# **EX NO: 1** Implementation of Lexical Analyzer

Date:

#### Aim:

To develop a lexical analyzer capable of tokenizing a subset of a programming language, identifying tokens such as identifiers, keywords, literals, and punctuation symbols.

#### Algorithm:

- 1. Start
- 2. **Input**: Take the source code as input.
- 3. **Initialize**: Initialize an empty list/array to store tokens.
- 4. **Tokenization Loop**: Iterate through each character in the source code.
  - If the character is a whitespace, skip it.
  - If the character is a digit, start collecting characters until a non-digit or non-decimal point character is encountered. Add the collected characters as a token of type "NUMBER" to the token list.
  - If the character is an alphabet or underscore, start collecting characters until a non-alphanumeric or underscore character is encountered. Add the collected characters as a token of type "IDENTIFIER" to the token list.
  - If the character is an operator or delimiter (e.g., +, -, \*, /, =, (, ), etc.), add it as a token to the token list.
- 5. **Output**: Return the list of tokens.
- 6. End

```
Source Code:
#include <stdio.h>
#include <stdib.h>
#include <string.h>
#include <ctype.h>

#define MAX_CODE_LENGTH 1000
#define MAX_TOKEN_LENGTH 100

// Token types
typedef enum {
    NUMBER,
    IDENTIFIER,
```

```
ASSIGN,
  PLUS,
  MINUS,
  MULTIPLY,
  DIVIDE,
  LPAREN,
  RPAREN,
  NEWLINE,
  WHITESPACE,
  END OF FILE
} TokenType;
// Token structure
typedef struct {
  TokenType type;
  char value[MAX TOKEN LENGTH];
} Token;
// Function to tokenize the input code
Token* tokenize(const char* code) {
  Token* tokens = malloc(strlen(code) * sizeof(Token));
  if (!tokens) {
    fprintf(stderr, "Memory allocation error\n");
    exit(EXIT FAILURE);
  }
  int i = 0;
  while (*code) {
    if (isspace(*code)) {
       code++;
       continue; // Skip whitespace
    if (isdigit(*code)) {
       char* start = code;
       while (isdigit(*code) || *code == '.') {
         code++;
       int len = code - start;
       tokens[i].type = NUMBER;
```

```
strncpy(tokens[i].value, start, len);
  tokens[i].value[len] = '\0'; // Null-terminate string
} else if (isalpha(*code) || *code == ' ') {
  char* start = code;
  while (isalnum(*code) \parallel *code == ' ') {
     code++;
  int len = code - start;
  tokens[i].type = IDENTIFIER;
  strncpy(tokens[i].value, start, len);
  tokens[i].value[len] = '\0'; // Null-terminate string
} else {
  switch (*code) {
     case '=':
       tokens[i].type = ASSIGN;
       strncpy(tokens[i].value, "=", 1);
       break;
     case '+':
       tokens[i].type = PLUS;
       strncpy(tokens[i].value, "+", 1);
       break:
     case '-':
       tokens[i].type = MINUS;
       strncpy(tokens[i].value, "-", 1);
       break;
     case '*':
       tokens[i].type = MULTIPLY;
       strncpy(tokens[i].value, "*", 1);
       break;
     case '/':
       tokens[i].type = DIVIDE;
       strncpy(tokens[i].value, "/", 1);
       break;
     case '(':
       tokens[i].type = LPAREN;
       strncpy(tokens[i].value, "(", 1);
       break;
     case ')':
       tokens[i].type = RPAREN;
```

```
strncpy(tokens[i].value, ")", 1);
            break;
         case '\n':
            tokens[i].type = NEWLINE;
            strncpy(tokens[i].value, "\n", 1);
            break;
         default:
            fprintf(stderr, "Illegal character: %c\n", *code);
            exit(EXIT FAILURE);
       code++;
       tokens[i].value[1] = '\0'; // Null-terminate string
    i++;
  tokens[i].type = END OF FILE;
  tokens[i].value[0] = '\0'; // Null-terminate string
  return tokens;
}
// Function to free memory allocated for tokens
void free tokens(Token* tokens) {
  free(tokens);
// Function to print token type
const char* token type to string(TokenType type) {
  switch(type) {
    case NUMBER: return "NUMBER";
    case IDENTIFIER: return "IDENTIFIER";
    case ASSIGN: return "ASSIGN";
    case PLUS: return "PLUS";
    case MINUS: return "MINUS";
    case MULTIPLY: return "MULTIPLY";
    case DIVIDE: return "DIVIDE";
    case LPAREN: return "LPAREN";
    case RPAREN: return "RPAREN";
```

```
case NEWLINE: return "NEWLINE";
    case WHITESPACE: return "WHITESPACE";
    case END OF FILE: return "END OF FILE";
    default: return "UNKNOWN";
}
// Test the lexer
int main() {
  char code[MAX CODE LENGTH];
  printf("Enter code (max %d characters):\n", MAX CODE LENGTH - 1);
  if (fgets(code, MAX CODE LENGTH, stdin) == NULL) {
    fprintf(stderr, "Failed to read input\n");
    return EXIT FAILURE;
  Token* tokens = tokenize(code);
  for (int i = 0; tokens[i].type != END OF FILE; i++) {
    printf("Token: Type=%s, Value=%s\n", token type to string(tokens[i].type),
tokens[i].value);
  free tokens(tokens);
  return EXIT SUCCESS;
}
```

#### Output:

```
Enter code (max 999 characters):

x = 10 + 20 y = x * 2

Token: Type=IDENTIFIER, Value=x

Token: Type=ASSIGN, Value==

Token: Type=NUMBER, Value=10

Token: Type=PLUS, Value=+

Token: Type=NUMBER, Value=20

Token: Type=IDENTIFIER, Value=y

Token: Type=ASSIGN, Value==

Token: Type=ASSIGN, Value=x

Token: Type=MULTIPLY, Value=*

Token: Type=NUMBER, Value=2
```

#### **Result:**

Thus, a program a lexical analyzer capable of tokenizing a subset of a programming language, identifying tokens such as identifiers, keywords, literals, and punctuation symbols was implemented

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# **EX NO: 2** Conversion from Regular Expression to NFA

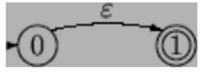
Date:

# Aim:

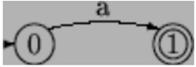
Program to convert Regular Expression (R.E.) to Non-Deterministic Finite Automata (N.F.A.)

# **Algorithm:**

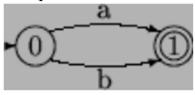
1. The NFA representing the empty string is:



2. If the regular expression is just a character, eg. a, then the corresponding NFA is :



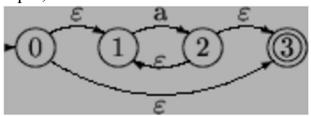
3. The union operator is represented by a choice of transitions from a node; thus a|b can be represented as:



4. Concatenation simply involves connecting one NFA to the other; eg. ab is:



5. The Kleene closure must allow for taking zero or more instances of the letter from the input; thus a\* looks like:



#### **Source Code:**

```
#include <stdio.h>
#include <string.h>
#include <ctype.h>
int ret[100];
static int pos = 0;
static int sc = 0;
void nfa(int st, int p, char *s) {
  int i,sp,fs[15],fsc=0;
 sp=st;pos=p;sc=st;
 while(*s!=NULL)
 if(isalpha(*s))
 ret[pos++]=sp;
 ret[pos++]=*s;
 ret[pos++]=++sc;}
 if(*s=='.')
 {
 sp=sc;
 ret[pos++]=sc;
 ret[pos++]=238;
 ret[pos++]=++sc;
 sp=sc;}
if(*s=='|')
 {
 sp=st;
 fs[fsc++]=sc;}
if(*s=='*')
 ret[pos++]=sc;
 ret[pos++]=238;
 ret[pos++]=sp;
ret[pos++]=sp;
 ret[pos++]=238;
 ret[pos++]=sc;
```

```
if (*s=='(')
 {char ps[50];
 int i=0,flag=1;
 s++;
   while(flag!=0)
  \{ps[i++]=*s;
  if (*s=='(')
  flag++;
  if (*s==')')
  flag--;
  s++;
  ps[--i]='\setminus 0';
  nfa(sc,pos,ps);
  s--;
 s++;
sc++;
 for(i=0;i<fsc;i++)
  \{ret[pos++]=fs[i];
  ret[pos++]=238;
  ret[pos++]=sc;
  }
  ret[pos++]=sc-1;
  ret[pos++]=238;
  ret[pos++]=sc;
}
int main() {
  int i;
  char inp[50]; // Changed from pointer to array
  printf("Enter the regular expression: ");
  fgets(inp, sizeof(inp), stdin);
  inp[strcspn(inp, "\n")] = '\0'; // Remove newline character if present
  nfa(1, 0, inp);
```

```
printf("\nState Input State\n");
for (i = 0; i < pos; i = i + 3)
    printf("%d --%c--> %d\n", ret[i], ret[i + 1], ret[i + 2]);
return 0;
}
```

# **Output:**

```
Enter the regular expression: a+b*
State Input State
     --a-->
1
     --b-->
                 3
                 1
3
     --2-->
                 3
1
     --ε-->
3
                 4
     ---->
Process returned 0 (0x0) execution time: 16.745 s
Press any key to continue.
```

## **Result:**

Thus, a program to convert Regular Expression (R.E.) to Non-Deterministic Finite Automata (N.F.A.) was implemented

## EX NO: 3 Conversion from NFA to DFA

Date:

#### Aim:

Program to convert NFA to Deterministic Finite Automata(D.F.A.)

