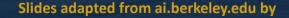
Introduction to AI – 236501 **Problem Solving and Blind Search** [Chapter 3]





and from Shaul Markovitz @ Technion





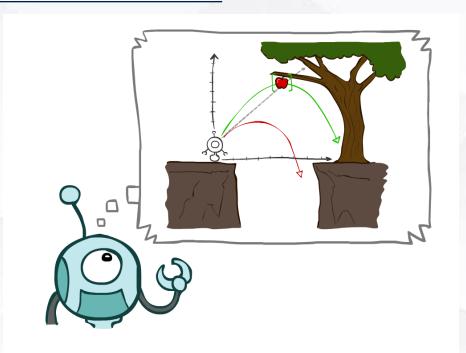


The Henry and Marilyn Taub Faculty of Computer Science



Agents that Plan Ahead

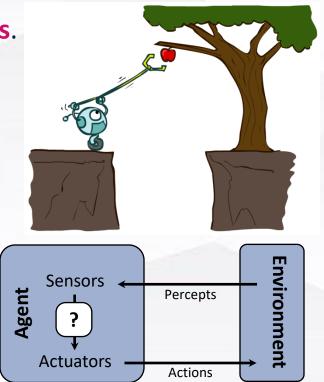
Search Problems





Designing Rational Agents

- ▶ An agent is an entity that perceives and acts.
- ▶ A rational agent selects actions that maximize its (expected) utility.
- Characteristics of the percepts,
 environment, and action space dictate
 techniques for selecting rational actions



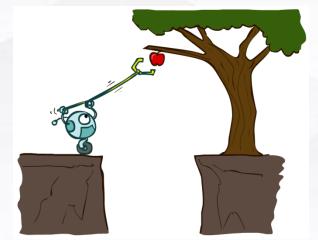


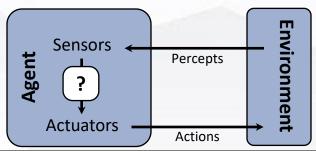
How do we design the rational agent's decision procedure?

- An agent's decision procedure maps a world's state and a desired goal to an action
- We can hardcode the mapping

$$State \times Goal \rightarrow Action$$

- May be an infinite table
- We can hardcode the mapping $\operatorname{Percept} \times \operatorname{Goal} \to \operatorname{Action}$
 - May (still) be an infinite table
 - This is called a reflex agent



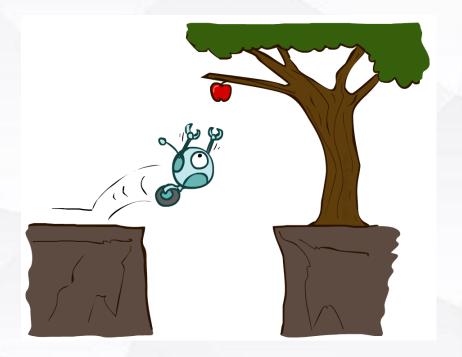




Reflex Agents

Reflex agents:

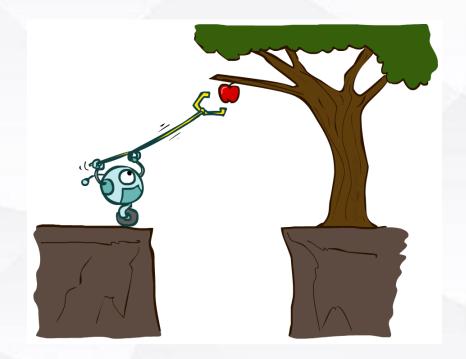
- Choose action based on current percept (and maybe memory)
- May have memory or a model of the world's current state
- Do not consider the future consequences of their actions
- Consider how the world IS





Planning Agents

- Planning agents:
 - Ask "what if"
 - Decisions based on (hypothesized) consequences of actions
 - Must have a model of how the world evolves in response to actions
 - Must formulate a goal (test)
 - Consider how the world WOULD BE
- Optimal vs. complete planning
- Planning vs. replanning

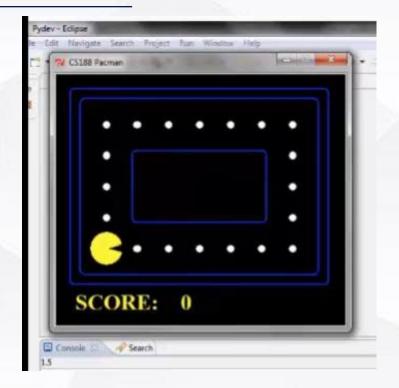




Video of Demo Reflex - Optimal

Reflex agent

- Move in the direction of the nearest uneaten dot
- You can think ahead all you want, still end up doing the same thing
- Can be rational

