

Today's Content :

↳ Back Tracking ↳ ↳ ↳

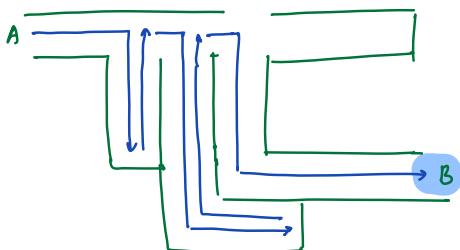
- 1) All numbers
- 2) Subset sum
- 3) General way

Back Tracking :

↳ Generating all possibilities to get correct ans

Maze:

↳ Harry Potter Goblet of fire:



Try out all combinations, if a particular is not working re-trace & pick a new combination.

Note: All Back tracking codes are written by Recursion

Q8) Given N digits print all N digit numbers formed only by 1 & 2
In increasing order of numbers

$$N=1: \begin{array}{c} 1 \\ 2 \end{array}$$

$$N=3: \begin{array}{c} 1 \ 1 \ 1 \\ 1 \ 1 \ 2 \\ 1 \ 2 \ 1 \\ 1 \ 2 \ 2 \\ 2 \ 1 \ 1 \\ 2 \ 1 \ 2 \\ 2 \ 2 \ 1 \\ 2 \ 2 \ 2 \end{array}$$

Idea: If we replace 1 → 0

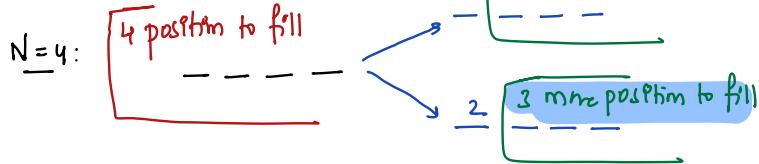
2 → 1, Can be solved

using bit-marking

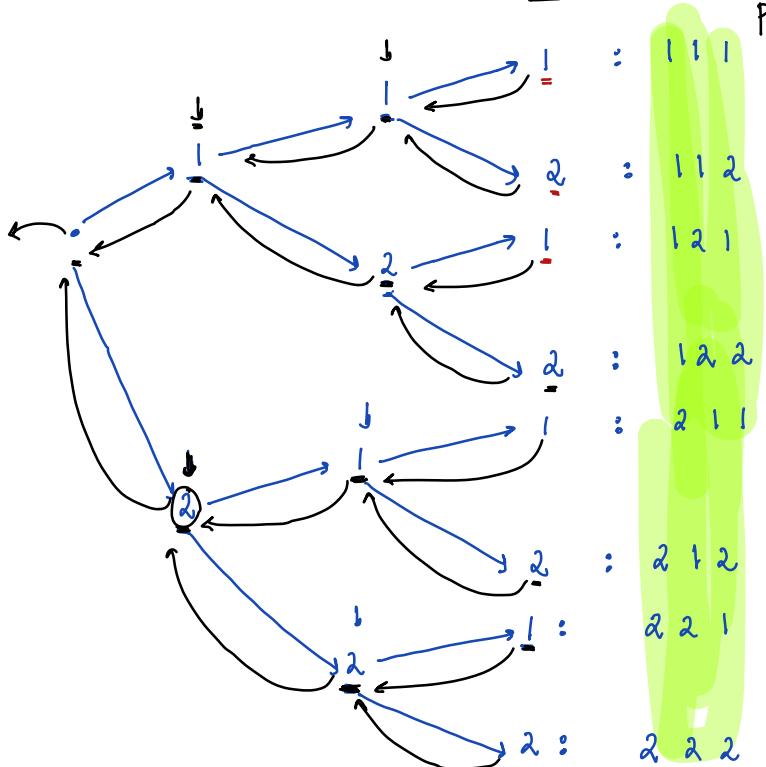
$$T C: O(2^n * n)$$

$$N=2: \begin{array}{c} 1 \ 1 \\ 1 \ 2 \\ 2 \ 1 \\ 2 \ 2 \end{array} \left[\begin{array}{c} 11 \times 12 \\ 12 \times 21 \\ 21 \times 22 \end{array} \right]$$

