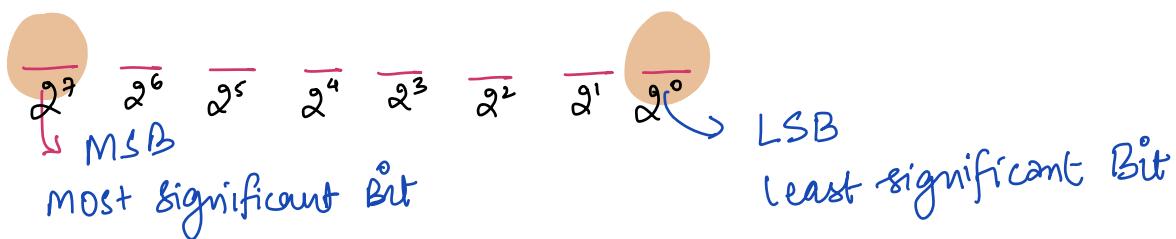


## Agenda

- 1) Negative Nos.
- 2) Ranges
- 3) Overflows
- 4) Questions



$$\begin{array}{cccc} 1 & 0 & 0 & 0 \\ \underbrace{\quad\quad\quad}_{8} \end{array} \rightarrow \begin{array}{cccc} 0 & 1 & 1 & 1 \\ \underbrace{\quad\quad\quad}_{7} \end{array}$$

$$\begin{array}{cccccccc} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ \hline 2^7 & 2^6 & 2^5 & 2^4 & 2^3 & 2^2 & 2^1 & 2^0 \end{array} \rightarrow \begin{array}{cccccccc} 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ \hline \underbrace{2^0 + 2^1 + 2^2 + \dots + 2^6}_{2^0(2^7 - 1)} \end{array}$$

$\frac{2^0(2^7 - 1)}{2-1} = 2^7 - 1$

$a=2^0$      $x=2^7$

$$\frac{1}{2^{N-1}} \frac{0}{2^{N-2}} \frac{0}{2^{N-3}} \frac{0}{2^{N-4}} \dots \frac{0}{2^2} \frac{0}{2^1} \frac{0}{2^0} > \frac{0}{2^{N-1}} \frac{1}{2^{N-2}} \frac{1}{2^{N-3}} \frac{1}{2^{N-4}} \dots \frac{1}{2^2} \frac{1}{2^1} \frac{1}{2^0}$$

$\underbrace{\hspace{10em}}$

$2^0 + 2^1 + 2^2 + \dots + 2^{N-2}$

$2^{N-1} - 1$

**MSB overpowers**

Negative No. Representation in Binary

10 : 0 0 0 0 1 0 1 0

↓  
assume  
MSB will be  
1 for ve- nos.  
& 0 " ve+ nos.

-10 : 1 0 0 0 1 0 1 0

8 Bit

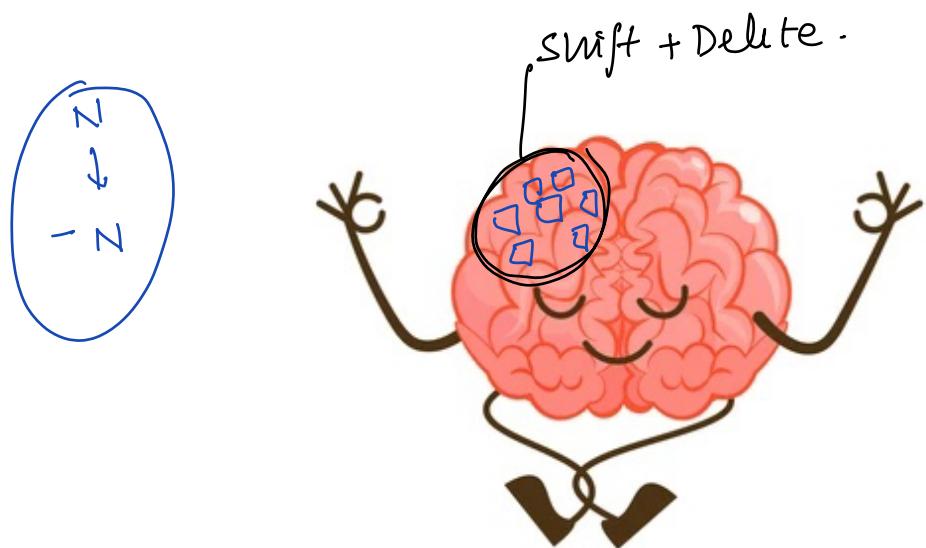
-3 : 1 0 0 0 0 0 1 1

-4 : 1 0 0 0 0 1 0 0

-7      0 0 0 0 0 1 1 1 X .

