

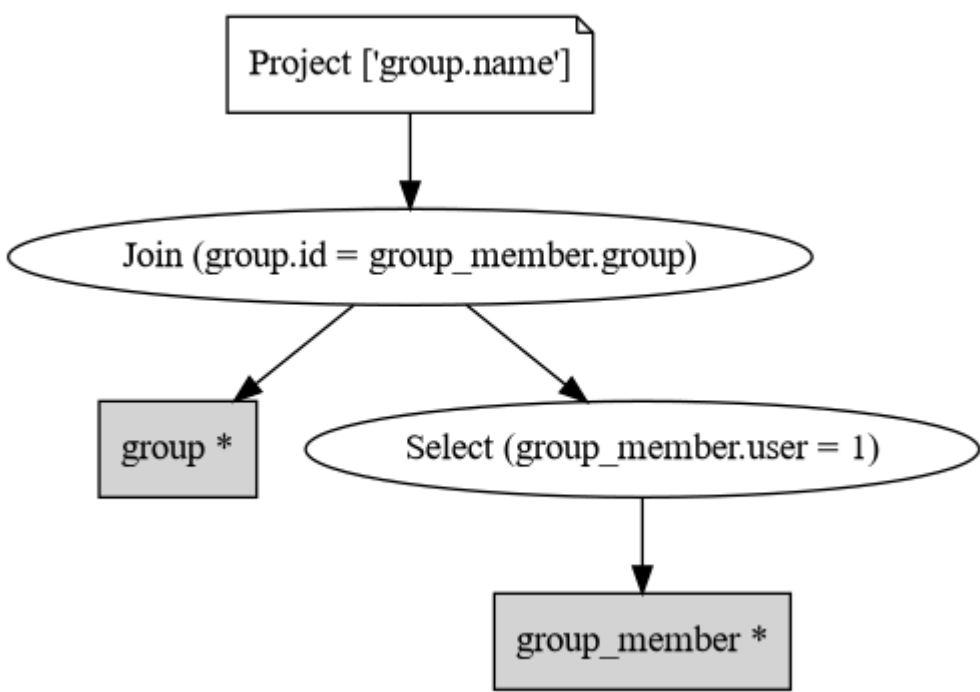
# Queries related to the application

## Get all group names the user is a part of

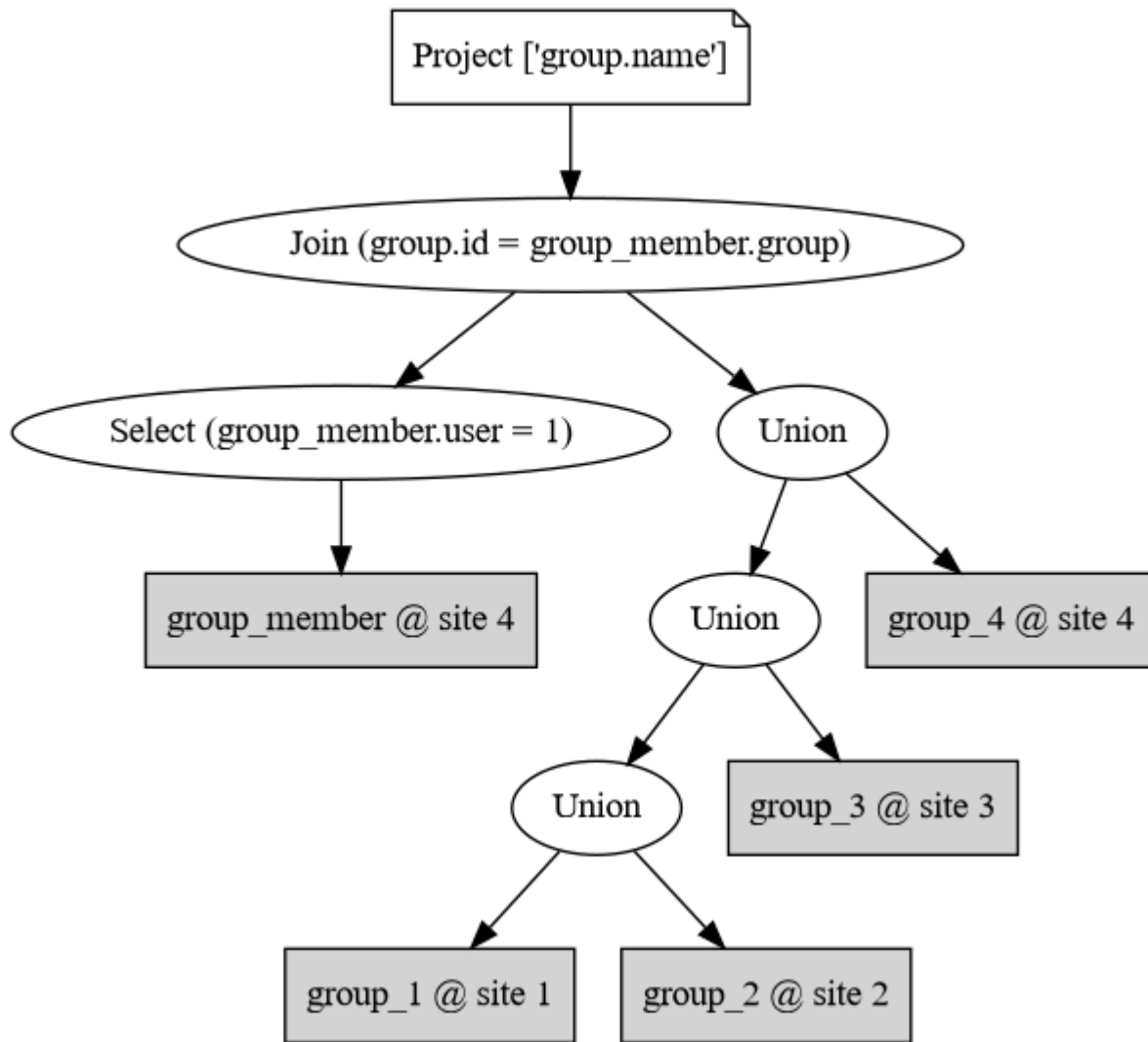
Given user id \$ID (= 1),

```
select G.`name`
  from `group` G, `group_member` GM
where
  GM.`user` = $ID
  and G.`id` = GM.`group`;
```