# Algorithm:

- 1. Convert into NFA using above rules for operators (union, concatenation and closure) and precedence.
- 2. Find E -closure of all states.
- 3. Start with epsilon closure of start state of NFA.
- 4. Apply the input symbols and find its epsilon closure. Dtran[state, input symbol] =  $\varepsilon$  closure(move(state, input symbol)) where Dtran transition function of DFA
- 5. Analyze the output state to find whether it is a new state.
- 6. If new state is found, repeat step 4 and step 5 until no more new states are found.
- 7. Construct the transition table for Dtran function.
- 8. Draw the transition diagram with start state as the E -closure (start state of NFA) and final state is the state that contains final state of NFA drawn.

#### **Source Code:**

```
void push(int val)
stack[++top]=val;
int pop()
return stack[top--];}
int priority(char op)
switch(op)
case '+': return 1;
case '.': return 2;
case '*': return 3;
}return 0;
void init nfa table()
int i;
for(i=0; i<20; i++)
 NFA[i].a = NFA[i].b = -1;
 NFA[i].eps1 = NFA[i].eps2 = -1;
void symbol(char c)
if(c=='a')
 NFA[cur].a = cur+1;
if(c=='b')
 NFA[cur].b = cur+1;push(cur);
push(cur+1);
```

```
cur += 2;
}//=====
void concat()
int first1, first2, last1, last2;
last2 = pop();
first2 = pop();
last1 = pop();
first1 = pop();
NFA[last1].eps1 = first2;
push(first1);
push(last2);
void parallel()
int first1, first2, last1, last2; last2 = pop();
first2 = pop();
last1 = pop();
first1 = pop();
NFA[cur].eps1 = first1;
NFA[cur].eps2 = first2;
NFA[last1].eps1 = cur+1;
NFA[last2].eps2 = cur+1;
push(cur);
push(cur+1);
cur += 2;
void closure()
int first,last;last = pop();
first = pop();
```

```
NFA[cur].eps1 = first;
NFA[cur].eps2 = cur+1;
NFA[last].eps1 = first;
NFA[last].eps2 = cur+1;
push(cur);
push(cur+1);
cur += 2;
void construct nfa(char *postfix)
int i=0;
top=-1;
for(i=0; postfix[i]!='\0'; i++)
 switch(postfix[i])
 case 'a':
 case 'b': symbol(postfix[i]);
   break;
 case '.': concat();
   break;
 case '+': parallel();
   break;
 case '*': closure(); }
final state = pop();
initial state = pop();
void disp NFA(){
int i;
printf("\nstate\ta\tb\t?");
for(i=0;i<cur;i++)
 if(i==initial state)
```

```
printf("n->\%d",i);
 else
 if(i==final_state)
 printf("\n* %d",i);
 else
 printf("\n %d",i);
 if(NFA[i].a==-1)
 printf("\t-");
 else
 printf("\t{%d}",NFA[i].a);
 if(NFA[i].b==-1)printf("\t-");
 else
 printf("\t{\mbox{$\%$d}}",NFA[i].b);
 if(NFA[i].eps1!=-1)
 printf("\t{%d",NFA[i].eps1);if(NFA[i].eps2!=-1)
  printf(",%d",NFA[i].eps2);
 printf("}");
 else
 printf("\t-");
void init dfa table()
int i,j;
for(i=0;i<20;i++)
 for(j=0;j<20;j++)
 DFA[i].state[j]=-1;
 DFA[i].a[j]=-1;
 DFA[i].b[j]=-1;
```

```
void print state(int t[])
int i=0;
printf("[");
for(i=0;t[i]!=-1;i++)
 printf("%d,",t[i]);
printf("\b]");
int isPresent(int T[], int v)
int i;
for(i=0;T[i]!=-1;i++)
 if(T[i]=v)
 return 1;
return 0;
} //====
void disp_DFA(int n)
int i;
printf("\nstate\t\ta\t\tb");
for(i=0;i \le n;i++)
 printf("\n");
 if(i==0)
 printf("->");
if(isPresent(DFA[i].state,final_state))
  printf("*");
 print state(DFA[i].state);
 printf("\t\t");
 if(DFA[i].a[0]!=-1)
 print_state(DFA[i].a);
 else
```

```
printf("\t-");
 printf("\t\t");
 if(DFA[i].b[0]!=-1)
 print_state(DFA[i].b);
 else
 printf("\t-");
void epsilon_closure(int T[], int t[])
int i,v;
top=-1;
for(i=0;t[i]!=-1;i++)
 push(t[i]);
i=0;
while(top!=-1)
 v = pop();
 if(isPresent(T,v)==0)
 T[i++]=v;
 if(NFA[v].eps1!=-1)
 push(NFA[v].eps1);
 if(NFA[v].eps2!=-1)
 push(NFA[v].eps2);
```

```
void init t(int t[])
{int i;
for(i=0;i<20;i++)
 t[i]=-1;
}
int search(int n,int t2[])
int i,j;
for(i=0;i<=n;i++)
 {
 for(j=0;t2[j]!=-1;j++)
 if(isPresent(DFA[i].state,t2[j])==0)
  break;
 if(t2[j]=-1)
 return 1;
return 0;
void copy(int t1[], int t2[])
int i;
for(i=0;t2[i]!=-1;i++)
 t1[i]=t2[i];
void main()
char postfix[20];
int t[20],v;
int n=0, i=0, j, k;
system("cls");
printf("\nEnter Regular Expression: ");
scanf("%s",postfix);
```

```
printf("\nPostfix Expression: %s",postfix);
getch();
init_nfa_table();
construct nfa(postfix);
system("cls");
disp NFA();
getch();
init dfa table();
init t(t);t[0]=initial state;
epsilon closure(DFA[0].state,t);
init t(t);
for(j=0,k=0; DFA[0].state[j]!=-1; j++)
v = DFA[0].state[j];
if(NFA[v].a!=-1)
if(isPresent(t,NFA[v].a)==0)
 t[k++]=NFA[v].a;
epsilon closure(DFA[0].a,t);
init t(t);
for(j=0,k=0;DFA[0].state[j]!=-1;j++)
v = DFA[0].state[j];if(NFA[v].b!=-1)
\{if(isPresent(t,NFA[v].b)==0)
 t[k++]=NFA[v].b;
```

```
epsilon closure(DFA[0].b,t);
for(i=0;i \le n;i++)
if( search( n, DFA[i].a)==0)
 n++;
 copy(DFA[n].state,DFA[i].a);
 init t(t);
 for( j=0,k=0; DFA[n].state[j]!=-1; j++)
 v = DFA[n].state[j];
 if(NFA[v].a!=-1)
  if(isPresent(t,NFA[v].a)==0)
  t[k++]=NFA[v].a;
 } epsilon closure(DFA[n].a,t);
 init_t(t);
 for(j=0,k=0;DFA[n].state[j]!=-1;j++)
 v = DFA[n].state[j];
 if(NFA[v].b!=-1)
  if(isPresent(t,NFA[v].b)==0)
  t[k++]=NFA[v].b;
 epsilon closure(DFA[n].b,t);
}
if( search( n, DFA[i].b ) == 0)
```

```
n++;
 copy(DFA[n].state,DFA[i].b);
 init t(t);
 for( j=0,k=0; DFA[n].state[j]!=-1; j++)
  v = DFA[n].state[j];
if( NFA[v].a!=-1)
   if(isPresent(t,NFA[v].a)==0)
   t[k++]=NFA[v].a;
 epsilon_closure(DFA[n].a,t);
 init t(t);
 for(j=0,k=0;DFA[n].state[j]!=-1;j++)
  v = DFA[n].state[j];
  if(NFA[v].b!=-1)
   if(isPresent(t,NFA[v].b)==0)
   t[k++]=NFA[v].b;
 epsilon_closure(DFA[n].b,t);}
disp DFA(n);
 getch();
```

#### **OUTPUT:**

```
Enter Regular Expression: ab+*

Postfix Expression: ab+*

state a b ?
0 {1} - - {5}
1 - {5}
2 - {3} - - {3}
3 - - - {4,7}
--6 - - {4,7}
--6 - - {4,7}
* 7 - - - *

state a b ?
0 {1} - 0 {5}
2 - {3} - 0 {5}
3 - 0 - 0 {6,2}
5 - 0 {6,2}
5 - 0 {6,2}
5 - 0 {6,2}
5 - 0 {6,2}
5 - 0 {6,7,4,2,0}
5 - 0 {6,7,4,2,0}
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8 - 0 {6,7,4,2,0}
8 - 0 {6,7,4,
```

# **Result:**

Thus, a Program to convert NFA to Deterministic Finite Automata(D.F.A.) was implemented.

