

SOEN331/W- Winter 2017
Introduction to Formal Methods for Software Engineering

Midterm 1 answers

Question 1 (1 points)

Given the following well-formed formulas:

- 1.1 $P \wedge (P \vee Q)$
- 1.2 $\neg P \wedge (P \vee (P \Rightarrow Q))$
- 1.3 $(P \Rightarrow Q) \Rightarrow (\neg P \vee Q)$

Which are tautologies? Use truth tables to prove your answer

Answer: 1.3

Question 2 (6 points)

Consider the argument given by the following sentences.

- P1. If the program does not terminate, then the alarm rings forever*
- P2. Either the computer is not intelligent or the program does not terminate.*
- P3. “the computer runs forever” is implied by the fact that it is not intelligent,*
- Q. Therefore, either the computer runs forever, or the alarm rings forever.*

- 2.1) **Formalize the statement in propositional logic**
- 2.2) **Prove that Q is a logical consequence of the premises P1, P2, P3 using proof by contradiction technique seen in class**

Answer:

T- program terminates

A-alarm rings forever

I-computer is intelligent

R-computer runs forever

P1. $\neg T \rightarrow A$

P2. $\neg I \vee \neg T$

P3. $\neg I \rightarrow R$

Q. $R \vee A$

Proof:

1. $T \vee A$ (from P1)
2. $\neg I \vee \neg T$ (P2)
3. $I \vee R$ (from P3)
4. $\neg R$ (from $\neg Q$.)
5. $\neg A$ (from $\neg Q$.)
6. T (from 1 and 5)
7. I (from 3 and 4)
8. $\neg T$ (from 2 and 6)
9. NIL ((from 6 and 8))

Question 3 (4 points)

Consider the following statements:

1.1 Every student owns some disk space.

1.2 A student may erase the disk space he/she owns.

1.3 A student who receives the disk space from another student cannot erase the contents of the disk space.

Formalize the statements in predicate logic.

Answer:

1.1 $\forall s \in \text{Students} \exists ds \in \text{DiskSpace} \bullet \text{owns}(s, ds)$

1.2 $\forall s \in \text{Students} \forall ds \in \text{DiskSpace} \bullet (\text{owns}(s, ds) \wedge \text{decides-erase}(s, ds) \rightarrow \text{erase}(s, ds))$

1.3 $\forall s, ss \in \text{Students} \forall ds \in \text{DiskSpace} \bullet (s \neq ss \wedge \text{owns}(s, ds) \wedge \text{receives}(s, ss, ds) \rightarrow \neg \text{erase}(ss, ds))$

Question 4 (2 points. Circle the right answer(s))

4.1 Which of the following predicates are valid formalizations of “*Somebody likes somebody*”?

(a) $\exists p : Person \bullet likes(p, q)$	(c) $\exists p : Person \bullet \exists q : Person \bullet likes(p, q)$	(f) None of the previous.
(b) $\exists q : Person \bullet likes(p, q)$	(d) $\forall p : Person \bullet \exists q : Person \bullet likes(p, q)$	
	(e) $\exists p : Person \bullet \neg \exists q : Person \bullet likes(p, q)$	

Answer: c)

4.2 Which of the following predicates are valid formalizations of “*some cats are sleepy*”?

Note: *isAcat(x)* is true if *x* is a cat.

(a) $\exists x : Animal \bullet isAcat(x) \rightarrow sleepy(x)$	(d) $isAcat(x) \& sleepy(x)$	(g) $\exists x : Animal \bullet sleepy(x)$
(b) $\exists x : Animal \bullet isAcat(x) \& sleepy(x)$	(e) $isAcat(x) \rightarrow sleepy(x)$	(h) $\exists c : Cat \bullet sleepy(c)$
(c) $\exists c : Cat \bullet isAcat(c) \rightarrow sleepy(c)$	(f) $isAcat(c) \& sleepy(c)$	(i) None of the previous.

Answer: b) and h)

Question 5 (2 points)

Is Propositional Logic sound (consistent)? complete? Explain your answer.

Propositional Logic is **Sound** —All provable statements are semantically true. That is, if a set of premises *S* syntactically entails a proposition *P*, then there is an interpretation in which *P* can be reasoned about from *S*.

if $S \vdash P$, then $S \models P$

Propositional Logic is **Complete**—All semantically true statements are provable. That is, if a set of premises *S* semantically entails a proposition *P*, then *P* can be derived formally (syntactically) within the formalism

if $S \models P$, then $S \vdash P$