$$\begin{array}{c} 1 \ 2 \ 1 \\ 1 \ 2 \ 2 \\ 2 \ 1 \ 1 \\ 2 \ 1 \ 2 \\ 2 \ 2 \ 1 \\ 2 \ 2 \ 2 \end{array}$$



$$N=3: \begin{array}{c} 0 \\ - \\ \downarrow \\ 1 \\ - \\ \downarrow \\ 2 \\ - \\ \downarrow \\ 3 \end{array} \rightarrow \text{we filled the number} \\ \text{print it & back track}$$



→ funcC: →

parameters: N, ar[N], i → at what index we are currently
function calls: 2 filling data

return type: void

pass by reference

void printall (int ar[], int N, int i) {

if (i == N) { // We have a number hence print & return
 print (ar)
 return;

pass by refena:

All function calls are going to
same array

TC: $[2^N * N]$

Sc: $[N + N]$

ar[] Call stack size

ar[i] = 1; } made

printall (ar, N, i+1);

ar[i] = 2; }

printall (ar, N, i+1);

main () {

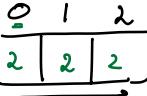
 int N;

 Read N

 int ar[N];

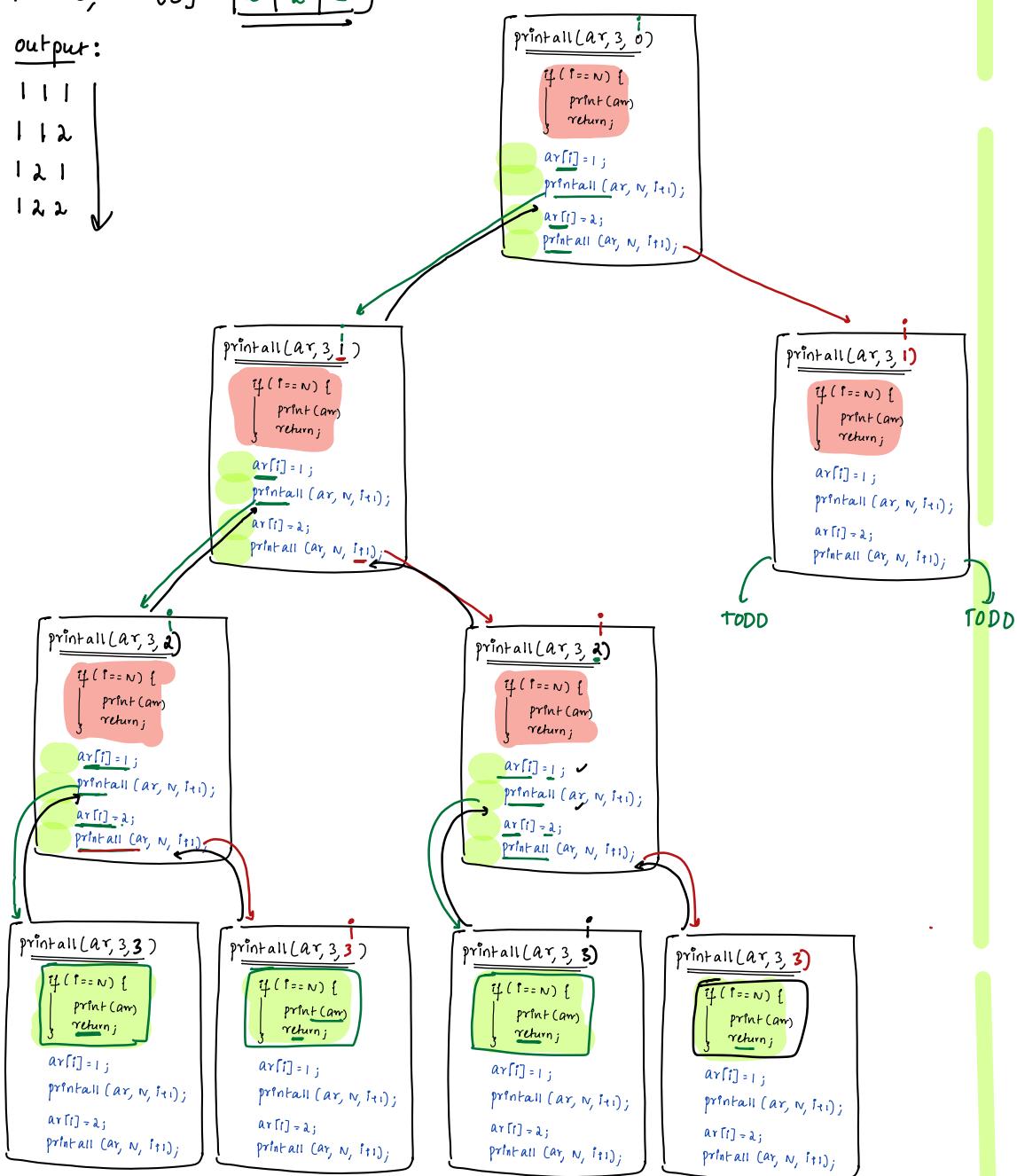
 printall (ar, N, 0)

