$\mathsf{Hedge}(\eta)$

Online learning using Bregman Divergences

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Outline

 $\mathbf{Hedge}(\eta) \mathbf{Algorithm}$

Bound on total loss

The hedging problem

- N possible actions
- At each time step t = 1, 2, ..., T:
 - Algorithm chooses a distribution p^t over actions.
 - ▶ Losses $0 \le \ell_i^t \le 1$ of all actions i = 1, ..., N are revealed.
 - Algorithm suffers expected loss p^t · l_t
- ► Goal: minimize total expected loss
- Here we have stochasticity but only in algorithm, not in outcome

The **Hedge**(η)Algorithm

Consider action *i* at time *t*

► Total loss:

$$L_i^t = \sum_{s=1}^{t-1} \ell_i^s$$

Weight:

$$\mathbf{w}_{i}^{t} = \mathbf{w}_{i}^{1} \mathbf{e}^{-\eta L_{i}^{t}}$$

Note freedom to choose initial weight $(w_i^1) \sum_{i=1}^n w_i^1 = 1$.

- $ightharpoonup \eta > 0$ is the learning rate parameter. Halving: $\eta \to \infty$
- Probability:

$$\rho_i^t = \frac{w_i^t}{\sum_{j=1}^N w_i^t}, \quad \mathbf{p}^t = \frac{\mathbf{w}^t}{\sum_{j=1}^N w_i^t}$$

Bound on the loss of $Hedge(\eta)$ Algorithm

Theorem (main theorem)

For any sequence of loss vectors ℓ_1, \dots, ℓ_T , and for any $i \in \{1, \dots, N\}$, we have

$$L_{\mathsf{Hedge}(\eta)} \leq \frac{-\ln(w_i^1) + \eta L_i}{1 - e^{-\eta}}.$$

- Proof: by combining upper and lower bounds on $\sum_{i=1}^{N} w_i^{T+1}$

Comparing with the best distribution

- Comparison class: single experts. hindsite.
- Does not take advantage of multiple good experts.
- We will get tighter bounds by increasing the comparison class to include all convex combinations of the experts.

Recall Single step bound for **Hedge**(η)

The total weight has to decrease if the loss is large

$$\sum_{i=1}^N w_i^{t+1} \leq \left(\sum_{i=1}^N w_i^t\right) \left(1 - (1 - e^{-\eta})\mathbf{p}^t \cdot \ell_t\right)$$

Enlarging the comparison set

- Bound compares cumulative loss to that of best expert in hindsite.
- Does not take advantage of multiple good experts.
- We will get tighter bounds by comparing to the best convex combination of experts.

Comparing with the best distribution

- ▶ Denote by \mathbf{q} an arbitrary distribution over \mathbf{N} experts. $\mathbf{q} \in \Delta^{\mathbf{N}}$. Distribution = convex combination.
- Compare loss of algorithm to loss of best convex combination of experts:

$$\sum_{t=1}^{T} L_A^t \le +a \min_{\mathbf{q} \in \Delta^N} \sum_{t=1}^{T} \mathbf{q} \cdot \ell_t + cX$$

- ▶ When comparing to single best expert $X = \log N$
- ► Intuition: X should be small if best distribution q* is close to initial distribution p⁰

Relative Entropy Bound

- KL-divergence or Relative Entropy: X
- ► For any distribution \mathbf{q} and any iteration of $\mathbf{Hedge}(\eta)$:

 $Hedge(\eta)$

Proof (from RE to ratio)

 $\mathsf{Hedge}(\eta)$

Proof (from ratio to bound)



 $\mathsf{Hedge}(\eta)$

Visual Intuition