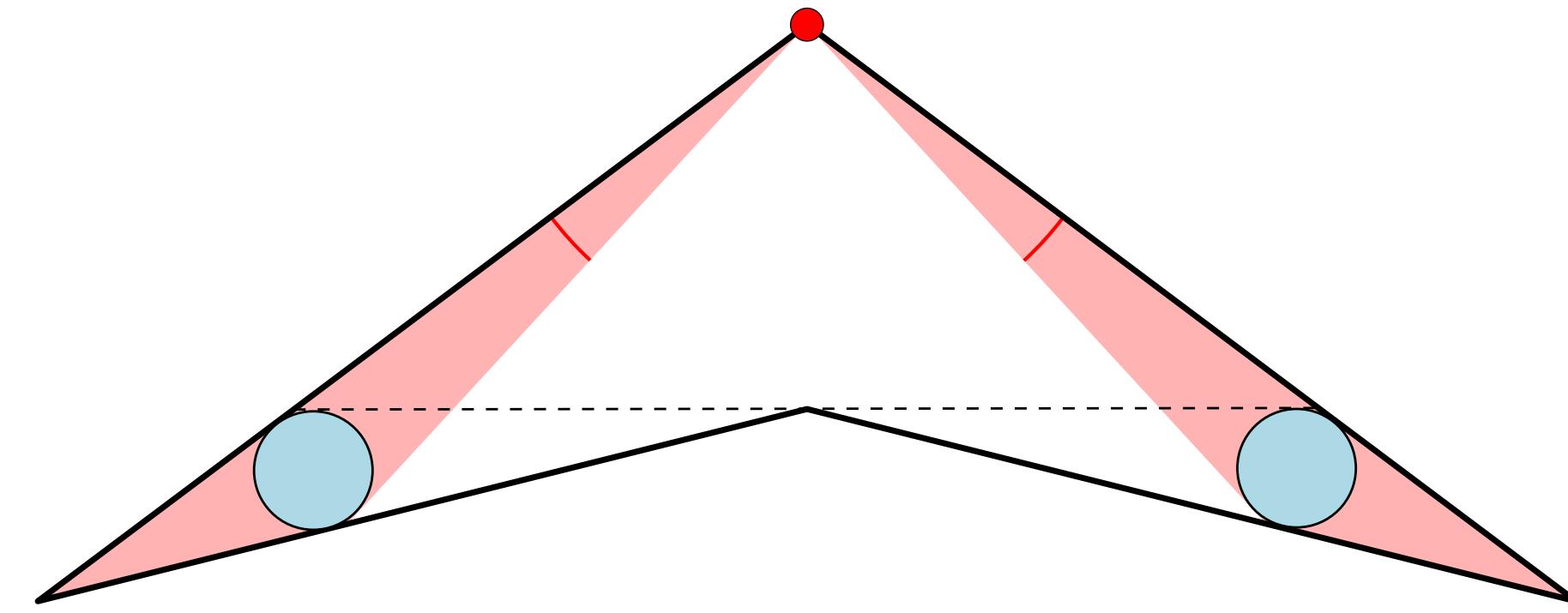
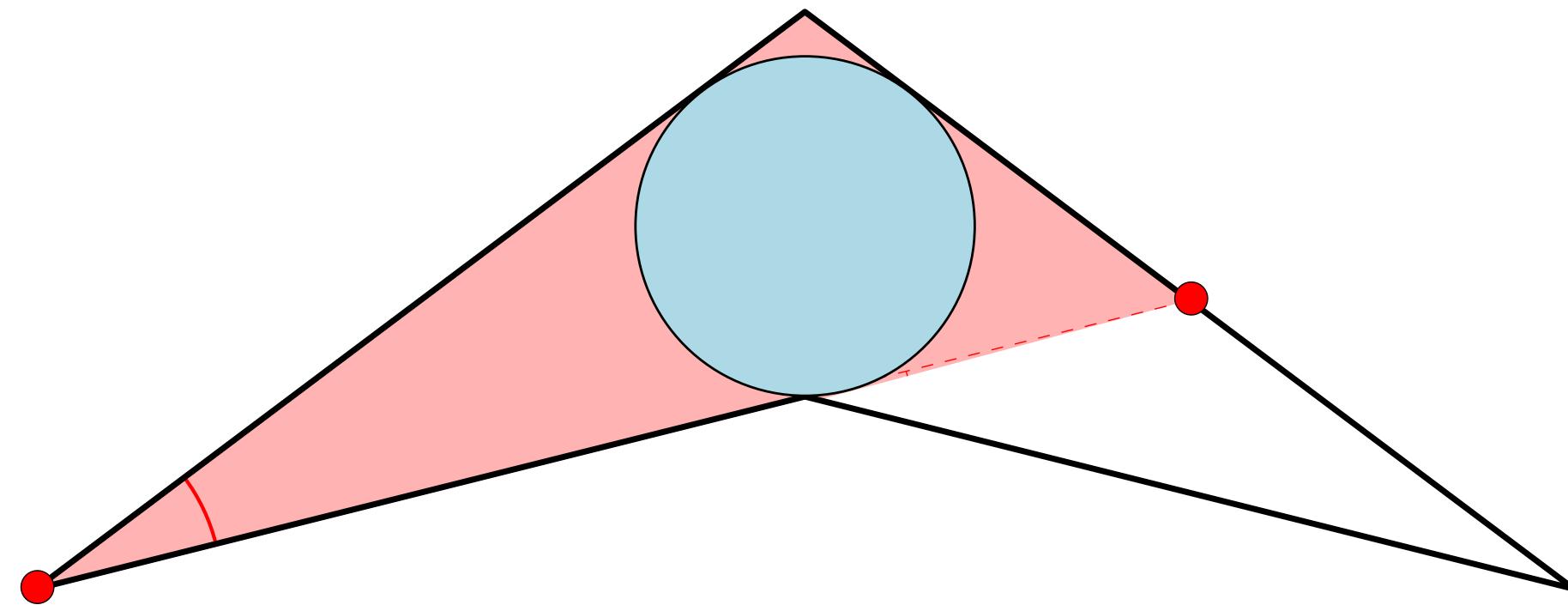
Duality to independent circle packing

Independent Circle Packing Problem

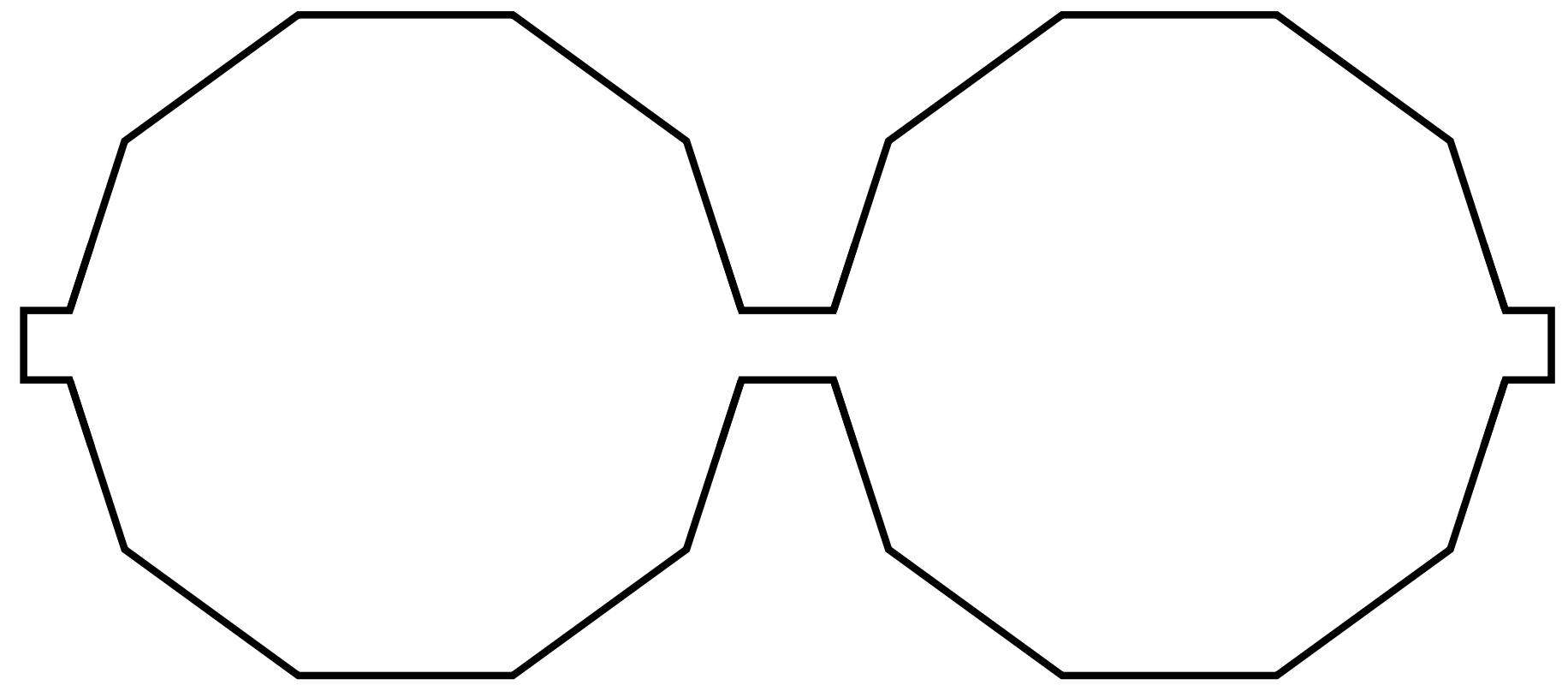
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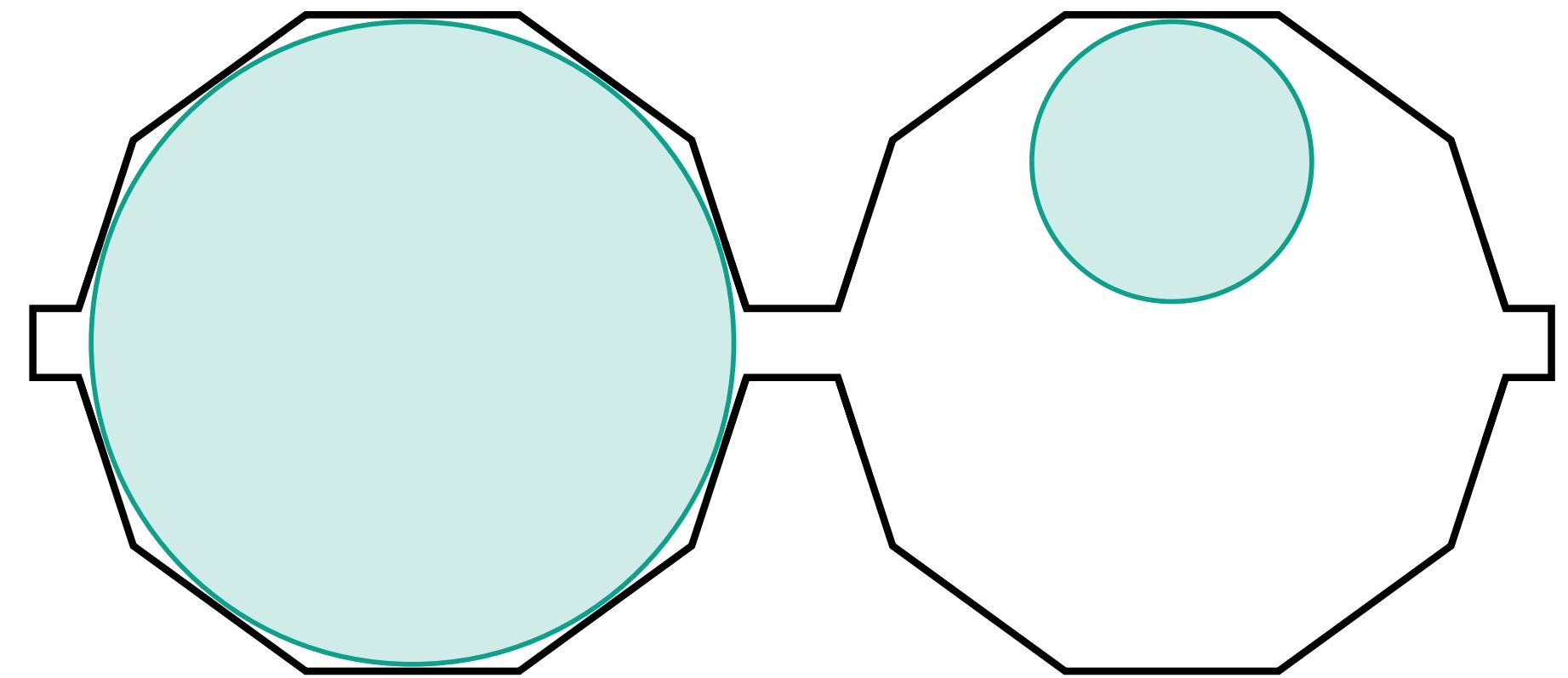
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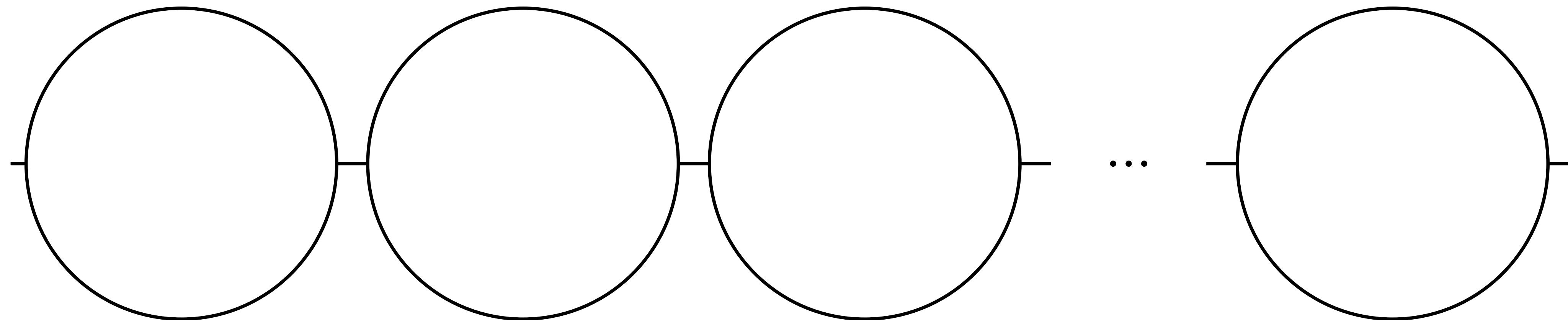
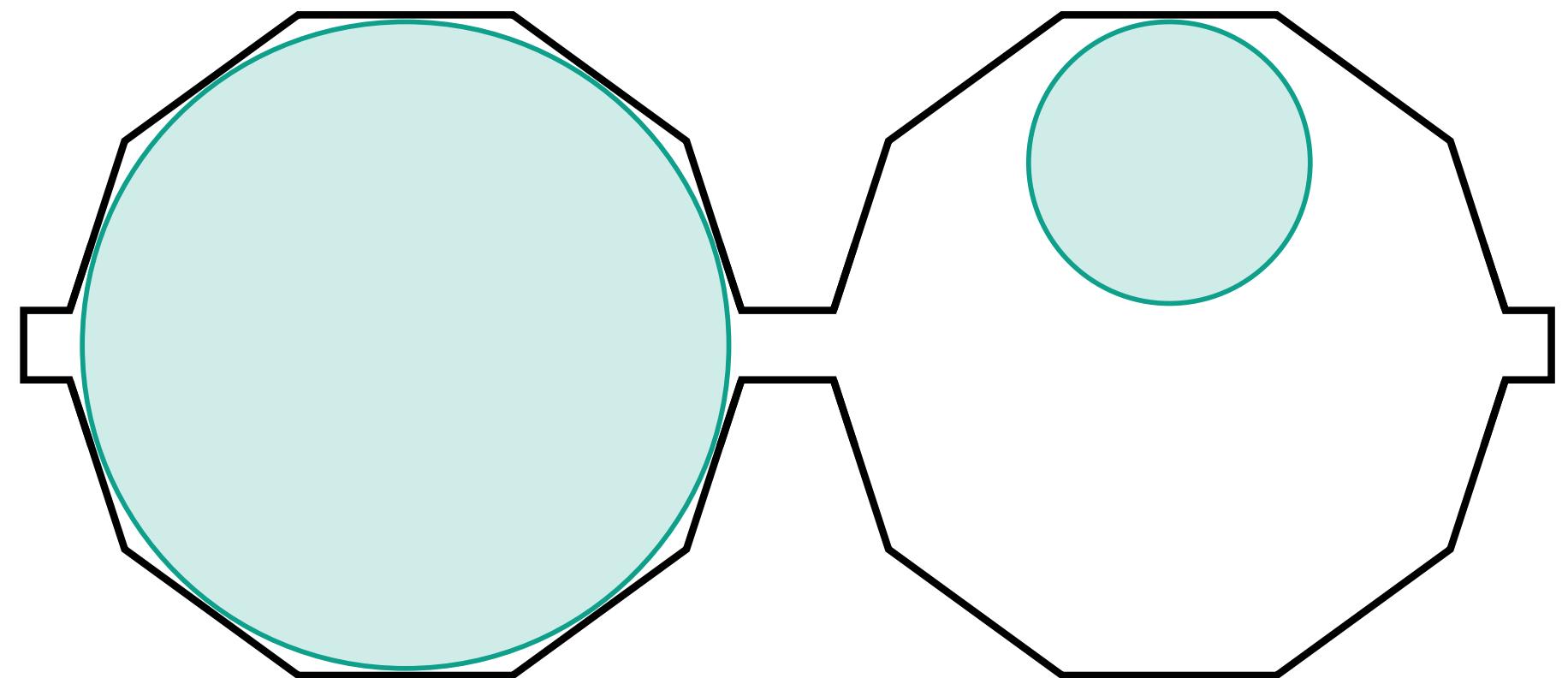
Duality Gap



