

# Introduction to VC-dimension and its applications in combinatorial problems

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2024.08.20

Def.  $(X, \mathcal{F})$  is called a set system if  $X$  is a set and  $\mathcal{F} \subseteq P(X)$ .

$$\{\mathcal{Y} : \mathcal{Y} \subseteq X\}$$

Def. Let  $(X, \mathcal{F})$  be a set system.

(1)  $A \subseteq X$  is shattered by  $\mathcal{F}$  if  $\forall A' \subseteq A, \exists F \in \mathcal{F}$  s.t.  $A \cap F = A'$ .

Equivalently,  $A \subseteq X$  is shattered by  $\mathcal{F}$  if  $\{A \cap F : F \in \mathcal{F}\} = P(A)$ .

(2)  $\mathcal{F}$  has VC-dimension  $\geq d$  if  $\exists A \subseteq X$  shattered by  $\mathcal{F}$  with  $|A| = d$ .

We write  $VC(\mathcal{F}) \geq d$ .

If  $VC(\mathcal{F}) \geq d$  but not  $VC(\mathcal{F}) \geq d+1$ , then  $VC(\mathcal{F}) = d$ .

If  $VC(\mathcal{F}) \geq d$  for all  $d \in \mathbb{N}$ , then  $VC(\mathcal{F}) = \infty$ .

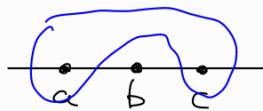
(3) For  $A \subseteq X$ ,  $A \cap \mathcal{F} := \{A \cap F : F \in \mathcal{F}\}$ .

Example.

(1) Let  $X = \mathbb{R}$ ,  $\mathcal{F} = \{(-\infty, a) : a \in \mathbb{R}\} \cup \{(b, \infty) : b \in \mathbb{R}\}$ .

Then  $VC(\mathcal{F}) = 2$ .

$$\therefore \{a, b\}$$



$\{a, b, c\}$  is not shattered.

$$(\{a, b, c\} \cap \mathcal{F} \not\Rightarrow \{a, c\}).$$

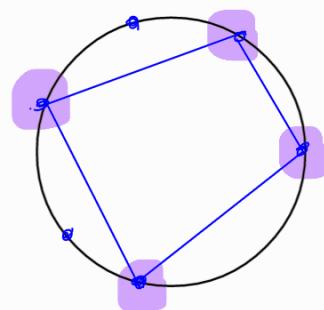
(2) Let  $X = \mathbb{R}^2$ ,  $\mathcal{F} = \{\text{convex polygons}\}$ .

Then  $VC(\mathcal{F}) = \infty$ .

$\therefore$  Let  $A \subset \{(x, y) : x^2 + y^2 = 1\}$ ,  $|A| = n$ .

$\forall A' \subseteq A$  with  $|A'| = m \leq n$ ,

$\exists$  convex  $m$ -gon  $G$  s.t.  $A \cap G = A'$ .



Def. Let  $(X, \mathcal{F})$  be a set system.

The shatter function  $\pi_{\mathcal{F}}(n) := \max \{|A \cap \mathcal{F}| : A \subseteq X \text{ and } |A| = n\}$ . ( $\leq 2^n$ )

Lemma. (Sauer-Shelah lemma)

Let  $(X, \mathcal{F})$  be a set system. Assume  $VC(\mathcal{F}) \leq d$ .

Then  $\forall n \geq d$ ,  $\pi_{\mathcal{F}}(n) \leq \sum_{i=0}^d \binom{n}{i}$ . ( $= O(n^d)$ , i.e. polynomial growth.)