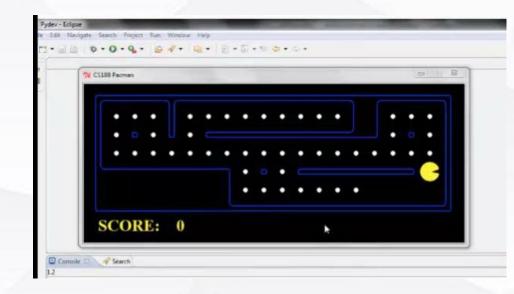






Video of Demo Reflex – incomplete

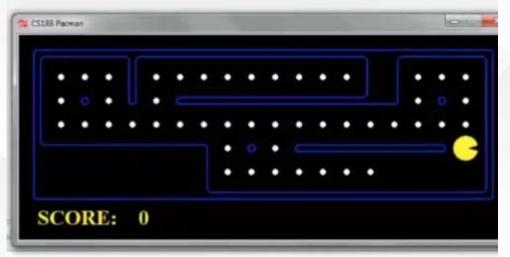
- Reflex agent
 - Move in the direction of the nearest uneaten dot
 - Can be incomplete
 - No planning
- Reflex agents not generally optimal, but they can be depending on circumstance.





Video of Demo replaning reflex agent – sub optimal

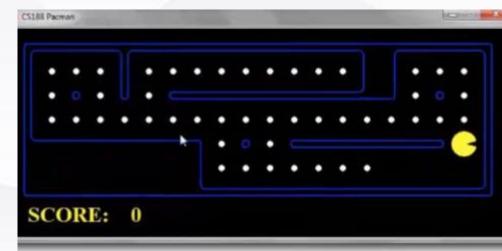
- Replaning reflex agent
 - Move in the direction of the nearest uneaten dot
 - Clears board
 - Non-optimal
 - Fast





Video of Demo mastermind- Optimal

- Mastermind agent
 - Thinking, thinking, thinking...
 - Not just clear NEXT dot, but all dots
 - Know to end with dead end so it doesnt have to back up



Problem-Solving Agents







How do we formalize a search problem?





Search Problems

- A search problem consists of:
 - A state space





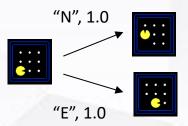








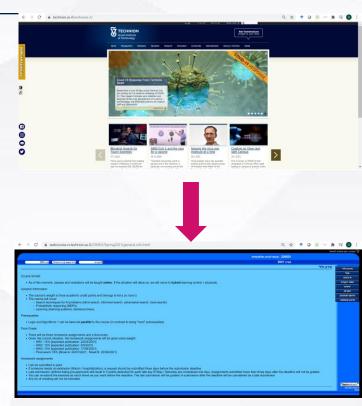
 A successor function (with actions, costs) or a set of operators



- A start state and a goal test
- A solution is a sequence of actions (a plan) which transforms the start state to a goal state

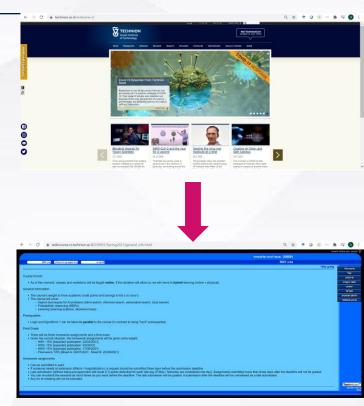


Problem: Given the Technion's homepage, find the series of clicks that will get you to the homepage of Intro To AI



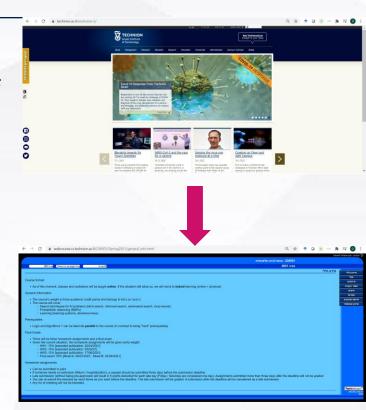


- ▶ **Problem**: Given the Technion's homepage, find the series of clicks that will get you to the homepage of Intro To AI
- State space:



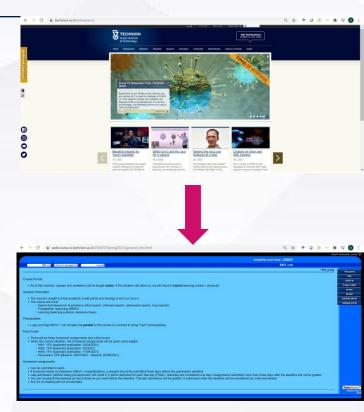


- **Problem**: Given the Technion's homepage, find the series of clicks that will get you to the homepage of Intro To Al
- **State space:** all URLs





- ▶ **Problem**: Given the Technion's homepage, find the series of clicks that will get you to the homepage of Intro To AI
- State space: all URLs
- Successor function:







- Problem: Given the Technion's homepage, find the series of clicks that will get you to the homepage of Intro To AI
- State space: all URLs
- Successor function: for each page the set of successors is defined by the links available at that page (cost = 1)...











- Problem: Given the Technion's homepage, find the series of clicks that will get you to the homepage of Intro To AI
- ▶ State space: all URLs
- Successor function: for each page the set of successors is defined by the links available at that page (cost = 1)...
- Start state:
- Goal state:







- ▶ **Problem**: Given the Technion's homepage, find the series of clicks that will get you to the homepage of Intro To AI
- ▶ State space: all URLs
- Successor function: for each page the set of successors is defined by the links available at that page (cost = 1)...
- Start state: https://www.cs.technion.ac.il/
- Goal state: https://webcourse.cs.technion.ac.il/236501/







Search Problems Are Models

- Models aren't perfect
 - Too detailed, you can't solve
 - Not detailed enough, doesn't solve

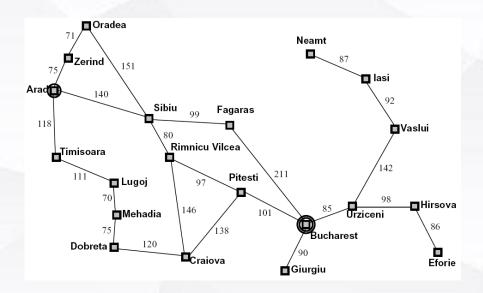




Example: Traveling in Romania

- State space:
 - Cities
- Successor function:
 - Roads: Go to adjacent city with cost = distance
- Start state:
 - Arad
- Goal test:
 - Is state == Bucharest?







What's in a State Space?

The world state includes every last detail of the environment



A search state keeps only the details needed for planning (abstraction)

- Problem: Path finding
 - States: (x,y) location
 - Actions: NSEW
 - Successor: update location only
 - Goal test: is (x,y)=END

- Problem: Eat-All-Dots
 - States: {(x,y), dot booleans}
 - Actions: NSEW
 - Successor: update location and possibly a dot boolean
 - Goal test: dots all false



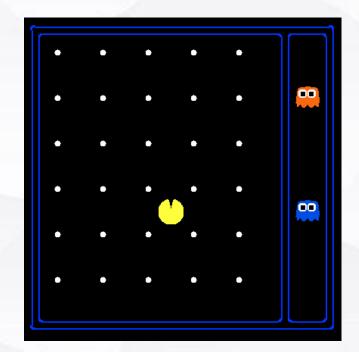
State Space Sizes?

World state:

- Agent positions: 12X10 = 120
- Food count: 30
- **Ghost positions: 12**
- Agent facing: NSEW

How many

- World states? $120x(2^{30})x(12^2)x4$
- States for path finding? 120
- States for eat-all-dots? $120x(2^{30})$





Some assumptions

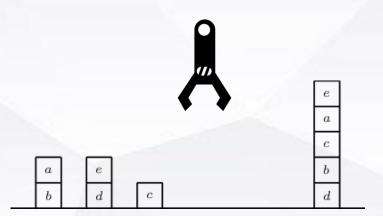
- Set of actions is finite and bounded by some constant
 - In practice, they are often parametrized and thus infinite (e.g., steering angle of a car)
 - In such settings we can discretize the set of actions
- State of agent is observable
 - i.e., agent knows its state
- Actions are deterministic
 - i.e., agent knows what will happen when it does something
- We will relax some of these assumptions along the course

State Space Search Problem: Classical Planning

- Formal definition:
 - a finite and discrete state space S
 - a known initial state $s_0 \in S$
 - a non-empty set of goal states $S_G \subseteq S$ (or a goal test)
 - set of actions/operators $A(s) \in A$ applicable in each state $s \in S$
 - A successor function a **deterministic** state transition function f(a, s) such that s' = f(s, a) stands for the state resulting from applying an applicable action a at state s
 - A cost function c(a, s) associating a positive cost to applying action a in s
- Objective
 - Find a plan, a sequence of applicable actions, that leads from the initial state to a goal state.
 - Optimality criteria typically interested in shortest/ cheapest plan (but other considerations are possible)

Factored State Space Definition: Using Features

- ▶ Typically, it is not possible to explicitly specify the **state space** *S*
- Instead, **feature set** X is used to represent the state space such that a state is described via a set of random variables $X = x_1, \ldots, x_n$ such that each variable takes on values in some finite domain $Dom(x_i)$.
- Can be Boolean or multi-valued
- Example:
 - Position of block a
 - Position of block c
 - ...
 - Is gripper empty



Formalizing a search problem







Problem:

- 3 missionaries and 3 cannibals must cross a river using a boat that carries at most 2 people
- If there are missionaries present on the bank, they cannot be outnumbered by cannibals (if they were, the cannibals would eat the missionaries)
- The boat cannot cross the river by itself with no people on board
- How do we represent the search problem?



"Hmm ... maybe I'll just have a salad."



► Take (1):

- A state is a 7-tuple $< L_1, L_2, L_3, L_4, L_5, L_6, B >$
- $L_i \in \{\text{Left}, \text{Right}, \text{Boat}\}$ location of person i
- $B \in \{\text{Left}, \text{Crossing}, \text{Right}\}$ location of boat
- Start state: < Right, Right, Right, Right, Right, Right >
- Goal state: < Left, Left, Left, Left, Left, Left, Left >
- How many states do we have?



"Hmm ... maybe I'll just have a salad."



► Take (1):

- A state is a 7-tuple $< L_1, L_2, L_3, L_4, L_5, L_6, B >$
- $L_i \in \{\text{Left}, \text{Right}, \text{Boat}\}$ location of person i
- $B \in \{\text{Left}, \text{Crossing}, \text{Right}\}$ location of boat
- Start state: < Right, Right, Right, Right, Right, Right >
- Goal state: < Left, Left, Left, Left, Left, Left, Left >
- How many states do we have? $3^7 = 2187$
- Unnecessary information
 - We don't care which missionary / cannibal is on which side
 - We don't care about the state of the boat while crossing



"Hmm ... maybe I'll just have a salad."

► Take (2):

- A state is a 3-tuple $\langle M_L, M_C, B \rangle$
- $M_L \in \{0, \dots, 3\}$ number of missionaries on left bank
- $M_C \in \{0, \dots, 3\}$ number of cannibals on left bank
- $B \in \{\text{Left}, \text{Right}\}\$ location of boat
- Start state: < 0, 0, Right >
- Goal state: < 3, 3, Left >
- How many states do we have?



"Hmm ... maybe I'll just have a salad."



► Take (2):

- A state is a 3-tuple $\langle M_L, M_C, B \rangle$
- $M_L \in \{0,\ldots,3\}$ number of missionaries on left bank
- $M_C \in \{0, \dots, 3\}$ number of cannibals on left bank
- $B \in \{\text{Left}, \text{Right}\}\$ location of boat
- Start state: < 0, 0, Right >
- Goal state: < 3, 3, Left >
- ▶ How many states do we have? $4 \times 4 \times 2 = 32$



"Hmm ... maybe I'll just have a salad."

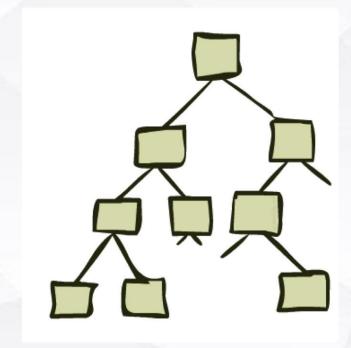
Search algorithms - preliminaries







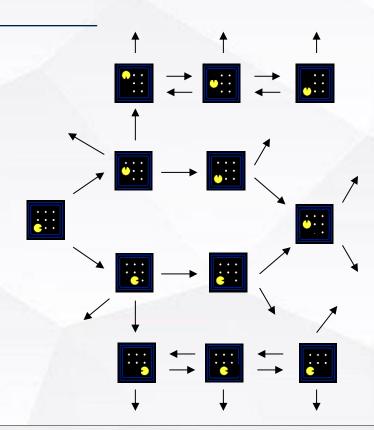
State Space Graphs and Search Trees





State-Space Graphs

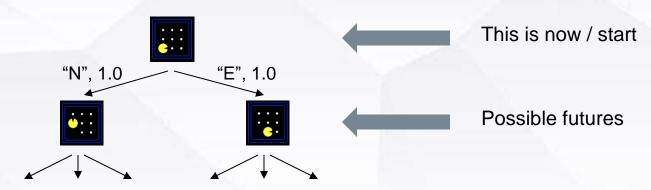
- State-space graph: A mathematical representation of a search problem
 - Nodes are (abstracted) world configurations
 - Arcs represent successors (action results)
 - The goal test is a set of goal nodes (maybe only one)
- In a state-space graph, each state occurs only once!
- We can rarely build this full graph in memory (it's too big), but it's a useful idea
 - The graph may be infinite
 - The graph is given implicitly







