

- 1) Queue Basics
- 2) Implementation of Queue
- 3) Problems
 - ↳ ① Nth No. using only 1, 2, 3
 - ↳ ② Find Nth perfect no.

Real world

- ATM
- Shaadi Ka Khana
- Cafeteria

Computers { Messaging Queue
Playlist
CPU Scheduling
Printer

FIFO : First In First Out

Data: 7 12 6 14 ↗ 8 3 ↗ 12 ↗

logically ↓

coffee ~~7~~ ~~12~~ ~~6~~ 14 8 3 12

Functionalities of Queue

- 1) Enqueue(x) : an ele is insert at the back of queue
- 2) dequeue() : delete an ele from front of the queue
- 3) front() : return us the ele present at front
- 4) back()/rear() : returns ele present at back
- 5) size()
- 6) isEmpty()

Eq: 7 8 9 dq() front() 10 19 back()

~~7~~ 8 9 10 19

Implementation

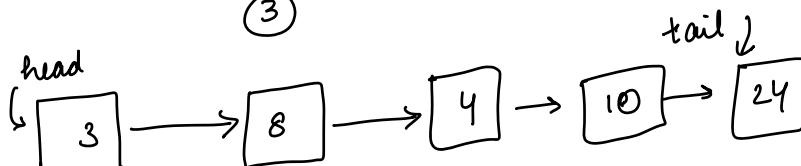
Linked List

Insert at head
Remove from tail
(Because we can't go back)
Deque will be
 $O(N)$ solution

Insert at Tail
Remove from Head

```
tail.next = newNode
tail = newNode
```

3 8 4 10 front() 24 back()
 ↓
 3



```
int front() {
    if (head == NULL) return -1;
    return head.data;
}
```

```
int back() {
    if (head == NULL) return -1;
    return tail.data;
}
```

```
void enqueue(x) {
    Node newNode = new Node();
    if (head == NULL) {
        head = newNode;
        tail = newNode;
        return;
    }
    tail.next = newNode;
    tail = newNode;
}
```

```
void dequeue() {
    if (head == NULL) return;
    temp = head;
    head = head.next;
    temp.next = NULL;
    free(temp);
}
```

* HW: If a single node is present set head & tail to NULL & free your temp

int size() : returns the size of the queue.

TC: $O(1)$ for each function

```
queue<int> q;
```

```
q.enqueue(x)
```

```
q.dequeue()
```

```
q.size()
```

```
q.isEmpty()
```

```
q.front()
```

```
q.back()
```

* Find it out what library functions are there for queue in your language.

Que: Given k .
There's a series made of only 1, 2 & 3.

↓ increasing order.

Return km no. in this series.

1 2 3 11 12 13 21 22 23 31 32 33 111 112

$$K=7 \quad : \quad 21$$
$$K=9 : 23$$
$$K = 4 : 11$$
$$K = 1 : 1$$

Brute Force

1

cnt = 0

cut=0
Iterate on nos. from [1 ..

check if all digits are 1, 2 & 3.

```
if yes cnt++
```

if yes cnt++
if (cnt == k) return mat no. & break.

1 2 3 11 12 13 21 22 23 31 32 33 111 112 113 121

1 2 3 11 12 13 21 22 23 31 32 33 111 112 113 121 122 12

Queue

~~1~~ ~~2~~ 3 11 12 13 21 22 23

How to append 1, 2 & 3 behind a no.?

① String \rightarrow str + '1' str + '2' str + '3'

② $\text{int} \times 10 + 1$
 $\text{int} \times 10 + 2$
 $\text{int} \times 10 + 3$

$K = 7$

1	2	3
---	---	---

~~1~~ ~~2~~ ~~3~~ 11 12 13 21 22 23 31 32 33 —

$$1 \times 10 + 1 = 11$$

$$1 \times 10 + 2 = 12$$

$$1 \times 10 + 3 = 13$$

$$2 \times 10 + 1 = 21$$

$$2 \times 10 + 2 = 22$$

$$2 \times 10 + 3 = 23$$

queue <int> q

q.enqueue(1)

q.enqueue(2)

q.enqueue(3)

cnt = 0

while (cnt != K-1) {

int ele = q.front()

q.dequeue()

cnt ++

q.enqueue(ele * 10 + 1)

q.enqueue(ele * 10 + 2)

q.enqueue(ele * 10 + 3)

}

return q.front()

K = 5

cnt = 0

~~1~~ ~~2~~ ~~3~~ ~~4~~ 12 13 21 22 23 31 32 33 111 112 113

cnt K-1

1 != 4

2 != 4

3 != 4

4 == 4

TC: $O(K * 3)$

$O(K * d)$

↳ only 3 digits
Hence d is constant

SC:
 $O(K * 3)$
 $O(K * d)$
 $O(K)$

$O(K)$

list <int> l;

l.add(1) }

l.add(2) }

l.add(3) }

cnt = 3

while (cnt < K) {

ele = l[P1]

l.add(ele * 10 + 1)

l.add(ele * 10 + 2)

l.add(ele * 10 + 3)

P1 ++

cnt += 3

}

return l[K-1]

TC: $O(K)$
SC: $O(K)$

K=7 → 6th index

cnt = 3
6
9

0	1	2	3	4	5	6	7	8	...
1	2	3	11	12	13	21	22	23	
P1	P1	P1							

Break: 10:50.