}

// N=3, ar[3] = 

output:

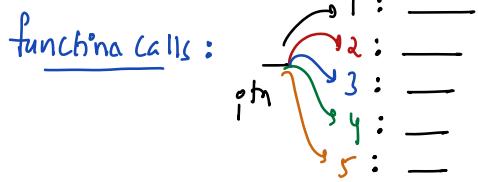
1 1 1
1 1 2
1 2 1
1 2 2



// Given N print all N digit numbers formed by digits + {1 2 3 4 5}
in increasing order

func:

parameters: ar[], N, i



void printall (int ar[], int N, int i) {

if (i == N) { // We have a number hence print & return
 print(ar),
 return;

ar[i] = 1;

printall (ar, N, i+1);

ar[i] = 2;

printall (ar, N, i+1);

ar[i] = 3;

printall (ar, N, i+1);

ar[i] = 4;

printall (ar, N, i+1);

ar[i] = 5;

printall (ar, N, i+1);

for (int v=1; v<=5; v++) {
 ar[i] = v;
 printall (ar, N, i+1);

$$TC: 5^N \times \{N\}$$

$$SC: N + N$$

arr ↘
 Call stack size

10: Q3 → 10: 80pm

Q8) Given N array elements count no:of subsets wthn sum==k

$$\begin{bmatrix} \text{arr} = & \{1, 2, 3, 6, 1, 5\} \\ & 0 \\ & -k \end{bmatrix}$$

