

→ Students Names

{
 → Marks : Total Marks
 → Rollno : —
 → Name : — }

→ Say N Students

	Index:	0 th	1 st	2 nd	3 rd
→ Marks [N]	—	97	65	35	75
→ Roll [N]	—	14	24	31	33
→ Name [N]	—	Rashika	Sardutt	Viraj	Piyali

→ 3 ways

→ Syntax of Class Creation

→ user defined datatype }

class student {

 int rollno ;

 String name :

 int marks ;

student(int ro, int ma, String na) {

 rollno = ro marks = ma

 name = na

// construct

① → construct name should

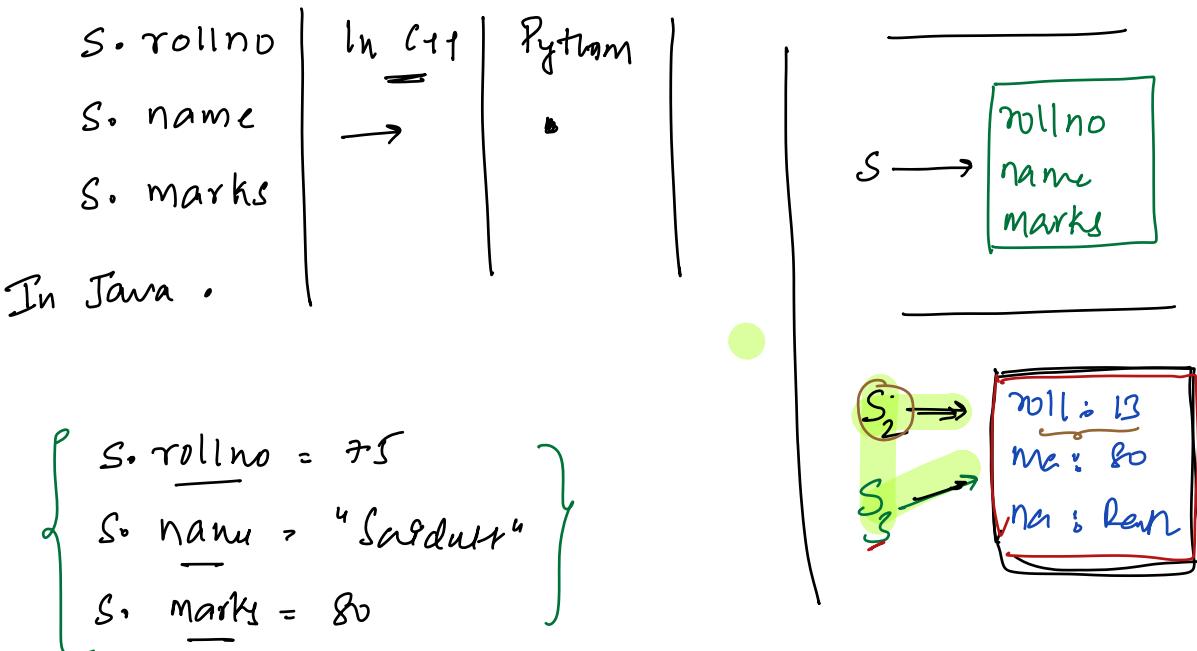
be exactly same as

class name

② Constructor is exactly like a function but no return type.

main() {

Student s = new Student()



Student (S₂) = new Student(13, 90, "Ravi")

print(S₂.marks) // Whether data initialized

Student S₃;

S₃ = S₂

print(S₃.marks)

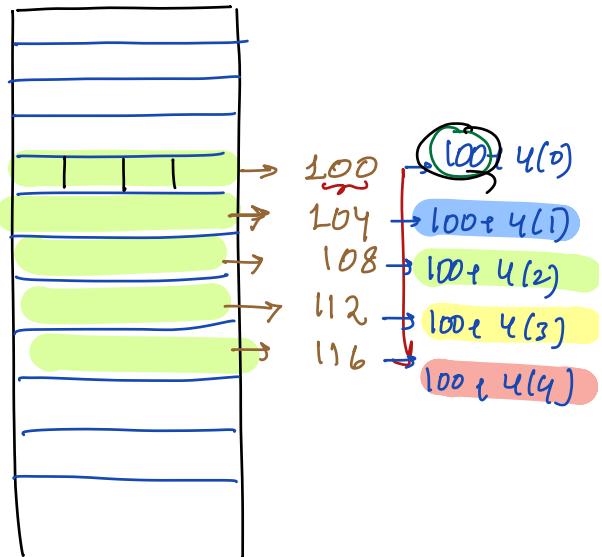
= Most commonly data structure

→ Arrays ?