$$\begin{array}{r}
 6 : \quad \underline{\textcircled{1}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{1}} \quad \underline{\textcircled{1}} \quad \underline{\textcircled{0}} \\
 -2 : \quad \underline{\textcircled{1}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{1}} \quad \underline{\textcircled{0}} \\
 \hline
 \underline{\textcircled{1}} \quad \underline{0} \quad \underline{0} \quad \underline{0} \quad \underline{1} \quad \underline{0} \quad \underline{0} \quad \underline{0}
 \end{array}$$

$$\begin{array}{r}
 0 : \quad \underline{\textcircled{0}} \quad \times \\
 -0 : \quad \underline{\textcircled{1}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}} \quad \underline{\textcircled{0}}
 \end{array}$$



# How to get Negative Representation of N in Binary.



Steps

- ① Toggle all the Bits of  $N$
- ② Add 1

Why  $2^{\text{'s complement}}$  works?

$$\begin{array}{r}
 N : 10 \\
 \underline{0} \quad \underline{0} \quad \underline{0} \quad \underline{0} \quad \underline{1} \quad \underline{0} \quad \underline{1} \quad \underline{0} \\
 | \quad | \\
 \underline{1} \quad \underline{1} \quad \underline{1} \quad \underline{1} \quad \underline{0} \quad \underline{1} \quad \underline{0} \quad \underline{1} \\
 + 1 \\
 \hline
 \end{array}$$

$1^{\text{'s complement}}$

Add +1

$$\begin{array}{r}
 | \quad | \\
 \hline
 -2^7 \quad 2^6 \quad 2^5 \quad 2^4 \quad 2^3 \quad 2^2 \quad 2^1 \quad 2^0
 \end{array}$$

$$-2^7 \times 1 + 2^6 \times 1 + 2^5 \times 1 + 2^4 \times 1 + 0 + 2^2 \times 1 + 2^1 \times 1 + 0$$

$$-2^7 + 2^6 + 2^5 + 2^4 + 2^2 + 2^1$$

$$-10$$

int N;

unsigned int N;