### Notation.

$$\text{Av}(x_1, \dots, x_n; A) := \frac{(1_A(x_1) + \dots + 1_A(x_n))}{n} = \text{relative frequency of } A$$

Theorem. (Weak law of large numbers) (WLLN)

$\forall$  event  $F, \forall \varepsilon > 0, \forall n > 0,$

$$P\left[|\text{Av}(x_1, \dots, x_n; F) - P[F]| \geq \varepsilon\right] \leq \frac{1}{4n\varepsilon^2} \quad (\rightarrow 0 \text{ as } n \rightarrow \infty).$$

Remark.

$$P(A \cup B) \leq P(A) + P(B)$$

Let  $\mathcal{F}$  be some family of events. By WLLN and union bound,

$$\begin{aligned} & P\left[\exists \text{event } F \in \mathcal{F} \text{ s.t. } |\text{Av}(x_1, \dots, x_n; F) - P[F]| \geq \varepsilon\right] \\ &= P\left[\sup\left\{|\text{Av}(x_1, \dots, x_n; F) - P[F]| : F \in \mathcal{F}\right\} \geq \varepsilon\right] \leq \frac{1}{4n\varepsilon^2} \cdot |\mathcal{F}|. \end{aligned}$$

This "uniform version" w.r.t.  $\mathcal{F}$  has dependency on  $\mathcal{F}$ .

Theorem. (VC-theorem) (Uniform version of the WLLN)

Let  $(X, \mathcal{F})$  be a set system with finite  $X$ .

Regard  $X$  as a probability space with some  $P$ .  $\forall \varepsilon > 0,$

(e.g.  $P(A) = |A|/|X|$ . Counting probability.)

$$P\left[\sup\left\{|\text{Av}(x_1, \dots, x_n; F) - P[F]| : F \in \mathcal{F}\right\} > \varepsilon\right] \leq 8\pi_{\mathcal{F}}(n) \exp\left(-\frac{n\varepsilon^2}{32}\right).$$

$(\rightarrow 0 \text{ as } n \rightarrow \infty \text{ if } \pi_{\mathcal{F}}(n) \text{ has polynomial growth.})$

Corollary.

(no dependency on  $\mathcal{F}$ )

$\forall d \in \mathbb{N}, \forall \varepsilon > 0, \exists N = N(d, \varepsilon) \text{ s.t.}$

$\forall$  set system  $(X, \mathcal{F})$  with some  $P$  on finite  $X$ ,

$\text{VC}(\mathcal{F}) \leq d \Rightarrow \exists \underline{\varepsilon\text{-approximation}}$  for  $\mathcal{F}$  of size  $\leq N$ .

$$\left( \exists \{x_1, \dots, x_N\} \subseteq X \text{ s.t. } \forall F \in \mathcal{F}, |\text{Av}(x_1, \dots, x_N; F) - P[F]| < \varepsilon. \right)$$

(possibly with repetitions)

Proof. By VC-theorem,

$$\mathbb{P} \left[ \sup \left\{ |A_v(x_1, \dots, x_n; F) - \mathbb{P}[F]| : F \in \mathcal{F} \right\} > \varepsilon \right] \leq 8\pi_{\mathcal{F}}(n) \exp \left( -\frac{n\varepsilon^2}{32} \right).$$

$$\Rightarrow \mathbb{P} \left[ \forall F \in \mathcal{F}, |A_v(x_1, \dots, x_n; F) - \mathbb{P}[F]| \leq \varepsilon \right] \geq 1 - 8\pi_{\mathcal{F}}(n) \exp \left( -\frac{n\varepsilon^2}{32} \right).$$

By SS lemma,  $8\pi_{\mathcal{F}}(n) \exp \left( -\frac{n\varepsilon^2}{32} \right) \rightarrow 0$  as  $n \rightarrow \infty$ .

In particular, with sufficiently large  $N$  depending only on  $d$  and  $\varepsilon$ ,

$$\mathbb{P} \left[ \forall F \in \mathcal{F}, |A_v(x_1, \dots, x_N; F) - \mathbb{P}[F]| < \varepsilon \right] \neq 0.$$

Thus  $\exists$  at least one  $N$ -tuple  $\{x_1, \dots, x_N\}$  s.t.

$$\forall F \in \mathcal{F}, |A_v(x_1, \dots, x_N; F) - \mathbb{P}[F]| < \varepsilon. \blacksquare$$

## Application to bounding transversal numbers

Def.

Let  $(X, \mathcal{F})$  be a set system.