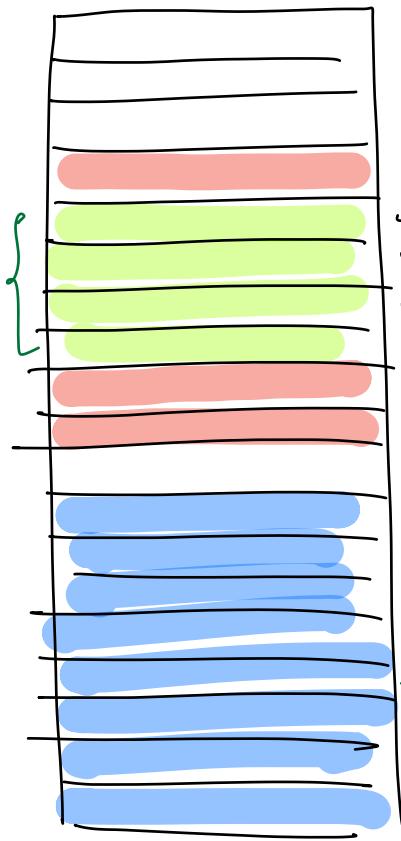
↳ i^{th} index $\rightarrow ar[i] \Rightarrow O(1)$?

→ $\text{int } ar[5] \Rightarrow$

It will allocate continuous
Blocks of memory



→ Insertion in a List \Rightarrow $O(n)$
↳ average



→ Dynamic arrays

→ q of arrays

int a[4] → 4 Elements filled

int a₁[8] =

→ copy a[4] → a₁[8] = O(N)

→ Insert new element in a₁[8]

int a[N]

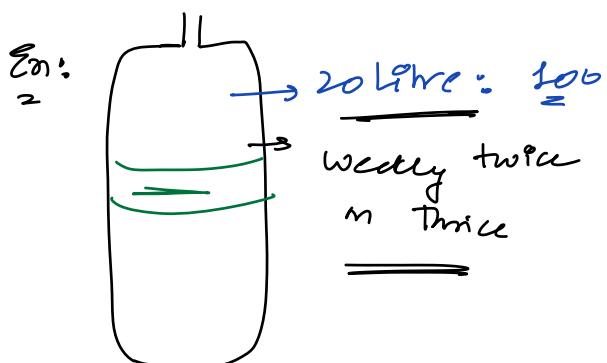
→ All N partitions

int a₁[2N]

N copy → a[] → a₁[] : O(N)

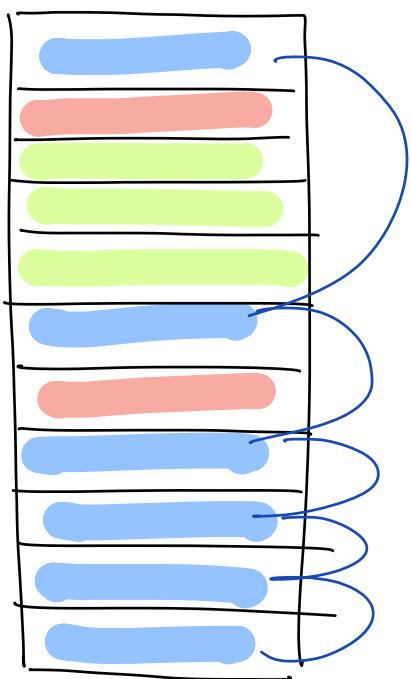
add new element a₁[]

Inbuilt
Implementation



1 time quarter

Aquaproof: 10,000 ?



