Complie and Execute

```
Q1: gcc A2Q1.c -o a2q1 -lpthread and ./a2q1 30 60 Q3: gcc A2Q3.c -o a2q3 -lpthread and ./a2q3 30 60
```

Question 1

Measurement

Observation

This implementation have a starvation problem that writer have to wait significantly longer than reader.

Reasoning

- All readers at a time slice should be considered to hold the rw_mutex lock if one of them actually holds, and therefore blocks all other pending writers
- Readers may give up rw_mutex lock only after read_count drops to 0, and a writer could write, or another reader could read, if they obtained the rw_mutex lock.
- Writer have to wait for readers if there are multiple readers, even if some readers arrives after the writer.
- For instance, if one reader is reading, and while the reader is reading, a writer followed by another n readers came, then the writers will wait for the (at least) n+1 readers to finish their reading (even if n reader came after the writer), since the writer will not give up file mutex lock until read_count is 0.

Question 2

Starvation already observerd in Q1

Question 3

See A2Q3.c

Question 4

Measurement

Since 4.680607 ms is not significantly higher than 7.408260 ms, there is no starvation occurred.

Proof

```
std_dev_read = range / 4 = 408.929 / 4 = 102.23225 (approx)
std_dev_write = range / 4 = 373.970 / 4 = 93.4925 (approx)
using Unpaired T-test, p-value is 0.6454, and even if alpha is 0.1, it is considered as not significant.
```

END of this assignment