# **EX NO: 4** Elimination of Ambiguity, Left Recursion and Left Factoring Date:

#### Aim:

A program for Elimination of Left Recursion and implementation Of Left Factoring

# Algorithm:

## **Left Recursion Elimination:**

- 1. Start the program.
- 2. Initialize the arrays for taking input from the user.
- 3. Prompt the user to input the no. of non-terminals having left recursion and no. of productions for these non-terminals.
- 4. Prompt the user to input the production for non-terminals.
- 5. Eliminate left recursion using the following rules:-

A->A
$$\alpha$$
1| A $\alpha$ 2 | . . . . . |A $\alpha$ m  
A-> $\beta$ 1|  $\beta$ 2| . . . . . |  $\beta$ n  
Then replace it by  
A->  $\beta$ i A' i=1,2,3,....m  
A'->  $\alpha$ j A' j=1,2,3,....n

- 6. After eliminating the left recursion by applying these rules, display the productions without left recursion.
- 7. Stop.

# **Left Factoring:**

- 1. Start
- 2. Ask the user to enter the set of productions
- 3. Check for common symbols in the given set of productions by comparing with:
  - 1. A->aB1|aB2
- 4. If found, replace the particular productions with:
  - a. A->aA'
  - b. A'->B1 | B2|ε
- 5. Display the output
- 6. Exit

#### **Source Code:**

# **<u>Left Recursion Elimination:</u>** Left Recursion Elimination.cpp

```
#include <iostream>
#include <vector>
#include <string>
using namespace std;
int main()
int n;
cout << "\nEnter number of non terminals: ";
cin>>n;
cout<<"\nEnter non terminals one by one: ";
int i;
vector<string> nonter(n);
vector\leqint\geq leftrecr(n,0);
for(i=0;i< n;++i) {
cout << "\nNon terminal " << i+1 << ": ";
cin>>nonter[i];
 }
vector<vector<string> > prod;
cout << "\nEnter '^' for null";
for(i=0;i< n;++i) {
cout << "\nNumber of "<< nonter[i] << " productions: ";
int k;
cin>>k;
int j;
cout<<"\nOne by one enter all "<<nonter[i]<<" productions";
vector<string> temp(k);
for(j=0;j< k;++j) {
cout << "\nRHS of production "<< j+1 << ": ";
string abc;
cin>>abc;
temp[j]=abc;
if(nonter[i].length()<=abc.length()&&nonter[i].compare(abc.substr(0,nonter[i].length()))==0)
leftrecr[i]=1;
prod.push back(temp);
for(i=0;i<n;++i) {
```

```
cout<<leftrecr[i];
for(i=0;i<n;++i) {
if(leftrecr[i]==0)
continue;
int j;
nonter.push_back(nonter[i]+""");
vector<string> temp;
for(j=0;jjprod[i].size();++j) {
if(nonter[i].length() \!\!<\!\!=\! prod[i][j].length() \& \& nonter[i].compare(prod[i][j].substr(0,nonter[i].length()) \\
h
()))==0) {
string
abc=prod[i][j].substr(nonter[i].length(),prod[i][j].length()-nonter[i].length())+nonter[i]+"";
temp.push back(abc);
prod[i].erase(prod[i].begin()+j);
--j;
else {
prod[i][j]+=nonter[i]+""";
temp.push back("^");
prod.push_back(temp);
cout << "\n\n";
cout << "\nNew set of non-terminals: ";
for(i=0;i<nonter.size();++i)
cout << nonter[i] << " ";
cout << "\n\nNew set of productions: ";
for(i=0;i<nonter.size();++i) {
int j;
for(j=0;jjprod[i].size();++j) {
cout<<"\n"<<nonter[i]<<" -> "<<pre>prod[i][j];
return 0;
```

## **Left Factoring:** Left Factoring.cpp

```
#include <iostream>
#include <math.h>
#include <vector>
#include <string>
#include <stdlib.h>
using namespace std;
int main()
cout<<"\nEnter number of productions: ";</pre>
int p;
cin>>p;
vector<string> prodleft(p),prodright(p);
cout << "\nEnter productions one by one: ";
int i;
for(i=0;i< p;++i) {
cout<<"\nLeft of production "<<i+1<<": ";</pre>
cin>>prodleft[i];
cout<<"\nRight of production "<<i+1<<": ";</pre>
cin>>prodright[i];
int j;
int e=1;
for(i=0;i< p;++i) {
for(j=i+1;j< p;++j)  {
if(prodleft[j]==prodleft[i]) {
int k=0;
string com="";
while(k<prodright[i].length()&&k<prodright[j].length()&&prodright[i][k]==prodright[j][k]) {
com+=prodright[i][k];
++k;
if(k==0)
continue;
char* buffer;
string comleft=prodleft[i];
if(k==prodright[i].length()) {
prodleft[i]+=string(itoa(e,buffer,10));
prodleft[i]+=string(itoa(e,buffer,10));
```

```
prodright[i]="^";
prodright[j]=prodright[j].substr(k,prodright[j].length()-k);
else if(k==prodright[j].length()) {
prodleft[i]+=string(itoa(e,buffer,10));
prodleft[i]+=string(itoa(e,buffer,10));
prodright[j]="^";
prodright[i]=prodright[i].substr(k,prodright[i].length()-k);
else {
prodleft[i]+=string(itoa(e,buffer,10));
prodleft[i]+=string(itoa(e,buffer,10));
prodright[j]=prodright[j].substr(k,prodright[j].length()-k);
prodright[i]=prodright[i].substr(k,prodright[i].length()-k);
int 1;
for(1=j+1;1< p;++1) {
if(comleft==prodleft[1]&&com==prodright[1].substr(0,fmin(k,prodright[1].length()))) {
prodleft[1]+=string(itoa(e,buffer,10));
prodright[1]=prodright[1].substr(k,prodright[1].length()-k);
prodleft.push back(comleft);
prodright.push back(com+prodleft[i]);
++p;
++e;
cout << "\n\nNew productions";
for(i=0;i<p;++i) {
cout<<"\n"<<pre>prodleft[i]<<"->"<<pre>prodright[i];
return 0;
```

#### **SAMPLE OUTPUT:**

#### **Left Recursion Elimination:**

```
Enter number of non terminals: 3
Enter non terminals one by one:
Non terminal 1 : E
Non terminal 2 : T
Non terminal 3 : F
Enter '^' for null
Number of E productions: 2
One by one enter all E productions
RHS of production 1: E+T
RHS of production 2: T
Number of T productions: 2
One by one enter all T productions
RHS of production 1: T*F
RHS of production 2: F
Number of F productions: 2
One by one enter all F productions
RHS of production 1: (E)
RHS of production 2: I
110
New set of non-terminals: E T F E' T'
New set of productions:
E -> TE'
T -> FT'
F -> (E)
F -> Î
E' -> +TE'
E' -> ^
T' -> *FT'
T' -> ^
Process returned 0 (0x0)
                           execution time: 40.299 s
Press any key to continue.
```

# **Left Factoring:**

```
Enter number of productions: 3

Enter productions one by one:
Left of production 1: A

Right of production 1: abcd

Left of production 2: B

Right of production 2: abef

Left of production 3: C

Right of production 3: abgh

New productions
A->abcd
B->abef
C->abgh
Process returned 0 (0x0) execution time: 197.964 s
Press any key to continue.
```

#### **Result:**

Thus, program for Elimination of Left Recursion and implementation Of Left Factoring was implemented.

## EX NO: 5 FIRST AND FOLLOW

Date:

#### Aim:

To write a program to perform first and follow using any language.

# Algorithm:

#### FIRST(X) for all grammar symbols X:

- 1. If X is terminal,  $FIRST(X) = \{X\}$ .
- 2. If  $X \to \varepsilon$  is a production, then add  $\varepsilon$  to FIRST(X).
- 3. If X is a non-terminal, and  $X \to Y1 \ Y2 \dots Yk$  is a production, and  $\varepsilon$  is in all of FIRST(Y1), ..., FIRST(Yk), then add  $\varepsilon$  to FIRST(X).
- 4. If X is a non-terminal, and  $X \to Y1 \ Y2 \dots Yk$  is a production, then add a to FIRST(X) if for some i, a is in FIRST(Yi), and  $\varepsilon$  is in all of FIRST(Y1), ..., FIRST(Yi-1).

#### FOLLOW(A) for all non-terminals A:

- 1. If \$ is the input end-marker, and S is the start symbol,  $\$ \in FOLLOW(S)$ .
- 2. If there is a production,  $A \to \alpha B\beta$ , then  $(FIRST(\beta) \epsilon) \subseteq FOLLOW(B)$ .
- 3. If there is a production,  $A \to \alpha B$ , or a production  $A \to \alpha B\beta$ , where  $\epsilon \in FIRST(\beta)$ , then  $FOLLOW(A) \subseteq FOLLOW(B)$ .