Dynamic array

`int arr[3]`

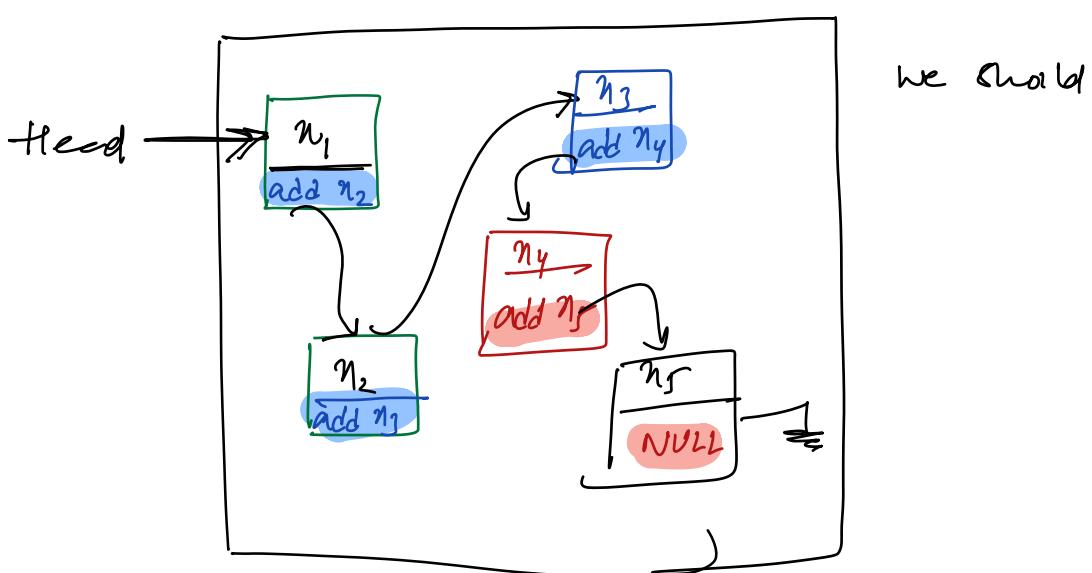
3 blocks

`Put b[6]` = 6 continuous blocks

⇒ You want to properly use memory,
if you comprise **old acc time**?

// of linked list

→ We can use random locations



```
class Node { }
```

int data

Node next // we will address of next node

Node (int n) { \Rightarrow constructor }

data = n

next = NULL

↳ keyword

```
main() {
```

Node head = new Node(5)

head.next = new Node(6)

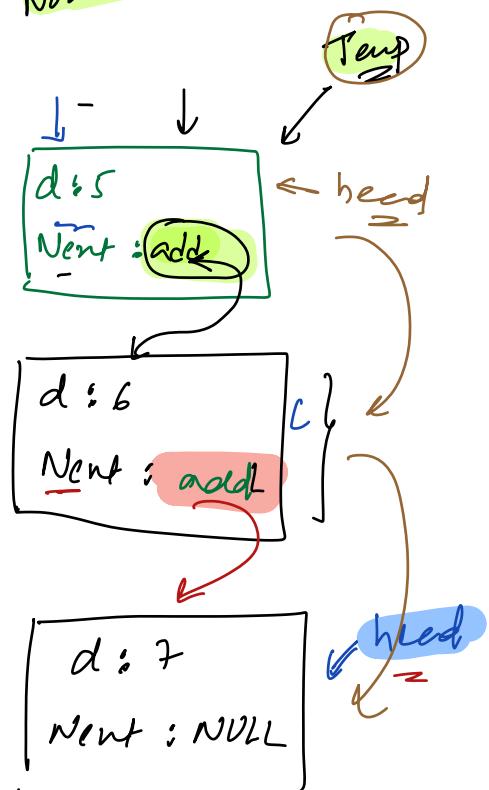
head = head.next

head.next = new Node(7)

