## Concordia University Dept. of Computer Science & Software Engineering COMP 353 – Databases SAMPLE MIDTERM EXAM

NOTE: The proportion points in parts A and B below may be different in your actual midterm.

- 1. Max time: 70 minutes.
- 2. Books, notes, sheets are not allowed.
- 3. Make sure you have **7 pages** excluding this cover page.
- 4. Provide your answer at the area below each question; Use other areas for rough work.
- 5. Please do not unstaple the exam booklet.

| Part A / 12 | Part B / 8 | Total / 20 |
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Part A [10 Points] This part has 10 multiple-choice questions.

Note: By "Key" we always mean "minimal key."

- \*\* Of the following statements, which one is NOT correct?
  - E/R is a language for describing database schemas.
  - DDL is a database design notation.
  - DML is a language for accessing the data.
  - ODL is a language for describing database structures.
- \*\* Which of the following statements is NOT correct?
  - A relation instance that has no tuples, may have no keys.
  - A relation with a set of trivial FD's has a unique key.
  - In ODL, it is optional to define a key for a class.
  - A relation with an empty set of FD's has just one key.
- \*\* Let  $R = \{A_1, \ldots, A_k\}$  be a relation schema, and r be an instance of R. Which of the following statements is NOT correct?
  - If r has n tuples, then r can be represented in  $k! \times n!$  ways.
  - Suppose  $\{A_1\}$  is a key of R. If r has n tuples, then the number of different values of  $A_1$  in r could be less than n.
  - If  $\{A_1\}$  is a key of R, then R has  $2^{k-1}$  superkeys.
  - In relational model, every relation has at least one key.

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- \*\* Consider the relations schemas R(A, B), S(A, B), and T(A, B). Suppose  $\{A\}$  is the key of all these relations. If  $U = R \cup S \cup T$ , that is U is the set union of these relations, then what is/are the key(s) of U?
  - {*A*}
  - {*B*}
  - {*A*, *B*}
  - $\{A\}$  and  $\{B\}$
- \*\* Which of the following statements is correct?
  - In the E/R-style method, we normally create more relations than the OO or nulls methods.
  - Taking the OO method, we create a class for all entity sets, even those that have no attributes.
  - Taking the nulls approach, we create a relation for an entity set that has no attributes.
  - Taking the E/R-style method to convert an *isa* hierarchy, we create a relation for an entity set that has no attributes.
- \*\* Consider a relation schema S(A, B, C, D) and the FD's  $F = \{A \to B, BC \to D\}$ . Which of the following statements is NOT correct?
  - F is a cover of itself.
  - The FD  $AC \to D$  follows from F.
  - The FD's  $B \to D$  and  $C \to D$  follow from F.
  - The FD  $AC \to AD$  follows from F.

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\*\* Consider the relation schema  $R = \{A, B, C, D\}$ . Let  $F = \{A \to B, B \to C\}$  and  $G = \{A \to B, B \to C, A \to C\}$ . Which of the following statements is NOT correct?

- $\bullet$  F is a minimal cover for itself.
- G is a cover for F.
- F is a minimal cover for G.
- G is a minimal cover for itself.

\*\* Consider the relation T and FD's F given in the previous question. Which of the following statements is correct?

- $\{A\}$  is the key of T.
- $\{B, C, D\}$  is the key of T.
- $\{A, D\}$  is the key of T.
- $\{A, C, D\}$  is the key of T.

\*\* Which of the following statements is correct?

- A relation instance is a set of attributes.
- A database instance is a set of tuples.
- A database schema is a set of attributes of the existing relations.
- A relation schema is a set of tuples.

\*\* Which of the following statements is NOT correct about meta-data?

- Meta-data describes the indexes to access data efficiently.
- A DBMS uses meta-data to locate data.
- Meta-data describes the structure of the database.
- Meta-data describes the constraints on the database.

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Part B [10 Points] This part has 2 questions, B1 (on this page) and B2 (the next page).

**B1.** [5 Points] Consider the relation schema S = (A, B, C, D) and the sets of FD's

$$F = \{C \to AB, \quad A \to D, \quad B \to D, \quad A \to B, \quad B \to A\}.$$

i. Give a relation instance of S with 2 tuples that satisfies the FD's in F.

ii. Prove or disprove that F covers G ( $F \models G$ ), where

$$G = \{C \to A, A \to D, B \to D, A \to B, B \to A\}.$$

iii. Obtain a minimal cover for F.

iv. How many minimal covers F has? Explain.

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**B2.** [5 Points] This question is on the database design. In a company, there are employees, which have Eid's, Ename's, and Eaddress's. Each employee has one position in his/her "home" department, and could have other positions in the other departments, as part-time. Each employee gets a contract for each part-time employment. Each contract has a salary. No employee gets a total salary which is more than the total salary of any of his/her boss(es). Each employee could have some dependents and the company keeps track of this information for the purpose of insurance benefit only.

(a) Express this information in E/R model. Be "reasonably" creative, but not too much! Be careful, some of this information may not be modeled. Identify such information as well.

(b) Express in ODL the relationship *contract*. For this, you need to identify which classes are involved and how/where to express the relationships.

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Part C [5 Points] For each of the 5 questions in this part, mark T if the given statement is **ALWAYS** true. Otherwise mark F and **justify** your answer.

- ullet Two sets of FD's are equivalent if they yield the same set of keys.  $\Box \mathbf{T}$
- Given any relation instance with 3 tuples, we can determine its FD's.  $\Box T$
- ullet ODL allows using complex attributes of types sets and bags.  $\Box \mathbf{T}$
- If F is a cover for itself, then F is in canonical form.  $\Box \mathbf{T} = \Box \mathbf{F}$
- ullet A "trivial" FD is called so because it is easy to check if it is violated.  $\Box \mathbf{T}$   $\Box \mathbf{F}$