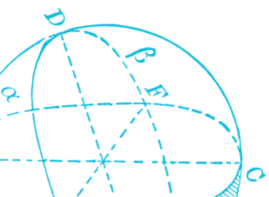
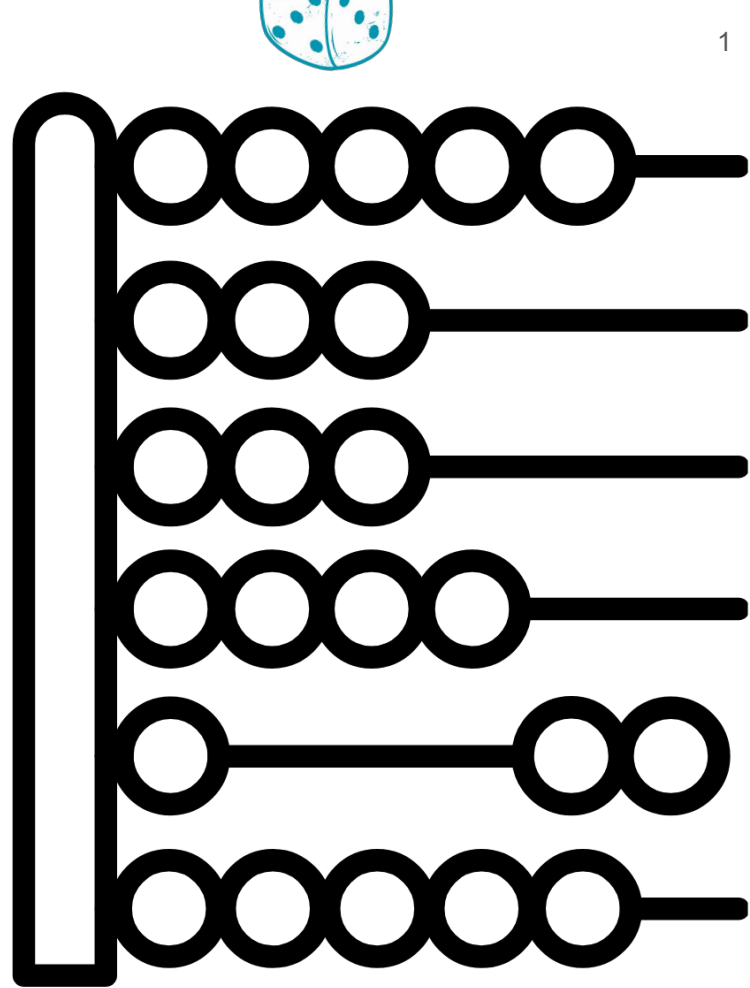
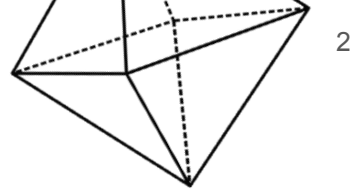


Numerics II

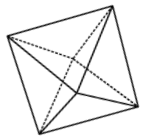
A Deeper Dive into CP Math



Lecture Objectives

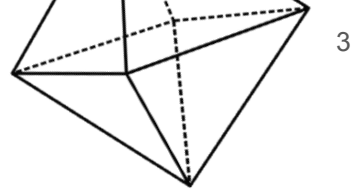


- Understanding the principles of hashing, its applications, and how hash collisions occur and are managed.
- Exploring advanced concepts in modular arithmetic in greater depth.
- Developing combinatorial reasoning skills for solving problems related to arrangements, selections, and distributions.
- Understanding Extended Euclidean Algorithm.

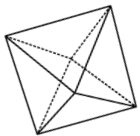


Lecture Outline

- Prerequisites
- Hashing and Hash Collisions
- Modular Arithmetic
- Combinatorics and Probability
- Extended Euclidean Algorithm
- Quote of the Day

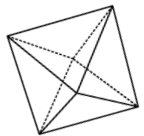
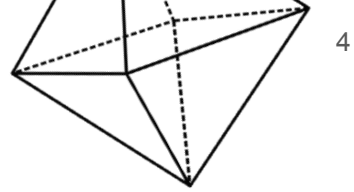


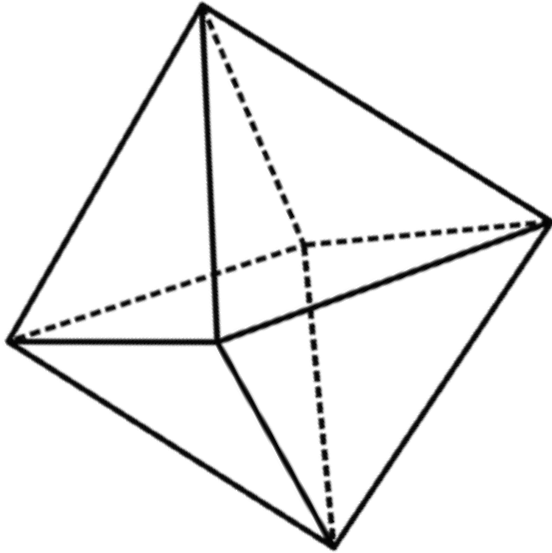
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Prerequisites

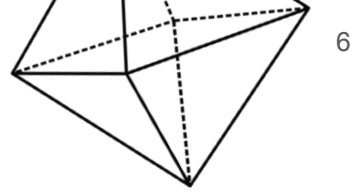
- Math I
 - Basic Understanding of Modular Arithmetic
 - Euclidean Algorithm
- Bit Manipulation
- Time and Space Complexity Analysis





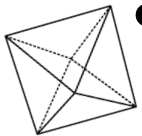
Hashing

What is Hashing?