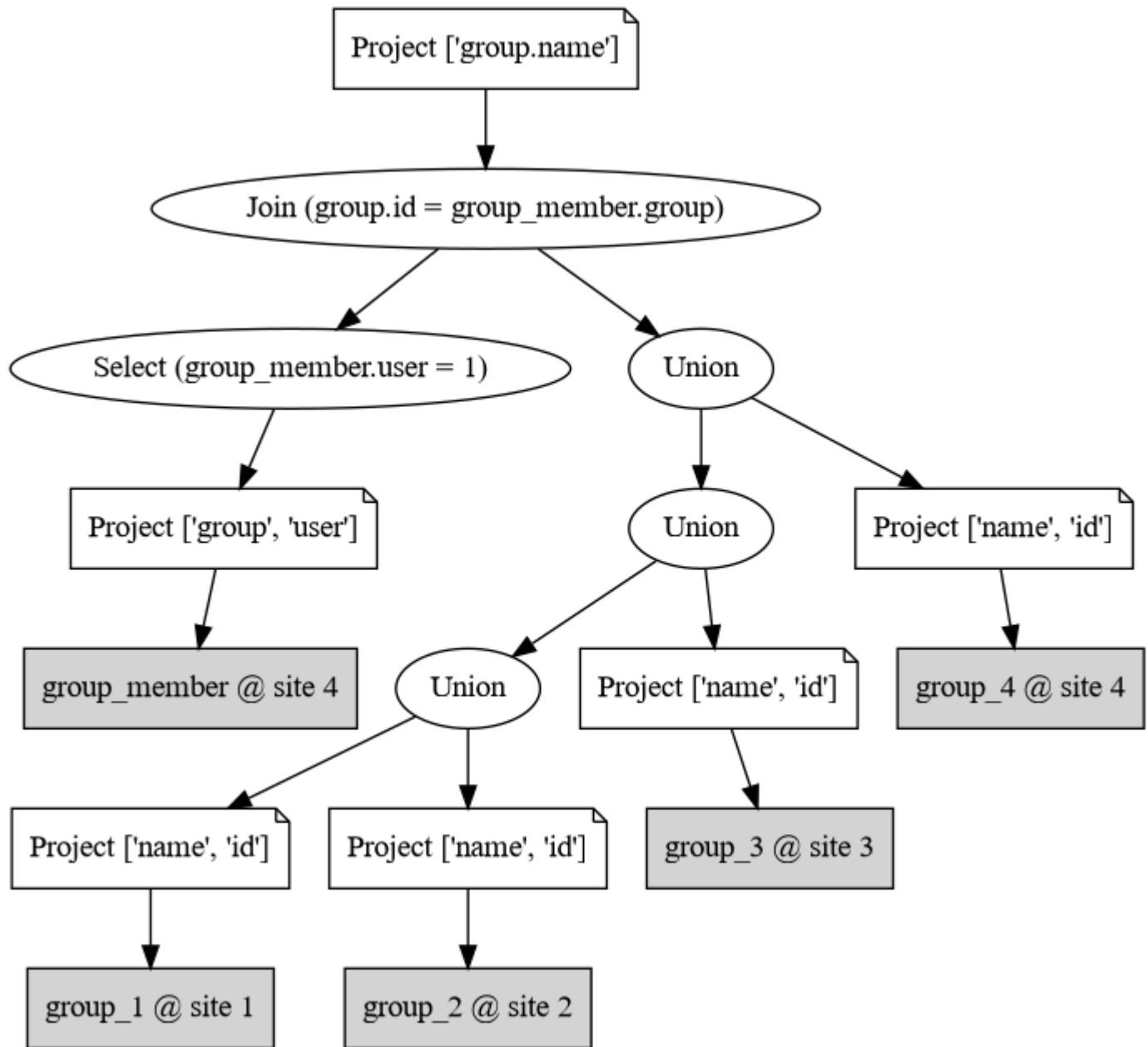
### Initial Query Tree



### Localized Query Tree



## Optimized Query Tree



## Output

```

g12
g13
g123

```

## Get all groups the user is an admin of

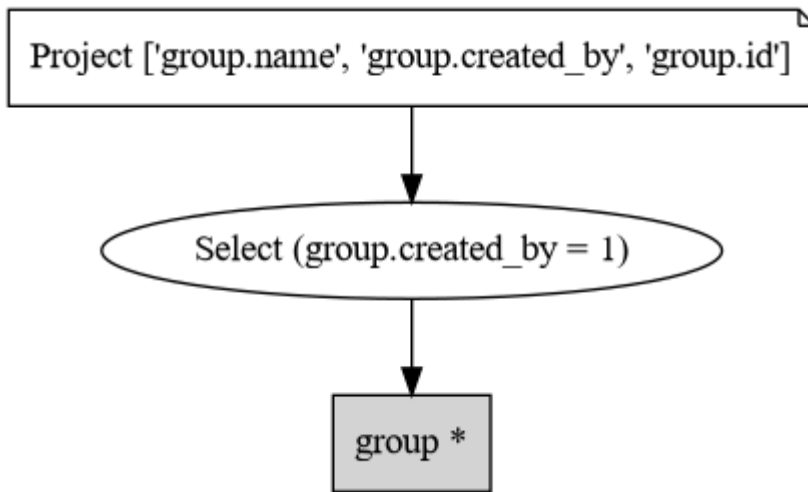
Given user id \$ID (= 1),

```

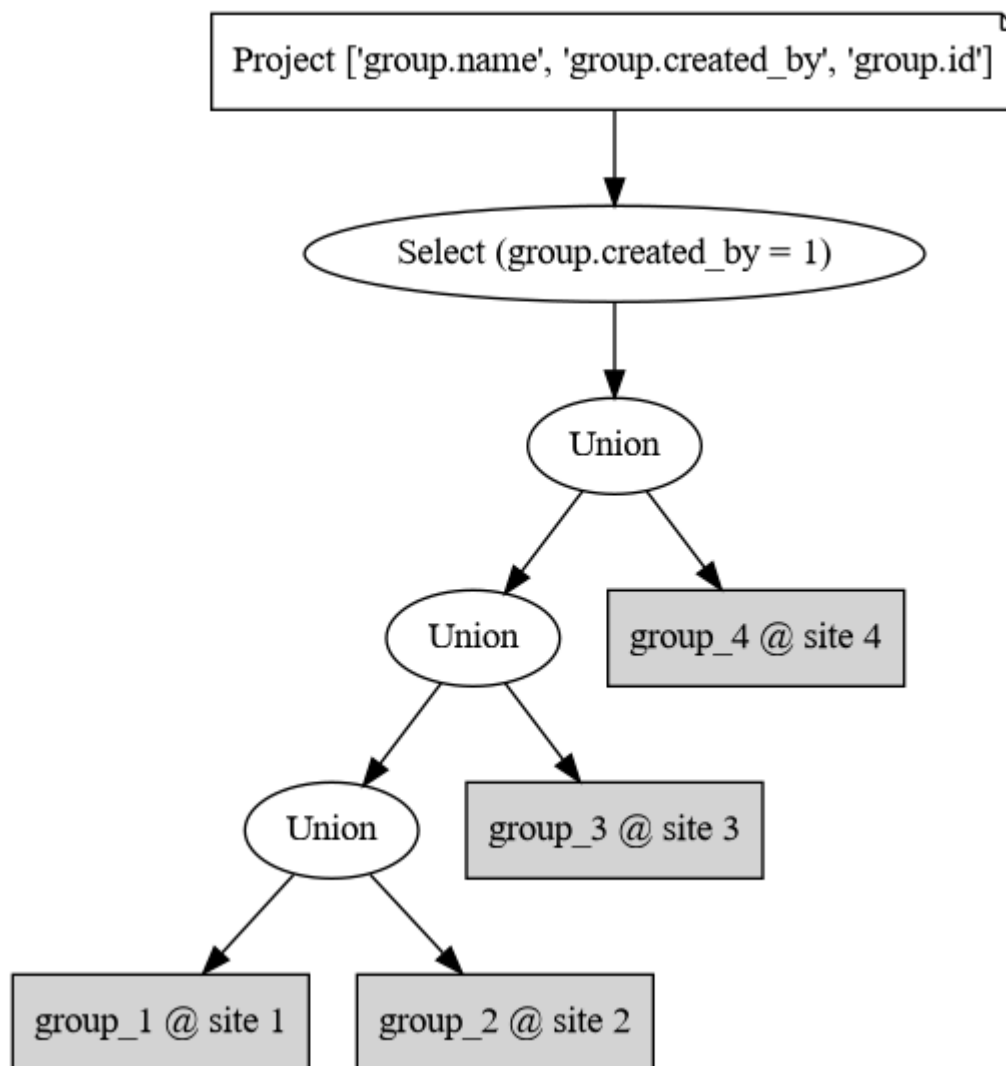
select * from `group` where `created_by` = $ID;

