## Introduction to AI – 236501 Local Search & Optimization [Chapter 4.1]

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Slides adapted from Shaul Markovitz @ Technion







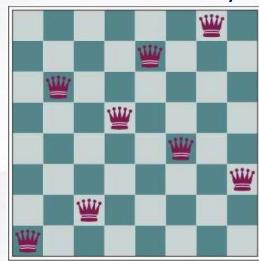
### Local search – when the goal is all that matters

In search problems we were interested in computing a path in the search space

Sometimes we care about only finding a state that maximizes some utility

function (or minimizes some cost function)

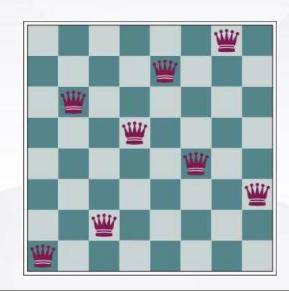
- 8-queens problem
- Integrated circuit design
- Floor layout
- Telecommunication network optimization





### Local search – when the goal is all that matters

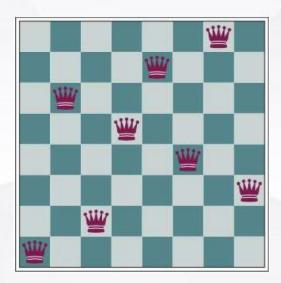
- In local search we typically store only one or a few states in memory (without the path to reach these states) nor the states that were already reached
  - No systematic coverage of the search space
  - Very low memory footprint
  - Can often find good solutions in reasonable times in huge search spaces
- Heuristics here don't measure the distance to the goal but how good a certain state is





### **Example 8-Queen problem**

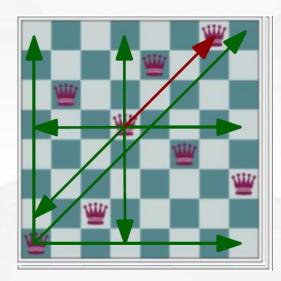
Given n=8 queens, set them up on a chess board so no queen can attack another





### **Example 8-Queen problem**

Given n=8 queens, set them up on a chess board so no queen can attack another

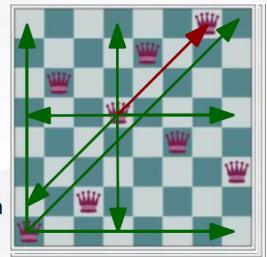




#### **Example 8-Queen problem**

▶ Given n=8 queens, set them up on a chess board so no queen can attack another

- Didn't we solve this in Intro to CS using recursion?
- Isn't that good enough?
  - Using recursion you can search up to n! states
  - OK when **n=8** and **n!** is around **40,000**
  - When **n=20**, **n!** is around **10^18**
  - Local algorithms like the one we will learn can handle in reasonable time instances with n=3,000,000

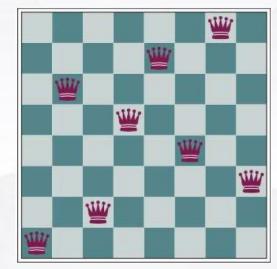




### 8-Queen problem - representation

- A state will be defined as an 8-tuple
- ▶ The i'th element will denote the row of the queen placed in the i'th column
- An operator moving from one state to another is a pair (i,j) stating that the i'th queen is moved to the j'th row
  - => branching factor is **7X8=56**
- A heuristic can count the number of pairs of queens attacking each another

(8, 3, 7, 4, 2, 5, 1, 6)





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(h=17: 3+4+2+3+2+2+1+0)

18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	<b>W</b>	13	16	13	16
事	14	17	15	w	14	16	16
17	W	16	18	15	W	15	W
18	14	<b>W</b>	15	15	14	<b>W</b>	16
14	14	13	17	12	14	12	18

## Hill climbing







### Steepest ascent hill-climbing search (SAHC)

- Search starts with a given start state
- Every step, expand all neighbors
- Choose neighbor with best heuristic value
- ▶ If several exist, choose one randomly
- ▶ Terminate when no neighbor can improve current state

