Lab 3: Constraint Satisfaction Problems

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1 Exercise 1: Rearrange the Seats

1.1 _classroom_conflict()

This function is to check if $var1 \leftarrow val1$ and $var2 \leftarrow val2$ will cause conflicts. There are two conditions to cause conflicts. One is that val1 == val2 and another is that var1 and var2 is adjacent to each other and they are friends. So my implementation is as follows:

1.2 backtracking()

I use nconflicts() to check consistency, assign() to assign a value to a variable and unassign() to unassign a value from a variable. My implementation is as follows:

```
def backtracking(
 1
2
        csp,
3
        select_unassigned_variable=mrv,
        order_domain_values=lcv
4
5
   ):
6
        def backtrack(assignment):
 7
            if len(assignment) == len(csp.variables):
8
                return assignment
9
            var = select_unassigned_variable(assignment, csp)
            for value in order_domain_values(var, assignment, csp):
10
                if csp.nconflicts(var, value, assignment) == 0:
11
12
                    csp.assign(var, value, assignment)
                    result = backtrack(assignment)
13
14
                    if result != None:
15
                        return result
                    csp.unassign(var, assignment)
16
17
            return None
        result = backtrack({})
18
        assert result is None or csp.goal_test(result)
19
20
        return result
```

1.3 result

The result of easy_classroom is as follows:

```
Time consumption: 0.0056s

Result: □ Success

Solution:

[[4 2]

[3 1]

[5 6]

[7 8]]
```

The result of fail classroom is as follows:

```
1 Time consumption: 0.0001s 2 Result: 	t B Fail (no solution found).
```

2 Exercise 2: Sudoku (Filtering)

2.1 forward_checking()

The forward_checking() is to prune neighbor values inconsistent with var=value. My implementation is as follows:

2.2 AC3()

The AC3() is to make sure all arcs are consistent after assigning a variable. My implementation is as follows:

```
def AC3(csp, removals=None):
1
2
       def revise(Xi, Xj):
           removed = False
3
           for i in csp.curr_domains[Xi]:
4
               consistent = False
               for j in csp.curr_domains[Xj]:
6
                   if csp.constraints(Xi, i, Xj, j):
                        consistent = True
8
               if not consistent:
9
```

```
csp.prune(Xi, i, removals)
10
11
                    removed = True
12
            return removed
13
14
        queue = {(Xi, Xk) for Xi in csp.variables for Xk in csp.neighbors[Xi]}
        csp.support_pruning() # It is necessary for using csp.prune()
15
16
       while queue:
17
            Xi, Xj = queue.pop()
            if revise(Xi, Xj):
18
19
                for Xk in csp.neighbors[Xi]:
20
                    queue.add((Xk, Xi))
21
22
        return True # CSP is satisfiable
```

2.3 backtracking_with_inference()

The backtracking_with_inference() is to use filtering strategies in backtracking algorithm. My implementation is as follows:

```
def backtracking_with_inference(
 1
2
        csp,
3
        inference,
 4
        select_unassigned_variable=mrv,
       order_domain_values=lcv
5
6
   ):
7
       def backtrack(assignment):
            if len(assignment) == len(csp.variables):
8
9
                return assignment
            var = select_unassigned_variable(assignment, csp)
10
11
            for value in order_domain_values(var, assignment, csp):
12
                if csp.nconflicts(var, value, assignment) == 0:
                    removals = csp.suppose(var, value)
13
14
                    csp.assign(var, value, assignment)
                    inference(csp, var, value, assignment, removals)
15
                    result = backtrack(assignment)
16
                    if result != None:
17
                        return result
18
19
                    csp.unassign(var, assignment)
20
                    csp.restore(removals)
21
            return None
22
23
        result = backtrack({})
24
        assert result is None or csp.goal_test(result)
25
        return result
```

2.4 result

The result of forward_checking() on easy_sudoku is as follows:

The result of AC3() on easy_sudoku is as follows:

The result of forward_checking() on harder_sudoku is as follows:

The result of AC3() on harder_sudoku is as follows:

We can find that forward_checking() is quicker than AC3() because forward_checking() just needs to check neighbours of this variable while AC3() needs to check all arcs.

3 Exercise 3: N-Queens (Hill Climbing)

3.1 min_conflicts()

I choose stochastic to implement min_conflicts() which means I randomly choose conflicted variables instead of choosing max_conflicts variable and assign min_conflicts_value to it. My implementation is as follows:

```
def min_conflicts(csp, max_steps=100000):
1
2
        assignment = {}
        for v in csp.variables:
3
            csp.assign(v, min_conflicts_value(csp, v, assignment), assignment)
4
5
6
       for i in range(max_steps):
            vars = csp.conflicted_vars(assignment)
            if not vars:
8
                return assignment
9
10
            var = random.choice(vars)
11
            val = min_conflicts_value(csp, var, assignment)
            csp.assign(var, val, assignment)
12
13
14
        return None
```

Its result on 8_nqueens is as follows:

3.2 min_conflicts() for Sudoku

The result of min_conflicts() on easy_sudoku is as follows:

```
1 Time consumption: 46.5223s
2 Result: ② Fail (no solution found).
```

It failed to find the solution because it will easily get stuck in a local optima because there are too many variables in sudoku problems and stochastic will not work well.