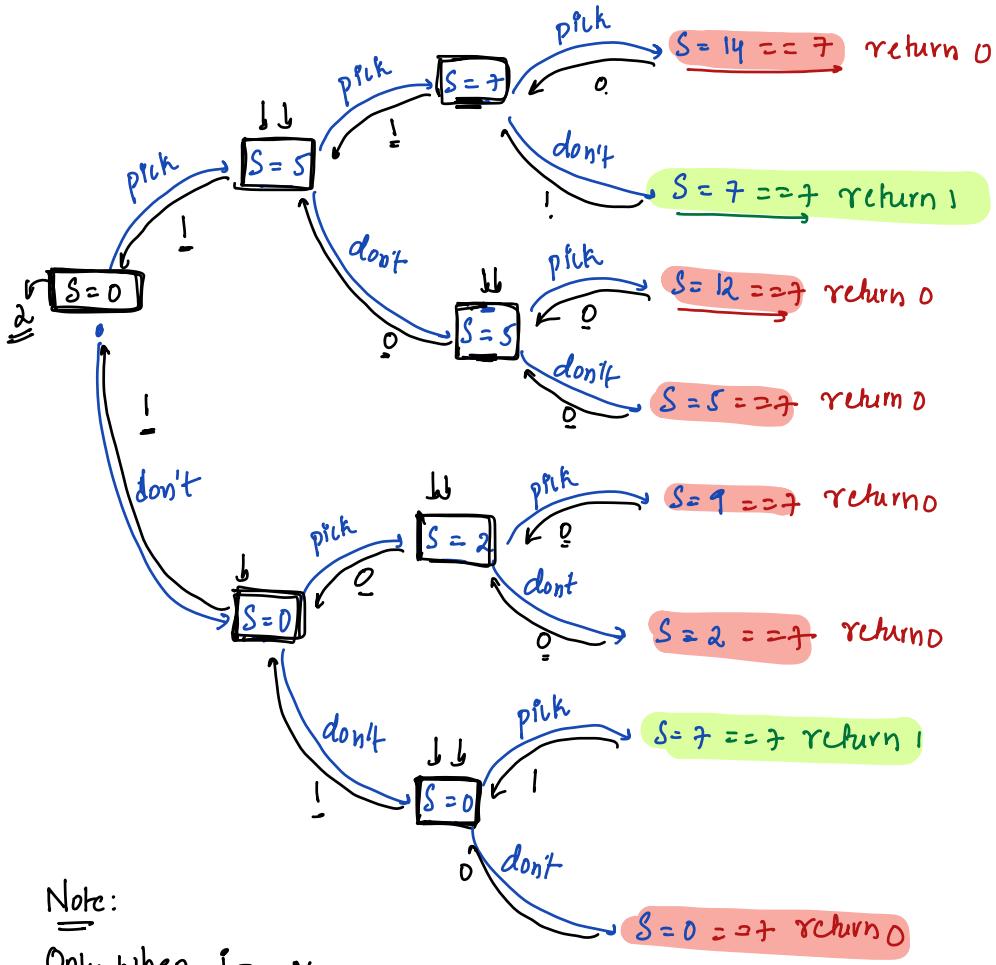
Ex: [1 0 2 3 6 1 5]

k=8 : {2 6} {1 7} {5 2 1} \Rightarrow ans=3

Ideal: Generate all subset sum & compare == k

a) Using Bit manipulations TC: O($2^N * N$)

$$\text{arr}[5] = \begin{bmatrix} 0 & 1 & 2 \\ 5 & 2 & 3 \\ 1 & 6 & 1 \\ 0 & & \end{bmatrix} \quad k = 7$$



Note:

Only when $i == N$,

We are checking if $\text{sum} == k$

func: given arr[], size, target sum
parameters: arr[], N, k, i, sum → current sum

function calls: 2 function calls

return type: int

```
int Subset(int arr[], int N, int k, int i, int sum){
```

```
    if (i == N) { if (sum == k) return 1 else return 0 }
```

```
    int c = 0
```

```
    sum = sum + arr[i] // updated
```

```
c = c + Subset(arr, N, k, i+1, sum)
```

} pick i^{th} element

```
sum = sum - arr[i] // subtract
```

```
c = c + Subset(arr, N, k, i+1, sum)
```

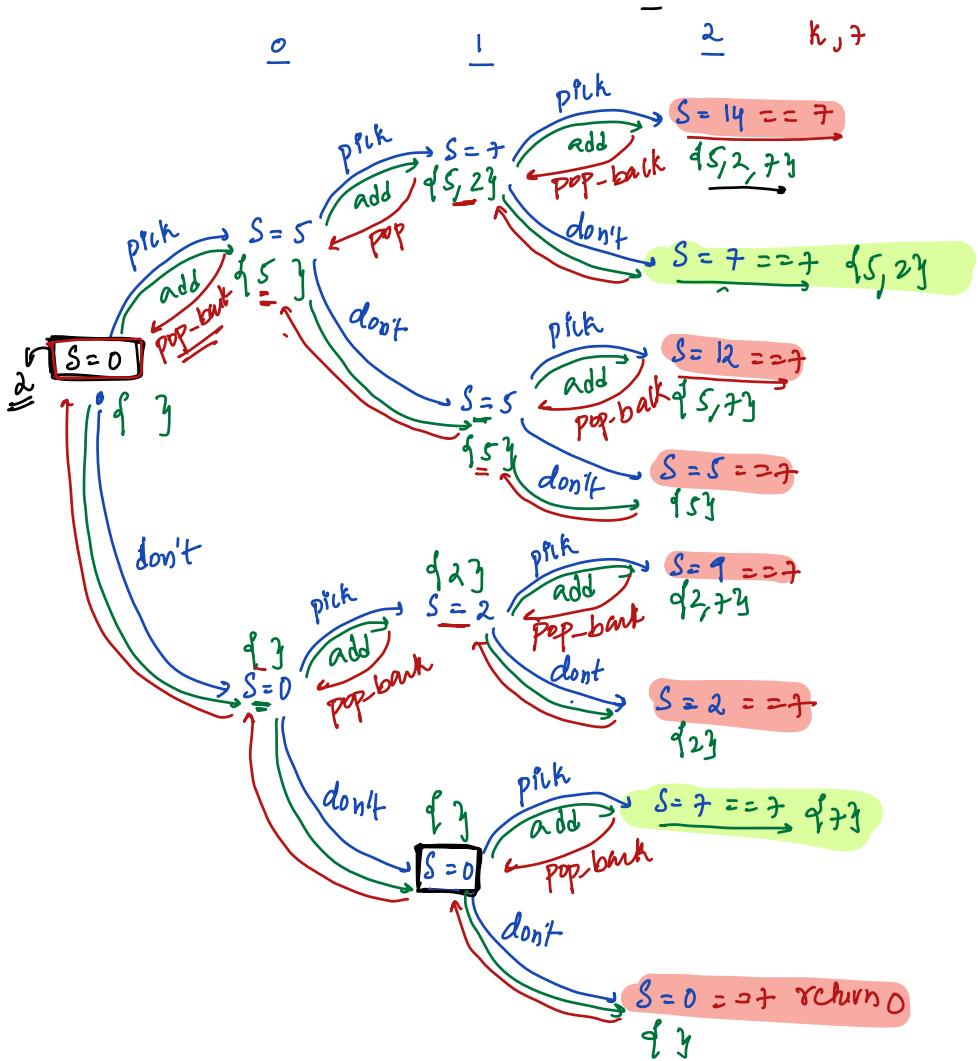
} leave i^{th} element

```
return c;
```

T_C: $2^N * \{1\}$ S_C: $O(N)$
 ↳ stack space

// 2Q) Print all subset with sum == k.

add: Will add element at back
pop back: Will remove last element
array-list is pass by reference



func: given arr[] size target sum
parameters: arr[], N, k, i, sum, current sum, current position

function calls: 2 function calls
return type: void

```
void printSubSum(int ar[], int N, int k, int i, int sum, list<int> l){
```

```
    if (i == N) {
```

```
        if (sum == k) { print(l); }
```

```
        return;
```

```
}
```

```
    sum = sum + ar[i]; } Made changes
```

```
l.add(ar[i]);
```

```
printSubSum(ar, N, k, i+1, sum, l);
```

```
sum = sum - ar[i]; } revert changes
```

```
l.pop_back();
```

```
printSubSum(ar, N, k, i+1, sum, l);
```

```
}
```

// In general: Backtracking *function calls*: parameter how many choices are present
return type.

```
fun( _____ )
```

Base Case:

```
    return;
```

At ith node:

Make changes

fun(_____)

revert changes

Make other calls: