FO? $(A) \rightarrow (B,C)$ FO? $\{C\} \rightarrow \{A,D\}$ FO? $\{D,E\} \rightarrow \{F\}$

a) $\{C_3 \rightarrow \{B_3\}$ we can derive $\{C_3 \rightarrow \{B_3\}\}$ using transitivity $\begin{cases} x \rightarrow Y \\ a \rightarrow b & Y \rightarrow Z \\ then & X \rightarrow Y \end{cases}$ FD2 says that $\{C_3 \rightarrow \{A_3\}, avd \quad C_{10} \sim$ FD1 gives us $\{A_3 \rightarrow \{B_3\}, avd \quad C_{10} \sim$ directly through trasitivity we derive $\{C_3 \rightarrow \{B_3\}, avd \quad C_{10} \sim$

b) {A,E} -> (F)

From FD1 and FD2 through trasistivity $A \rightarrow C \rightarrow D$ we get $A \rightarrow D$ using augmentation (if $X \rightarrow Y$, then $XZ \rightarrow YZ$ for any Z)

we get $\{A,E\} \rightarrow \{D,E\}$ using transistivity on the new FD1 and FD3 $A,E \rightarrow D,E \rightarrow F$ we get $\{A,E\} \rightarrow \{F\}$

FD1 (A) -> (B. () Task 2 FO2 (C) -> [A,D] a) X = (A) Apply FD1 we get X+ = (A,B,C) we have C in X+, we can apply FO2 (A) Already in X+, we add (D) so X+= (A,B,C,D) we don't have (F) : x+ the attribute closure is X+=(A,B,C,D) b) X= (C, E) we have (C3, so we can apply FD2 we get X+ = { C, E, A, D} we have (D,E), we apply FD3 x+ = (C, E, A, D, F) lasticy apply FD1 so the authibute closure is X+= (C, E, A, D, F, B)

FO1 [A,B] -> (.D.E.F) Task 3 FOR (E3 -> (F3 a) determine candidate key: FO3 (D) > (B) (A.B) attribute closure is X+ = (A, B, C, D, E, F, B] so it functionally determines all attributes in R. we need to check it for minimalisty neither subsets (A) or (B) can determine R alone . so {A,B} is minimal. The other attributes (E 3 and (D) can't determine 12 alone. So {A,B} is the candidate key. b) we check that for X->4, X is a super key for R. [A.B] is a super key, [E] and [D] are not Superkeys, they don't determine all other attibutes in R. Both FD2 and FD3 violate BCNF. FDZ and FD3 violates BCNF as mentioned in question to. Decomposition of relation that violate BCNF start with FD2, decompose into R1, R2 R,(E,F) , R, (A,B,C,D,E) , decompose FD3 into R3, R4 R3 (D, B) , R4 (A, C, D, E) check for candidate key in all new relations R.(E,F) with FD2, candidate Ken: E R3(D,B) with FD3, cardidate key: D R4, FDI is superkey, but we weed D to link it to B, so we choose candidate (14,D)

FDI {A,B,C} → (D,E) Task 4 FOR (B, C, D) - (A, E) a) we need to check every FD3 [C] - [0] FD X > Y where X should be a superkey. FD1: X is a superkey, it doesn't violate BEMF FD2: X is also a supervey, doesn't violate BCNF FD3: (C3 only determinen D, it's not asuper key, violates DCOF R is not in BCNF, because FD3 violates its condition bl-we determind that FD3 violates DCHF - Decompose R into R, and R_L, preserving FD3:

R, (C, D) from FD3, C :s key

R₂(A,B,C,E) for the remaining attributes

-R₂ related new FDS: FD, (A,B,C) \rightarrow (E)

FD₂ [B,C] \rightarrow (A,E) (D determined by C) - check if X in FD1 and FD2 are superkeys. 19 FD. is a super key FD, is not a supertey (missing D) - decompose Rz based on FD2: R3 (B,C,A,E) from FD2 R4 (B,C,E) for rest of attributes eq(B,C,c).

- check candidate key and BCNF:

R,(C,D), {C} is condidate key, {C} -> {D}

R3 (B,C,M,E), {B,C} is candidate key, {B,C} -> {A,E}

R4 (B,C,E), {B,C} is and ideate key, it determines {A,E}

R4 (B,C,E), {B,C} is and ideate key, in R3, {A,E} determines

all other attributes in R2