



IEEE/ACM CCGRID 2023

The 23rd IEEE/ACM International Symposium on Cluster, Cloud and Internet Computing

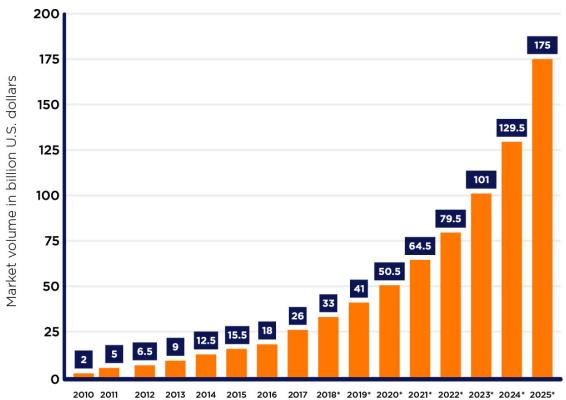
Highly Scalable Large-Scale Asynchronous Graph Processing using Actors

Youssef Elmougy*, Akihiro Hayashi, and Vivek Sarkar
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* Corresponding Author and Presenting Author

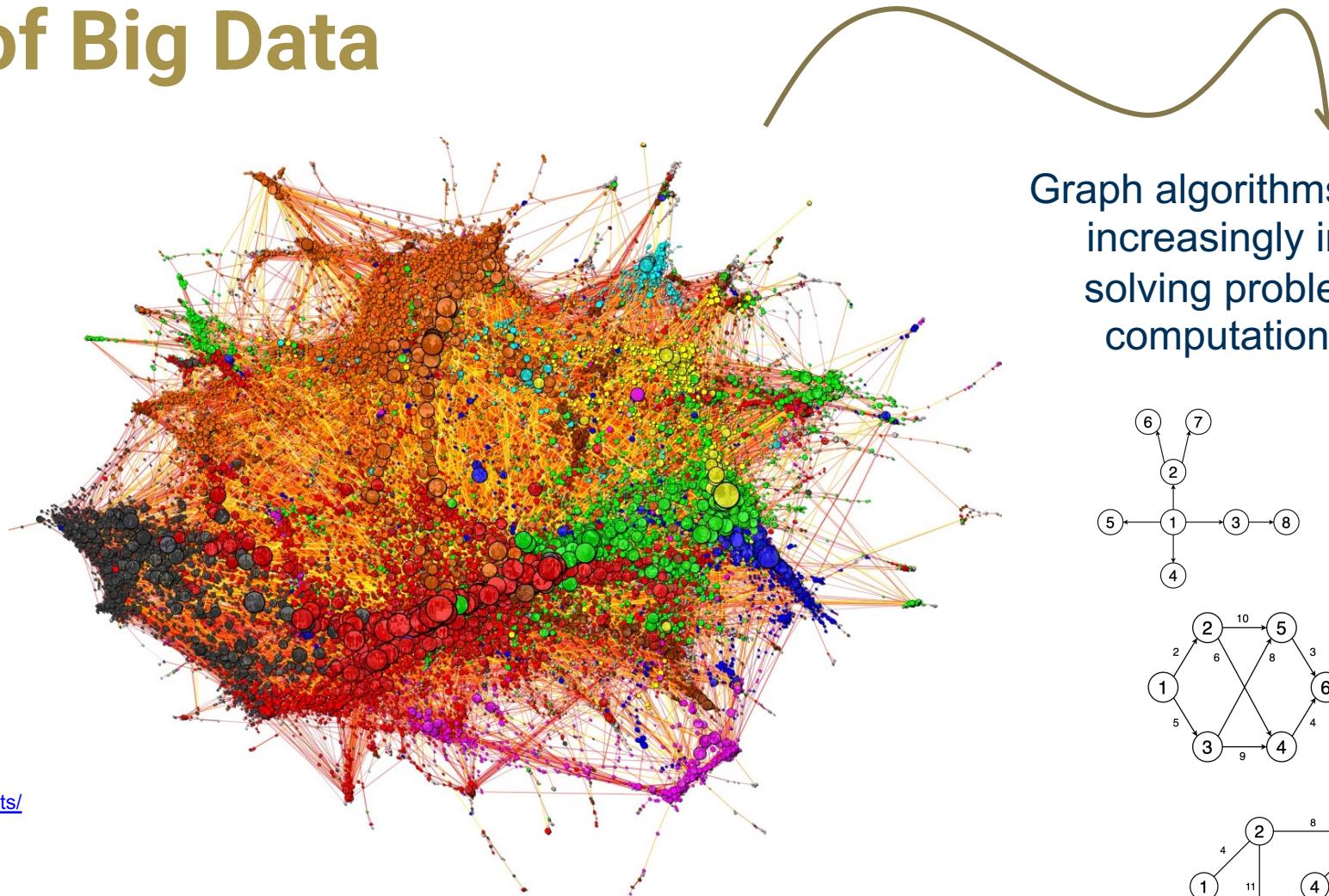
The Growth of Big Data

INFORMATION CREATED GLOBALLY 2010-2025

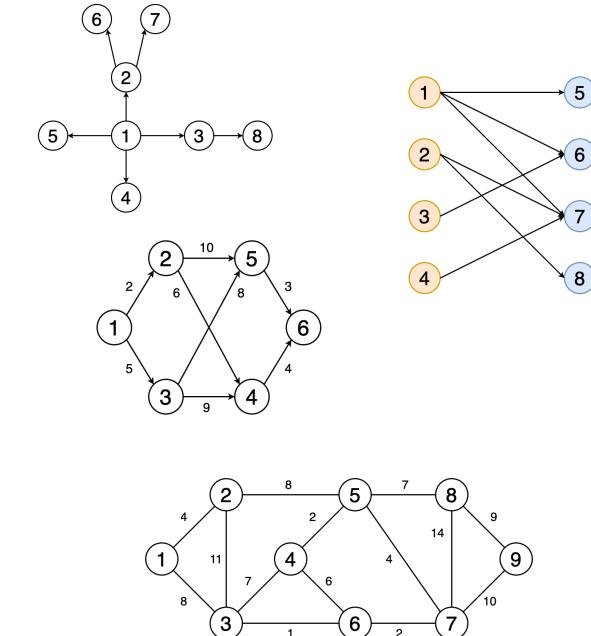


Picture borrowed from:

<https://www.iteratorshq.com/blog/big-data-business-impacts/>

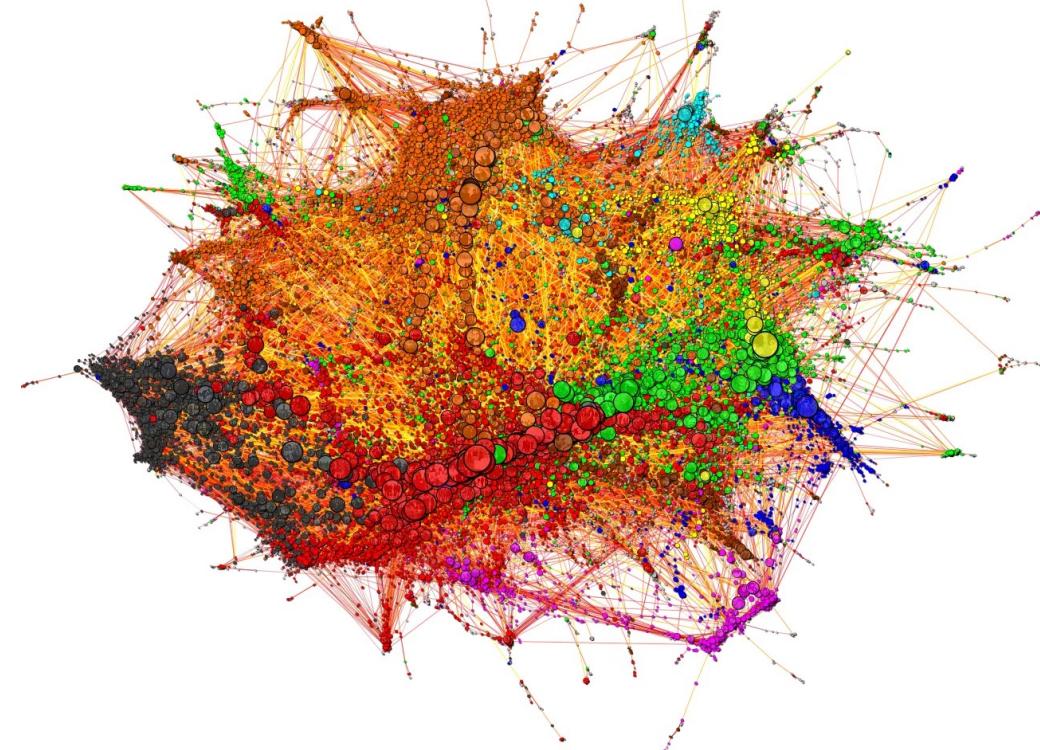
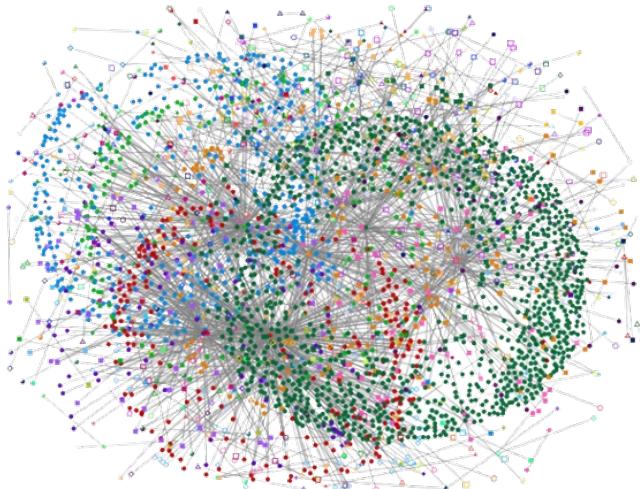
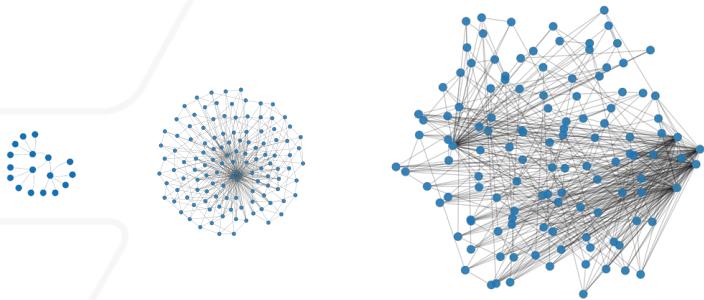


Graph algorithms have become increasingly important for solving problems in many computational domains



The scale of these graphs present difficulties to their processing and analysis!

The Scalability Problem



increasing graph size and complexity



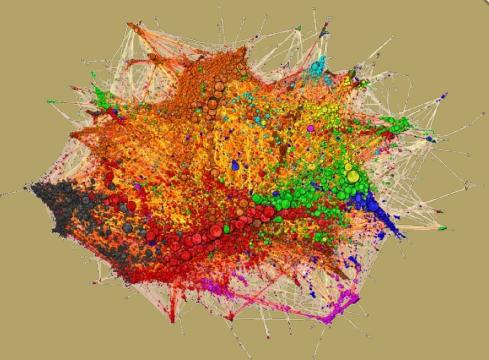
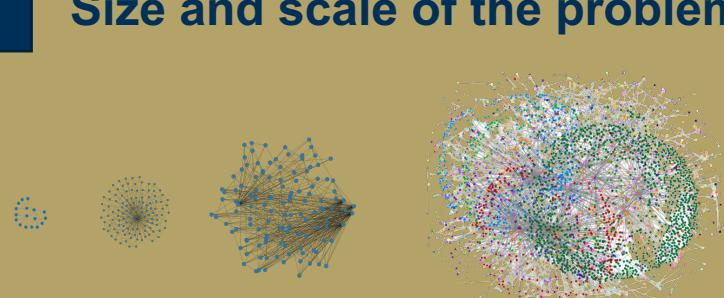
- There is a need for parallel computing resources to meet the computational and memory requirements
- Existing algorithms and software that perform well for mainstream parallel scientific applications are not necessarily efficient for large-scale graph applications

It is critical to develop lightweight and scalable systems to efficiently process large-scale graphs!

The 4 Overarching Challenges in Parallel Graph Processing

1

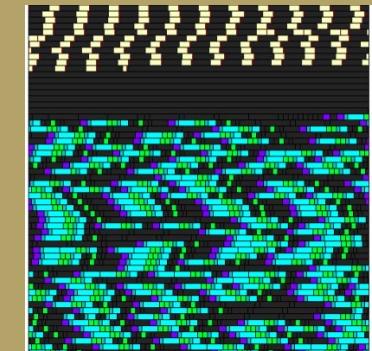
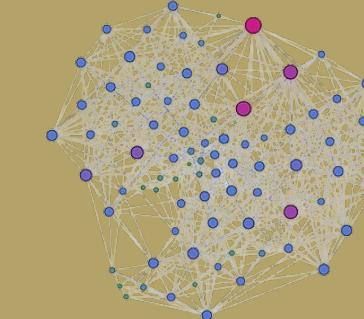
Size and scale of the problem



increasing graph size and complexity

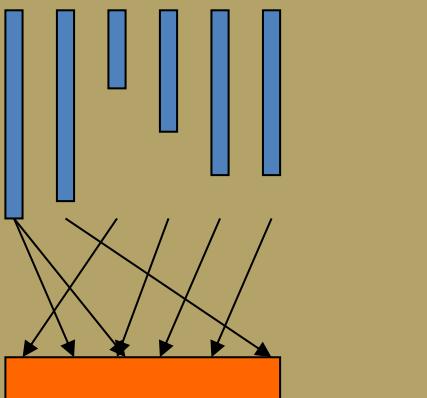
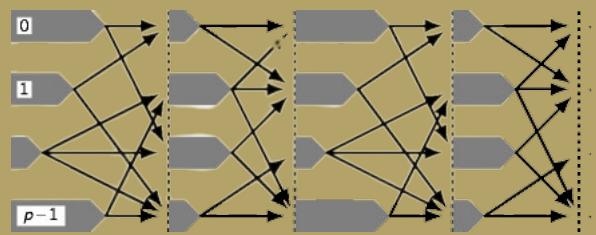
3

Poor locality of memory access



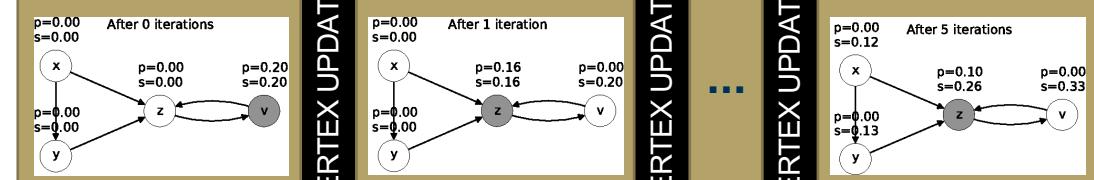
2

Synchronization overhead



4

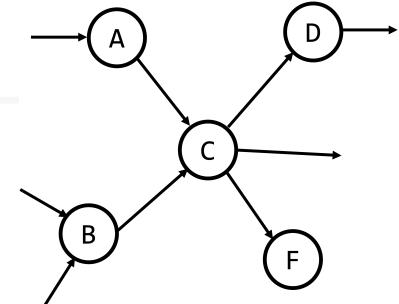
Slow convergence



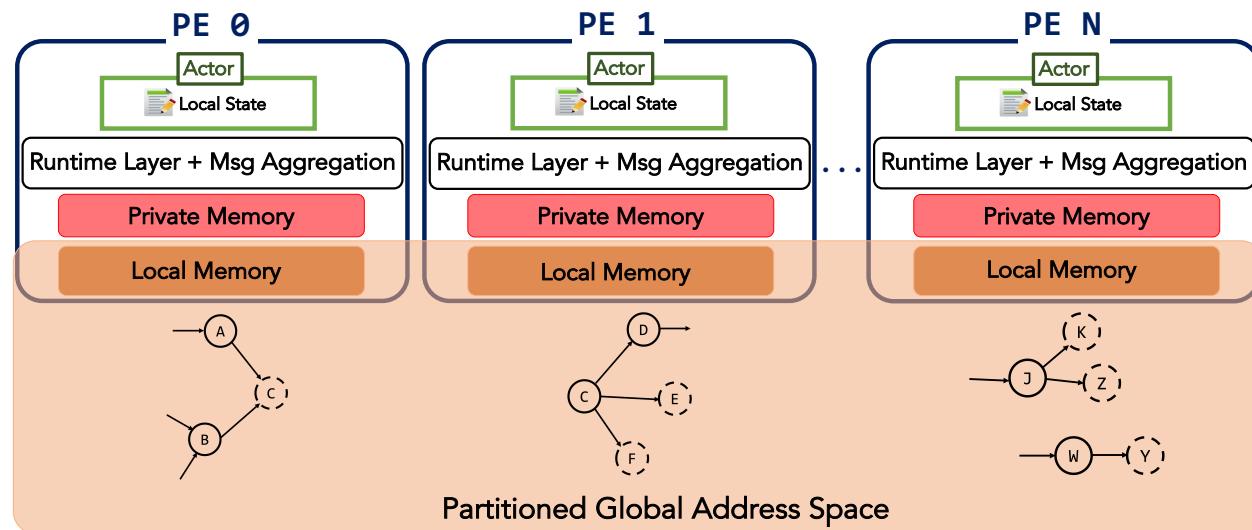
This paper studies the scalability of our novel actor-based programming system to overcome the inherent challenges of large-scale graph processing!

Actor-based Scalable Architecture for Graph Processing

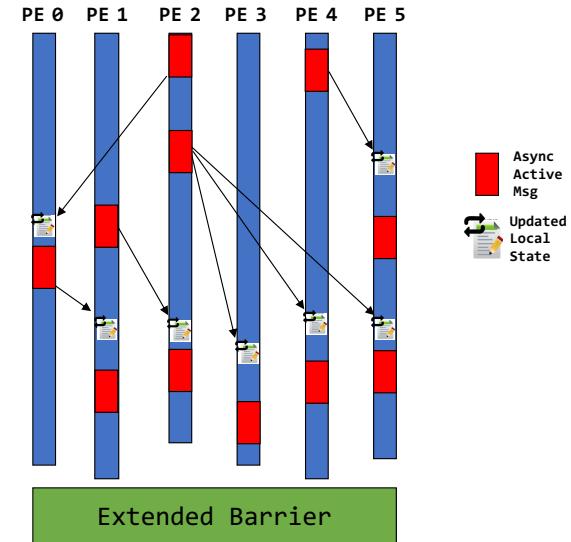
Sample Graph



Distributed Actor Runtime



Execution Model

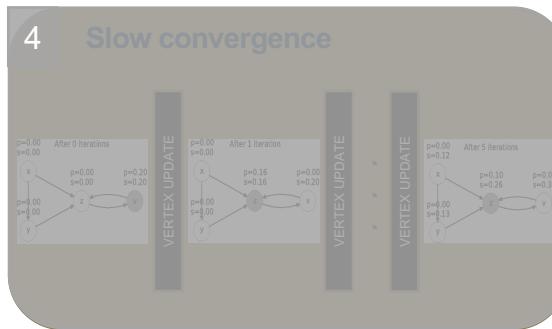
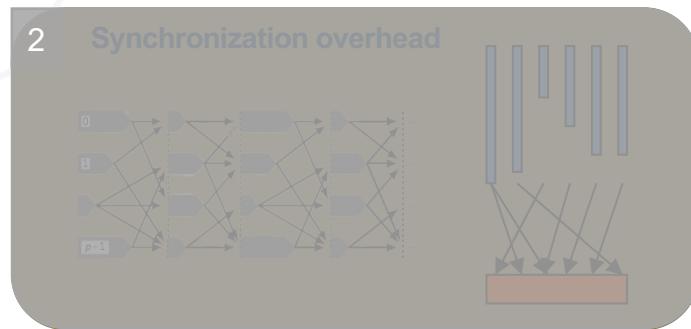
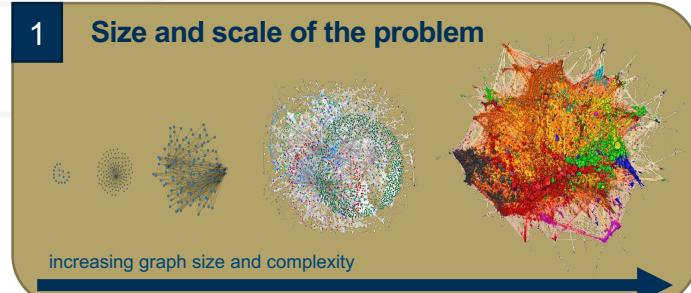


- Presents a lightweight, asynchronous computation model
- Utilizes fine-grained asynchronous actor messages to express point-to-point remote operations
- Treats actors as primitives of computation, where actors are inherently isolated and share no mutable state
- Actors process messages sequentially within its mailbox, thereby avoiding data races and synchronization

NOTE: “Actor” and “Selector” will be used interchangeably

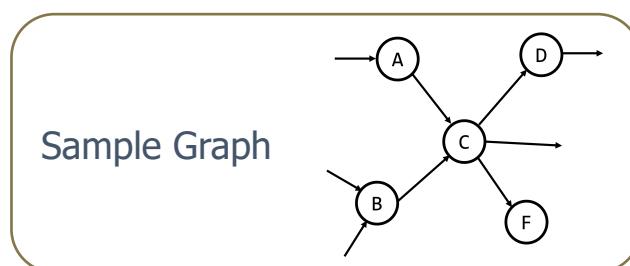
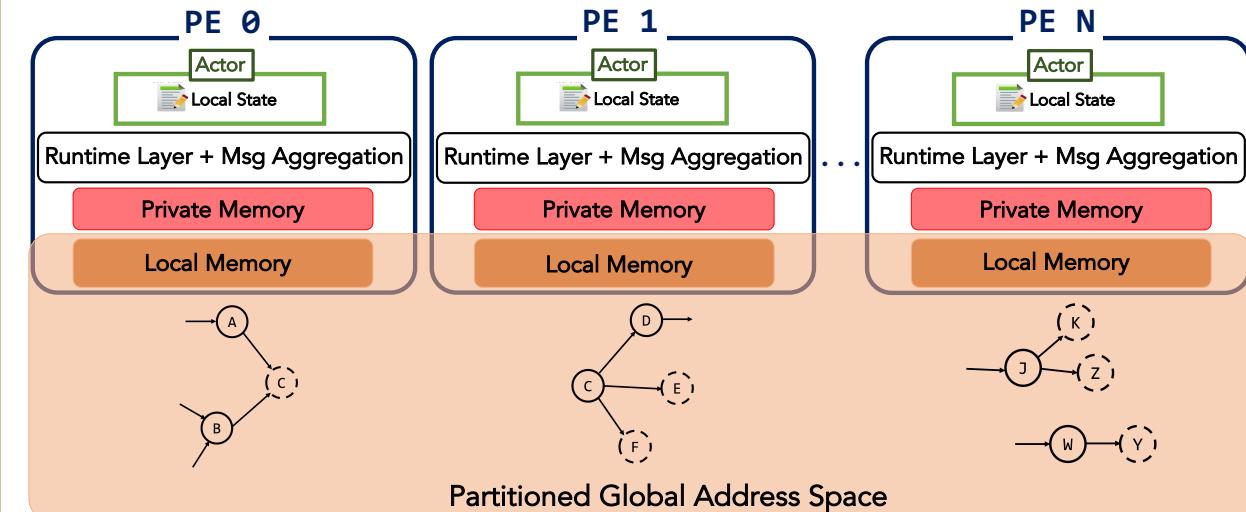
Actor-based Scalable Architecture for Graph Processing

The 4 overarching challenges of parallel graph processing:



SOLUTION

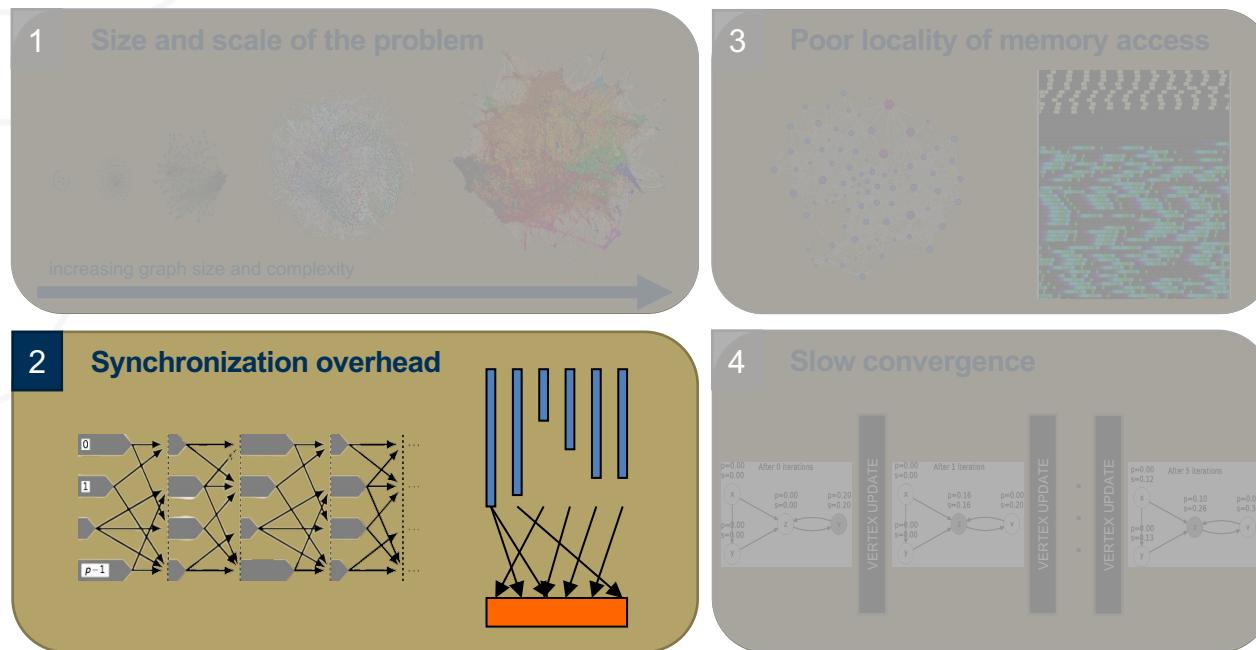
Distributed Actor Runtime



- Large-scale graph can be distributed across multiple PEs, where the local partition per PE is small enough to fit in its local memory

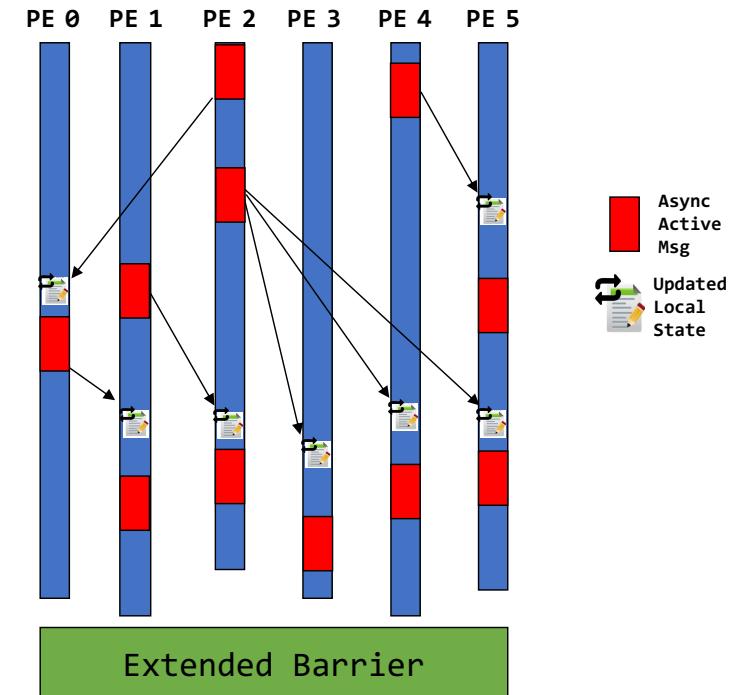
Actor-based Scalable Architecture for Graph Processing

The 4 overarching challenges of parallel graph processing:



SOLUTION

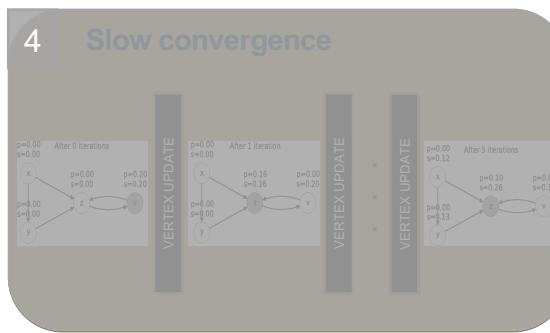
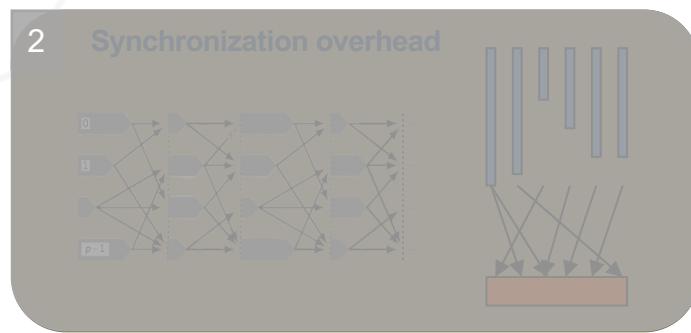
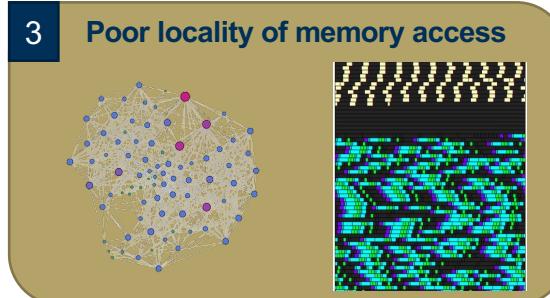
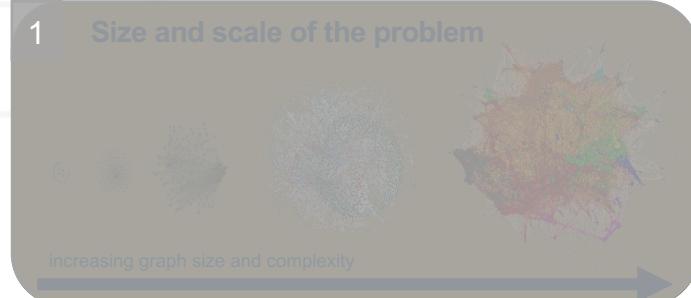
Execution Model



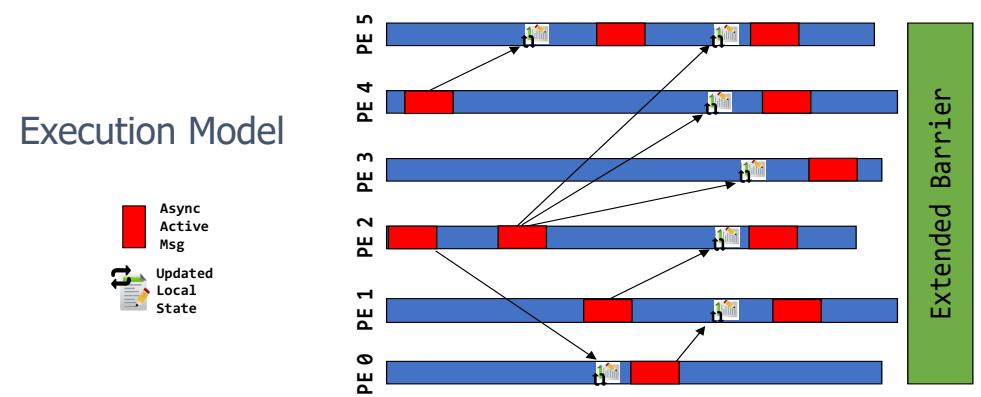
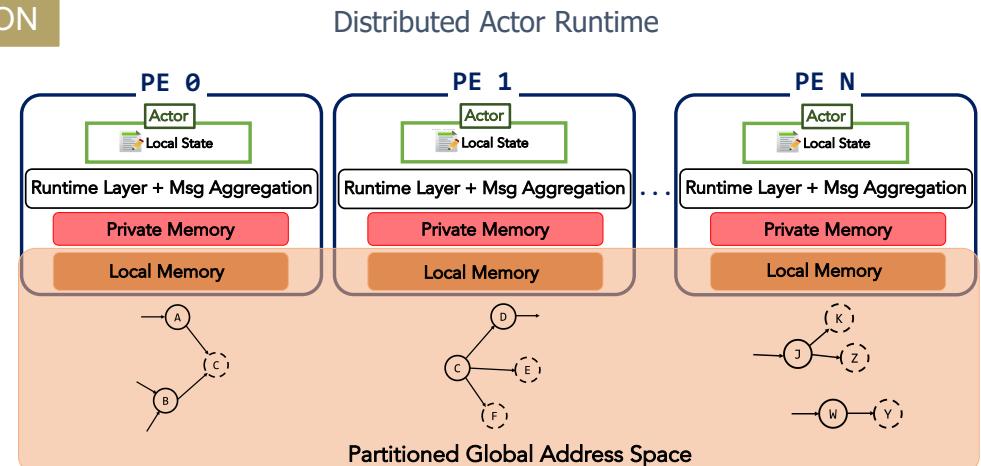
- Fine-grained Asynchronous Bulk-Synchronous (FABS) Parallelism model
- Reduces barriers and time spent idling at barriers, further reducing stall cycles

Actor-based Scalable Architecture for Graph Processing

The 4 overarching challenges of parallel graph processing:



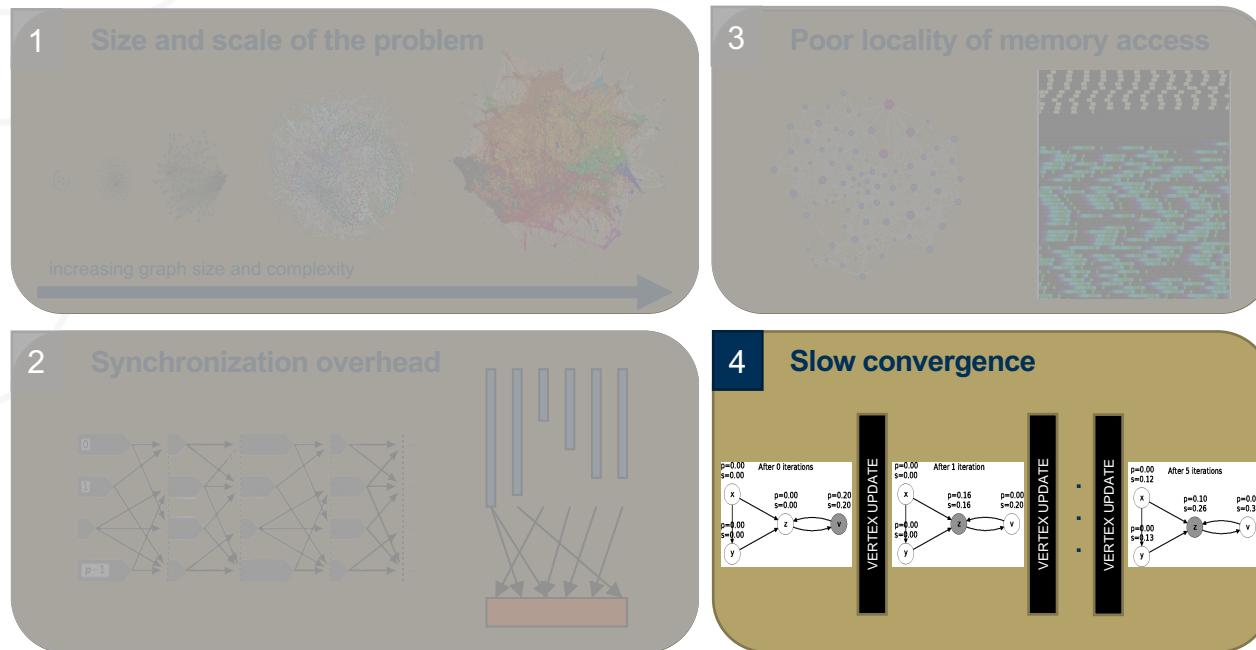
SOLUTION



- Well suited for irregular applications due to the Partitioned Global Address Space (PGAS)
- Expresses point-to-point remote operations as short, one-sided fine-grained async messages
- Message aggregation allows for low overhead and high network utilization
- Computation is migrated to where the data is located (moving compute to data)

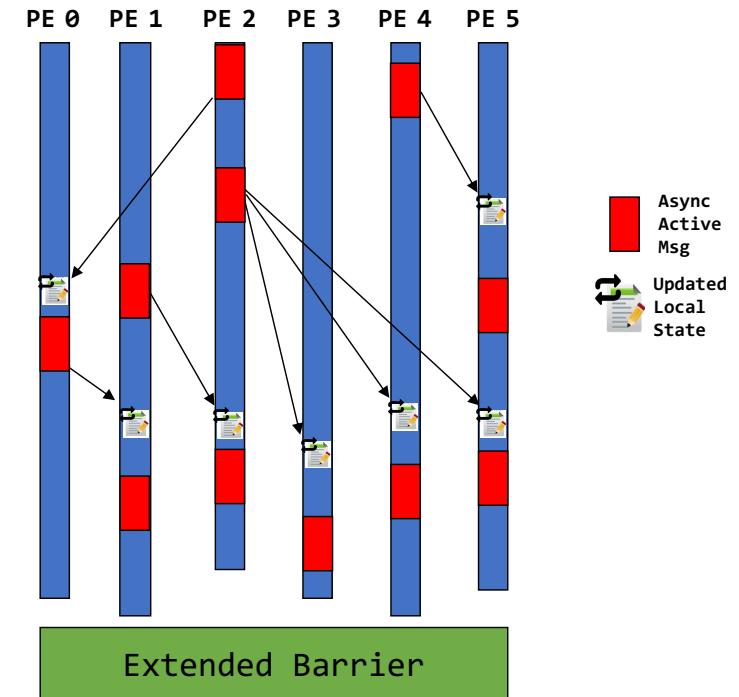
Actor-based Scalable Architecture for Graph Processing

The 4 overarching challenges of parallel graph processing:



SOLUTION

Execution Model



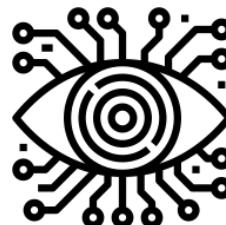
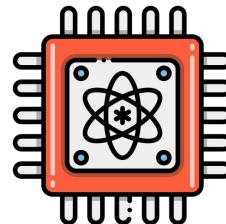
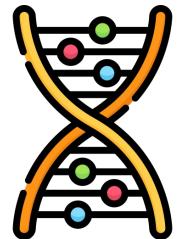
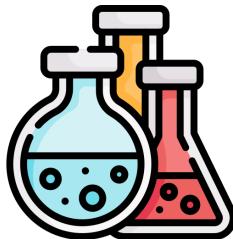
- As actor messages are executed, the local state of the actor is updated, allowing the next received actor message to utilize the updated state within the same super-step
- Updates the neighboring vertices with most recent values within the same iteration

Showing the System's Extreme Scalability with PageRank & Jaccard Index

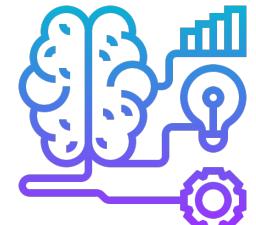
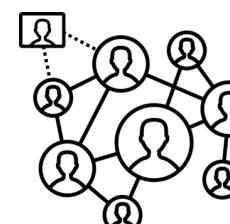
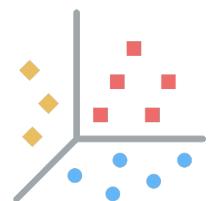
We focus on PageRank and Jaccard Index due to two reasons:

1. They show iterative vs non-iterative application types
2. They have been applied to many real-world problems with social impact

PageRank

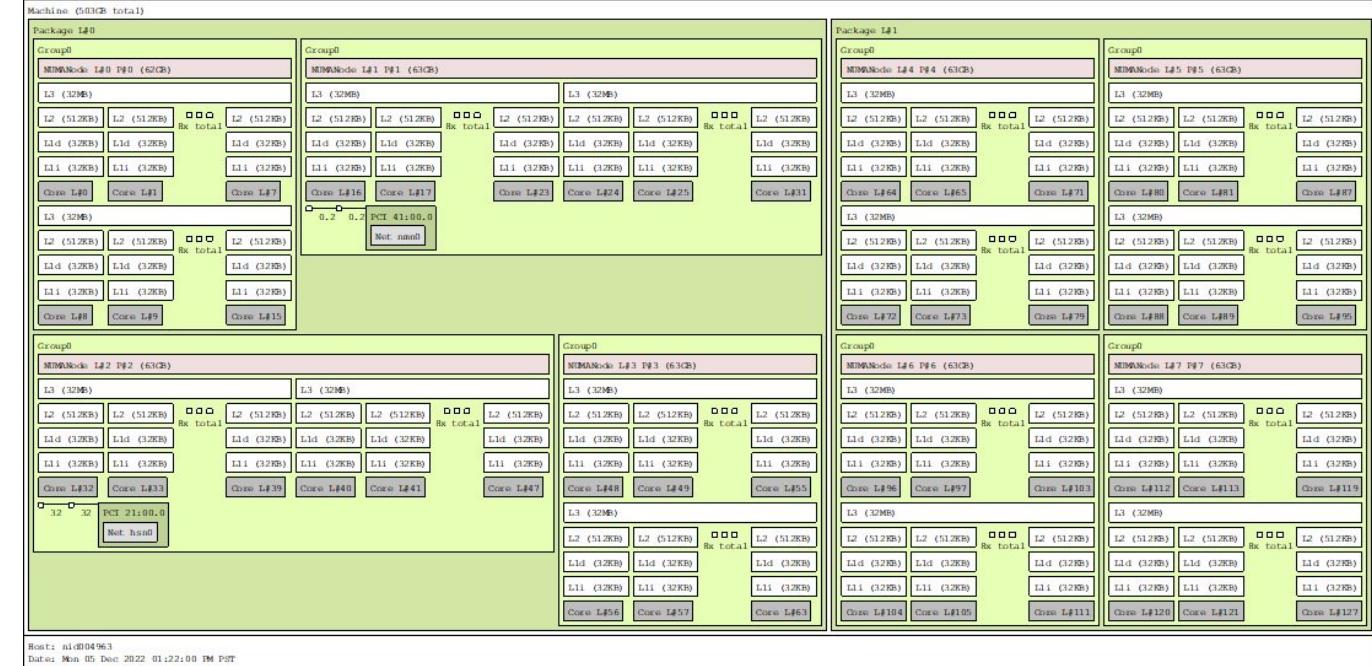
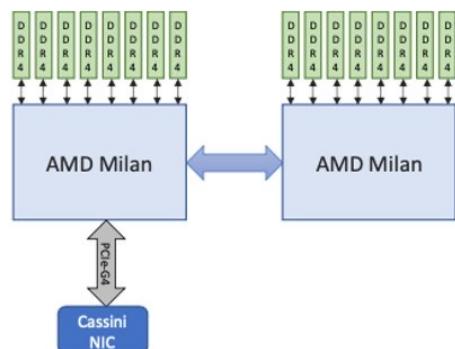


Jaccard Index



Experimental Setup and Architecture

- **Metrics:** execution time (in seconds) and traversed edges per second (TEPS)
- Experiments conducted on the CPU nodes of the **Perlmutter supercomputer** at the National Energy Research Scientific Computing Center (NERSC)
 - 2x AMD EPYC 7763 (Milan) CPUs
 - 64 physical cores per CPU
 - 512 GB memory
 - 1x HPE Cray Slingshot Interconnect
- Results for **different dimensions of scalability** are presented

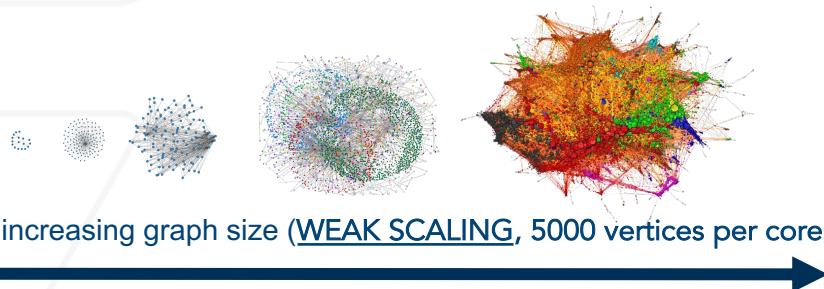


Picture borrowed from: <https://docs.nersc.gov/systems/perlmutter/architecture/>

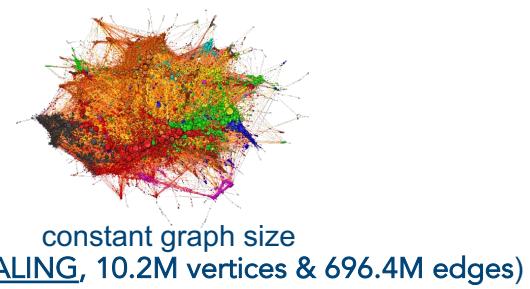
Dimensions of Scalability

Scaling performance is shown using three experiment types:

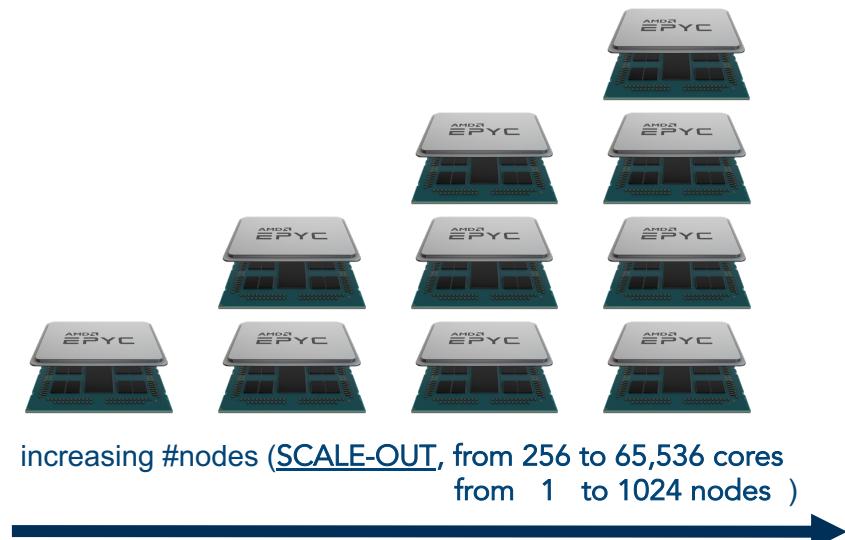
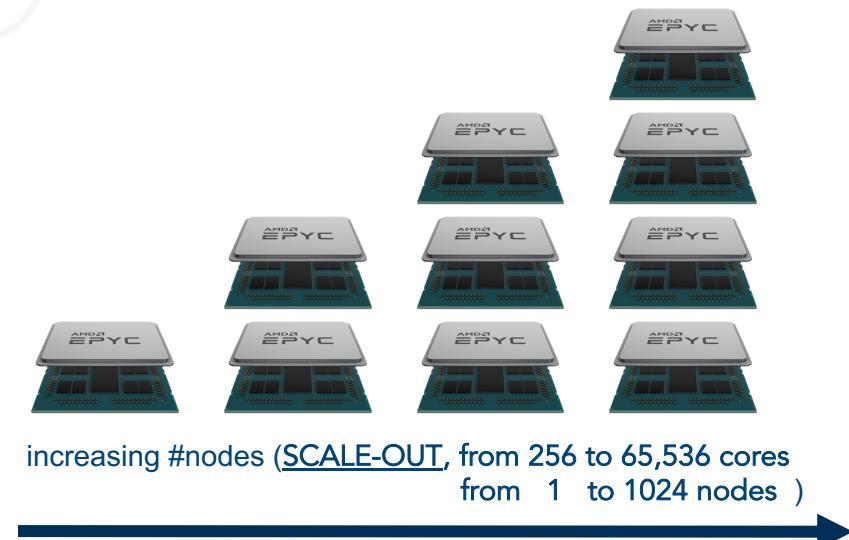
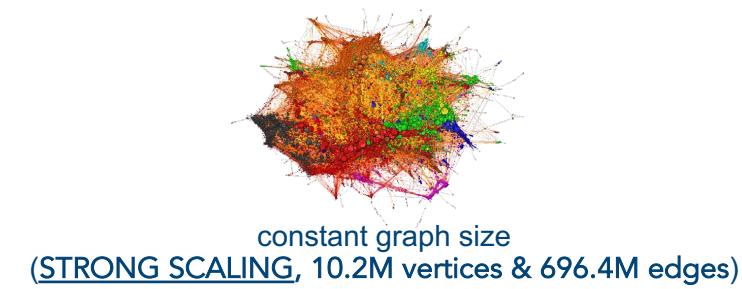
(1) SCALE1