Search Trees

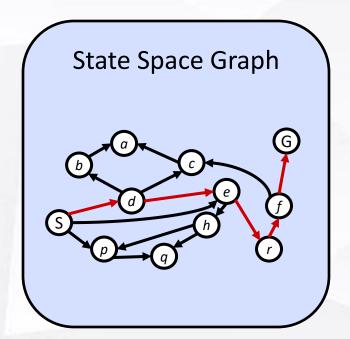


A search tree:

- A "what if" tree of plans and their outcomes
- The start state is the root node
- Children correspond to successors
- Nodes show states, but correspond to PLANS that achieve those states
- For most problems, we can never actually build the whole tree

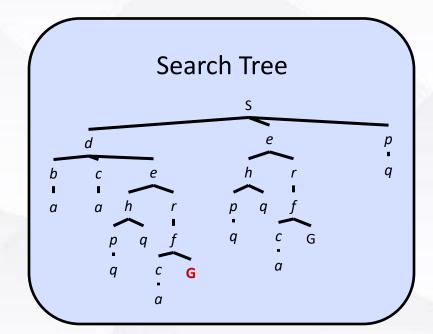


State Space Graphs vs. Search Trees



Each **NODE** in the search tree is an entire PATH in the state space graph.

We construct both on demand – and we construct as little as possible.

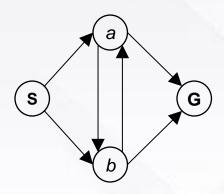




State Space Graphs vs. Search Trees

Consider this 4-state graph:

How big is its search tree (from S)?



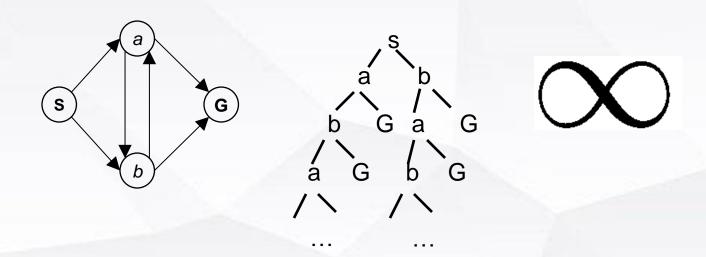




State Space Graphs vs. Search Trees

Consider this 4-state graph:

How big is its search tree (from S)?



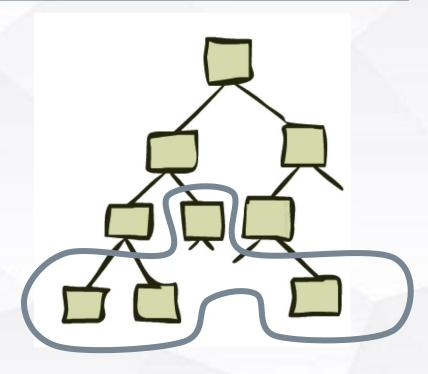
Important: Lots of repeated structure in the search tree!





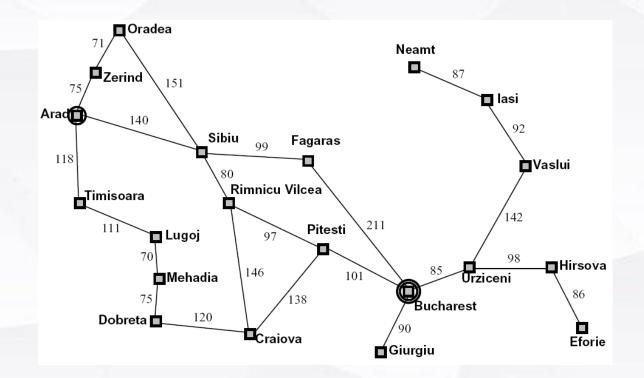


Tree Search



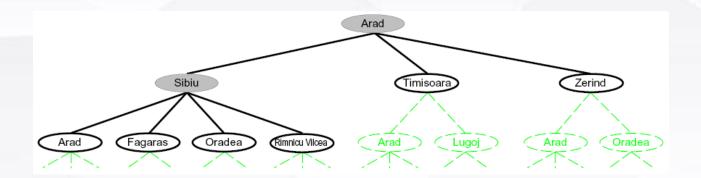


Search Example: Romania





Searching with a Search Tree



Search:

- Expand out potential plans (tree nodes)
- Maintain a fringe of partial plans under consideration
- Try to expand as few tree nodes as possible



General Tree Search

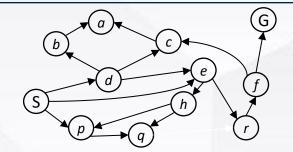
function TREE-SEARCH (problem, strategy) returns a solution, or failure initialize the search tree using the initial state of *problem* loop do if there are no candidates for expansion then return failure choose a leaf node for expansion according to strategy if the node contains a goal state then return the corresponding solution else expand the node and add the resulting nodes to the search tree end

- Important ideas:
 - Fringe
 - **Expansion**
 - **Exploration strategy**
- Main question: which fringe nodes to explore?



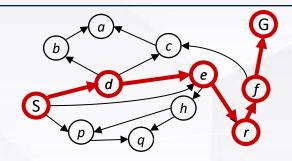


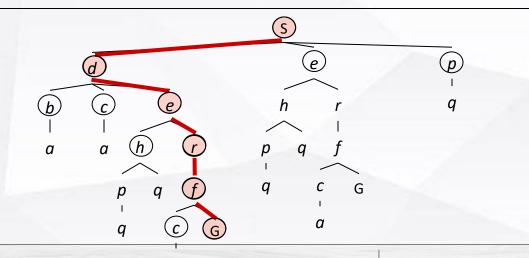
Example: Tree Search





Example: Tree Search





 $\begin{array}{c} s \\ s \rightarrow d \\ s \rightarrow e \\ s \rightarrow p \\ s \rightarrow d \rightarrow b \\ s \rightarrow d \rightarrow c \\ s \rightarrow d \rightarrow e \\ s \rightarrow d \rightarrow e \rightarrow h \\ s \rightarrow d \rightarrow e \rightarrow r \\ s \rightarrow d \rightarrow e \rightarrow r \rightarrow f \rightarrow c \\ s \rightarrow d \rightarrow e \rightarrow r \rightarrow f \rightarrow c \\ \end{array}$



Measuring performance of a search algorithm

- ▶ We care about two main (conflicting) metrics
 - Algorithm's resources: running time, memory consumption
 - Quality of solution: Optimal, bounded suboptimal, unbounded
- ▶ Typically, the more resources we invest, the higher quality the solution we can obtain



Algorithm's resources

Running time

- What we actually care about
- Problematic metric sensitive to hardware, compiler, optimization, implementation language etc.
- Alternative measure: The number of nodes expanded



Algorithm's resources

Running time

- What we actually care about
- Problematic metric sensitive to hardware, compiler, optimization, implementation language etc.
- Alternative measure: The number of nodes expanded

Memory consumption

- Critical resource may deem a search algorithm worthless
- An algorithm that expands 1,000,000 nodes / second and uses 10 bytes / node can run out of memory in less than 2 minutes



Quality of solution

- We will want to minimize the cost of a solution
- ▶ E.g., plan the path for a robot; action costs correspond to the robot's energy consumption; minimize total energy consumption





General Tree Search (recap)

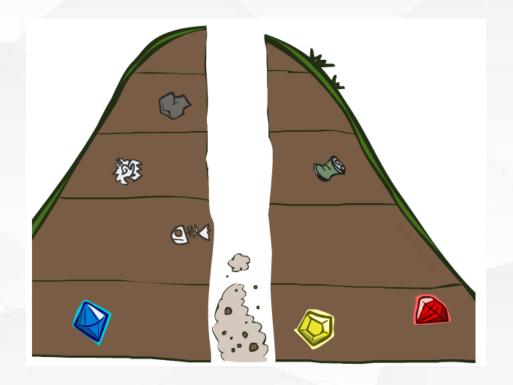
```
function TREE-SEARCH (problem, strategy) returns a solution, or failure
initialize the search tree using the initial state of problem
loop do
    if there are no candidates for expansion then return failure
     choose a leaf node for expansion according to strategy
    if the node contains a goal state then return the corresponding solution
    else expand the node and add the resulting nodes to the search tree
end
```

- Important ideas:
 - Fringe
 - **Expansion**
 - **Exploration strategy**
- Main question: which fringe nodes to explore?