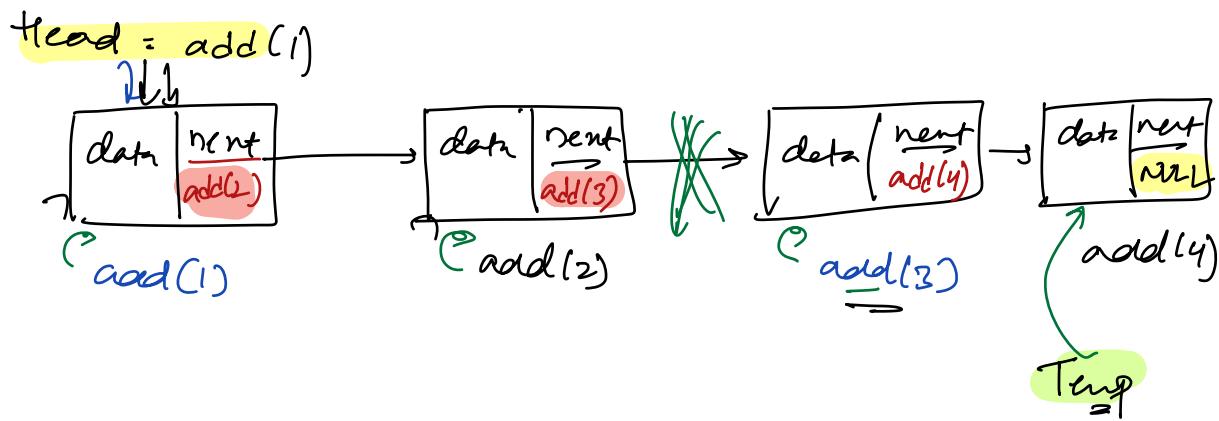
head = head.next

Note: Don't look
head node address

Node temp = head



1 way



Temp = add(1)

Temp = Temp.next] add(2)

Temp = add(2)

Temp = Temp.next] add(3)

Temp = add(3)

Temp = Temp.next] add(4)

Temp = add(4)