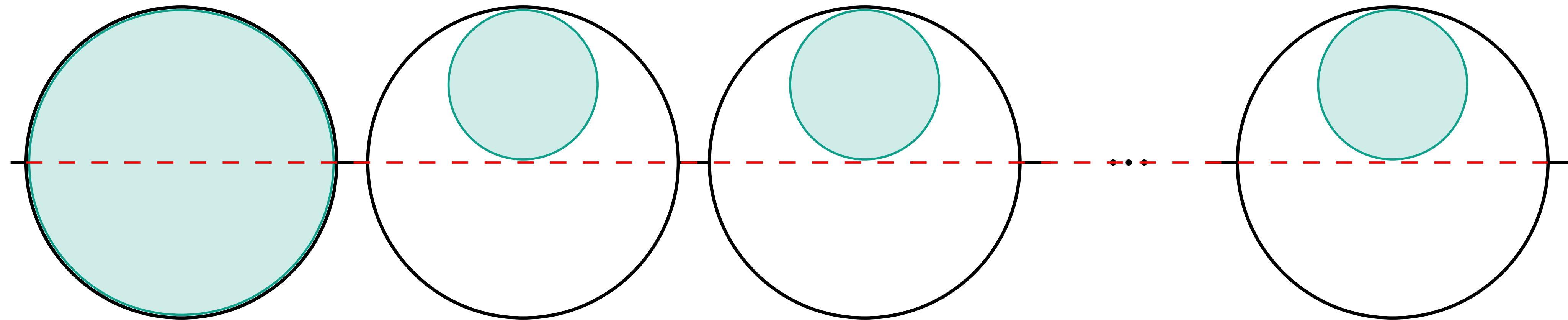
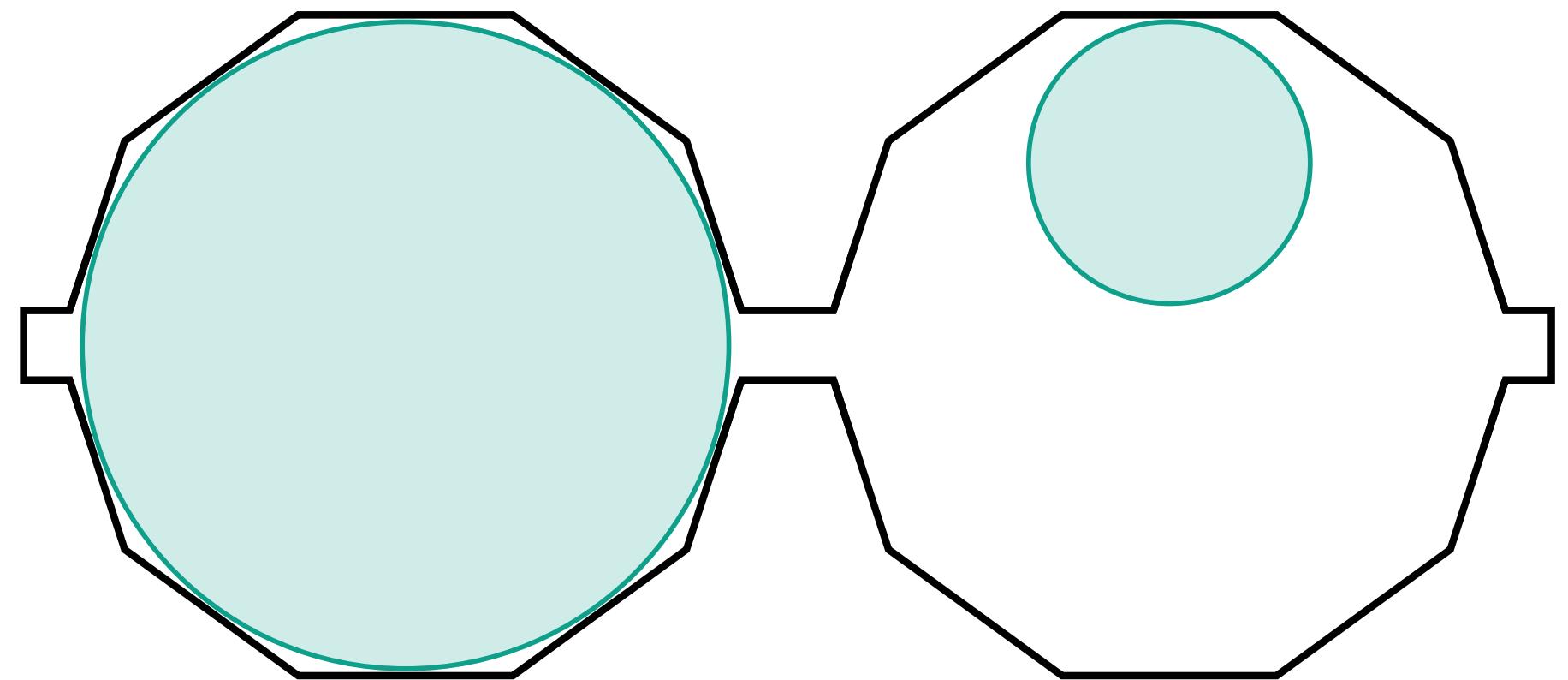
Duality Gap