(1)  $T \subseteq X$  is a transversal of  $\mathcal{F}$  if  $\forall F \in \mathcal{F}, T \cap F \neq \emptyset$ .

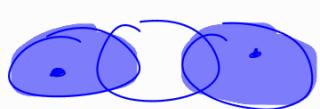
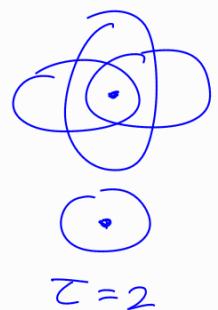
The transversal number of  $\mathcal{F}$ ,

$\tau(\mathcal{F})$  is the smallest size of a transversal.

(2)  $G \subseteq \mathcal{F}$  is a packing if  $\forall F_1 \neq F_2 \in G, F_1 \cap F_2 = \emptyset$ .

The packing number of  $\mathcal{F}$ ,

$\nu(\mathcal{F})$  is the largest size of a packing.



Remark.

(1)  $\nu(\mathcal{F}) \leq \tau(\mathcal{F})$ .

( $\because$  Let  $G \subseteq \mathcal{F}$  be a packing with the largest size.

Then  $\forall F \in G$ , any transversal  $T \subseteq X$  should contain one element from  $F$ .)

(2) It is difficult to calculate the actual transversal/packing number.  
(integer programming is NP-hard.)

(3) Very little is known about the reverse direction:

$\tau \leq f(z)$  for some fct  $f$  under "reasonable" conditions?

(4) If  $H$  is a hypergraph, then we have a set system

$$X = V(H), \mathcal{F} = E(H) \subseteq P(V(H)),$$

and a transversal of  $\mathcal{F}$  is called a vertex cover of  $H$ ,  $\tau(H)$ ,  
a packing of  $\mathcal{F}$  is called a matching of  $H$ ,  $z(H)$ .

(5) If  $H$  is an  $r$ -uniform hypergraph, i.e.  $\forall e \in E(H), |e|=r$ , then  
 $\tau(H) \leq r \cdot z(H)$ .

( $\because$  The union of edges from a maximal matching is a vertex cover.)



Ryser's Conjecture (1971)

If  $H$  is an  $r$ -uniform  $r$ -partite hypergraph, then  $\tau(H) \leq (r-1) \cdot z(H)$ .

(Solved only for  $r=2$ : König's theorem (1931).)

and  $r=3$  by Ron Aharoni (2001)

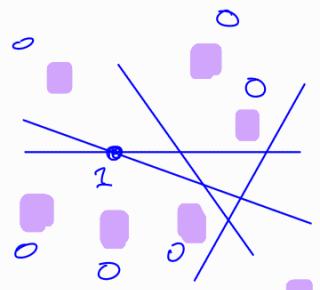
Example.

Let  $X = \mathbb{R}^2$  and  $\mathcal{F}$  be a set of  $n$  lines  $l_1, \dots, l_n$  s.t.

$\forall 1 < i < j, |l_i \cap l_j| = 1$  but  $|l_i \cap l_j \cap l_k| = \emptyset$ .

Then  $z(\mathcal{F}) = 1$  ( $\because$  every two lines intersect),

but  $\tau(\mathcal{F}) \geq \frac{n}{2}$  ( $\because$  any point is contained in at most 2 lines).



Def.

(1) Let  $(X, \mathcal{F})$  be a set system with  $X$  finite.

A function  $\phi: X \rightarrow [0, 1]$  is a fractional transversal for  $\mathcal{F}$

if  $\forall F \in \mathcal{F}, \sum_{x \in F} \phi(x) \geq 1$ .

The size of  $\phi$  is  $\sum_{x \in X} \phi(x)$ .

The fractional transversal number  $\tau^*(\mathcal{F})$  is

the infimum of the sizes of fractional transversals for  $\mathcal{F}$ .

(2) Let  $(X, \mathcal{F})$  be a set system with  $X$  finite.