4 Bits

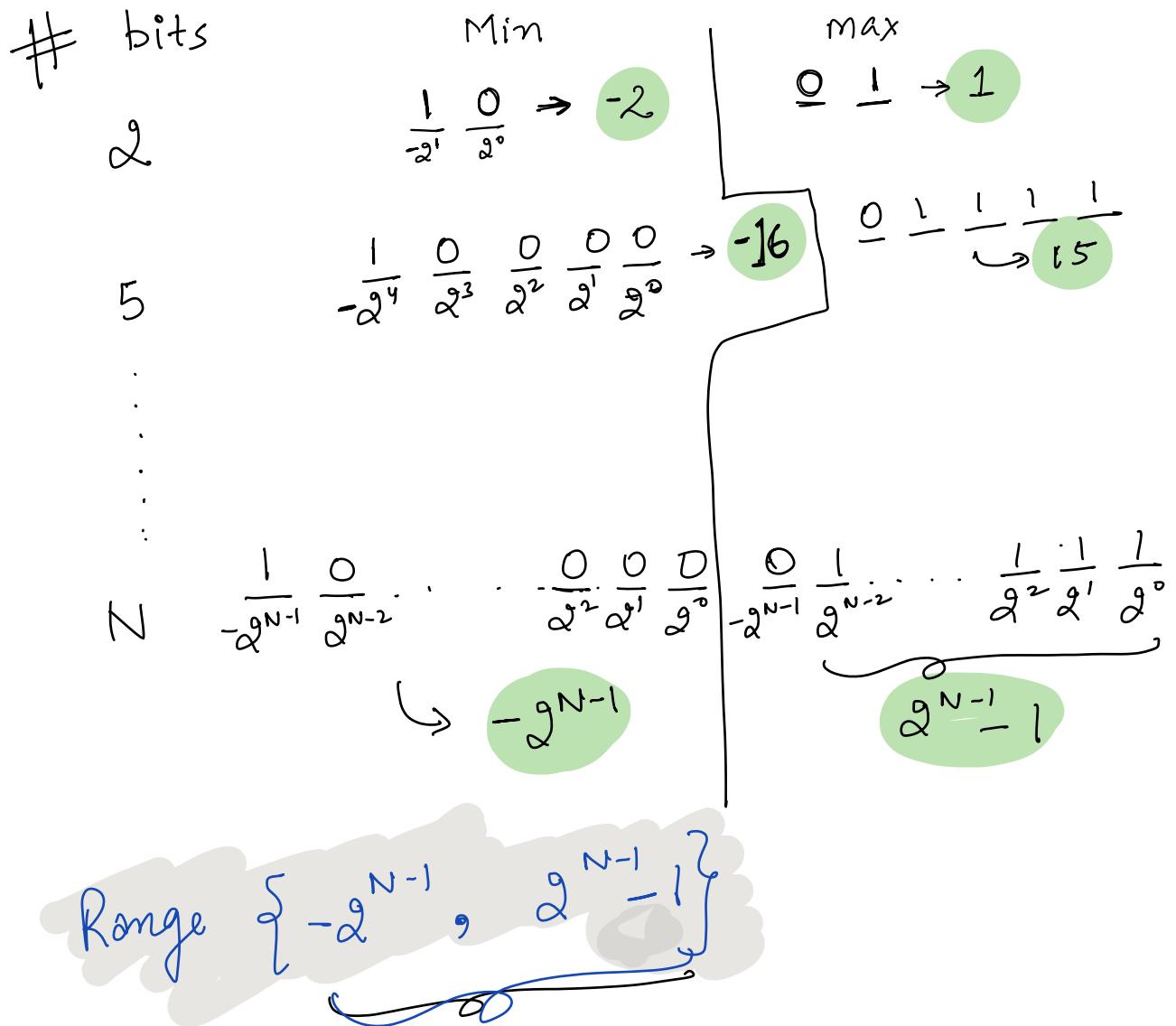
$\frac{0}{2^3} \quad \frac{1}{2^2} \quad \frac{1}{2^1} \quad \frac{1}{2^0}$

max  $\rightarrow 7$

$\frac{1}{2^3} \quad \frac{0}{2^2} \quad \frac{0}{2^1} \quad \frac{0}{2^0}$

min  $\rightarrow -8$

4 Bit  $\quad [-8 \quad 7]$



1 Byte  $\Rightarrow 8$  Bits  $\{ -2^7, 2^7 - 1 \}$

short int & Bytes  $\Rightarrow 16$  Bits  $\{ -2^{15}, 2^{15} - 1 \} \rightarrow \{ -32768, 32767 \}$

int 4 Bytes  $\Rightarrow 32$  Bits  $\{ -2^{31}, 2^{31} - 1 \}$   
 $\{ -2 \times 10^9, 2 \times 10^9 - 1 \}$

$$\text{long} \quad 8 \text{ Bytes} \Rightarrow 64 \text{ bits} , \left\{ -2^{63} \quad 2^{63}-1 \right\} \\ \left\{ -8 \times 10^{18} \quad 8 \times 10^{18} \right\}$$