(2) SCALE2



(3) SCALE3



Performance Results: SCALE1

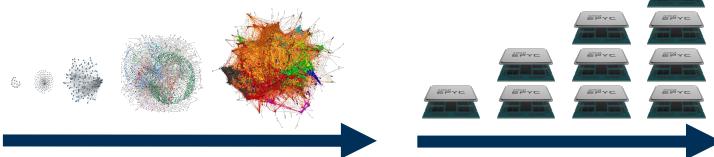
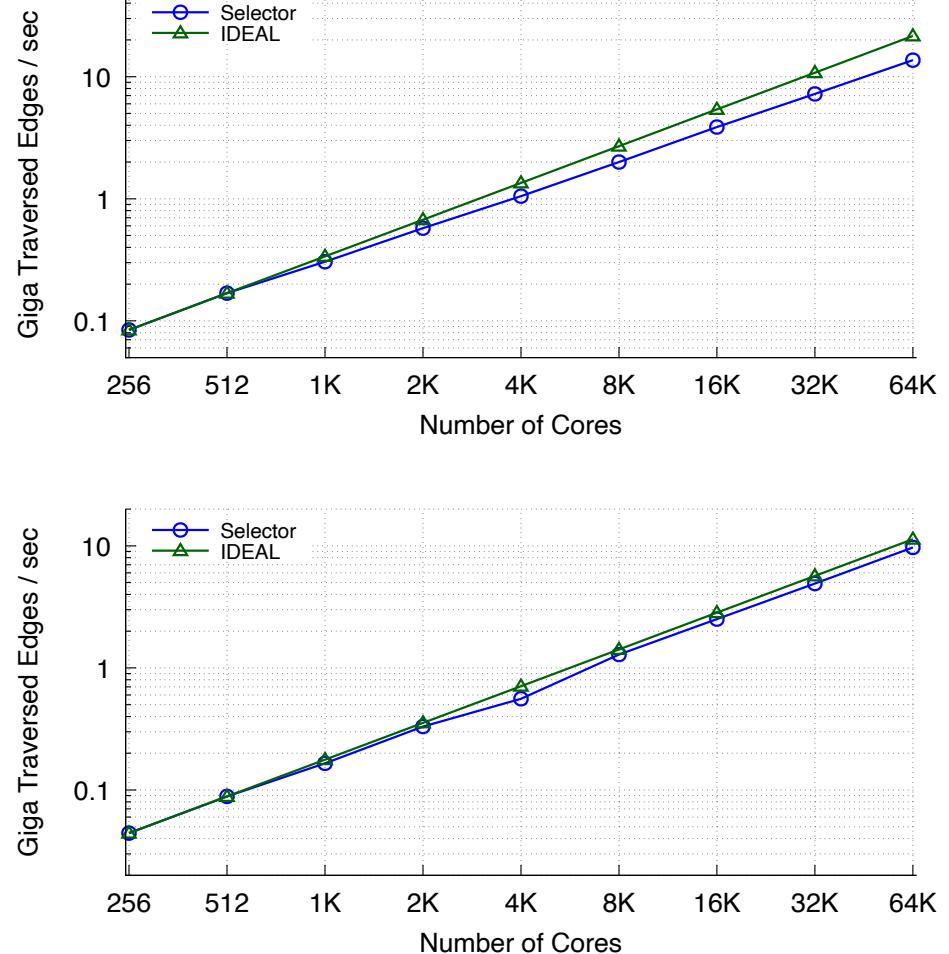
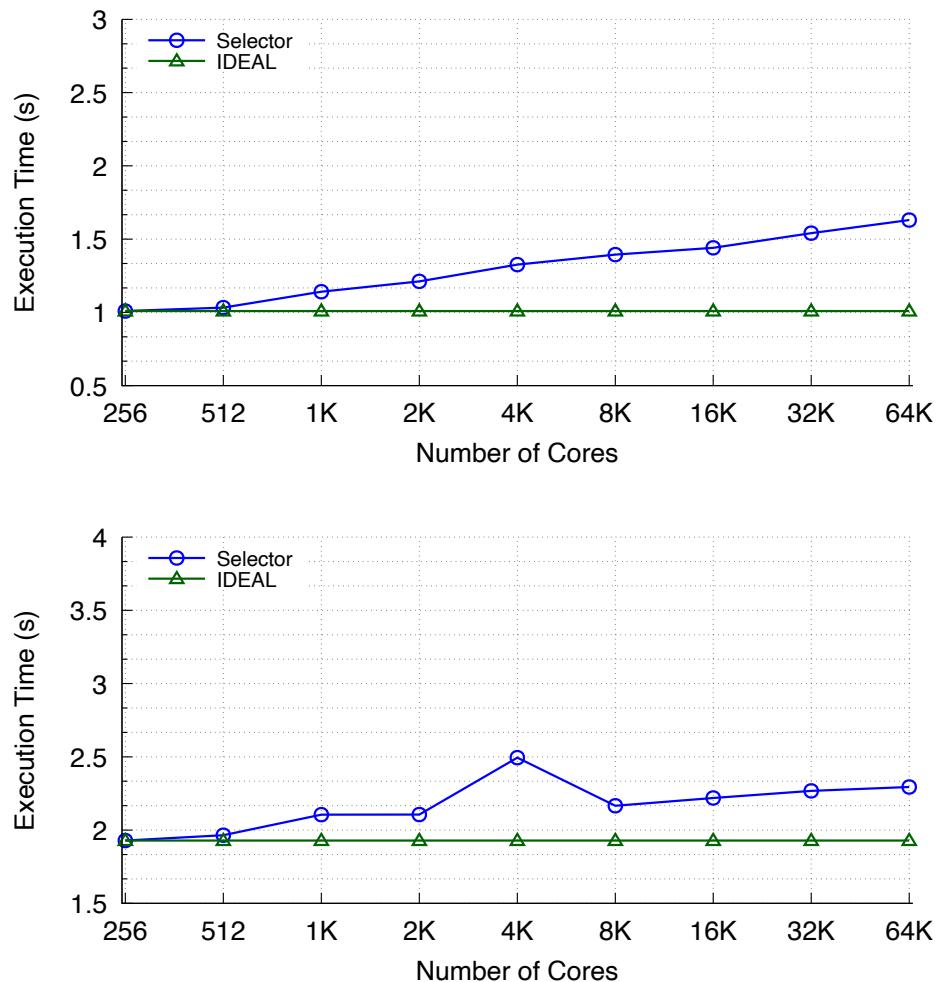
PageRank

13B TEPS

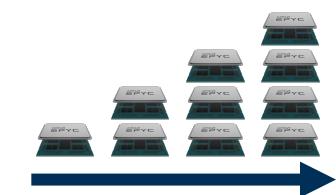
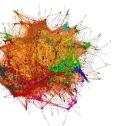
Jaccard Index

85.7% parallel
efficiency

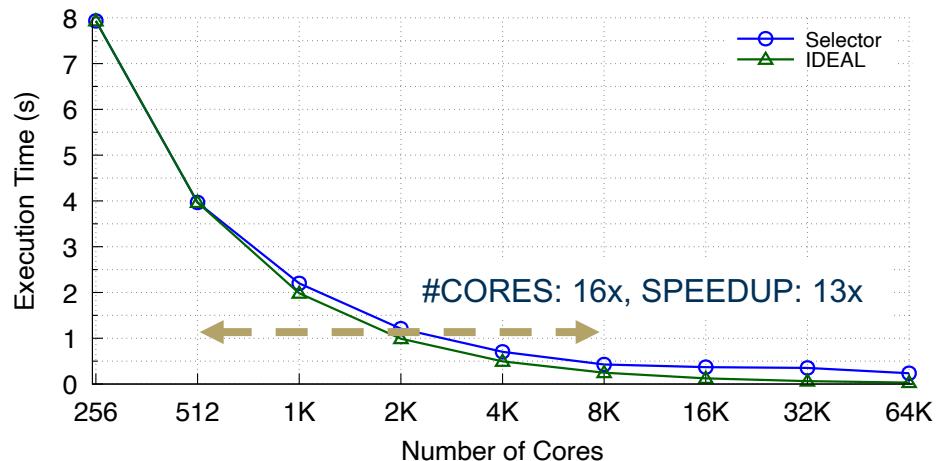
10B TEPS



Performance Results: SCALE2

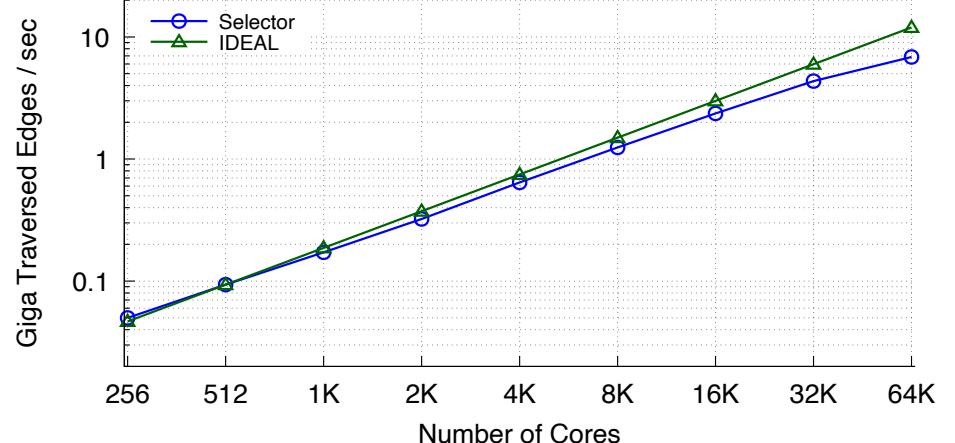
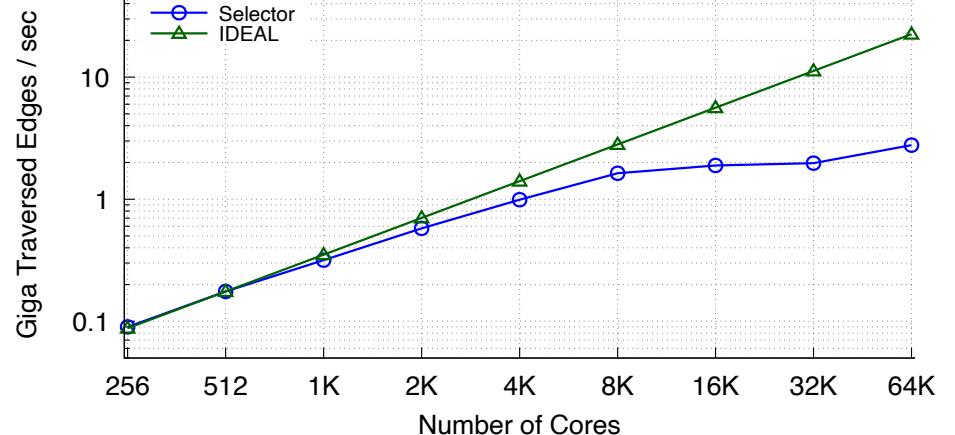
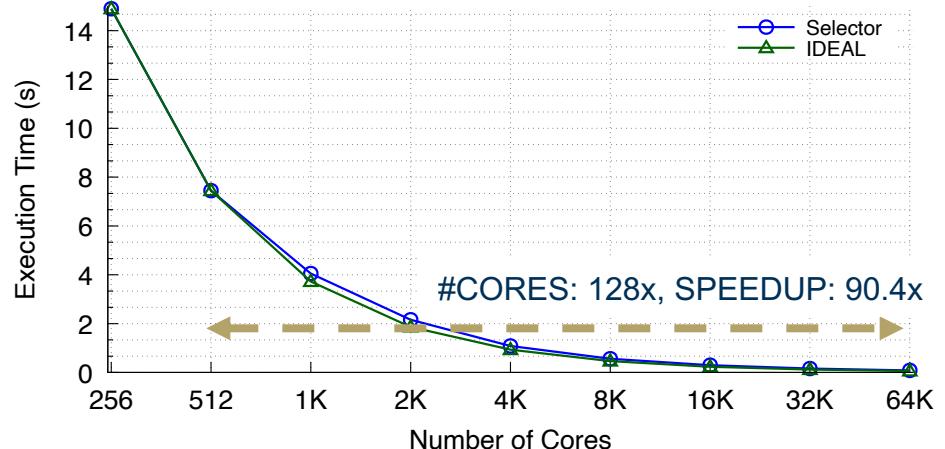


PageRank

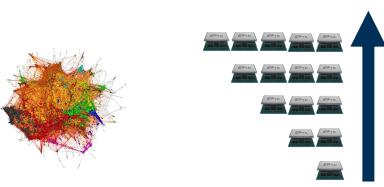


Jaccard Index

70.6% parallel efficiency

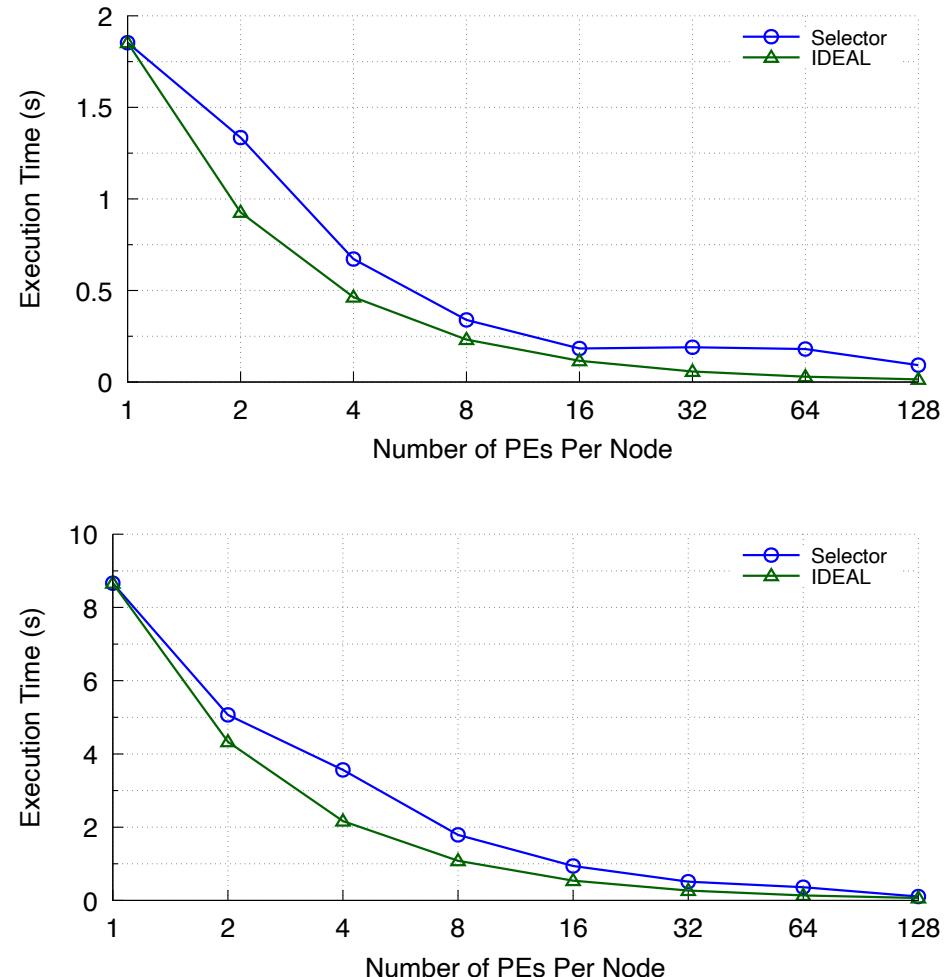


Performance Results: SCALE3



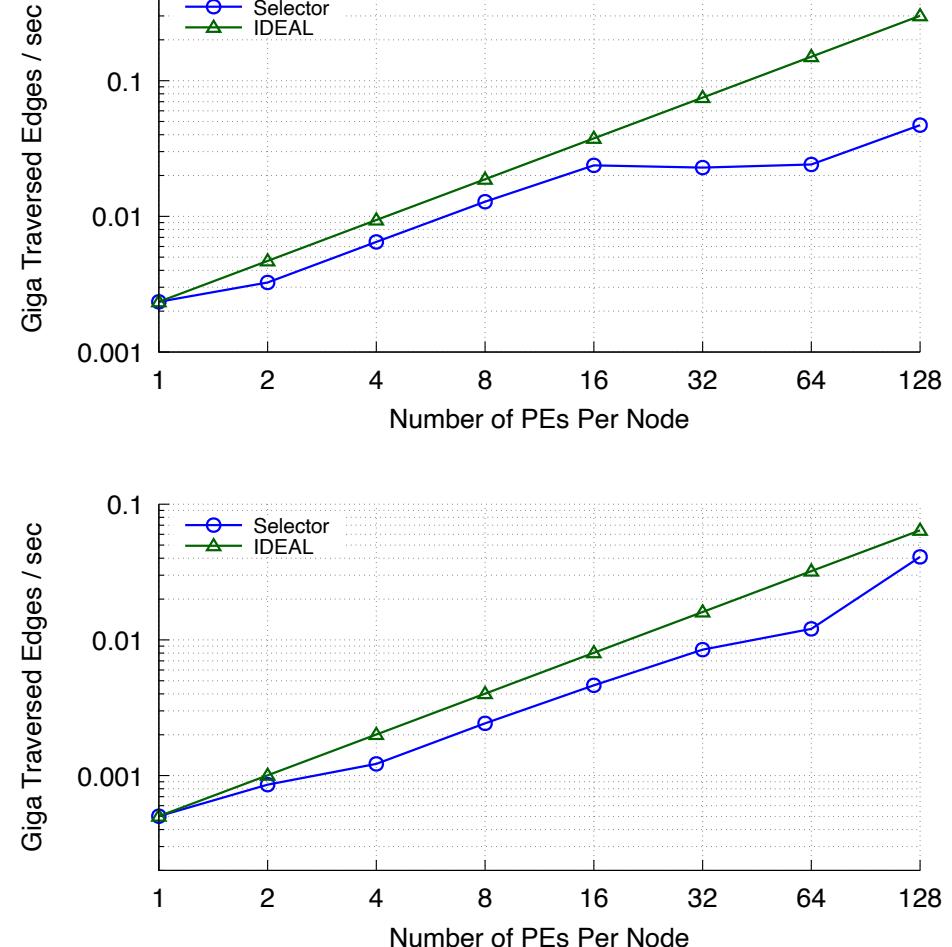
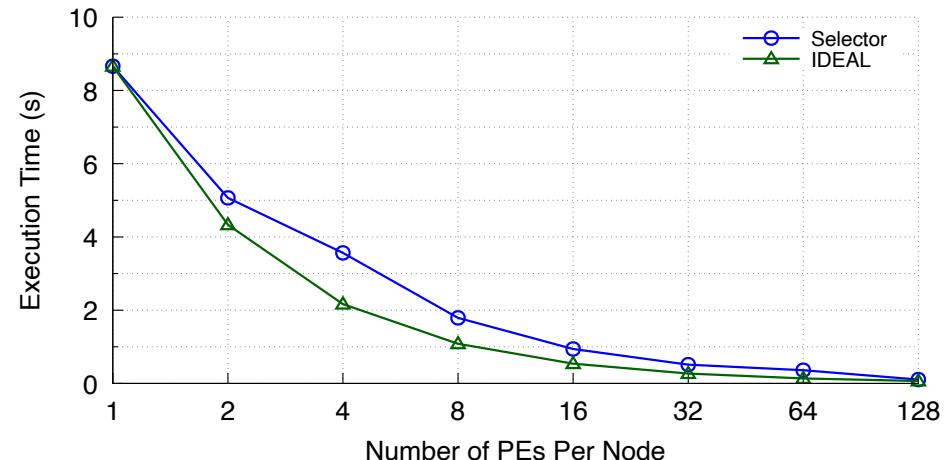
PageRank

20.1x ↑



Jaccard Index

81.5x ↑



Contrasting to Related Approaches

- We contrast with respect to **remote atomics performance** and **graph application performance**
- Related approaches: OpenSHMEM, UPC, MPI3-RMA, YGM

Remote
Atomics

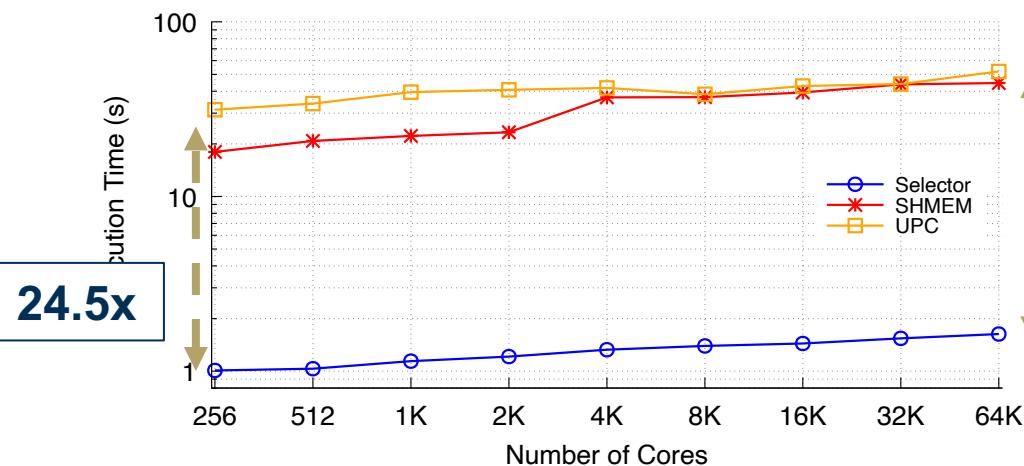
Communication System	Non-Blocking	Read / Get	Update / Set
OpenSHMEM (cray-shmem 7.7.19)	N	40.06	N.A.
OpenSHMEM NBI (cray-shmem 7.7.19)	Y	4.79	4.99
UPC (Berkley-UPC 2022.5.0)	N	30.37	30.03
UPC NBI (Berkley-UPC 2022.5.0)	Y	20.58	N.A.
MPI3-RMA (OpenMPI 4.1.2)	1.8x	25.44	142.04
MPI3-RMA (cray-mpich 7.7.19)	Y	9.67	59.47
YGM	Y	N.A.	> 600
Actors (cray-shmem 7.7.19)	Y	2.70	0.66

7.6x

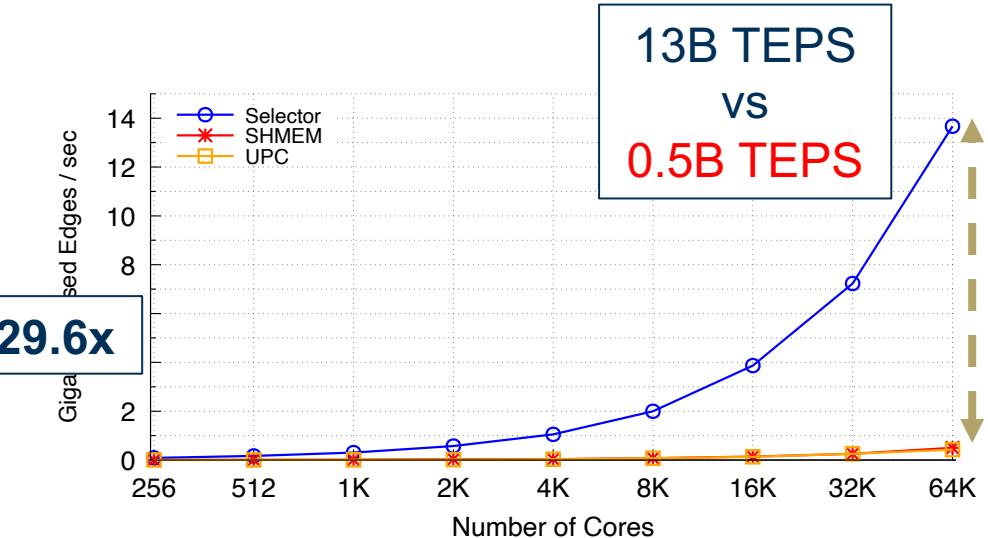
Contrasting to Related Approaches

PageRank

Weak Scaling



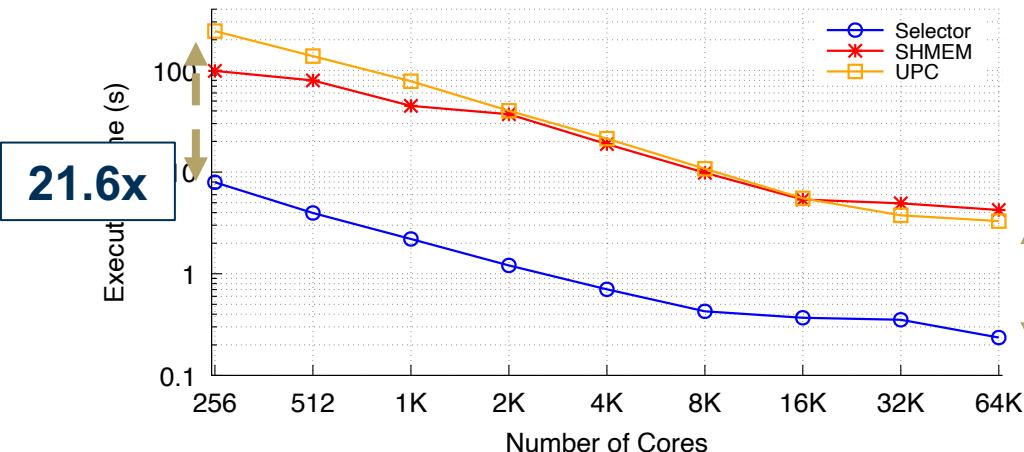
24.5x



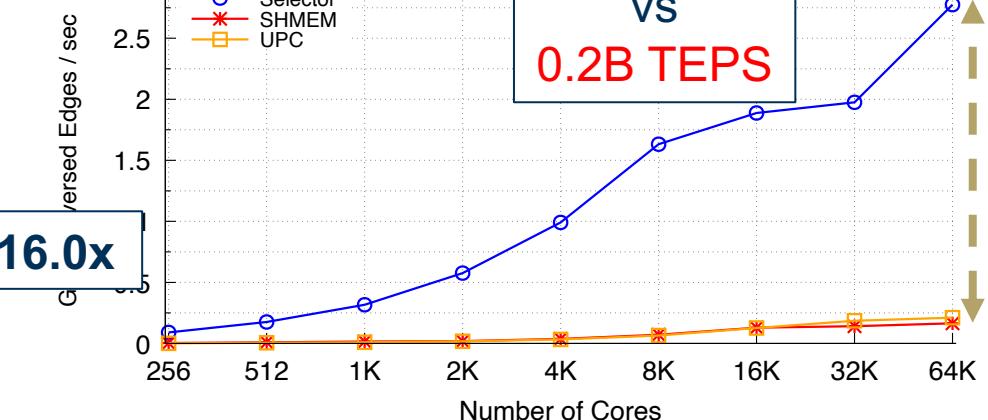
29.6x

13B TEPS
vs
0.5B TEPS

Strong Scaling



21.6x



16.0x

2.8B TEPS
vs
0.2B TEPS

Impact of the Solution

- The actor-based system has shown **scalability, productivity, and performance**
- The extensibility of this system can front four impacts:

1

The system can be used on graphs of higher scale as well as systems of higher scaled-up/-out hardware resources

2

The system can be extended to other large-scale (iterative/non-iterative) graph applications

3

The system can be applied to structured, regular applications

4

The system can be expanded and compared to other related PGAS and non-PGAS approaches

You can try this at home... Just visit hclib-actor.com !

Bulk Synchronous Parallel - HC x +

hclib-actor.com/background/bsp/

HClib-Actor Documentation Search hclib_actor

Table of contents

What is the bulk synchronous parallel model?

Single Program Multiple Data (SPMD) Programming

Further Readings

Bulk Synchronous Parallel

What is the bulk synchronous parallel model?

The Bulk Synchronous Parallel (BSP) model is one of the most popular parallel computation models.

The model consists of:

- A set of processor-memory pairs.
- A communication network that delivers messages in a point-to-point manner.
- Efficient barrier synchronization for all or a subset of the processes.

Diagram illustrating the BSP model:

- Virtual Processors:** PE₀, PE₁, PE₂, ...
- Local Computation:** Each processor performs local computation.
- SUPERSTEP:** A bracket indicates the duration of a superstep.
- Inter-processor Communications:** Arrows show communication between different processors.
- Barrier Synchronization:** A horizontal bar at the bottom represents the barrier synchronization step.

ACKNOWLEDGEMENT

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Thank you for your attention!