```

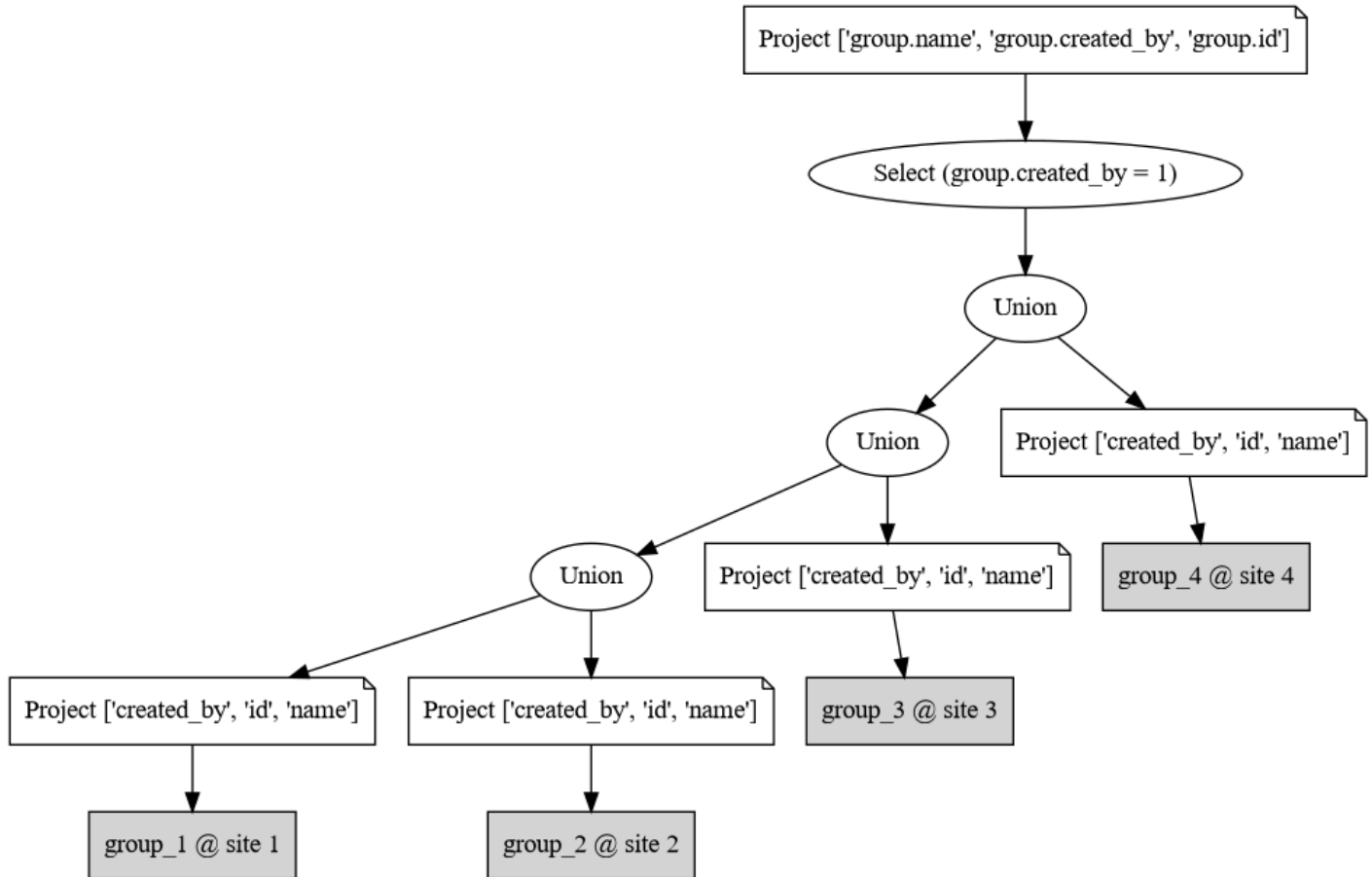
## Initial Query Tree



## Localized Query Tree



## Optimized Query Tree



## Output

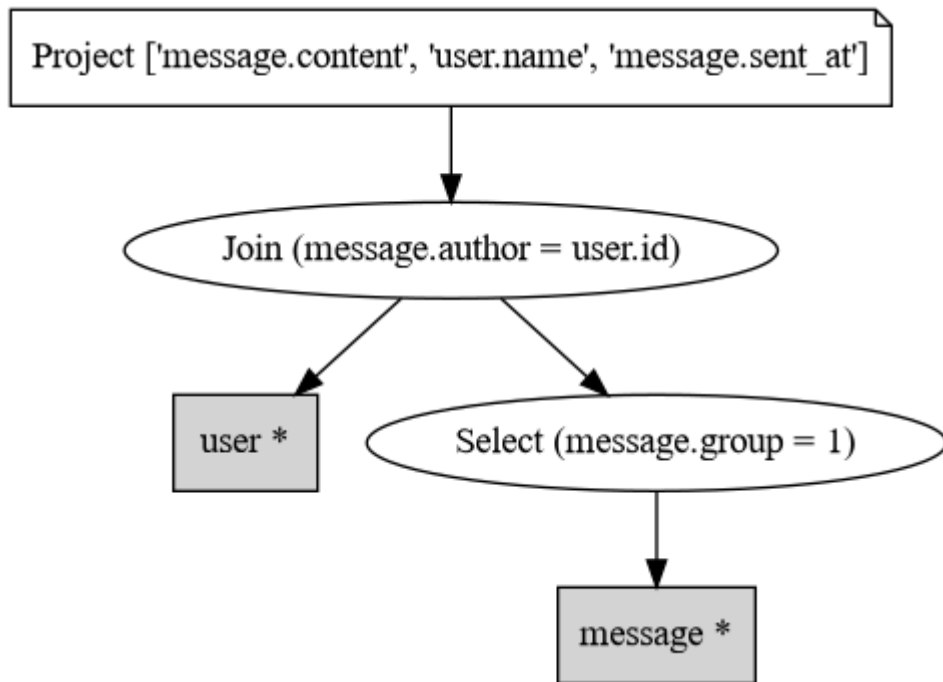
```
1,g12,1
2,g13,1
```

## Get author name, time and content of messages of a group

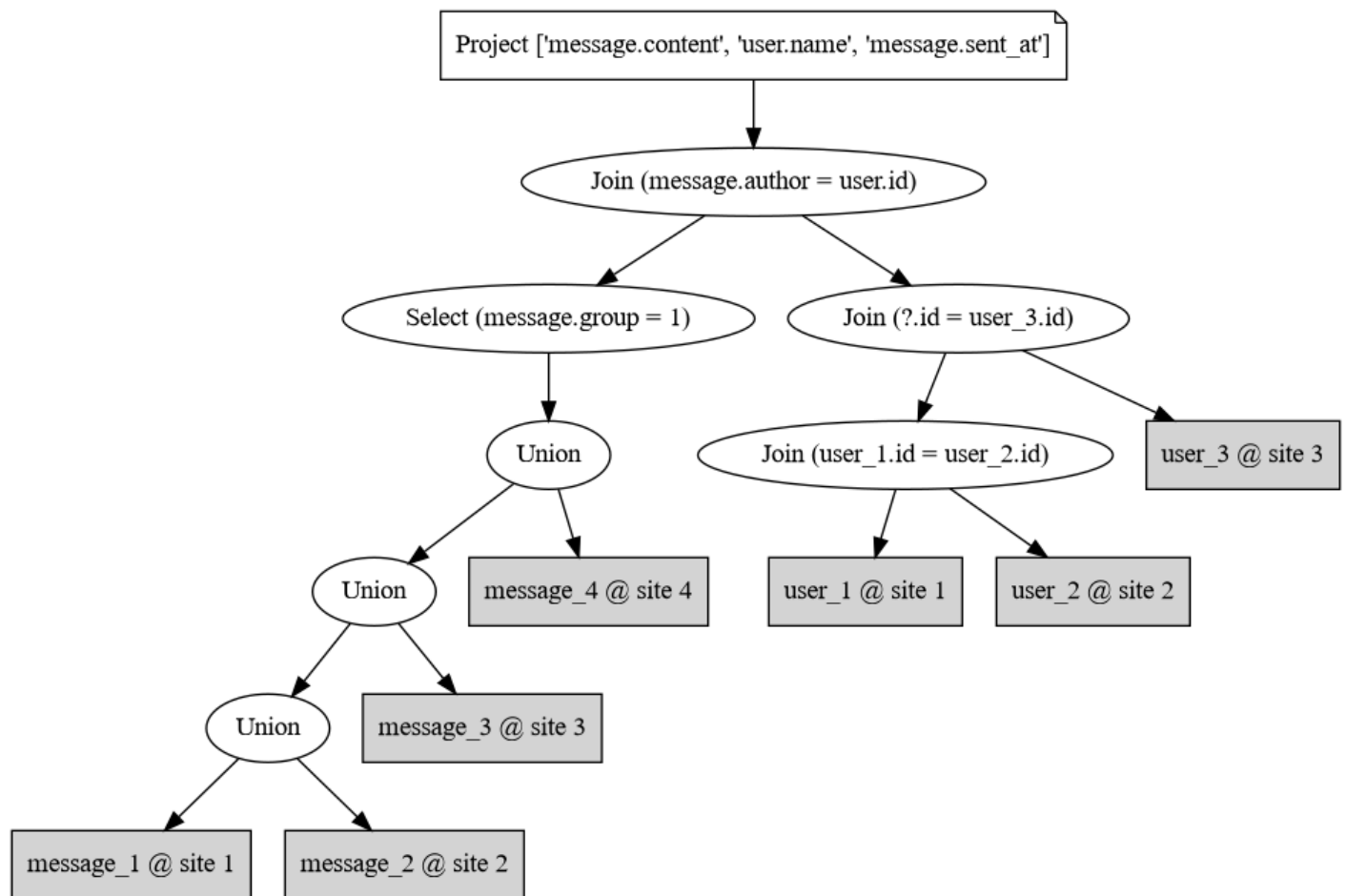
Given group id \$ID (= 1),

```
select U.`name`, M.`sent_at`, M.`content`
  from `message` M, `user` U
 where
   M.`group` = $ID and
   M.`author` = U.id;
```

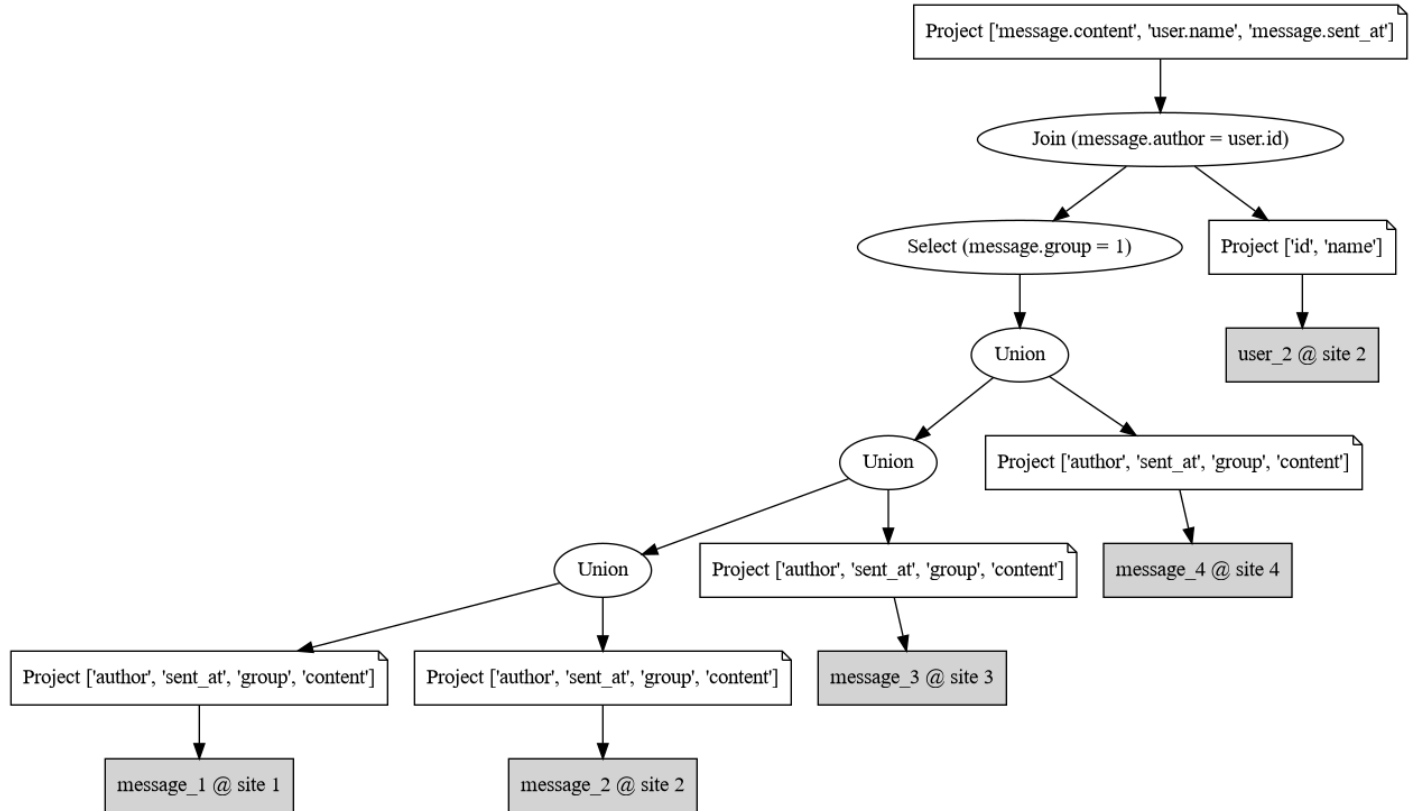
## Initial Query Tree



## Localized Query Tree



## Optimized Query Tree



## Output

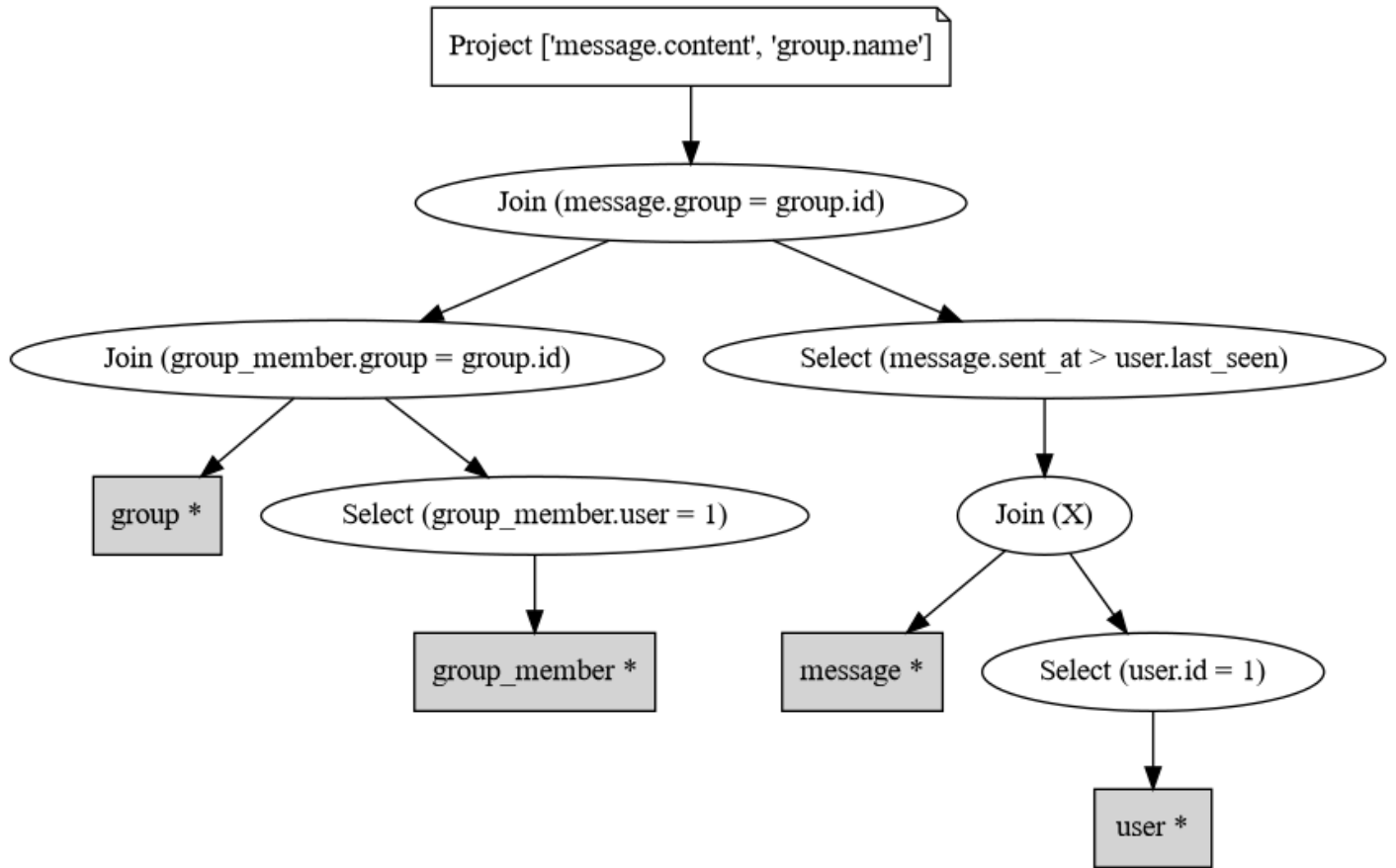
```
yk1,1647254210,hello 2
yk2,1647254211,hello 1
yk2,1647254212,what are you doing 1
yk1,1647254213,nothing
```