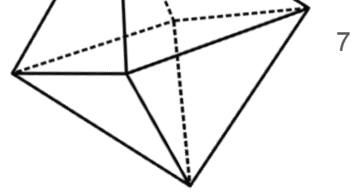


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- Hashing is a process that converts **input data** of **any size** into a **fixed-size** string of characters/letters known as **hash value**.
- Examples:
 - Password Storage
 - Python's built-in hash function
- What is the purpose then?

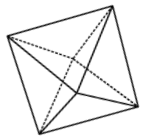


What is Hashing? - Hash Functions

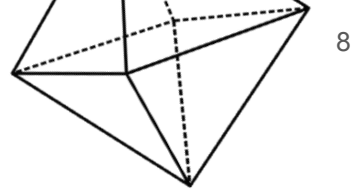


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- In hashing, keys of **arbitrary size** are converted into keys of **fixed-size** by using **hash function**.
- What is a hash function? A **good** hash function has the following properties
 - Deterministic
 - Fixed-Size Output
 - Efficient Computation
 - Preimage Resistant
 - Second Preimage Resistance
 - Avalanche Effect
 - Uniform Distribution
 - Collision-Resistant

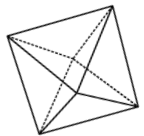


Hashing



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- Hashing is implemented in two steps:
 - **Hash Generation:** A hash function is used to generate a **fixed-size integer or string** for a particular input
 - **Storage:** The generated hash is stored in a **large table** as a key along with the input as the value



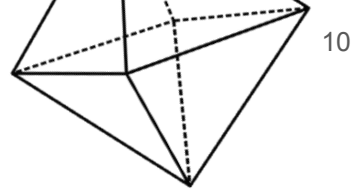
Hash Collision

- Let hash be a hashing function, and let x_1 and x_2 be two **different** inputs. A hash collision happens when

$$\text{hash}(x_1) = \text{hash}(x_2)$$

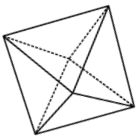
- Example: if `hash = lambda x: sum(ord(c) for c in x)`, then `hash('abc') = hash('cab')`.

Hash Collision

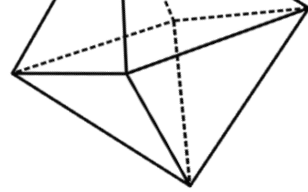


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- How can we handle hashing collisions?
- Python uses hashing for basic data structures like **sets** and **dictionaries**. What are some implications of hashing collisions in competitive programming?



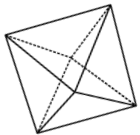
Modular Hashing



- Modular Hashing is mapping a key k into one of m slots by taking the remainder of k divided by m .

$$\text{hash}(k) = k \% m$$

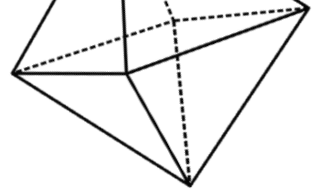
- Modular Hashing is a **too simple** to be practical for real world applications. Which properties of a good hashing function does it satisfy and fail?



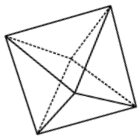


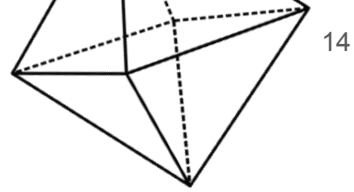
Modular Arithmetic

A Recap on Modular Arithmetic



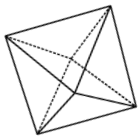
- Also called Clock Arithmetic
 - $(a + b) \% m = (a \% m + b \% m) \% m$
 - $(a - b) \% m = (a \% m - b \% m) \% m$
 - $(a * b) \% m = (a \% m * b \% m) \% m$
 - If $(a - b) \% m = 0$, then $a \% m = b \% m$
 - Division is a bit more complicated.



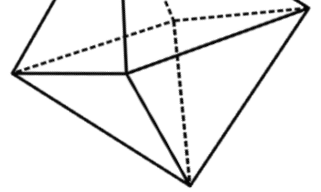


A Recap on Modular Arithmetic

- Mathematically, we can interpret two numbers that have the same remainder as the same number and have a working algebra. This is called **Modular Congruence**.
- For example: 17 is congruent to 8 modulo 9.
- We say $17 \% 9 = 8 \% 9$ or $17 = 8 \bmod 9$



Modular Division



- First of all, not all division is well-defined under modular arithmetic.

