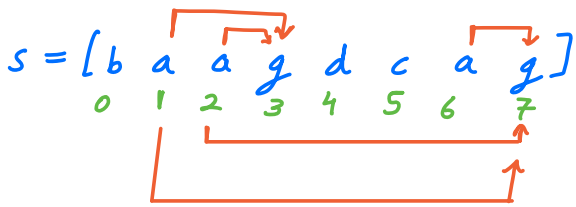
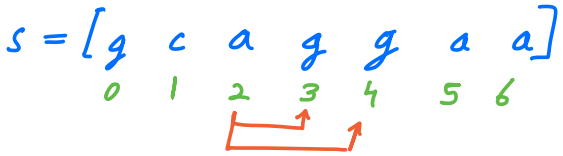
Q → Given a character array, calculate no. of pairs s.t. $(i<j)$ & $s[i] = 'a'$ & $s[j] = 'g'$.
In input all characters are lowercase. (a, b, c, d ... z)



(1, 3) (2, 3) (6, 7)
(1, 7) (2, 7)
Ans = 5



(2, 3) (2, 4) Ans = 2

Bruteforce → ∀ i, j pairs check if $(i<j)$ & $s[i] = 'a'$ & $s[j] = 'g'$.
TC = $O(N^2)$ SC = $O(1)$

```
ans = 0
for i → 0 to (N-2)
    for j → (i+1) to (N-1)
        if (s[i] == 'a' && s[j] == 'g')
            ans++
return ans
```

i = N-1
j → (N-1) + 1 = N Error! ✗

Observations → 1) ∵ given i < j ∴ start j from (i+1). → TC = $O(N^2)$

```
for (j = i+1 ; j <= (N-1) ; j++) {
    ....
}
```

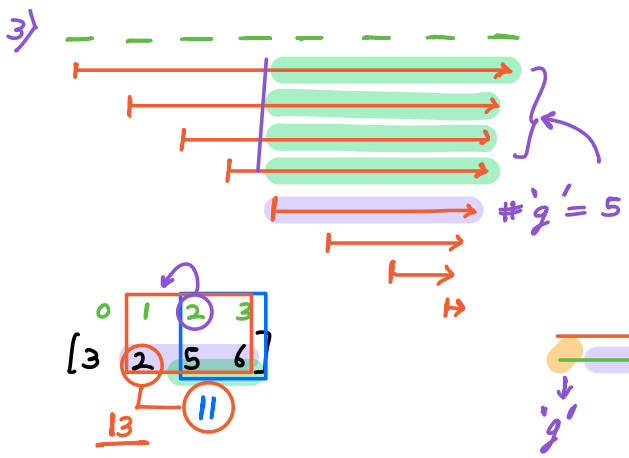
2) ∵ s[i] == 'a' ∴ only loop j when i^{th} char is 'a'.

```
ans = 0
for i → 0 to (N-2)
    if (s[i] == 'a')
        for j → (i+1) to (N-1)
            if (s[j] == 'g')
                ans++
return ans
```

TC = O(N²) SC = O(1)

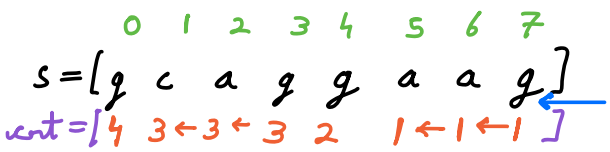
s = [a a a a a ... a]

→ counting # 'g' from index (i+1) to (N-1). ✓



Sum of elements from i to (N-1) ∀i.


```
cnt[i] = cnt[i+1]
if (s[i] == 'g')  cnt[i] += 1
```



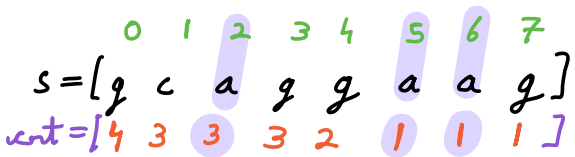
∀i cnt[i] = 0

```
if (s[N-1] == 'g')
    cnt[N-1] = 1
for i → (N-2) to 0
    cnt[i] = cnt[i+1]
    if (s[i] == 'g')
        cnt[i] += 1
```

TC = O(N)
SC = O(N)

```
if (s[i] == 'g')
    cnt[i] = 1 + cnt[i+1]
else
    cnt[i] = cnt[i+1]
```

// cnt[i] → count of 'g' from i to (N-1).



Ans = 3+1+1 = 5 ✓
(2, 3), (2, 4), (2, 7), (5, 7) (6, 7)

```
ans = 0
for i → 0 to (N-1)
    if (s[i] == 'a')
        ans += cnt[i]
return ans
```

TC = O(N)
SC = O(1)

$TC = O(N+N) = O(N)$

$SC = O(N+1) = \boxed{O(N)}$

```
ans = 0
cnt = 0
for i → (N-1) to 0

    if (s[i] == 'g')
        cnt += 1
    else if (s[i] == 'a')
        ans += cnt
return ans
```

	0	1	2	3	4	5	6	7
	✓	✓	✓	✓	✓	✓	✓	✓
s =	g	c	a	g	g	a	a	g

cnt = ~~0+1+1+2~~ 4

ans = 0 + 1 + 1 + 3

= 5

Carry Forward

R to L

$TC = O(N)$

$SC = O(1)$

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Q→ Given an integer array A, count the number of leaders in the array.

Leader → An element which is greater than all elements on right of it.

→ Note → A[N-1] is always a leader.