## Get all unseen messages of a user

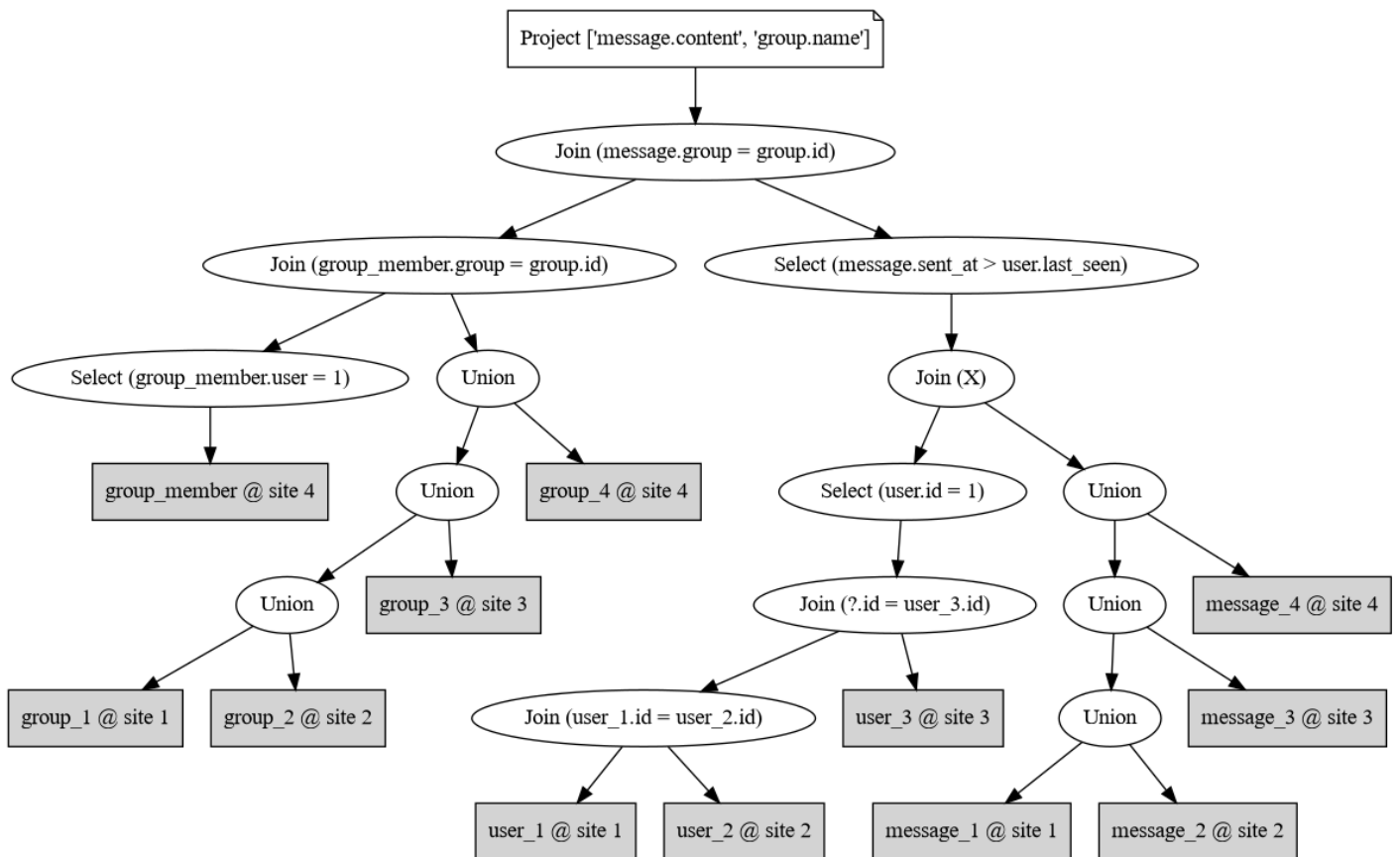
Given user id \$ID (= 1),

```
select G.`name`, M.`content`
  from `group` G, `message` M, `group_member` GM, `user` U
 where
   GM.`user` = $ID and
   U.`id` = $ID and
   GM.`group` = G.`id` and
   M.`group` = G.`id` and
   M.`sent_at` > U.`last_seen`;
```

## Initial Query Tree

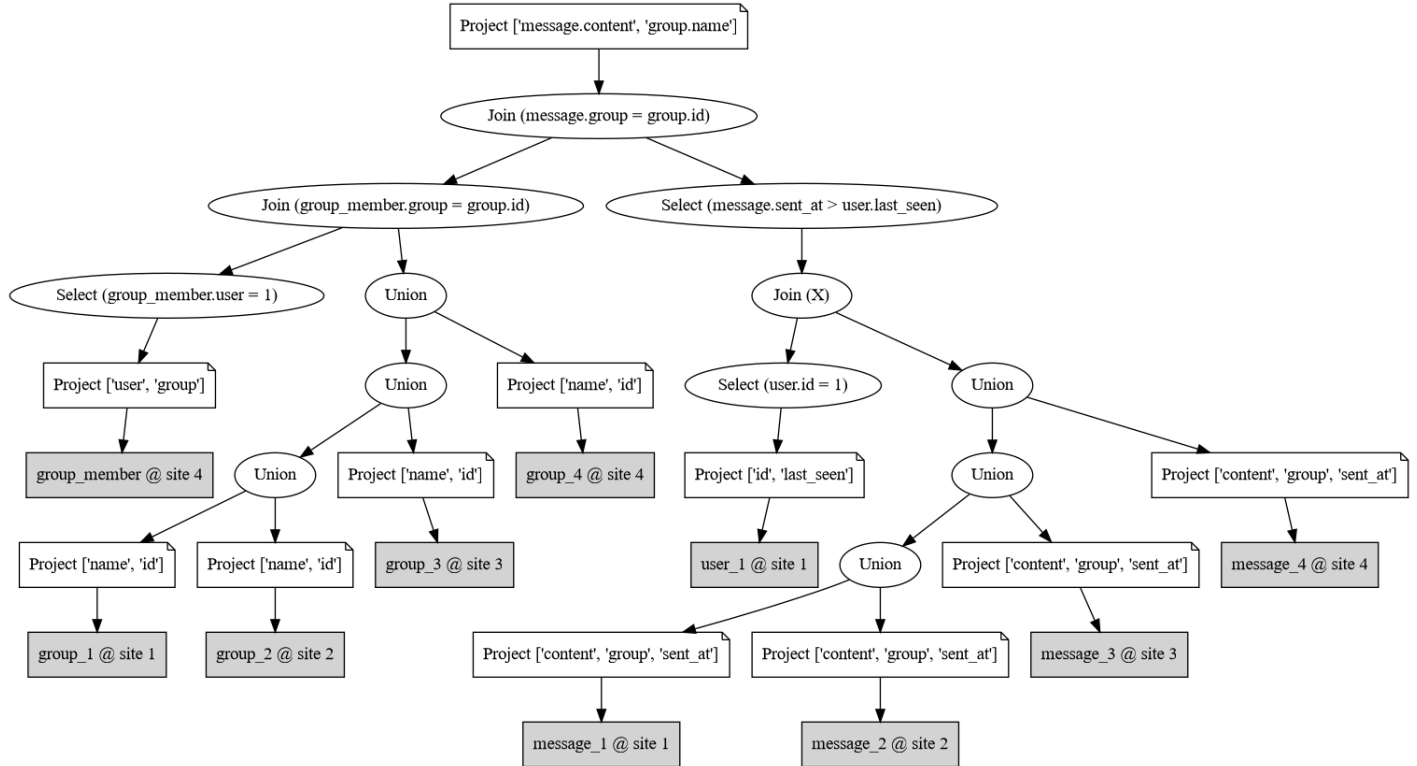


## Localized Query Tree

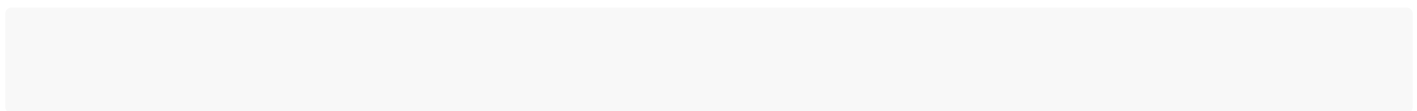


## Optimized Query Tree





## Output



## Data structures used

- Custom `dataclass` was made for each of the relations and system catalog tables (created dynamically from system catalog or csv files).
- Implemented custom wrapper over `list : PyQL` which allowed me to use sql-like queries on previously mentioned `dataclass` objects.
- SQL query was parsed into a custom `SelectQuery` object, which separates out everything that is needed for budding the query tree into individual components. Joins (only inner) are also translated to `where` conditions. Conditions are represented by three objects (which can be nested in each other, except `Condition`): `Condition`, `ConditionAnd`, `ConditionOr`.
- For query tree, each type of node: `ProjectionNode`, `SelectionNode`, `UnionNode`, `JoinNode`, `RelationNode` are custom hashable classes. The graph itself is stored as `networkx.DiGraph`.

## Algorithms used

- Parsing the select query: bunch of if conditions, storing stuff in lists and dictionaries for future ease.
- `where` clause conditions are reduced to simplest form.  $((A \ \&\& \ B) \ \&\& \ C \rightarrow A \ \&\& \ B \ \&\& \ C)$ .  
Check if child conditions are of the same type as parent and extract the conditions into parent

recursively.

- Building the query tree: sort by `where` condition clauses, start with the most specific. Traverse graph from leaf nodes ( `RelationNode` s) and find the root; add condition node there. If the condition refers to two or more columns, add `JoinNode` s before them.
- Localization: break the `RelationNode` s into fragments using information from system catalog
- Optimization: Traverse the tree from each `RelationNode` to find out which columns from that relation where used, add a project on all `RelationNode` s. For vertical fragments, remove the fragments from which no columns are used. For horizontal fragments, use `z3` solver to check if the `SelectionNode` condition and fragment predicate are compatible.