#### Code:

```
#include<stdio.h>
#include<stdlib.h>
#include<conio.h>
#include<String.h>
int n,m=0,p,i=0,j=0;
char a[10][10],f[10];
void follow(char c);
void first(char c);
int main(){
int i,z;
char c,ch;
printf("Enter the no of productions:\n");
scanf("%d",&n);
printf("Enter the productions:\n");
for(i=0;i< n;i++)
scanf("%s%c",a[i],&ch);
do{
m=0;
```

```
printf("Enter the elements whose first & follow is to be found:");
scanf("%c",&c);
first(c);
printf("First(%c)={",c);
for(i=0;i<m;i++)
printf("%c",f[i]);
printf("}\n");
strcpy(f," ");
fflush(stdin);
m=0;
follow(c);
printf("Follow(%c)={",c);
for(i=0;i<m;i++)
printf("%c",f[i]);
printf("}\n");
printf("Continue(0/1)?");
scanf("%d%c",&z,&ch);
\}while(z==1);
return(0);
void first(char c)
int k;
if(!isupper(c))
f[m++]=c;
for(k=0;k< n;k++)
if(a[k][0]==c)
if(a[k][2]=='\$')
follow(a[k][0]);
else if(islower(a[k][2]))
f[m++]=a[k][2];
else first(a[k][2]);
void follow(char c)
```

```
if(a[0][0]==c)
f[m++]='$';
for(i=0;i<n;i++)
{
  for(j=2;j<strlen(a[i]);j++)
  {
   if(a[i][j]==c)
  {
   if(a[i][j+1]!='\0')
   first(a[i][j+1]);
   if(a[i][j+1]=='\0' && c!=a[i][0])
  follow(a[i][0]);
  }
}
}</pre>
```

#### **SAMPLE OUTPUT:**

```
Enter the no of productions:
Enter the productions:
S=AbCd
A=Cf
A=a
C=gE
E=h
Enter the elements whose first & follow is to be found:S
First(S)={ga}
Follow(S)={$}
Continue(0/1)?1
Enter the elements whose first & follow is to be found:A
First(A)={ga}
Follow(A)={b}
Continue(0/1)?1
Enter the elements whose first & follow is to be found:C
First(C)={g}
Follow(C)={df}
Continue(0/1)?1
Enter the elements whose first & follow is to be found:E
First(E)={h}
Follow(E)={df}
Continue(0/1)?0
```

#### **Result:**

Thus, program to perform first and follow using any language was implemented

# EX NO: 6 PREDICTIVE PARSING

Date:

# Aim:

Program for the construction of predictive parsing table.

# Algorithm:

- 1. Start the program.
- 2. Initialize the required variables.
- 3. Get the number of coordinates and productions from the user.
- 4. Perform the following

```
for (each production A \to \alpha in G) {
for (each terminal a in FIRST(\alpha))
add A \to \alpha to M[A, a];
if (\epsilon is in FIRST(\alpha))
for (each symbol b in FOLLOW(A))
add A \to \alpha to M[A, b];
```

- 5. Print the resulting stack.
- 6. Print if the grammar is accepted or not.
- 7. Exit the program.

#### Code:

```
PREDICTIVE PARSING: predictive_pasing.c
#include<stdio.h>
#include<string.h>
#include<ctype.h> // Include ctype.h for isupper()

void main() {
    char st[10][20], ft[20][20], fol[20][20];
    int n, i, j, k, l;

printf("Enter the number of non-terminals: ");
    scanf("%d", &n);
    printf("Enter the productions in the grammar:\n");
    for (i = 0; i < n; i++)
        scanf("%s", st[i]);

for (i = 0; i < n; i++)</pre>
```

```
fol[i][0] = '\0';
// Compute FIRST sets
for (i = 0; i < n; i++)
  i = 3;
  1 = 0;
  while (st[i][i] != '\0') 
     if (!isupper(st[i][j])) {
        ft[i][l++] = st[i][j];
        ft[i][1] = '\0'; // Terminate the string
        break; // Break after encountering a terminal symbol
     } else {
        k = 0;
        while (st[i][j] != st[k][0])
          k++;
        strcat(ft[i], ft[k]);
        if (strchr(ft[k], '@') == NULL) // If '@' (epsilon) is not in the FIRST set
          break; // Break if epsilon is not in the FIRST set of current non-terminal
     j++;
// Compute FOLLOW sets
for (i = 0; i < n; i++)
  for (i = 0; i < n; i++)
     1 = strlen(st[i]);
     for (k = 3; k < 1; k++)
        if(st[i][k] == st[i][0]) {
          if (st[i][k+1] != '\0') {
             if (!isupper(st[i][k+1])) {
                fol[i][strlen(fol[i])] = st[i][k+1];
                fol[i][strlen(fol[i])] = '\0'; // Terminate the string
             } else {
                int x = 0;
                while (st[j][k+1] != st[x][0])
                   X++;
                strcat(fol[i], ft[x]);
             }
```

```
} else {
        int x = 0;
        while (st[j][0] != st[x][0])
            x++;
        strcat(fol[i], fol[x]);
}

// Print FIRST and FOLLOW sets
printf("\nFIRST:\n");
for (i = 0; i < n; i++)
        printf("FIRST[%c] = %s\n", st[i][0], ft[i]);

printf("\nFOLLOW:\n");
for (i = 0; i < n; i++)
        printf("FOLLOW[%c] = %s\n", st[i][0], fol[i]);
}</pre>
```

```
Enter the number of non-terminals: 2
Enter the productions in the grammar:
S->aAb
A->b

FIRST:
FIRST[S] = a
FIRST[A] = b

FOLLOW:
FOLLOW[S] =
FOLLOW[A] = b

Process returned 2 (0x2) execution time : 23.452 s
Press any key to continue.
```

## Result:

```
EX NO: 7
                          Implementation of Shift Reduce Parsing
Date:
Aim:
      To write a program to Implementation of Shift Reduce Parsing
Algorithm:
      Repeat: For each production A \in \in of the grammar do
             For each terminal in FIRST(\in)
                   add A \in \in \text{to } M[A, a]
             if FIRST(\in) contains \in
                   add A \in \in \text{to } M[A, b] for each bin FOLLOW(A)
             if \in is in FIRST(\in) and $ is in FOLLOW(A)
                   add A \in \in \text{ to } M[A,\$]
             make each undefined entry of M be error
Code:
#include<stdio.h>
#include<conio.h>
#include<string.h>
struct prodn
      char p1[10];
      char p2[10];
};
void main()
      char input[20],stack[50],temp[50],ch[2],*t1,*t2,*t;
      int i,j,s1,s2,s,count=0;
      struct prodn p[10];
      FILE *fp=fopen("sr input.txt","r");
      stack[0]='\0';
      clrscr();
      printf("\n Enter the input string\n");
      scanf("%s",&input);
      while(!feof(fp))
       {
             fscanf(fp,"%s\n",temp);
             t1=strtok(temp,"->");
             t2=strtok(NULL,"->");
```

```
strcpy(p[count].p1,t1);
      strcpy(p[count].p2,t2);
      count++;
i=0;
while(1)
      if(i<strlen(input))</pre>
             ch[0]=input[i];
             ch[1]='\0';
             i++;
             strcat(stack,ch);
             printf("%s\n",stack);
      for(j=0;j<count;j++)
             t=strstr(stack,p[j].p2);
             if(t!=NULL)
                    s1=strlen(stack);
                    s2=strlen(t);
                    s=s1-s2;
                    stack[s]='\0';
                    strcat(stack,p[j].p1);
                    printf("%s\n",stack);
                    j=-1;
             }
      if(strcmp(stack,"E")==0&&i==strlen(input))
             printf("\n Accepted");
             break;
      if(i==strlen(input))
       {
             printf("\n Not Accepted");
             break;
       }
```

```
}
getch();
}
Input File: sr_input.txt

E->E+E
E->E*E
E->i
```

```
Enter the input string
i*i+i
i
E
E*
E*i
E*E
E
E+
E+i
E+E
E
T
Accepted
Process returned 13 (0xD) execution time : 9.468 s
Press any key to continue.
```

## Result:

**EX NO: 8 Computation of Leading and Trailing** Date: Aim: To write a program to Computation of Leading and Trailing Algorithm: Leading sets: Repeat: For each production  $A \rightarrow \alpha \beta$  in the grammar: Compute leading set for non-terminal A if  $\alpha$  is a terminal or  $\alpha$  is  $\epsilon$ : Add α to leading[A] else if  $\alpha$  is a non-terminal: Add leading  $[\alpha]$  to leading [A] // Add leading of non-terminal  $\alpha$  to leading of A If  $\alpha$  derives  $\varepsilon$ , compute leading set for  $\beta$  and add it to A's leading set if FIRST( $\alpha$ ) contains  $\epsilon$ : Add FIRST( $\beta$ ) - { $\epsilon$ } to leading[A] Trailing sets Repeat: For each production  $A \rightarrow \alpha\beta$  in the grammar: Compute trailing set for non-terminal A if  $\beta$  is a terminal or  $\beta$  is  $\epsilon$ : Add  $\beta$  to trailing[A] else if  $\beta$  is a non-terminal: Add trailing  $[\beta]$  to trailing [A] // Add trailing of non-terminal  $\beta$  to trailing of A If  $\beta$  derives  $\epsilon$ , compute trailing set for  $\alpha$  and add it to A's trailing set if FIRST( $\beta$ ) contains  $\epsilon$ : Add trailing  $[\alpha]$  to trailing [A]

