コード生成 + Shift0/Reset0 の型システム

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answer type は考えていない.
後で，answer type を加えたやつを考える.
answer type modification については考えない

1 Syntax

$$v ::= x \mid \lambda x.e \mid \texttt{<}e\texttt{>}$$

$$\begin{aligned} e ::= {} & x \mid \lambda x.e \mid e\ e \\ & \mid \underline{\lambda} x.e \mid \underline{\mathbf{reset0}}\ e \mid \underline{\mathbf{shift0}}\ k \to e \mid \underline{\mathbf{throw}}\ k\ e \\ & \mid \underline{\mathbf{clet}}\ x = e\ \underline{\mathbf{in}}\ e \mid \texttt{<}e_1\texttt{>}\ \underline{@}\ \texttt{<}e_2\texttt{>} \mid \cdots \end{aligned}$$

2 Semantics

2.1 Evaluation Context

$$E ::= [] \mid E\ e \mid v\ E \mid \mathbf{reset0}\ E \mid \underline{\lambda} x.E$$

2.2 Operation Semantics

$$\begin{aligned} E[(\lambda x.e)\ v] &\rightsquigarrow E[e\{x := v\}] \\ E[\mathbf{reset0}\ v] &\rightsquigarrow E[v] \\ E[(\underline{\lambda} x.e)v] &\rightsquigarrow E[\underline{\lambda} y.e\{x := \texttt{<}y\texttt{>}\}] \\ E[\underline{\lambda} y.\texttt{<}e\texttt{>}] &\rightsquigarrow E[\texttt{<}\lambda y.e\texttt{>}] \\ E[\mathbf{reset0}(E'[\mathbf{shift0}\ k \to E''[\underline{\mathbf{throw}}\ k\ e]])] &\rightsquigarrow E[E''[e\{k := \underline{\lambda} x.\mathbf{reset0}\ (E'[x])\}]] \\ E[\texttt{<}e_1\texttt{>}\ \underline{@}\ \texttt{<}e_2\texttt{>}] &\rightsquigarrow E[\texttt{<}e_1\ e_2\texttt{>}] \\ E[\underline{\mathbf{clet}}\ x = e_1\ \underline{\mathbf{in}}\ e_2] &\rightsquigarrow E[\underline{\lambda} x.e_2\ \underline{@}\ e_1] \end{aligned}$$

3 Type System

$$t ::= \text{BasicType} \mid t \to t \mid \langle t \rangle^{\gamma}$$

Typing rule for code-level lambda:

$$\frac{\Gamma,\ \gamma_1 \geq \gamma,\ x : \langle t_1 \rangle^{\gamma_1} \vdash e\ :\ \langle t_2 \rangle^{\gamma_1}}{\Gamma \vdash \underline{\lambda} x.e\ :\ \langle t_1 \to t_2 \rangle^{\gamma}}\ (\gamma_1 \text{ is eigen var})$$

Typing rule for code-level let (derived rule):

$$\frac{\Gamma \vdash e_1\ :\ \langle t_1 \rangle^{\gamma} \quad \Gamma,\ \gamma_1 \geq \gamma,\ x : \langle t_1 \rangle^{\gamma_1} \vdash e_2\ :\ \langle t_2 \rangle^{\gamma_1}}{\Gamma \vdash \underline{\mathbf{clet}}\ x = e_1\ \underline{\mathbf{in}}\ e_2\ :\ \langle t_2 \rangle^{\gamma}}\ (\gamma_1 \text{ is eigen var})$$

Typing rule for code-level reset0:

$$\frac{\Gamma \vdash e\ :\ \langle t \rangle^{\gamma}}{\Gamma \vdash \underline{\mathbf{reset0}}\ e\ :\ \langle t \rangle^{\gamma}}$$

Typing rule for code-level shift0:

$$\frac{\Gamma,\ k : (\langle t_1 \rangle^{\gamma_1} \Rightarrow \langle t_0 \rangle^{\gamma_0}) \vdash e\ :\ \langle t_0 \rangle^{\gamma_0} \quad \Gamma \models \gamma_1 \geq \gamma_0}{\Gamma \vdash \underline{\mathbf{shift0}}\ k \to e\ :\ \langle t_1 \rangle^{\gamma_1}}$$

Typing rule for code-level throw:

$$\frac{\Gamma,\ \gamma_3 \geq \gamma_1,\ \gamma_3 \geq \gamma_2 \vdash e\ :\ \langle t_1 \rangle^{\gamma_3} \quad \Gamma \models \gamma_2 \geq \gamma_0}{\Gamma,\ k : (\langle t_1 \rangle^{\gamma_1} \Rightarrow \langle t_0 \rangle^{\gamma_0}) \vdash \underline{\mathbf{throw}}\ k\ e\ :\ \langle t_0 \rangle^{\gamma_2}}\ (\gamma_3 \text{ is eigen var})$$

4 Example

$$\begin{aligned} e_1 = {} & \underline{\mathbf{reset0}}\ \ \underline{\mathbf{clet}}\ x_1 = \%3\ \underline{\mathbf{in}} \\ & \underline{\mathbf{reset0}}\ \ \underline{\mathbf{clet}}\ x_2 = \%5\ \underline{\mathbf{in}} \\ & \underline{\mathbf{shift0}}\ k\ \to\ \underline{\mathbf{clet}}\ y = t\ \underline{\mathbf{in}} \\ & \underline{\mathbf{throw}}\ k\ (x_1 \mathbin{\underline{+}} x_2 \mathbin{\underline{+}} y) \end{aligned}$$

If $t = \%7$ or $t = x_1$, then e_1 is typable.
If $t = x_2$, then e_1 is not typable.

$$\begin{aligned} e_2 = {} & \underline{\mathbf{reset0}}\ \ \underline{\mathbf{clet}}\ x_1 = \%3\ \underline{\mathbf{in}} \\ & \underline{\mathbf{reset0}}\ \ \underline{\mathbf{clet}}\ x_2 = \%5\ \underline{\mathbf{in}} \\ & \underline{\mathbf{shift0}}\ k_2\ \to\ \underline{\mathbf{shift0}}\ k_1\ \to\ \underline{\mathbf{clet}}\ y = t\ \underline{\mathbf{in}} \\ & \underline{\mathbf{throw}}\ k_1\ (\underline{\mathbf{throw}}\ k_2\ (x_1 \mathbin{\underline{+}} x_2 \mathbin{\underline{+}} y)) \end{aligned}$$

If $t = \%7$, then e_1 is typable.
If $t = x_2$ or $t = x_1$, then e_1 is not typable.