Temp = Temp.next] NULL

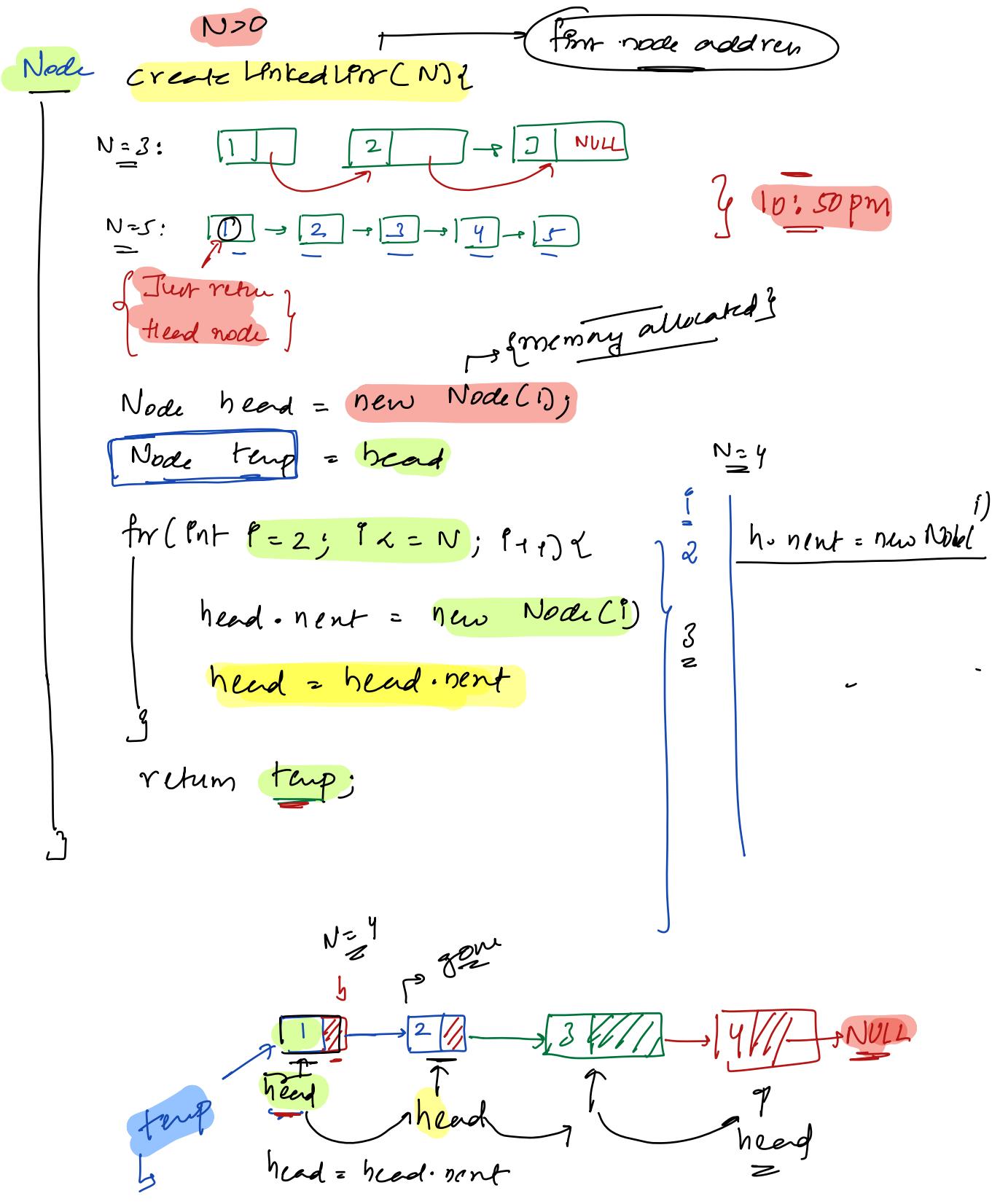
Temp = NULL

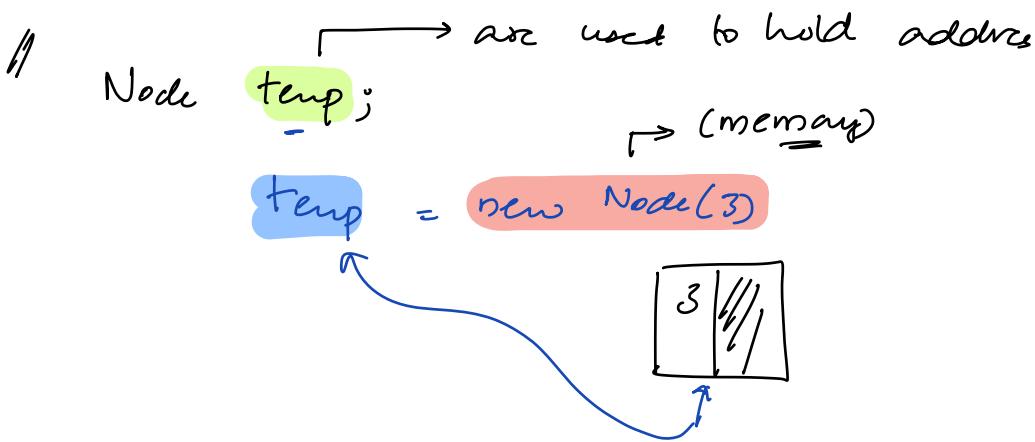
End game
no more nodes

```

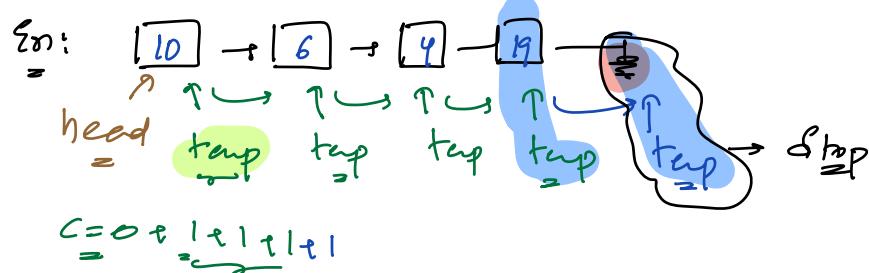
class Node {
    int data
    Node next // we hold address of next node
}
Node(int n) { // constructor
    data = n
    next = NULL // keyword
}
  
```

data can be anything
float/char/String/arr[]/





// print_sqz (Node head)



Node temp = head // good practice

while (temp != NULL) { if temp == NULL: break

temp = temp.next // if update temp }

C++

}

return c;

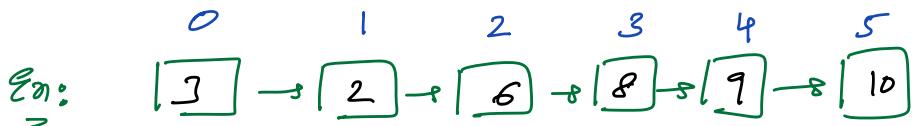
temp = NULL → { Nothing }

$\left\{ \begin{array}{l} \text{temp.data} \\ \text{temp.next} \end{array} \right\} \Rightarrow \begin{array}{l} \text{NULL (Pointer Encapsulation)} \\ \underline{\text{Error}} \end{array}$