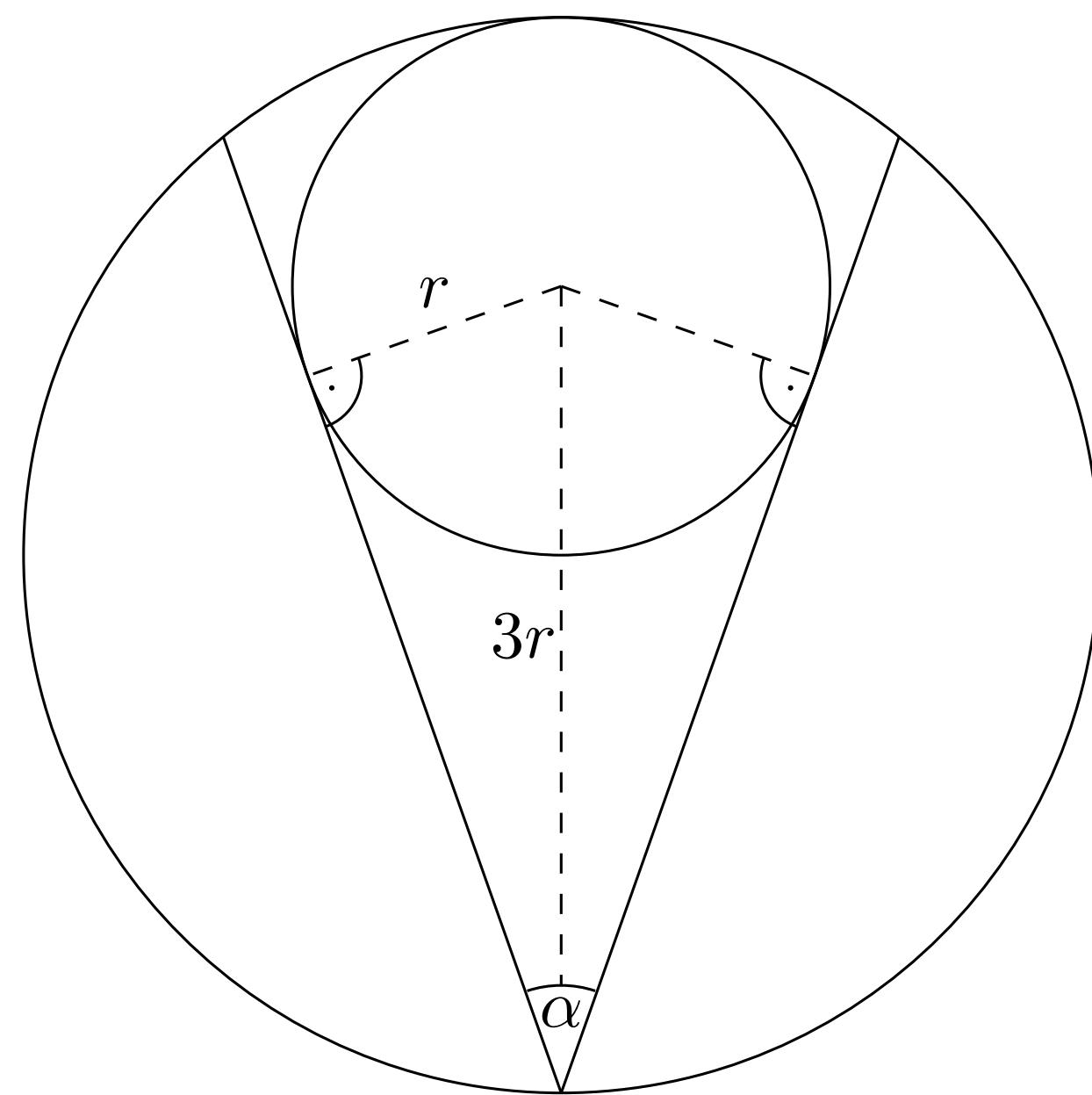
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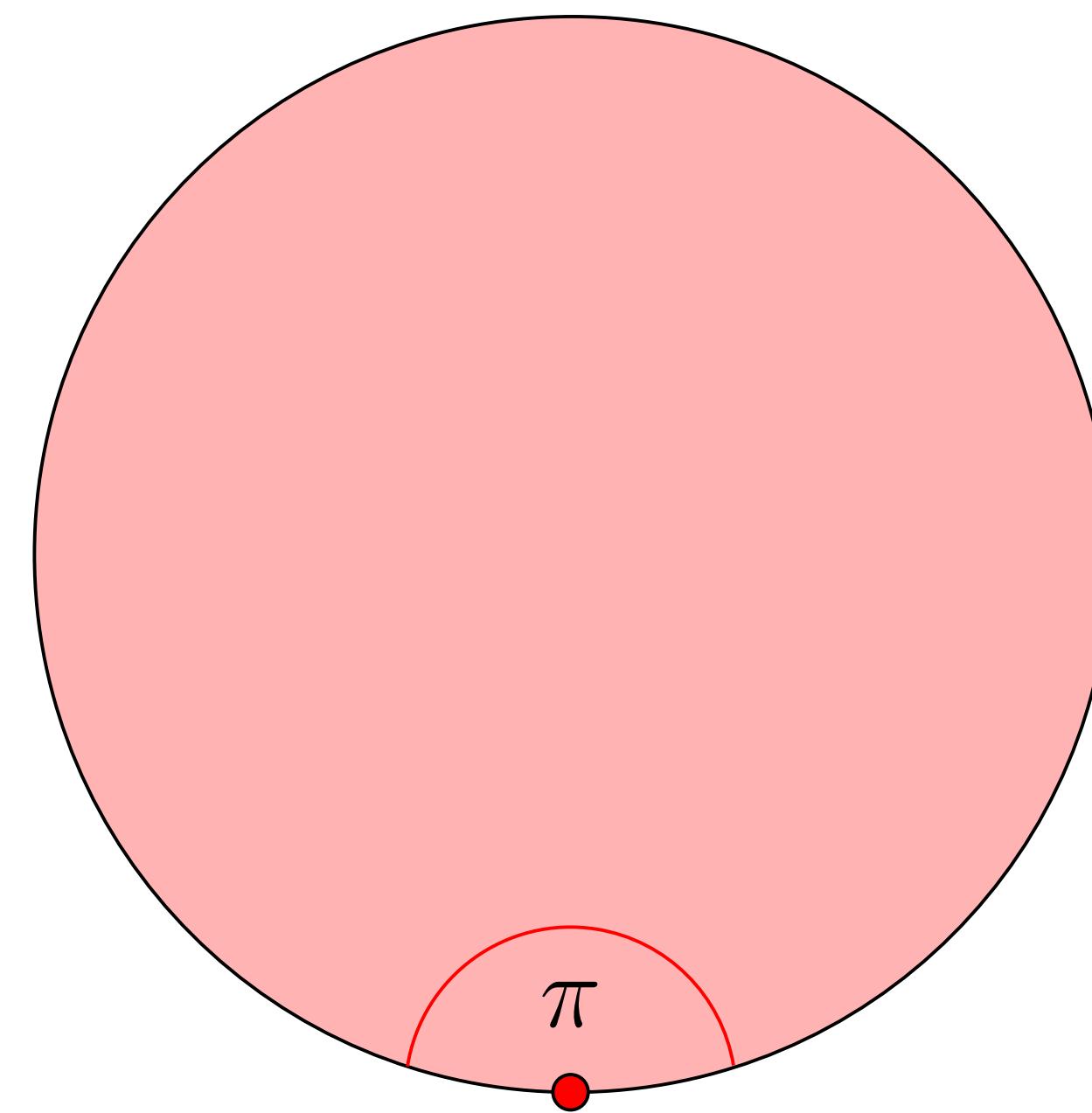
Duality Gap



Duality Gap



$$\text{ICP} = \alpha = 2 \arcsin \frac{1}{3}$$



$$\text{AAGP} = \pi$$

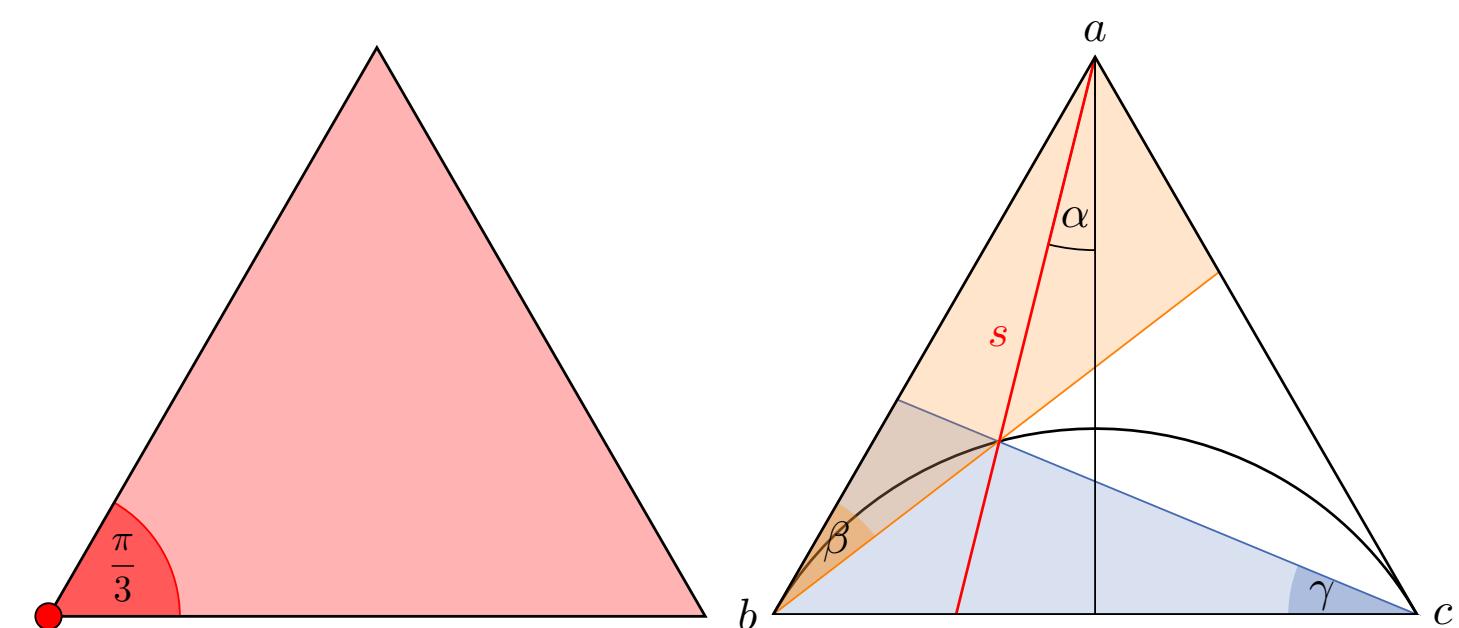
$$\frac{\text{AAGP}}{\text{ICP}} = \frac{\pi}{2 \arcsin \frac{1}{3}} \approx 4.622$$

Conclusion

Conclusion

An optimal covering of an equilateral triangle has an angle of

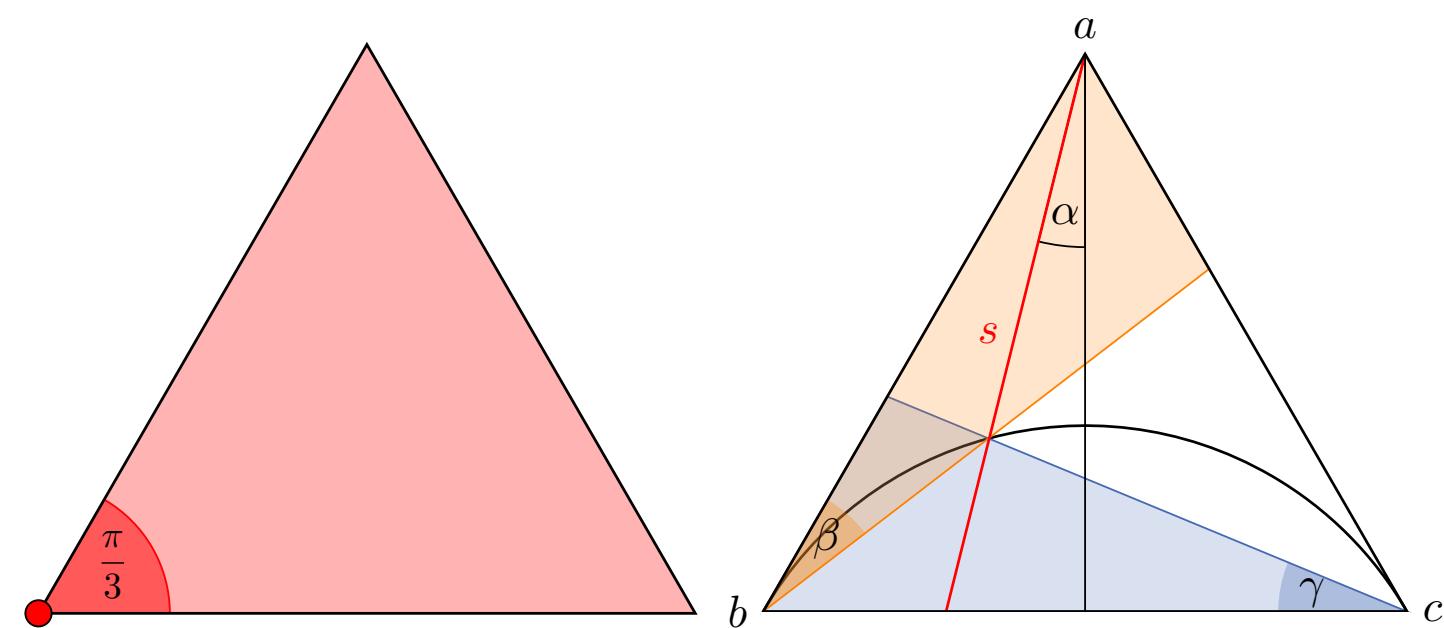
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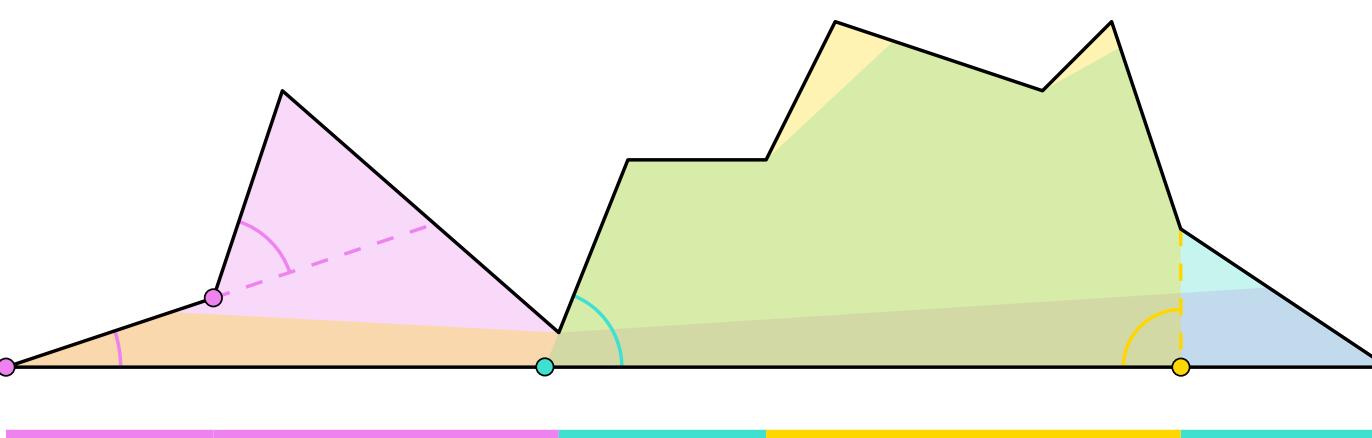
$$\frac{\pi}{3}$$