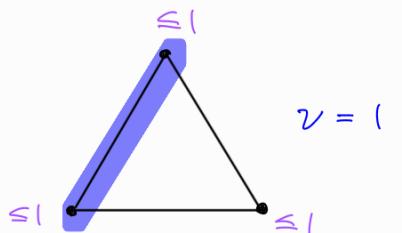
A function  $\psi: \mathcal{F} \rightarrow [0, 1]$  is a fractional packing for  $\mathcal{F}$  if  $\forall x \in X, \sum_{\{F \in \mathcal{F}: x \in F\}} \psi(F) \leq 1$ .

The size of  $\psi$  is  $\sum_{F \in \mathcal{F}} \psi(F)$ .

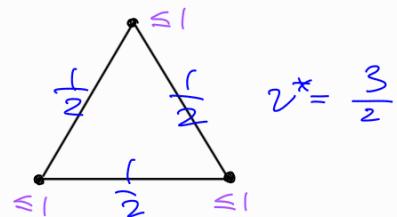
The fractional packing number  $\nu^*(\mathcal{F})$  is the supremum of the sizes of fractional packings for  $\mathcal{F}$ .

Example.

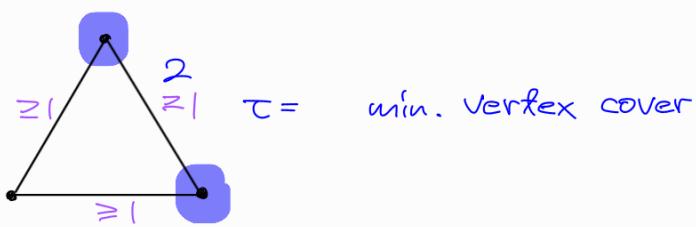
Triangle.  $X = \{a_1, a_2, a_3\}, \mathcal{F} = \{\{a_1, a_2\}, \{a_2, a_3\}, \{a_3, a_1\}\}$ .



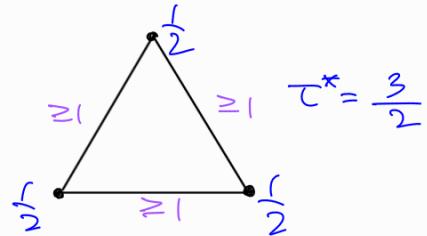
$$\nu = 1 \text{ matching number}$$



$$\nu^* = \frac{3}{2}$$



$$\tau = \text{min. vertex cover}$$



$$\tau^* = \frac{3}{2}$$

Remark.

$$\nu^* \geq \nu \leq \tau \geq \tau^*$$

max min

$$\begin{array}{c} \tau \\ \diagdown \quad \diagup \\ \nu^* \quad \nu \end{array}$$

Remark.

(1)  $X = [2n], \mathcal{F} = \binom{X}{n}$ . Then  $\tau = n+1$  and  $\tau^* = 2$ .

( $\because$  If  $n$  points  $a_1, \dots, a_n$  are chosen, then  $X \setminus \{a_1, \dots, a_n\} \in \binom{X}{n}$  and  $\{a_1, \dots, a_n\} \cap (X \setminus \{a_1, \dots, a_n\}) = \emptyset \Rightarrow \tau > n$ .

Assign  $\phi(x) = \frac{1}{n}$  for all  $x \in X \Rightarrow \tau^* \leq 2$ .

$$\sum_{x \in X} \phi(x) = \sum_{x \in F} \phi(x) + \sum_{x \in X \setminus F} \phi(x) \geq ( + 1 = 2)$$

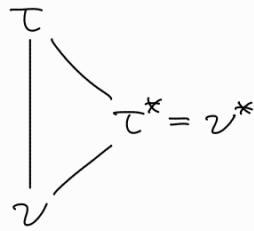
(2)  $X = \binom{[n]}{2}, \mathcal{F} = \left\{ \{1, 2\}, \{1, 3\}, \{1, 4\}, \dots, \{2, 3\}, \{2, 4\}, \dots \right\}$

Then  $\nu = 1, \nu^* = \frac{n}{2}$ .