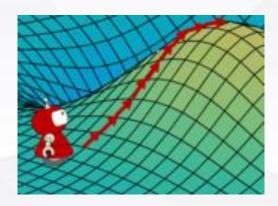
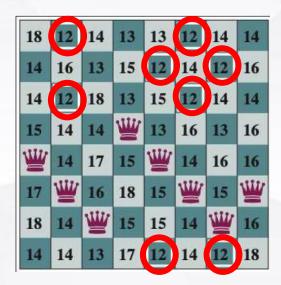


Figure adapted from https://tinyurl.com/2wbyrbta



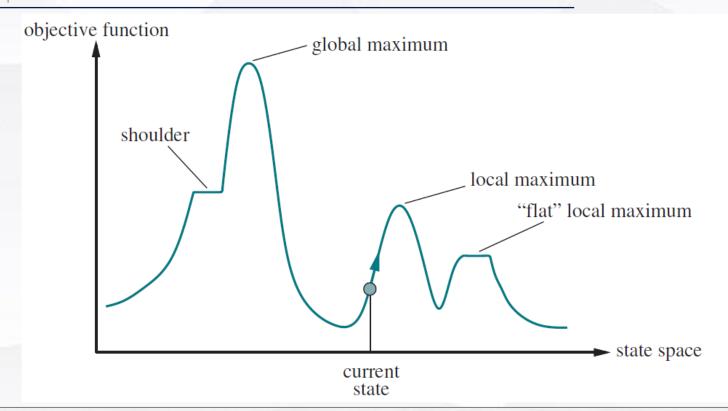
### Steepest ascent hill-climbing search

- In this example, we have
  - Heuristic value of 17
  - Branching factor is 56
  - 8 neighbors with a value of 12 one of them will be picked randomly





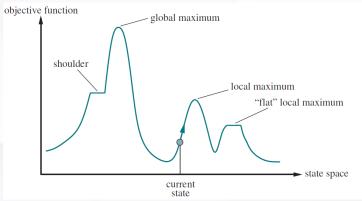
### Topography of the objective function

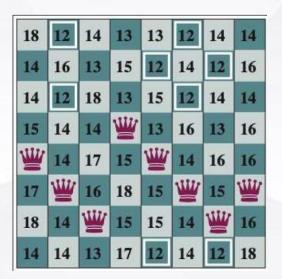




### How will SAHC work on the 8-queen puzzle?

- ▶ After running the algorithm **1,000** times from random start locations
  - Success rate of 14% (rest get stuck in a local minimum)
  - 4 iterations on average when succeeds
  - **3** iterations on average when fails

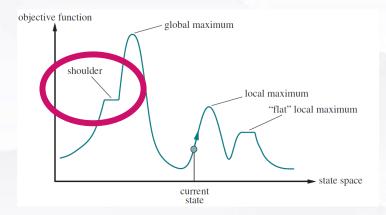






### SAHC search with sideway moves

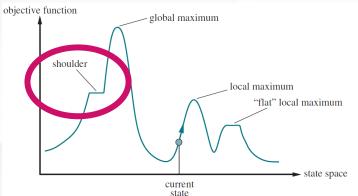
- ▶ In many settings, SAHC fails because it gets stuck on a shoulder where neighboring states have identical values
- ▶ We can overcome this by allowing to transition to a state with the same value
- Typically we add a limit to the number of times this is allowed

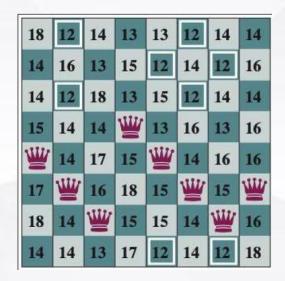




### How will SAHC + sideway moves work on the 8-queen puzzle?

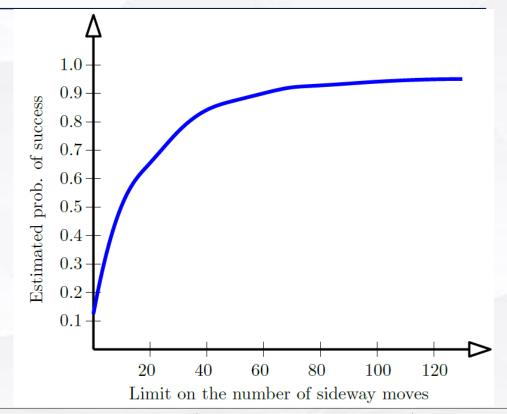
- After running the algorithm 1,000 times from random start locations
  - Success rate of 94% (instead of 14%)
  - 21 iterations on average when succeeds (instead of 4)
  - 64 iterations on average when fails (instead of 3)