There is an upper bound of

$$(n - 1) \frac{\pi}{6}$$

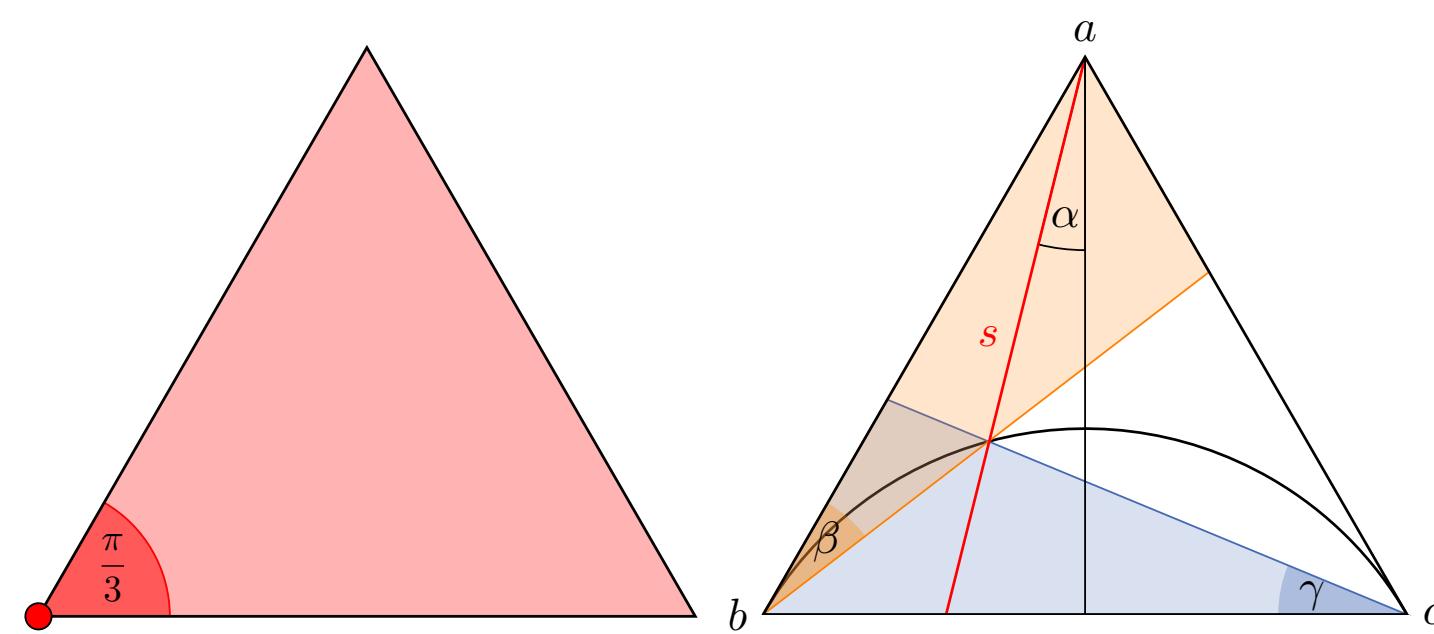
sufficient to cover any given histogram polygon.



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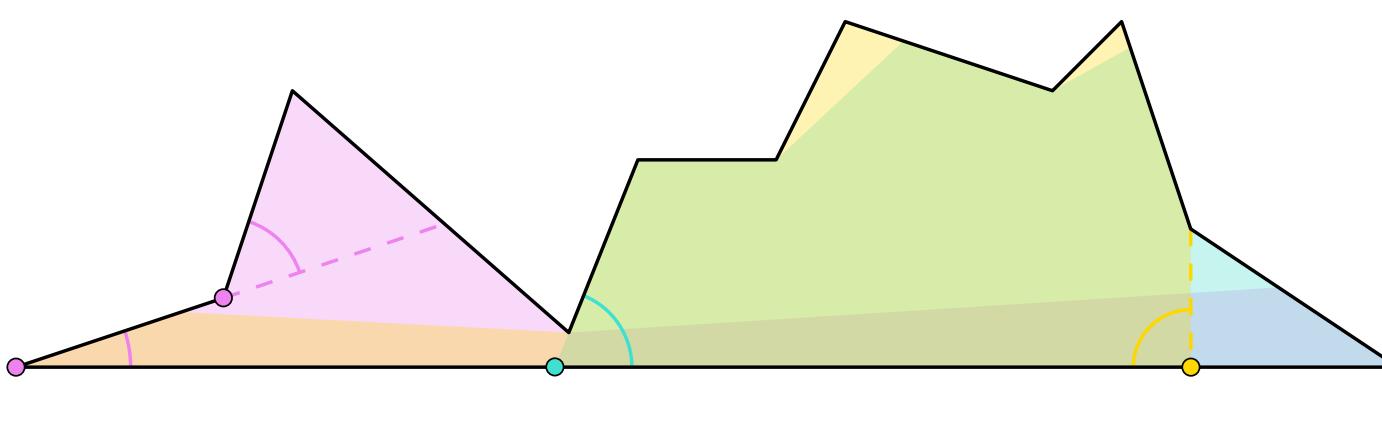
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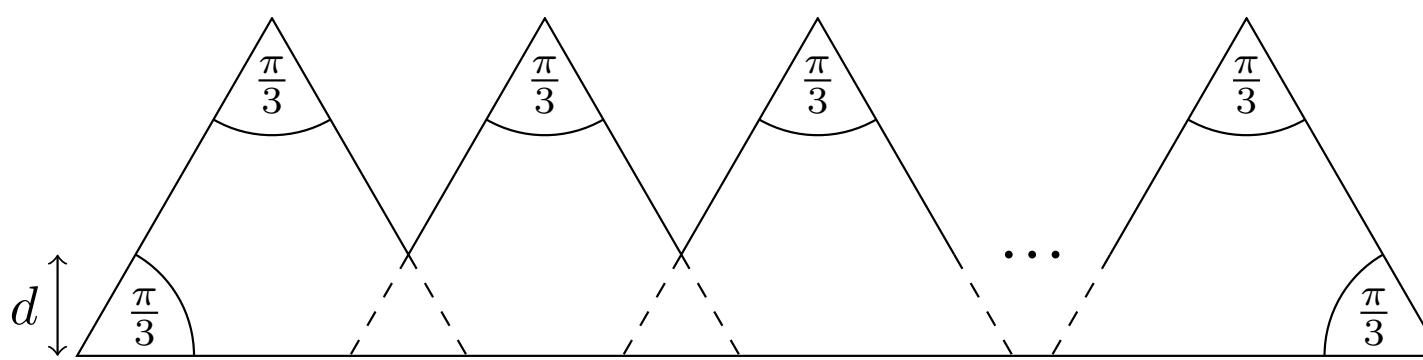
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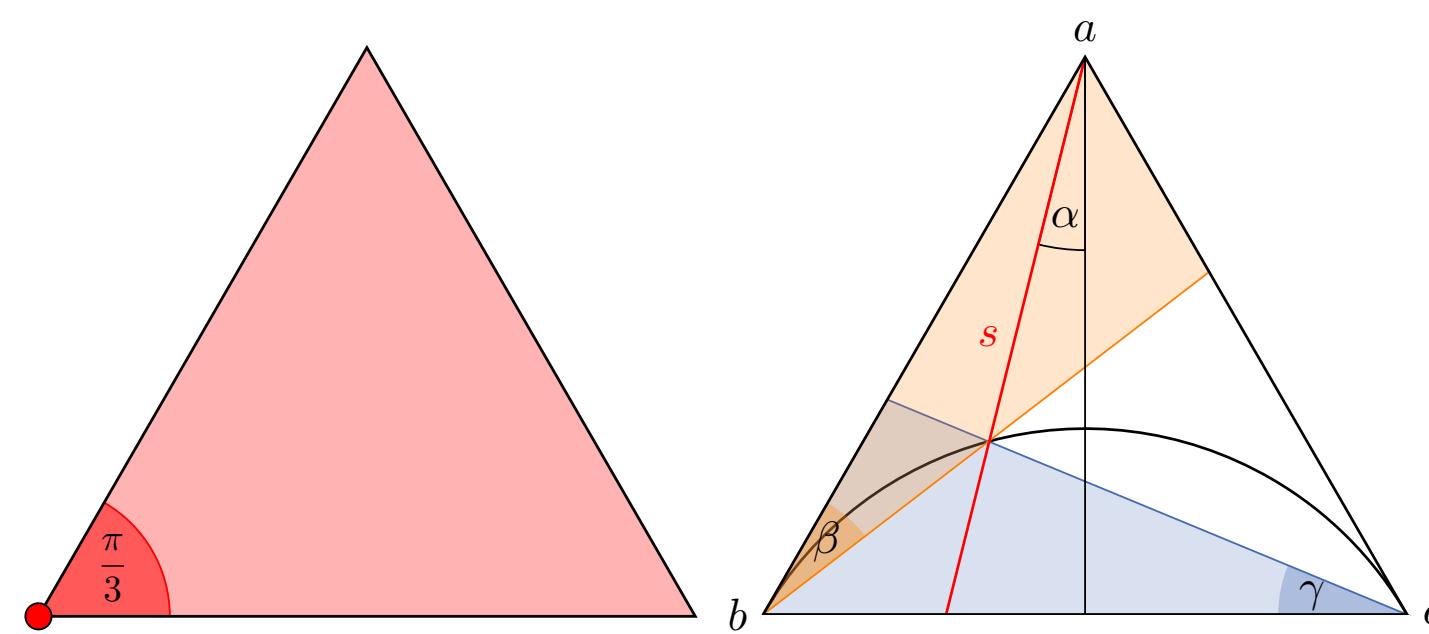
This bound is worst case optimal.



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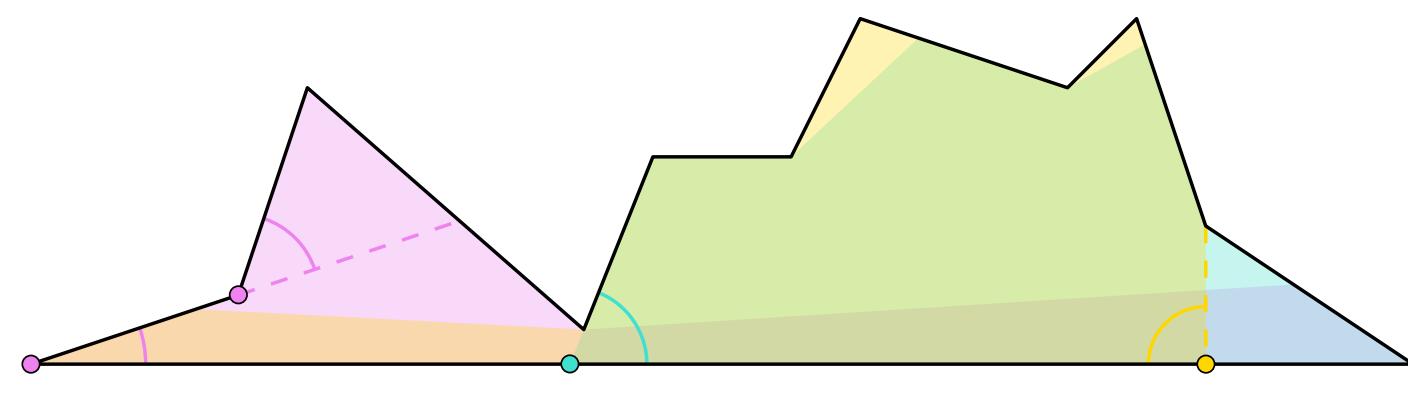
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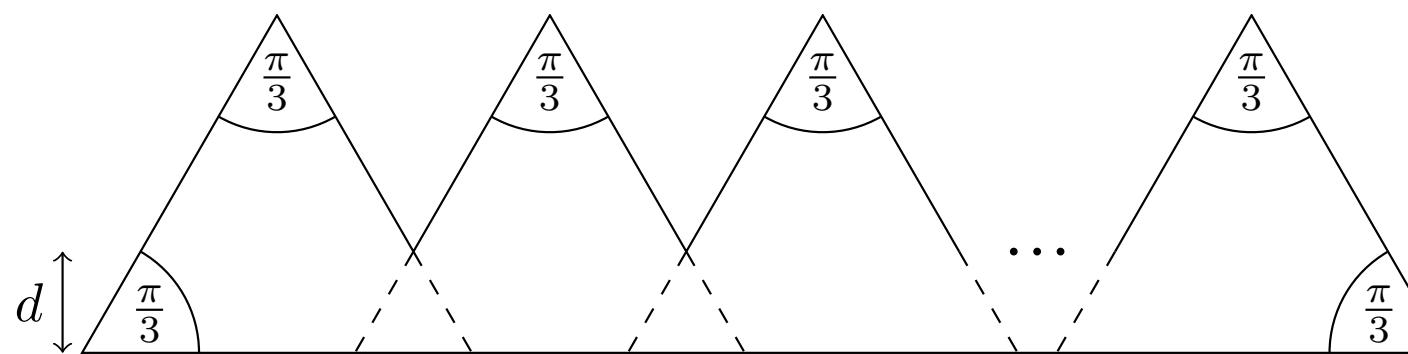
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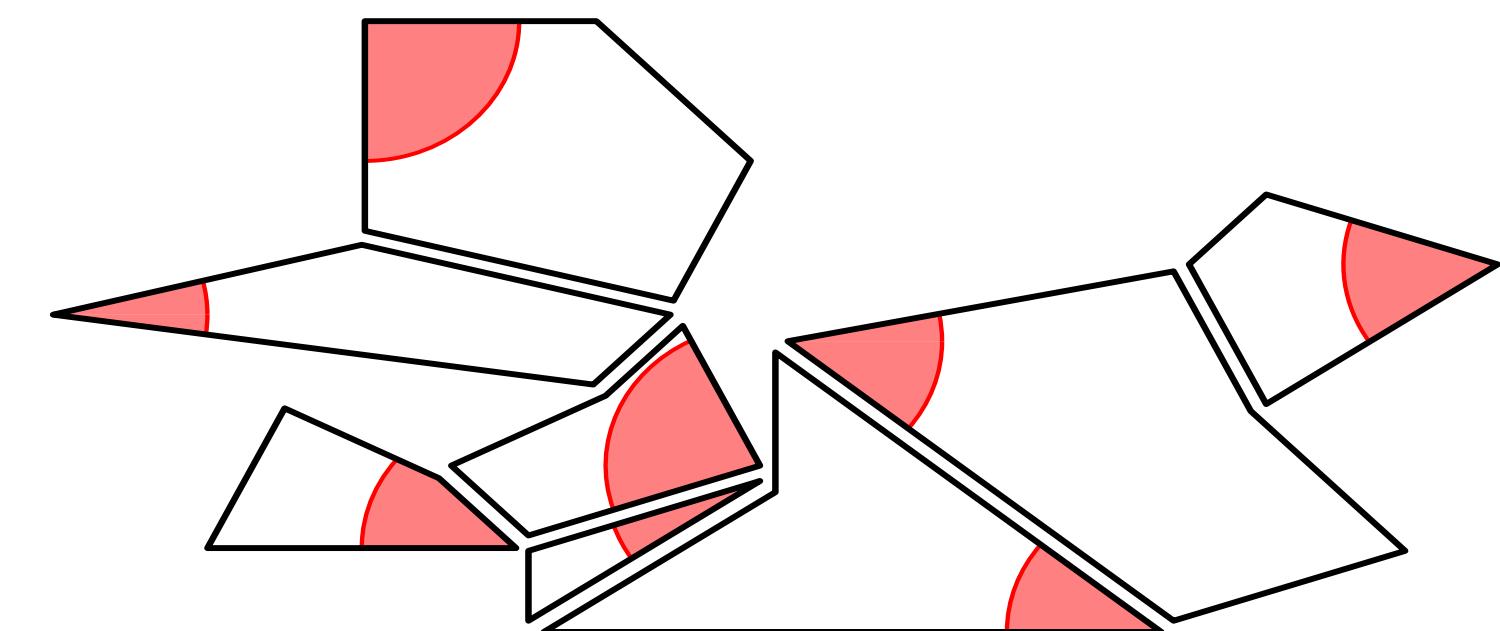
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For simple polygons, we presented an upper bound of

$$(n - 2) \frac{\pi}{4}$$

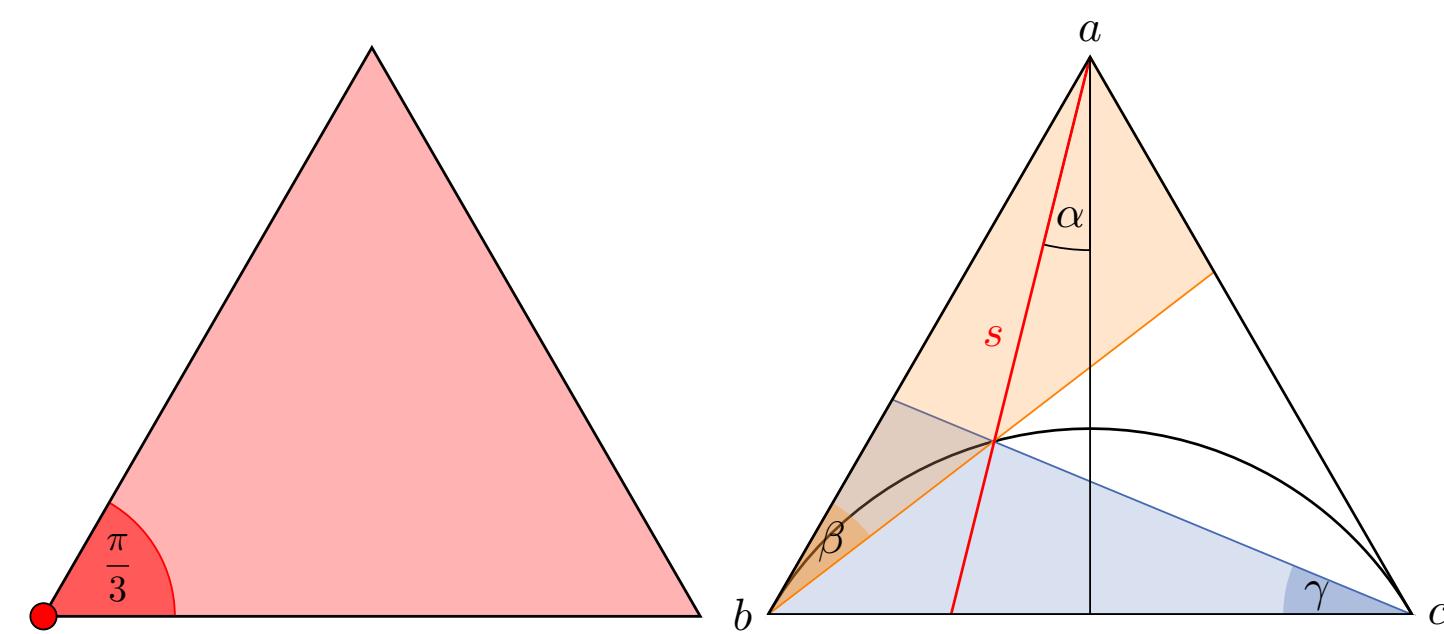
which may be reduced by proving a tight bound for polygons with a limited size.



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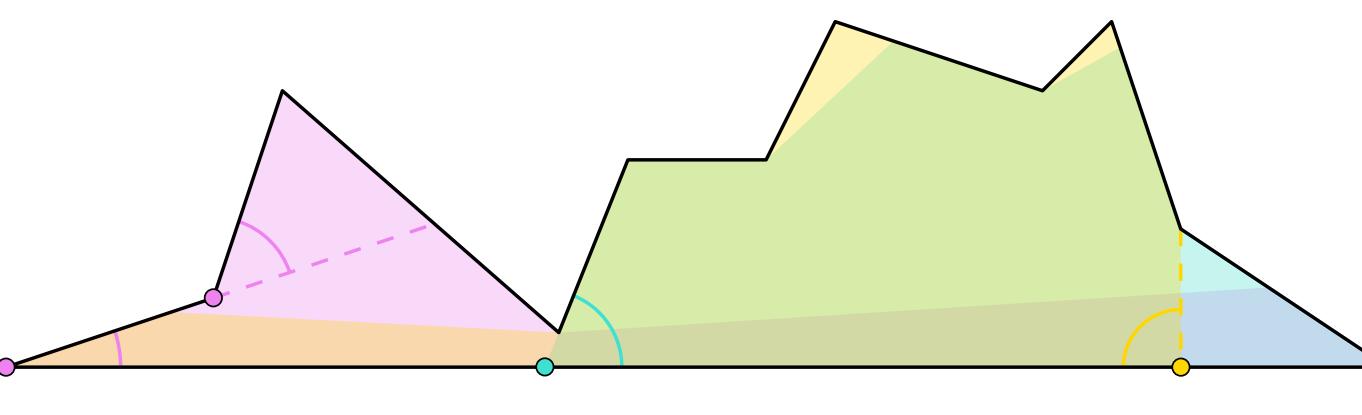
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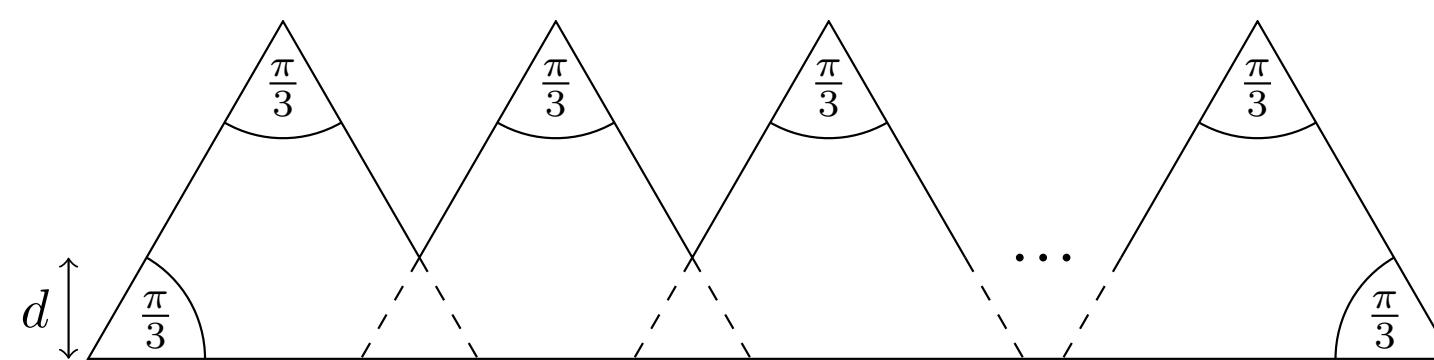
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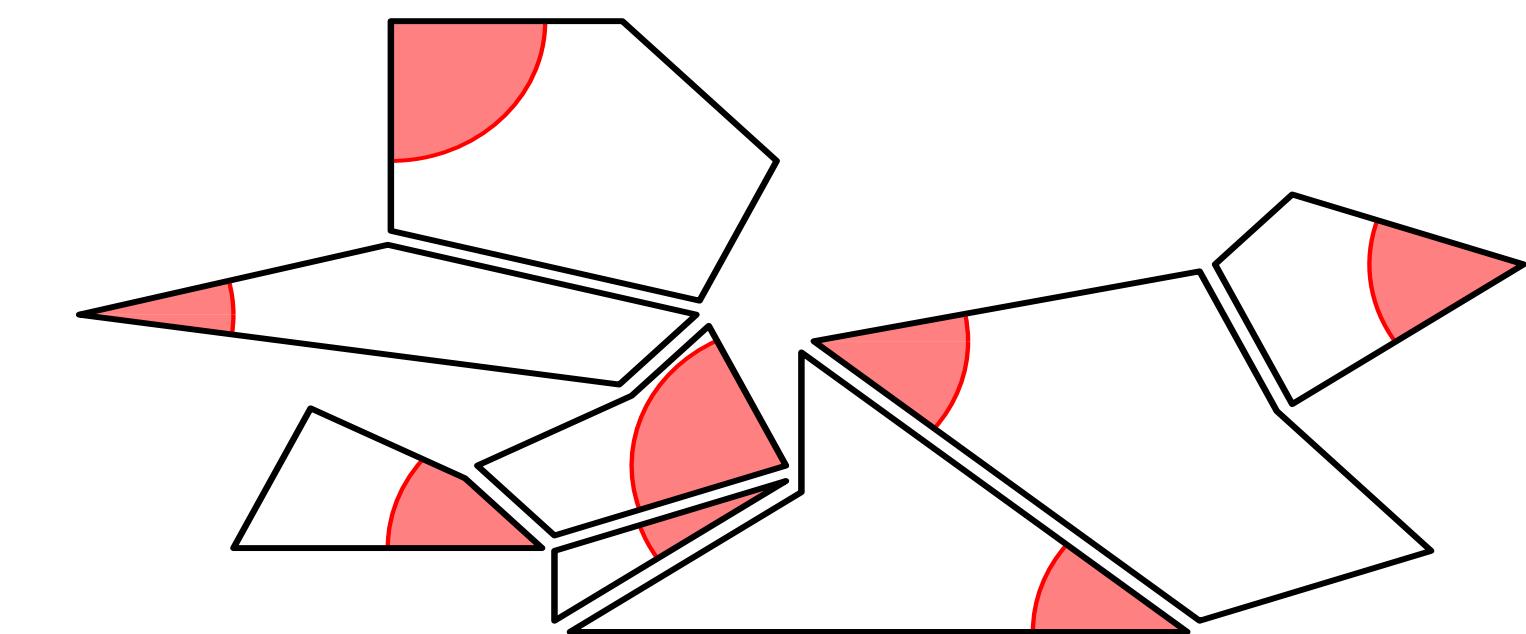
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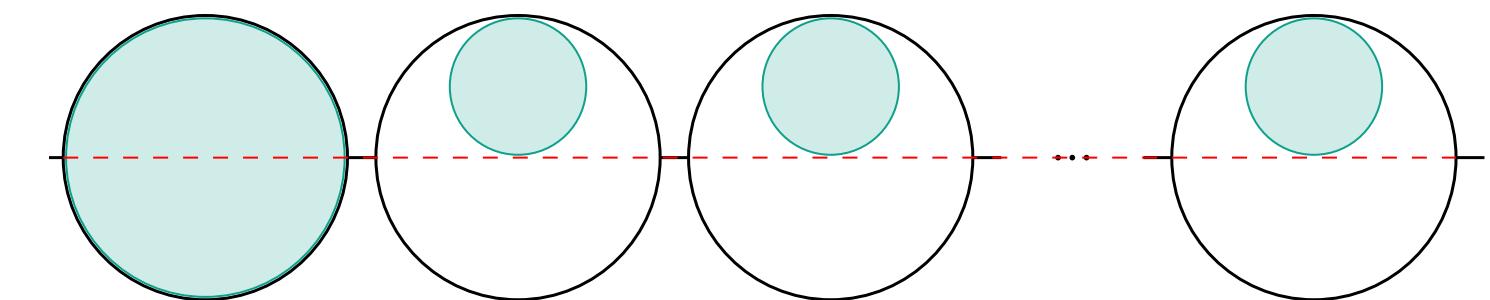


A dual problem deals with independent circle packing.

We determined a duality gap of

$$\approx 4.662$$

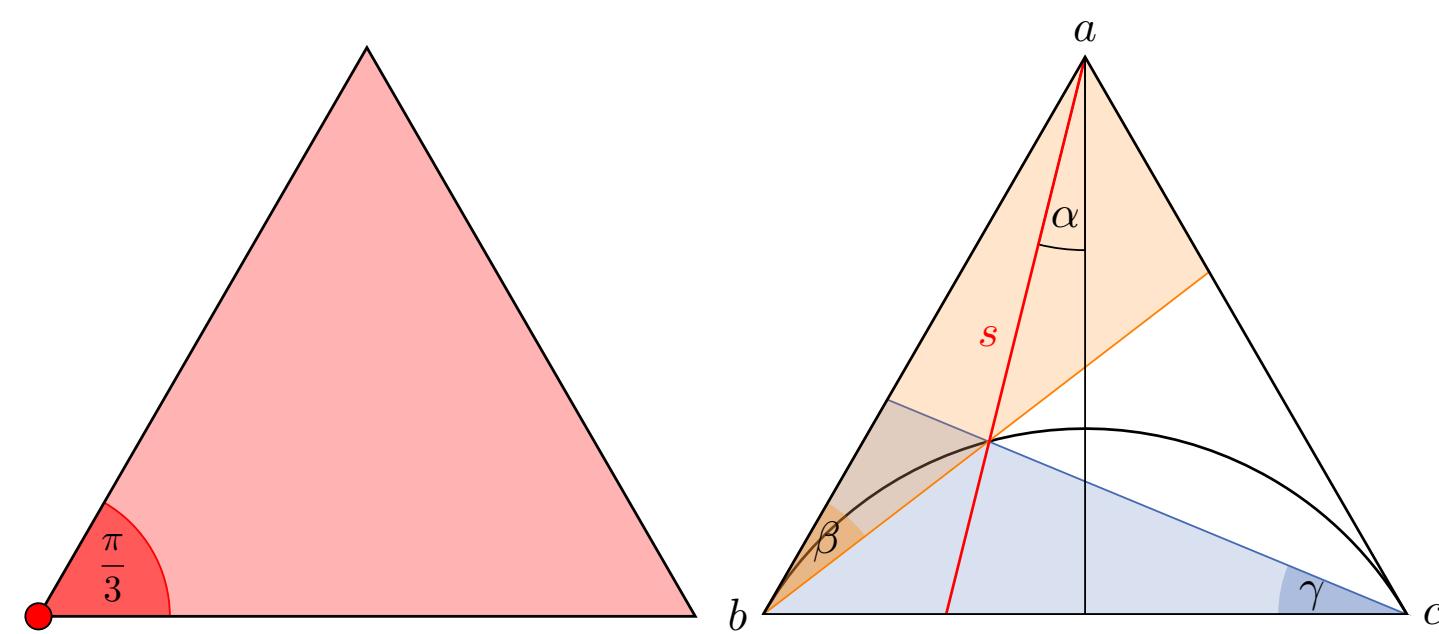
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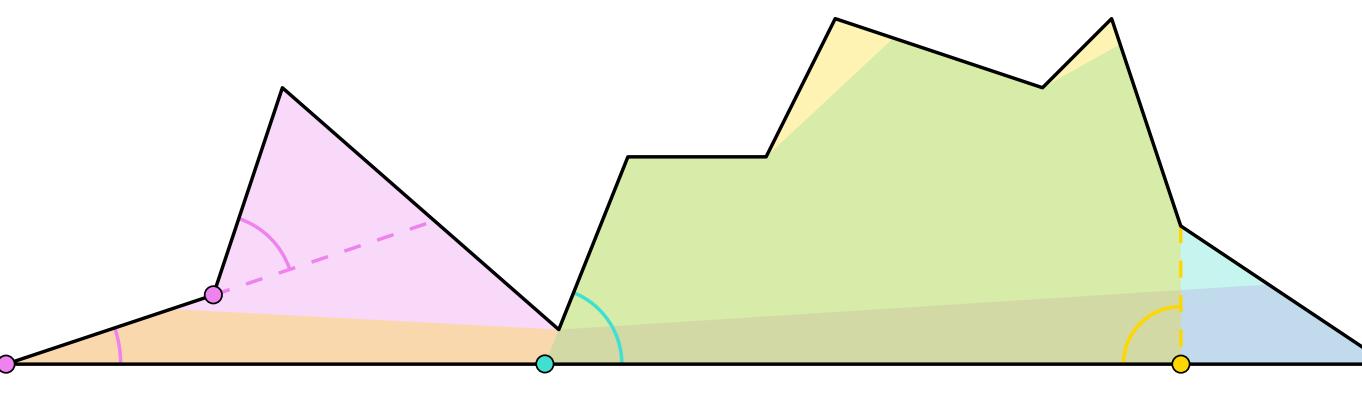
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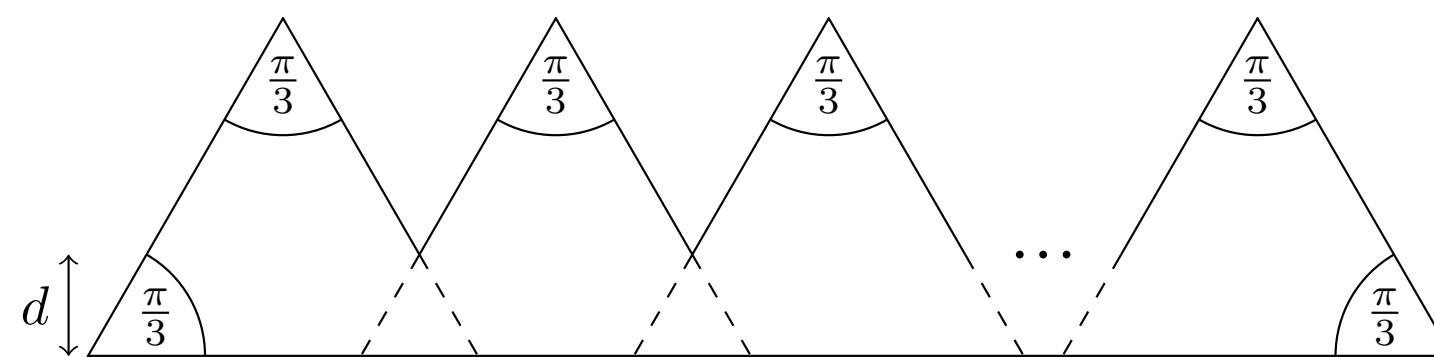
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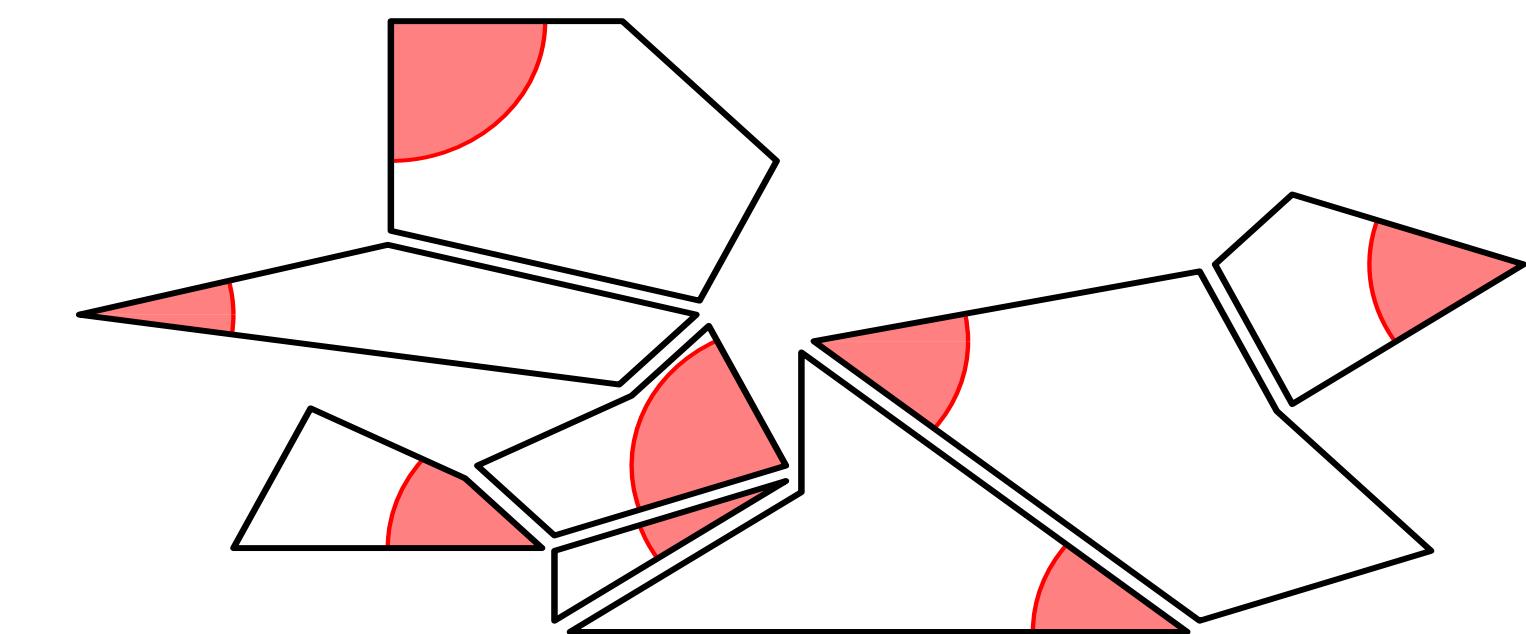
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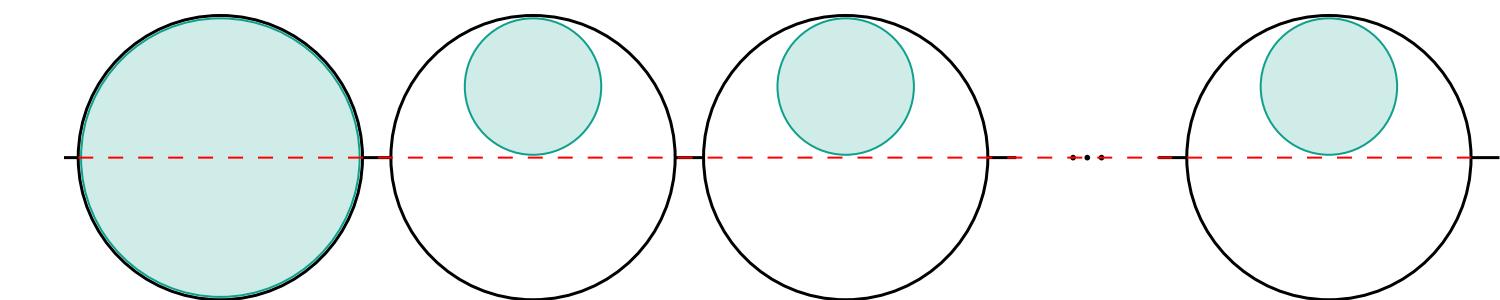


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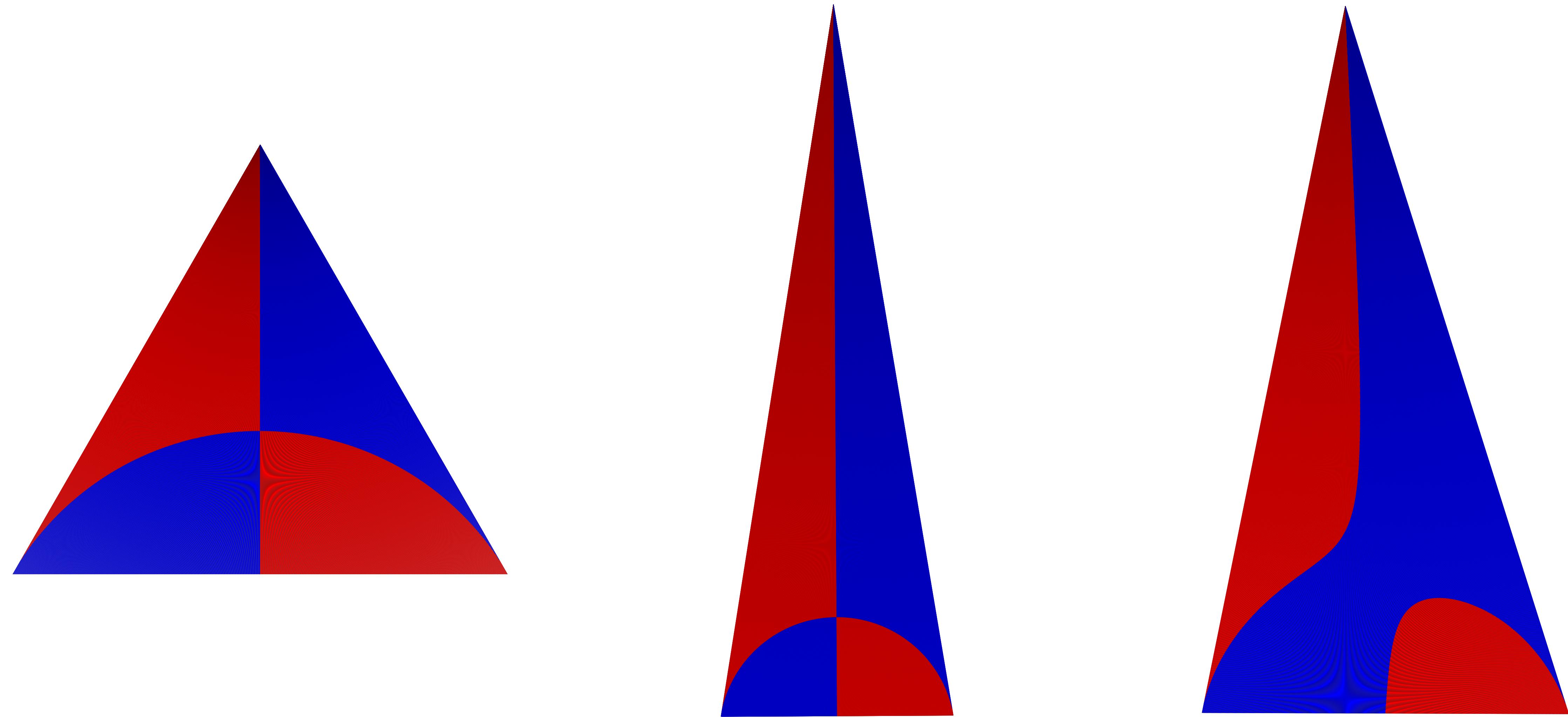
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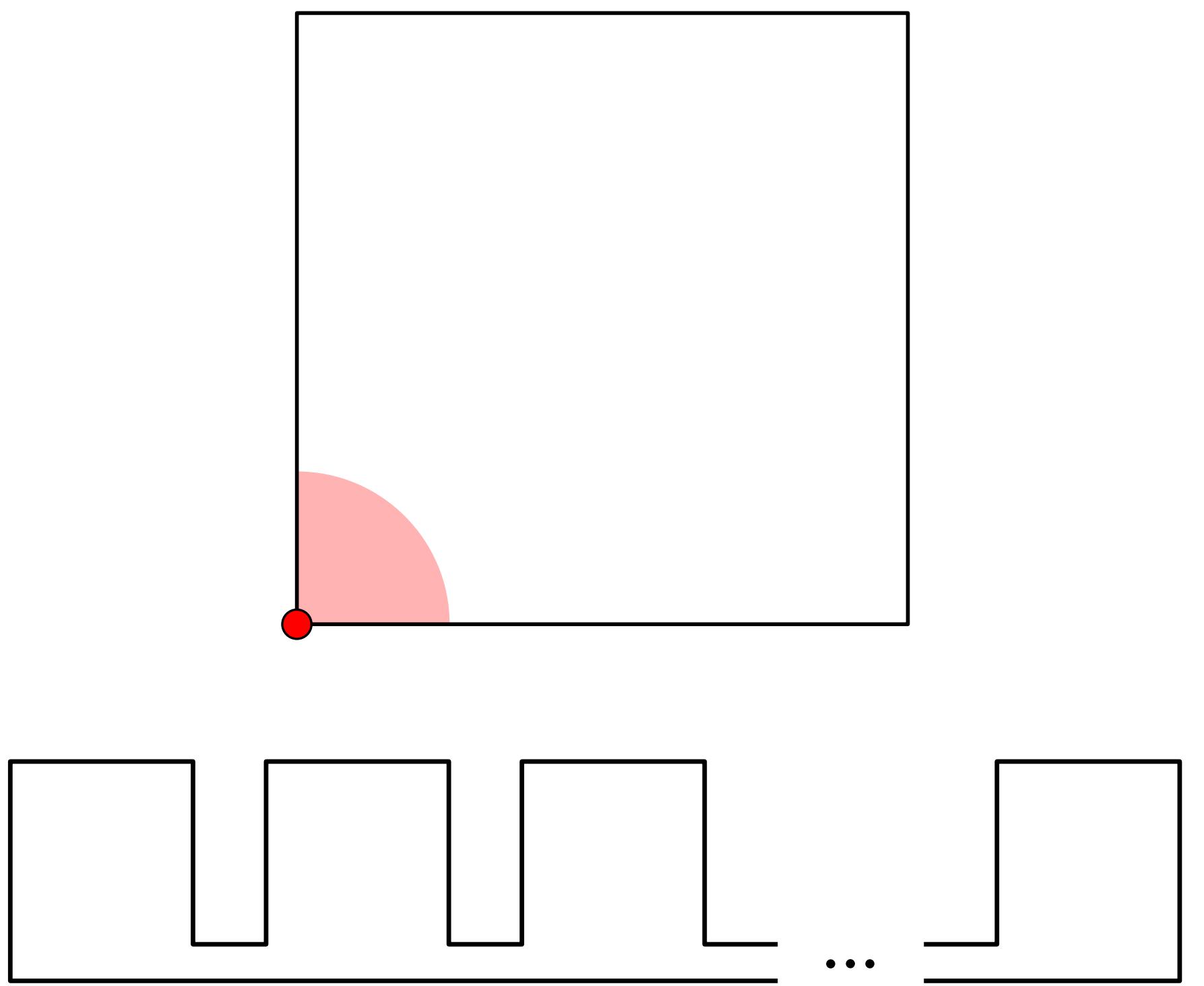
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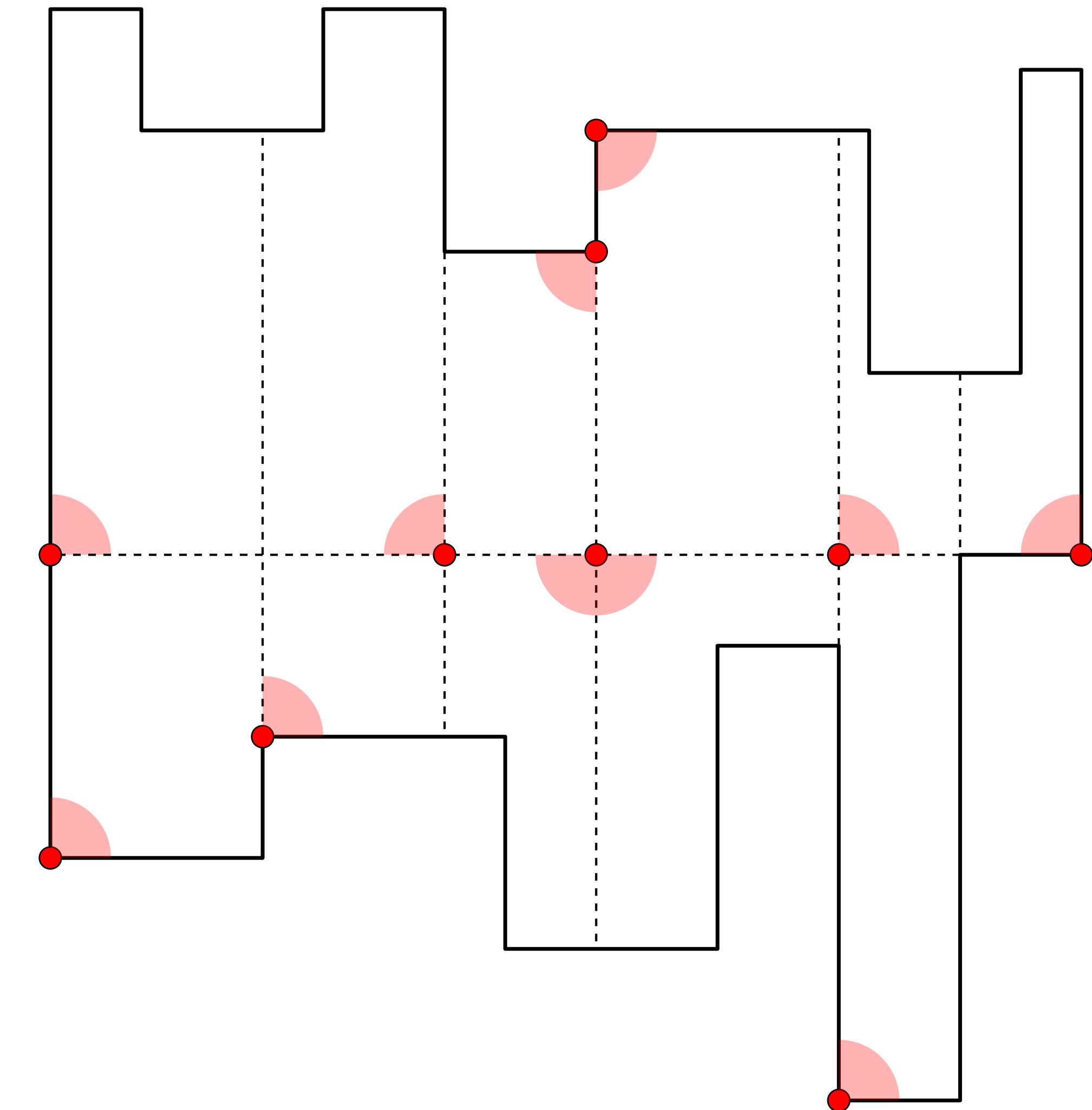
Future Work



Future Work



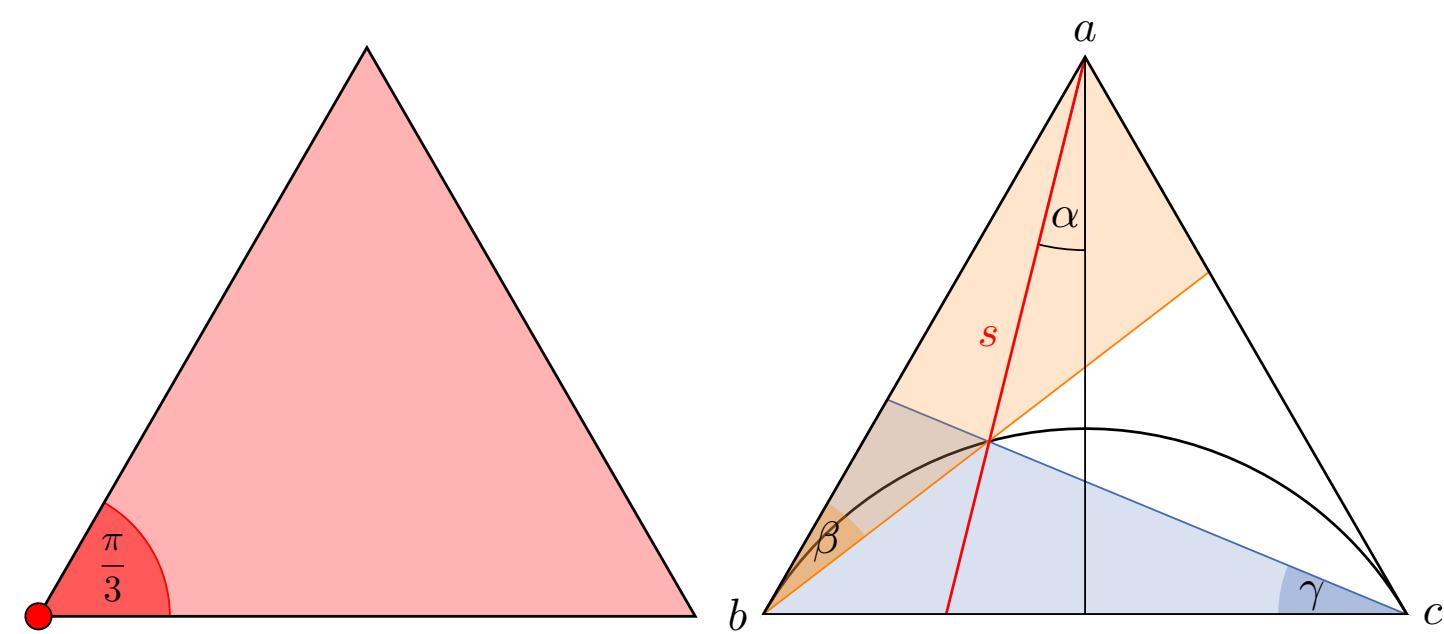
Assumed tight upper bound: $\left\lfloor \frac{n}{4} \right\rfloor \frac{\pi}{2}$



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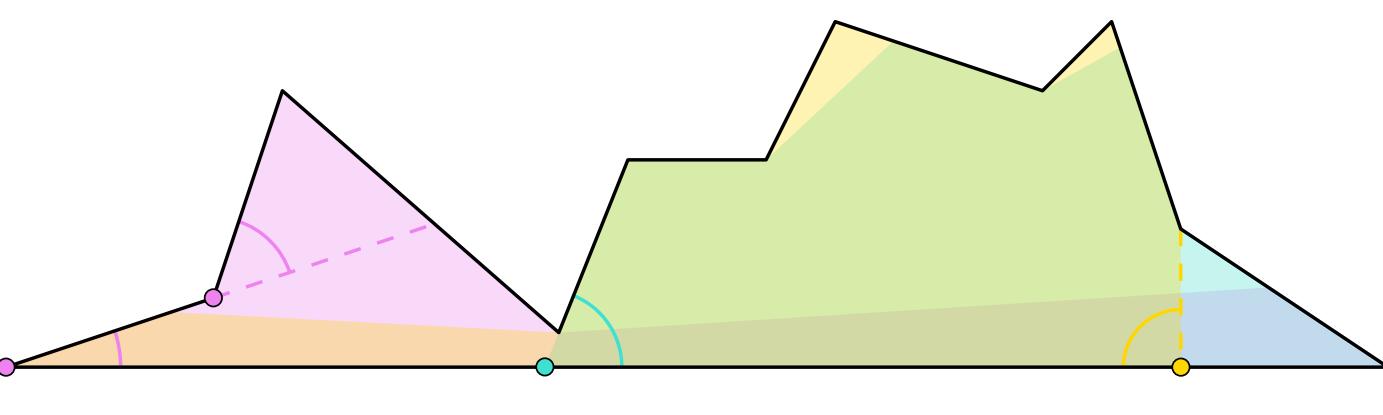
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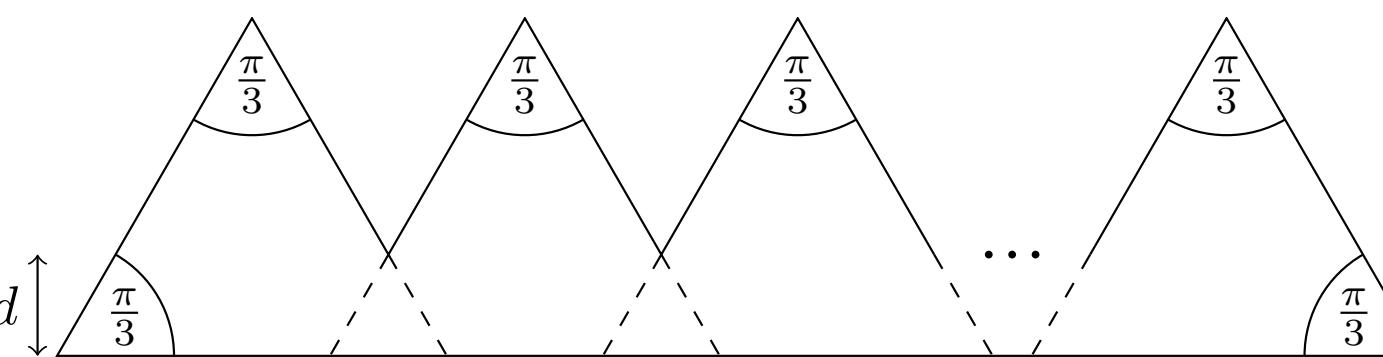
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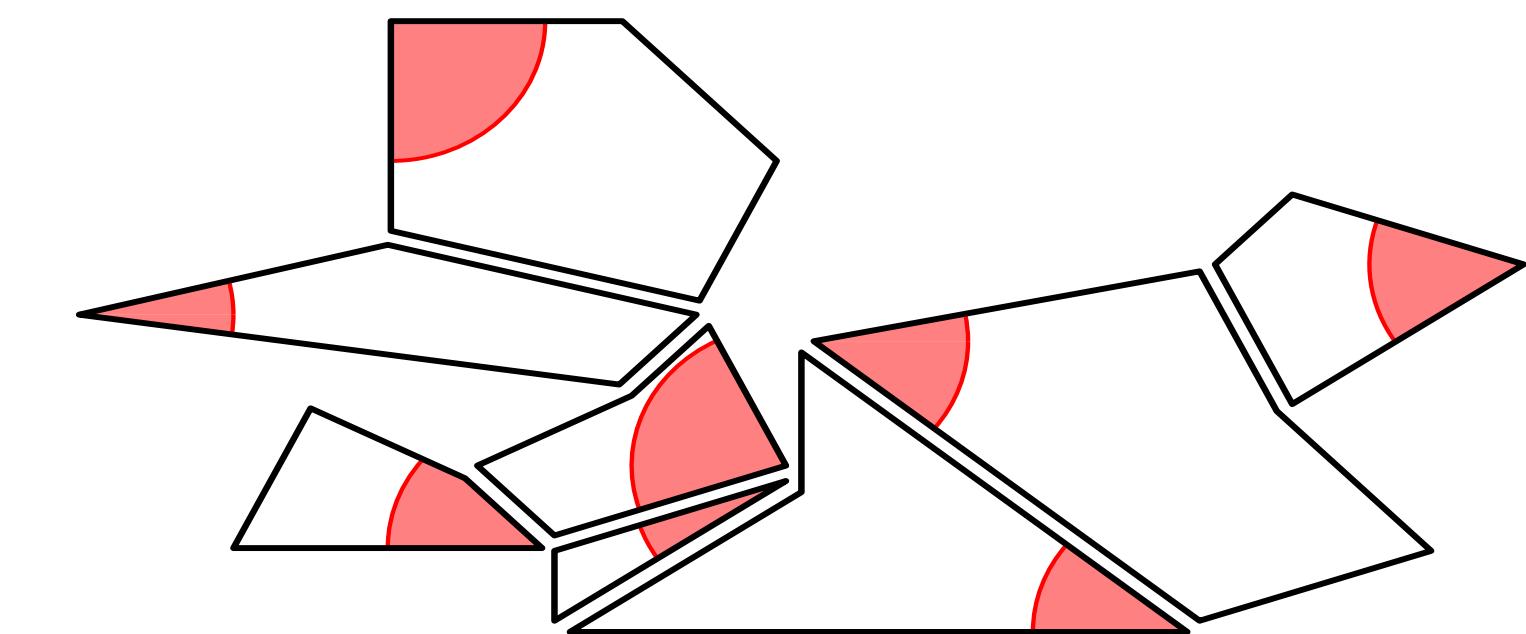
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