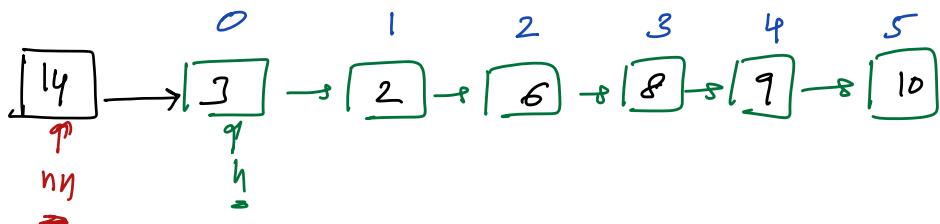
→ Edge Case //

↑ of board analogy
at that node

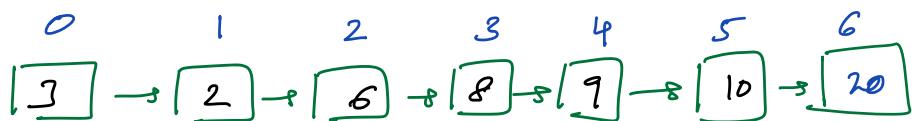
— Insert at k (Node head, Put n, int k)



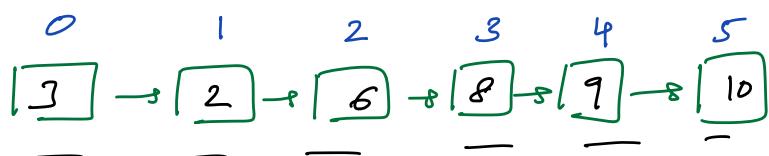
Case : $k=0, n=14$



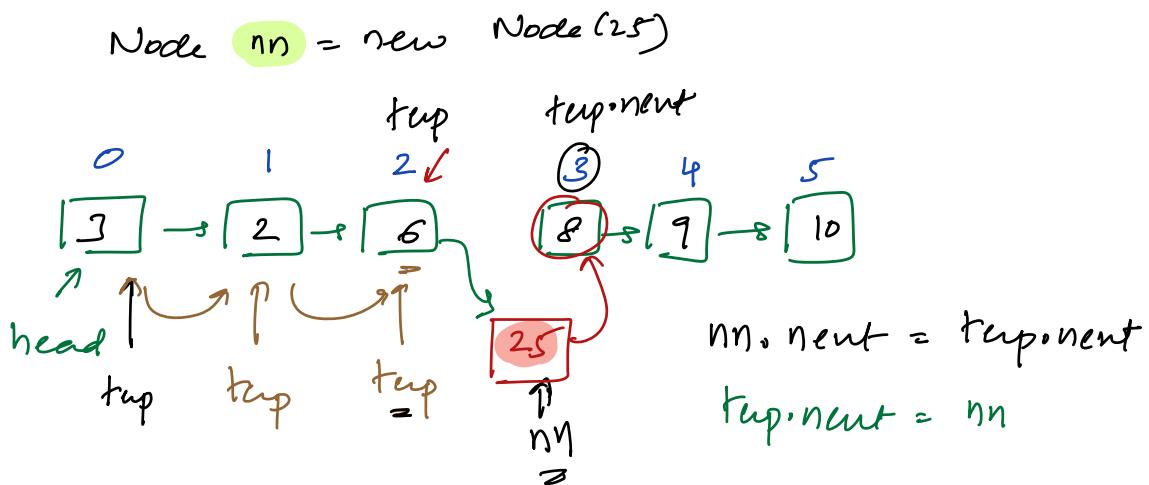
Case 2 : $k=6, n=20$



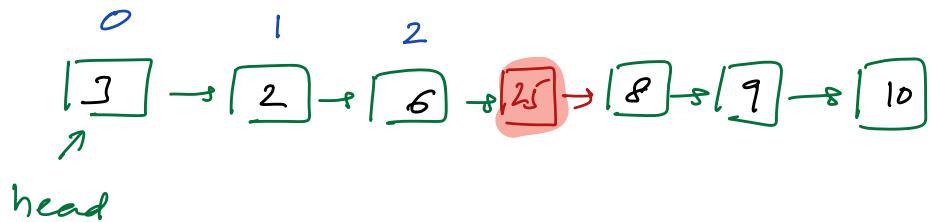
Case 3 : $k=7, n=14$: If $k > (\text{len of linked list})$ → cannot