### How will SAHC + sideway moves work on the 8-queen puzzle?





### **Stochastic Hill Climbing**

- ▶ The SAHC is a greedy algorithm that chooses at each step the mostpromising neighbor
- In Stochastic Hill Climbing we allow other moves as well: choose a neighbor proportionally to the improvement it offers
  - Usually converges slower than SAHC but in some settings finds better solutions (when?)
  - If  $\Delta_i$  denotes the improvement of moving to neighbor i
  - Stochastic Hill Climbing randomly picks neighbor i with probability

$$P_i = \frac{\Delta_i}{\sum_{i \ s.t.\Delta_i > 0} \Delta_i}$$

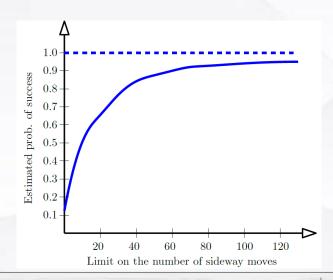


### **First-choice Hill Climbing**

- In certain settings the branching factor may be huge (e.g., thousands of successors)
- In **first-choice hill climbing** we randomly generate successors and pick the first one that is better than the current state

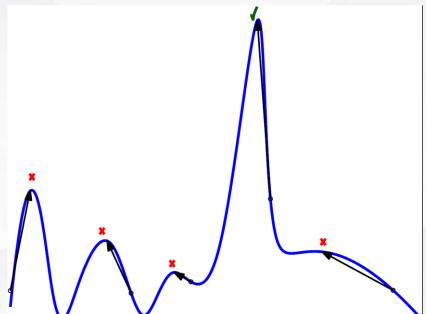


- Conduct a series of hill climbing searches from randomly—generated initial states
- Complete with prob. 1 (why?)



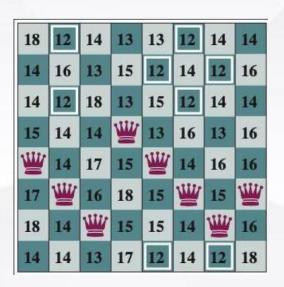


 Conduct a series of hill climbing searches from randomly –generated initial states





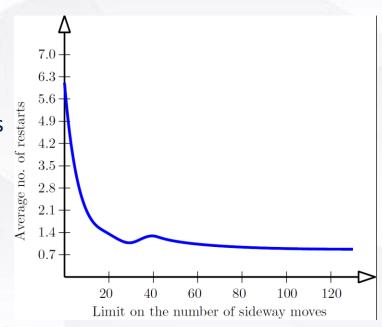
- Conduct a series of hill climbing searches from randomly –generated initial states
- ▶ On the 8-queen puzzle
  - Success rate of 100% (instead of 14% / 94%)
  - 22 iterations on average w.o. sideway moves (instead of 4)
  - 25 iterations on average w. sideway moves (instead of 21)





 Conduct a series of hill climbing searches from randomly –generated initial states

- ▶ On the 8-queen puzzle
  - Success rate of 100% (instead of 14% / 94%)
  - 22 iterations on average w.o. sideway moves (instead of 4)
  - 25 iterations on average w. sideway moves (instead of 21)



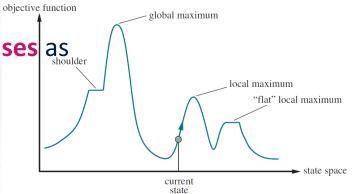






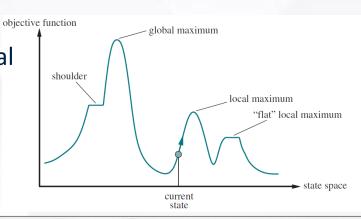
- All hill-climbing algorithms never makes "down-hill moves" towards states with lower value
- Can get stuck in a local minimum
- Simulated annealing allows moving to states with a lower heuristic value
- The probability to choose such a state decreases the search progresses

How?





