

## Strongly Connected Components

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## Strongly Connected Components

**Definition:** a strongly connected component (SCC) of a directed graph  $G=(V,E)$  is a **maximal** set of vertices  $U \subseteq V$  such that

- For each  $u, v \in U$  we have both  $u \mapsto v$  and  $v \mapsto u$   
i.e.,  $u$  and  $v$  are **mutually reachable** from each other ( $u \rightleftarrows v$ )

Let  $G^T=(V,E^T)$  be the **transpose** of  $G=(V,E)$  where

$$E^T = \{(u,v) : (v,u) \in E\}$$

- i.e.,  $E^T$  consists of edges of  $G$  with their directions reversed

Constructing  $G^T$  from  $G$  takes  $O(V+E)$  time (adjacency list rep)

Note:  $G$  and  $G^T$  have the same SCCs ( $u \rightleftarrows v$  in  $G \Leftrightarrow u \rightleftarrows v$  in  $G^T$ )

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## Strongly Connected Components

### Algorithm

- Run **DFS**( $G$ ) to compute finishing times for all  $u \in V$
- Compute  $G^T$
- Call **DFS**( $G^T$ ) processing vertices in main loop in decreasing  $f[u]$  computed in Step (1)
- Output vertices of each **DFT** in **DFF** of Step (3) as a separate **SCC**

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## Strongly Connected Components

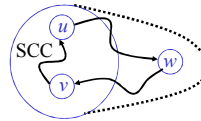
**Lemma 1:** no path between a pair of vertices in the same SCC, ever leaves the SCC

**Proof:** let  $u$  and  $v$  be in the same SCC  $\Rightarrow u \rightleftarrows v$

let  $w$  be on some path  $u \mapsto w \mapsto v \Rightarrow u \mapsto w$

but  $v \mapsto u \Rightarrow \exists$  a path  $w \mapsto v \mapsto u \Rightarrow w \mapsto u$

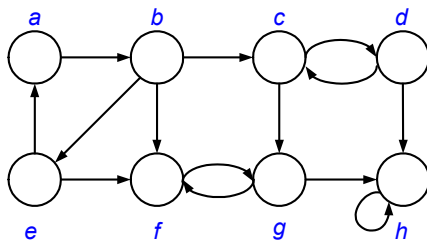
therefore  $u$  and  $w$  are in the same SCC



QED

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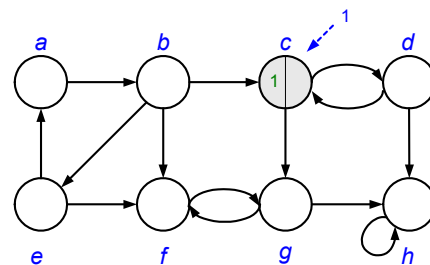
## SCC: Example



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## SCC: Example

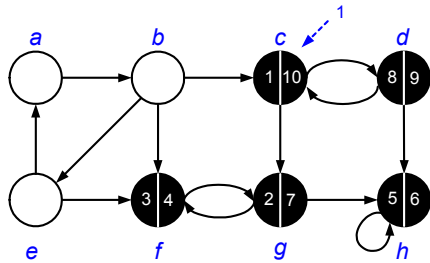
- Run **DFS**( $G$ ) to compute finishing times for all  $u \in V$



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### SCC: Example

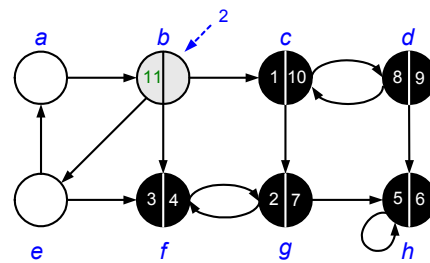
(1) Run  $\text{DFS}(G)$  to compute finishing times for all  $u \in V$



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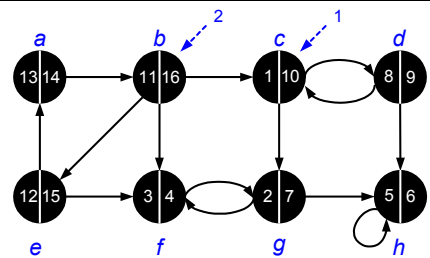
### SCC: Example

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### SCC: Example



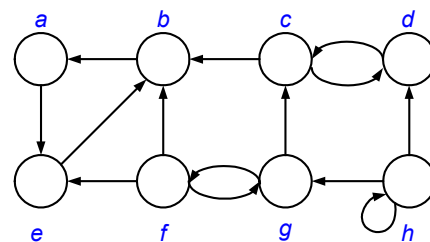
Vertices sorted according to the finishing times:

$\langle b, e, a, c, d, g, h, f \rangle$

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### SCC: Example

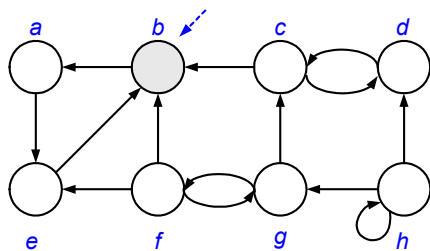
(2) Compute  $G^T$



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### SCC: Example

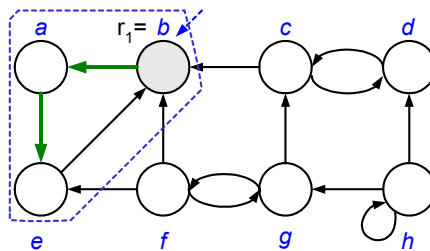
(3) Call  $\text{DFS}(G^T)$  processing vertices in main loop in decreasing  $f[u]$  order:  $\langle b, e, a, c, d, g, h, f \rangle$



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### SCC: Example

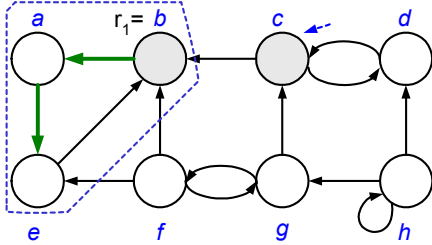
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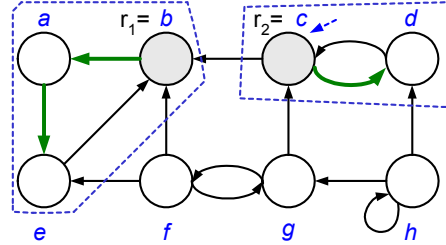
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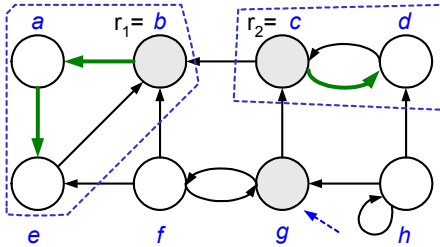
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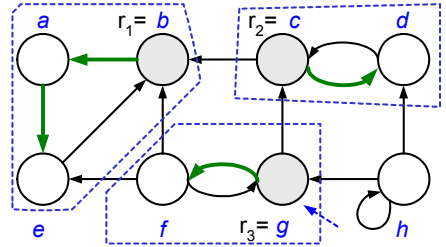
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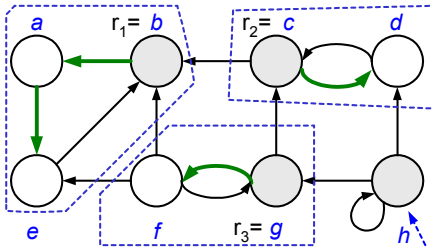
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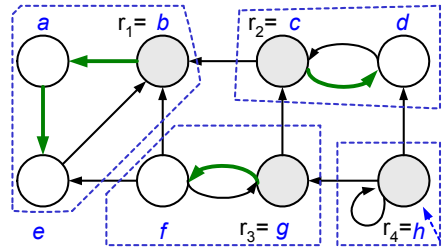
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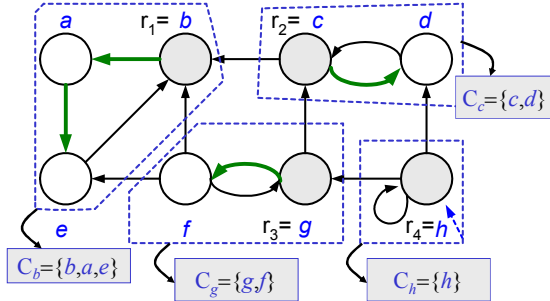
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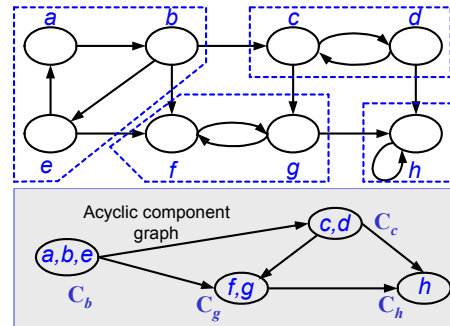
## SCC: Example

(4) Output vertices of each DFT in DFF as a separate SCC



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## SCC: Example



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