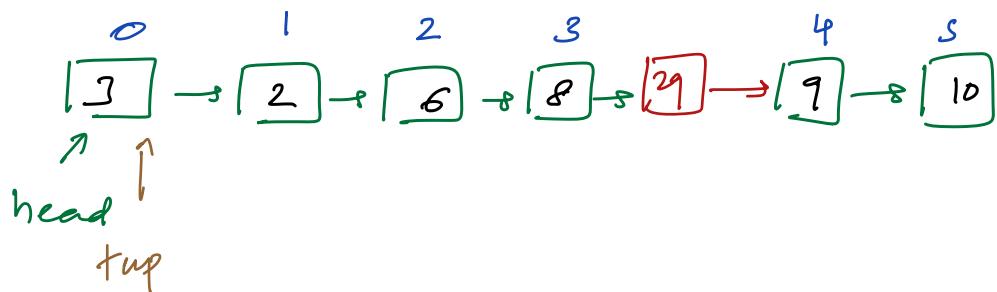
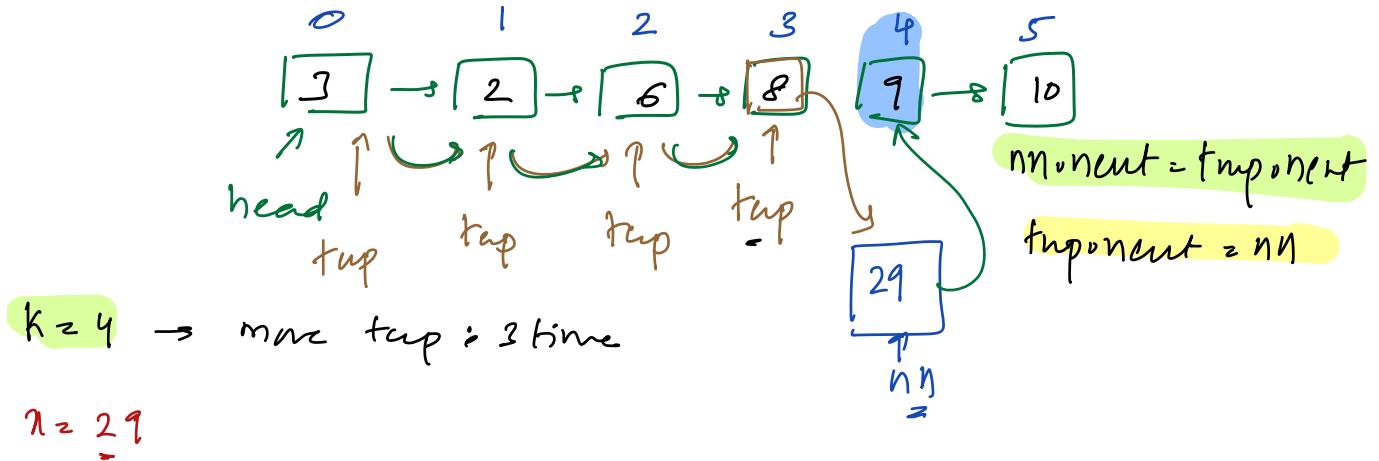
Case 4: $k = 3$, $n = 25$



Node: Take temp till prev address



$nn = \text{new Node}(29)$



InsertKpos(Node head, int clc, int k) {

If (k > size(head)) { return head; }

Node nn = new Node(clc);

If (k == 0) {

nn.next = h; h = nn; return h;

Node temp = head;

for (int i = 1; i <= (k - 1); i++) {

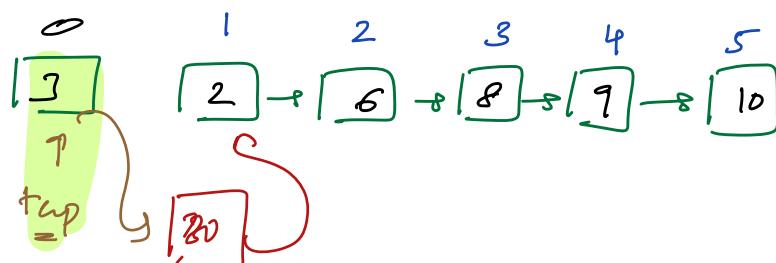
temp = temp.next;

nn.next = temp.next;

temp.next = nn;

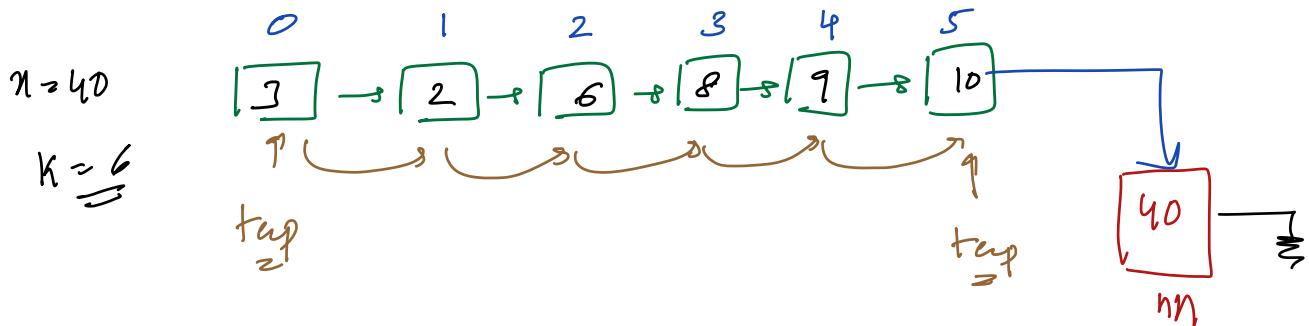
return head;

}



n = 20, k = 1

Iterate (k-1)



Moment = Response

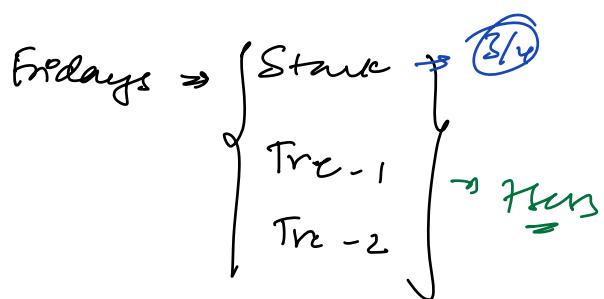
Nonzero = NUL

Response = nn

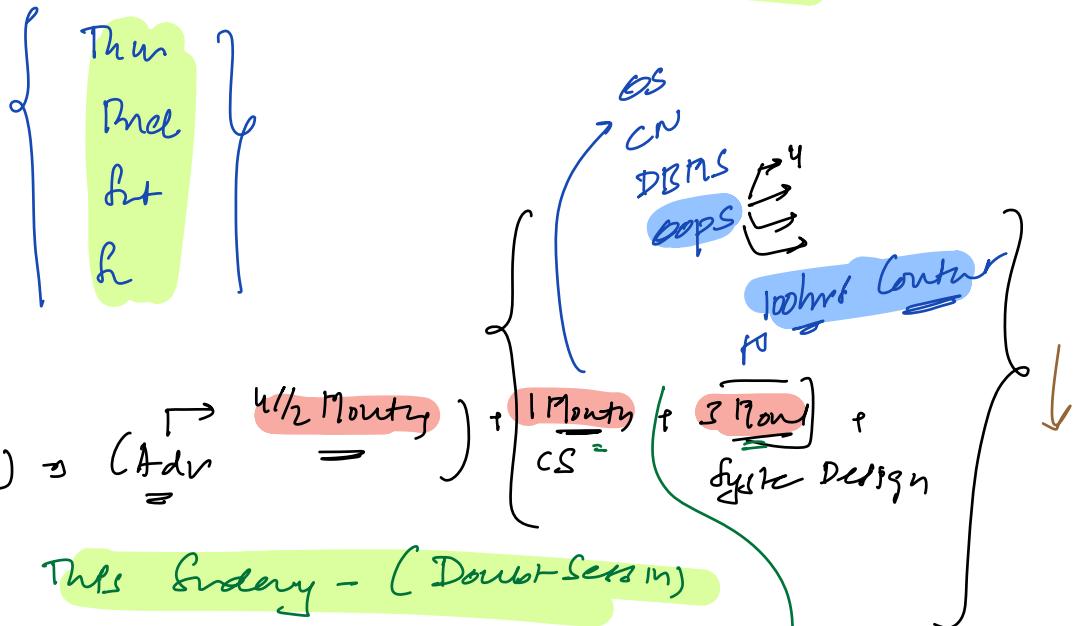
3 Delete k^M pos \Rightarrow CTODD

Advanced Content \Rightarrow 4 Scans

- 1) \rightarrow
- 2) insert
- 3) del



⇒ Now Wednesday → (Last Intermediate Session)



⇒ (Reverse Topics)

- 1) 2 weeks Plan
- 2) Intensive Plan