- Consider the fraction $1 / 2$ under modulo 6.
- We know $2 * (1 / 2) = 1 \bmod 6$.
- However . . .

$$\blacksquare 2 * 0 = 0$$

$$\blacksquare 2 * 3 = 0$$

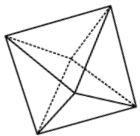
$$\blacksquare 2 * 1 = 2$$

$$\blacksquare 2 * 4 = 2$$

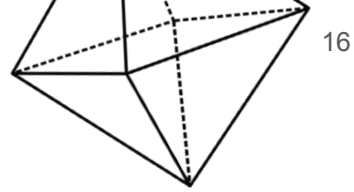
$$\blacksquare 2 * 2 = 4$$

$$\blacksquare 2 * 5 = 4$$

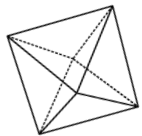
- Therefore $1 / 2$ is not defined under modulo 6.



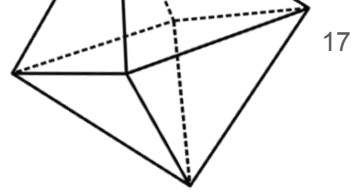
Modular Division



- For a division to be well-defined in modular arithmetic, **the divisor has to be relatively prime with the modulo.**
- This is also a sufficient condition for the division to be well-defined.

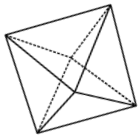


Modular Division - Naive Way

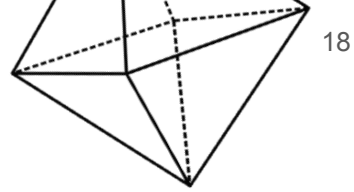


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- If a is relatively prime with m , we can find the $(1/a) \% m$ by iterating a variable inv_a from 0 to $m - 1$ and checking if $(inv_a * a) \% m == 1$. Why?



Modular Division - Fermat's Way



- When m is prime, we have Fermat's little theorem:

$$1/a = a^{m-2} \bmod m$$

Modular Division - Euler's Way (generalization)

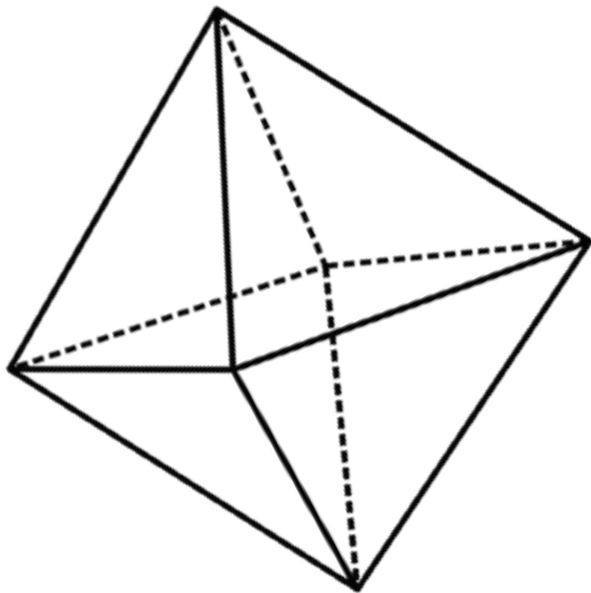
- If m is not a prime, there is still a way. We can use the following fact in number theory as follows:

$$1/a = a^{**}(\text{phi}(m) - 1) \bmod m$$

- Where $\text{phi}(n)$ is the [Euler's totient function](#).

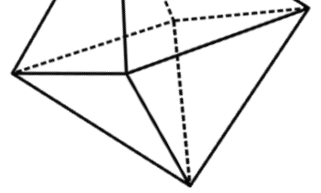
[Proof](#)

[Implementation of Euler's totient function](#)

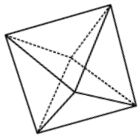


Binary Exponentiation

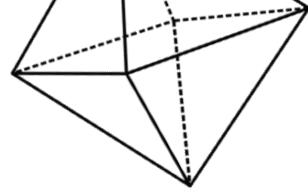
Binary Exponentiation



- Binary exponentiation (also known as exponentiation by squaring) is a trick which allows to calculate $a^{**}n$ using only $O(\log n)$ multiplications.
- Let's, calculate $3^{**}(13)$.



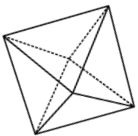
Binary Exponentiation



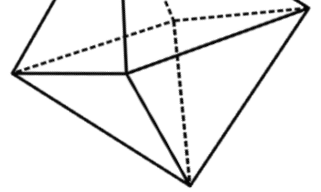
- The idea of binary exponentiation is, that we split the work using the binary representation of the exponent.

$$3^{13} = 3^{1101_2} = 3^8 \cdot 3^4 \cdot 3^1$$

- In general, since the number n has exactly $\log_2 n + 1$ digits in base 2, we only need to perform $O(\log_2 n)$ multiplications, if we know the powers a^{**0} , a^{**1} , a^{**2} , a^{**4} etc.



Binary Exponentiation



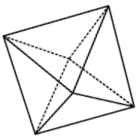
- Luckily, this is very easy. Since an element in the sequence is just the square of the previous element.

$$3^1 = 3$$

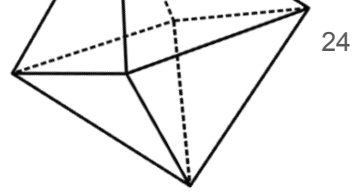
$$3^2 = (3^1)^2 = 3^2 = 9$$

$$3^4 = (3^2)^2 = 9^2 = 81$$

$$3^8 = (3^4)^2 = 81^2 = 6561$$

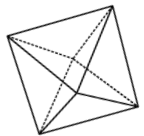


Binary Exponentiation

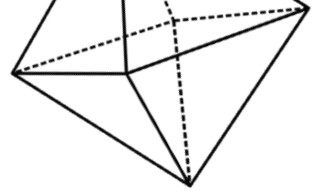


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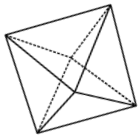
- Implementation Steps
 1. Initialize three variables:
 - a. `result` to 1
 - b. `base` to the given base
 - c. `exponent` to the given exponent

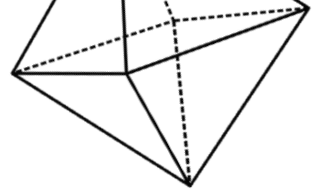


Binary Exponentiation



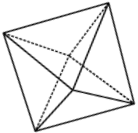
- Implementation Steps
 2. While the exponent is greater than 0:
 - a. If the **least significant bit (LSB)** of the exponent is 1, multiply result by base.
 - b. Square the base.
 - c. Right-shift exponent by 1.
 3. Return result as the final result.





Binary Exponentiation - Iterative

```
def binary_exponentiation(base, exponent):  
    result = 1  
    while exponent > 0:  
        if exponent & 1:  
            result *= base  
        base *= base  
        exponent >>= 1  
    return result
```



Modular Arithmetic - Template

```
class ModularArithmetic:
    def add(self, a, b, p):
        return ((a % p) + (b % p)) % p

    def subtract(self, a, b, p):
        return ((a % p) - (b % p)) % p

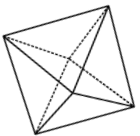
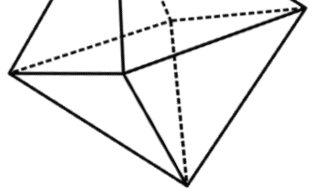
    def multiply(self, a, b, p):
        return ((a % p) * (b % p)) % p
```

Modular Arithmetic - Template

```
def binary_exponentiation(self, base, exponent, p):  
    result = 1  
    while exponent > 0:  
        if exponent & 1:  
            result = self.multiply(base, result, p)  
        base = self.multiply(base, base, p)  
        exponent >>= 1  
    return result
```

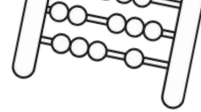
```
def inverse(self, a, p):  
    return self.binary_exponentiation(a, p - 2, p)
```

```
def division(self, a, b, p):  
    return self.multiply(a, self.inverse(b, p), p)
```



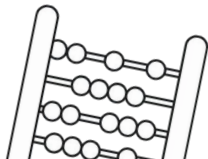
Combinatorics

“Combinatorics is the art of counting without counting.” - Andrzej Schinzel

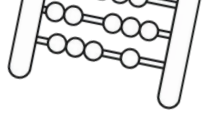


The Addition Rule

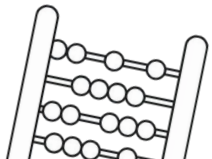
- Suppose two experiments are to be performed.
 - If experiment **E1** has **n_1** possible outcomes
 - And experiment **E2** has **n_2** possible outcomes,
- then, when ***performed independently***, the total possible outcomes is **$n_1 + n_2$** .



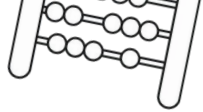
The Addition Rule



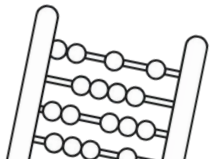
- Example:
 - Suppose you have 6 roads and 3 railways from Addis Ababa to Berland. How many possible ways do you have to go Berland?
 - Number of ways = 6 roads + 3 railways = 9.

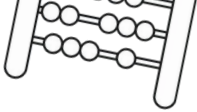


The Multiplication Rule



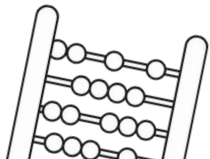
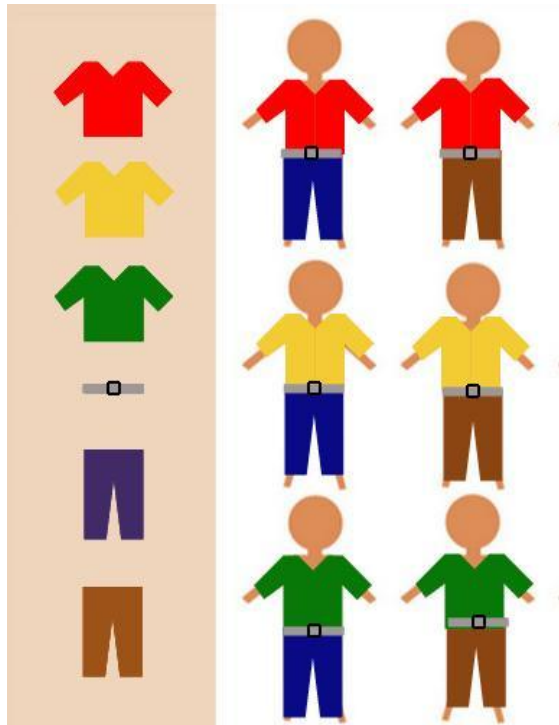
- Suppose that two experiments are to be performed.
 - If there are m outcomes in experiment $E1$.
 - And for each outcome of experiment $E1$, there are n outcomes in experiment $E2$.
- Then, the total possible outcomes of **running the experiments back to back** is $m * n$.





The Multiplication Rule

- Example:



Example: Vowels of All Substrings

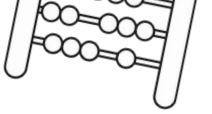
- Given a string word, return the sum of the number of vowels ('a', 'e', 'i', 'o', and 'u') in every substring of word.
- Example:

Input: “aba”

Output: 6

Substrings: “a”, “ab”, “aba”, “ba”, “a”

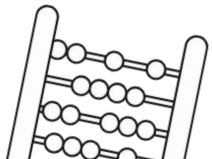
$$1 + 1 + 2 + 1 + 1 = 6$$

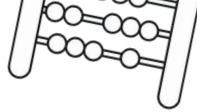


Example: Vowels of All Substrings


d b a c c d

- The only vowel we have is 'a'. How do we count the number of substrings 'a' appears in?

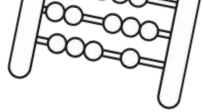




Basic rules of counting

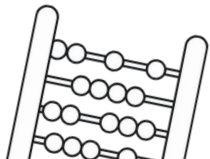

d b a c c d

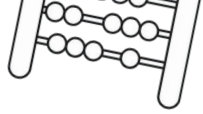
The substring should start from 'd', 'b', 'a' inorder to include 'a'. In our case, we have 3 possible starting points.



Basic principle of counting

- For any given index i , there are $i + 1$ possible starting indices that allow us to include the character at index i .



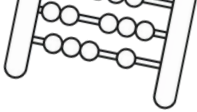


Basic rule of counting

↓ ↓ ↓ ↓ ↓
d b a c c d

dba
dbac
dbacc
dbaccd

If we start from 'd' we have 4 possible ending indices. This means we can form 4 substrings starting from 'd'.



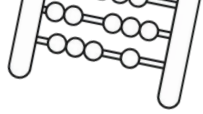
Basic rule of counting

d b a c c d

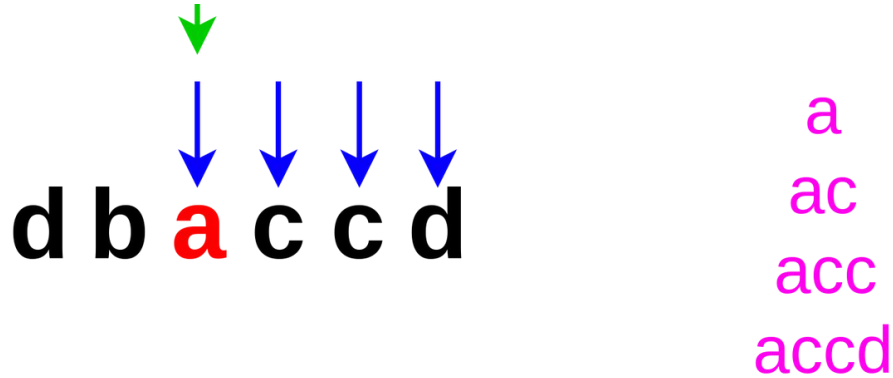
Diagram illustrating the basic rule of counting substrings starting from 'b'. Arrows point to the characters: a green arrow points to 'b', and four blue arrows point to 'a', 'c', 'c', and 'd' respectively.

ba
bac
bacc
baccd

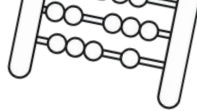
Starting from 'b' we can also form 4 substrings.



Basic rule of counting

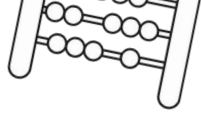


Finally starting from 'a' we have 4 possible substrings.

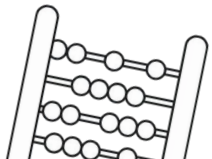


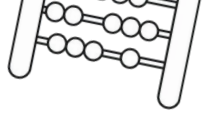
Basic rule of counting

- In general,
 - For any vowel at index i (0-indexed), there are $i + 1$ possible starting indices.
 - For each of these starting indices j , there are $n - i$ possible ending indices.



Which counting principle is applicable in this case?



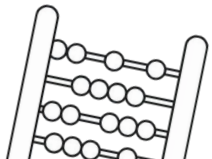


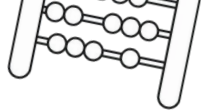
Basic rule of counting

Hence for one vowel, there are

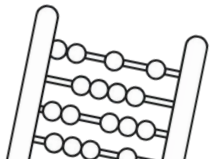
$$(i + 1) * (n - i)$$

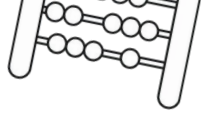
substrings that include the vowel.



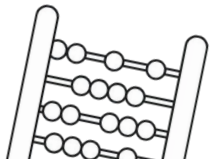


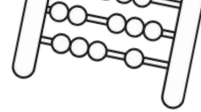
What if we have more than
one vowel in our string?



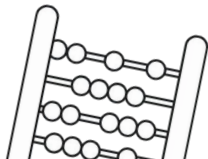


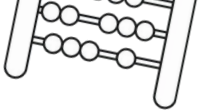
We use the **addition rule**.





Permutations



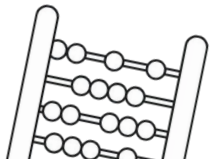


Permutation

- A permutation is an **arrangement of objects without repetition** where **order is important**.

Examples

- Ways to arrange 3 letters: ABC, ACB, BCA, BAC, CAB, CBA
- Ways to arrange x bricks where consecutive bricks have different colors



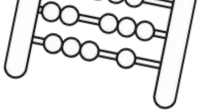
Permutation

- In general, if we have n objects and want to arrange, then we will have a total of $n!$. If $n!$ is positive integer, then

$$n! = n * (n - 1) * (n - 2) * \dots * (1)$$

$$n! = n * (n - 1) !$$

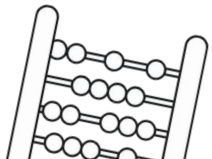
$$1! = 1$$



Permutation

- If we need to arrange r objects ($r \leq n$) among n distinct objects, we will have a total of $P(n, r)$ different ways.

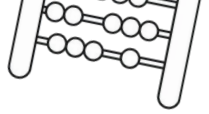
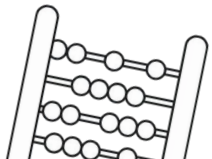
$$P(n, r) = n! / (n - r)!$$

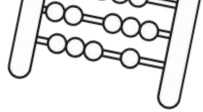


Permutation



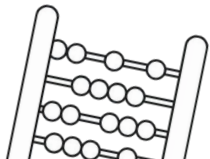
In how many ways can we
arrange the letters in PEPPER?





Let's reverse the problem

- How many ways can we arrange $P_1E_1P_2P_3E_2R_1$ (with indices)? $6!$ ways.
- Let's assume there are x ways to arrange PEPPER (without the indices).
- For a single arrangement, in how many ways can we index the repeated letters?

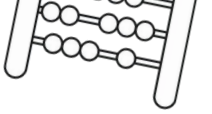


Answer

$3! * 2! * 1!$ ways

Therefore: $6! = x * 3! * 2!$

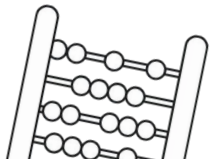
$$\Rightarrow x = 6! / (3! * 2!)$$

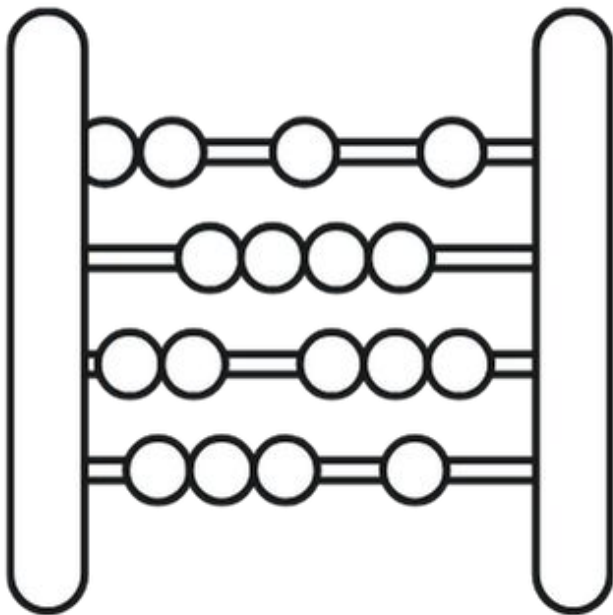


Permutation

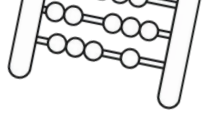
If we need to arrange n objects where $n_1, n_2, n_3 \dots n_r$ are alike. Then we will have

$$\frac{n!}{n_1! \times n_2! \times \dots \times n_r!}$$



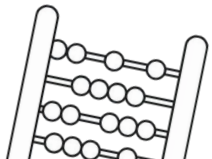


Combinations



Combinations

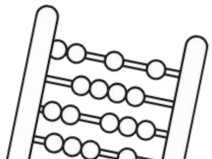
Combination is **a way of selecting items** from a collection, such that **(unlike permutations) the order of selection does not matter.**





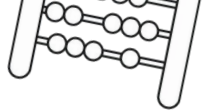
Combinations

- Determine the number of different groups of 3 objects from 5 items A, B, C, D and E.



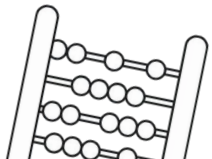
Combinations

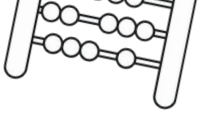
- Every group of 3 will be counted $3! = 6$ times. Example: ABC
 - ABC
 - ACB
 - BAC
 - BCA
 - CAB
 - CBA



Combinations

- Initially, there are 5 ways to select the first item.
- Then, there are 4 ways to select the next item.
- Finally, there are 3 ways to select the last item.
- There are thus $5 \times 4 \times 3$ ways of selecting the group(if order matters).

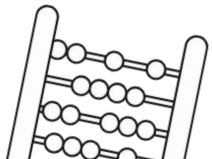


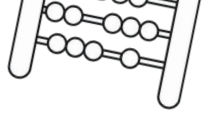


Combinations

- The total number of groups that can be formed is:

$$\frac{5 \cdot 4 \cdot 3}{3 \cdot 2 \cdot 1} = 10$$

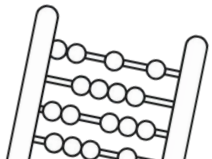


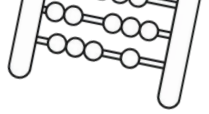


Combinations - Formula

Selecting r objects ($r \leq n$) among n objects can be done by

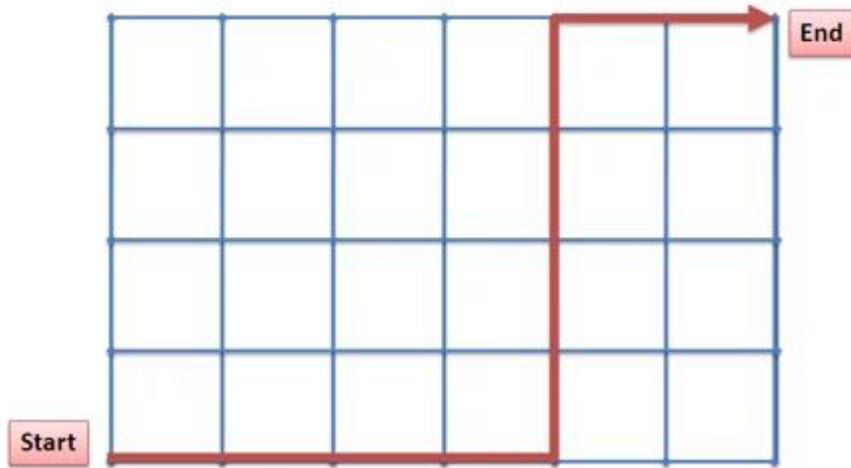
$${}^nC_r = \binom{n}{r} = \frac{n!}{r!(n-r)!}$$





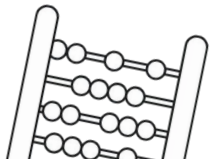
Combinations - Lattice Paths

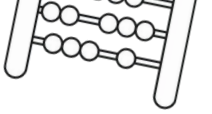
How Many Paths?



Suppose you're on a 4×6 grid, and want to go from the bottom left to the top right.

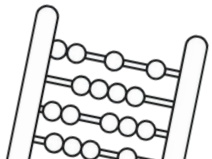
How many different paths can you take? We can only move **right** or **up**.





Combinations - Lattice Paths

- Let's write down a path using “U” (for up) and “R” for (right).
- Path: **R R R R R R U U U U**
- That is, go all the way right (6 R'S) then all the way up (4 U's)

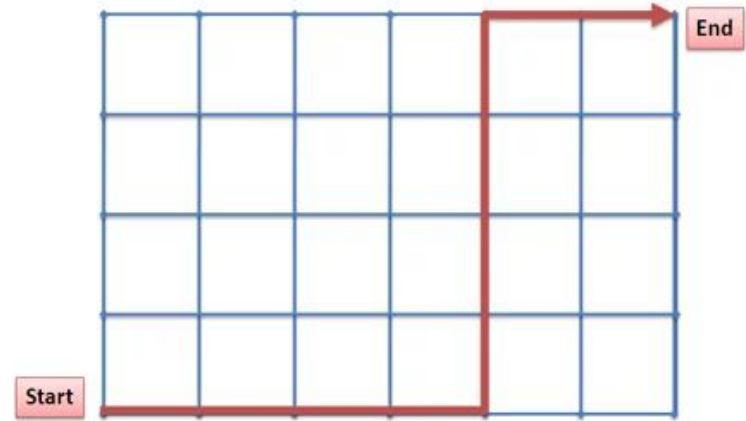


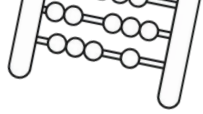
Combinations - Lattice Paths

- For the diagram the path is

R R R R U U U U R R

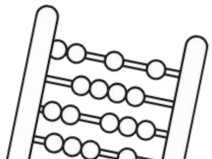
How Many Paths?

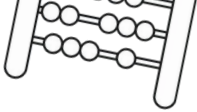




Combinations - Lattice Paths

- The question now becomes “In how many ways can we rearrange 4 U’s and 6 R’s.”



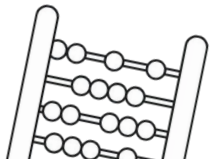


Combinations - Approach

Imagine we start with **10** R's

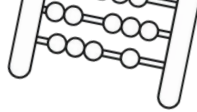
Path: **R R R R R R R R R R**

Clearly, we need to change **4** of those **R**'s into **U**'s.





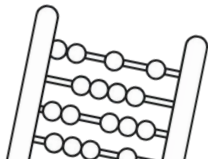
In how many ways can
we pick the 4R's
to change them to U's?



Combinations - Lattice Paths

- We have **10** choices for the first **R** to convert
- Then **9** for the second
- Then **8** for the third
- Finally **7** for the last

There are **$10 * 9 * 8 * 7 = 5040$** possibilities



But Wait!

- We need to remove the redundancies:
- Converting **#1, #2, #3 and #4** is the same as converting **#4, #3, #2, and #1** or **#3, #1, #2 and #4** and soon.
- We have **4!** ways to rearrange the **U**'s positions we picked, so finally we get:
- $5040 / 24 = 210 \text{ ways} = C(10, 4) = C(10, 6)$



Probability



Probability

- **Sample space:**
 - The set of **all possible outcomes** of an experiment.
 - Denoted by S .
- Example: If the experiment consists of flipping two coins then

$$S = \{(H, H), (H, T), (T, H), (T, T)\}$$





Probability

- **Event Space:**
 - The set of **desired** outcomes of the experiment.
 - Denoted by E .
- Example: if E is the event that a head appears on the first coin, then $E = \{(H, H), (H, T)\}$





Probability

The probability of an event E is the number of outcomes favourable to E divided by the total number of outcomes.

$$\text{Probability of an event happening} = \frac{\text{Number of ways an event can occur}}{\text{Total number of possible outcomes}}$$





Rules of Probability

- Rule 1: The probability $P(A)$ of any event A satisfies $0 \leq P(A) \leq 1$
- Rule 2: If S is the sample space in a probability model, then $P(S) = 1$
- Rule 3: If A and B are disjoint. $P(A \text{ or } B) = P(A) + P(B)$
- Rule 4: For any event A , $P(A \text{ does not occur}) = 1 - P(A)$





Probability - Exercise

- Suppose you throw two dice. We are interested in the sum of the upper face of the dice. Let E be the event that the sum of the dice is odd. Find $P(E)$





Conditional Probability

- The conditional probability of an event A is the probability that the event will occur given the knowledge that an event B has already occurred.

*Probability of event A
and event B*

$$P(A|B) = P(A \text{ and } B) / P(B)$$

*Probability of event A
given B has occurred*

*Probability
of event B*





Conditional Probability

- The probability that both A and B have occurred together is

$$P(A \text{ and } B) = P(A|B) P(B)$$





Extended Euclidean Algorithm

Extended Euclidean Algorithm

- The Extended Euclidean Algorithm is an extension of the Euclidean Algorithm.
- It finds the greatest common divisor (gcd) of two integers and also finds coefficients (often denoted as x and y) such that

$$ax + by = \gcd(a, b)$$

- It is also true that $\gcd(a, b)$ is the smallest positive integer for any values of x and y .

Practice Problems

☒ [Pow\(x, n\)](#)

☒ [Count Good Numbers](#)

☒ [Unique Paths](#)

☒ [Selection of Personnel](#)

☒ [Make it Alternating](#)

☒ [Number of Ways to Reach a Position After Exactly k Steps](#)

☒ [Vowels of All Substrings](#)

☒ [Make Sum Divisible by P](#)

☒ [Password](#)

☒ [Dictionary](#)

☒ [Beautiful Numbers](#)

References

- ✓ [CP-Algorithms](#)
- ✓ [Numerics Lecture - A2SV Gen 3 Camp 22](#)
- ✓ [CP-Algorithms](#)

**He who is not courageous
enough to take risks will
accomplish nothing in life.**

Muhammad Ali

quotefancy