$\mathcal{F}$ : finite

Fact.

$$(1) \mathcal{F} : \text{finite} \Rightarrow \tau^*(\mathcal{F}) = \nu^*(\mathcal{F}).$$



(2) It is easy to calculate or bound the fractional transversal number (linear programming, P.)

Thus it is very useful to bound  $\tau$  in terms of  $\tau^*$ .

Def.

Let  $(X, \mathcal{F})$  be a set system with some  $P$  on  $X$ .

$A \subseteq X$  is called an  $\varepsilon$ -net if  $\forall F \in \mathcal{F}$  s.t.  $P[F] \geq \varepsilon$ ,  $A \cap F \neq \emptyset$ .

Remark.

Note that every  $\varepsilon$ -approximation is an  $\varepsilon$ -net.

( $\because$  Let  $\{x_1, \dots, x_N\}$  be an  $\varepsilon$ -approximation.

$$\forall F \in \mathcal{F}, |Av(x_1, \dots, x_N; F) - P[F]| < \varepsilon$$

$\Rightarrow$  If  $P[F] \geq \varepsilon$  but  $\{x_1, \dots, x_N\} \cap F = \emptyset$ , then  $Av(x_1, \dots, x_N; F) = 0$

$$\text{and } |Av(x_1, \dots, x_N; F) - P[F]| \geq \varepsilon, \text{ w.r.t.}$$

Fact.

$$VC(\mathcal{F}) \leq d \Rightarrow \exists \varepsilon\text{-net for } \mathcal{F} \text{ of size } \leq Cd \frac{1}{\varepsilon} \log \frac{1}{\varepsilon}.$$

$\mathcal{F}$ : finite,  $VC(\mathcal{F}) \leq d$

$O(d\tau^* \log \tau^*)$

Proposition.

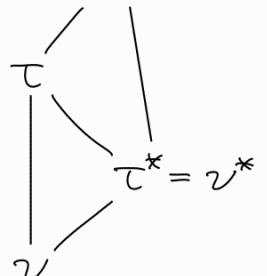
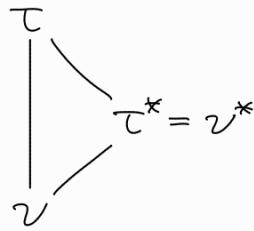
Let  $(X, \mathcal{F})$  be a set system with finite  $X$  and  $VC(\mathcal{F}) \leq d$ .

Then  $\tau(\mathcal{F}) \leq Cd\tau^*(\mathcal{F}) \log \tau^*(\mathcal{F})$ , where  $C$ : constant.

Proof.

Let  $\phi: X \rightarrow [0, 1]$  be an optimal fractional transversal, so that

$$\tau^*(\mathcal{F}) = \sum_{x \in X} \phi(x).$$



A function  $P$  defined by  $P(\{x\}) = \frac{\phi(x)}{\tau^*(f)}$  for all  $x \in X$  defines a probability on  $X$ .

$$\text{Note that } \forall F \in \mathcal{F}, P(F) = \frac{\sum_{x \in F} \phi(x)}{\tau^*(f)} \geq \frac{1}{\tau^*(f)}.$$

Let  $A \subseteq X$  be a  $\frac{1}{\tau^*(f)}$ -net for  $f$  of size  $\leq C d \tau^*(f) \log \tau^*(f)$ .

Then  $\forall F \in \mathcal{F}$ ,  $A \cap F \neq \emptyset$ , i.e.,  $A$  is a transversal for  $\mathcal{F}$  of desired size.  $\blacksquare$

### Def.

Let  $(X, \mathcal{F})$  be a set system.

$(X^*, \mathcal{F}^*)$  is called the dual set system, where

$X^* = \mathcal{F}$  and  $\mathcal{F}^* = \{\mathcal{F}_x : x \in X\}$  with  $\mathcal{F}_x := \{F \in \mathcal{F} : x \in F\}$ .

The dual VC-dimension of  $\mathcal{F}$ ,  $VC^*(\mathcal{F})$  is  $VC(\mathcal{F}^*)$ .