- ▶ The probability to choose such a state decreases as the search progresses
- ▶ This is done by a parameter called the **temperature**
- ▶ The algorithm starts with a high temperature and decreases it gradually
- Multiple approaches on how to choose the initial temperature and how to schedule the decrease



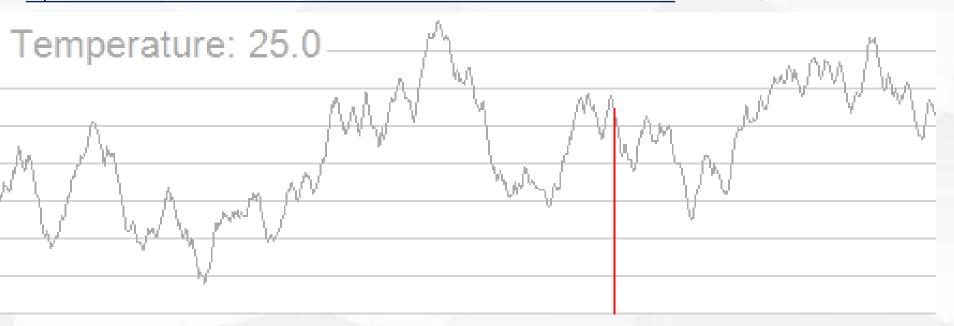


#### Simulated-annealing (state)

- curr <- state; curr-val <- h(state); T<- T<sub>0</sub>
- loop until resources exhausted
  - new <- random-neighbor(curr)</p>
  - new-val<-h(new)</p>
  - if goal(new) then return (new)
  - $\Delta F$  <- new-val curr-val
  - if  $\Delta E > 0$  then curr <- new; curr-val <- new-val
  - else with prob.  $e^{-|\Delta E|/T}$  curr <- new; curr-val <- new-val
  - T<- 0.95 \* T</p>
    //simple scheduling scheme
- return curr



### Simple 1-D cost map (find max)



By Kingpin13 - Own work, CCO, https://commons.wikimedia.org/w/index.php?curid=25010763



**Task**: find the shortest path to visit all cities exactly once

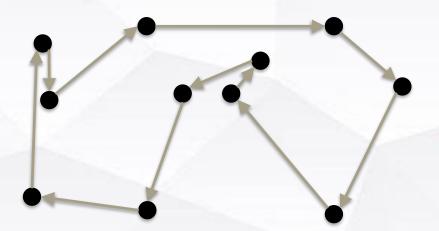
- Hamiltonian cycle in a graph
- Simplification: cities are points in 2D





Task: find the shortest path to visit all cities exactly once

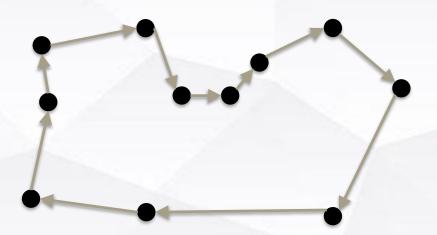
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**Task**: find the shortest path to visit all cities exactly once

- Hamiltonian cycle in a graph
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Task: find the shortest path to visit all cities exactly once

- Hamiltonian cycle in a graph
- Simplification: cities are points in 2D

#### **Decision variables:**

Where to go next after every city

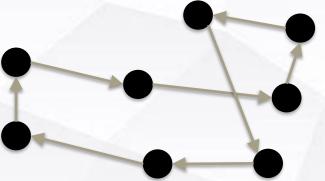
TSP is probably the most well-studied combinatorial problem!



**Task**: find the shortest path to visit all cities exactly once

- Hamiltonian cycle in a graph
- simplification: cities are points in 2D

**Local move**: select two edges and replace them by two other edges

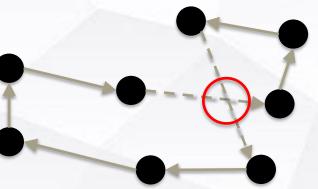




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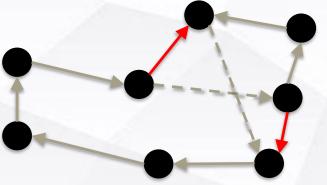
Crossings are bad!



**Task**: find the shortest path to visit all cities exactly once

- Hamiltonian cycle in a graph
- simplification: cities are points in 2D

**Local move**: select two edges and replace them by two other edges

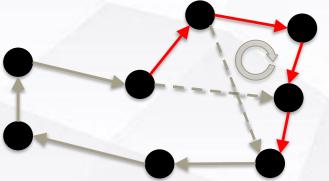




Task: find the shortest path to visit all cities exactly once

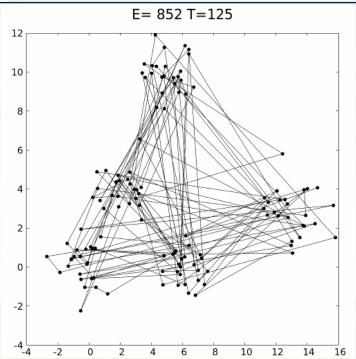
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### **Application to TSP**



By Geodac - Own work, CCO, https://commons.wikimedia.org/w/index.php?curid=67988888



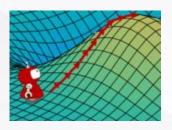
#### **Application to TSP**

Adapted from https://www.youtube.com/watch?v=NPE3zncXA5s

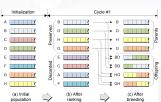


#### Local search: conclusion

- We examined several approach for local that start from an initial solution and iteratively improve the solution.
- It is not a systematic search and therefore compromises completeness, but it is very efficient and practical.
- In the appendix you can find information about genetic\
  evolutionary algorithms that use biologically inspired ideas such as mutation, crossover and selection to find solutions. They are beyond scope for the course.







# Evolutionary algorithms (not part of the course material)







#### **Genetic algorithm**

- Local search algorithm that is inspired by evolution
- Main idea: take "good parts" from different states and "fuse" them together to obtain new states
- There are endless forms of genetic algorithms. We will only give a flavor of the approach

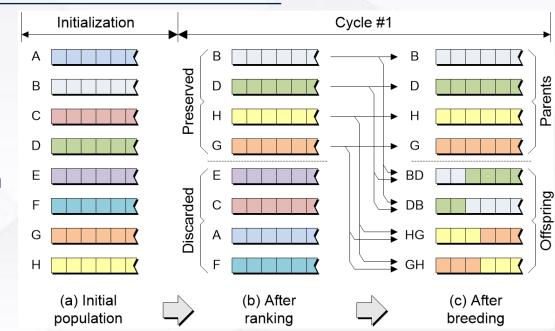


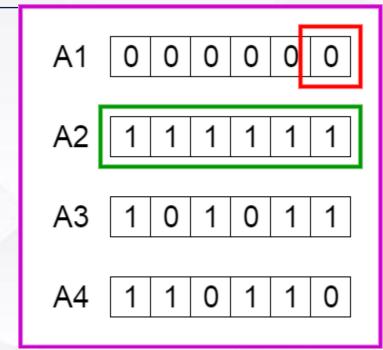
Image source: Max Maxfield





#### **Genetic algorithm (cont.)**

- The algorithm stores a set of solutions called the population
- Each solution, sometimes referred to as a "chromosome" is typically represented as a vector
- ► Each element in a vector is called a "gene"



Gene

Chromosome

Population

Image source: https://tinyurl.com/2n7nssbh

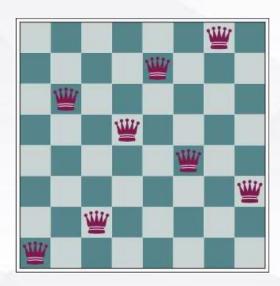


### **Genetic algorithm (cont.)**

- We typically have a fixed populations size
- The "mixing number"  $\rho$  defines how many parent genes are used to create a new child node. Typically is 1 or 2
- ▶ The selection process defines how we choose chromosomes that will
  - Move to the next iteration or
  - Serve as "parents" to create "children" for the next generation
     Often implemented by selecting individuals according to their fitness score
- The recombination process (assuming  $\rho \geq 2$  ) typically chooses a random crossover point to split each parent string and then combine the two parts of each string
- ▶ The "mutation rate" defines how often we randomly flip a gene



