

Asgmt#1- algorithm2

Ln#	Algorithm algortihm2	Cost	Times	Comments
1	num=0,num1=1,fibo=0	C1=3	1	3 variables assigned
2	For(int i=0;i<n;i++)	C2=1	1	Assigning of variable i
		C3=1	n+1	n successes and one failure
		C4=2	n	Incrementation and assigning of i
3	fibo=num+num1	C5=2	n	Addition and assigning repeated n times
4	num=num1	C6=1	n	Assignment n times
5	num1=fibo	C7=1	n	Assignment n times
6	Return num	C8=1	1	return

$$f(n)=C1*1+C2*1+C3(n+1)+C4(n)+C5(n)+C6(n)+C7(n)+C8*1=6n+6$$

Algorithm2, with $f(n)=6n+6$:

Requires $6n+6$ time units to solve a problem of size n

It requires time proportional to $6n$

Its growth rate is linear

F(n)	Growth Rate Analysis			Asymptotic Analysis
	Time Units	Prop to	Rate	Big-Oh
$F(n)=6n+6$	$6n+6$	$6n$	linear	$O(n)$

algorithm3

Ln#	Function Fibonacci	Cost	Times	Comments
1	If(n=0) return 0	C1=1	1	We take the max cost, which is the matrix initialization and calling matrixpower, so we neglect cost of return
2	Initialize FM	C2=12	1	Assign 4 matrix elements, each costs 3
3	Call MatrixPower(n-1)	C3=1	1	Function called once
4	Return fm[0][0]	C4=3	1	Accessing two array elements, and returning a value

$$F(n) \text{ Fibonacci} = C1*1 + C2*1 + C3*1 + C4*1 = 1 + 12 + 1 + 3 = 17$$

Ln#	Procedure MatrixPower	Cost	Times	Comments
1	If(n>1)	C1=1	1	Comparison done one time before recursively called again + failure
2	Call MatrixPower(n/2)	C2=1	f(n/2)	Recursive calls
3	Update FM=FM*FM	C3=56	1	
4	if(n is odd)	C4=1	1	comparison
5	Update FM=FM*{{1,1},{1,0}}	C5=56	1	

$$F(n) \text{ MatrixPower} = C1*1 + C2*f(n/2) + C3*1 + C4*1 + C5*1 = 4 + f(n/2) = 4\log_2 n + 1$$

$$F(n) \text{ of algorithm3} = f(n) \text{ Fibonacci} + f(n) \text{ MatrixPower} = 17 + 4\log_2 n + 1 = 18 + 4\log_2 n$$

Algorithm3, with $f(n) = 18 + 4\log_2 n$:

Requires $18 + 4\log_2 n$ time units to solve a problem of size n

It requires time proportional to $4\log_2 n$

Its growth rate in logarithmic

F(n)	Growth Rate Analysis			Asymptotic Analysis
	Time Units	Prop to	Rate	Big-Oh
$f(n) = 18 + 4\log_2 n$	$18 + 4\log_2 n$	$4\log_2 n$	logarithmic	$O(\log_2 n)$

References: CSPC 319 lecture slides

- "matrix operation costs from lecture material"

- "01 - Algorithm Design Patterns - (1) Recursion - soln-HW"