### Remark.

(1) Given  $(X, \mathcal{F})$ , the incidence matrix  $M$  is defined as an

$$|X| \times |\mathcal{F}| - \text{matrix s.t. } M_{x,F} = \begin{cases} 1 & \text{if } x \in F, \\ 0 & \text{if } x \notin F. \end{cases}$$

(2) If the incidence matrix of  $(X', \mathcal{F}')$  is the same with  $M$  of  $(X, \mathcal{F})$ ,  
(up to permutations of rows/columns),

then  $(X, \mathcal{F})$  and  $(X', \mathcal{F}')$  are "essentially" the same.

(3) Given  $(X, \mathcal{F})$  with incidence matrix  $M$ ,

the incidence matrix of  $(X^*, \mathcal{F}^*)$  is  $M^T$ .

( $\because x \in F \iff F \in \mathcal{F}_x$ .)

### Example.

Let  $X = \mathbb{R}$ ,  $\mathcal{F} = \{(-\infty, a) : a \in \mathbb{R}\} \cup \{(b, \infty) : b \in \mathbb{R}\}$ .

Then  $VC(\mathcal{F}) = 2$ .

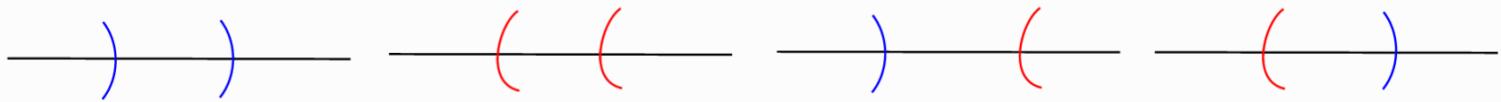


$$X^* = \{(-\infty, a) : a \in \mathbb{R}\} \cup \{(b, \infty) : b \in \mathbb{R}\},$$

$$\mathcal{F}^* = \left\{ \{(-\infty, a) : c \in (-\infty, a)\} \cup \{(b, \infty) : c \in (b, \infty)\} : c \in \mathbb{R} \right\}.$$

$$\exists \mathcal{F}_1 = \left\{ \dots, (-\infty, 0), \dots, (-\infty, -\sqrt{2}), \dots, (-\infty, 1.01), \dots, (-10^{10}, \infty), \dots, (0.9999, \infty), \dots \right\}$$

$$VC^*(\mathcal{F}) = 1.$$



Proposition.

Let  $(X, \mathcal{F})$  be a set system.

$$\text{Then } VC^*(\mathcal{F}) < 2^{VC(\mathcal{F})+1} \text{ and } VC(\mathcal{F}) < 2^{VC^*(\mathcal{F})+1}.$$

(In particular,  $VC(\mathcal{F}) < \infty$  iff  $VC^*(\mathcal{F}) < \infty$ .)

Proof.

Assume  $VC(\mathcal{F}) \geq 2^n$ . Then  $\exists A \subseteq X$  shattered by  $\mathcal{F}$  and  $|A| = 2^n$ .

Write  $A = \{\alpha_C : C \subseteq [n]\}$ . possible

$\forall k \in [n]$ , let  $F_k \in \mathcal{F} = X^*$  s.t.  $A \cap F_k = \{\alpha_C : k \in C \subseteq [n]\}$ .

Then  $\{F_1, \dots, F_n\} \subseteq X^*$  is shattered by  $\mathcal{F}^*$ .

( $\because \forall 1 \leq i_1 < \dots < i_d \leq n, \{F_{i_1}, \dots, F_{i_d}\} \cap \underbrace{\mathcal{F}_{\alpha_{\{i_1, \dots, i_d\}}}}_{\text{collection of } F\text{'s s.t. } F \ni \alpha_{\{i_1, \dots, i_d\}}} = \{F_{i_1}, \dots, F_{i_d}\}.$ ) ■

collection of  $F$ 's s.t.  $F \ni \alpha_{\{i_1, \dots, i_d\}}$

Fact. (Helly's theorem)

Let  $X = \mathbb{R}^d$ ,  $\mathcal{F}$  be some finite family of convex subsets of  $\mathbb{R}^d$ .