## Approximation

$$2^{10} = 1024 \approx 10^3$$

$$(2^{10})^3 \approx (10^3)^3$$

$$2^{30} \approx 10^9$$

$$2^{10} \approx 10^3$$

$$(2^{10})^6 \approx (10^3)^6$$

$$2^{60} \approx 10^{18}$$

$$8 \times 2^{60} \approx 8 \times 10^{18}$$

$$2^{63} \approx 8 \times 10^{18}$$

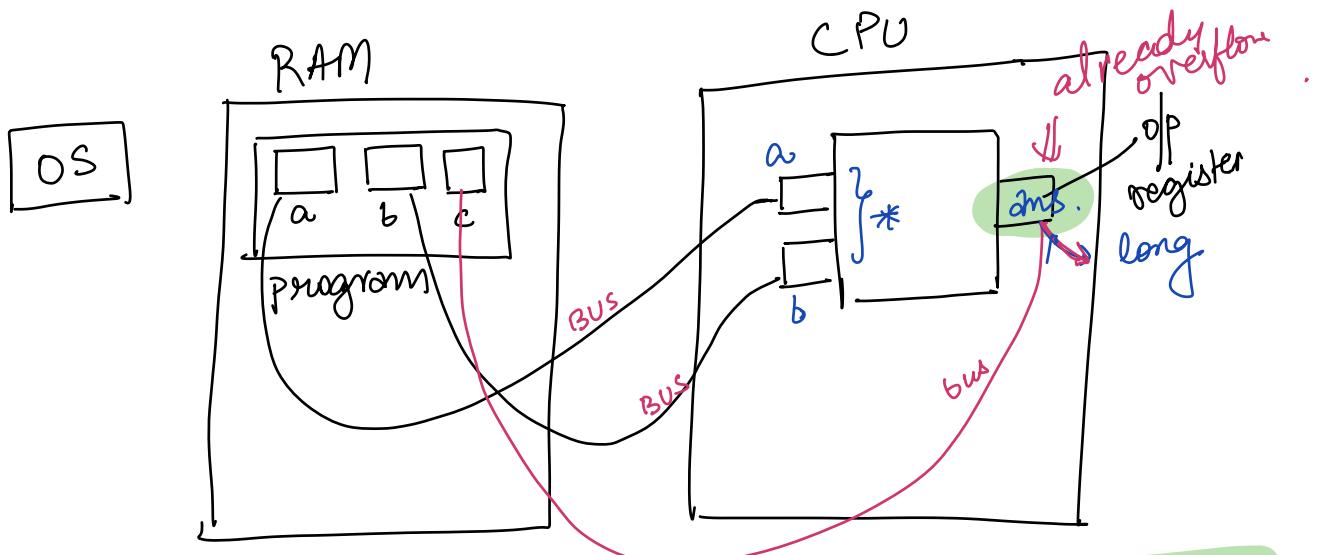
## Overflow

$$\text{int } a = 10^5, b = 10^6$$

~~int~~  $c = \overbrace{a \times b}^{10^5 \times 10^6 = 10^{11}}$  X → ①

~~long~~  $c = a \times b$  X → ②

~~long~~  $c = \underline{\text{long}}(a \times b)$  X → ③



$\text{long } c = (\text{long})(a) * b$

$\text{long } c = a * \text{long}(b)$

Ques. Given N array ele, calc sum of array elements

~~long~~  
int sum = 0  
for(i=0; i<N; i++) {  
 sum += A[i]  
}  
return sum.

### Constraints

$$1 \leq N \leq 10^5$$

$$1 \leq arr[i] \leq 10^6$$

$$1 \leq sum \leq \dots$$

$$\{10^6, 10^6, \dots, 10^6\}$$

$10^5$   
 $10^6 \times 10^5 = 10^{11}$

Break

10:40

Ques. Given N array elements, each ele is either 0 or 1.

Calc no. of subarrays whose OR of all elements = 1

OR (1)

Ex: {0, 1, 0, 0, 1} }<sup>0 1 2 3 4</sup>

[0 0] → 0

[1 1] → 1

[2 2] → 0

[0 1] → 1

[1 2] → 1

[2 3] → 0

[0 2] → 1

[1 3] → 1

[2 4] → 1

[0 3] → 1

[1 4] → 1

[3 3] → 0

[0 4] → 1

ans: 11

[3 4] → 1

[4 4] → 1

Brute Force

TC → O(N<sup>2</sup>)

SC → O(1)

Try this out yourself.

Observation

If in a subarray even if there is a single 1, then its OR = 1 How to check?

If we find out in how many subarray there are all 0s., then we are done?

ans : #totals - # subarrays having all 0s.

$$\frac{n(n+1)}{2}$$

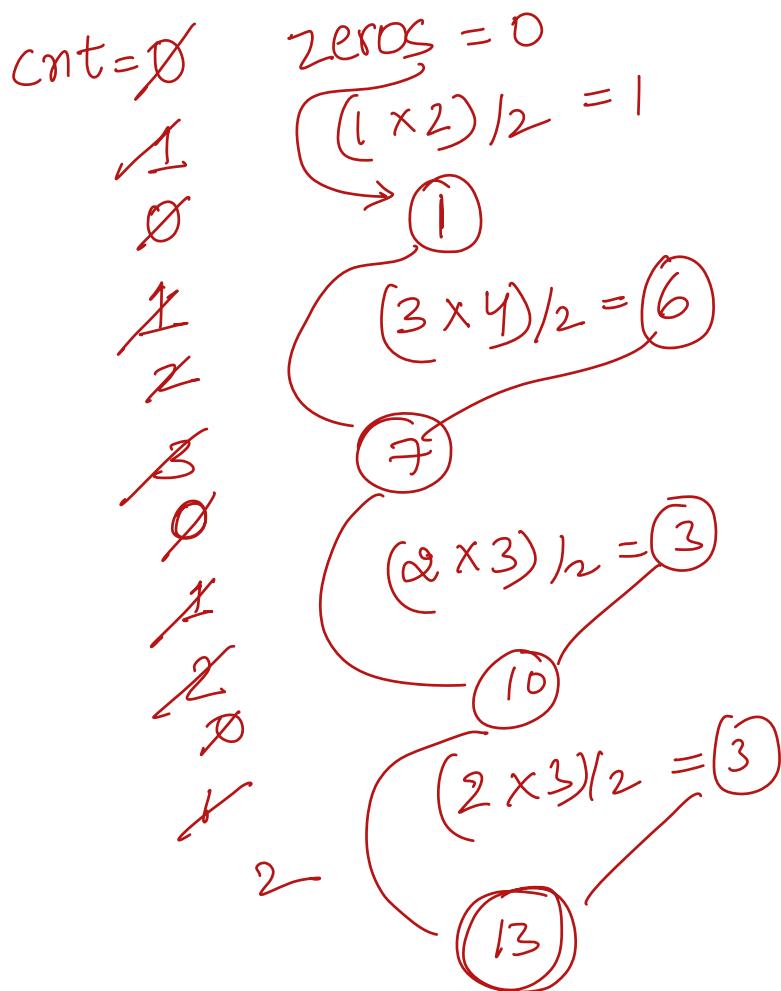
$$\left\{ \begin{array}{ccccccccccccc} \downarrow & \downarrow \\ 0, & 0, & 0, & 1, & 0, & 0, & 1, & 0, & 0, & 1, & 0, & 1, & 0, & 0 \\ \text{cnt=3} & \text{cnt=0} & & \underbrace{1 \quad 2}_{0} & & \text{cnt} & \underbrace{1 \quad 2}_{0} & & \text{cnt} & \underbrace{1 \quad 2}_{0} & & \text{cnt} & \underbrace{1 \quad 2}_{0} \\ \frac{3(4)}{2} = 6 & & & \frac{2(3)}{2} = 3 & & & & & & \frac{2(3)}{2} = 3 & & & \frac{2(3)}{2} = 3 \end{array} \right.$$