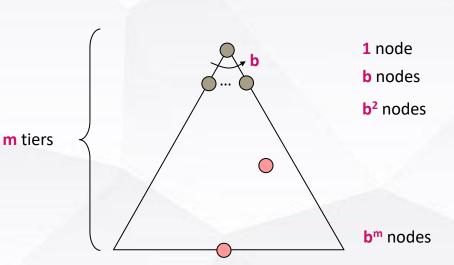
Search Algorithm Properties



.ill

Search Algorithm Properties

- Complete: Guaranteed to find a solution if one exists?
- Optimal: Guaranteed to find the least cost path?
- Time complexity?
- Space complexity?
- Cartoon of search tree:
 - b is the branching factor
 - m is the maximum depth
 - solutions at various depths
- Number of nodes in entire tree?
 - $1 + b + b^2 + \ldots + b^m = O(b^m)$



Breadth-First Search







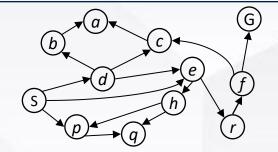
The Henry and Marilyn Taub **Faculty of Computer Science**

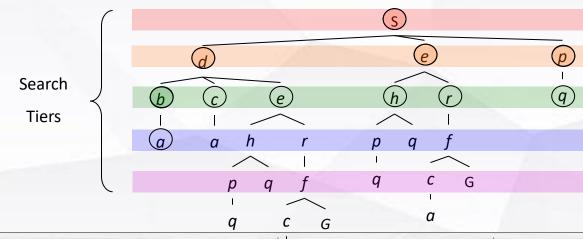


Breadth-First Search

Strategy: expand a shallowest node first

Implementation: Fringe is a FIFO queue





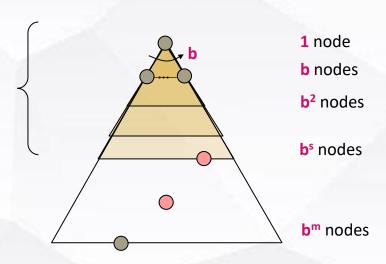


Breadth-First Search (BFS) Properties

- What nodes does BFS expand?
 - Processes all nodes above shallowest solution
 - Let depth of shallowest solution be s
 - Search takes time O(b^s)

s tiers

- How much space does the fringe take?
 - Has roughly the last tier, so O(b^s)
- Is it complete?
 - s must be finite if a solution exists
- Is it optimal?
 - Only if costs are all 1 (more on costs later)





BFS Memory requirements

▶ Assume that BFS expands 1,000,000 nodes / sec for a problem with branching factor b=10 and each node uses 1000 bytes node

Depth	No. of nodes	Time	Memory
2	111	0.1 ms	11 KB
4	11,111	11ms	11 MB
6	~10^6	1 sec	~1 GB
8	~10^8	100 sec	~100 GB
10	~10^10	3 hours	~10 TB ←
12	~10^12	11 days	~1 PB
14	~10^14	3 years	~100 PB

Depth-First Search





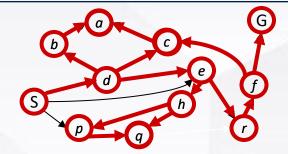


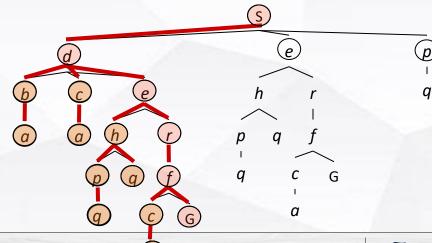
The Henry and Marilyn Taub **Faculty of Computer Science**



Strategy: expand a deepest node first

Implementation: Fringe is a LIFO stack

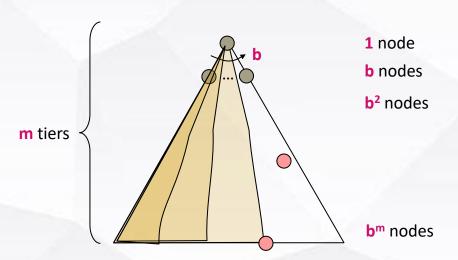






Depth-First Search (DFS) Properties

- What nodes DFS expand?
 - Some left prefix of the tree.
 - Could process the whole tree!
 - If m is finite, takes time O(b^m)
- How much space does the fringe take?
 - Only has siblings on path to root, so O(bm)
- Is it complete?
 - m could be infinite
 - To avoid this, we can call the algorithm with a bound L on the maximum depth – this is called DFS-L
- Is it optimal?
 - No, it finds the "leftmost" solution, regardless of depth or cost





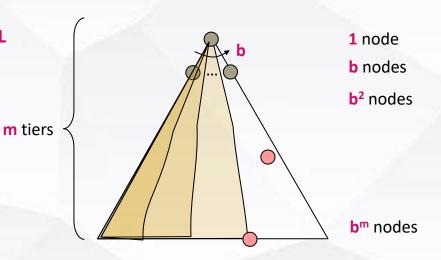
Depth-First Search Variants

▶ DFS-L

Same as DFS but with a maximum tree depth L

▶ DFS-L-BT

- Same as **DFS-L** but with the following modification:
- At every node, we run an operator O over the same node
- When we backtrack, we undo the previous operator
- We only need to store one node in memory (with some additional book keeping)





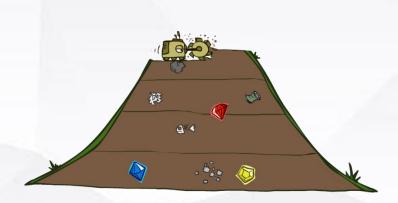






Quiz: DFS vs BFS







Video of Demo Maze Water DFS/BFS (part 1)





Video of Demo Maze Water DFS/BFS (part 2)





	BFS	DFS	DFS-L
Complete	YES	NO*	NO**
Optimal#	YES	NO	NO
Memory	Exponential in solution length	Infinite	Linear in L



	BFS	DFS	DFS-L
Complete	YES	NO*	NO**
Optimal#	YES	NO	NO
Memory	Exponential in solution length	Infinite	Linear in L

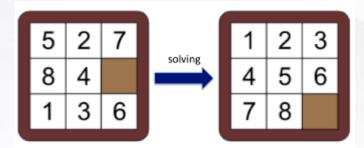
Assuming uniform cost edges

^{*} If tree is unbounded, can run infinitely even if a solution exists

^{**} If L is less than the solution depth, no solution is returned even though one exists

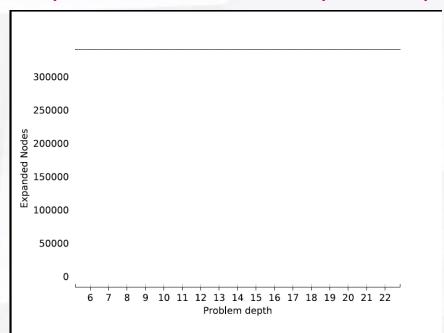


- ▶ Python implementation by Shaul Markovitz of BFS, DFS-L and DFS-L-BT
- Random 3X3 puzzles whose solution-lengths range between 6 and 22 (20 problem instances per solution length)
- ▶ DFS-L was run with L=22 (for all instances)



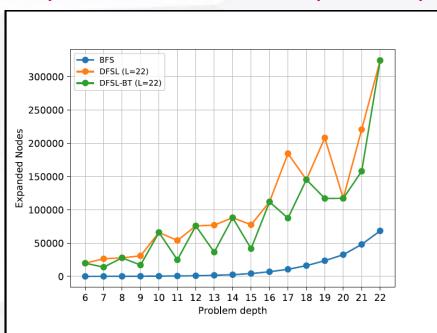


Expanded nodes as a function of **problem depth**

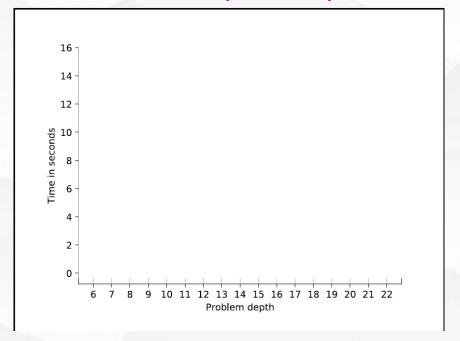




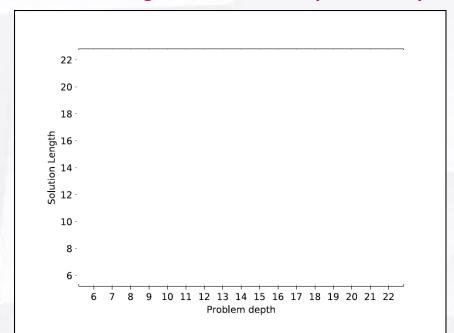
Expanded nodes as a function of **problem depth**



▶ Time as a function of problem depth



Solution length as a function of problem depth





Memory footprint as a function of problem depth

