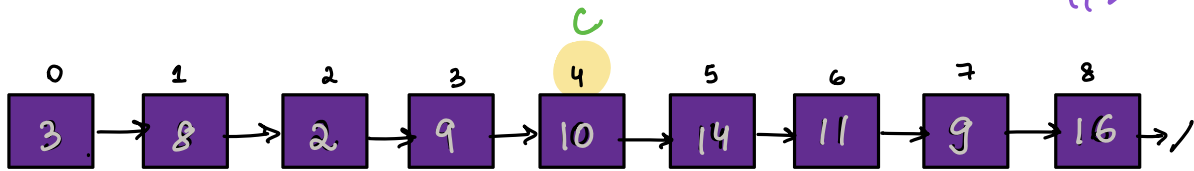


Find Middle

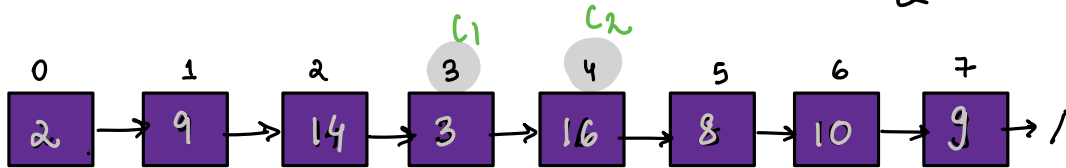
$$\frac{N-1}{2} = \frac{9-1}{2} = \frac{8}{2} = 4 //$$

$$9/2 = 4$$



$$\frac{8-1}{2} = \frac{7}{2} = 3$$

$$8/2 = 4$$



C & C2
(odd) (even)

① Find size of the LL

② $K = \text{size}/2$

getKthEle(head, size/2)

C & C1
(odd) (even)

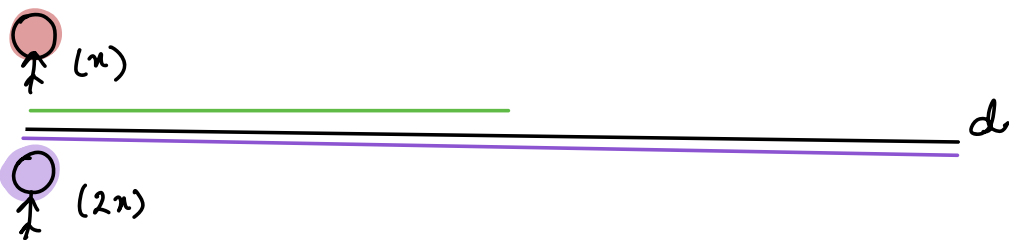
① Find size of LL

② $K = (\text{size}-1)/2$

getKthEle(head, (size-1)/2)

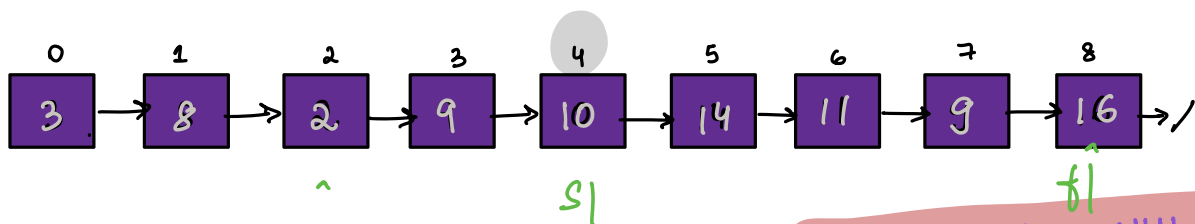
TC: $O(N)$

Q. If odd \rightarrow center
If even \rightarrow C1
In single iteration

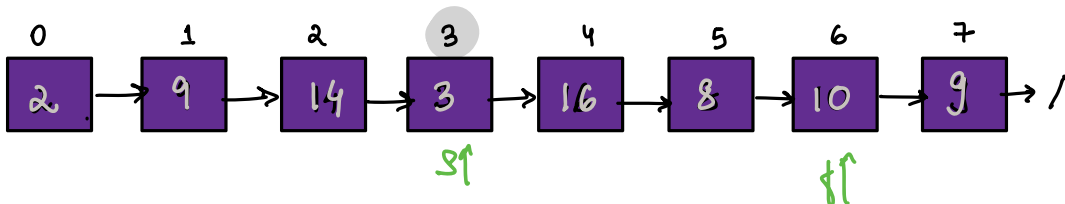


slow
 \downarrow
by 1

fast
 \downarrow
by 2 steps



if ($f \cdot \text{next} == \text{NULL}$)



if ($f \cdot \text{next} \cdot \text{next} == \text{NULL}$)