15

$$\text{ans} = \frac{13(14)}{2} - 15 = 76$$

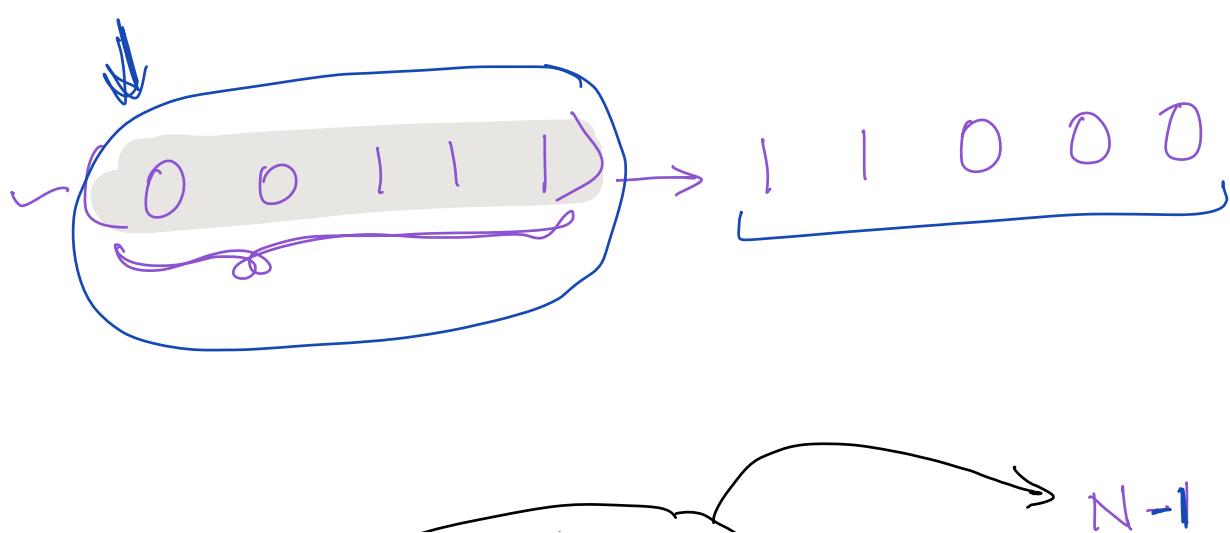
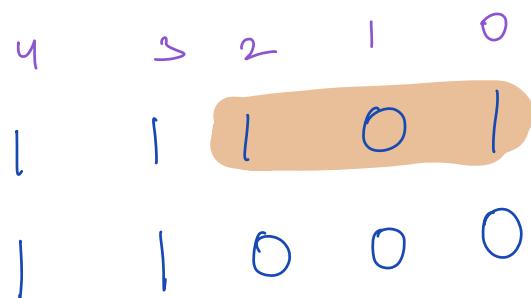
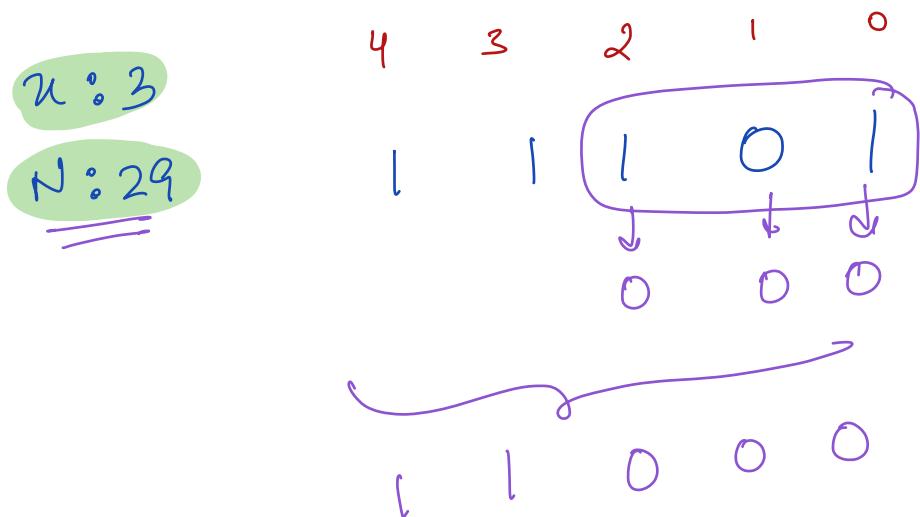
. 1 1 1 1