#### Code:

```
Leading sets:
#include <stdio.h>
#include <stdlib.h>
#include <stdbool.h>
#include <string.h>
#define MAX NONTERMINALS 10
#define MAX TERMINALS 10
typedef struct {
  char symbol;
  int numProductions;
  char productions[10][10];
} NonTerminal;
void computeLeadingSet(NonTerminal nonTerminals[], char
leadingSet[][MAX TERMINALS], int numNonTerminals, int numTerminals);
bool isTerminal(char symbol);
bool is NonTerminal (char symbol);
bool addToLeadingSet(char leadingSet[][MAX TERMINALS], char symbol, char terminal, int
nonTerminalIndex);
int main() {
  int numNonTerminals, numTerminals;
  printf("Enter the number of non-terminals: ");
  scanf("%d", &numNonTerminals);
  printf("Enter the number of terminals: ");
  scanf("%d", &numTerminals);
  NonTerminal nonTerminals[MAX NONTERMINALS];
  for (int i = 0; i < numNonTerminals; i++) {
    printf("Enter non-terminal %d and its productions separated by space: ", i + 1);
    scanf(" %c", &nonTerminals[i].symbol);
    scanf("%d", &nonTerminals[i].numProductions);
    for (int j = 0; j < nonTerminals[i].numProductions; <math>j++) {
       scanf("%s", nonTerminals[i].productions[i]);
```

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```
char leadingSet[MAX NONTERMINALS][MAX TERMINALS];
  computeLeadingSet(nonTerminals, leadingSet, numNonTerminals, numTerminals);
  printf("\nLeading Sets:\n");
  for (int i = 0; i < numNonTerminals; i++) {
    printf("Leading(%c): { ", nonTerminals[i].symbol);
    for (int j = 0; j < numTerminals; j++) {
       if (leadingSet[i][j] != '\0') {
         printf("%c ", leadingSet[i][j]);
    printf("\n");
  return 0;
}
void computeLeadingSet(NonTerminal nonTerminals[], char
leadingSet[][MAX TERMINALS], int numNonTerminals, int numTerminals) {
  for (int i = 0; i < numNonTerminals; i++) {
    memset(leadingSet[i], '\0', sizeof(leadingSet[i]));
  }
  bool changed = true;
  while (changed) {
     changed = false;
    for (int i = 0; i < numNonTerminals; i++) {
       for (int j = 0; j < nonTerminals[i].numProductions; <math>j++) {
          char firstSymbol = nonTerminals[i].productions[i][0];
          if (isTerminal(firstSymbol)) {
            changed |= addToLeadingSet(leadingSet, nonTerminals[i].symbol, firstSymbol, i);
          } else if (isNonTerminal(firstSymbol)) {
            for (int k = 0; k < numNonTerminals; k++) {
              if (nonTerminals[k].symbol == firstSymbol) {
                 for (int l = 0; l < numTerminals; l++) {
                   if (leadingSet[k][1] != '\0') {
```

```
changed |= addToLeadingSet(leadingSet, nonTerminals[i].symbol,
leadingSet[k][l], i);
                if (!nonTerminals[k].productions[0][0]) {
                   changed |= addToLeadingSet(leadingSet, nonTerminals[i].symbol, 'e', i);
                }
   }
}
}
}
bool isTerminal(char symbol) {
  return symbol >= 'a' && symbol <= 'z';
}
bool isNonTerminal(char symbol) {
  return symbol >= 'A' && symbol <= 'Z';
}
bool addToLeadingSet(char leadingSet[][MAX TERMINALS], char symbol, char terminal, int
nonTerminalIndex) {
  for (int i = 0; i < MAX TERMINALS; i++) {
    if (leadingSet[nonTerminalIndex][i] == terminal) {
       return false;
    if (leadingSet[nonTerminalIndex][i] == '\0') {
       leadingSet[nonTerminalIndex][i] = terminal;
       return true;
  return false;
```

```
Trailing sets:
#include <stdio.h>
#include <stdlib.h>
#include <stdbool.h>
#include <string.h>
#define MAX NONTERMINALS 10
#define MAX TERMINALS 10
typedef struct {
  char symbol;
  int numProductions;
  char productions[10][10];
} NonTerminal;
void computeTrailingSet(NonTerminal nonTerminals[], char
trailingSet[][MAX TERMINALS], int numNonTerminals, int numTerminals);
bool isTerminal(char symbol);
bool isNonTerminal(char symbol);
bool addToTrailingSet(char trailingSet[][MAX TERMINALS], char symbol, char
terminal, int nonTerminalIndex);
int main() {
  int numNonTerminals, numTerminals;
  printf("Enter the number of non-terminals: ");
  scanf("%d", &numNonTerminals);
  printf("Enter the number of terminals: ");
  scanf("%d", &numTerminals);
  NonTerminal nonTerminals[MAX NONTERMINALS];
  for (int i = 0; i < numNonTerminals; i++) {
    printf("Enter non-terminal %d and its productions separated by space: ", i + 1);
```

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```
scanf(" %c", &nonTerminals[i].symbol);
     scanf("%d", &nonTerminals[i].numProductions);
     for (int j = 0; j < nonTerminals[i].numProductions; <math>j++) {
       scanf("%s", nonTerminals[i].productions[j]);
    }
  }
  char trailingSet[MAX NONTERMINALS][MAX TERMINALS];
  computeTrailingSet(nonTerminals, trailingSet, numNonTerminals,
numTerminals);
  printf("\nTrailing Sets:\n");
  for (int i = 0; i < numNonTerminals; i++) {
     printf("Trailing(%c): { ", nonTerminals[i].symbol);
     for (int j = 0; j < numTerminals; j++) {
       if (trailingSet[i][j] != '\0') {
         printf("%c ", trailingSet[i][j]);
     printf("\n");
  return 0;
}
void computeTrailingSet(NonTerminal nonTerminals[], char
trailingSet[][MAX TERMINALS], int numNonTerminals, int numTerminals) {
  for (int i = 0; i < numNonTerminals; i++) {
    memset(trailingSet[i], '\0', sizeof(trailingSet[i]));
  }
  bool changed = true;
  while (changed) {
     changed = false;
```

```
for (int i = 0; i < numNonTerminals; i++) {
       for (int j = 0; j < nonTerminals[i].numProductions; <math>j++) {
          int len = strlen(nonTerminals[i].productions[j]);
          char lastSymbol = nonTerminals[i].productions[j][len - 1];
          if (isTerminal(lastSymbol)) {
            changed |= addToTrailingSet(trailingSet, nonTerminals[i].symbol,
lastSymbol, i);
          } else if (isNonTerminal(lastSymbol)) {
            for (int k = 0; k < numNonTerminals; k++) {
               if (nonTerminals[k].symbol == lastSymbol) {
                 for (int l = 0; l < numTerminals; l++) {
                    if (trailingSet[k][1] != '\0') {
                      changed |= addToTrailingSet(trailingSet,
nonTerminals[i].symbol, trailingSet[k][l], i);
                 if (!nonTerminals[k].productions[0][0]) {
                    changed |= addToTrailingSet(trailingSet,
nonTerminals[i].symbol, 'e', i);
bool isTerminal(char symbol) {
  return symbol >= 'a' && symbol <= 'z';
}
bool isNonTerminal(char symbol) {
```

```
return symbol >= 'A' && symbol <= 'Z';
}
bool addToTrailingSet(char trailingSet[][MAX_TERMINALS], char symbol, char
terminal, int nonTerminalIndex) {
    for (int i = 0; i < MAX_TERMINALS; i++) {
        if (trailingSet[nonTerminalIndex][i] == terminal) {
            return false;
        }
        if (trailingSet[nonTerminalIndex][i] == '\0') {
            trailingSet[nonTerminalIndex][i] = terminal;
            return true;
        }
    }
    return false;
}</pre>
```

## Leading set:

```
Enter the number of non-terminals: 2 Enter the number of terminals: 3 Enter non-terminal 1 and its productions separated by space: S 2 Ab c Enter non-terminal 2 and its productions separated by space: A 2 d \varepsilon Leading Sets: Leading (S): { c d } Leading (A): { d }
```

## **Trailing Set:**

```
Enter the number of non-terminals: 2
Enter the number of terminals: 3
Enter non-terminal 1 and its productions separated by space: S 1 Ab
Enter non-terminal 2 and its productions separated by space: A 3 aA bA c

Trailing Sets:
Trailing(S): { b }
Trailing(A): { c }
```

#### Result:

```
EX NO: 9
                        Computation of LR (0) items
Date:
Aim:
      To write a program to Computation of LR (0) items
Algorithm:
ComputeLR0ItemsAndPrint():
  Input:
    numRules <- Prompt user to enter the number of rules
    for i from 0 to numRules - 1:
       rules[i] <- Prompt user to enter the rule in the format 'S -> AB'
  LR0ItemComputationPhase(rules, numRules)
  PrintLR0Items()
Algorithm LR0ItemComputationPhase(rules[], numRules):
  numLR0Items <- 0
  for i from 0 to numRules - 1:
    lr0Items[numLR0Items].ruleIndex <- i</pre>
    lr0Items[numLR0Items].dotPosition <- 0
    numLR0Items++
Algorithm PrintLR0Items():
  for each item in lr0Items:
    Print item with dot before the first symbol of the right-hand side
Code:
#include <stdio.h>
#include <string.h>
#include <stdbool.h>
#define MAX RULES 10
#define MAX SYMBOLS 10
typedef struct {
  char left;
  char right[MAX SYMBOLS];
} ProductionRule;
```

```
typedef struct {
  int ruleIndex;
  int dotPosition;
} LR0Item;
void printLR0Items(ProductionRule rules[], int numRules) {
  int i, j;
  printf("LR(0) Items:\n");
  for (i = 0; i < numRules; i++) {
     printf("[%d] %c -> ", i, rules[i].left);
     for (j = 0; rules[i].right[j] != '\0'; j++) {
       if (i == 0) printf(".");
       printf("%c", rules[i].right[j]);
     if (j == 0) printf(".");
     printf("\n");
void computeLR0Items(ProductionRule rules[], int numRules, LR0Item lr0Items[], int
*numLR0Items) {
  int i;
  *numLR0Items = 0;
  for (i = 0; i < numRules; i++)
     lr0Items[*numLR0Items].ruleIndex = i;
     lr0Items[*numLR0Items].dotPosition = 0;
     (*numLR0Items)++;
}
int main() {
  ProductionRule rules[MAX RULES];
  LR0Item lr0Items[MAX RULES];
  int numRules, numLR0Items;
  printf("Enter the number of rules: ");
  scanf("%d", &numRules);
  printf("Enter the rules in the format 'S -> AB' (use '->' for epsilon):\n");
```

```
for (int i = 0; i < numRules; i++) {
    scanf(" %c -> %s", &rules[i].left, rules[i].right);
}

computeLR0Items(rules, numRules, lr0Items, &numLR0Items);
printLR0Items(rules, numRules);
return 0;
```

```
Enter the number of rules: 3  
Enter the rules in the format 'S \rightarrow AB' (use '\rightarrow' for epsilon): S \rightarrow AB  
A \rightarrow aA \mid \rightarrow B \rightarrow bB \mid \rightarrow LR(0) Items: [0] S \rightarrow . AB  
[1] A \rightarrow . aA  
[2] \mid \rightarrow . B
```

## Result:

**EX NO: 10a** Intermediate core generation - Postfix

Date:

### Aim:

To write a program to generate Intermediate core generation – Postfix

## Algorithm:

infixToPostfix(infix, postfix):

Initialize an empty stack

Initialize an empty string for the postfix expression

For each character 'ch' in the infix expression:

If 'ch' is an operand:

Append 'ch' to the postfix expression

Else If 'ch' is '(':

Push 'ch' onto the stack

Else If 'ch' is ')':

While the stack is not empty and the top of the stack is not '(':

Pop an operator from the stack and append it to the postfix expression

Pop '(' from the stack

Else If 'ch' is an operator:

While the stack is not empty and precedence of 'ch' is less than or equal to precedence of the top of the stack:

Pop an operator from the stack and append it to the postfix expression Push 'ch' onto the stack

While the stack is not empty:

Pop an operator from the stack and append it to the postfix expression

Return the postfix expression

# Main Function:

Input infix expression from the user

Convert infix expression to postfix using infixToPostfix function

Output the resulting postfix expression

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#### Code:

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <ctype.h>
#define MAX EXPR SIZE 100
typedef struct {
  int top;
  char items[MAX EXPR SIZE];
} Stack;
void initialize(Stack *s) {
  s->top = -1;
}
int isEmpty(Stack *s) {
  return s->top == -1;
}
void push(Stack *s, char c) {
  s->items[++(s->top)] = c;
char pop(Stack *s) {
  return s->items[(s->top)--];
}
int isOperator(char c) {
  return (c == '+' || c == '-' || c == '*' || c == '/');
}
int precedence(char op) {
  if (op == '+' || op == '-')
     return 1;
  if (op == '*' || op == '/')
     return 2;
  return 0;
```

```
}
void infixToPostfix(char *infix, char *postfix) {
  Stack stack;
  initialize(&stack);
  int i, j = 0;
  for (i = 0; infix[i]; i++) {
     char ch = \inf x[i];
     if (isalnum(ch)) {
       postfix[j++] = ch;
     } else if (ch == '(') {
       push(&stack, ch);
     } else if (ch == ')') {
       while (!isEmpty(&stack) && stack.items[stack.top] != '(') {
          postfix[j++] = pop(\&stack);
       if (!isEmpty(&stack) && stack.items[stack.top] != '(') {
          printf("Invalid expression\n");
          exit(1);
       } else {
          pop(&stack);
     } else {
       while (!isEmpty(&stack) && precedence(ch) <= precedence(stack.items[stack.top])) {
          postfix[j++] = pop(\&stack);
       push(&stack, ch);
  }
  while (!isEmpty(&stack)) {
     postfix[j++] = pop(\&stack);
  postfix[j] = '\0';
int main() {
  char infix[MAX EXPR SIZE], postfix[MAX EXPR SIZE];
```

```
printf("Enter infix expression: ");
  fgets(infix, sizeof(infix), stdin);
  infix[strcspn(infix, "\n")] = "\0';
  infixToPostfix(infix, postfix);
  printf("Postfix expression: %s\n", postfix);
  return 0;
}
Output:
```

```
Enter infix expression: (A+B)*(C-D)
Postfix expression: AB+CD-*
```

#### **Result:**

## **EX NO:10b** Intermediate Code Generation - Prefix

Date:

## Aim:

To write a program to generate Intermediate Code Generation - Prefix

## Algorithm:

infixToPrefix(infix, prefix):

Initialize an empty stack

Initialize an empty string for the postfix expression

For each character 'ch' in the infix expression:

If 'ch' is an operand:

Append 'ch' to the postfix expression

Else If 'ch' is '(':

Push 'ch' onto the stack

Else If 'ch' is ')':

While the stack is not empty and the top of the stack is not '(':

Pop an operator from the stack and append it to the postfix expression Pop '(' from the stack

Else If 'ch' is an operator:

While the stack is not empty and precedence of 'ch' is less than or equal to precedence of the top of the stack:

Pop an operator from the stack and append it to the postfix expression Push 'ch' onto the stack

While the stack is not empty:

Pop an operator from the stack and append it to the postfix expression

Reverse the postfix expression to obtain the prefix expression

## Main Function:

Input infix expression from the user

Convert infix expression to prefix using infixToPrefix function

Output the resulting prefix expression

## Code:

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```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <ctype.h>
#define MAX EXPR SIZE 100
typedef struct {
  int top;
  char items[MAX EXPR SIZE];
} Stack;
void initialize(Stack *s) {
  s->top = -1;
}
int isEmpty(Stack *s) {
  return s->top == -1;
}
void push(Stack *s, char c) {
  s->items[++(s->top)] = c;
char pop(Stack *s) {
  return s->items[(s->top)--];
}
int isOperator(char c) {
  return (c == '+' \parallel c == '-' \parallel c == '*' \parallel c == '/');
}
int precedence(char op) {
  if (op == '+' || op == '-')
     return 1;
  if (op == '*' || op == '/')
     return 2;
  return 0;
```

```
void infixToPostfix(char *infix, char *postfix) {
  Stack stack;
  initialize(&stack);
  int i, j = 0;
  for (i = 0; infix[i]; i++) {
     char ch = infix[i];
     if (isalnum(ch)) {
       postfix[j++] = ch;
     } else if (ch == '(') {
       push(&stack, ch);
     } else if (ch == ')') {
       while (!isEmpty(&stack) && stack.items[stack.top] != '(') {
          postfix[j++] = pop(\&stack);
       if (!isEmpty(&stack) && stack.items[stack.top] != '(') {
          printf("Invalid expression\n");
          exit(1);
       } else {
          pop(&stack);
     } else {
        while (!isEmpty(&stack) && precedence(ch) <= precedence(stack.items[stack.top])) {
          postfix[j++] = pop(\&stack);
       push(&stack, ch);
  }
  while (!isEmpty(&stack)) {
     postfix[j++] = pop(\&stack);
  postfix[j] = '\0';
```

```
void reverse(char *str) {
  int length = strlen(str);
  for (int i = 0; i < length / 2; i++) {
     char temp = str[i];
     str[i] = str[length - i - 1];
     str[length - i - 1] = temp;
}
int main() {
  char infix[MAX EXPR SIZE], postfix[MAX EXPR SIZE], prefix[MAX EXPR SIZE];
  printf("Enter infix expression: ");
  fgets(infix, sizeof(infix), stdin);
  \inf[x[strcspn(infix, "\n")] = '\0';
  infixToPostfix(infix, postfix);
  reverse(postfix);
  printf("Prefix expression: %s\n", postfix);
  return 0;
```

```
Enter infix expression: (A+B) * (C-D)
Prefix expression: -DC *+BA
```

#### **Result:**

# EX NO:11 Intermediate Code Generation - Quadruple, Triple, Date: Indirect triple

## Aim:

To write a program to generate Intermediate Code Generation - Quadruple

## Algorithm:

The algorithm takes a sequence of three-address statements as input. For each three address statements of the form a: = b or c performs the various actions. These are as follows:

- 1. Invoke a function getreg to find out the location L where the result of computation b op c should be stored.
- 2. Consult the address description for y to determine y'. If the value of y currently in memory and register both then prefer the register y'. If the value of y is not already in L then generate the instruction MOV y', L to place a copy of y in L.
- 3. Generate the instruction OP z', L where z' is used to show the current location of z. if z is in both then prefer a register to a memory location. Update the address descriptor of x to indicate that x is in location L. If x is in L then update its descriptor and remove x from all other descriptors.
- 4. If the current value of y or z have no next uses or not live on exit from the block or in register then alter the register descriptor to indicate that after execution of x := y op z those register will no longer contain y or z.

#### Code:

```
#include<stdio.h>
#include<stdlib.h>
#include<stdlib.h>
#include<string.h>
void small();
void dove(int i);
int p[5]={0,1,2,3,4},c=1,i,k,l,m,pi;
char sw[5]={'=','-','+','/','*'},j[20],a[5],b[5],ch[2];
void main()
{
    printf("Enter the expression:");
    scanf("%s",j);
    printf("\tThe Intermediate code is:\n");
    small();
}
void dove(int i)
```

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```
a[0]=b[0]='\0';
if(!isdigit(j[i+2])&&!isdigit(j[i-2]))
a[0]=j[i-1];
b[0]=j[i+1];
if(isdigit(j[i+2])){
a[0]=j[i-1];
b[0]='t';
b[1]=j[i+2];
if(isdigit(j[i-2]))
b[0]=j[i+1];
a[0]='t';
a[1]=j[i-2];
b[1]='\0';
if(isdigit(j[i+2]) &&isdigit(j[i-2]))
a[0]='t';
b[0]='t';
a[1]=j[i-2];
b[1]=j[i+2];
sprintf(ch,"%d",c);
j[i+2]=j[i-2]=ch[0];
}
if(i[i]=='*')
printf("tt\%d=\%s*\%s\n",c,a,b);
if(i[i]=='/')
printf("tt\%d=\%s/\%s\n",c,a,b);
if(j[i]=='+')
printf("\tt\%d=\%s+\%s\n",c,a,b);if(i[i]=='-')
printf("\tt%d=%s-%s\n",c,a,b);
if(j[i]=='=')
printf("\t%c=t%d",j[i-1],--c);
sprintf(ch,"%d",c);
j[i]=ch[0];
```

```
c++;
small();
}
void small()
{
pi=0;l=0;
for(i=0;i<strlen(j);i++)
{
for(m=0;m<5;m++)
if(j[i]==sw[m])
if(pi<=p[m])
{
pi=p[m];
l=1;
k=i;
}
}
if(l==1)
dove(k);
else
exit(0);}</pre>
```

```
Enter the expression:a=b+c-d
The Intermediate code is:
t1=b+c
t2=t1-d
a=t2
Process returned 0 (0x0) execution time : 17.410 s
Press any key to continue.
```

## **Result:**

# **EX NO:12** A Simple Code Generator

Date:

## Aim:

To write a program to create a Simple Code Generator

## **Algorithm:**

- 1. Start
- 2. Define a function generate\_code(expr) that takes an arithmetic expression as input.
- 3. Within generate code(expr):
  - a. Tokenize the expression using strtok() function.
  - b. Initialize a register counter reg to 1.
  - c. Loop through the tokens:
  - i. Print "MOV R%d, token" where token is the current token and %d represents the current register.
  - ii. Get the next token.
  - iii. If there's a next token:
    - i. Determine the operation (+, -, \*, /) based on the operator preceding the token.
    - ii. Print the corresponding assembly-like operation (ADD, SUB, MUL, DIV) using the current register and the token.
  - iv. Increment the register counter.
- 4. End generate\_code function.
- 5. In the main function:
  - a. Declare a character array expr to store the input expression.
  - b. Prompt the user to input an arithmetic expression and store it in expr using scanf().
  - c. Call generate\_code(expr) to generate the assembly-like code for the input expression.
- **6.** End

## Code: (Python)

```
class CodeGenerator:
    def __init__(self):
        self.code = []

    def emit(self, instruction):
        self.code.append(instruction)

    def generate_code(self, ast):
```

```
if isinstance(ast, int):
       self.emit(f"MOV R1, {ast}")
     elif isinstance(ast, str):
       self.emit(f"MOV R1, {ast}")
     elif isinstance(ast, tuple):
       op, left, right = ast
       self.generate code(left)
       self.emit(f"MOV R2, R1")
       self.generate code(right)
       if op == '+':
          self.emit("ADD R1, R2, R1")
       elif op == '-':
          self.emit("SUB R1, R2, R1")
       elif op == '*':
          self.emit("MUL R1, R2, R1")
       elif op == '/':
          self.emit("DIV R1, R2, R1")
# Function to parse the arithmetic expression
def parse expression(expression):
  # Implementation of your parsing logic goes here
  # For simplicity, let's assume the expression is already parsed into an AST
  return ('+', 3, ('*', 4, 5)) # Example AST for "3 + 4 * 5"
expression = input("Enter an arithmetic expression: ")
ast = parse expression(expression)
generator = CodeGenerator()
generator.generate code(ast)
print("Generated Code:")
for instruction in generator.code:
  print(instruction)
```

```
Enter an arithmetic expression: 3 + 4 * 5
Generated Code:
MOV R1, 3
MOV R2, R1
MOV R1, 4
MOV R2, R1
MOV R2, R1
MOV R1, 5
MUL R1, R2, R1
ADD R1, R2, R1
```

## **Result:**

## EX NO:13 Implementation of DAG

Date:

#### Aim:

To write a program to create a Simple Code Generator

## Algorithm:

- 1. The leaves of a graph are labeled by a unique identifier and that identifier can be variable names or constants.
- 2. Interior nodes of the graph are labeled by an operator symbol.
- 3. Nodes are also given a sequence of identifiers for labels to store the computed value.
- 4. If y operand is undefined then create node(y).
- 5. If z operand is undefined then for case(i) create node(z).
- 6. For case(i), create node(OP) whose right child is node(z) and left child is node(y).
- 7. For case(ii), check whether there is node(OP) with one child node(y).
- 8. For case(iii), node n will be node(y).
- 9. For node(x) delete x from the list of identifiers. Append x to attached identifiers list for the node n found in step 2. Finally set node(x) to n.

#### Code:

```
#include <stdio.h>
#include <ctype.h>
#include <stdlib.h>
#include <string.h>
void small();
void dove(int i);
int p[5] = \{0, 1, 2, 3, 4\}, c = 1, i, k, l, m, pi;
char sw[5] = \{'=', '-', '+', '/', '*'\}, j[20], a[5], b[5], ch[2];
int main() {
  printf("Enter the expression:");
  scanf("%s", j);
  printf("\tThe Intermediate code is:\n");
  small();
  return 0;
}
void dove(int i) {
```

```
a[0] = b[0] = '\0';
if (!isdigit(j[i + 2]) && !isdigit(j[i - 2])) {
  a[0] = i[i - 1];
  b[0] = j[i + 1];
}
if (isdigit(j[i+2])) {
  a[0] = j[i - 1];
  b[0] = 't';
  b[1] = j[i + 2];
}
if (isdigit(j[i - 2])) {
  b[0] = i[i + 1];
  a[0] = 't';
  a[1] = j[i - 2];
  b[1] = '\0';
if (isdigit(j[i+2]) && isdigit(j[i-2])) {
  a[0] = 't';
  b[0] = 't';
  a[1] = j[i - 2];
  b[1] = i[i + 2];
  sprintf(ch, "%d", c);
  j[i+2] = j[i-2] = ch[0];
if(i[i] == '*')
  printf("\tt\%d=\%s*\%s\n", c, a, b);
if([i] == '/')
  printf("tt\%d=\%s/\%s\n", c, a, b);
if(i[i] == '+')
  printf("\tt%d=%s+%s\n", c, a, b);
if(i[i] == '-')
  printf("\tt%d=%s-%s\n", c, a, b);
if (j[i] == '=')
  printf("\t%c=t%d\n", j[i - 1], --c);
sprintf(ch, "%d", c);
j[i] = ch[0];
c++;
small();
```

}

