자료구조

L09 Sorting (3)

2022년 1학기

국민대학교 소프트웨어학부

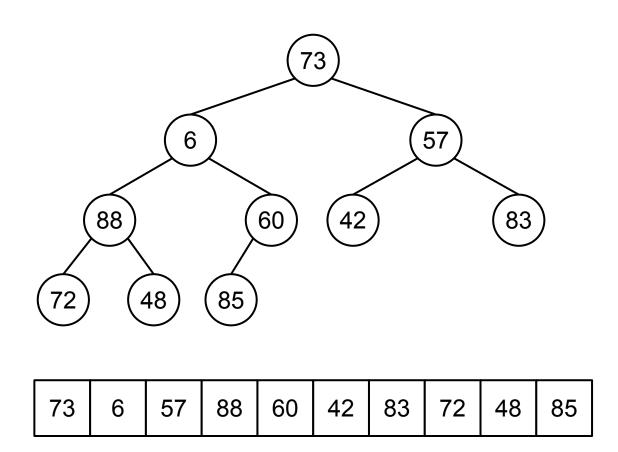
Overview

- Heapsort
- Binsort
- Radix Sort

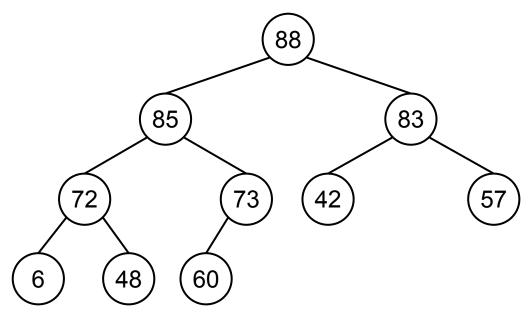
Heapsort

- Heap 자료구조 활용
- 크기 n의 배열이 입력으로 주어지면
 - max-heap 을 만들고
 - max값을 n번 꺼낸다

Original numbers



- Build heap
 - Time complexity?

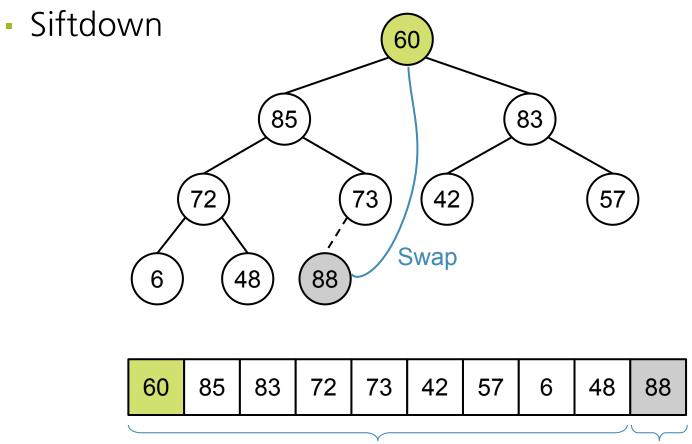


88 85 83 72 73 42 57 6 48 60

public E removemax() {
 swap(0, --n);
 if (n != 0) siftdown(0);
 return Heap[n];

Remove 88

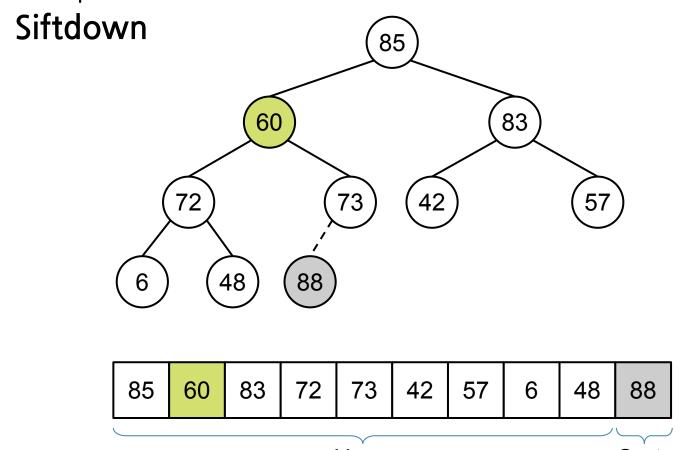
Swap the root and the last node



public E removemax() {
 swap(0, --n);
 if (n != 0) siftdown(0);
 return Heap[n];

Remove 88

Swap the root and the last node



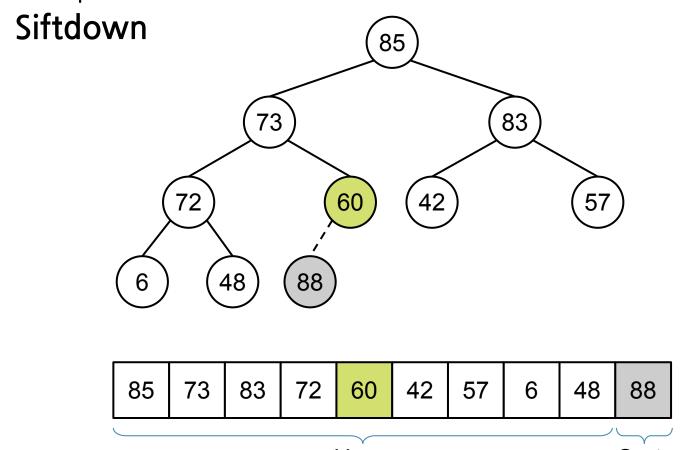
Heap

Sorted

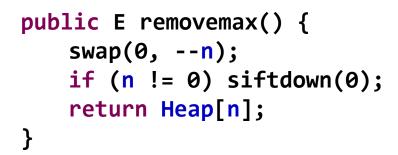
public E removemax() {
 swap(0, --n);
 if (n != 0) siftdown(0);
 return Heap[n];

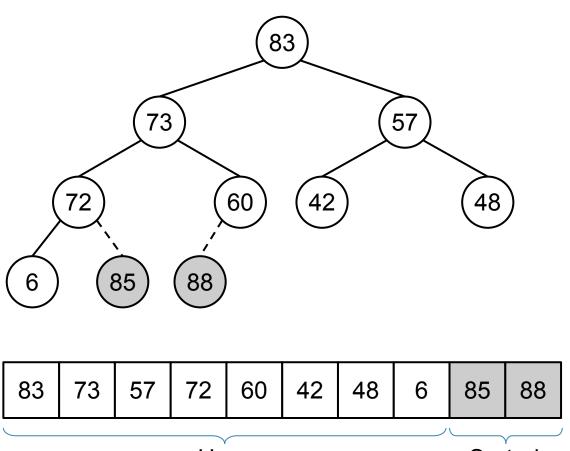
Remove 88

Swap the root and the last node

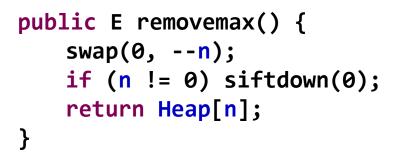


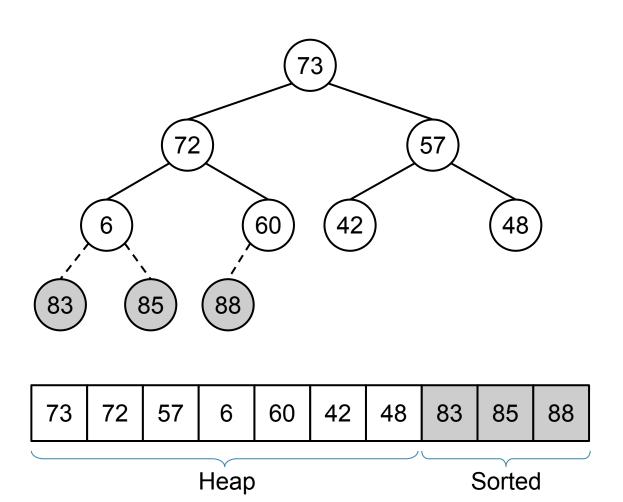
Remove 85





Remove 83





Heapsort

```
static <E extends Comparable<? super E>> # items max size
void heapsort(E[] A) {
   MaxHeap<E> H = new MaxHeap<E>(A, A.length, A.length);
   for (int i = 0; i < A.length; i++)
        H.removemax();
}</pre>
```

- heapsort는 in-place sort이다
 - In-place sort: 추가 메모리 사용 없이 정렬함
- Cost of heapsort:
- Cost of finding k largest elements:

Heapsort

```
static <E extends Comparable<? super E>> # items max size
void heapsort(E[] A) {
   MaxHeap<E> H = new MaxHeap<E>(A, A.length, A.length);
   for (int i = 0; i < A.length; i++)
        H.removemax();
}</pre>
```

- heapsort는 in-place sort이다
 - In-place sort: 추가 메모리 사용 없이 정렬함
- Cost of heapsort: $\Theta(n + n \log n) = \Theta(n \log n)$
- Cost of finding k largest elements: $\Theta(n + k \log n)$

If k is small, Heapsort is very fast!

Overview

- Heapsort
- Binsort
- Radix Sort

Binsort Motivation

- 1부터 n까지의 key를 갖는 개체가 랜덤하게 섞여있는 배열이 있다.
- 가장 빠르게 정렬하는 방법은?

Binsort

- 단순하지만 효율적인 알고리즘:
 - 하지만, key가 0부터 n-1까지 순서대로 있을 때만 정렬가능

```
for (i=0; i<n; i++)
B[A[i].key()] = A[i];</pre>
```

- 중복된 키가 있거나, key의 범위가 n보다 큰 경우에도 처리 가능하게 하려면?
- Main Idea:
 - 중복된 키 허용: linked list를 담은 array 사용하기.
 - 각각의 bin이 linked list를 담는다
 - 배열 B를 더 크게 만든다 (크기가 MaxKeyValue+1이 되도록)

The max key value from the input

Binsort

```
static void binsort(Integer A[]) {
   List<Integer>[] B =
             (LList<Integer>[]) new LList[MaxKey];
   Integer item;
   for (int i = 0; i < MaxKey; i++)</pre>
      B[i] = new LList<Integer>();
   for (int i = 0; i < A.length; i++)
      B[A[i]].append(A[i]);
   for (int i = 0; i < MaxKey; i++)</pre>
      for (B[i].moveToStart();
         (item = B[i].getValue()) != null; B[i].next())
         output(item);
Cost: \Theta(n + MaxKeyValue)
```

Binsort

- Binsort의 강점
 - MaxKeyValue 가 작으면 빠르고, 공간 효율적이다.

- Binsort의 약점
 - MaxKeyValue가 크면, 느리고 공간 비효율적이다.

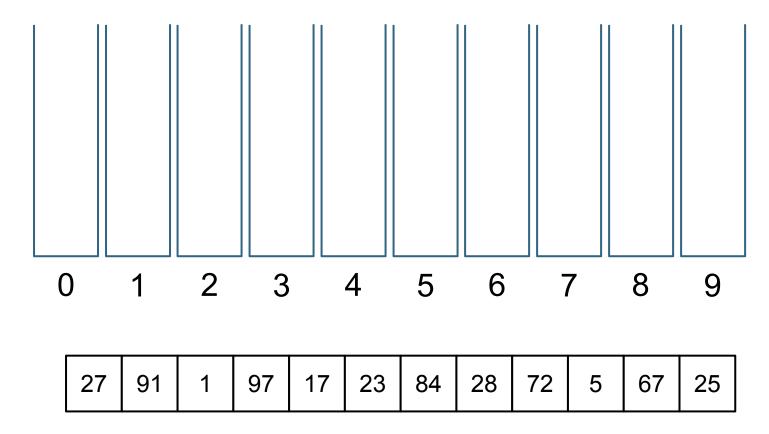
Overview

- Heapsort
- Binsort
- Radix Sort

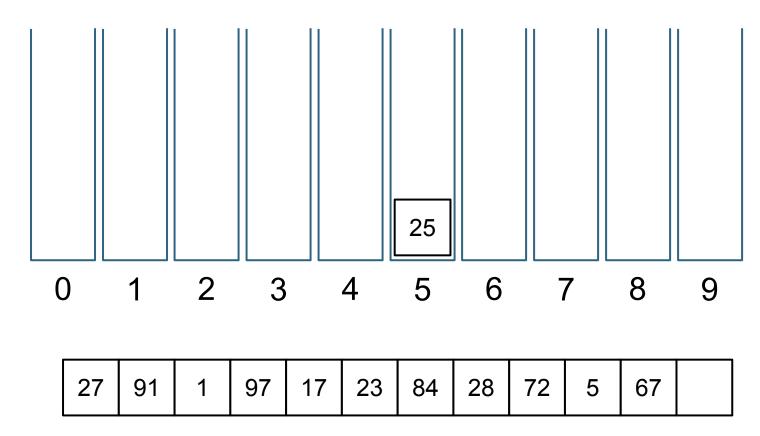
Radix Sort

- Binsort를 확장한 버전
 - radix sort는 binsort보다 더 공간 효율적
- key가 숫자나 짧은 문자열인 경우에만 사용 가능
- k자리 숫자들을 정렬하려면 k번 반복하여 연산
 - i번째 연산은 i번째 자릿수에 대해서 binsort를 수행
 - 첫번째 연산: 가장 낮은 자릿수 (LSD Least Significant Digit) 정렬
 - 각 반복 연산마다 bin에 있는 아이템들을 순서대로 조회