Q. Nth perfect No.

- ↳ even length
- ↳ palindrome
- ↳ digit (1, 2)

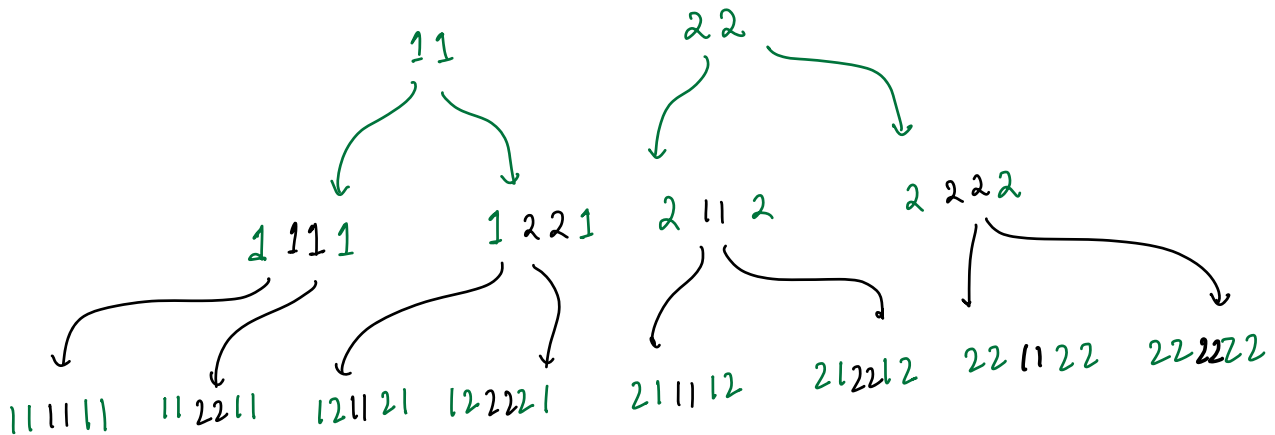
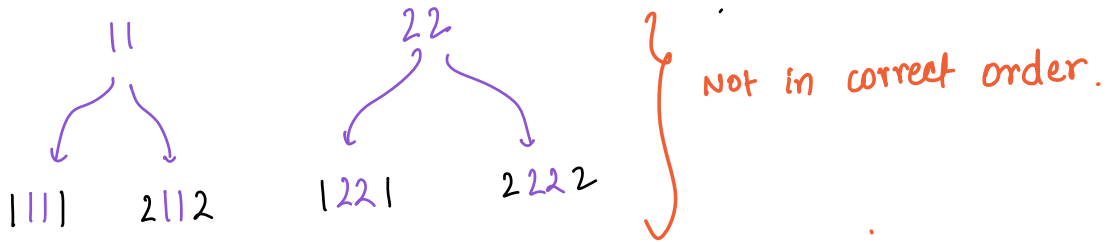
eg: 11 ✓ 22 ✓ 121 ✗ 343 ✗ 1221 ✓

11 22 1111 1221 2112 2222 111111 112211

N=5 2112

when can we say something is palindrome?

.....



```

queue<string> q;
q.enqueue("11")
q.enqueue("22")

```

11 22

How to insert in b/w ?

0 1 2 3 4 5 6 7
 1 1 2 2 2 2 1 1

str(0-3) + "11" + str(4-7)

check the
 string operation
 in your
 language

Que. Implement Queues using stacks (only)

enqueue(x)

dequeue()

↳ push(x)

↳ pop()

↳ top()

↳ size()

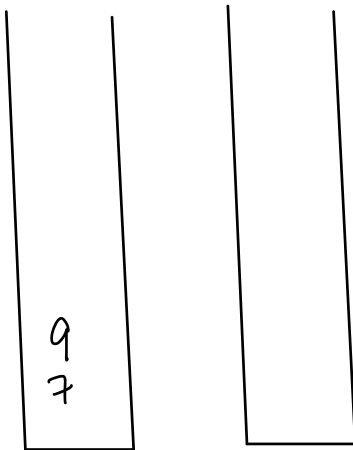
↳ isEmpty()

Imaginary

~~7~~ 9

One way

eq(5) eq(4) eq(7) eq(9) dq() dq()



TC:

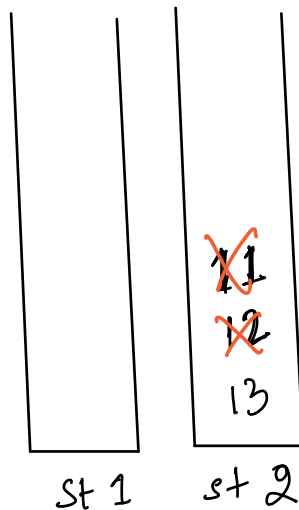
enqueue(x) $\Rightarrow O(1)$

dequeue() $\Rightarrow O(N)$

another way.