```
Enter the expression:a=b+c-d
The Intermediate code is:
a=t0
t1=t0+c
t2=t1-d
```

## **Result:**

# EX NO:14 Implementation of global data flow analysis Date:

## Aim:

To write a C program to implement Data Flow Analysis

## Algorithm:

- 1. Start the Program Execution.
- 2. Read the total Numbers of Expression
- 3. Read the Left and Right side of Each Expressions
- 4. Display the Expressions with Line No
- 5. Display the Data flow movement with Particular Expressions
- 6. Stop the Program Execution.

## **Code:** (c programming)

```
#include <stdio.h>
#include <string.h>
struct op {
  char 1[20];
  char r[20];
} op[10];
int main() {
  int i, j, n, lineno = 1;
  char *match;
  printf("Enter the number of statements: ");
  scanf("%d", &n);
  for (i = 0; i < n; i++) {
     printf("Statement %d (left): ", i + 1);
     scanf("%s", op[i].l);
     printf("Statement %d (right): ", i + 1);
     scanf("%s", op[i].r);
  }
  printf("\nIntermediate Code:\n");
  for (i = 0; i < n; i++)
     printf("Line No=%d\n", lineno++);
```

```
printf("\t\t\%s = %s\n", op[i].l, op[i].r);
}

printf("\n*** Data Flow Analysis for the Above Code ***\n");
for (i = 0; i < n; i++) {
    for (j = 0; j < n; j++) {
        match = strstr(op[j].r, op[i].l);
        if (match != NULL) {
            printf("\n%s is live at %s\n", op[i].l, op[j].r);
        }
    }
}

return 0;</pre>
```

```
Enter the number of statements: 3
Statement 1 (left): x
Statement 1 (right): y+z
Statement 2 (left): y
Statement 2 (right): a*b
Statement 3 (left): z
Statement 3 (right): p-q
Intermediate Code:
Line No=1
                        x = y+z
Line No=2
                        v = a*b
Line No=3
                        z = p-q
*** Data Flow Analysis for the Above Code ***
y is live at y+z
z is live at y+z
```

#### **Result:**

# EX NO:15 Implement any one storage allocation strategies Date: (heap, stack, static)

## Aim:

To implement Stack storage allocation strategies using C program.

## Algorithm:

- 1. Initially check whether the stack is empty
- 2. Insert an element into the stack using push operation
- 3. Insert more elements onto the stack until stack becomes full
- 4. Delete an element from the stack using pop operation
- 5. Display the elements in the stack
- 6. Top the stack element will be displayed

## **Code:** (C programming)

```
#include<stdio.h>
#include<stdlib.h>
#define TRUE 1
#define FALSE 0
typedef struct Heap
  int data;
  struct Heap *next;
} node;
node *create();
void display(node *);
node *search(node *,int);
node *insert(node *);
void dele(node **);
int main()
{
  int choice, val;
  char ans;
  node *head = NULL;
  do
```

```
{
     printf("\nprogram to perform various operations on heap using dynamic memory
management");
     printf("\n1.create");
     printf("\n2.display");
     printf("\n3.insert an element in a list");
     printf("\n4.delete an element from list");
     printf("\n5.quit");
     printf("\nenter your choice(1-5): ");
     scanf("%d", &choice);
     switch(choice)
       case 1:
          head = create();
          break;
       case 2:
          display(head);
          break;
       case 3:
          head = insert(head);
          break;
       case 4:
          dele(&head);
          break;
       case 5:
          exit(0);
          break;
       default:
          printf("invalid choice, try again\n");
          break;
  while(choice != 5);
  return 0;
node* create()
```

```
node *temp, *New, *head;
int val;
char ans='y';
temp = NULL;
head = NULL;
do
  printf("\n enter the element: ");
  scanf("%d", &val);
  New = (node*)malloc(sizeof(node));
  if(New == NULL)
    printf("\nMemory is not allocated\n");
    return NULL;
  New->data = val;
  New->next = NULL;
  if(head == NULL)
    head = New;
    temp = head;
  }
  else
    temp->next = New;
    temp = New;
  }
  printf("\nDo you want to enter more elements? (y/n): ");
  scanf(" %c", &ans);
while(ans == 'y' \parallel ans == 'Y');
printf("\nThe list is created\n");
return head;
```

```
void display(node *head)
  node *temp;
  temp = head;
  if(temp == NULL)
    printf("\nThe list is empty\n");
    return;
  }
  while(temp != NULL)
    printf("%d -> ", temp->data);
    temp = temp->next;
  printf("NULL\n");
node *search(node *head, int key)
  node *temp;
  temp = head;
  if(temp == NULL)
    printf("The linked list is empty\n");
    return NULL;
  }
  while(temp != NULL)
    if(temp->data == key)
       printf("\nThe element is present in the list\n");
       return temp;
    temp = temp->next;
```

```
}
  printf("The element is not present in the list\n");
  return NULL;
}
node *insert(node *head)
  int choice;
  node *insert head(node *);
  void insert last(node *);
  void insert after(node *);
  printf("\n1. Insert a node as a head node");
  printf("\n2. Insert a node as the last node");
  printf("\n3. Insert a node at an intermediate position in the list");
  printf("\nEnter your choice for insertion of node: ");
  scanf("%d", &choice);
  switch(choice)
     case 1:
       head = insert head(head);
        break;
     case 2:
       insert last(head);
       break;
     case 3:
       insert after(head);
       break;
     default:
       printf("Invalid choice\n");
        break;
  return head;
}
node *insert head(node *head)
{
```

```
node *New, *temp;
  int val;
  New = (node*)malloc(sizeof(node));
  if(New == NULL)
  {
    printf("Memory allocation failed\n");
    return head;
  printf("\nEnter the element which you want to insert: ");
  scanf("%d", &val);
  New->data = val;
  New->next = head;
  head = New;
  return head;
}
void insert last(node *head)
{
  node *New, *temp;
  int val;
  New = (node*)malloc(sizeof(node));
  if(New == NULL)
  {
    printf("Memory allocation failed\n");
    return;
  }
  printf("\nEnter the element which you want to insert: ");
  scanf("%d", &val);
  New->data = val;
  New->next = NULL;
  temp = head;
```

```
while(temp->next != NULL)
    temp = temp->next;
  temp->next = New;
}
void insert after(node *head)
{
  int key;
  node *New, *temp;
  New = (node*)malloc(sizeof(node));
  if(New == NULL)
    printf("Memory allocation failed\n");
    return;
  }
  printf("\nEnter the element which you want to insert: ");
  scanf("%d", &New->data);
  printf("\nEnter the element after which you want to insert the node: ");
  scanf("%d", &key);
  temp = head;
  while(temp != NULL)
    if(temp->data == key)
       New->next = temp->next;
       temp->next = New;
       return;
    temp = temp->next;
  printf("Element %d not found in the list\n", key);
```

```
}
void dele(node **head)
  node *temp, *prev;
  int key;
  temp = *head;
  if(temp == NULL)
     printf("\nThe list is empty\n");
     return;
  }
  printf("\nEnter the element you want to delete: ");
  scanf("%d", &key);
  temp = search(*head, key);
  if(temp != NULL)
     prev = *head;
     if(prev == temp)
       *head = temp->next;
       free(temp);
       printf("\nThe element is deleted\n");
       return;
     while(prev->next != temp)
       prev = prev->next;
     prev->next = temp->next;
     free(temp);
     printf("\nThe element is deleted\n");
```

}

#### **Output:**

```
program to perform various operations on heap using dynamic memory management
1.create
2.display
3.insert an element in a list
4.delete an element from list
5.quit
enter your choice(1-5): 1
enter the element: 10
Do you want to enter more elements? (y/n): y
enter the element: 20
Do you want to enter more elements? (y/n): y
enter the element: 30
Do you want to enter more elements? (y/n): n
The list is created
program to perform various operations on heap using dynamic memory management
1.create
2.display
3.insert an element in a list
4.delete an element from list
5.quit
enter your choice(1-5): 2
10 -> 20 -> 30 -> NULL
program to perform various operations on heap using dynamic memory management
1.create
2.display
3.insert an element in a list
4.delete an element from list
5.quit
enter your choice(1-5): 3

    Insert a node as a head node

2. Insert a node as the last node
3. Insert a node at an intermediate position in the list
Enter your choice for insertion of node: 2
Enter the element which you want to insert: 40
```

```
program to perform various operations on heap using dynamic memory management
1.create
2.display
3.insert an element in a list
4.delete an element from list
5.quit
enter your choice(1-5): 2
10 -> 20 -> 30 -> 40 -> NULL
program to perform various operations on heap using dynamic memory management
1.create
2.display
3.insert an element in a list
4.delete an element from list
5.quit
enter your choice(1-5): 5
Process returned 0 (0x0)
                           execution time : 113.236 s
Press any key to continue.
```

#### **Result:**