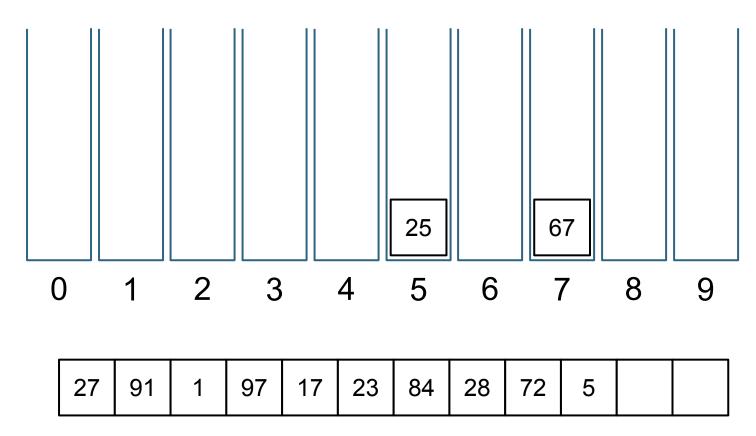
- First pass (on the 1st digit from right)
 - Put the elements to bins



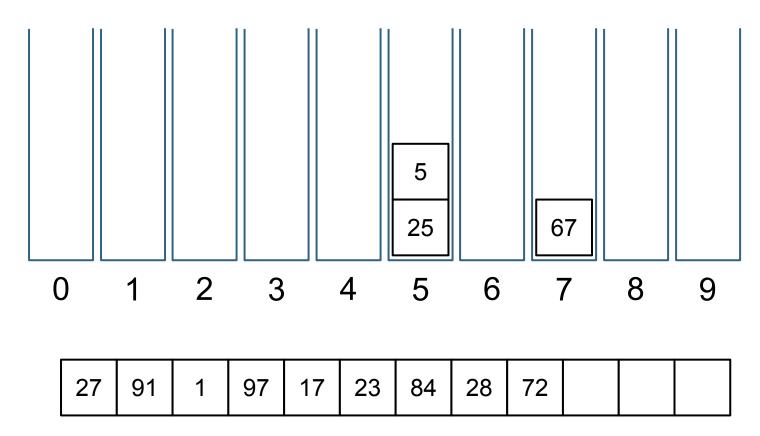
- First pass (on the 1st digit from right)
 - Put the elements to bins



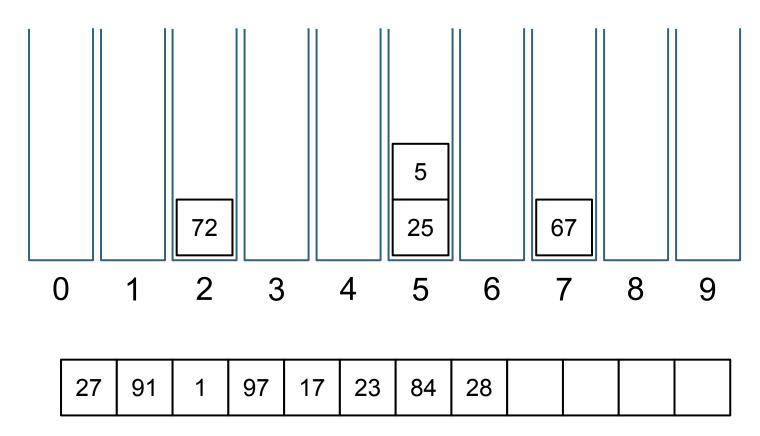
- First pass (on the 1st digit from right)
 - Put the elements to bins



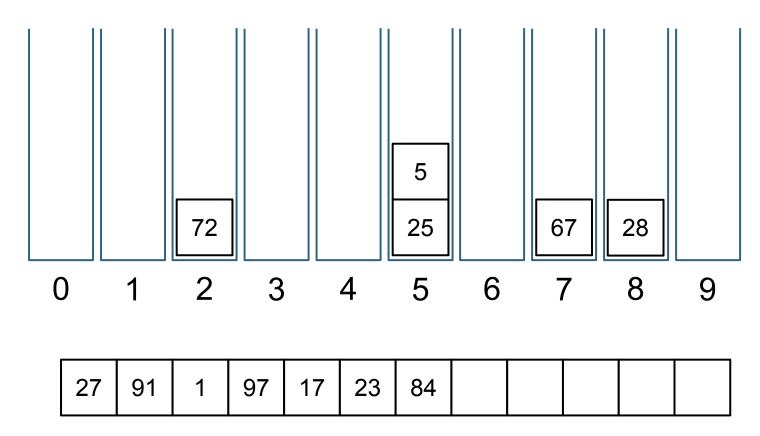
- First pass (on the 1st digit from right)
 - Put the elements to bins



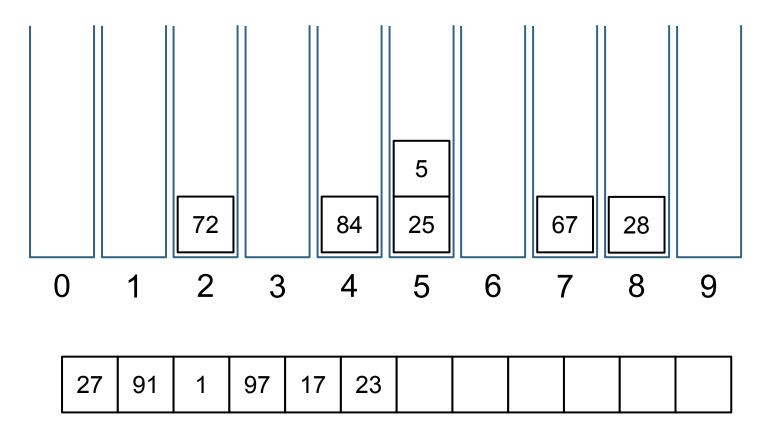
- First pass (on the 1st digit from right)
 - Put the elements to bins



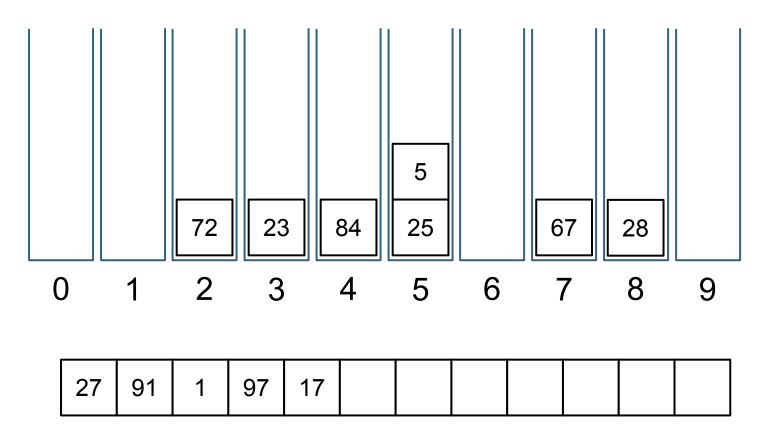
- First pass (on the 1st digit from right)
 - Put the elements to bins



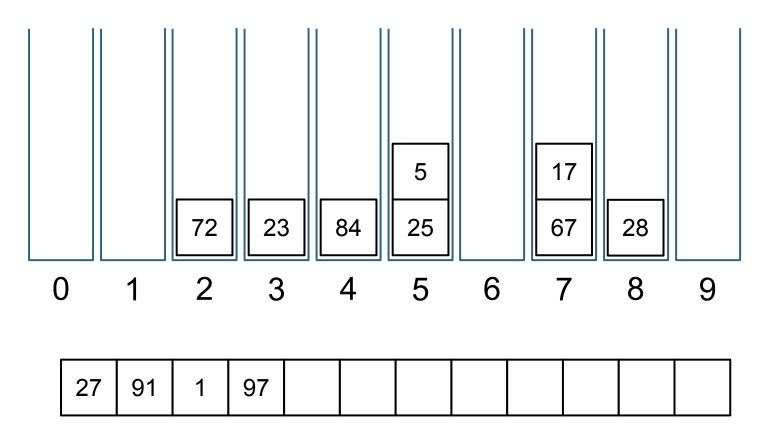
- First pass (on the 1st digit from right)
 - Put the elements to bins



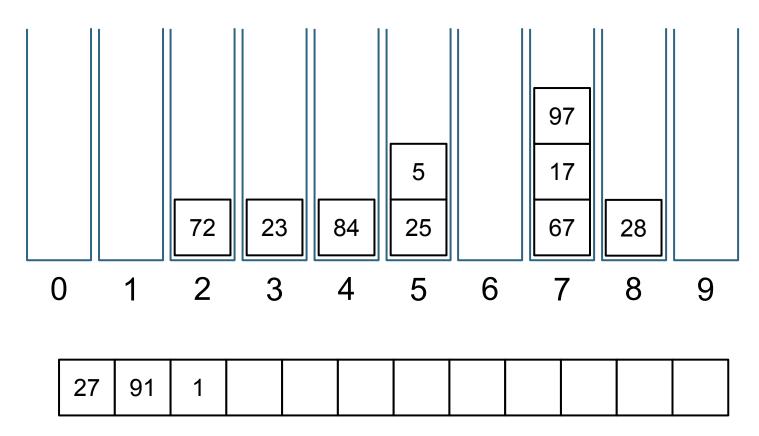
- First pass (on the 1st digit from right)
 - Put the elements to bins



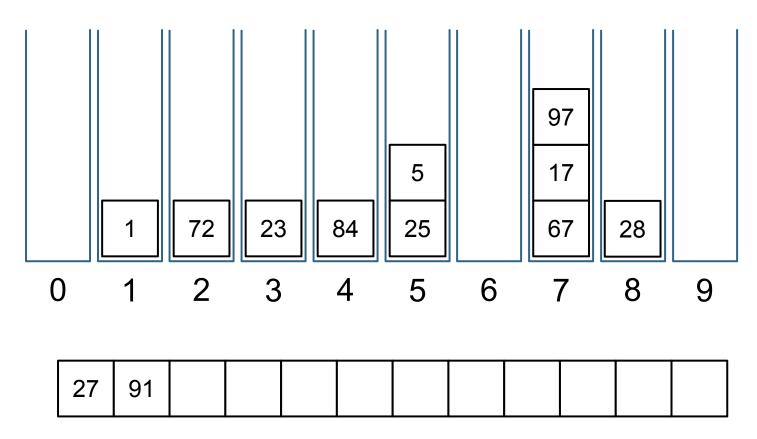
- First pass (on the 1st digit from right)
 - Put the elements to bins



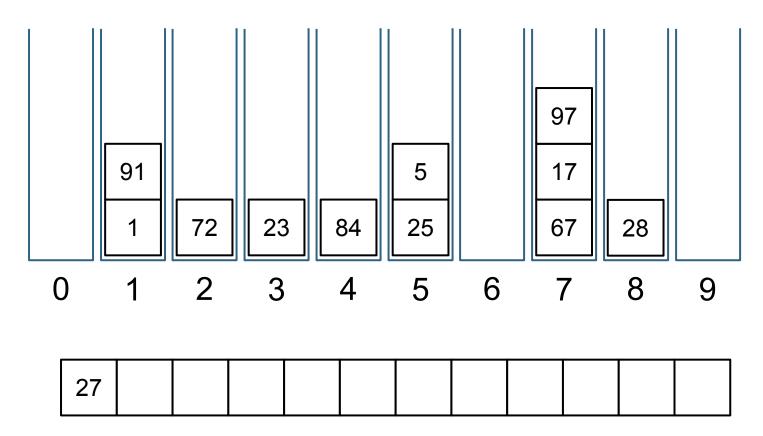
- First pass (on the 1st digit from right)
 - Put the elements to bins



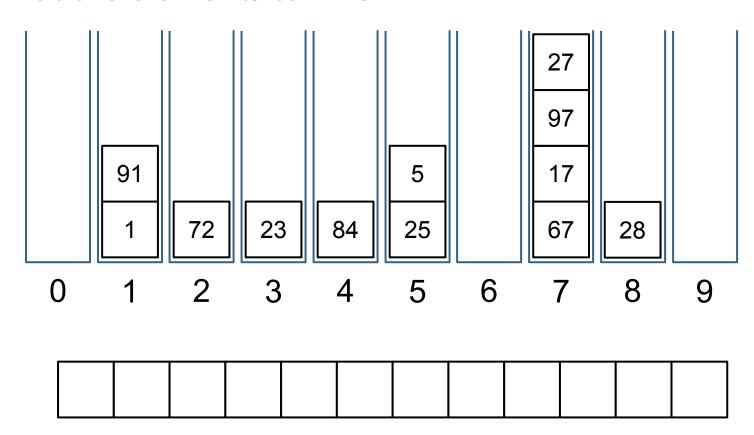
- First pass (on the 1st digit from right)
 - Put the elements to bins



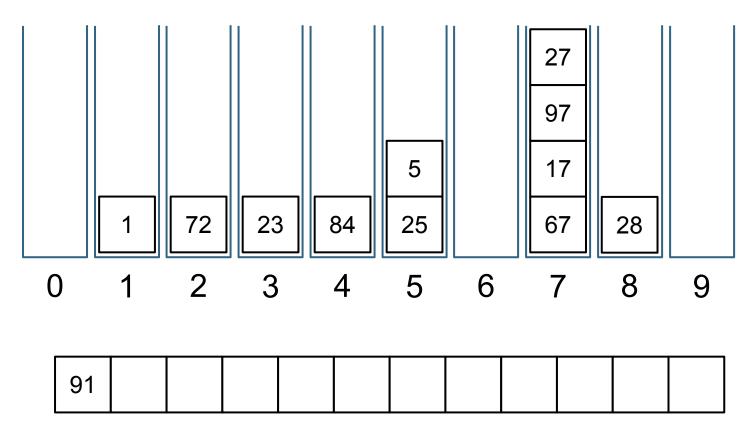
- First pass (on the 1st digit from right)
 - Put the elements to bins



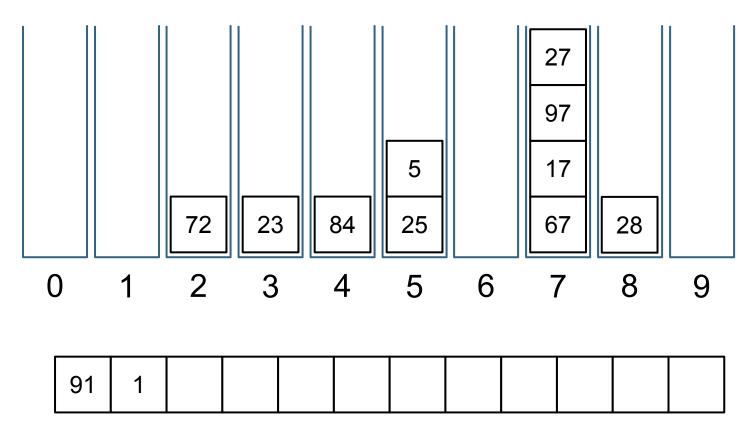
- First pass (on the 1st digit from right)
 - Put the elements to bins



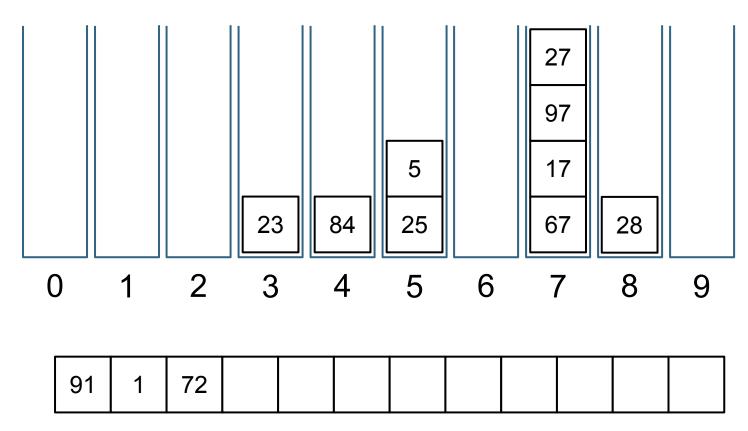
- First pass (on the 1st digit from right)
 - Put the elements back into the array



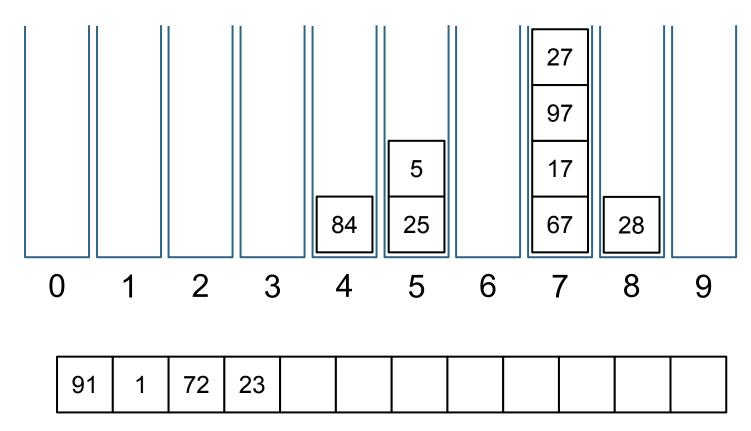
- First pass (on the 1st digit from right)
 - Put the elements back into the array



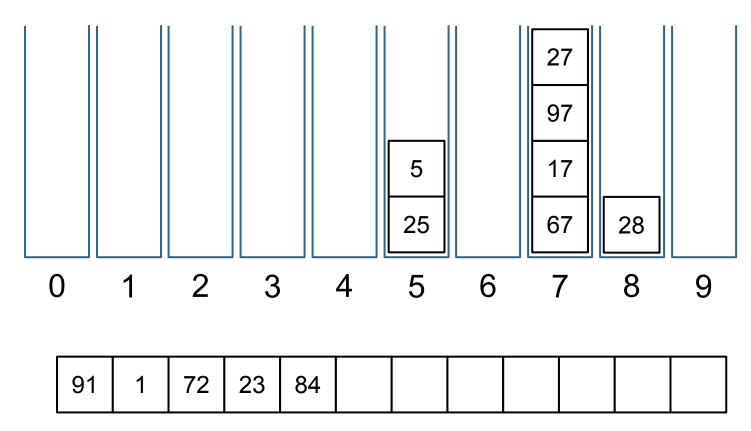
- First pass (on the 1st digit from right)
 - Put the elements back into the array



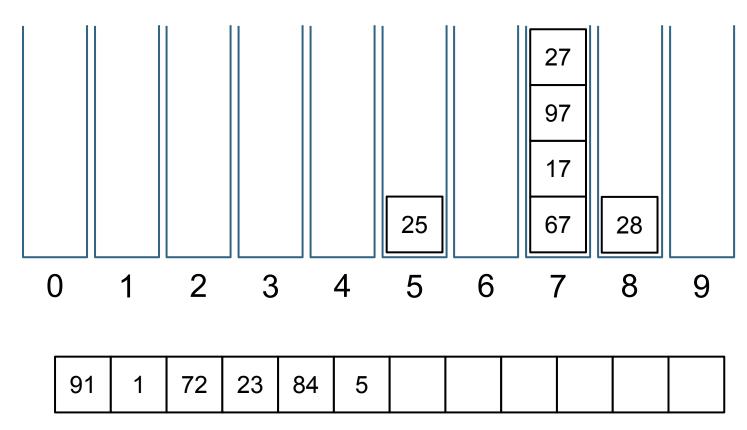
- First pass (on the 1st digit from right)
 - Put the elements back into the array



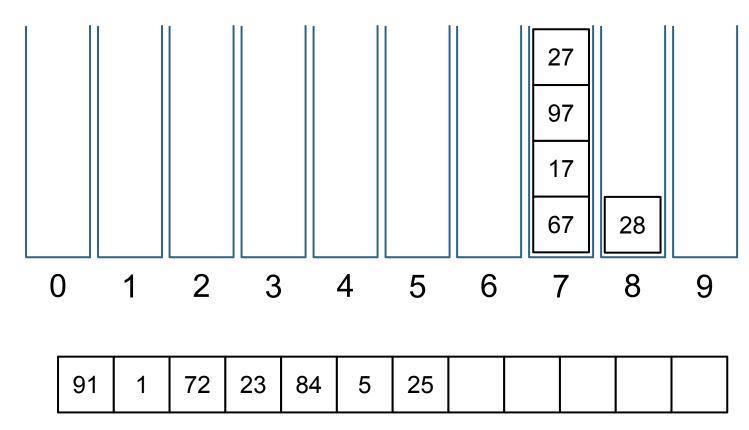
- First pass (on the 1st digit from right)
 - Put the elements back into the array



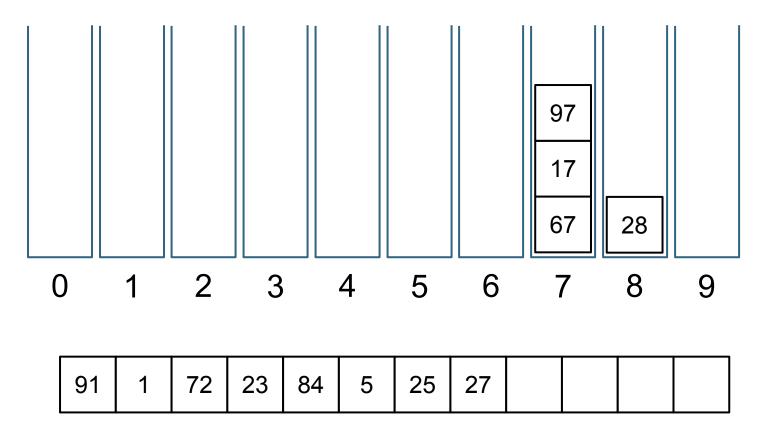
- First pass (on the 1st digit from right)
 - Put the elements back into the array



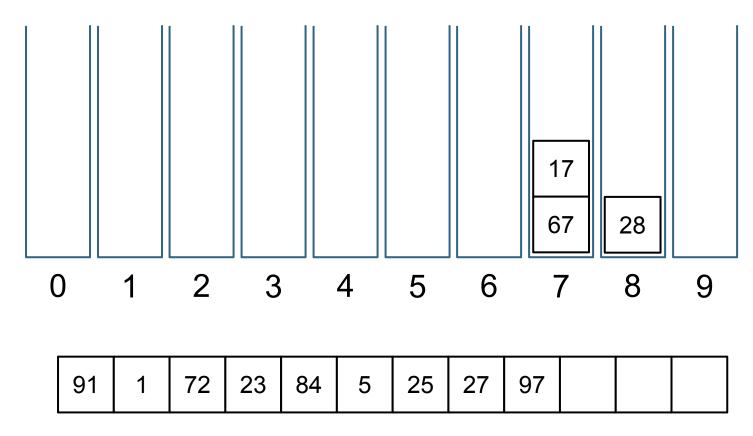
- First pass (on the 1st digit from right)
 - Put the elements back into the array



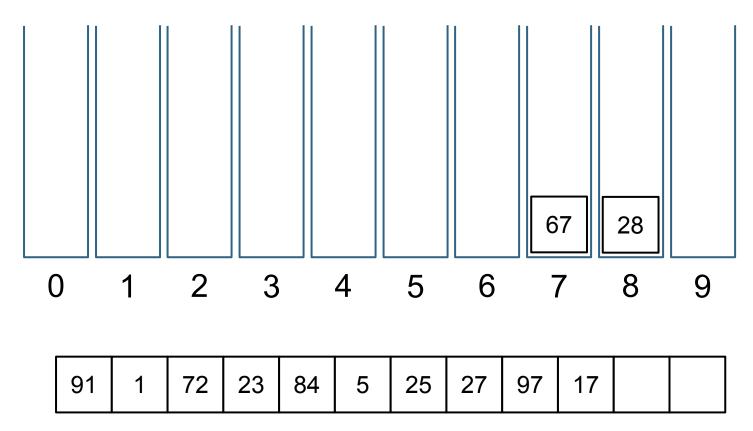
- First pass (on the 1st digit from right)
 - Put the elements back into the array



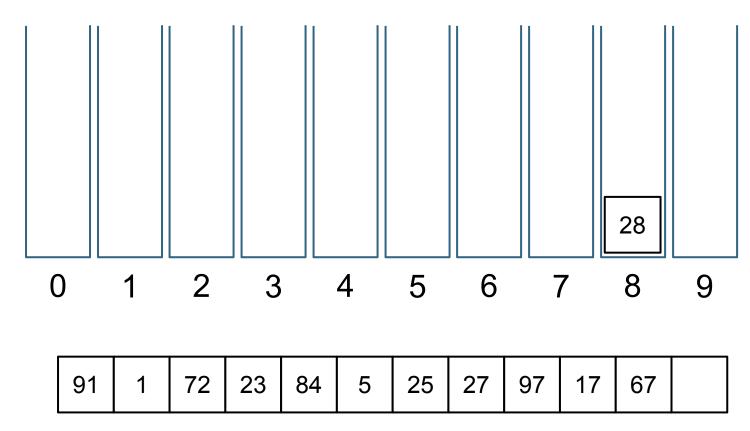
- First pass (on the 1st digit from right)
 - Put the elements back into the array



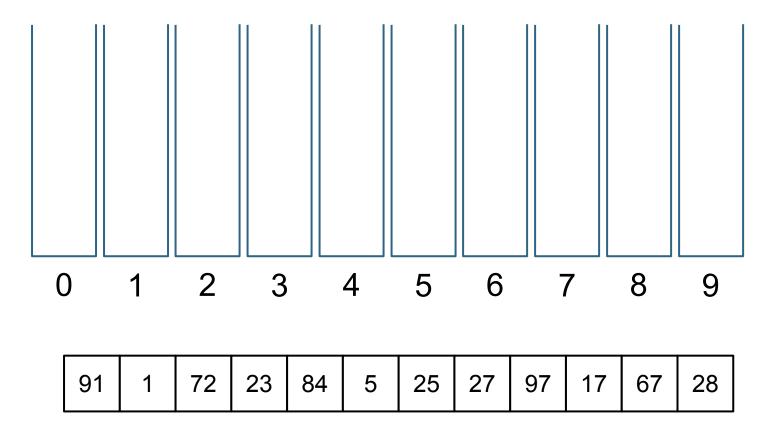
- First pass (on the 1st digit from right)
 - Put the elements back into the array



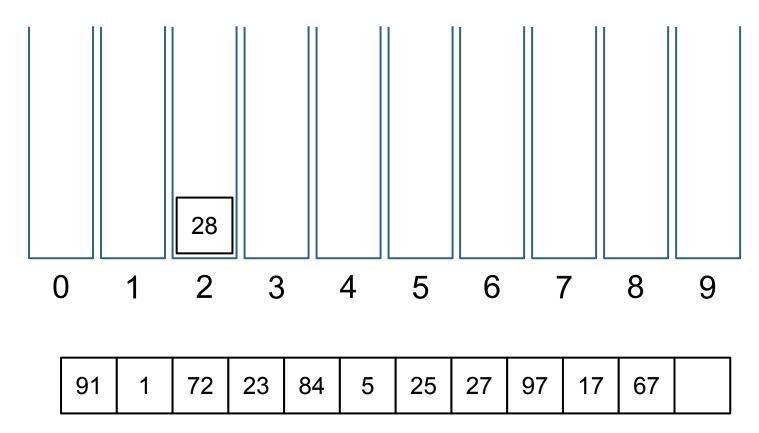
- First pass (on the 1st digit from right)
 - Put the elements back into the array



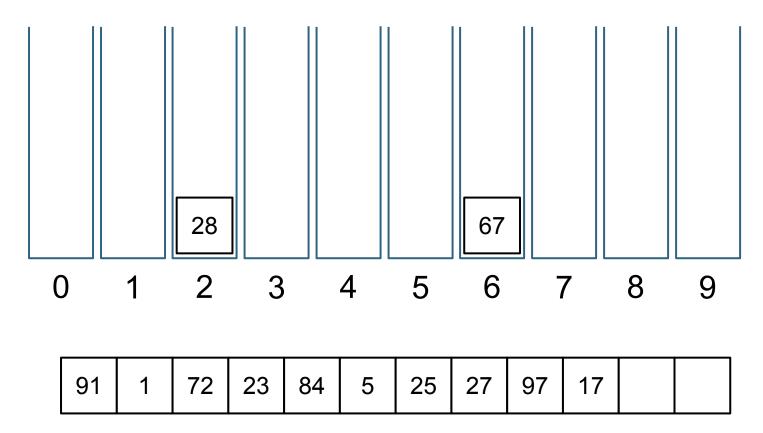
- First pass (on the 1st digit from right)
 - Put the elements back into the array



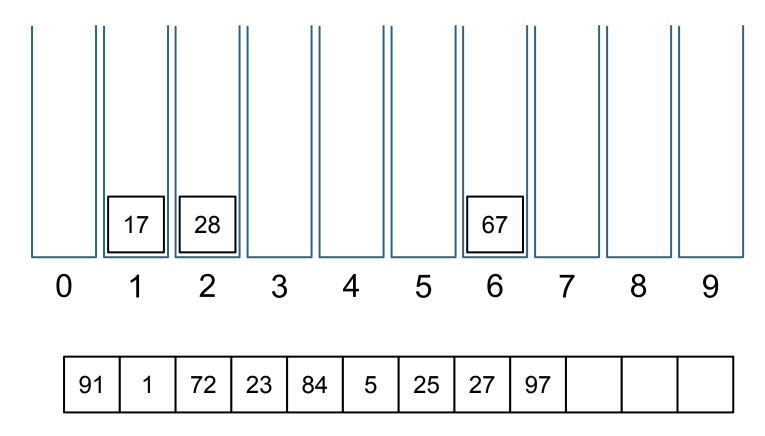
- Second pass (on the 2nd digit from right)
 - Put the elements to bins



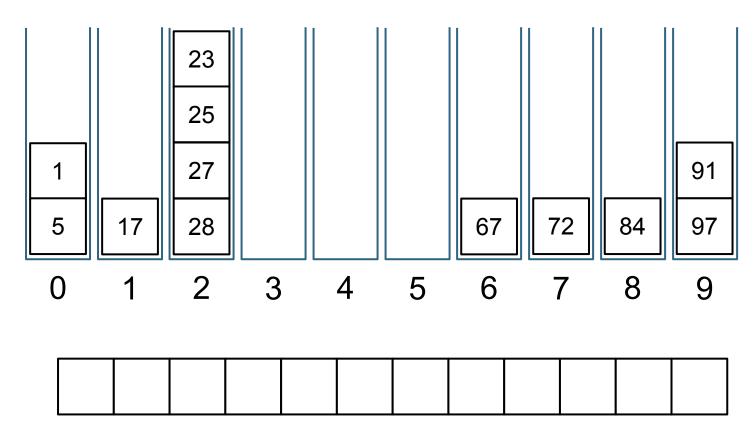
- Second pass (on the 2nd digit from right)
 - Put the elements to bins



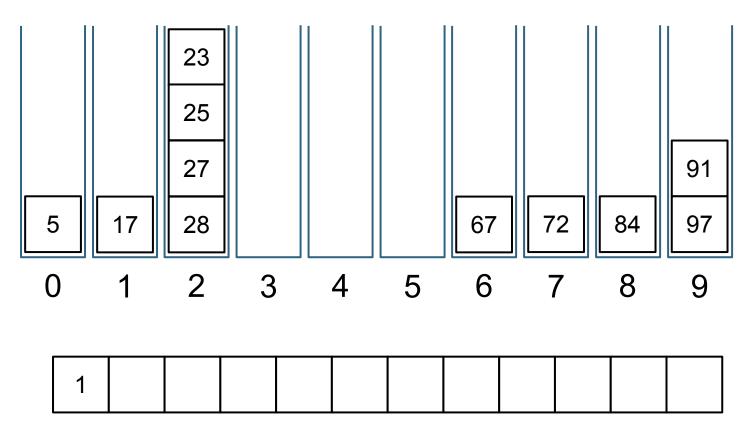
- Second pass (on the 2nd digit from right)
 - Put the elements to bins



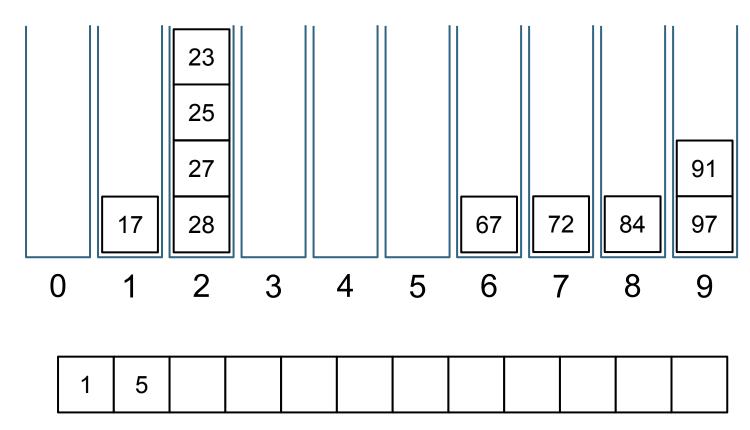
- Second pass (on the 2nd digit from right)
 - Put the elements to bins (some steps are skipped)



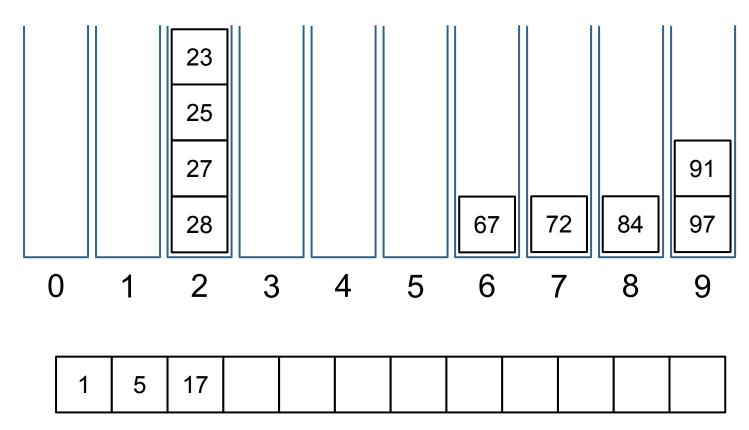
- Second pass (on the 2nd digit from right)
 - Put the elements back into the array



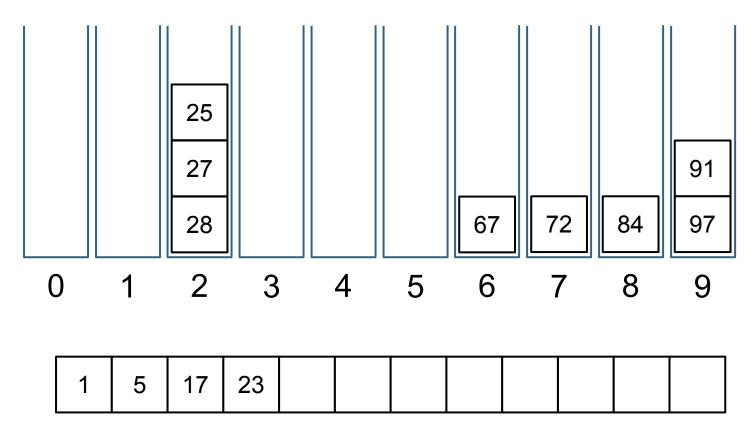
- Second pass (on the 2nd digit from right)
 - Put the elements back into the array



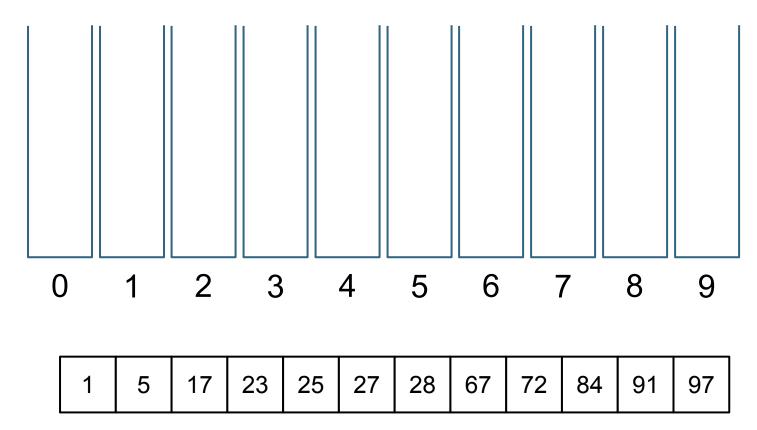
- Second pass (on the 2nd digit from right)
 - Put the elements back into the array



- Second pass (on the 2nd digit from right)
 - Put the elements back into the array

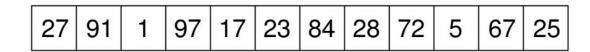


- Second pass (on the 2nd digit from right)
 - Put the elements back into the array (some steps are skipped)



- 이 방법은 배열과 (linked) stack을 필요로 함
- radix sort를 크기가 n인 배열을 이용해서 구현 가능?
 - 추가 배열을 사용해도 된다고 하면...
- Yes!
 - Main idea: 2-step algorithm on the input data
 - Step 1: 각 bin에 몇 개의 item이 들어가는지 개수를 셈
 - Step 2: step 1 에서 센 것을 사용하여 각 item을 적절한 위치에 놓음

Initial Input: Array A

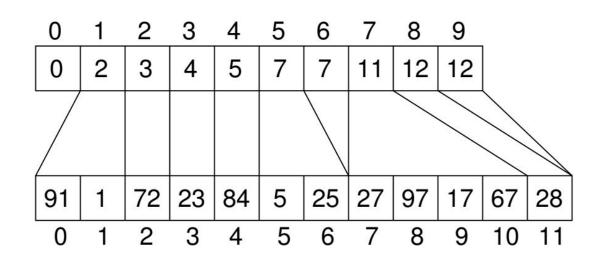


First pass values for Count. rtoi = 1.

0	1	2	3	4	5	6	7	8	9
0	2	1	1	1	2	0	4	1	0

Count array: Index positions for Array B.

End of Pass 1: Array A.



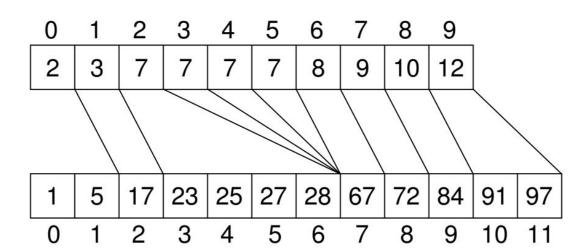
End of Pass 1: Array A.

91	1	72	23	84	5	25	27	97	17	67	28
		2									

Second pass values for Count. rtoi = 10.

0						1355			100000
2	1	4	0	0	0	1	1	1	2

Count array: Index positions for Array B.



End of Pass 2: Array A.

```
static void radix(Integer[] A, Integer[] B,
                  int k, int r, int[] count) {
   int i, j, rtok;
   for (i = 0, rtok = 1; i < k; i++, rtok *= r) {
       for (j = 0; j < r; j++) count[j] = 0;
       // Count # of elements for each bin on this pass
       for (j = 0; j < A.length; j++) count[(A[j] / rtok) % r]++;
       // count[j] is index in B for last slot of j
       for (j = 1; j < r; j++) count[j] = count[j - 1] + count[j];
       for (j = A.length - 1; j >= 0; j--)
           B[--count[(A[j] / rtok) % r]] = A[j];
       for (j = 0; j < A.length; j++) A[j] = B[j];
```

Radix Sort Cost

n: input size

k: length of each input key

r: radix

Radix Sort의 시간복잡도?

• 모든 key들이 서로 다르다면, 키 길이 k의 최소값은?

Radix Sort Cost

n: input size

k: length of each input key

r: radix

• Radix Sort의 시간복잡도? $\Theta(nk + rk)$

- 모든 key들이 서로 다르다면, 키 길이 k의 최소값은?
 - 최소 log_rn
 - 따라서, 일반적인 경우 radix sort는 0(nlogn) 알고리즘임

Overview

- Heapsort
- Binsort
- Radix Sort

Questions?