Assume that  $\forall \mathcal{F}' \subseteq \mathcal{F}$  s.t.  $|\mathcal{F}'| = d+1, \cap \mathcal{F}' \neq \emptyset$ .

Then  $\cap \mathcal{F} \neq \emptyset$  ( $\Leftrightarrow \tau(\mathcal{F}) = 1$ ).

Def.

Let  $(X, \mathcal{F})$  be a set system.  $\mathcal{F}$  has fractional Helly number  $k$  if

$\forall \alpha > 0, \exists \beta > 0$  s.t.  $\forall F_1, \dots, F_n \in \mathcal{F}$ ,

if at least  $\alpha$ -fraction of elements  $I \in \binom{[n]}{k}$  satisfy  $\bigcap_{i \in I} F_i \neq \emptyset$ ,

then  $\exists J \subseteq [n]$  s.t.  $|J| \geq \beta n$  and  $\bigcap_{s \in J} F_s \neq \emptyset$ .

### Remark.

- (1) Helly's theorem does not hold if convexity is replaced with  $\text{VC}(\mathcal{F}) \leq 2$ .  
 (2) If  $X = \mathbb{R}^d$ ,  $\mathcal{F}$  is some family of convex subsets of  $\mathbb{R}^d$ ,  
 then  $\mathcal{F}$  has fractional Helly number  $d+1$  (Katchalski, Liu).

### Theorem (Matousek)

Let  $(X, \mathcal{F})$  be a set system with  $\text{VC}^*(\mathcal{F}) \leq k-1$ .

Then  $\mathcal{F}$  has fractional Helly number  $k$ .

### Def.

The VC-density of  $\mathcal{F}$ ,  $\text{vc}(\mathcal{F}) = \limsup_{n \rightarrow \infty} \frac{\log(\pi_{\mathcal{F}}(n))}{\log n}$ .

That is, infimum of  $r$ 's s.t.  $\pi_{\mathcal{F}}(n) = O(n^r)$ .

(By SS Lemma,  $\text{vc}(\mathcal{F}) \leq \text{VC}(\mathcal{F})$ .)

VC-codensity  $\text{vc}^*(\mathcal{F}) = \text{vc}(\mathcal{F}^*)$ .

### Def.

Let  $p \geq q$  be natural numbers and  $(X, \mathcal{F})$  be a set system.

$(X, \mathcal{F})$  satisfies  $(p, q)$ -property if

$\forall \mathcal{F}_0 \subseteq \mathcal{F}$  s.t.  $|\mathcal{F}_0| = p$ ,  $\exists \mathcal{F}_1 \subseteq \mathcal{F}_0$  s.t.  $|\mathcal{F}_1| = q$  and  $\cap \mathcal{F}_1 \neq \emptyset$ .

### Remark.

Helly's theorem is equivalent to

$(X = \mathbb{R}^d, \mathcal{F}: \text{finite family of convex subsets satisfies } (d+1, d+1)\text{-property})$   
 $\Rightarrow \cap \mathcal{F} \neq \emptyset (\Leftrightarrow \tau(\mathcal{F}) = 1)$ .

### Fact. (Alon, Kleitman)

Let  $p \geq q \geq d+1$  be natural numbers. Then  $\exists N$  s.t.

$\mathcal{F}: \text{finite family of convex subsets of } \mathbb{R}^d \text{ satisfying } (p, q)\text{-property}$   
 $\Rightarrow \tau(\mathcal{F}) \leq N$ .

Theorem. (Alon, Kleitman + Matousek)

Let  $p \geq q \geq d+1$  be natural numbers,  $\mathcal{F}$  satisfy  $(p, q)$ -property.

Then  $\exists N = N(p, q, d)$  s.t.  $(\mathcal{F} : \text{finite and } \text{vc}^*(\mathcal{F}) \leq d \Rightarrow \tau(\mathcal{F}) \leq N)$ .

Main references:

- Artem Chernikov, Model theory and Combinatorics (draft)
- Pierre Simon, A Guide to NIP Theories