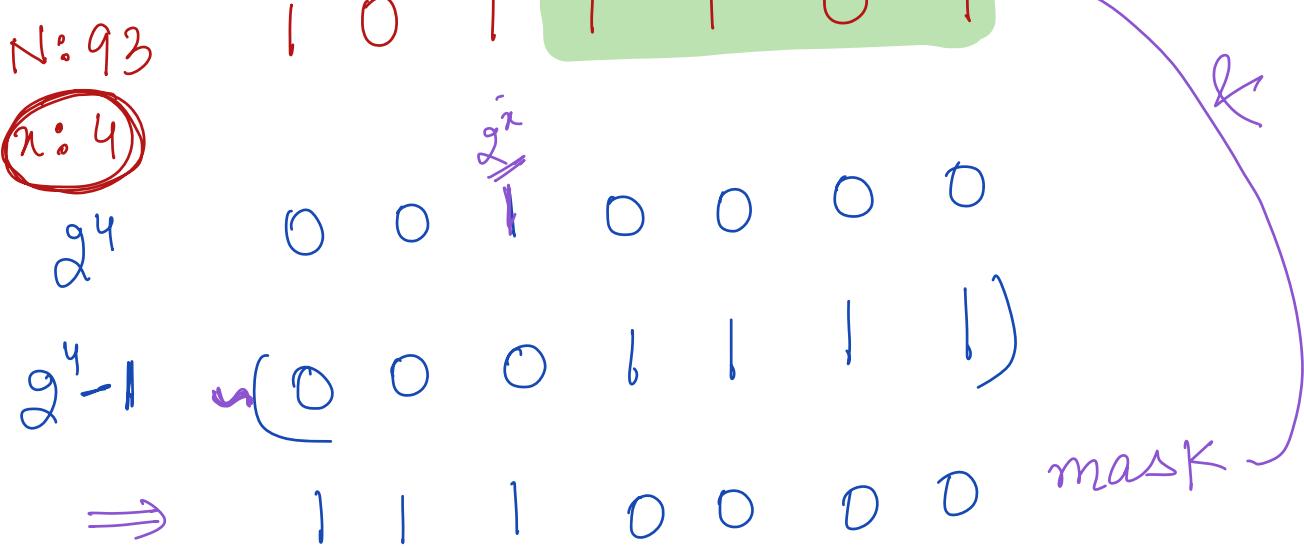
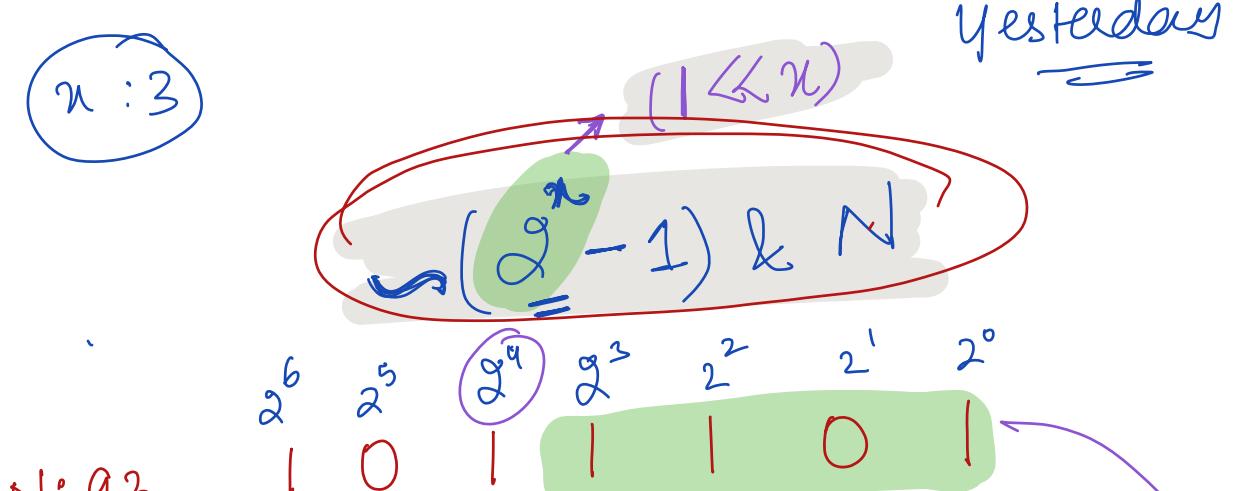
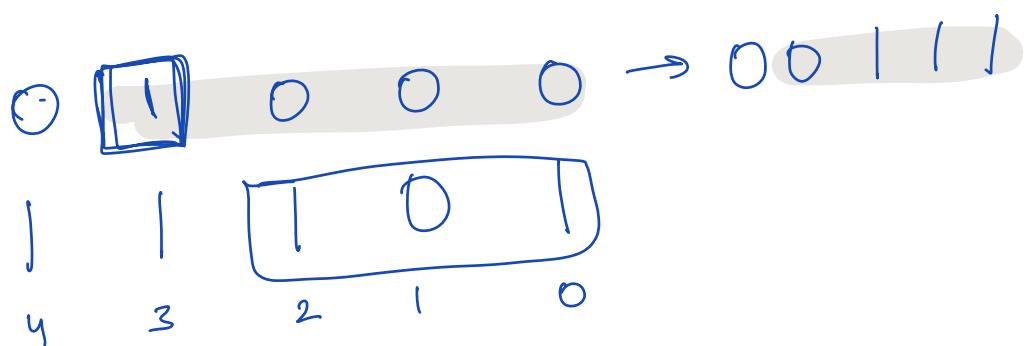
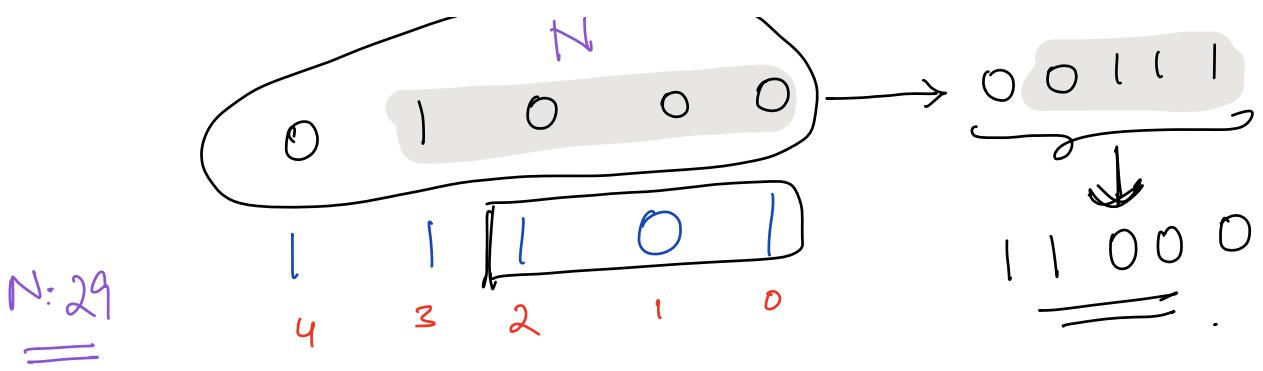
$\downarrow$     $\downarrow$     $\downarrow$     $\downarrow$     $\downarrow$     $\downarrow$     $\downarrow$     $\downarrow$     $\downarrow$   
 0   1   0   0   0   1   0   0   1   0   0