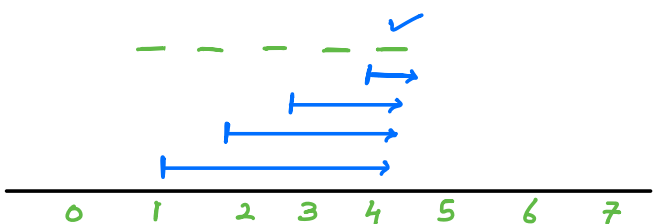
A = [15 -1 7 2 5 4 -2 3] Ans = 5

	0	1	2	3	4	5	6	7
	✓	X	✓	X	✓	✓	X	✓

Leader = A[i] if (A[i] > max element from (i+1) to (N-1))

Bruteforce → ∀i check if A[i] > max element on right.

$TC = O(N^2)$ $SC = O(1)$



	0	1	2	3	4	5	6	7
A =	15	-1	7	2	5	4	-2	3
max →	15	7	7	5	5	4	3	3
	✓		✓		✓	✓		✓

ans = ~~1+2+3~~ 4 5

```
ans = 1
m = A[N-1]                          TC = O(N)
for i → (N-2) to 0                  SC = O(1)
    if(A[i] > m)
        ans += 1
        m = A[i]
return ans
```

Subarray → continuous part of array.

	0	1	2	3	4	5	6	7
A =	15	-1	7	2	5	4	-2	3

Single Element → ✓
Complete Array → ✓

Empty Array → empty subarray

Q → Given an integer array, find the length of smallest subarray which contains both min & max of array.

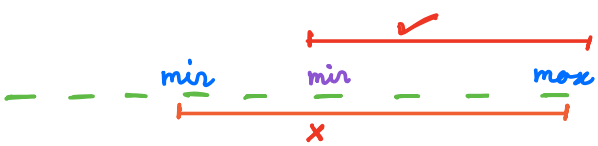
	0	1	2	3	4	5	6	7	8	9	10
A =	1	2	3	1	3	4	6	4	6	3	5

min = 1
max = 6
Ans = 4 (length)

A = [8 8 8 8 8]
min = 8
max = 8
Ans = 1

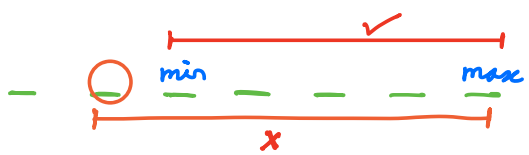
Observation → 1) Can smallest subarray containing min & max have more than 1 min element or max element?

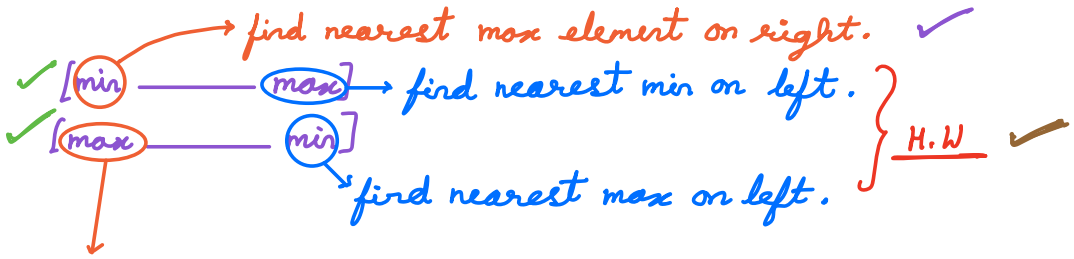
No



2) Can the ans subarray have element in corner which is neither min nor max?

No





A = [2 2 6 4 5 1 5 2 3 6 4 3 1 2 6 3 6]
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16

min = 1
max = 6
TC = O(N)
SC = O(1)

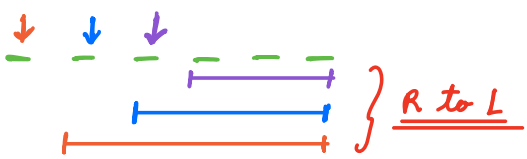
[L R] = R-L+1

(i, j)
(2, 5)
(5, 9)
(9, 12)
(12, 14)

Ans (length) → 5-2+1 = 4
9-5+1 = 5
12-9+1 = 4
14-12+1 = 3 ✓

A = [2 6 1 1 3 6]
0 1 2 3 4 5
min = 1
max = 6
Ans = 2

Sol1 → ∀i if A[i] is min or max travel right from i to get closest max or min. ✓
TC = O(N²) SC = O(1)



A = [2 6 4 1 5 2 6 3 4 6 5 3 2 1 1]
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14

min = 1
max = 6
minId = ~~1 11 13~~ 3
maxId = ~~1 9 6~~ 1

(9, 13) 13-9+1 = 5
(6, 13) 13-6+1 = 8
(3, 6) 6-3+1 = 4
(1, 3) 3-1+1 = 3
length
Ans

TC = O(N) SC = O(1) ✓
ans = N → 1
→ minId = -1 maxId = -1
// calculate minA , maxA ✓ → 10

```
for i ⟶ (N-1) to 0
→ if(A[i] == minA) {
        minId = i   → 0
        if(maxId != -1)  ✗
            ans = min(ans, maxId - minId + 1)
} → if(A[i] == maxA) {
        maxId = i   → 0
        if(minId != -1)  ✓
            ans = min(ans, minId - maxId + 1)
                  1   ,   0 - 0 + 1  = 1  ✓
}
return ans
```

[8 8 8 8 8] ✓

[10] i = 0
 0