Node getMiddleEle(head){

if(head == NULL) return head

s = f = head

while(f.next != NULL && f.next.next != NULL){

 s = s.next

 f = f.next.next

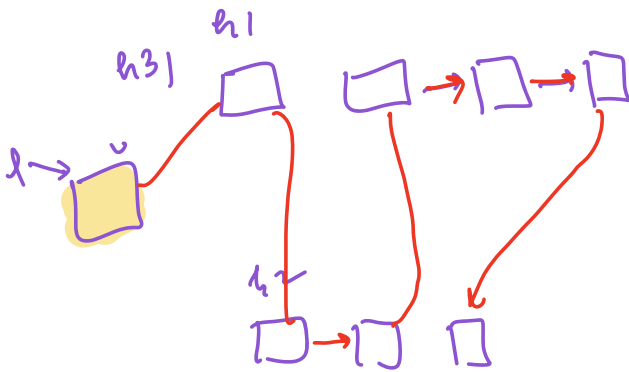
}

return s

head == NULL ✓

□ → ✓

□ → □ → / ✓



Que. Merge two sorted L.L.

h3
-1

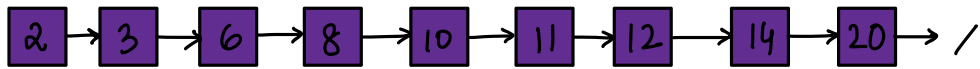
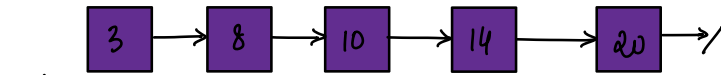
h3 &

2

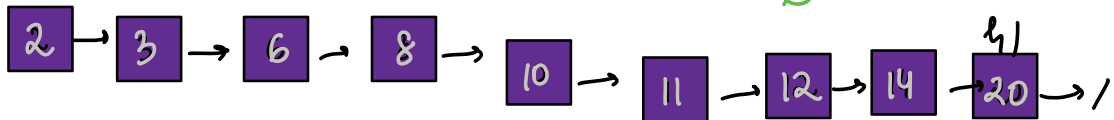
h1

h2

o/p



h3



l

h1

```
Node merge ( h1, h2 ) {
```

```
    if ( h1 == NULL ) return h2
```

```
    if ( h2 == NULL ) return h1
```

```
    if ( h1->data < h2->data ) {
```

```
        h3 = h1;          h1 = h1->next
```

```
    } else { h3 = h2;          h2 = h2->next }
```

```
    while ( h1 != NULL && h2 != NULL ) {
```

```
        if ( h1->data < h2->data ) {
```

```
            l->next = h1
```

```
            h1 = h1->next
```

```
        } else {
```

```
            l->next = h2
```

```
            h2 = h2->next
```

```
        } l = l->next
```

```
    } if ( h1 == NULL ) {
```

```
        l->next = h2
```

```
    }
```

```
    if ( h2 == NULL ) {
```

```
        l->next = h1
```

```
    }
```

```
    return h3
```

```
}
```

Node h3 = new Node (-1)
l = h3

TC :

$O(N+M)$

SC :

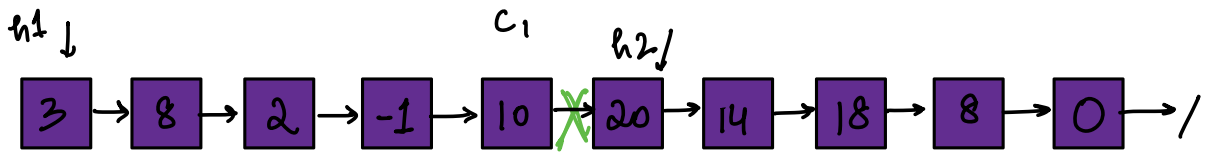
$O(1)$

h3 . next

* If online contests,
okay to use
dummy node
+ else avoid using it

MergeSort

Assumption: your function will return the sorted LL



```
Node mergeSort ( head ) {
```

```
    if ( head == NULL || head.next == NULL ) return head .
```

```
    h1 = head
```

```
    Node c1 = getMiddleEle ( head ) O(N)
```

```
    h2 = c1.next
```

```
    c1.next = NULL
```

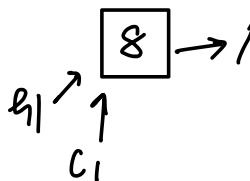
```
    h1 = mergeSort ( h1 )
```

```
    h2 = mergeSort ( h2 )
```

```
    return merge ( h1 , h2 ) O(N)
```

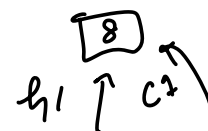
```
}
```

head == NULL



h2 = NULL

mergeSort (h1)



h2 = NULL

$$T(N) = 2T(N/2) + N$$

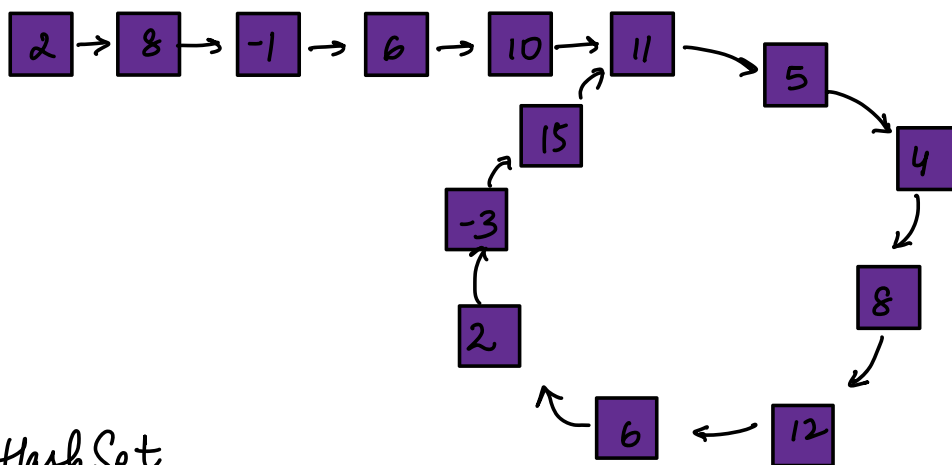
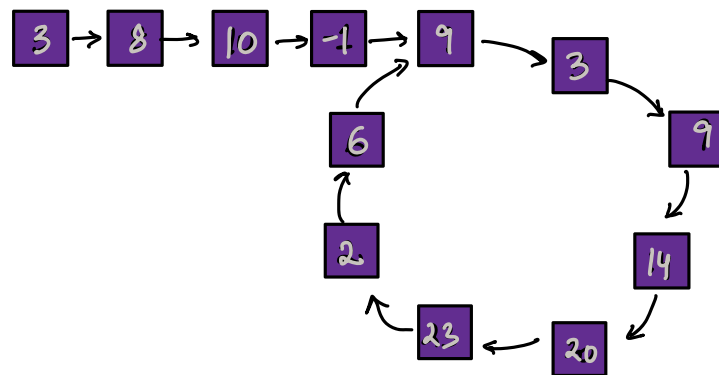
TC: $O(N \log N)$

SC: $O(\log N)$ stack space

Break : 10:35

Detect Cycle

HashSet<Node> s;

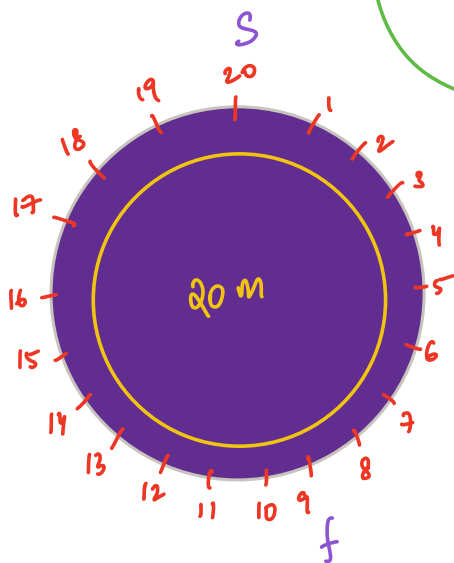
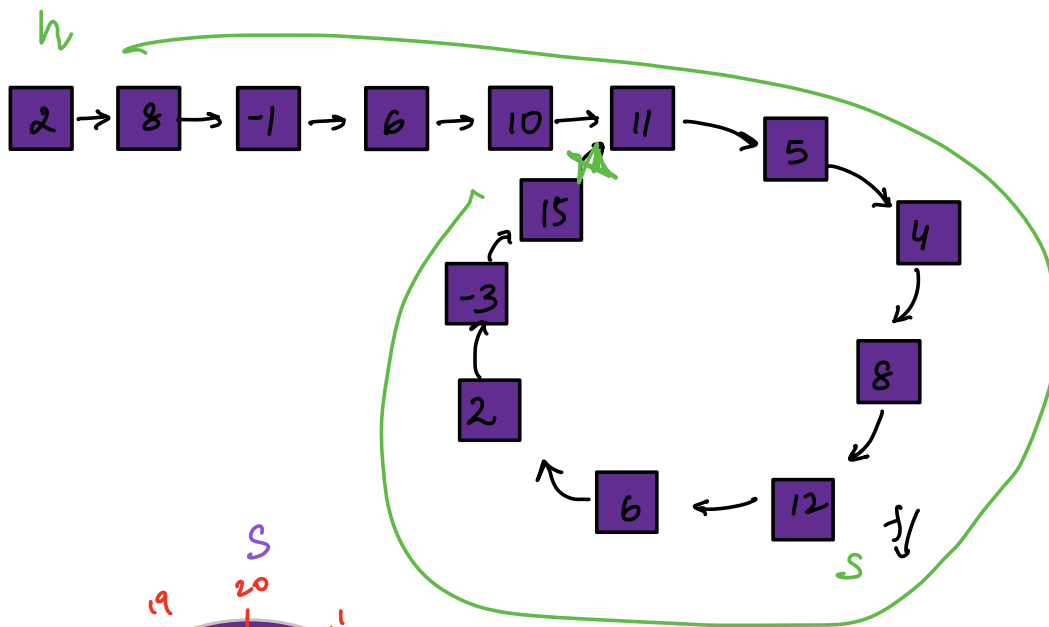
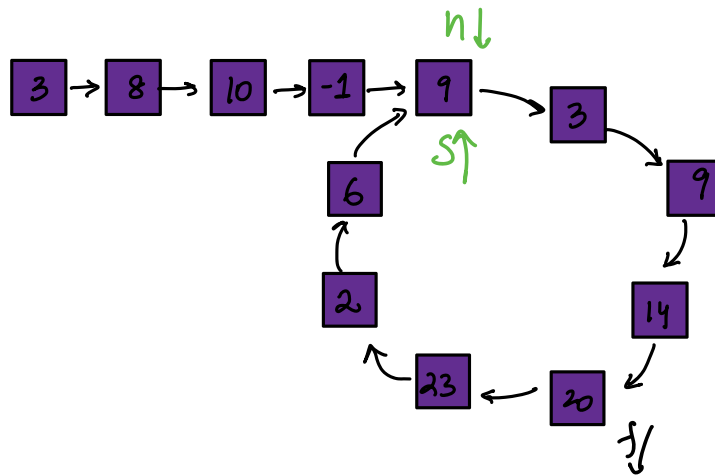


HashSet

Keep iterating on list & putting nodes in set if not already present
if present return true
if you found a NULL in LL, return false

Detect Cycle

Floyd cycle



hl.next = hl

```
s = head
f = head
while(f.next != NULL && f.next.next != NULL){
    s = s.next
    f = f.next.next
    if (s == f) return true
}
return false
```

1) Follow up question : Find start of the cycle

put h on head

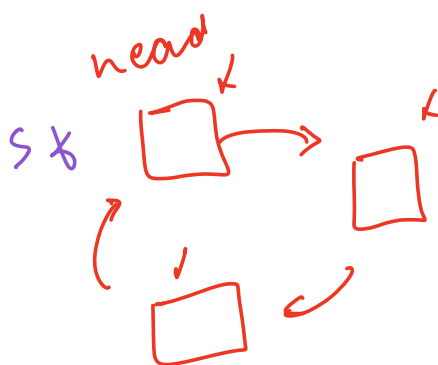
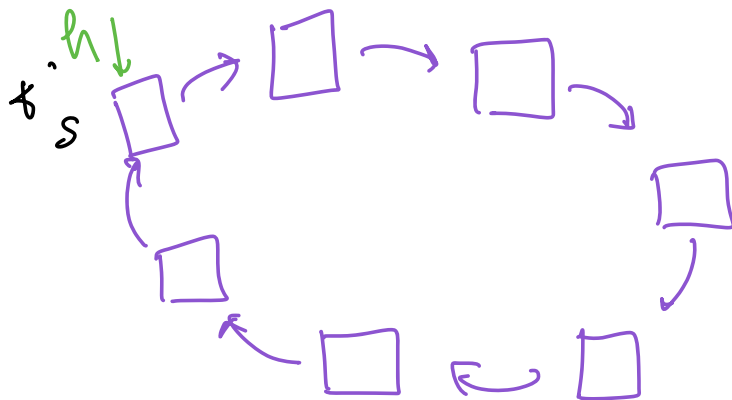
// s is present where s & f met in the
// cycle

move h & s by 1 step

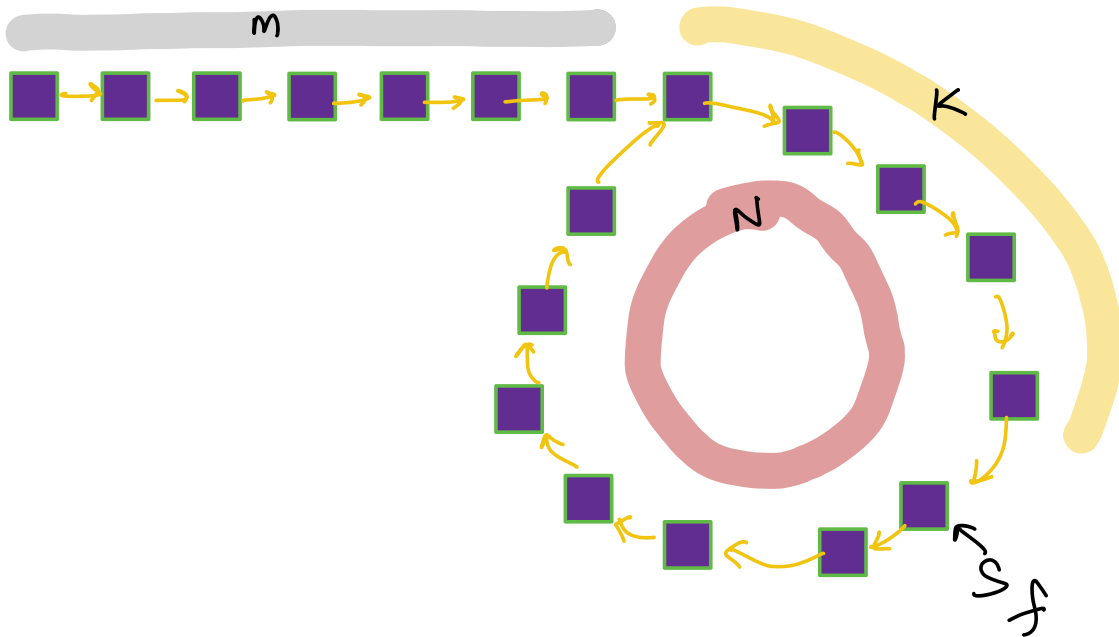
where they meet is the start of the cycle.

2) Break the cycle

start pointer from start of the
cycle & keep moving while hl.next != start
is true



Proof



$$D_s = m + k + N \times C_1$$

constants

$$D_f = m + k + N \times C_2$$

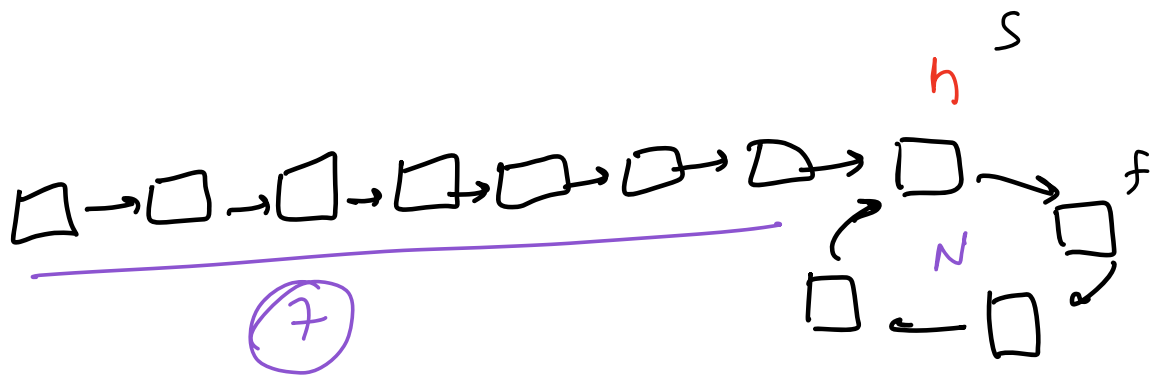
$$D_f = 2 \times D_s$$

$$m + k + N \times C_2 = 2(m + k + N \times C_1)$$

$$m + k + N \times C_2 = \underbrace{2m}_{\text{}} + \underbrace{2k}_{\text{}} + \underbrace{2N \times C_1}_{\text{}}$$

$$k + N \times C_2 - 2k - 2N \times C_1 = 2m - m$$

$$N(C_2 - 2C_1) - k = m$$



$$\underline{\underline{4 \times 2 = 8}}$$