- A chromosome here will be a vector of 8 genes (digits)
- ▶ The i'th digit represents the row number of the queen in column i
- The fitness is the number of nondominating pairs of queens





Assume we start with a population size of 4

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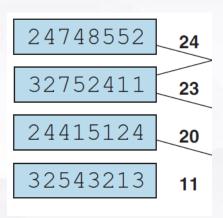
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24415124

32543213

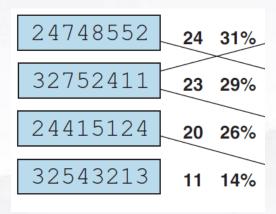


▶ We compute the **fitness** of each state





▶ The **fitness** induces the probability to choose that state for the next step of the process (this is the **selection** process)





- In this **selection** process, we choose **4** states
- Notice that one state was chosen twice and one was not chosen





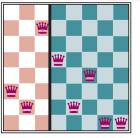
For each pair of parents we randomly choose a crossover point and use it to generate a pair of new child states

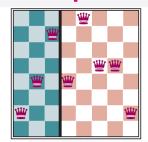


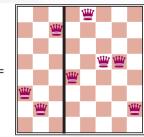


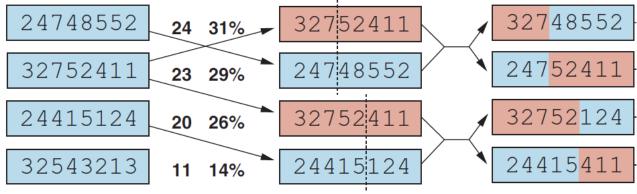
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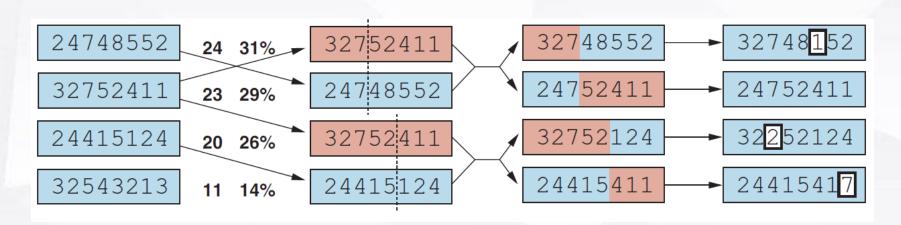








▶ These new child states are subject to mutation with some (small) probability





#### When to use GA?

▶ GA attracted a lot of attention in the 1990's and early 200's

Not clear if this is because of the superiority of the method or the appealing

metaphor of evolution

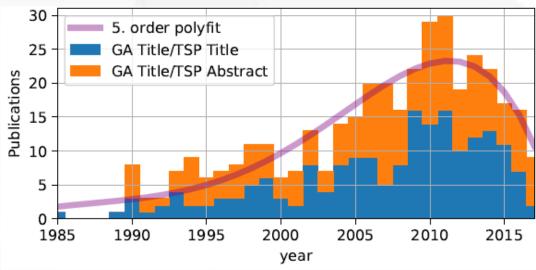


Figure adapted from Scholz 2018





## When to use GA?

- When we can represent a solution as a vector of values
- Works well when parts of a good solution also exhibit good performance
- Especially works well on variants of the problem of choosing a good subset
  - Given a set of objects S
  - Given a function U that gives a score to every subset of S
  - Find a subset S' with maximal U-value



#### When to use GA? Examples

- ▶ Given some text T with N sentences. Find a set of sentences that will summarize T in the best way
- Given a query Q and N search results. Find a subset of 10 search results that maximize the prob. that a user will click on one.
- Given N courses a student can take. Find a subset for next semester that maximize some utility (diversity, workload, good exam period etc.)