Todays Content

- → Cycle detection in Undirected Graph
- → Minimum Spanning Tree
- → Kruskals
- → Union find Algorithm

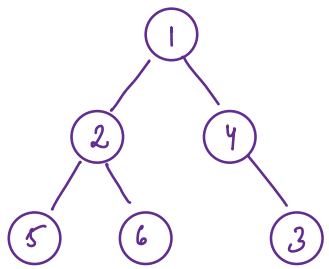
CP4

- → graphs
- → dp
- → Hashing
- → Trees
- → Bs

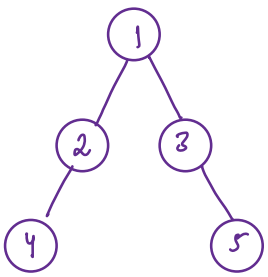
Cycle detection in Undirected Graph

→ In Tree with N Nodes, how many edges = N-1 Edges & 1 Component

Eg1:



5 Edges



4 Edges

// Given graph with N Nodes, & components?

# Nodes	# Components	# Edges it should have so that there is no cycle
N	1	N-1
N	2 → C_1 : x : x-1 Edges; C_2 : y : y-1 Edges	= x+y-2 = N-2
N	3 → C_1 : x : x-1 Edges; C_2 : y : y-1 Edges; C_3 : z : z-1 Edges	= x-1+y-1+z-1 = x+y+z-3 = N-3

Obs

N — C ⟶ # Edges we need to have so that there is no cycle = N-C

Col: Given a undirected graph with N Nodes & E Edges, check cycle, Calculate no: of components in graph = C

No Cycle if Edges E = (N-C) ; TC: O(N+E)

↳ Obs1: If Total Edges E >= N, 100% then is cycle

Pseudocode:

Step1: If E >= N {return True}

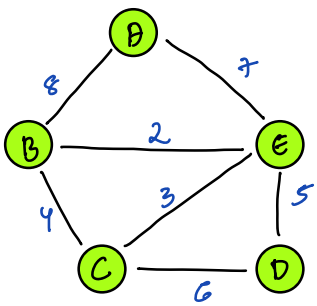
Step2: Calculate no: of components, (E == N-C)

↳ TC: O(N+E), E ≈ N TC: O(N)

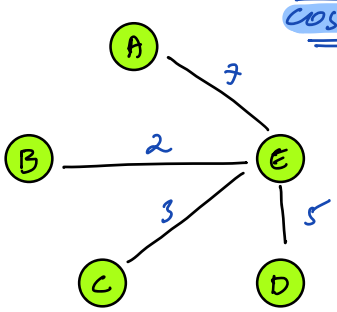
Minimum Spanning Tree

Given a **undirected weighted connected graph**, Convert into a tree with **Minimum Weight**,
{Sum of overall weights should be min},
Above tree is called Minimum Spanning Tree

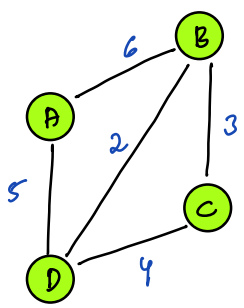
Eg1:



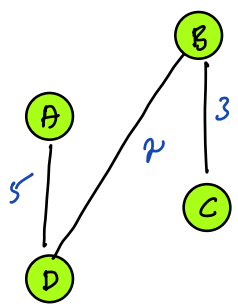
cost = 17



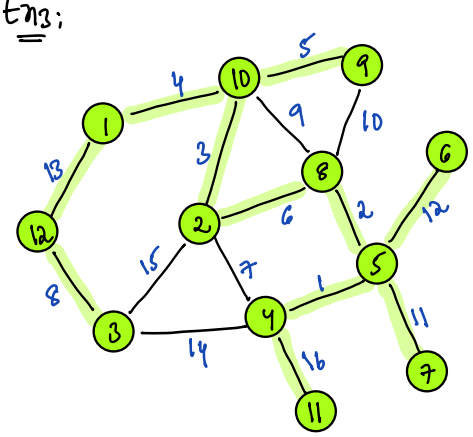
Eg2:

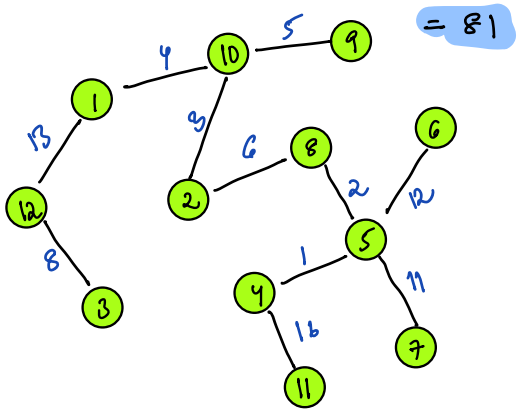


cost = 10



Eg:





= 81

Idea: {kruskals}

Step 1: Sort all Edges based on weight → TC: $E \log E$

Step 2: [Add Edges 1 by 1 to graph → TC: $E * N$

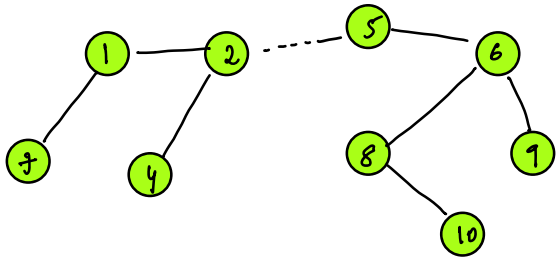
Note: If a particular edge forms a cycle, skip edge, don't add