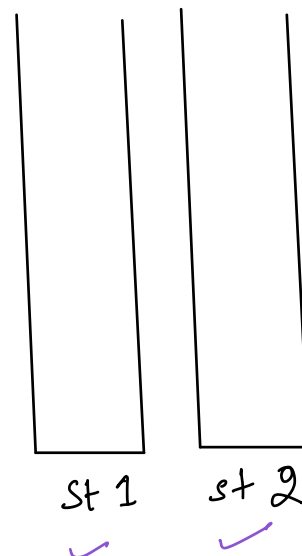
5 4 7 9 10 dq() 11 12 13 dq() dq() dq() dq()
dq() dq()

~~5~~ ~~4~~ ~~7~~ ~~9~~ ~~10~~ 11 12 13



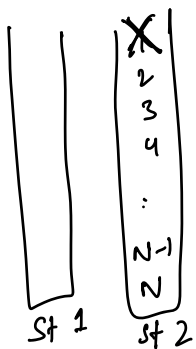
eq(7) eq(9) eq(11) eq(14) eq(17) eq(19)
eq(1) eq(7) dq() dq() dq() dq()
eq(2) dq() dq()
dq() dq() dq() dq()
dq() dq() dq() dq()

Annotations: Purple checkmarks above eq(7), eq(9), eq(11), eq(14), eq(17), eq(19). Red arrows pointing down from dq() to 17 and 19. Red arrows pointing down from dq() to 1, 7, and 7. A red bracket under the last dq().



```
void enqueue (x) {
    s1.push(x)
}
```

```
void dequeue () {
    if (s2.size == 0) {
        while (s1.size() > 0) {
            s2.push(s1.top())
            s1.pop()
        }
    }
    if (s2.size() > 0) {
        s2.pop()
    }
}
```

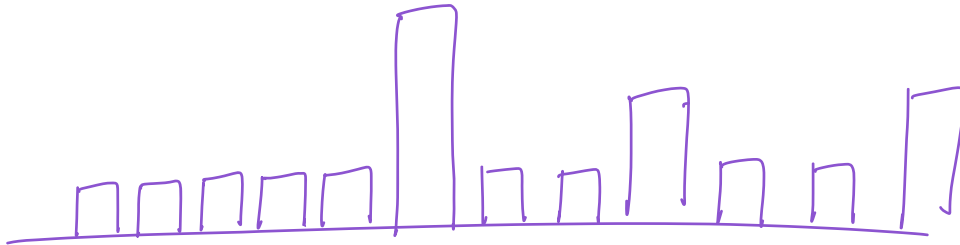


for 1 dq \rightarrow N iterations $\rightarrow 1 \times N$
 for N-1 dq \rightarrow TC: $O(1)$ $\rightarrow (N-1) \times 1$

Total iterations:

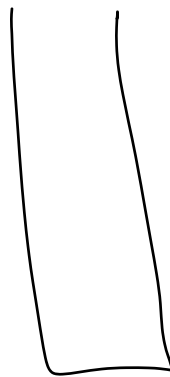
N dq $\rightarrow N + N - 1 = 2N - 1$
 $\approx O(N)$

1 dq $\rightarrow O(1)$ Amortized



x 2 3 4

back()



S1



S2

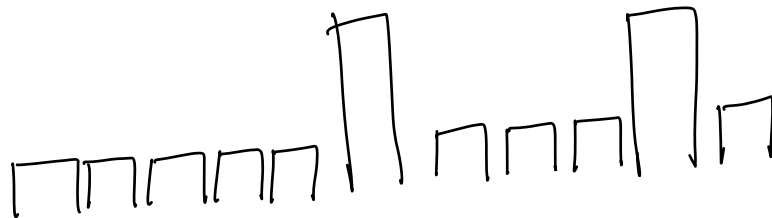
front()

HW ↑

Doubts

$O(N) + O(Q)$

queries



$Q \rightarrow O(1)$ amortized

queue<int> q

q.enqueue(1)

q.enqueue(2)

q.enqueue(3)

cnt = 3

while (cnt < K) {

int ele = q.front()

q.dequeue()

if (cnt < K) {

q.enqueue(ele * 10 + 1)

cnt++

}

if (cnt < K) {

q.enqueue(ele * 10 + 2)

cnt++

}

if (cnt < K) {

q.enqueue(ele * 10 + 3)

cnt++

}

return q.back()

K = 7

~~1~~ 2 3 11 12 13 21

cnt		K
3	<	7
4	<	7
5	<	7
6	<	7