$$\begin{aligned}
 \text{Ans: } & \frac{12 \times 13}{2} - 13 \\
 & 78 - 13 = 65
 \end{aligned}$$

Unset  $x$  continuous bits in a no.  $N$





$$\begin{array}{r}
 \text{1 0 1} \quad \text{1 1 0 1} \\
 \times \text{1 1 1 0 0 0 0} \\
 \hline
 \text{1 0 1 0 0 0 0}
 \end{array}$$

~~TG~~ O(1)  
~~SC~~ O(1)

Nº 0 1 0 1 0 0 → N-1

0 1 0 0 1 1

$$\begin{array}{r}
 \text{1 0 1 1 1 0 1} \\
 \times \text{1 1 1 0 0 0 0} \\
 \hline
 \text{1 0 1 0 0 0 0}
 \end{array}$$

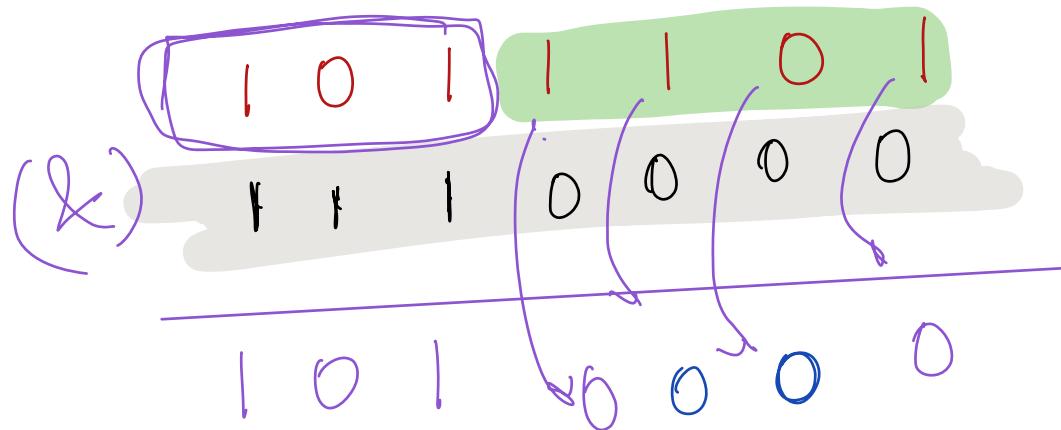
1 0 1 1 1 0 1

0 0 1 0 0 0 0 →

0 0 0 1 1 1 1  
||

ψ

F F I 0 0 0 0



N: 1 0 4 3 2 1 0  
0 0 | 0 0 0

$x = 3$

| | 0 1 1

(LLR2)