→ [After adding we need to check for cycle] → TC: $E * [N]$ → [We need to optimize]

final TC: $[E \log E + E * N]$

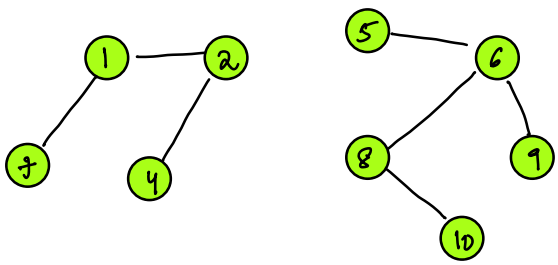
After adding 1 Edge Optimization Cycle detection

Ex1.



obs1: When 2 nodes of 2 different components are connected, it forms a single component

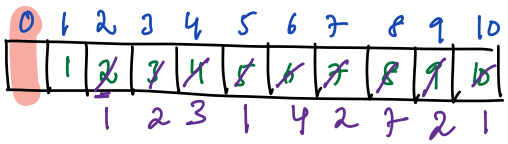
Ex2:



obs2: When 2 nodes of same components are connected, it forms cycle

// Given N = 10
W Edges

cmp[11] =



1 4-6 : cmp[6] = cmp[4]

3 3-4 : cmp[4] = cmp[3]

5 7-8 : cmp[8] = cmp[7]

5 2-7 : cmp[7] = cmp[2]

6 3-2 : cmp[3] = cmp[2]

8 6-8 : 2 : 2

10 9-6 : 9 : 2 : cmp[9] = cmp[2]

11 1-5 : 1 : 5 : cmp[5] = cmp[1]

12 7-9 : 2 : 2

14 5-10 : 1 : 10 : cmp[10] = cmp[1]

20 6-10 : 2 : 1 : cmp[2] = cmp[1]

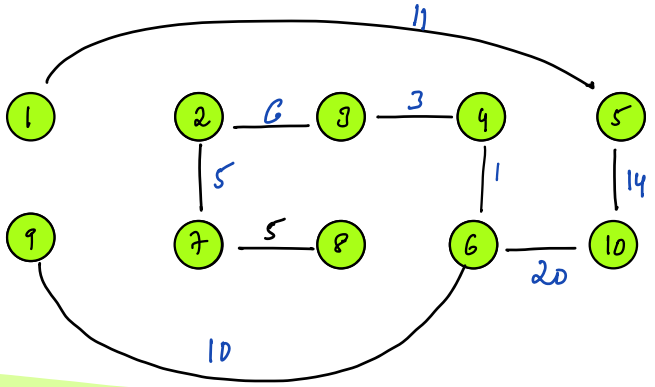
25 6-8 : 1 : 1

30 6-5 : 1 : 1

35 9-6 : 1 : 1

// WC can skip

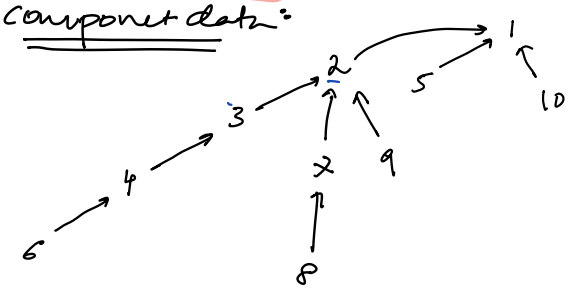
rule: assign smaller component → larger

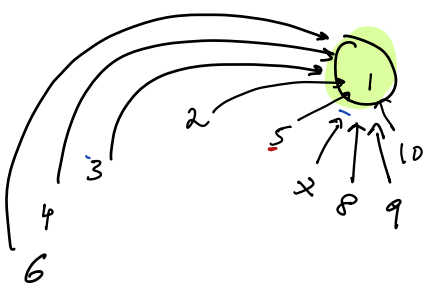


cmp[9] → cmp[2] → cmp[1]

1 1

component data:





cmp[6] → cmp[4] → cmp[3] → cmp[2] → cmp[1] : return 1

1 1 1 1 1

```
int kruskals (list< pair<int, pair<int,int>> > edges, int N){
    sort (edges) // sort edges based on weight → ElogE
    int cmp[N+1];
    int ans = 0;
    for(int i=0; i<=N; i++){ cmp[i] = i }
    for(int i=0; i<edges.size(); i++){
        pair<int, pair<int,int>> data = edges[i]
        int w = data.first
        int u = data.second.first , v = data.second.second
        // union find algo → detecting cycle in optimized manner
        //if u & v belong to different comp
        int cu = findc(u, cmp)   // find super component of
        int cv = findc(v, cmp)   // given node
                                 // TC: O(N)
        if(cu != cv){
            // assign lower component to higher component
            cmp[max(cu,cv)] = cmp[min(cu,cv)].
            ans = ans + w // Edge from u→v considered
        }
    }
    return ans;
}

int findc(int x, int cmp[]){ TC→ O(N) ——optimized——→ O(1)
    if(x == cmp[x]) return x
    cmp[x] = findc(cmp[x], cmp)   // before return update cmp[x]
                                  // find component of parent
    return cmp[x]
}
```

Overall TC → $E \log E + E * N$

↳ After changing find C

Overall TC → $E \log E + E * O(1)$

Minimum Spanning Tree → prims ismins TODO

↳ kruskals : 1) Sort all Edges

2) Add Edge by Edge, after adding an edge check, if you have a not

↳ To optimize union / find

- A single union : O(1)
- A find union : O(1)

Use Case: